

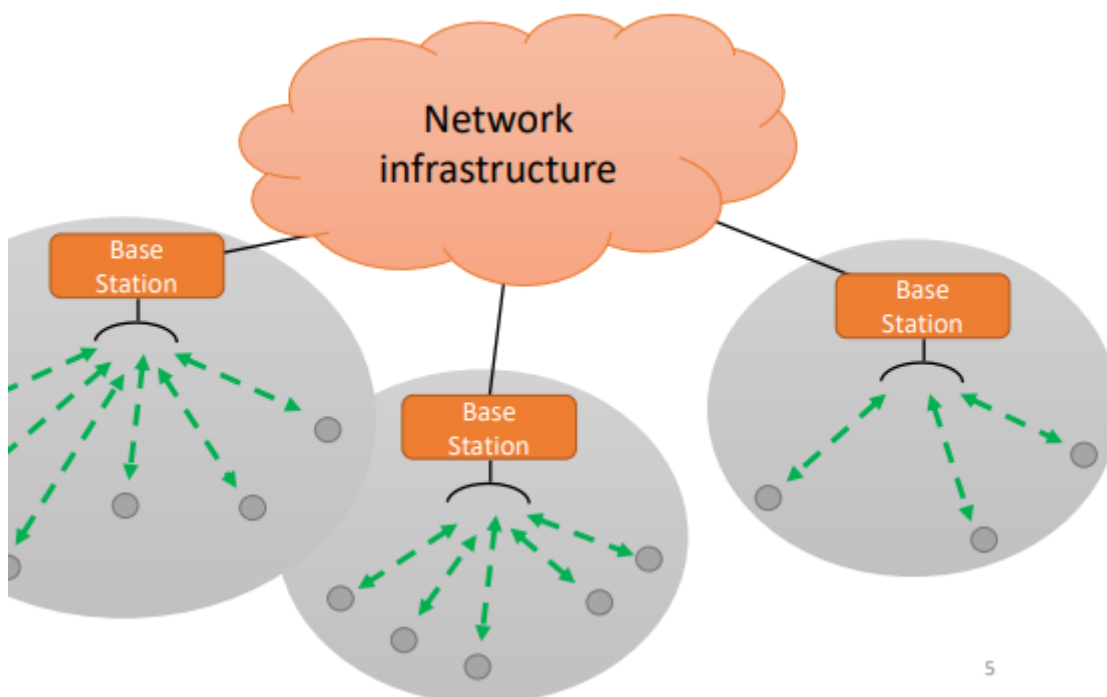
B 01 - Wireless Networks

Why wireless networks?

- need cables in computer to
 - communicate: transfer data between nodes
 - provide power: between devices
- Cyber-physical systems embed computers in physical objects
 - they're everywhere: free placement, mobility, exposed to the wild → go wireless (can't wire everything)
 - hence: wireless to replace cable in communications and batteries to replace cables in power supply

Wireless networks

- network of **hosts** connected by **wireless links**
 - **hosts**: end-system devices that runs apps
 - mobile (often, but not necessary)
 - battery-powered (typically)
- 2 modes of operations
 1. **infrastructure**: BS(base stations) || wired access points
 2. **ad hoc networking**: no centralised coordinators (that coordinate shared medium), links connected wired
- **elements of wireless network**



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- **wireless hosts:**

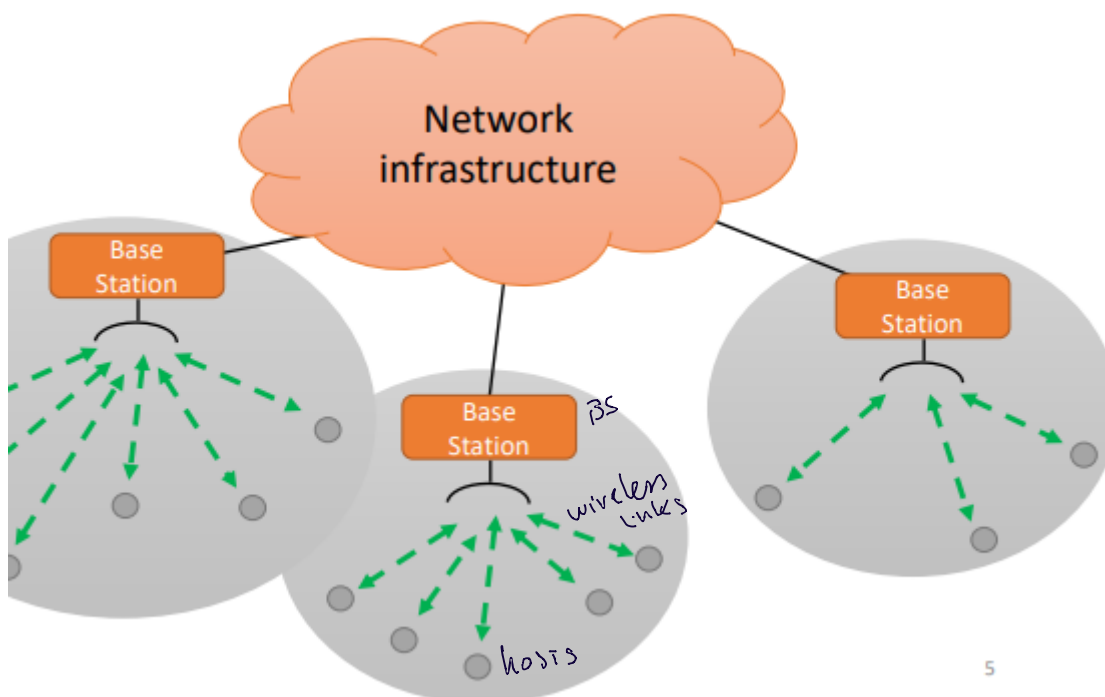
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- **wireless hosts:**

