

*Note:* Your TA probably will not cover all the problems. This is totally fine, the discussion worksheets are deliberately made long so they can serve as a resource you can use to practice, reinforce, and build upon concepts discussed in lecture, readings, and the homework.

## 1 Asymptotics and Limits

If we would like to prove asymptotic relations instead of just using them, we can use limits.

**Asymptotic Limit Rules:** If  $f(n), g(n) \geq 0$ :

- If  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} < \infty$ , then  $f(n) = \mathcal{O}(g(n))$ .
- If  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c$ , for some  $c > 0$ , then  $f(n) = \Theta(g(n))$ .
- If  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} > 0$ , then  $f(n) = \Omega(g(n))$ .

Note that these are all sufficient (and not necessary) conditions involving limits, and are not true definitions of  $\mathcal{O}$ ,  $\Theta$ , and  $\Omega$ . We highly recommend checking on your own that these statements are correct!)

(a) Prove that  $n^3 = \mathcal{O}(n^4)$ .

(b) Find an  $f(n), g(n) \geq 0$  such that  $f(n) = \mathcal{O}(g(n))$ , yet  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \neq 0$ .

(c) Prove that for any  $c > 0$ , we have  $\log n = \mathcal{O}(n^c)$ .

*Hint:* Use L'Hôpital's rule: If  $\lim_{n \rightarrow \infty} f(n) = \lim_{n \rightarrow \infty} g(n) = \infty$ , then  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{f'(n)}{g'(n)}$  (if the RHS exists)

- (d) Find an  $f(n), g(n) \geq 0$  such that  $f(n) = \mathcal{O}(g(n))$ , yet  $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$  does not exist. In this case, you would be unable to use limits to prove  $f(n) = \mathcal{O}(g(n))$ .

*Hint: think about oscillating functions!*

## 2 Recurrence Relations

Solve the following recurrence relations, assuming base cases  $T(0) = T(1) = 1$ :

(a)  $T(n) = 2 \cdot T(n/2) + O(n)$

(b)  $T(n) = 2 \cdot T(n/2) + O(n \log n)$

$$(c) \ T(n) = T(n - 1) + n$$

$$(d) \ T(n) = 3 \cdot T(n - 2) + 5$$

$$(e) \ T(n) = 3T(n^{1/3}) + O(\log n)$$

$$(f) \ T(n) = T(n - 1) + T(n - 2)$$