

Wireless networks & IEEE 802.11

Mobile and Cyber Physical Systems

Overview



Wireless networks basics



Wireless networks architectures



IEEE 802.11 (Wi-Fi) a/b/g/n/...

Why Wireless Networks?

1. **Cables in computers** to:
 - communicate → Transfer data between nodes
 - provide power → between devices
 2. **Cyber-physical systems embed computers in (any kind of) physical objects**
 - Cyber-physical systems integrate sensing, computation, control and networking into physical objects and infrastructure, connecting them to the Internet and to each other.
 - can't wire everything...
- They are
everywhere
free placement
mobility
Exposed into
the wild
↓
need
wires*

Hence:

- wireless to replace cables in communications...
- ... & batteries to replace cables in power supply

Wireless Networks

WIRELESS NETWORKS

- Networks of **hosts** connected by **wireless links**



- **Hosts**: end-system devices that run applications

- often mobile, but not necessarily
- **battery-powered** (typically)
- e.g. smartphone, tablet, sensor, home appliance, vehicle, etc...

- Two modes of operations:

- 1) • **Infrastructure**

BS

- with **base station(s)** or

wired AC_{SS}P

- with **wired access points**

→ we have nodes that must be coordinated

→ core network provides: AUTHENTICATION
- BILLING

- 2) • **ad hoc networking**

- **no centralized coordinators**

linked wireless nodes (adhoc)
→ coordinate shared network

Elements of a Wireless network (I)

Sure → mechanism

Wireless hosts

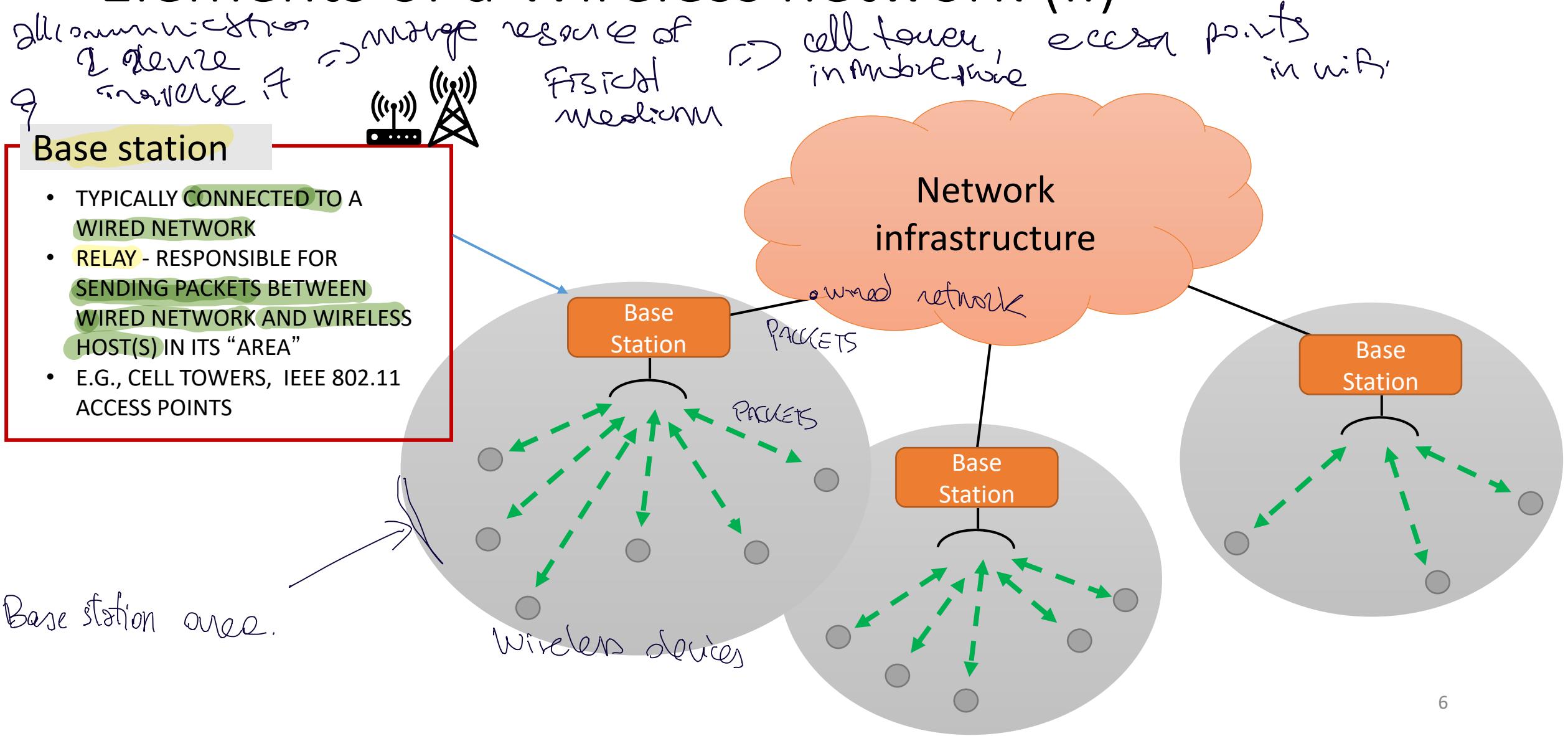
- LAPTOPS, SMARTPHONES
 - RUN APPLICATIONS
 - MAY BE STATIONARY (NON-MOBILE) OR MOBILE
 - WIRELESS DOES NOT
ALWAYS MEAN MOBILITY

- device went
to connect
to internet
 \Rightarrow traverses
internet (wind information)
 \Rightarrow like fire - carbon - atoms



BASE STATION:
cover geographic
Areas

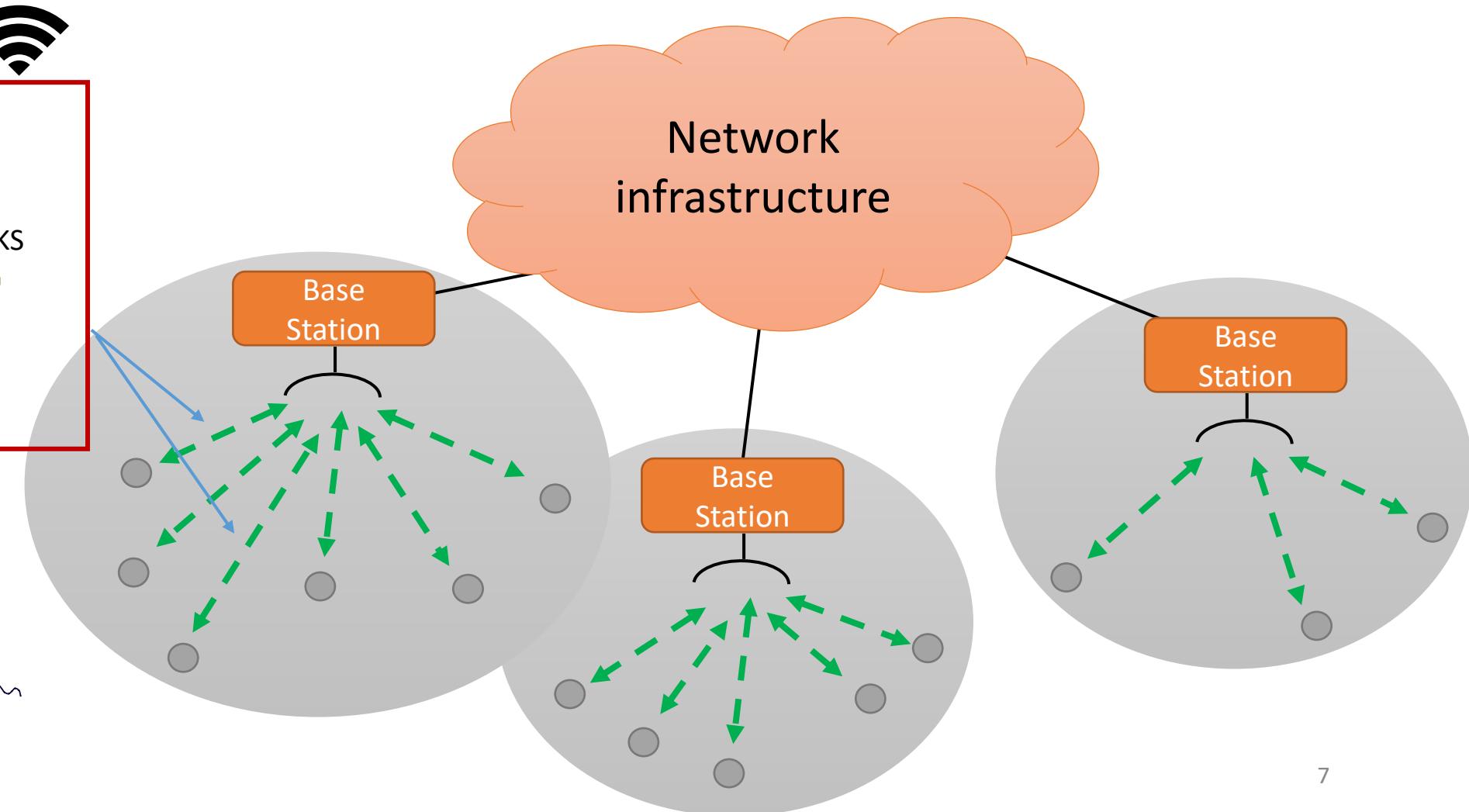
Elements of a Wireless network (II)



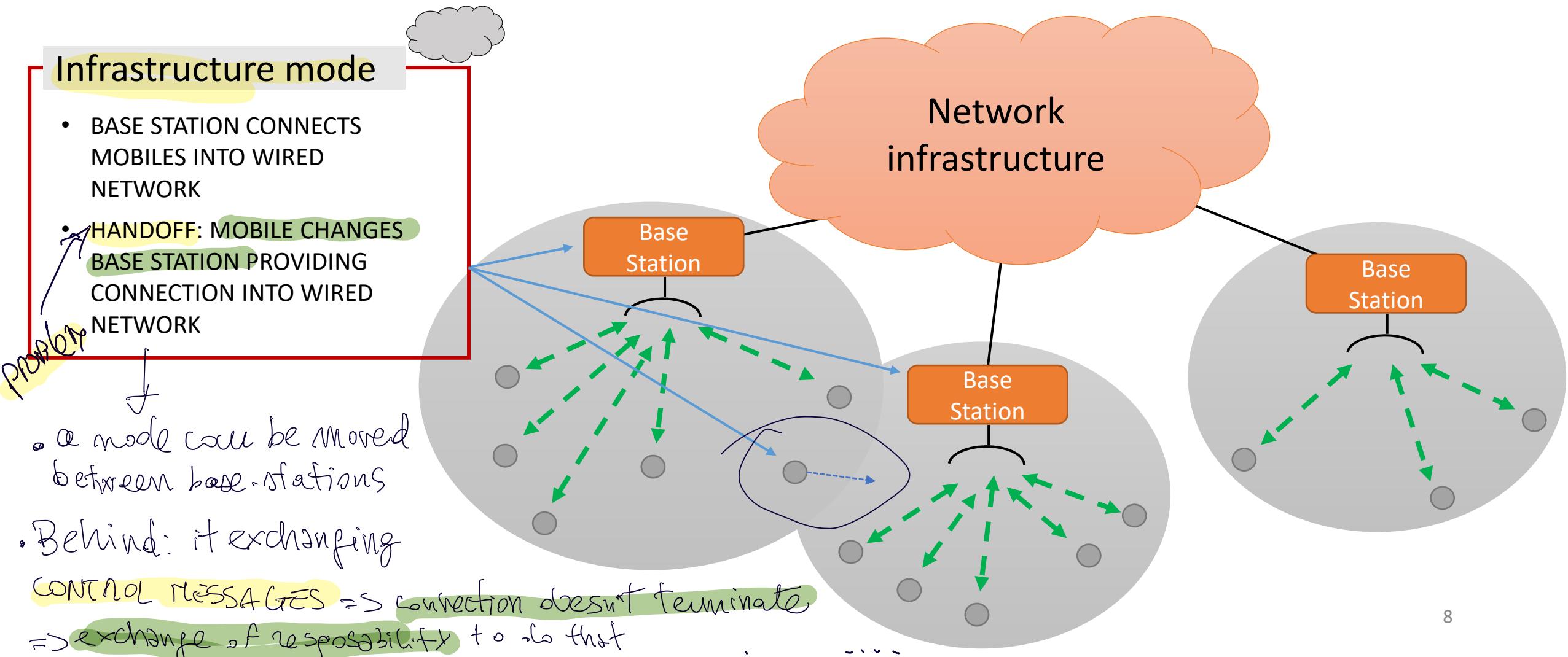
Elements of a Wireless network (III)

Wireless links

- CONNECT MOBILE(S) TO BASE STATION
- ALSO USED AS BACKBONE LINKS
- MULTIPLE ACCESS PROTOCOL COORDINATES LINK ACCESS
- VARIOUS DATA RATES, TRANSMISSION RANGE,...

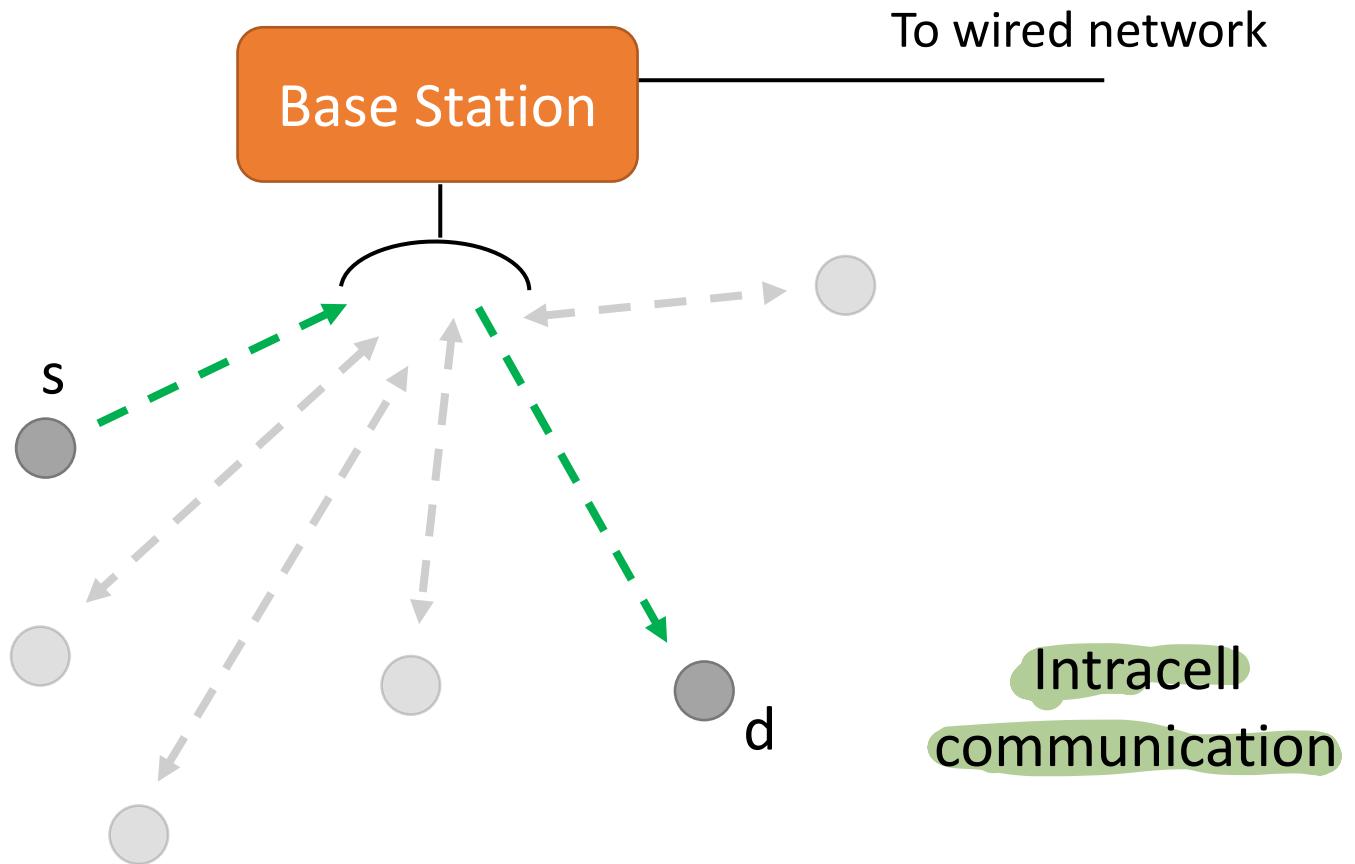


Elements of a Wireless network (IV)



↳ (example: Cellphones move on same/different operators/providers → AUTHENTICATION
BILLING)

Elements of a Wireless network (V)