- run applications
- stationary (fixed location), wireless not always mean mobility || mobile
- devices want to communicate to internet → traverse internet
- are connected wireless to BS
- **base stations:**
 - cover geographic areas
 - all communications to devices traverse BS
 - manage resource of physical medium
 - connected to a wired network
 - *relay*: responsible to send packets between wired network and wireless hosts (in its area)
 - e.g.: access points (in wi-fi) / cell towers (in mobile phone)...
 - **problem** to solve for **access points**: new accesses, coordination medium, moving devices
- **wireless links:**
 - connect mobile to base stations
 - *Multiple Access Protocol* to coordinate *link access*
 - various *data rates*, *transmission range*
- **infrastructure mode** (network infrastructure):
 - BS connects mobiles into wired network
 - problem: nodes can move between BSs → exchanging **control message**: connection doesn't terminate → **change of responsibility**
 - devices can move between BS of different/same **providers**:
 - they provide authentication, billing
- **intracell communication**: two hosts (end devices) communicate inside the same BS's area
- **intercell communication**: two hosts communicate through different BS's areas

wireless network taxonomy

	Single hop	Multiple hops
infrastructure (e.g., APs)	host connects to base station (WiFi, WiMAX, cellular 3G,4G,5G) which connects to larger Internet	host may have to relay through several wireless nodes to connect to larger Internet: MESH networks
no infrastructure	no base station, not necessarily connection to larger Internet (e.g. Bluetooth)	no base station, no connection to larger Internet. May have to relay on other nodes to reach a given wireless node (ZigBee, ad hoc, VANET)

Wireless links (characteristics) vs wired links

- electromagnetic waves in the air (tonight, oh lord, I can feel it comin' in the air tonight, Phil Collins)

1) • decreased signal strength:

- problem
- signal decreases with length
- radio signal attenuates as it propagates thru matter
 - obstacles can block electromagnetic waves

2) • interference with other sources:

- standardized wireless network frequency shared by other devices (e.g.: 2,4Ghz)
- microwaves ovens, mobile phones, engine, appliances may *interference as well
- receivers detect the sum of signal in the environment

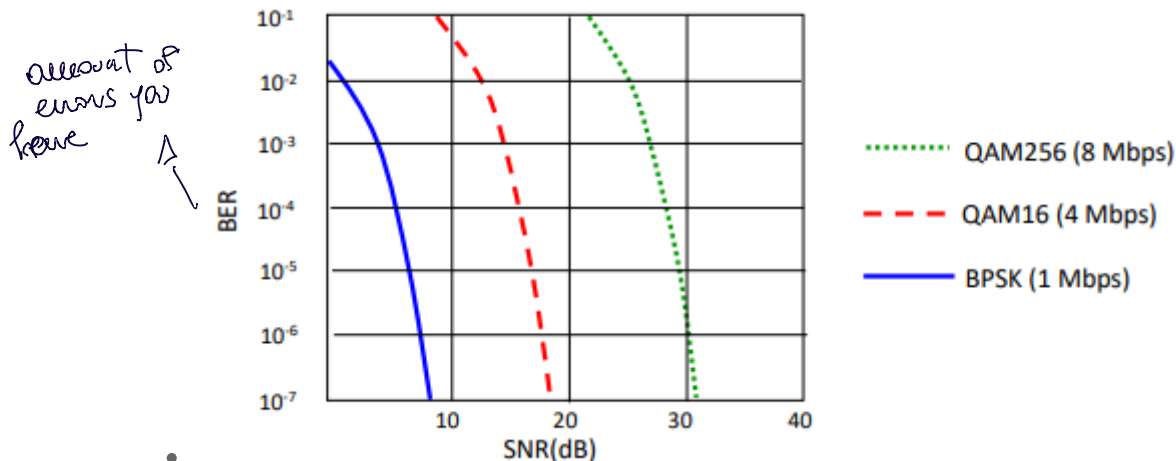
3) • multipath propagation*:

- radio signal (electro waves) reflect on objects or ground, arriving at destination at different times

Wireless tools

- **SNR: signal to noise ratio** → evaluate *quality of the signal*
 - $SNR = (\text{signal power} / \text{noise power})$
 - id est (meaningful input / meaningless or unwanted input (power of background noise)) such that $0 < SNR < +\infty$, (watt or dBm)
 - SNR in dB (10th part of bel (B symbol): $10 \text{ dB} = 1\text{B}$, *logarithmic unit ratio over two homogeneous quantities*) = $20 * \log(\text{signal/noise})$
 - *larger is SNR → easier is to extract signal from noise (cool thing)*
- **BER: Bit Error Rate**: probability that a transmitted bit is received in error at the receiver

- **SNR vs BER tradeoff**: how implicit is information in electro waves
 - **given a physical layer**: increase power ^{→ SAFER WAVE & CLEAR} → increase SNR (quality of the signal) → decrease BER (send bits → electro waves to the medium)
- **given SNR**: choose a physical layer that meets BER requirement, **giving highest throughput**
- **SNR** may change with *mobility*:
 - dynamically adapt physical layer (modulation technique)



- also **signal attenuation** or **obstacles limit transmission range**

Wireless network challenges

1. **Limited knowledge**:
 - a node detect a situation, but not all communications
 - a terminal cannot hear all the other

⇒ hidden/exposed terminal problem
2. **Mobility/failure of terminals**:
 - terminals move (join/exit) in the range of different BS
 - terminals move away from each other
3. **Limited terminals**:
 - battery life, memory, processing, trasmission range
4. **Privacy**:
 - eavesdropping of ongoin communications

