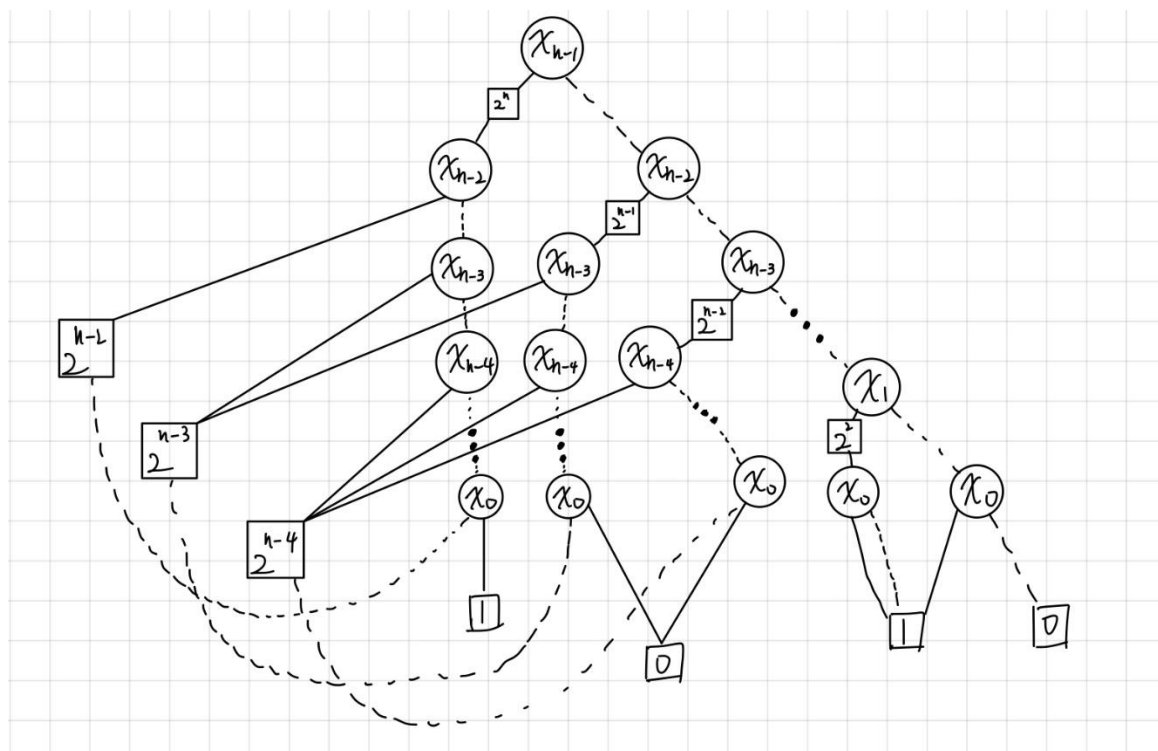


Problem 5

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(a)



(b) First, define the branch of i -th bit x_i is a sub-*BMD consists of the positive sub-*BMD of node x_i and node x_i itself. Here the node x_i should stay in the negative sub-*BMD of every node with higher index. The weight of a set of branch is the sum of all weight of branches in the set. After some definition above, we can dive into the pseudo part:

Def. Function SubRoot(Y, A, t, k, w_0, s):

```

If( $k - w_0 \geq 0$ ) then
    If( $Y \leq t + 2^{(k-w_0)} * \text{weight}(A) + 2^{2(k-w_0)}$ ) then
         $t = t + 2^{(k-w_0)} * \text{weight}(A) + 2^{2(k-w_0)}$ 
         $s = k - w_0$ 
        add the branch of  $x_s$  to set  $A$ 
        return ( $A, t, s, \text{weight}(A)$ )
    else
        ( $A, t, s, \text{weight}(A)$ ) = SubRoot ( $Y, A, t, k-1, w_0, s$ )
        return ( $A, t, s, \text{weight}(A)$ )
    end If
else
     $s = k - w_0$ 
    return ( $A, t, s, \text{weight}(A)$ )
end If

```

End function

The SubRoot function find the partial root of Y and help us solve the problem recursively. The main algorithm is presented below:

Initialization

```

assume  $y_{k_0}$  is the first not zero significant bit in  $Y$ 
let  $s = \lfloor k_0/2 \rfloor$ ,  $A = \{\}$ ,  $t = 2^{2s}$ ,  $s_0 = \lfloor k_0/2 \rfloor$ ,  $w_0 = \text{weight}(\text{the branch of } x_{s_0})$ , and  $w = 0$ 

```

add the branch of x_s to set A

end Initialization

While($s \geq 0$) **do**

find the first non-zero significant bit y'_k for $y' = Y - t$.

(A, t, s, w) = **SubRoot** (Y, A, t, k, w_0, s)

End While

Then the integer part of the square root of Y would be $w/2$.