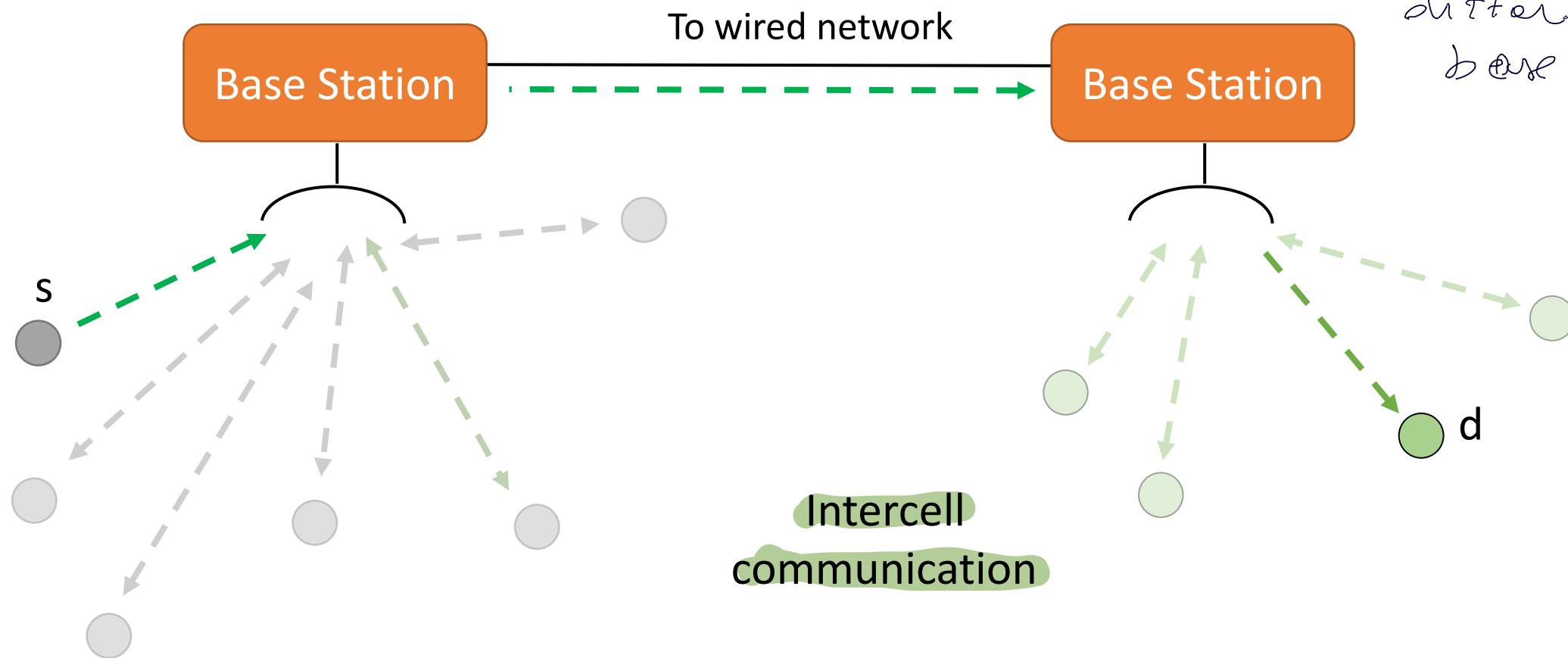


- the two host (end-devices) communicate on the same base station

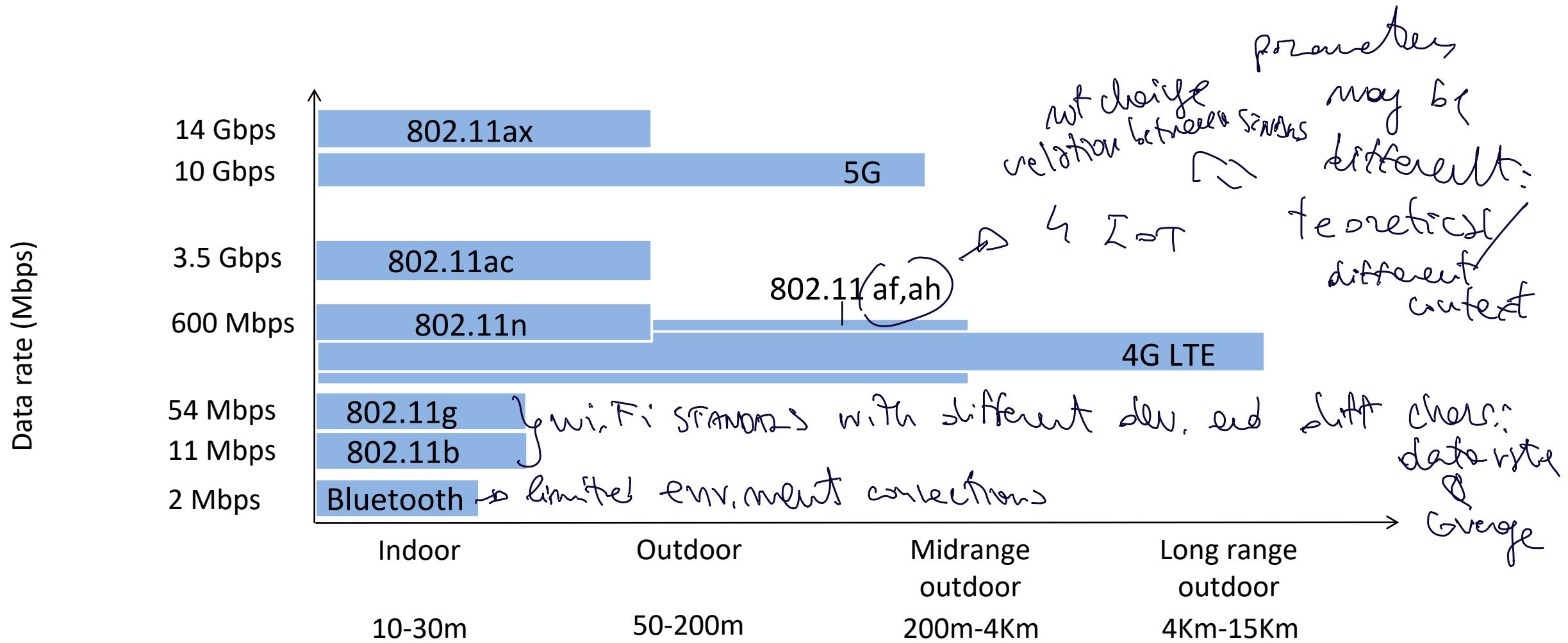
Intracell
communication

Elements of a Wireless network (VI)

the two hosts
communicates
between 2
different
base stations



Characteristics of selected wireless links



Wireless network taxonomy

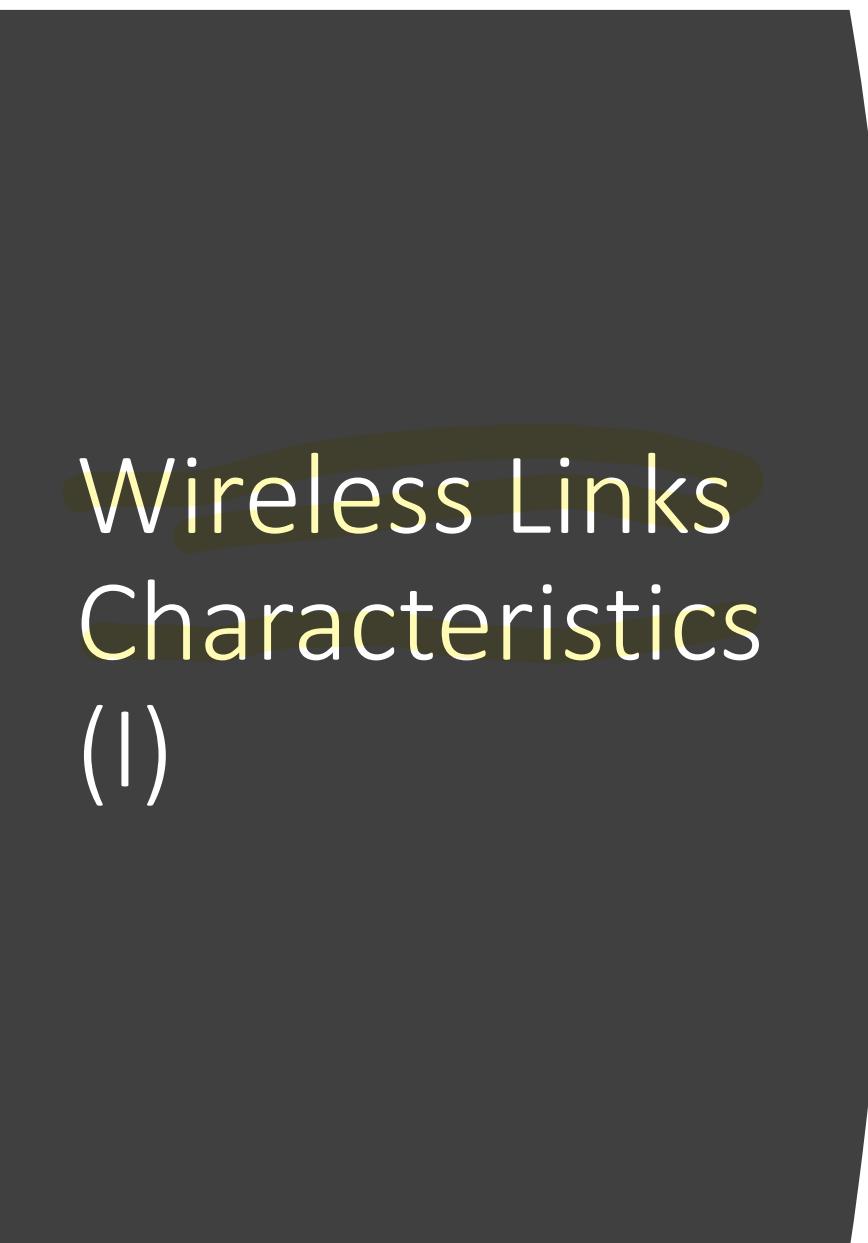
Distinguish with

① infrastructure (e.g., APs)	② Single hop: host connects to base station (WiFi, WiMAX, cellular 3G,4G,5G) which connects to larger Internet	③ Multiple hops: host may have to relay through several wireless nodes to connect to larger Internet: MESH networks
② no infrastructure	<ul style="list-style-type: none">• no base station, not necessarily connection to larger Internet (e.g. Bluetooth)	<ul style="list-style-type: none">• no base station, no connection to larger Internet. May have to relay on other nodes to reach a given wireless node (ZigBee, ad hoc, VANET)

different
sense of
a hop

Vehicular networks
allows end-to-end communication

Wireless Links Characteristics (I)



that coordinates each other \Rightarrow need a gateway to extend network coverage

important differences from wired link

- Electromagnetic waves in the air
 - \Rightarrow Problem: decreases with distance
- decreased signal strength: radio signal attenuates as it propagates through matter (path loss)
 - OBSTACLES: electromagnetic can be blocked
- interference from other sources: standardized wireless network frequencies (e.g., 2.4 GHz) shared by other wireless devices (e.g., mobile phones). Engines, appliances (microwave ovens...), ... may interfere as well
 - interference \Rightarrow mobile phone \Rightarrow microwave have same frequency
- multipath propagation: radio signal reflects off objects or ground, arriving at destination at slightly different times
 - reciever detect signal \Rightarrow sum of signals that are in the environment [noise or sources was seen first]

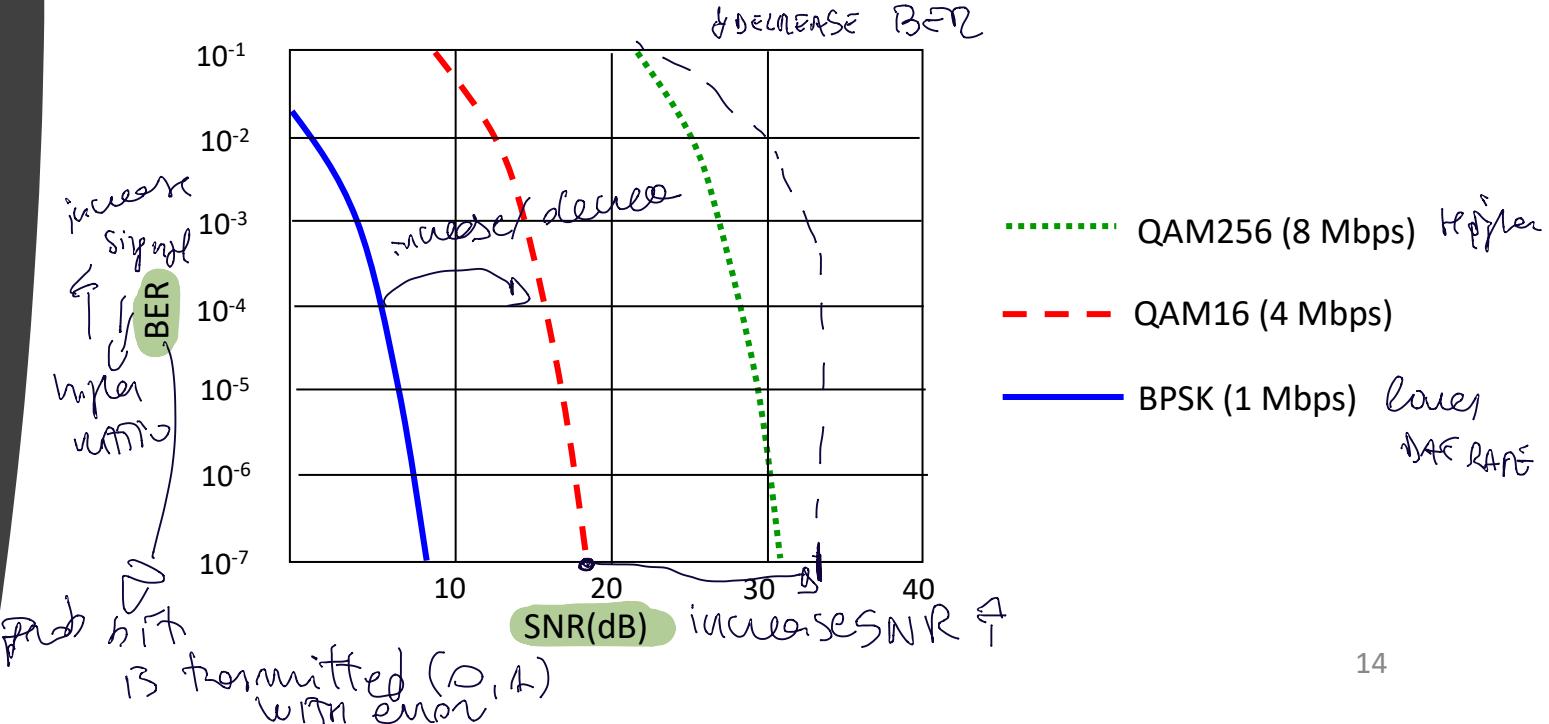
.... make communication across wireless links much more "difficult" (even a point to point)

 - Electromagnetic waves propagate in the environment \Rightarrow waves are reflected on objects
 - No specific signal, but \sum of different signals reflected
 - \Rightarrow reconstruct each signal by the \sum

Wireless Links Characteristics (II)

→ Evaluate quality of the signal $\Rightarrow \text{SNR} = \frac{\text{SIGNAL power}}{\text{noise power}} \text{ SNR}$

- **SNR: signal-to-noise ratio**
 - larger SNR – easier to extract signal from noise (a “good thing”)
 - SNR in dB = $20 \log(\text{signal/noise})$
 - **SNR versus BER tradeoffs** (modulation: how important info in electric waves)
 - bit error rate (BER): probability that a transmitted bit is received in error at the receiver
 - given physical layer: increase power -> increase SNR->**decrease BER** using
 - given SNR: choose physical layer that meets BER requirement, giving highest throughput
 - SNR may change with mobility: dynamically adapt physical layer (modulation technique, rate)

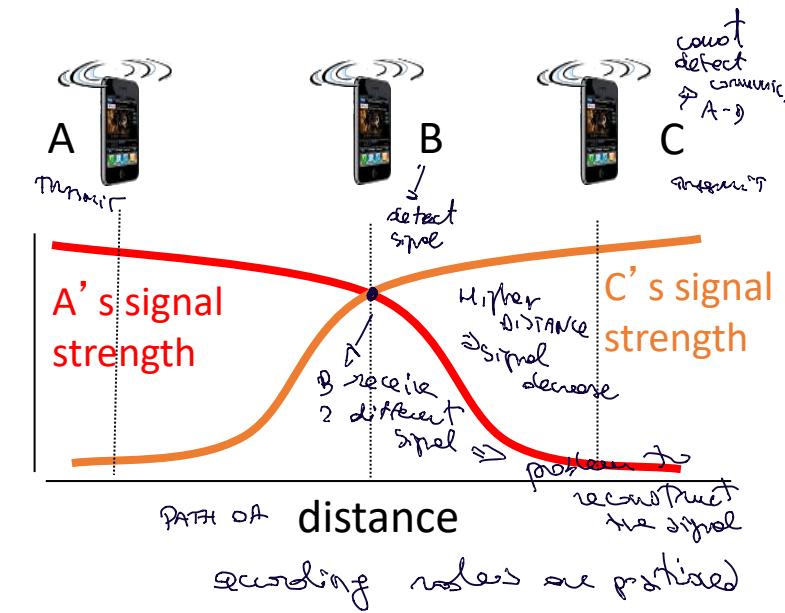
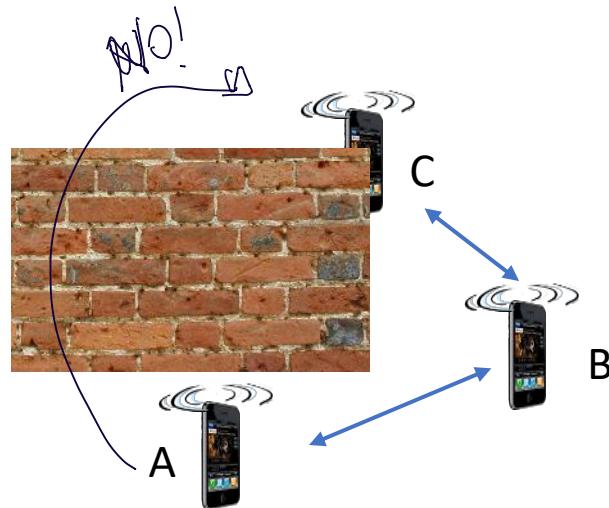


Wireless Links Characteristics (III)

Different: most in env \Rightarrow situations is different while nodes send stuff

Signal attenuation or obstacles limit transmission ranges

SIGNAL STRENGTH: potencia del se˜nale



- B, A hear each other
- B, C hear each other
- A, C can not hear each other

Wireless networks challenges

WIRELESS NETWORKS PROBLEMS

1) Limited knowledge

- a terminal cannot hear all the others
DISPAPARE
- hidden/exposed terminal problems

Join / Disappears

2) Mobility/Failure of terminals

- terminals move in the range of different BS
Join / EXIT
- terminals move away from each other