INTERCEPTION

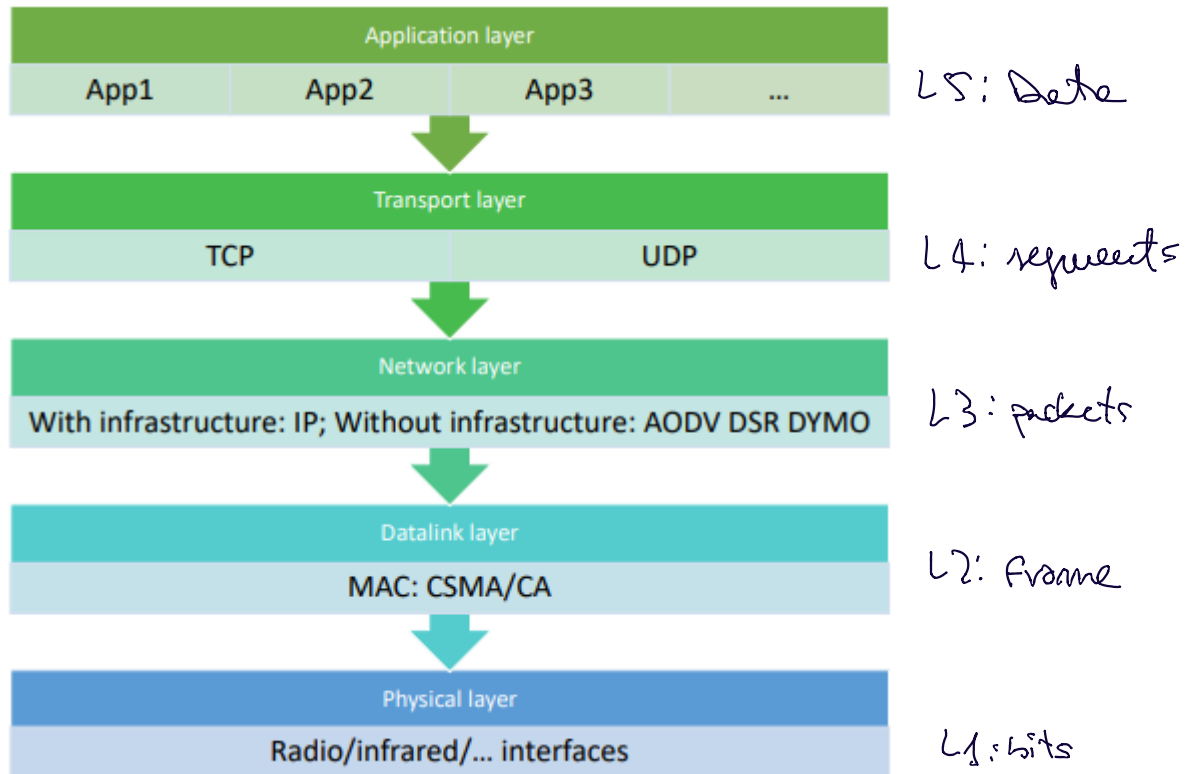
Wireless network required mechanism

1. **Access** to a shared wireless channel
 - CSMA/CD (Carrier Sense Multiple Accesses with Collision Detection) multi access protocol CANNOT be used.
2. **Hand-off** ("goin' away") (network with infrastructure):
 - moving terminal into range of different BS
3. **Routing** (multi-hop ad hoc networks)
 - finding a path form **source** to **destination** in multi-hops networks

- dealing with arbitrary changes in neighborhood

Wireless networks protocol stack

- | bits | frames | packets | segments | data |

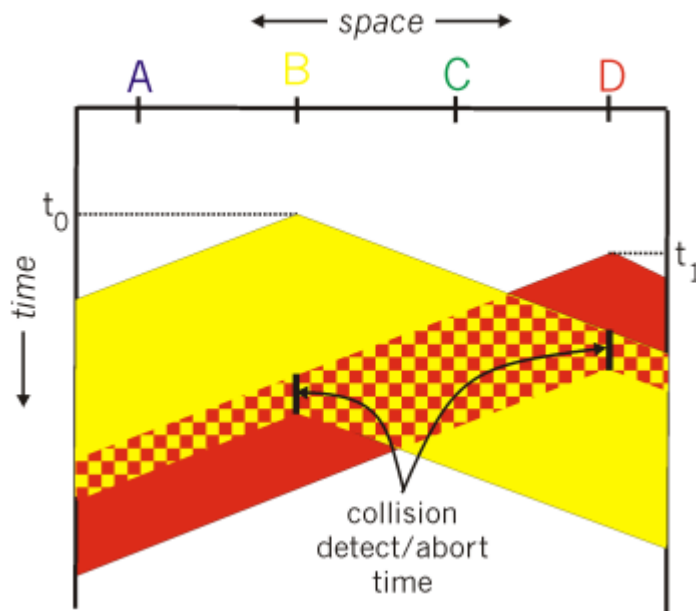


Recap: MAC protocol for wired networks

- a single channel available for **all** the communications
- **all** stations can transmit and send on the channel
- if frames are sent simultaneously on the channel
 - → resulting signal is **garbled (a collision occurs*)**
- CSMA, CSMA/CD, ALHOA, slotted ALOHA

CSMA

- Carrier Sense Multiple Accesses
- when a station has a frame to send → listen to the channel to see if anyone else is transmitting
- if channel is busy → station waits until it becomes idle
- when channel is idle → station transmits the frame
- if a collision occurs → station waits a random amount of time and repeats
-



