

Assignment10 – Knowledge Representation

Problem 10.1 (Modeling in Description Logic)

Consider the following situation:

- Some beings are persons, some are animals.
- Persons and animals may like other persons or animals.
- Alice is a person, and she likes the animal Bubbles.

1. Model this situation as a semantic network. Explain the different kinds of nodes and edges occurring in your network.

Solution: Kinds of nodes:

- represent concepts: being, person, animal
- represent individuals: Alice, Bubbles

Kinds of edges:

- assert that one concept is a subconcept of another (is-a): person-being, animal-being
- assert that one individual is an instance of a concept (inst): Alice-person, Bubbles-animal
- represent a relation between two concepts: an edge labeled 'like' relation person-animal (These edges can often be omitted.)
- assert that two individuals are in a relation: an edge labeled 'like' Alice-Bubbles

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2. Model the same situation in first-order logic and compare the results.

Solution: Nodes (by kind):

- For each concept-node c , we need a unary predicate symbol: $\text{being} \in \Sigma_1^p, \text{person} \in \Sigma_1^p, \text{animal} \in \Sigma_1^p$.
- For each individual-node i , we need a nullary function symbol: $\text{Alice} \in \Sigma_0^f, \text{Bubbles} \in \Sigma_0^f$.

Edges (by kind):

- For each subconcept-edge (is-a) from c to d , we need an axiom $\forall x.c(x) \Rightarrow d(x)$. Here: $\forall x.\text{person}(x) \Rightarrow \text{being}(x)$ and $\forall x.\text{animal}(x) \Rightarrow \text{being}(x)$.
- For each instance-edge (inst) from i to c , we need an axiom $c(i)$. Here: $\text{person}(\text{Alice})$ and $\text{person}(\text{Bubbles})$.

- For each relation-representing-edge r from c to d , we need a binary predicate symbol. Here: $like \in \Sigma_2^P$. To represent the subject and object concepts, we can add axioms $\forall x, y. r(x, y) \Rightarrow c(x)$ and $\forall x, y. r(x, y) \Rightarrow d(y)$. Here: $\forall x, y. like(x, y) \Rightarrow person(x)$ and $\forall x, y. like(x, y) \Rightarrow animal(y)$.
 - For each relation-asserting-edge for r from i to j , we need axiom $r(i, j)$. Here: $like(Alice, Bubbles)$.
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3. Explain the difference between inst and is-a edges.
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Solution: inst edges are between an individual and a concept. They correspond to \in math.

is-a edges are between two concepts. They correspond to \subseteq in math.

4. Explain the difference between having a relation edge between two concepts vs. asserting a relation between two individuals.
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Solution: The former models that a relation r exists between individuals of the respective concepts. That corresponds to a FOL-predicate symbol in Σ_2^P . For each pair of individuals the relation may be true or false.

The latter asserts that that relation holds for two individuals. For example, an r -edge from i to j corresponds to the FOL-axiom $r(i, j)$.

Problem 10.2 (ALC)

Consider the following description logic signature

- **concept** symbols: i (for instructor), s (for student), c (for **course**), p (for program)
- **role** symbol m (for is-member-of) used for
 - **instructors** giving a **course**
 - **students** taking a **course**
 - **students** being enrolled in a **degree program**
 - **courses** being part of a **degree program**

We use an extension of \mathcal{ALC} , in which there are dual roles: there is a role m^{-1} that captures the relation has-as-member, e.g., $MK m AI$ iff $AI m^{-1} MK$.

1. For the **signature** above, give a **concept axiom** that captures that instructors can only be members of **courses**.
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Solution: $i \sqsubseteq \forall m.c$

2. Give a **concept axiom** for the above **signature** that captures: **courses** that are taken by a **student**, must be given by an **instructor**.

Solution: $c \sqcap \exists m^{-1}.s \sqsubseteq \exists m^{-1}.i$

3. Calculate the translation to **first-order logic** of $s \sqsubseteq \forall m.\exists m.p$.

Solution: $\forall x.s(x) \Rightarrow (\forall y.m(x, y) \Rightarrow (\exists z.m(y, z) \wedge p(z)))$

4. Given a **first-order model** $\langle \mathcal{D}, \mathcal{I} \rangle$, define an appropriate case of the **interpretation** mapping for the formula $\forall r^{-1}.C$.

Solution: $\llbracket \forall r^{-1}.C \rrbracket = \{u \in \mathcal{D} \mid \text{for } v \in \mathcal{D}, \text{ if } (v, u) \in \llbracket r \rrbracket, \text{ then } v \in \llbracket C \rrbracket\}$

Problem 10.3 (ALC Semantics)

Consider the **ALC** concepts $\forall R.(C \sqcap D)$ and $\forall R.C \sqcap \forall R.D$.

1. By applying the semantics of **ALC**, show that the two are equivalent.

Solution: We have:

$$\begin{aligned} & \llbracket \forall R.(C \sqcap D) \rrbracket \\ &= \{x \in \mathcal{D} \mid \text{for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in \llbracket C \sqcap D \rrbracket\} \\ &= \{x \in \mathcal{D} \mid \text{for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in (\llbracket C \rrbracket \cap \llbracket D \rrbracket)\} \end{aligned}$$

$$\begin{aligned} & \llbracket \forall R.C \sqcap \forall R.D \rrbracket \\ &= \llbracket \forall R.C \rrbracket \cap \llbracket \forall R.D \rrbracket \end{aligned}$$

$$= \{x \in \mathcal{D} \mid \text{for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in (\llbracket C \rrbracket \cap Q)\}$$

where

$$Q = \{x \in \mathcal{D} \mid \text{for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in \llbracket D \rrbracket\}$$

Now to prove that sets are equal, consider an $x \in \mathcal{D}$ and see that both conditions are equivalent to

$$\text{for all } y \in \mathcal{D}, \text{ if } (x, y) \in \llbracket R \rrbracket, \text{ then } y \in \llbracket C \rrbracket \text{ and } y \in \llbracket D \rrbracket$$

2. Translate both formulas to first-order logic and state which FOL formula we would need to prove (e.g., with the ND calculus) to show that the two are equivalent.

Solution: The translation yields

$$C_1(x) = \forall y. R(x, y) \Rightarrow C(y) \wedge D(y)$$

$$C_2(x) = (\forall y. R(x, y) \Rightarrow C(y)) \wedge (\forall y. R(x, y) \Rightarrow D(y))$$

We need to show

$$\forall x. C_1(x) \Leftrightarrow C_2(x)$$

Problem 10.4 (ALC TBox)

Consider *ALC* with the following

- primitive concepts: woman, man
- roles: has_child, has_parent, has_sibling, has_spouse

Give an *ALC* TBox that defines the concepts person, parent, mother, father, grandmother, aunt, uncle, sister, brother, onlychild, cousin, nephew, niece, fatherinlaw, motherinlaw.

Solution: person = man \sqcup woman

parent = person \exists has_child. person

mother = woman parent

father = man parent

grandmother = woman \exists has_child. parent

aunt = woman \exists has_sibling. parent

uncle = man \exists has_sibling. parent

sister = woman \exists has_sibling. person

brother = man \exists has_sibling. person

onlychild = person brother \sqcup sister

cousin = person \exists has_parent. \exists has_sibling. parent

nephew = man \exists has_parent. \exists has_sibling. person

niece = woman \exists has_parent. \exists has_sibling. person

fatherinlaw = man \exists has_child. \exists has_spouse. person

motherinlaw = woman \exists has_child. \exists has_spouse. person