: base station

3) Limited terminals

- battery life, memory, processing and transmission range

limited battery power

4) Privacy

- eavesdropping of ongoing communications

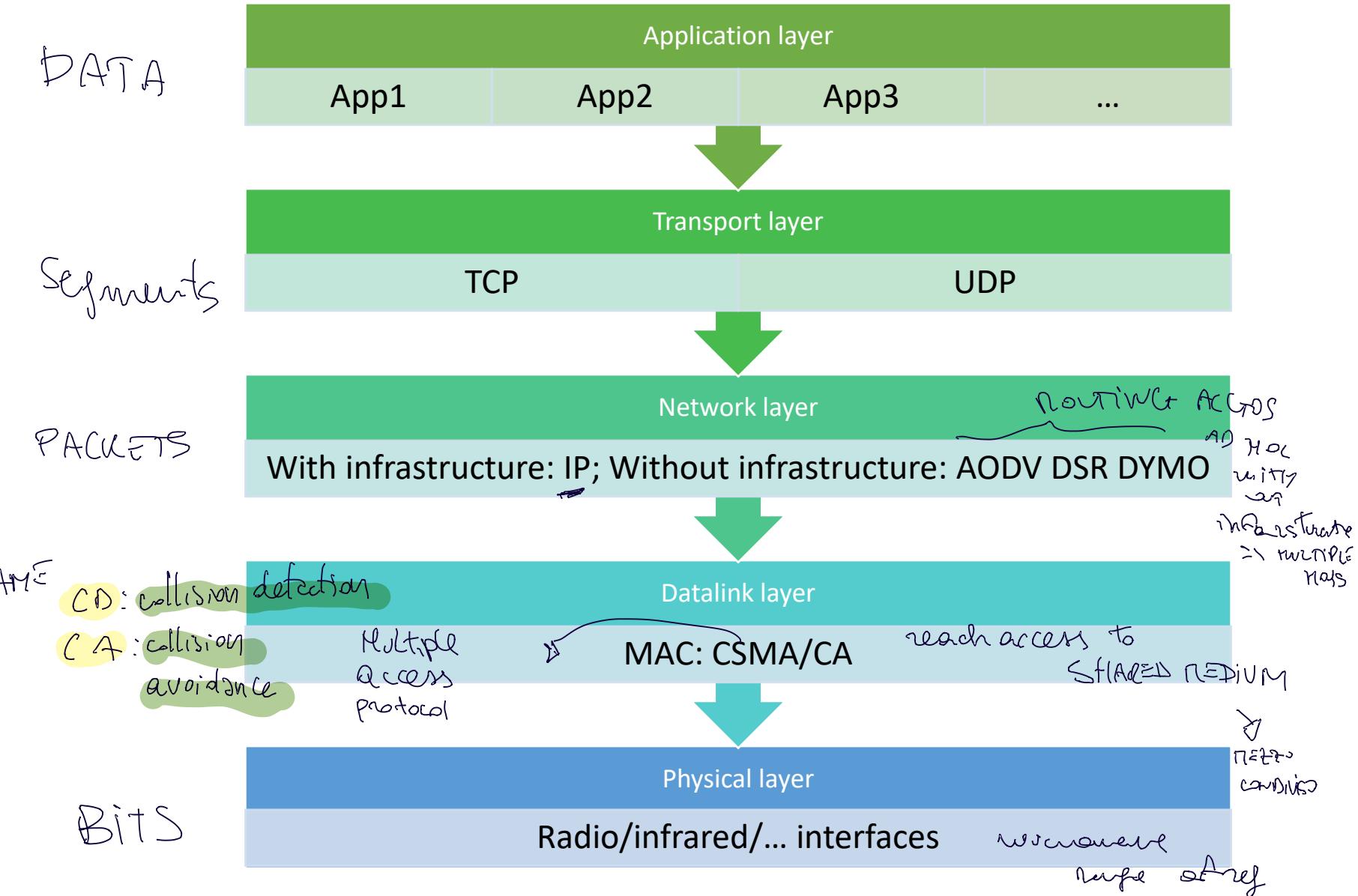
↳ intercettazioni

Wireless networks: required mechanisms

WIRELESS NETWORKS REQUIREMENTS

- 1) • Access to a shared wireless channel
 - CSMA/CD cannot be used...
MULTIPLE ACCESS PROTOCOL
- 2) • Hand-off (Networks with infrastructure)
TRAPASSO
 - moving a terminal into the range of a different BS
- 2) • Routing (multi-hop ad hoc networks)
 - finding a path from source to destination in multi hop networks
 - dealing with arbitrary changes in neighborhood

Wireless networks protocol stack



RECAP: MAC protocols for wired networks

MAC PROTOCOL = wired networks
↳ frames

- Basic assumptions:
 - a single channel is available for all communications
 - all stations can transmit on it and receive DATA from it
 - if frames are sent simultaneously on the channel the resulting signal is garbled (a collision)
 - all stations can detect collisions
- Different protocols
 - ALOHA, slotted ALOHA, CSMA, CSMA/CD, ...

↑
Collision detection

Carrier Sense Multiple Accesses with Collision Detection (CSMA/CD)

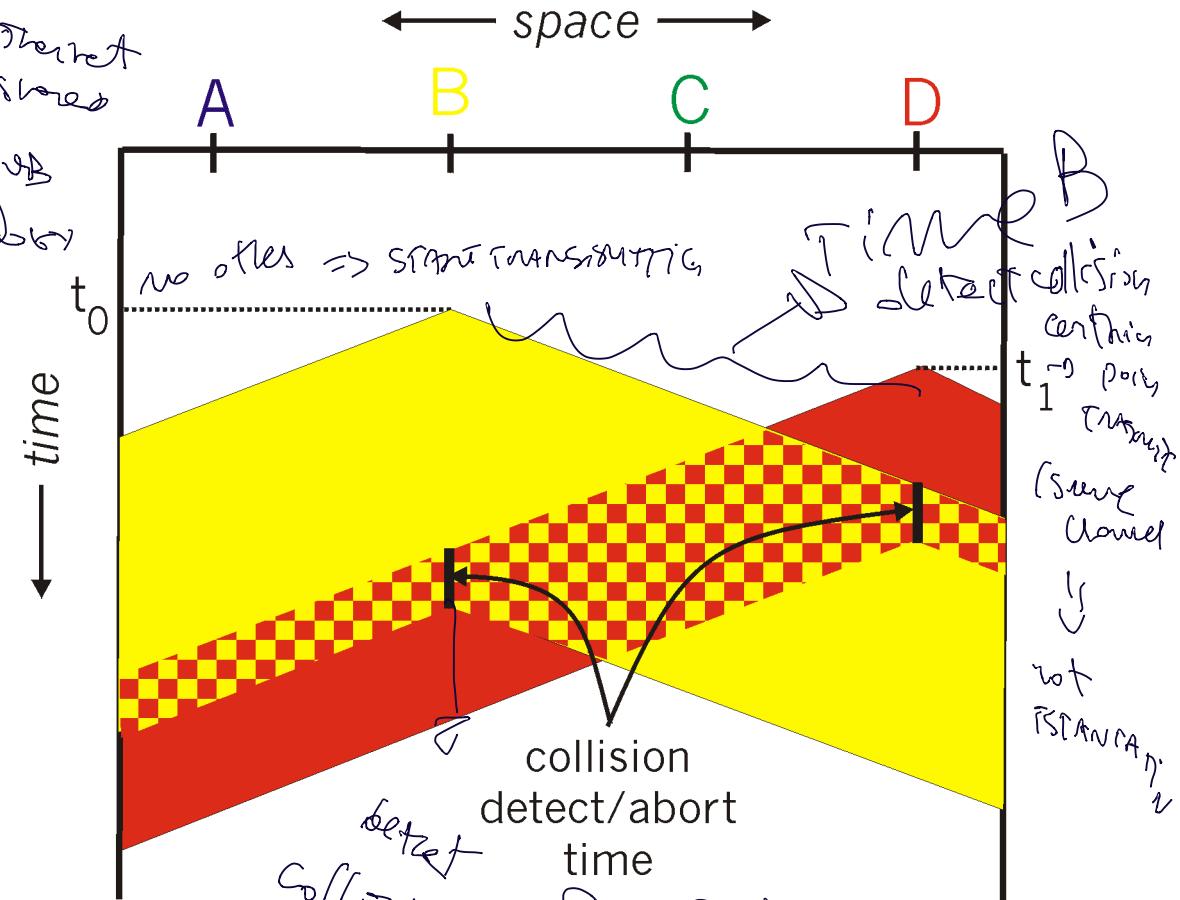
- Basic idea of CSMA:

- When a station has a frame to send it listens to the channel to see if anyone else is transmitting
- if the channel is busy, the station waits until it becomes idle
- when channel is idle, the station transmits the frame
- if a collision occurs the station waits a random amount of time and repeats the procedure.

if there is 2 transmission,
wait channel is free

SNR
For Ethernet
with stored access
to revB
(SIPAC) probably

if short
wait
enough time



20
=> also defines duration

CSMA/CD (2)

prob to have another collision
interval → after retransmit
(like 2 person one call, if)

- CSMA with *Collision Detection*

- a station aborts its transmission as soon as it detects a collision
 - if two stations sense the channel idle simultaneously and start transmitting, they quickly abort the frame as soon as collision is detected
- it is widely used on LANs in MAC sub-layer
- IEEE 802.3 Ethernet (look see network)

ALGO:

CSMA/CD (3)

How A node is using a channel

⇒ transmission of a frame

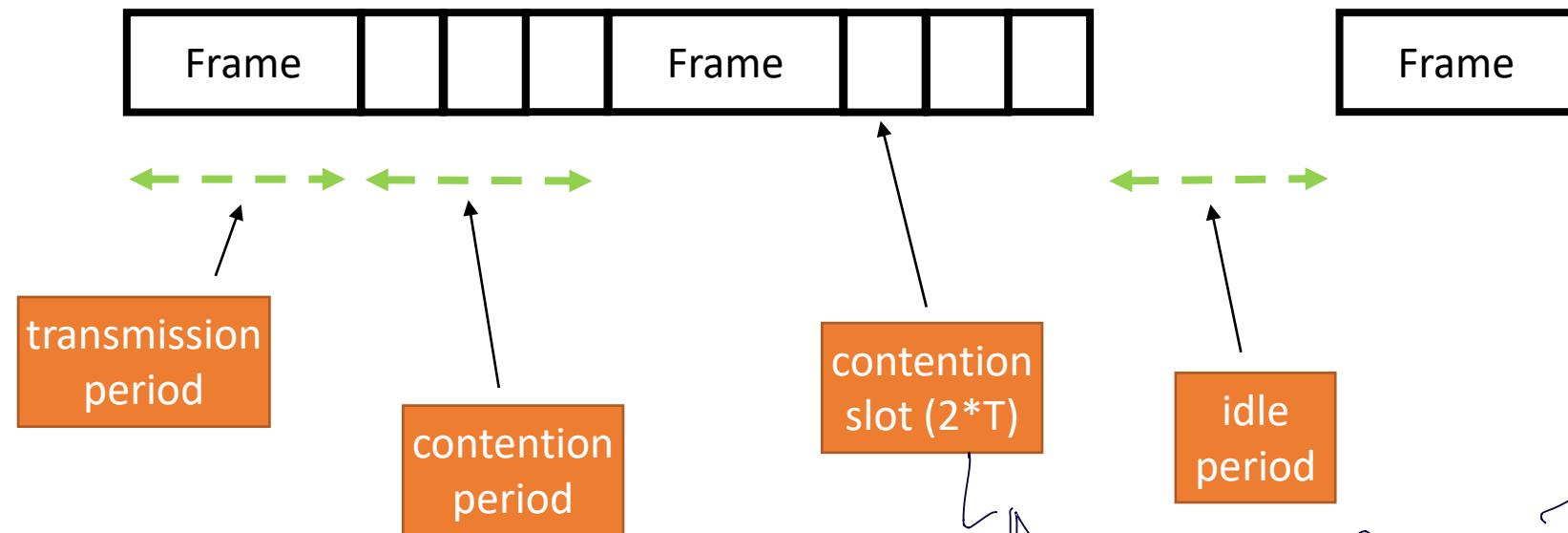
⇒ maybe collision ⇒ SITUATION WHICH

→ PIN' contention

+ minimum time

CSMA/CD behavior

- T is the time required to reach the farthest station
- It takes minimum of RTT time ($2*T$) to detect collision.



RTT:

tempo fra
invio sepele
e ricezione
della conferma

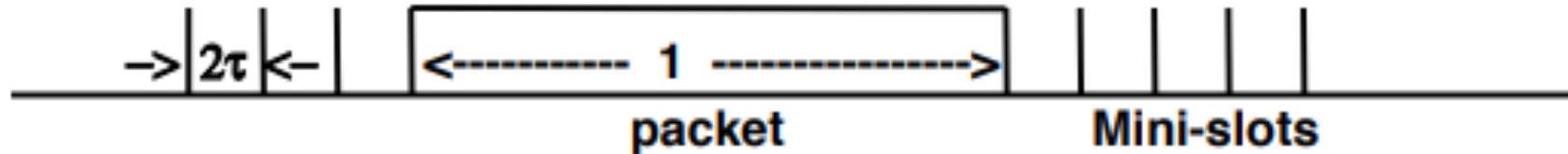
di invio della frame

~ V (RTT)
DIM =
A & B 22

minimum
idle
time

CSMA/CD (4)

- Consider a slotted system with “mini-slots” of duration 2τ
- If a node starts transmission at the beginning of a mini-slot, by the end of the mini-slot either
 - No collision occurred and the rest of the transmission will be uninterrupted
 - A collision occurred, but by the end of the mini-slot the channel would be idle again
- Hence a collision at most affects one mini-slot



Binary Exponential Backoff

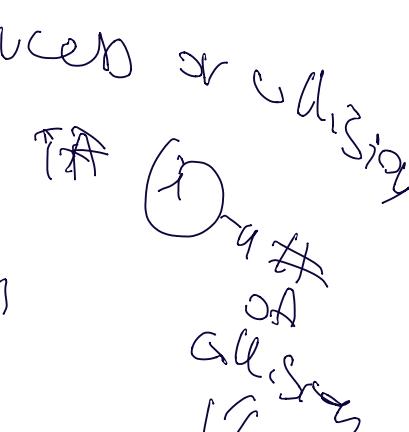
anetvane

as with wired networks

ALGO

+ to define time station wait to transmit again

- Time after a collision is divided in contention slots
 - length of a contention slot is equal to the worst case round propagation time ($2T$ if T is the time to reach the most distant station)
- After the first collision
 - each station waits 0 or 1 slot before trying again
- After collision i
 - chooses x at random in $[0, 2^i - 1]$
 - skips x slots before retrying
- After 10 collisions:
 - the randomization interval is frozen at 0..1023
- After 16 collisions
 - failure is reported back to upper levels



To exp: more collisions
have time
may increase

eventually slots

between
a wide range

WIFI don't work with wireless



Wireless networks: MAC

CSMA/CD detects interference while transmitting!

node see or detection & sita
re are is receiving

Hidden terminal problem
 \Rightarrow exposed terminal problem

Exposed terminal problem

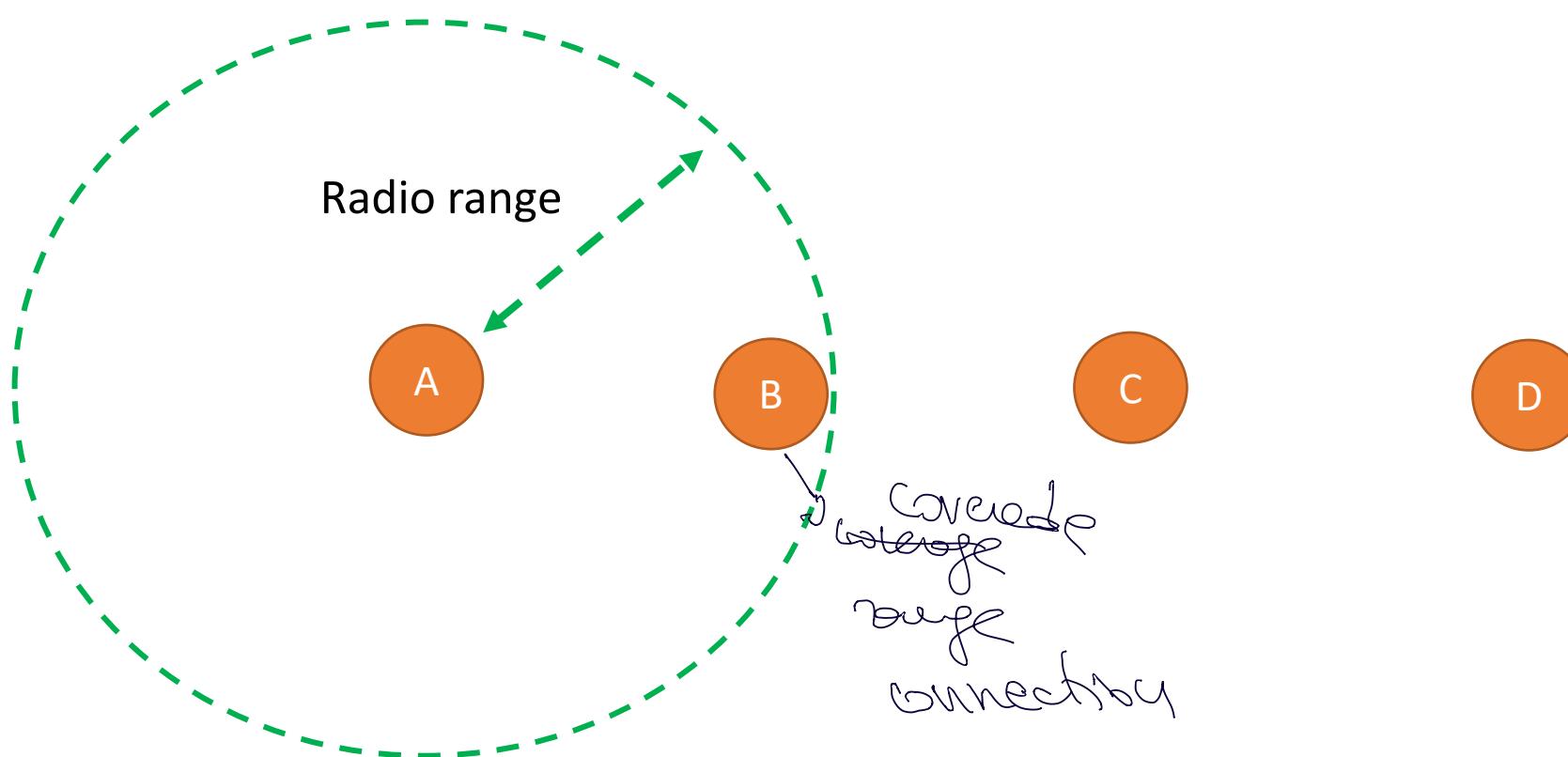
multiple antennas $\xrightarrow{\text{antennas}} \text{very low energy}$
cannot do that in wireless

in wired networks it's OK, but not in wireless!