CSMA/CD

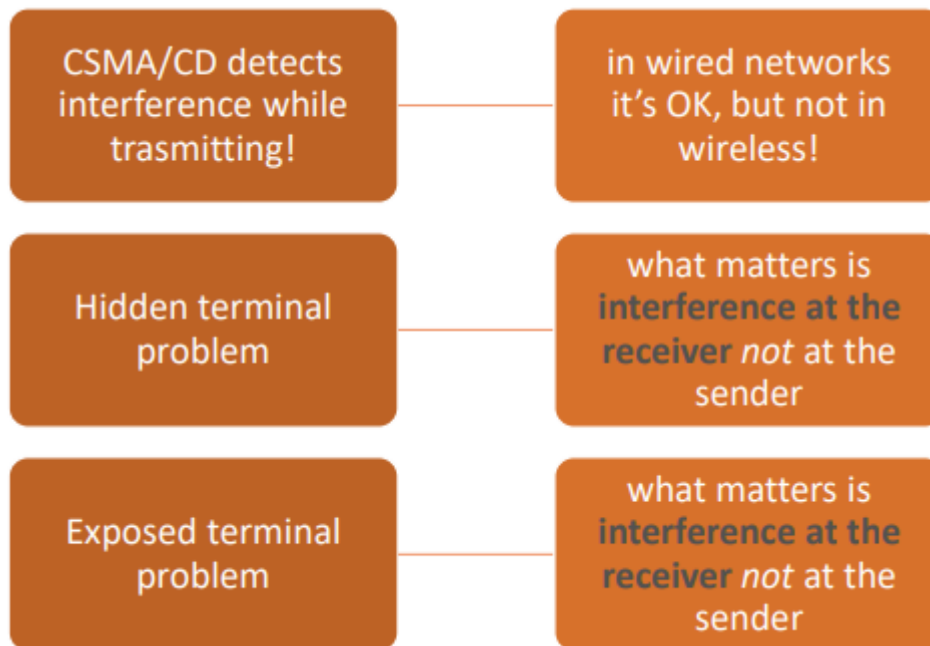
- a station **aborts its transmission** as soon as it **detects a collision**
 - if **two stations sense channel idle** at the same time and **start transmitting**
 - → stations **quickly abort the frame** as soon as **collision is detected**
- widely used in LANs in MAC sub-layer
- IEEE 802.3
- **T** is the time required to reach the farthest station
- it takes minimum RTT (# time between sending/receiving the signal) ($2 * T$) to **detect collision**

Binary Exponential Backoff

- defines **time** that station has to wait for transmitting again a frame
- algorithm:
- **time after a collision** is divided in **contention slots**:
 - **length of a contention slot** is equal to the **worst case round propagation time** ($2T$ if T time to reach farthest station)
- after **1st collision**:
 - each station waits 0 or 1 slot before trying again (success or collision)
- after **collision i**:
 - choose x at random in $(0, 2^i - 1)$ → exp: more collision, more time, may increase exponentially
- after **10 collisions**:
 - choose x at random in a frozen interval $(0..1023)$
- after **16 collisions**:
 - failure is reported back to upper levels

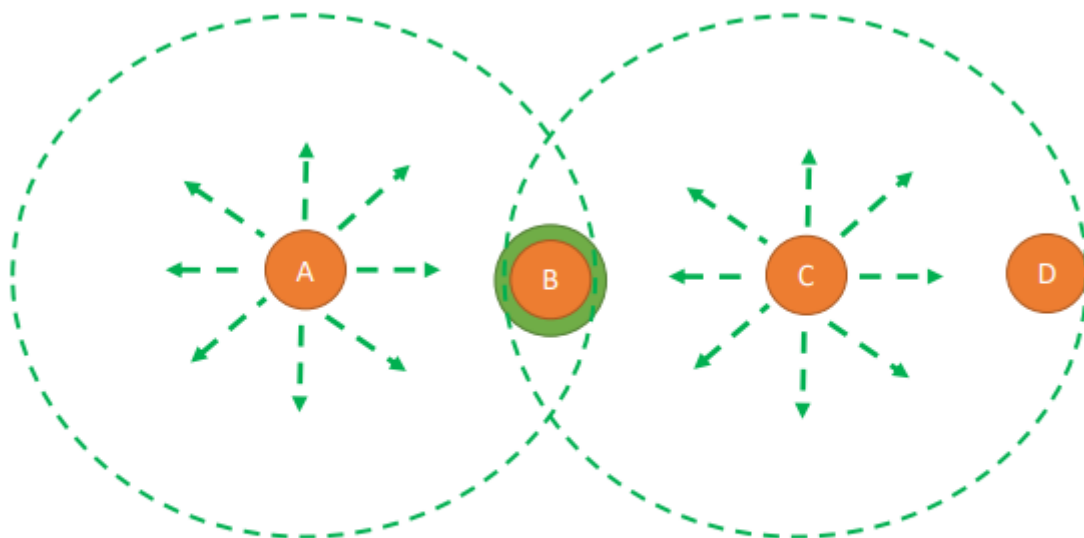
Why MAC doesn't work on wireless

- station transmit and send if there is another communication in channel → multiple antennas



Hidden terminal problem

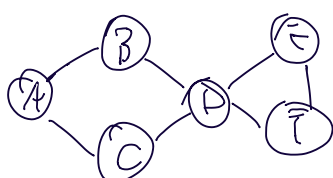
- two or more stations which are out of the range of each other transmit simultaneously to a comme recipient



A is sending to B

C senses the medium: it will NOT hear A, out of range

- C transmits to anybody (either B or to D) : **COLLISION at B!**



1) Hidden terminal $E \rightarrow D$: can interfere: $B \rightarrow D$ \Rightarrow F also in range of E (send channel) \Rightarrow no problem
 2) $D \rightarrow C$: A doesn't receive why from $D \Rightarrow$ hidden
 3) Exposed $D \rightarrow B$: C, E, F will not make the interfere with B , but send channel as $E \rightarrow D$
 4) $B \rightarrow A$: D

will not sense anything when T, destroy comm.