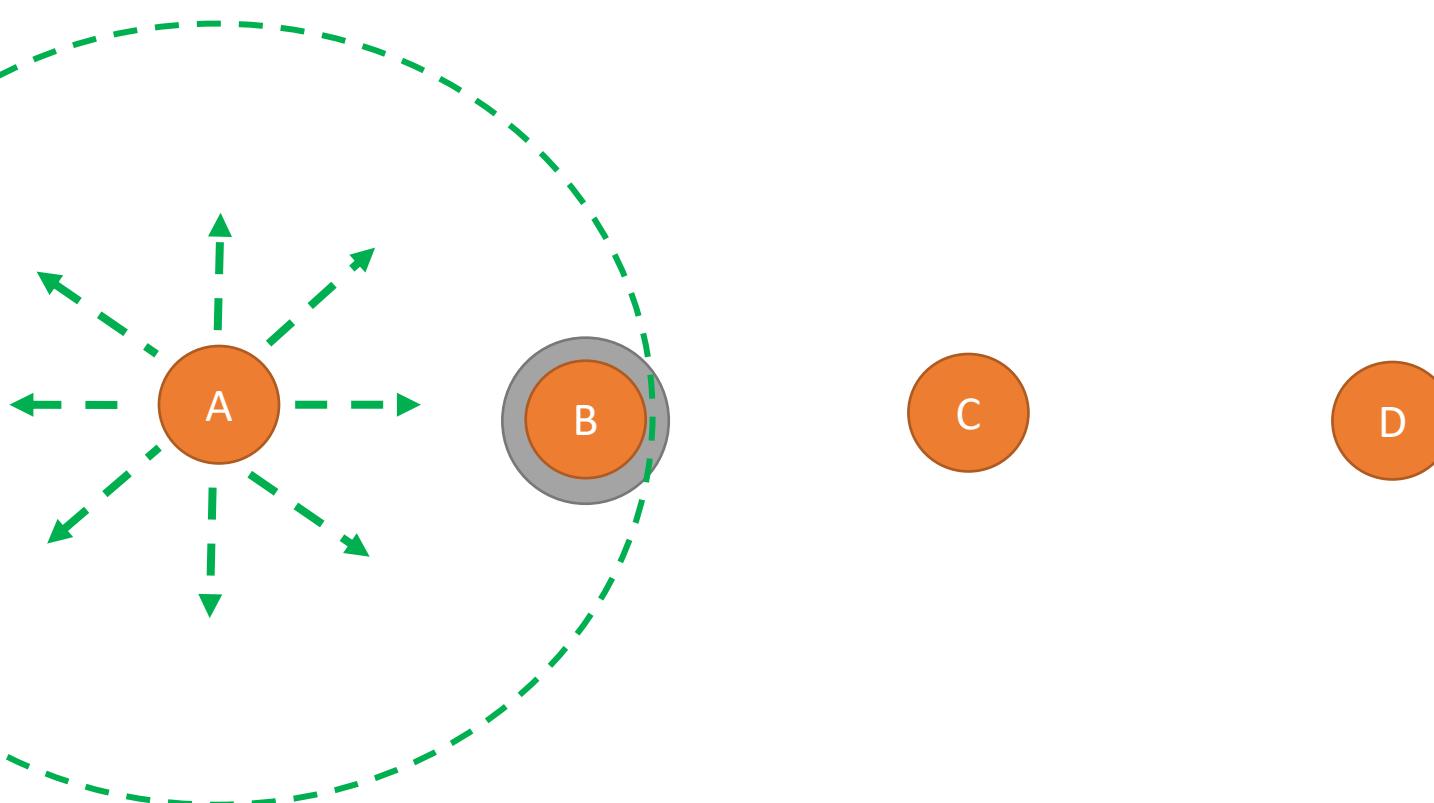
what matters is
◦ interference at the receiver *not at the sender*

what matters is
interference at the receiver *not at the sender*

The hidden terminal problem

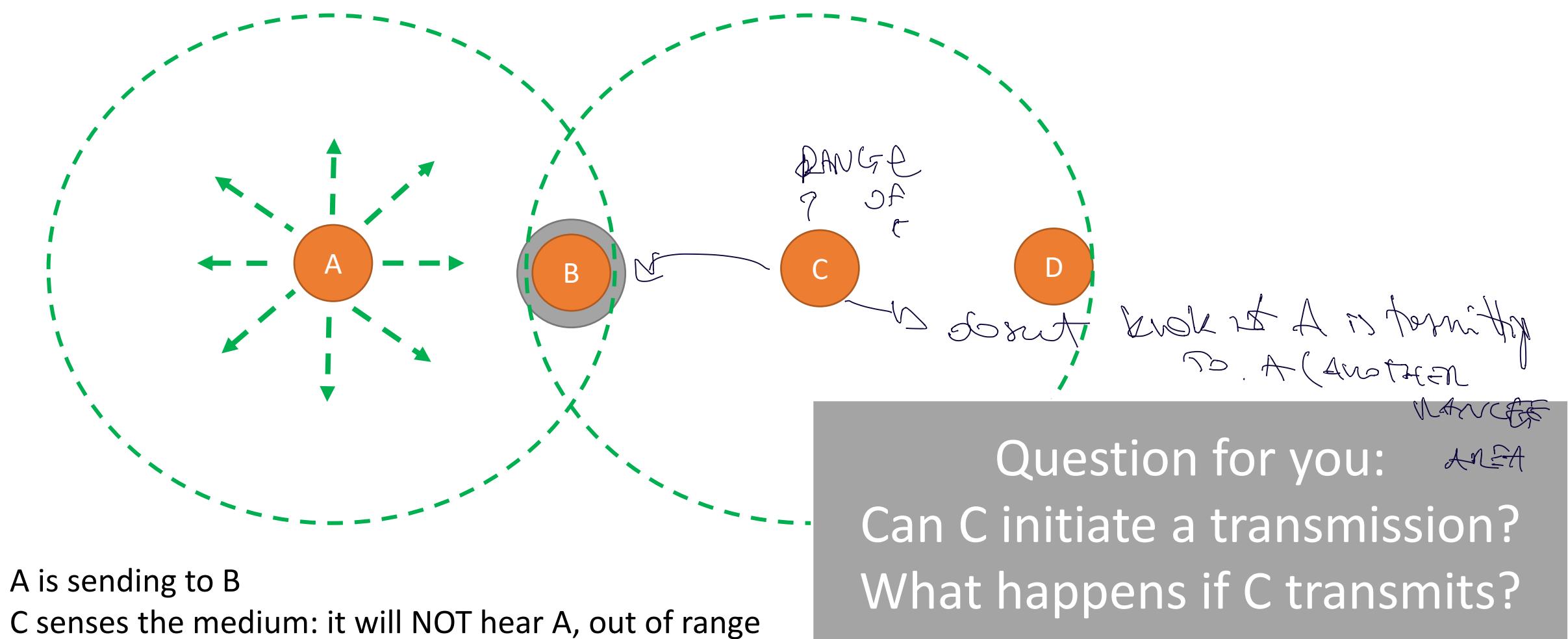


The hidden terminal problem (2)

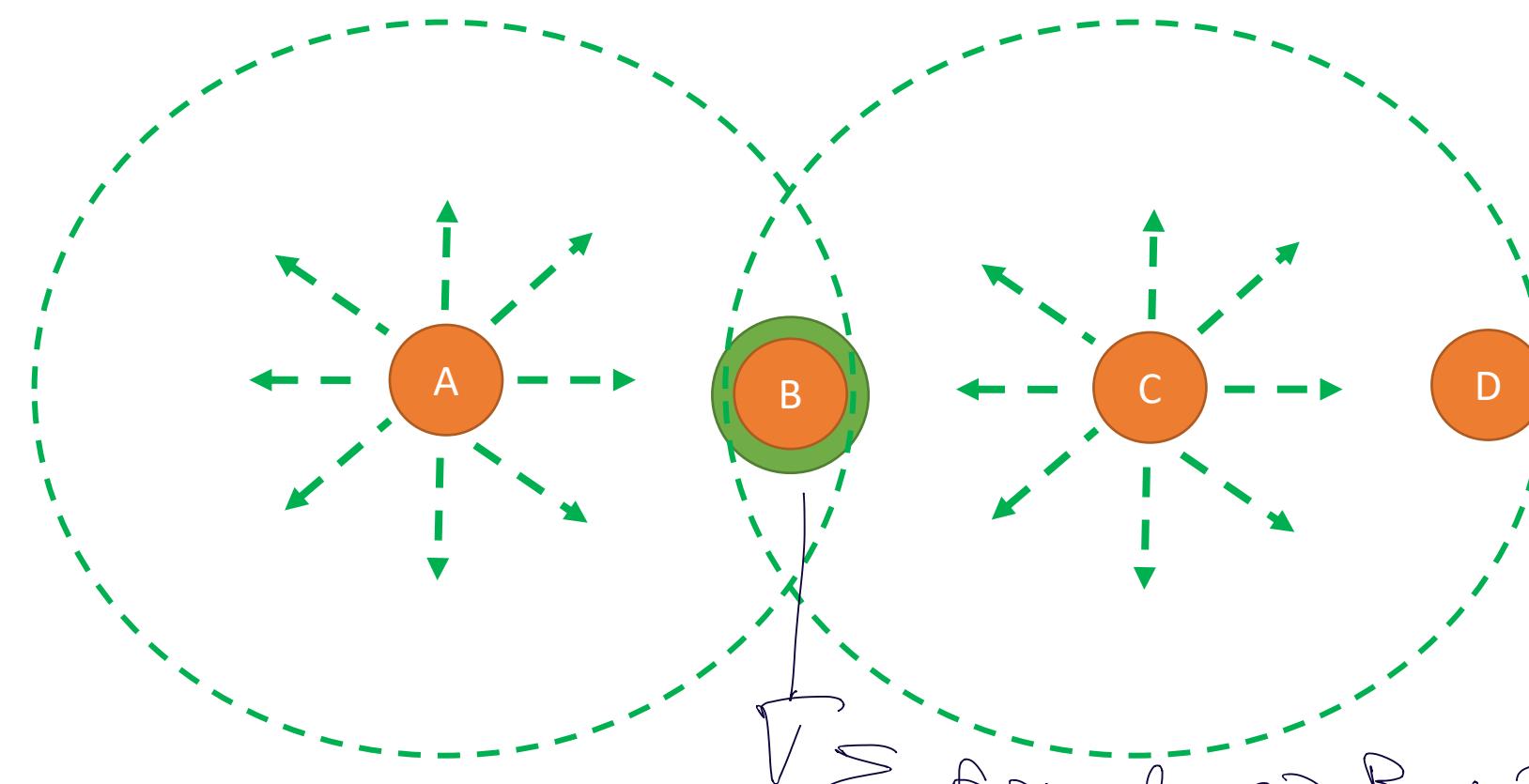


A is sending to B

The hidden terminal problem (3)



The hidden terminal problem (4)



A is sending to B

C senses the medium: it will NOT hear A, out of range

C transmits to anybody (either B or to D) : **COLLISION at B!**

Σ of signal → B not recognize
original signal
from A/C

Hidden terminal problem

C is not able to detect a potential competitor because it is out of range:
a collision happens at B (the receiver)
For the same reason A does not detect the collision
C is **hidden** with respect to the communication from A to B

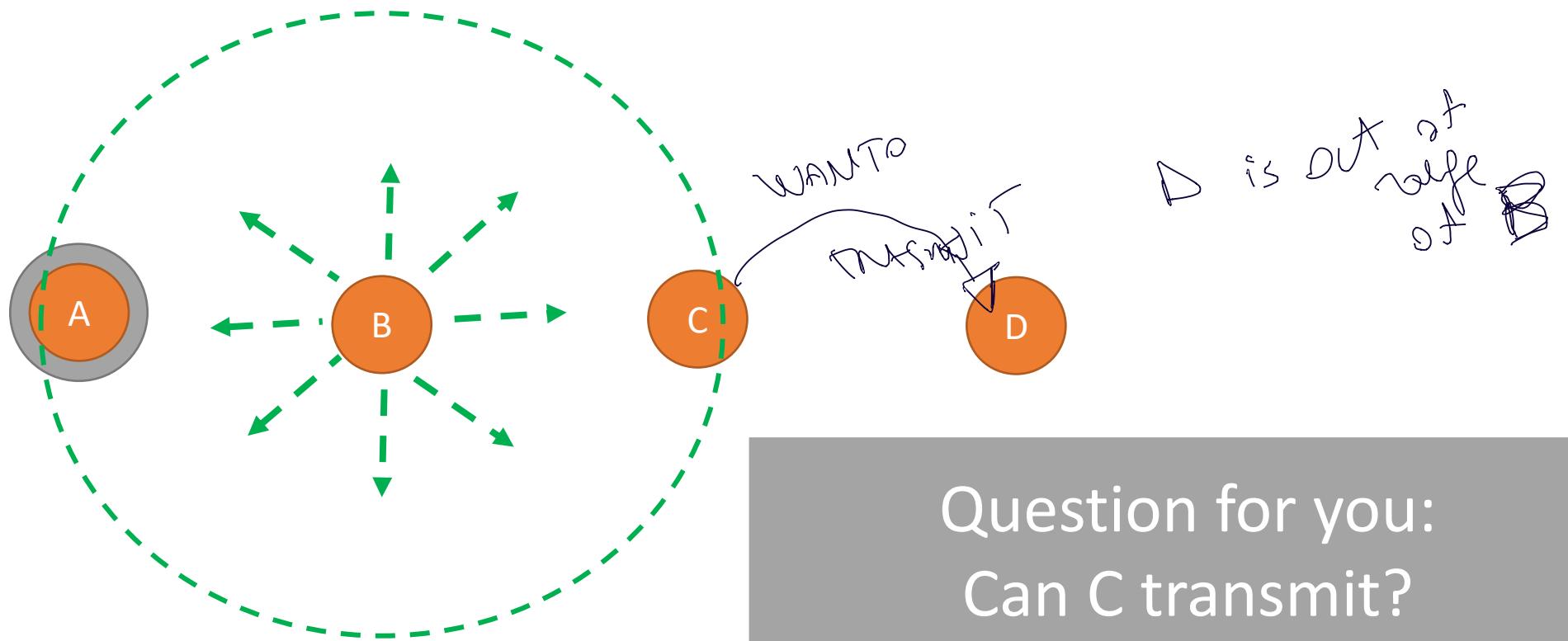
or already established

Hidden terminal problem

two or more stations which are out of range of each other transmit simultaneously to a common recipient (B)

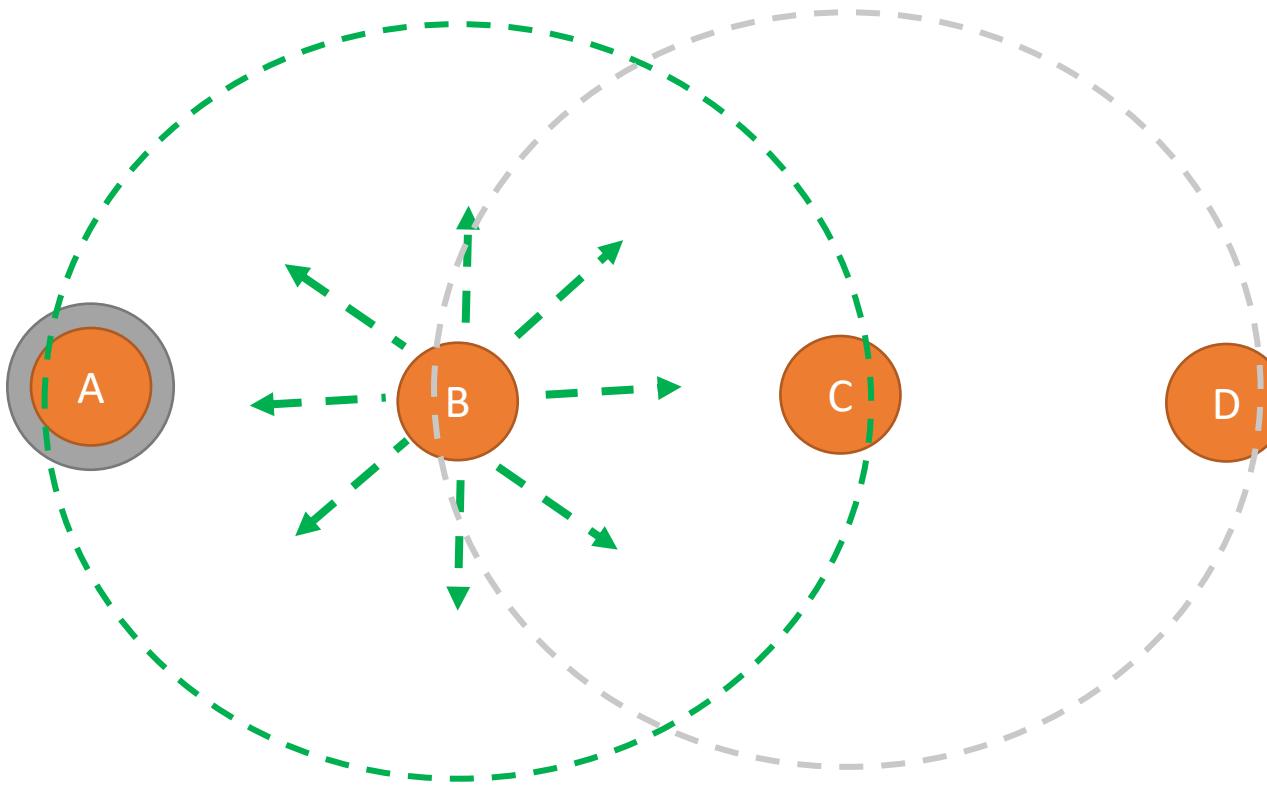
$\rightarrow A, C$

The exposed terminal problem



1. B is transmitting to A, C wants to transmit to D

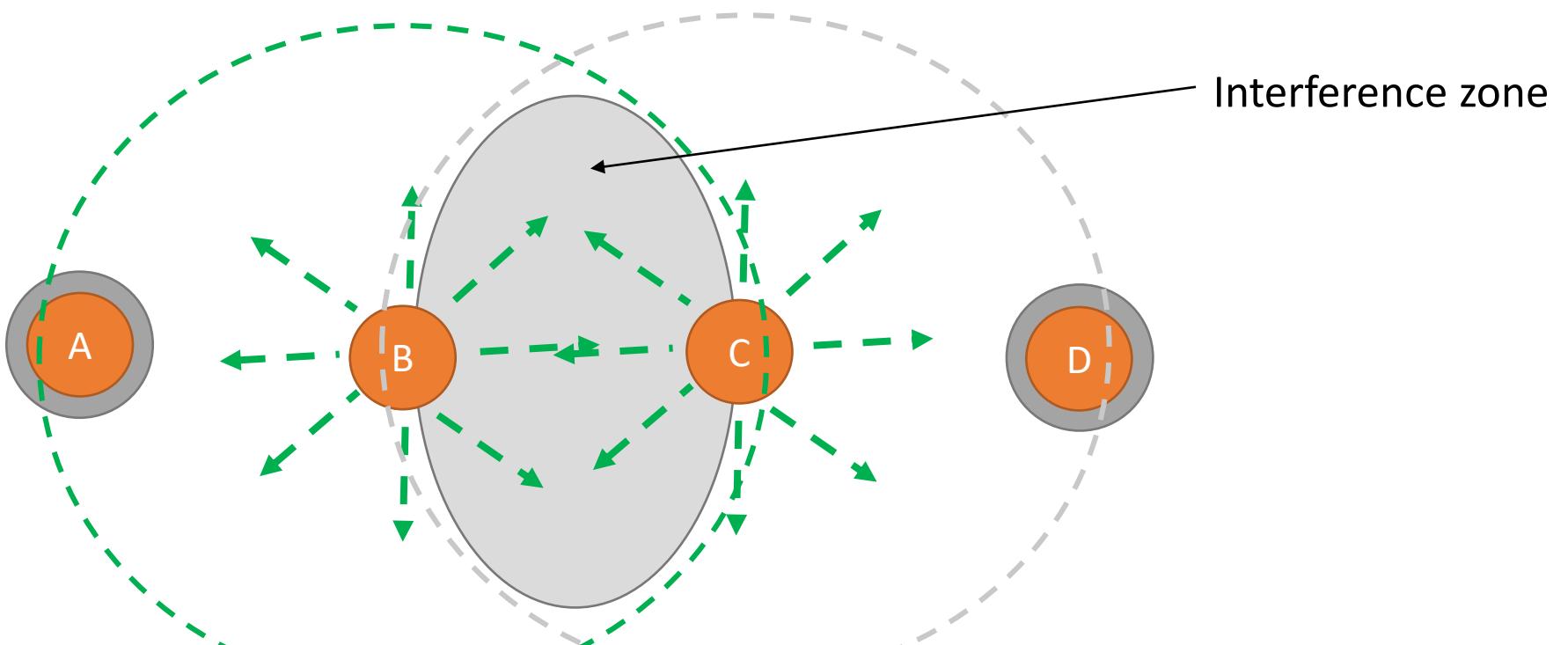
The exposed terminal problem (2)



1. B is transmitting to A, C wants to transmit to D
2. C senses the medium, concludes: **cannot transmit** to D

↳ Asciulta che B sta inviando frame ad A \Rightarrow I soli canale di comunicazione

The exposed terminal problem (3)



1. B is transmitting to A, C wants to transmit to D
2. C senses the medium, concludes: **cannot transmit** to D
3. The two transmissions can actually happen in parallel.

We have
Exposed
terminal
problem

C hears a transmission from B to A

C does not send to D although its transmission would be OK

C is **exposed** with respect to the communication from B to A

Exposed terminal problem

a transmitting station is prevented from sending frames due to interference with another transmitting station

Wireless networks

What matters is interference at the receiver not at the sender

- this cannot be checked by sensing the carrier at the sender

Multiple transmissions can occur simultaneously if destinations are out of range of each other

- a station may hear a transmission and be able to transmit without interfere with it

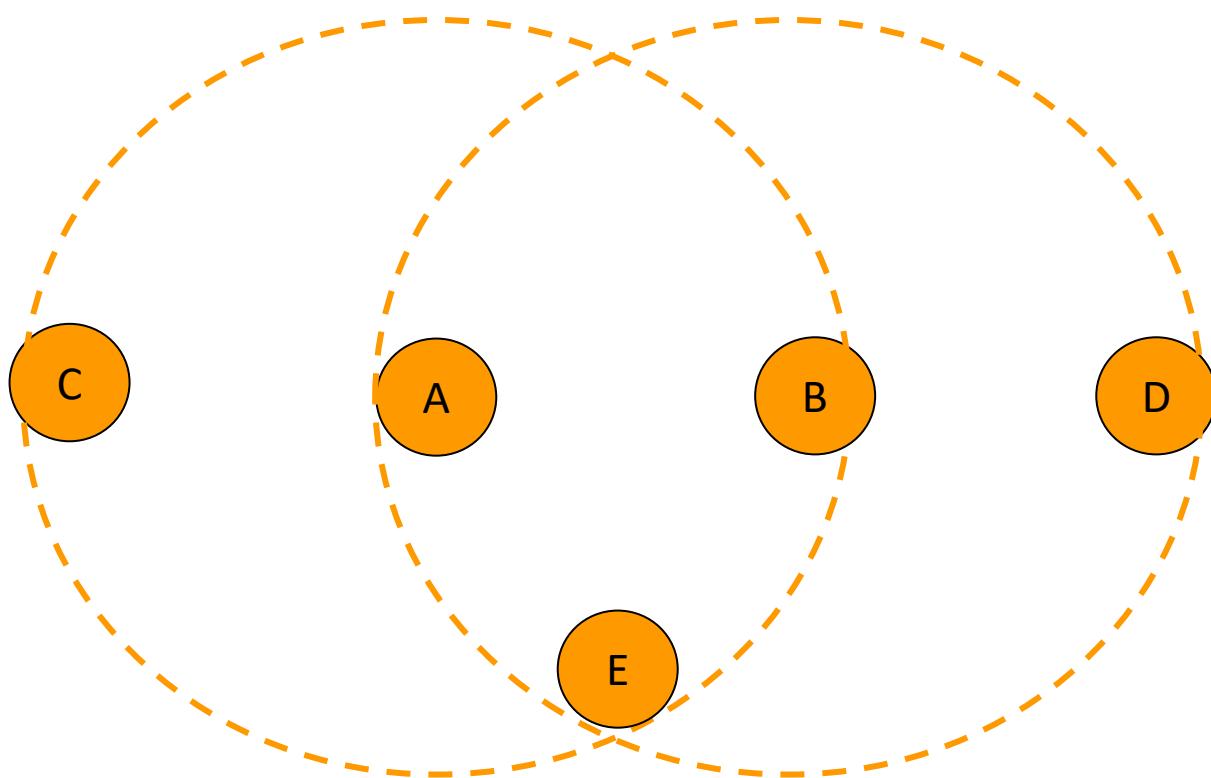
Need sensibly different MAC protocols from wired LANs

The MACA protocol

- Multiple Accesses with Collision Avoidance *Avoid collisions*
- Basic idea:
trigger SENDER
 - stimulate the receiver into transmitting a short frame first (*From sender REC.*)
 - then transmit a (long) data frame
 - stations hearing the short frame **refrain from transmitting** during the transmission of the subsequent data frame

RTT CTS

The MACA protocol



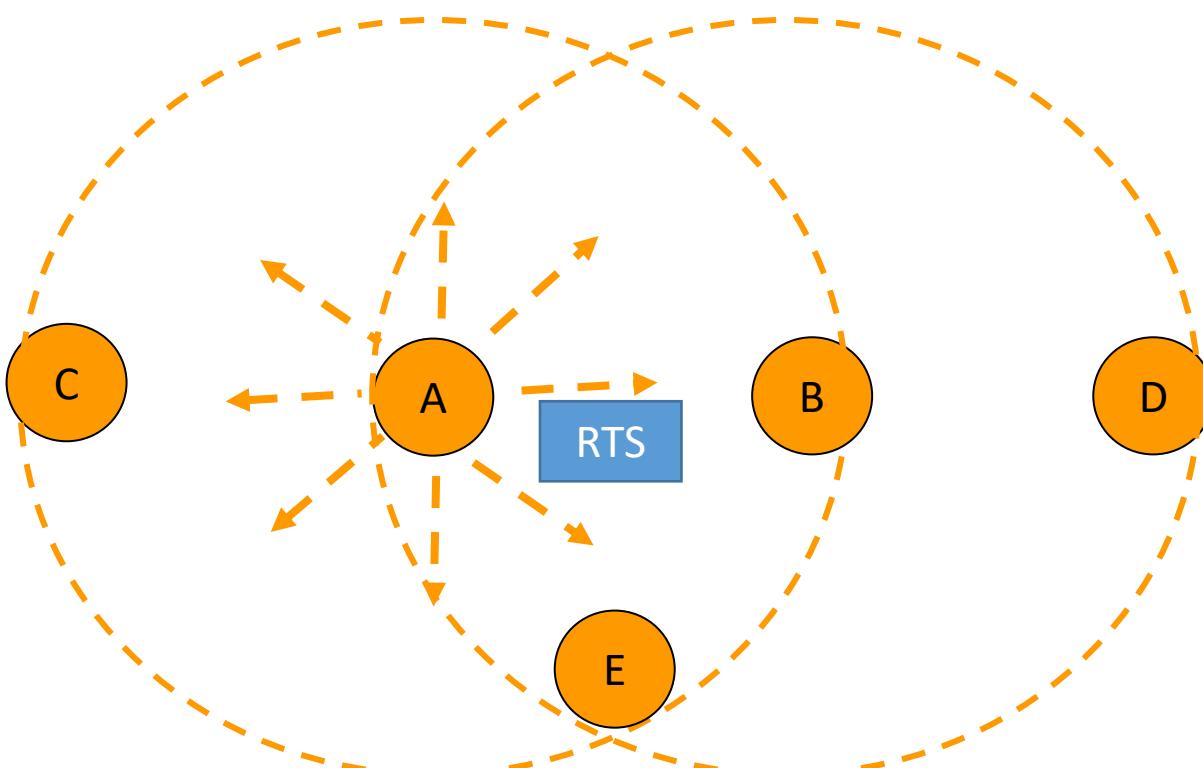
MANCÉE
Hi Seer
et d'ordre
Terrible
Punker

C is within range of A and *out of range of B and D*

D is within range of B and *out of range of A and C*

E is within range of both A and B

The MACA protocol (2)



1. A wants to transmit to B, sends a **Request To Send (RTS)** to B

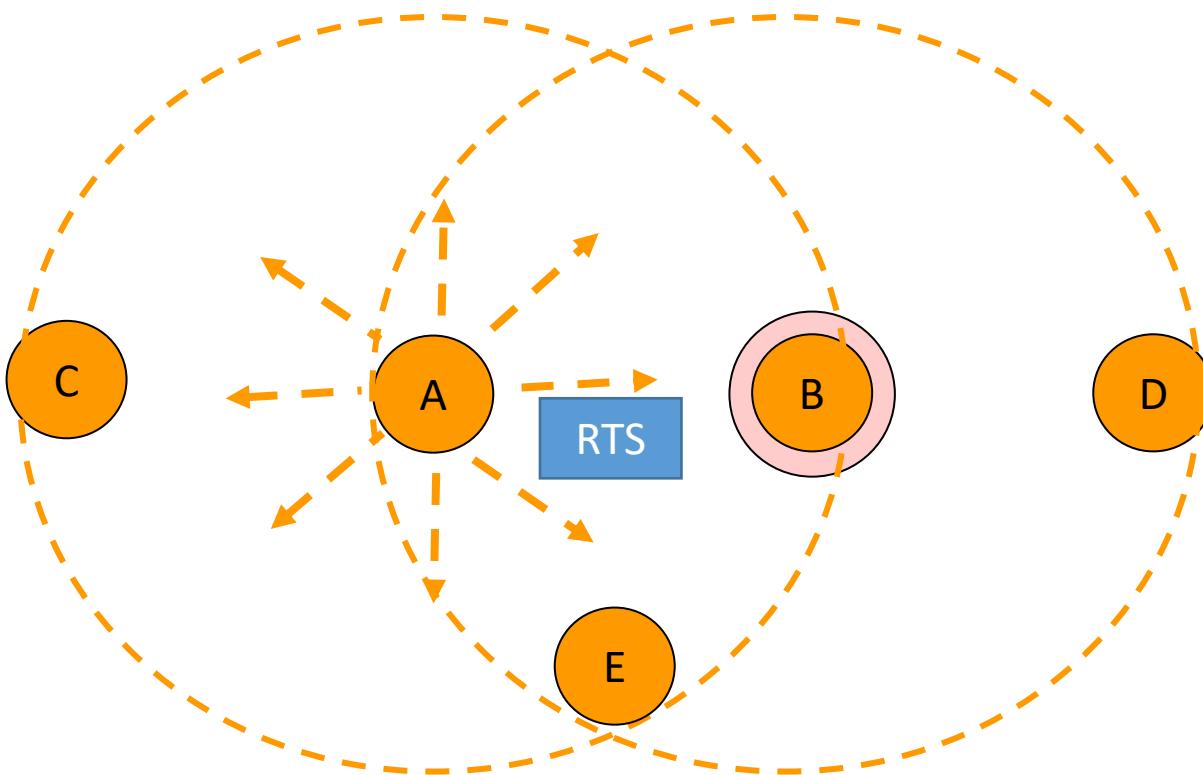
be here to send real payload



infra like
depth of
the frame
seed in 2'

The MACA protocol (3)

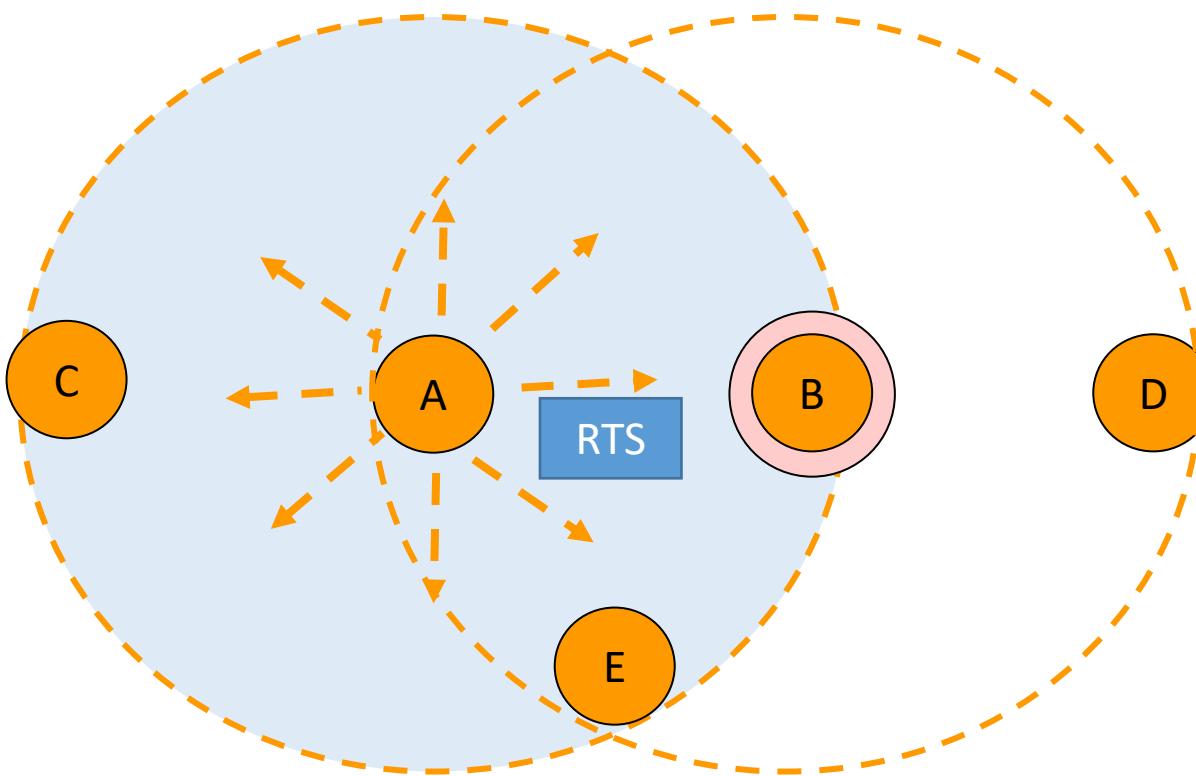
all the nodes in range of A receive



1. A wants to transmit to B, sends a **Request To Send** to B

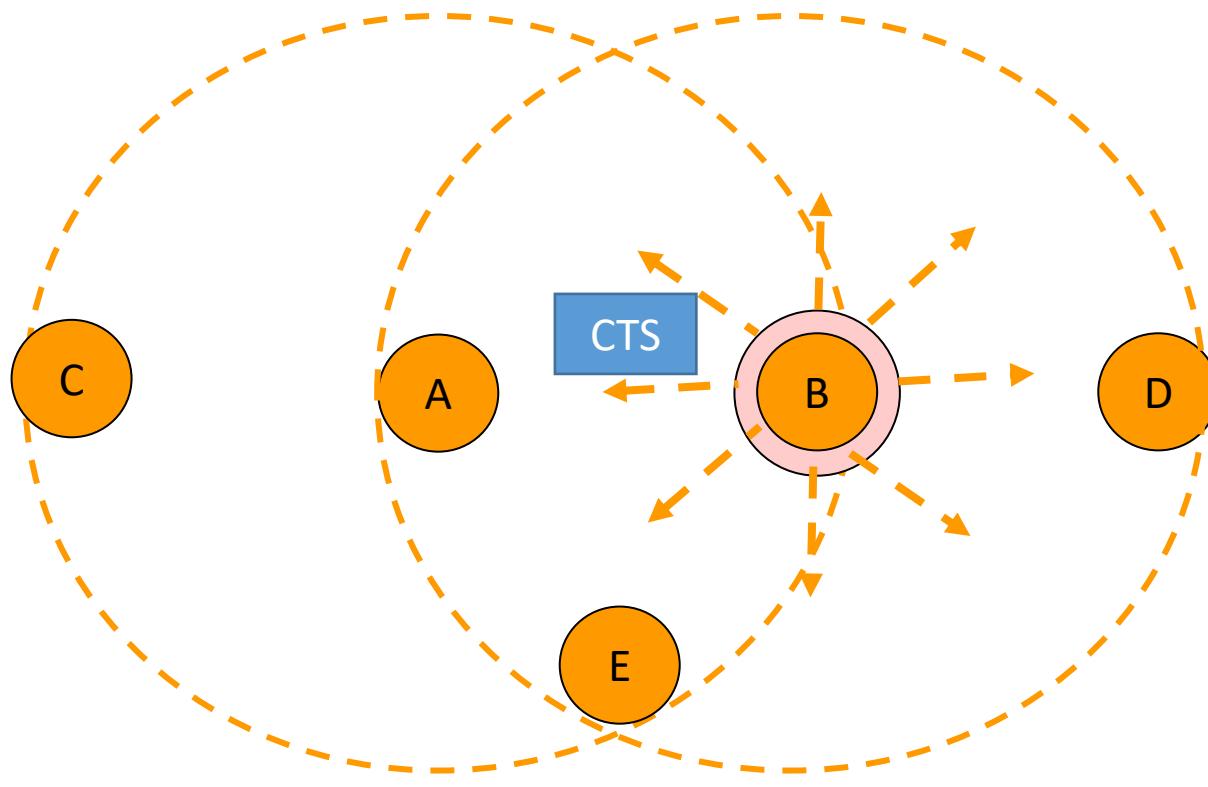
RTS is a short frame including the length of the data frame
that will eventually follow