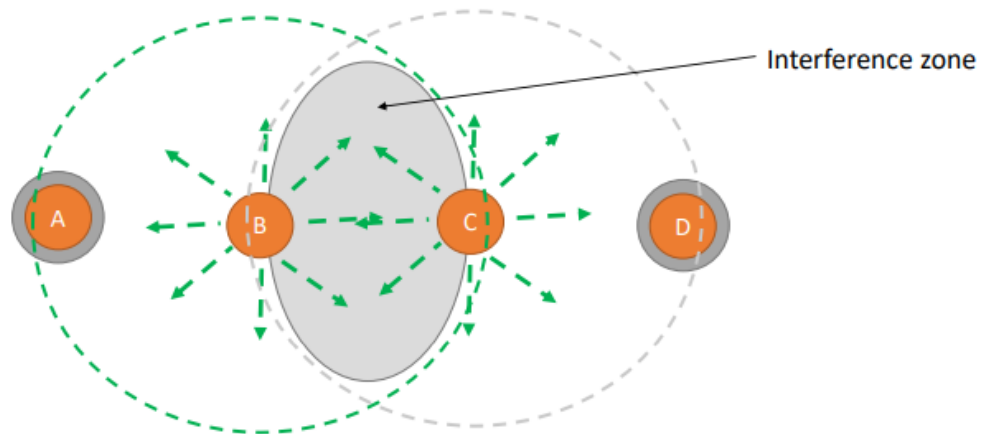
Hidden terminal problem

C is not able to detect a potential competitor because it is out of range:
a collision happens at B (the receiver)
For the same reason A does not detect the collision
C is **hidden** with respect to the communication from A to B

•

Exposed terminal problem

- a transmitting station is prevented from sending frames due to interference with another transmitting station



1. B is transmitting to A, C wants to transmit to D
2. C senses the medium, concludes: **cannot transmit** to D
3. The two transmissions can actually happen in parallel.

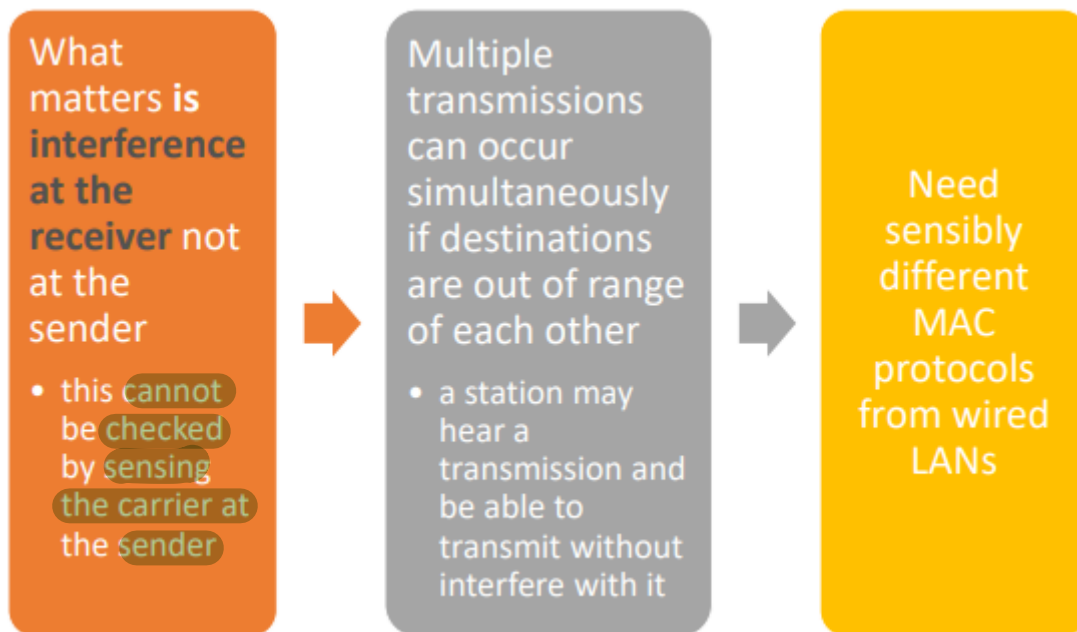
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Exposed terminal problem

C hears a transmission
C does not send to D although its transmission would be OK
C is **exposed** with respect to the communication from B to A

• •

- so:



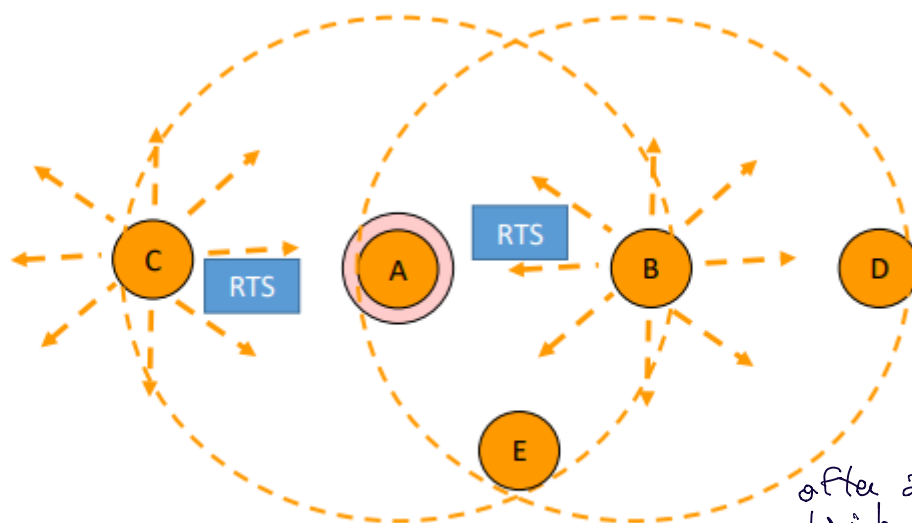
MACA protocol

CSMA/CA

- Multiple Accesses with Collision Avoidance:
 - stimulate the receiver into trasmittig to it a short frame first by sender
 - RTS**: Request To Send, includes length of the longer frame
 - if the station receiver wants to recive the message, replies with another frame
 - CTS**: Clear To Send short frame with data lenght copied from RTS frame
 - then transmit a (long data frame)
 - stations hearing the short frame **refrain from transmitting** during the transmission of the subsequent data frame