The MACA protocol (4)



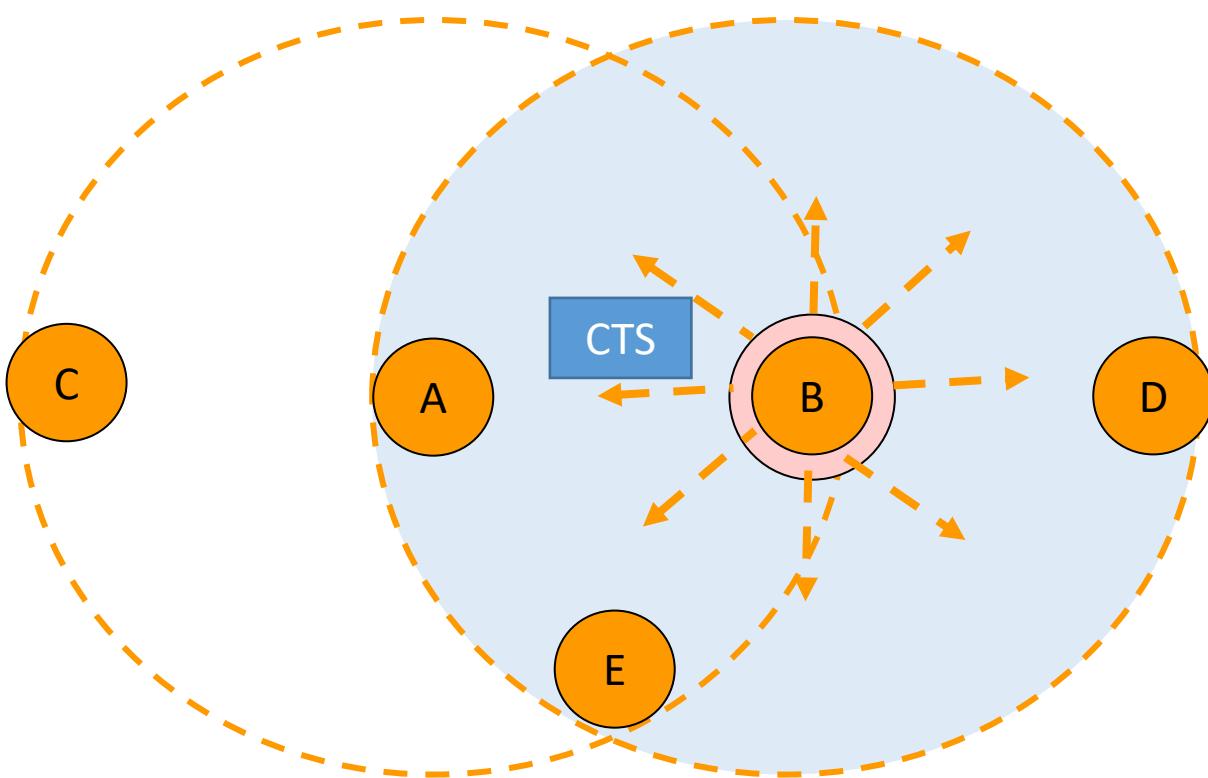
1. A wants to transmit to B, sends a **RTS** to B
B, C and E receive the **RTS** from A

The MACA protocol (5)



1. A wants to transmit to B, sends a **RTS** to B
2. If B wants to receive the message, it replies with a **Clear To Send**
CTS is a short frame with data length copied from RTS

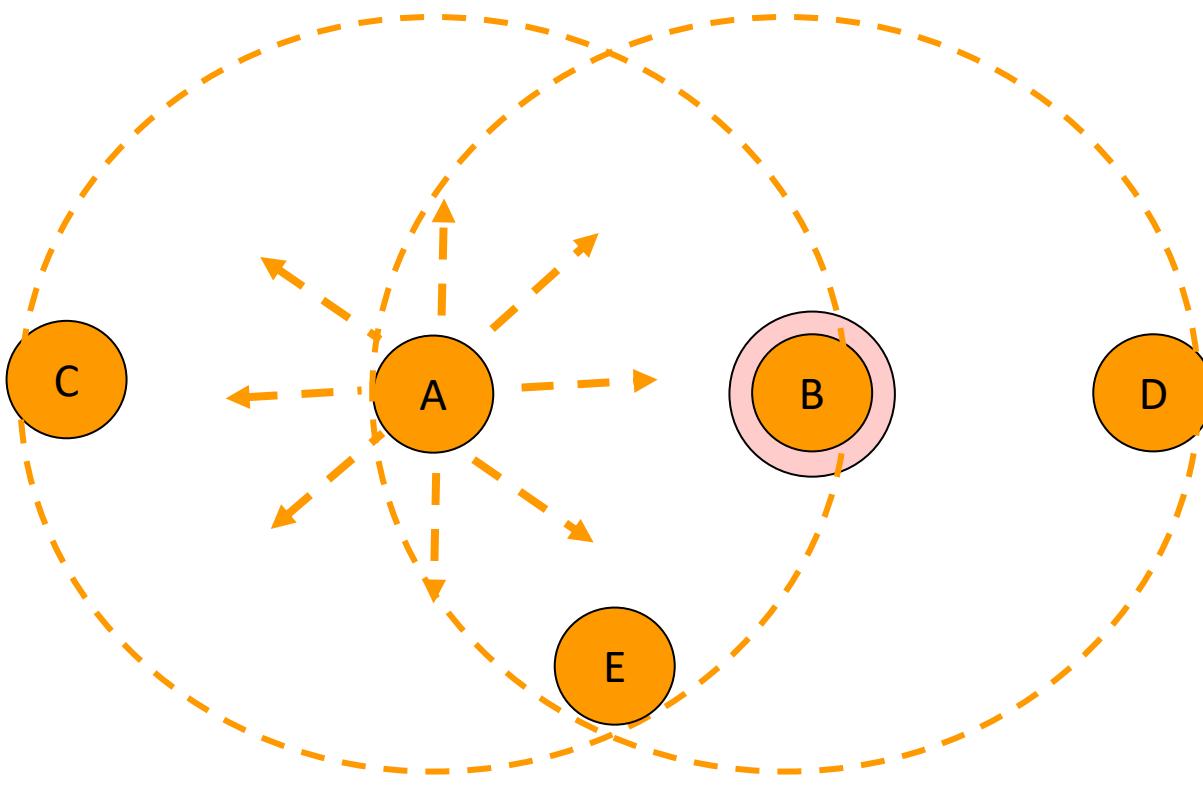
The MACA protocol (6)



1. A wants to transmit to B, sends a **RTS** to B
2. If B wants to receive the message replies with a **CTS**
CTS received by A, D, E

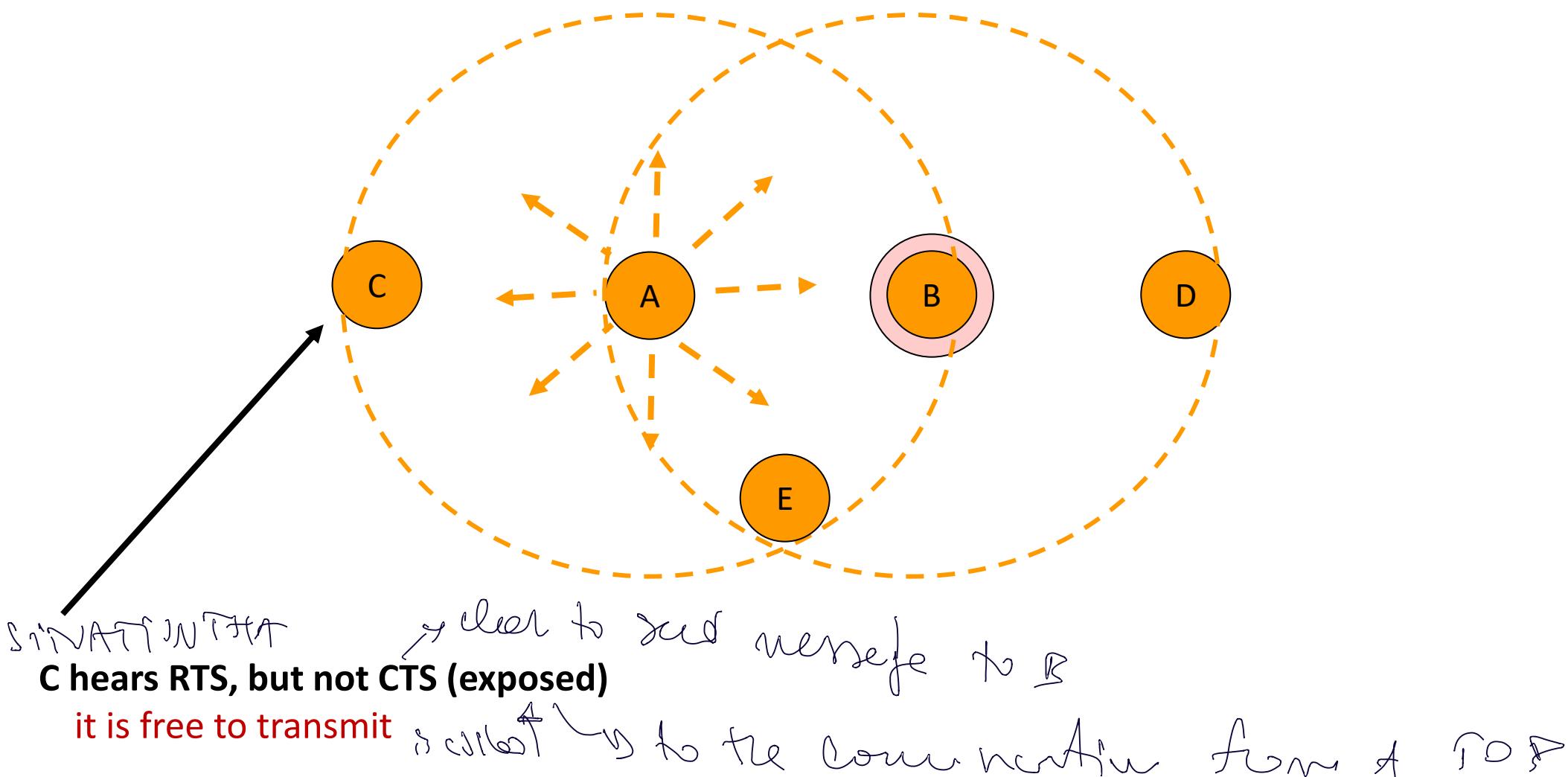
CTS receive
by A D
E

The MACA protocol (7)



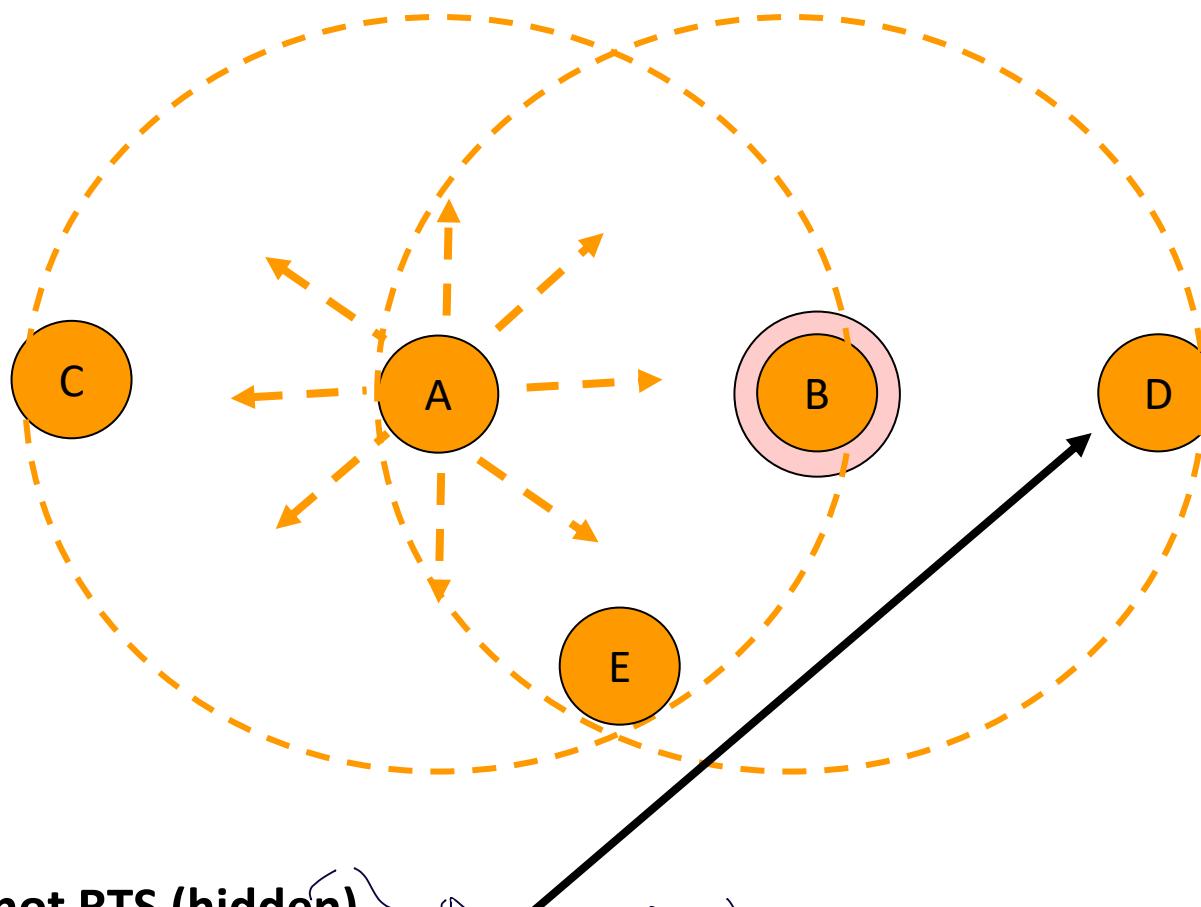
1. A wants to transmit to B, sends a **RTS** to B
2. If B wants to receive the message replies with a **CTS**
3. Upon reception of the CTS frame, A transmits the data frame

The MACA protocol (8)



⇒ C in principle is able
to communicate

The MACA protocol (9)

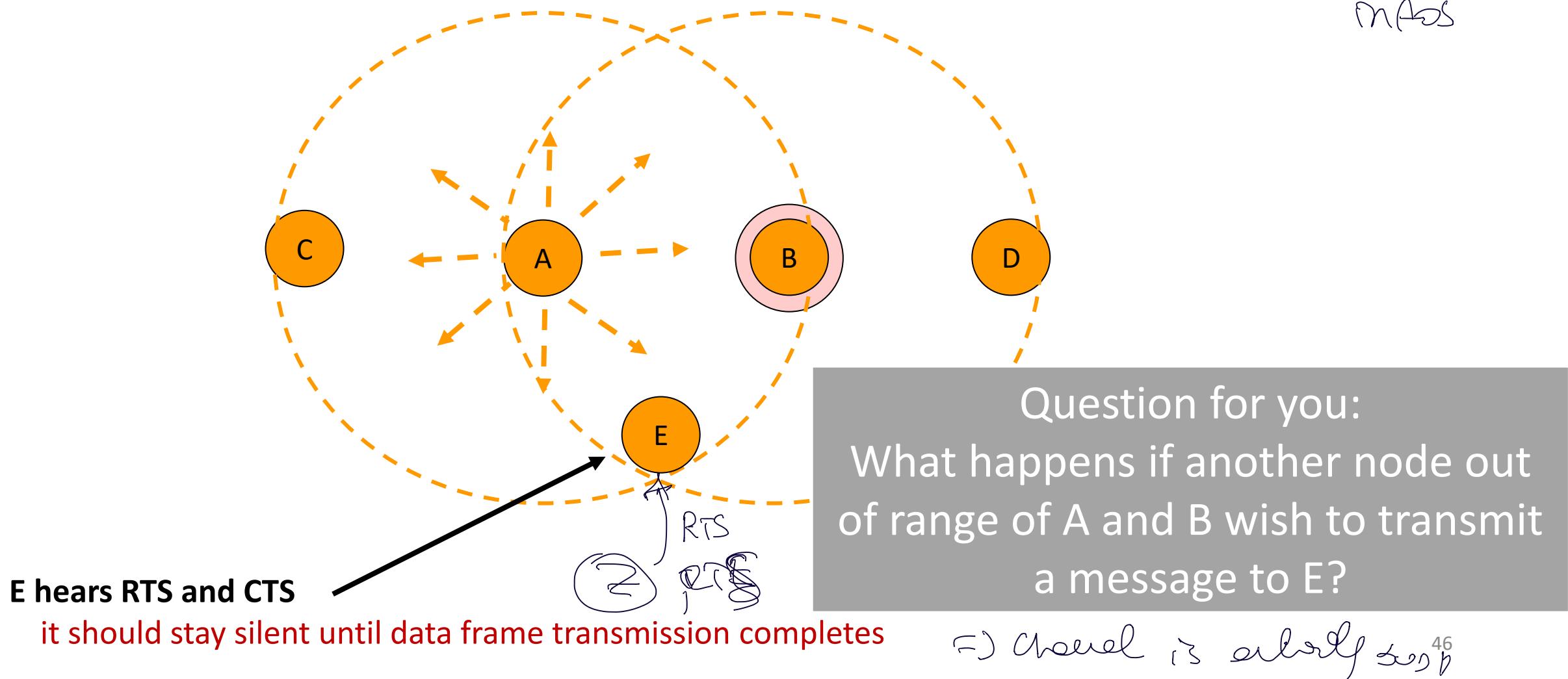


D hears CTS, but not RTS (hidden)
station

it should stay silent until data frame transmission completes (just waits for the required time)

The MACA protocol (10)

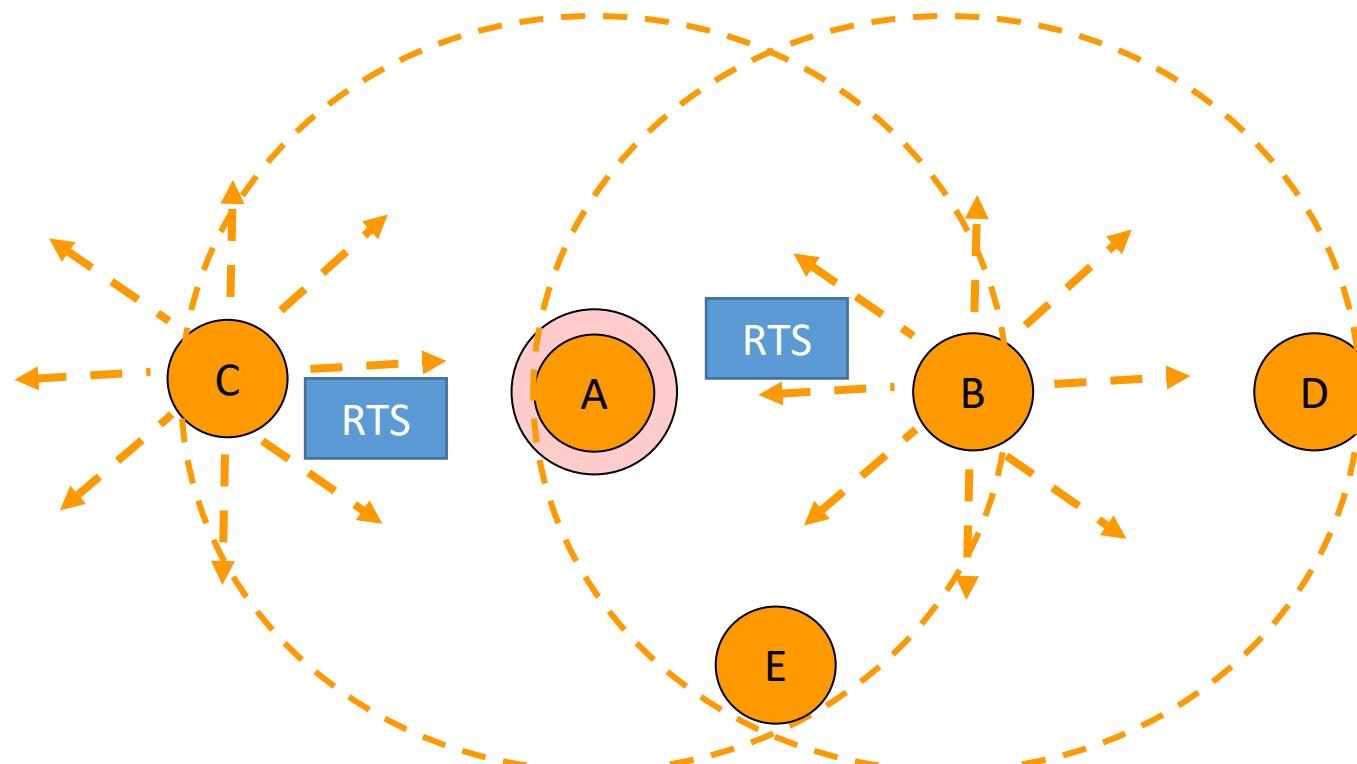
E not able to receive
CTS
MACS



→ channel is already busy

From lone
From A, B

The MACA protocol: collisions

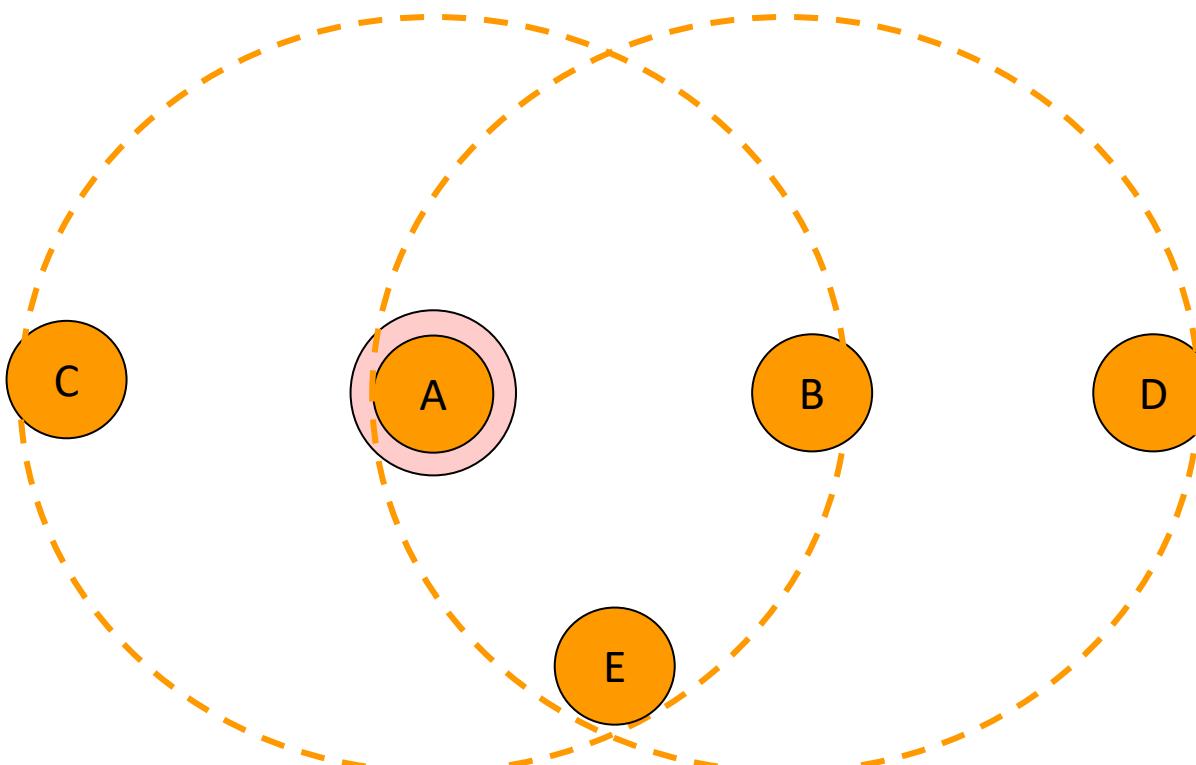


C and B send RTS simultaneously to A

The MACA protocol: collisions (2)

RTS
CTS

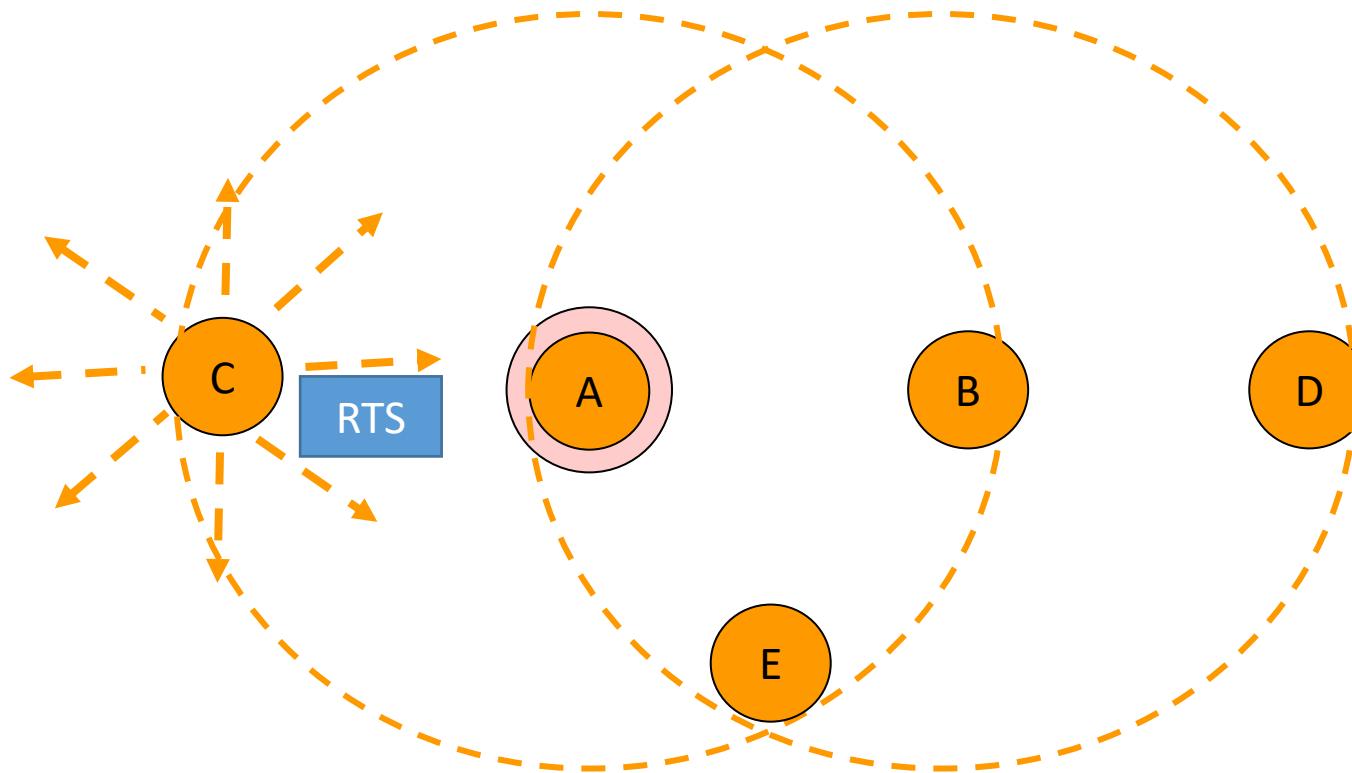
⇒ binary exponential backoff self-timing by eggy



C and B send RTS simultaneously to A

The two messages collide: No CTS is generated

The MACA protocol: collisions (3)



C and B use *Binary Exponential Backoff* (same as Ethernet) to retry RTS

MACAW: MACA for Wireless networks

resolving and RTS

Fine tunes MACA to improve performance:

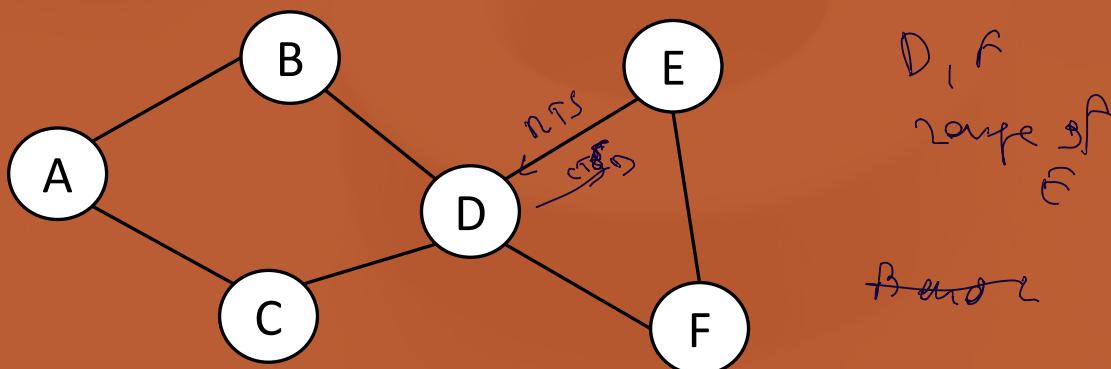
- introduces an **ACK frame** to acknowledge a successful data frame
- added **Carrier Sensing** to keep a station from transmitting RTS when a nearby station is also transmitting an RTS to the same destination
- exponential backoff is run for each separate pair source/destination and not for the single station
- mechanisms to exchange information among stations and recognize temporary congestion problems
- CSMA/CA used in IEEE 802.11 is based on MACAW

Question

Given the network in the figure, assume that the MAC protocol uses the RTS/CTS mechanism for the channel access.

Discuss which nodes detect themselves as hidden or exposed as consequence of the RTS/CTS handshake in the following cases:

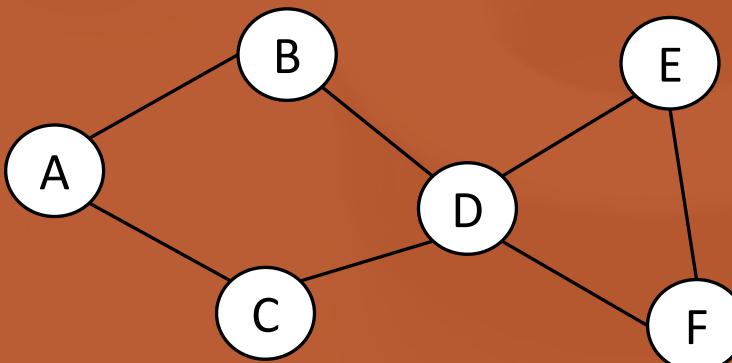
1. Hidden terminals with respect to a transmission from E to D
2. Hidden terminals with respect to a transmission from D to C
3. Exposed terminals with respect to a transmission from D to B
4. Exposed terminals with respect to a transmission from B to A



Question

Given the network in the figure:

1. assume that D hears the RTS sent by E but it does not hear the corresponding CTS. What does D can do?
2. assume that B hears the CTS sent by D but it does not hear the corresponding RTS. What does B can do? ~~Handover~~ → already RTS/CTS frame or valid
3. Assume D is receiving a communication from a node, and B did not receive the corresponding RTS & CTS and it does not hear the signal transmitted to D. If B wishes to transmit to D what happens?



until now

Should hear CTS
B will start transmission \Rightarrow Collision



Uggi MAB

Given by

only 12-13

Play more

2-3

IEEE 802.11

SIM 2 wireless

The IEEE 802.11 family

- IEEE 802.11 (Legacy mode)
 - First released in 1997 and clarified in 1999
 - rarely used today
- START WITH 1-2 Mbps data rate implemented via: physical link,
 - infrared (IR) signals,
 - radio frequencies in the 2.4GHz band (ISM -- Industrial Scientific Medical Frequency band) 1995: 802.11.a
- many degrees of freedom: interoperability among different products was a challenge
- rapidly supplemented (and made popular) by 802.11b
- most used today 802.11a/b/g/n

IEEE 802.11 family (2)

- IEEE 802.11a
 - Released in 1999
 - • Operating frequency: 5 GHz band (Unlicensed National Information Infrastructure U-NII band)
higher
 - Throughput (typ): 23 Mbps
 - Data rate (max): 54 Mbps
- IEEE 802.11b
 - Released in 1999
 - Operating frequency: 2.4GHz band (ISM band)
 - potential interference with other appliances : cordless telephones, microwave ovens etc
 - Throughput (typ): 4.3 Mbps
 - Data rate (max): 11 Mbps
 - problem: domestic env., interference
with telephone, microwave ovens

IEEE 802.11 family (3)

- IEEE 802.11g
 - Released in 2003
 - Operating frequency: 2.4GHz band (ISM band)
 - Throughput (typ): 19 Mbps and more... $> \beta$
 - Data rate (max): 54 Mbps
- IEEE 802.11n \rightarrow
 - Released in 2009
 - Operating frequency: 2.4GHz band and 5GHz band
 - Throughput (typ): 74 Mbps
 - Data rate (max): 248 Mbps
 - Support of MIMO technologies for using multiple antennas at the transmitter and the receiver

better user resources

\rightarrow a node can use multiple antennas transmitting/receiving

IEEE 802.11 family (4)

- Now available newer versions:
- WiFi 5 - IEEE 802.11ac
 - Released in 2013
 - Operating frequency: 2.4GHz and 5 GHz bands
 - Data rate (max): 1.3 Gbps at 5 GHz
 - Data rate (max): 450 Mbps at 2.4 GHz
- WiFi 6 - IEEE 802.11ax
 - Released in 2019
 - reaches up to 10 Gbps
 - improvements in power consumption and security