- collision:

time to keep farthest station T



C and B send RTS simultaneously to A

- the two messages collinde: **no CTS generated**

- C and B uses **Binary exponential backoff** (same as ethernet) to retry RTS

after 2 collisions,
divide frame in contention
slots = 2T
- after 1st collision
wait 0,1 contention slots
- after collision 2
wait 0,1,2 contention slots
- after collision 3
wait 0,1,2,3 contention slots
- after collision 4
wait 0,1,2,3,4 contention slots
- after collision 5
wait 0,1,2,3,4,5 contention slots

- after collision
→ report
error to
the upper
layers

MACAW: MACA for Wireless network

Fine tunes MACA to improve performance:

- introduces an **ACK frame** to acknowledge a **successful data frame**
- added **Carrier Sensing** to keep a station from transmitting **RTS** when a nearby station is also transmitting an **RTS** to the same destination
- **exponential backoff** is run for each separate pair **source/destination** and not for the single station
- mechanisms to exchange information among stations and recognize temporary congestion problems
- **CSMA/CA** used in IEEE 802.11 is based on **MACAW**

IEEE 802.11 FAMILY OF STANDARDS

- IEEE 802.11 (Legacy mode, 1st version)
 - released in '97, clarified in '99
 - rarely used today → evolving in a/b/g/n (these most used today)
 - **data rate**: started with 1.2 Mbps *data rate* implemented via:
 - infrared signals (IR)
 - **radio frequencies** in the 2.4 Ghz band (Industrial Scientific Medical Freq.)
 - many **degrees** of freedom:
 - **interoperability** among different products was challenging
 - became popular with version 802.11b
- IEEE 802.11a
 - '99
 - **operating (radio) frequency**: 5 Ghz band
 - **throughput**(typ): 23 Mbps
 - **data rate**(max): 54 Mbps
 - lot of difference, exploit max transmission capacity of technologies
- IEEE 802.11.b
 - '99
 - **operating (radio) frequency**: 2,4 Ghz band (ISM band)

- problem: in some environments (domestic) interference (microwave ovens, cordless tel)
 - **throughput** (typ): 4.3 Mbps (<A)
 - **data rate** (max): 11 Mbps
- IEEE 802.11g
 - 2003
 - **operating (radio) frequency**: 2,4 Ghz band (SM)
 - **throughput**(typ): 19 Mbps and more (>B)
 - **data rate**(max): 54 Mbps
- IEEE 802.11n
 - 2009
 - **operating (radio) frequency**: 2,4 Ghz and 5 Ghz band
 - **throughput**(typ): 74 Mbps and more (>G)
 - **data rate**(max): 248 Mbps
 - supports MIMO technologies:
 - a node can use multiple antennas at the transmitter/reciver
- **Wi-Fi 5** - IEEE 802.11ac
 - 2013
 - **operating (radio) frequency**: 2,4 Ghz and 5 Ghz band
 - **data rate** (max):
 - at **5 Ghz**: 1.3 Gbps
 - at **2.4 Ghz**: 450 Gbps
- **Wi-Fi 6** IEEE 802.11ax
 - 2019
 - **data rate**: reaches up 10 Gbps
 - improvements
 - power consumption and security
 -
- IEEE 802.11 Wireless LAN

IEEE 802.11 standard	Year	Max data rate	Range	Frequency
802.11b	1999	11 Mbps	30 m	2.4 Ghz
802.11g	2003	54 Mbps	30m	2.4 Ghz
802.11n (WiFi 4)	2009	600	70m	2.4, 5 Ghz
802.11ac (WiFi 5)	2013	3.47Gpbs	70m	5 Ghz
802.11ax (WiFi 6)	2020 (exp.)	14 Gbps	70m	2.4, 5 Ghz
802.11af	2014	35 – 560 Mbps	1 Km	unused TV bands (54-790 MHz)
802.11ah	2017	347Mbps	1 Km	900 Mhz

- all use CSMA/CA for multiple access, and have base-station and ad-hoc network versions

• 802.11 af-ah:

- sending data for measurement
- **low frequency → higher range** (higher coverage)
 - cool for IoT: sensor in fields need coverage

• 802.11 stack

- all these protocols are at L2 (physical, frame shift)
 - different releases → different implementations
 - MAC and logically controls
 -

network

MAC:
CSMA/CA

DATA LINK

802.11 Architecture

- a group of stations (AP) has **coordination function**
 - to transmit data send/rec

802.11 family

physical

- 1) • AC can be used in between nodes/other AC (link)/group of stations **via fixed infrastructure**

→ can used!