IEEE 802.11 Wireless LAN

overall view

know min used freq,
covered area

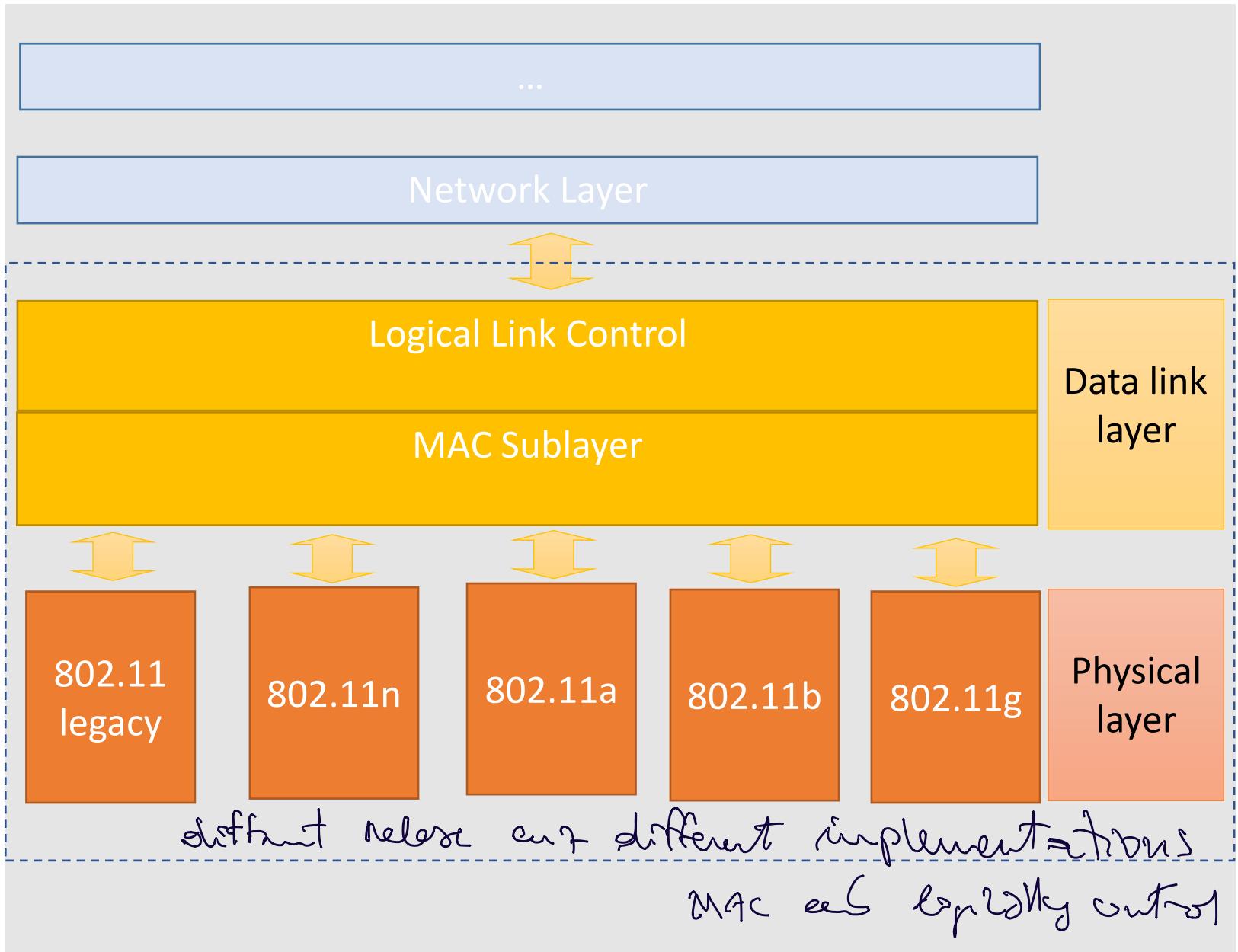
IMPORTANT
↓

IEEE 802.11 standard	Year	Max data rate	Range	Frequency
802.11b	1999	11 Mbps	30 m	2.4 Ghz
802.11g	2003	54 Mbps	30m	2.4 Ghz
802.11n (WiFi 4)	2009	600	70m	2.4, 5 Ghz
802.11ac (WiFi 5)	2013	3.47Gpbs	70m	5 Ghz
802.11ax (WiFi 6)	2020 (exp.)	14 Gbps	70m	2.4, 5 Ghz
802.11af	2014	35 – 560 Mbps	1 Km	unused TV bands (54-790 MHz) before TV
802.11ah	2017	347Mbps	1 Km	900 Mhz low freq allows larger coverage

- all use CSMA/CA for multiple access, and have base-station and ad-hoc network versions
 - lower freq, higher range → good for sensor in fields

IEEE 802.11: protocol stack

for standards



IEEE 802.11 Architecture

- A group of stations operating under a given coordination function
- ACCESSIONS
to users
- may (or may not) use a base station (Access Point - AP)
transmitter to receiver
 - if using AP a station communicates with another by channeling all the traffic through a centralized AP
 - AP can provide connectivity with other APs and other groups of stations via fixed infrastructure
- AC can be used
communication b/w nodes & AC
(can't make associate w/ AC)
- all traffic thru a centralised

IEEE 802.11 Architecture

- Supports ad hoc networks, which are, in the IEEE 802.11 view :

a group of stations that are under the direct control of a single coordination function without the aid of an infrastructure network

- a station can communicate directly with another without channeling all the traffic through AP

⇒ AC ~~must contain~~ ^{adjust}
can talk each other

no info

no centralized
or AC

Distributed
among nodes

C
KOF

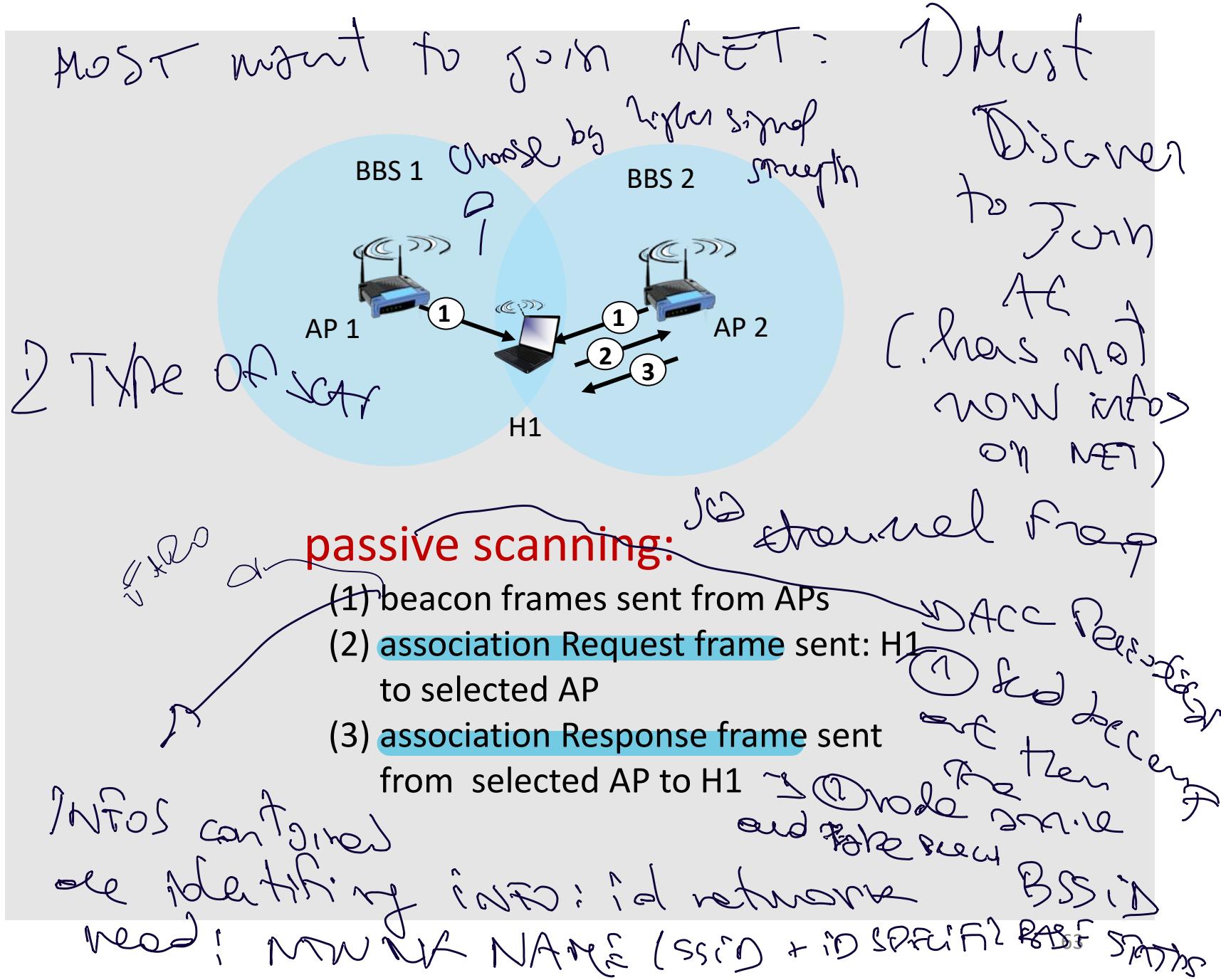
MW

IEEE 802.11 Channels, association (in MAST).

3 ways

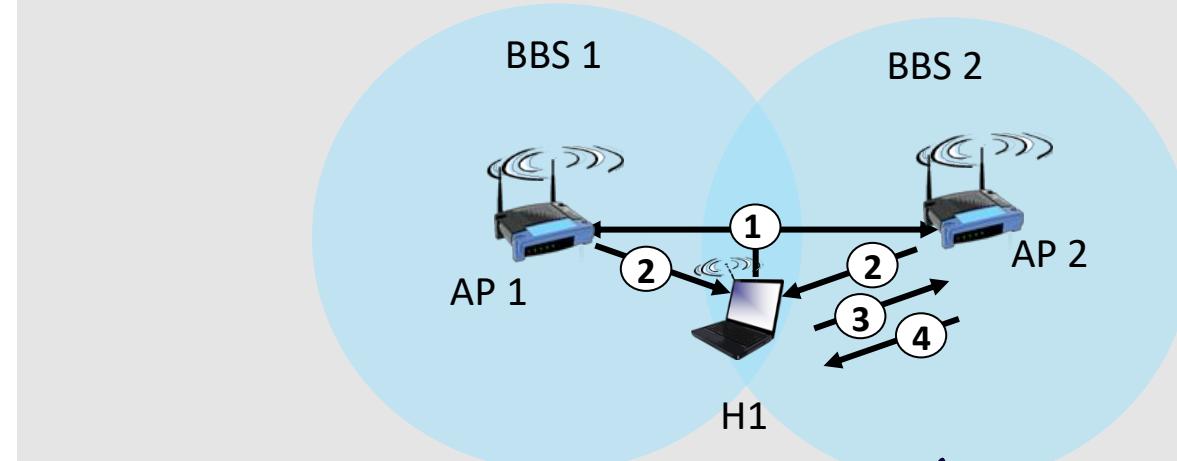
- node joins to a trunk that have AC
- range of freq divided in $\frac{1}{n}$ channels
- spectrum divided into channels at different frequencies
 - AP admin chooses frequency for AP
 - interference possible: channel can be same as that chosen by neighboring AP!
 - arriving host: must **associate** with an AP
 - scans channels, listening for *beacon frames* containing network's name (SSID) and MAC address of the AP (BSSID)
 - selects AP to associate with
 - then may perform authentication
 - then typically run DHCP to get IP address in AP's subnet
- AC communicates the channel using that freq with 'division'
- ②
- 3 ways

IEEE 802.11 passive/activ e scanning



IEEE 802.11 passive/activ e scanning

node select 1 AC, perform AUTH. if needed.



inter
frame
space

active scanning:

- (1) Probe Request frame broadcast from H1
- (2) Probe Response frames sent from APs
- (3) Association Request frame sent: H1 to selected AP
- (4) Association Response frame sent from selected AP to H1

one needed cut to have more ACs.

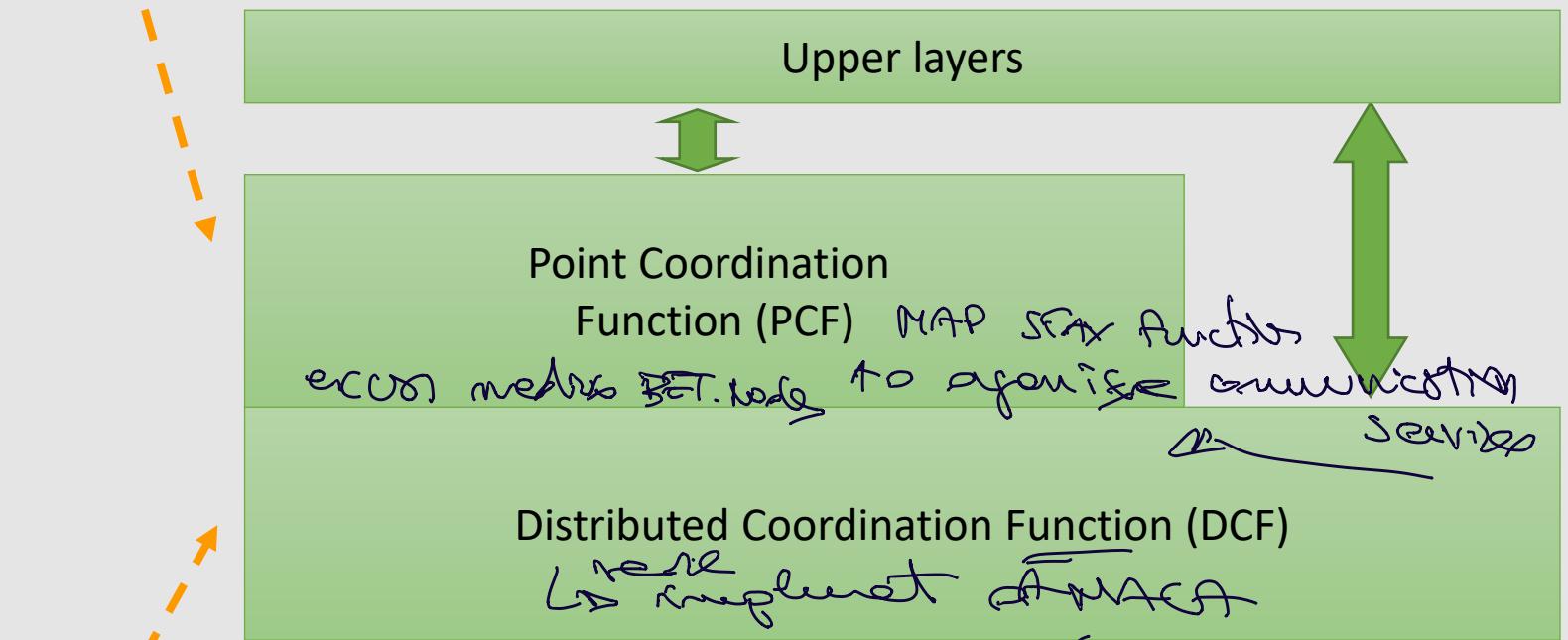
→ ~~send~~
Send
beacon
frames.
Waiting
replies
of AC

IEEE 802.11 Architecture: MAC Sublayer

Used for contention-free services and based on DCF

Used for contention services

But what about in the AC (offload)
so wait for ~~the~~ offload
private 802.11



distributed access
to the shared
medium

IEEE 802.11 Architecture: MAC Sublayer

Two modes of operations:

- DCF : Distributed Coordination Function
 - completely decentralized
 - thought for best effort asynchronous traffic
- PCF : Point Coordination Function *(contention free)*
 - uses base station to control all activity in its cell
 - thought for delay-sensitive traffic
 - AP polls stations for transmissions
 - based on DCF

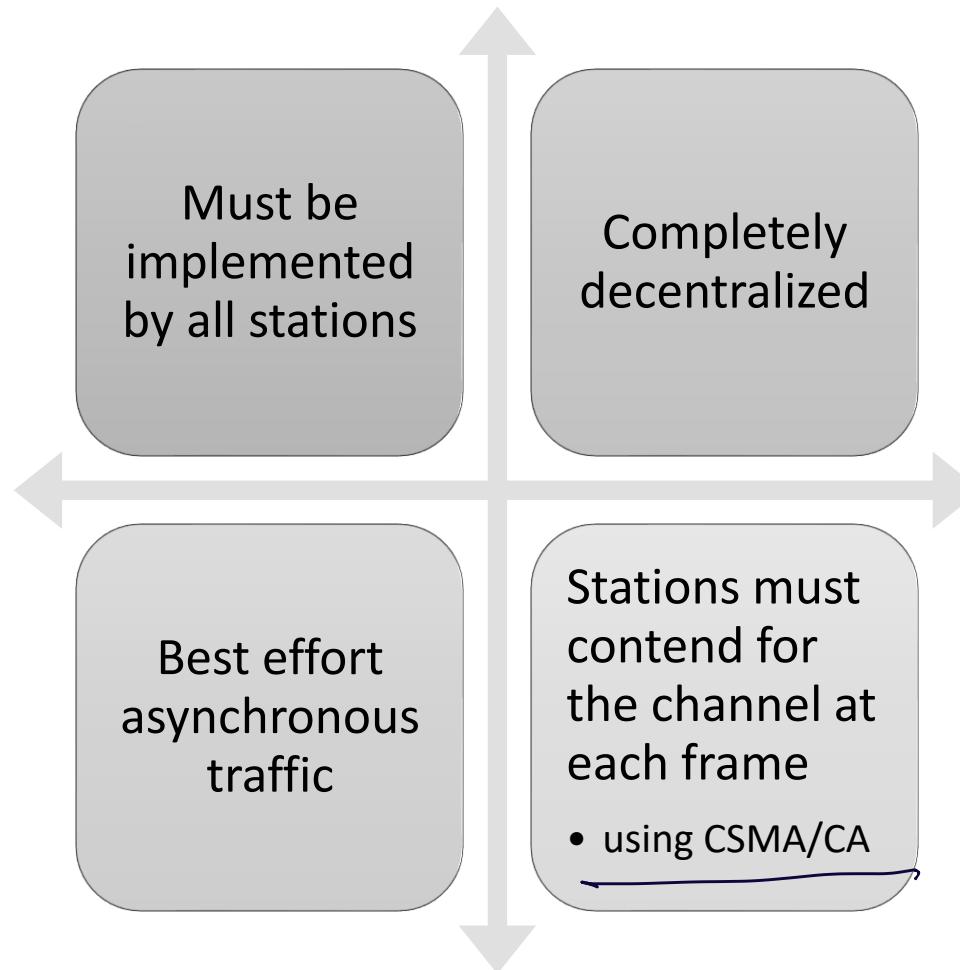
DCF

contention free
calling mech
(com. mech)
& # devices

DCF **must** be implemented by all stations

DCF and PCF can be active at the same time in the same cell

IEEE 802.11:
Distributed
Coordination
Function
– DCF (1)



Based on IEEE
try to send

IEEE 802.11:

Distributed
Coordination
Function
- DCF (2)

node 2
keeps RFT & CTS
from other stations
(A, B)

want avoid collision (TME)
end if know when
is surely short

Carrier sensing is performed at two levels:

physical CS

before frame: check if another
station is transmitting

- checking the frequency to determine whether the medium is in use or not

- physical carrier sense to detect an incoming signal

- detects any activity in the channel due to other sources

virtual CS

- performed sending duration information in the header of an RTS, CTS and data frame

- Keeps the channel "virtually busy" up to the end of a data frame transmission

- A channel is marked busy if either the physical or the virtual CS indicate busy

IEEE 802.11:

Distributed Coordination Function – DCF (3)

- Priority access to the medium is controlled through the use of interframe space (IFS) time intervals

- IFS: mandatory periods of idle time on the transmission medium

- Three IFS specified by the standard:

- short IFS (SIFS)
- point coordination function IFS (PIFS)
- Distributed coordination function IFS (DIFS)

- SIFS < PIFS < DIFS

- stations only required to wait a SIFS have the highest priority

microsec

shorter \rightarrow micro sec, shorter

bigger

two higher

wait less often
other priorities

min time ned
zone wait +
sub run free

Time
Between
2 frames

Summary



INFRASTRUCTURED AND
INFRASTRUCTURE-LESS
WIRELESS NETWORKS



PROPERTIES OF
WIRELESS CHANNELS



HIDDEN AND EXPOSED
TERMINALS



MACA PROTOCOL



IEEE 802.11