1) infrastructure wireless net
2) No infra wireless net

- if using AP, a station communicates with another
 - challenging all traffic thru a **centralised AP**

- may or not use a BS

- 2) • supports **ad hoc network** (no infrastructure):

- a group of stations that are under the **direct control of** a single channel **coordination function**, **without** the aid of an **infrastructure network**
- **no centralised role of AP**
- distributed among nodes belong ad hoc network

MAC SUBLAYER of 802.11

2 Modes of operations

1) DCF: Distributed Coordination Function \Rightarrow MUST BE PRESENT in ALL the DEVICES

- all stations have to run CSMA/CA to see if they can access the channel:
 - sense channel
 - if free \rightarrow send RTS
 - if receive CTS \Rightarrow start transmit
- } all stations

2) PCF: Point Coordination Function \Rightarrow OPTIONAL IMPLEMENT PCF

- MAC is coordinated by access point: ABOVE DCF

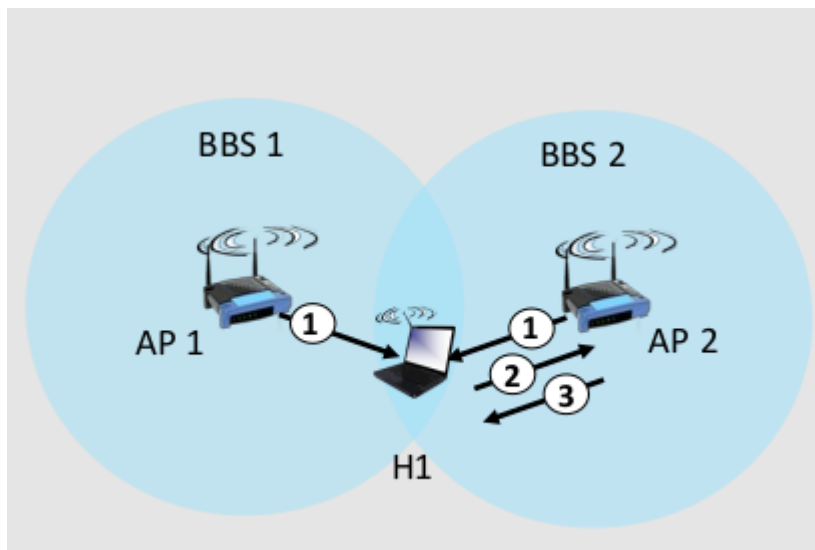
\Rightarrow AP may reserve part of the time and tell to other stations that that time slot is reserved to let 1 station to transmit

- used in video streaming transmission this will ~~not~~ carrier sense & send RTS \leftarrow CTS

- a station can communicate directly with another without channelling all traffic thru AP
- spectrum divided into channels at different frequencies
 - → range of frequency divided in # of channels (usually 11)
 - → AP admin choose frequency for AP
 - → AP communicate thru channel using that frequency with the device
 - interference possible: channel can be the same as that choose by neighboring AP
- a node join to a network must associate with AP
 - scans channels
 - listen for beacon frames
 - containing SSID(net name) && (BSSID) MAC of the AP
 - host selects AP to associate with
 - they can perform auth
 - typically they run DHCP to get IO address in P's subnet

Scanning

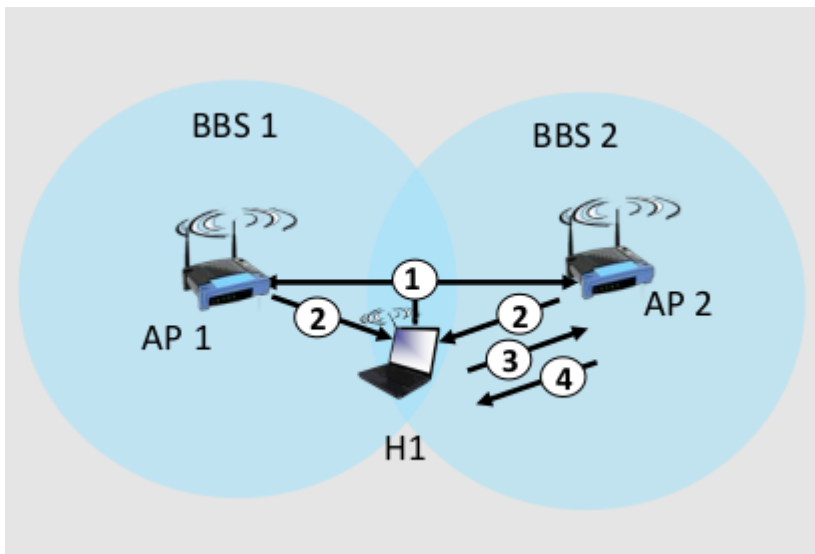
- host want to join the network
 - must discover infos abot AC to join to the network
- passive scanning



1. beacon frames sent from APs
 - AC periodically sends beacon frames
2. association request frame sent:
 - H1 to select AC (based on BSSID)
 - choosing the AP according higher signal strength (depends on vendor implementation really)

3. **association response** frame sent from AP to H1

- **active scanning**



- a **node arrives** in network → **send beacon** → want the **reply to an AP**

1. **Probe request** frame broadcast from H1

2. **probe response** frames sent from APs

3. **association request** frame sent:

- H1 to selected AP

4. **Association response** frame sent:

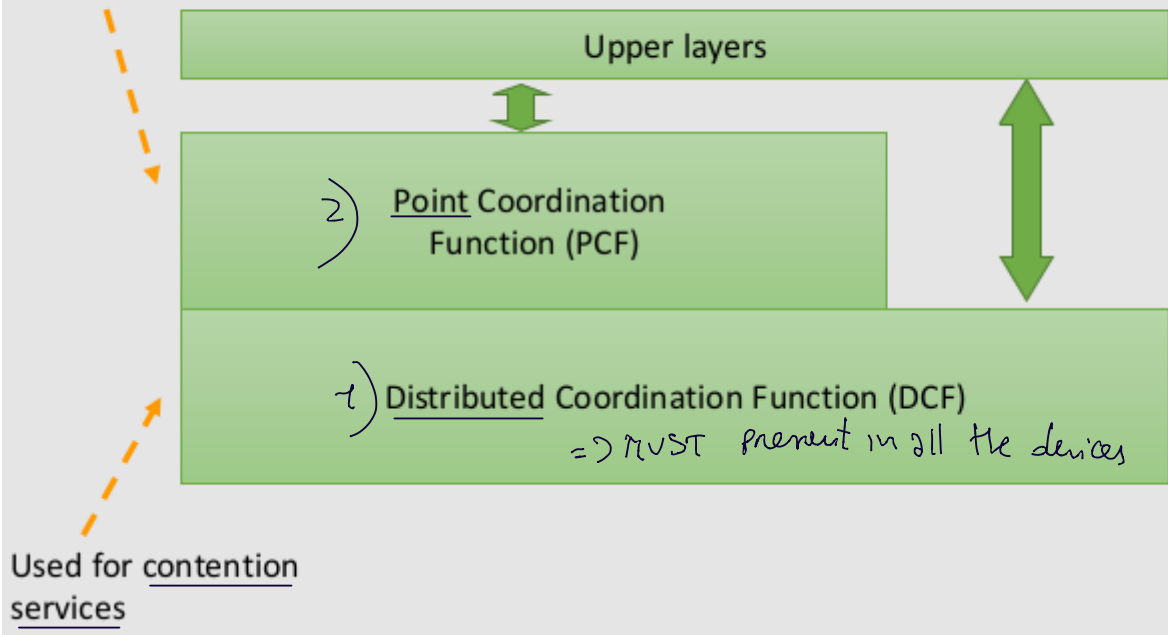
1. from selected AP to H1

- **association request** and **association response** **needed** cuz can have **multiple AP**, but not aware if AP replies

MAC Sublayer

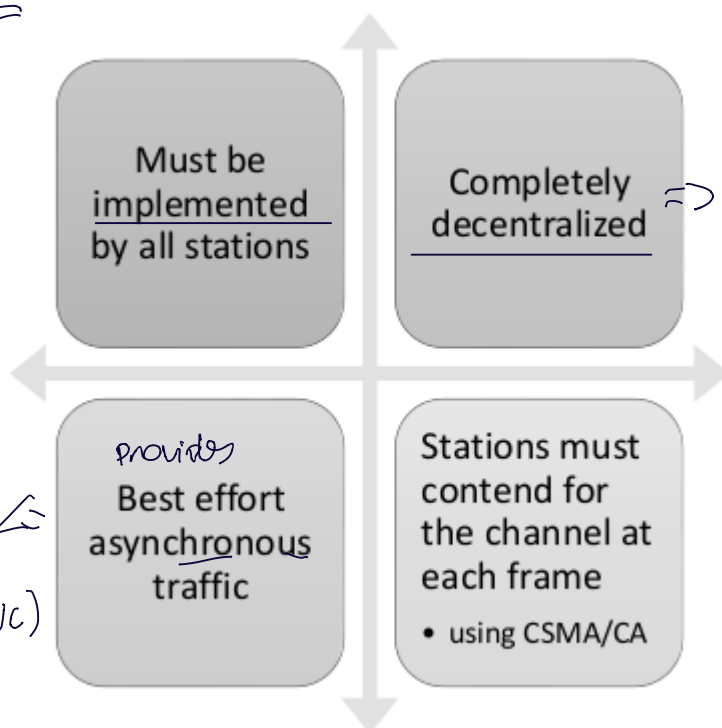
Used for contention-free services and based on DCF

MAC SUBLAYER



- **DCF:** Distributed Coordination Function, here implements MACA → distributed access to shared medium
 - completely decentred
 - thought for best effort asynchronous traffic
 - **must be implemented by all stations**

DCF



⇒ no central point that tells to BS when to talk

traffic async
 ↓
 can transmit when you want
 (no time sync)

best effort

↓
 you don't know if you will be able to transmit
 send RTS & will not have back a CTS

- Carrier sensing is performed at two levels:

- physical CS:** *Carrier Sense implemented in Physical: sense to channel before transmitting*
 - informal: before transmission if another station is transmitting → detect incoming signal → if channel is free → transmit
 - formal:
 - checking the frequency to determine whether the medium is in use or not
 - physical carrier sense to detect incoming signal
 - detects any activity in the channel due to other sources
- virtual CS:** *: if you receive a CTS → you set up virtual carrier sense (as if channel is busy for you)*
 - performed sending duration information in the header of an RTS, CTS and data frame
 - keep channel virtually busy up to the end of frame transmission
 - a channel is marked busy if either the physical or the virtual CS indicate busy
- a channel is marked busy if either the physical or the virtual CS indicate busy

- always **about coordination function:** *to make things work*

- priority access to the medium is controlled thru the use of **interframe space (IFS) time interval:** *min time need to wait to send new frame (time between 2 frames)* *↳ when transmit a frame, before transmit a new frame*
 - formal IFS: mandatory periods of idle time on the transmission medium *⇒ wait a fixed time interval*
- Three IFS specified by standard (always talking about IEEE 802.11)
 - Short IFS (SIFS)
 - Priority Coordination Function IFS (PIFS)
 - Distributed Coordination Function (DIFS)
 - $SIFS < PIFS < DIFS$
 - stations only require to wait SIFS have the higher priority:
 - node with higher priority wait less than others with smaller one

- PCF: Point Coordination Function: access medium between nodes to organise communication service

- contention free
- uses BS to control all activity in its cell
- though for delay-sensitive traffic
- AP polls station for transmissions
- based on DCF

*Turns a frame
→ Turns the 2nd*

*wait time before transmit
⇒ if you want to give priority to another station, other station will wait a smaller time interval
⇒ will transmit first
⇒ PCF*

IFS also to give priority

fixed if don't want overlap

- DCF || PCF can be active at same time in same cell

- MANCA ULTIMA SLIDE