

# 1 November - Midterm Review

## Logistics

- Hw 4 extended to next Friday: implies there will be no tree coding problems on exam
- You will take your mid-term in your lab section you are registered in
- Linux computers in lab, Windows computers in DAC testing center
  - Both have VS Code and OpenJDK
  - Also OK to use your own laptop
- Exam will be published on Canvas
  - Let me know when DAC exam is scheduled so I can set up your exam
  - Can take exam in DAC testing center any day next week

## Allowed materials

- Official docs allowed: <https://docs.oracle.com/en/java/javase/17/docs/api/help-doc.html>
- Recommend navigating the docs before exam
- One, single-sided sheet of notes (handwritten or typed) to be turned in when you leave the exam
- Blank scratch paper
- ANY other accessed web materials or local materials (phones, etc) will be grounds for a zero and referral to Academic Integrity.
- Coding problem will have starter code and tests for download, then you will re-upload code to Canvas. Programming, running and testing code will all happen locally. Can use any editor, but no AI / coPilot / etc.
- Coding problem scope: implementing one or two methods for insertion/deletion in data structure we have covered.
- Not testing on Git / Linux skills

## Exam Topics

Topic recap:

- Java understanding, reading and analyzing code
- Runtime Analysis
- Stacks, Queues, Linked Lists (“Linear” data structures)
- Hash Tables, Hashing
- Binary Search
- Sorting
- Binary Search Trees
- More general trees

## Open floor for questions

## Practice problems

### Code interpretation and runtime

Consider the following code. What is the output of `fn1(8)` and `fn2(7)`?

```
fn1(int n) {  
    if (n >= 0) {  
        System.out.print(n);  
        fn1(n-2);  
    }  
}
```

Output: 8 6 4 2 0

```
fn2(int n) {  
    if (n >= 0) {  
        fn2(n-2);  
        System.out.print(n);  
    }  
}
```

Output: 1 3 5 7

### Runtime

Consider the following psuedocode for computing the  $n$ th Fibonacci number  $F_n$ . Recall that  $F_n = F_{n-1} + F_{n-2}$  and  $F_0 = 0$  and  $F_1 = 1$ .

```
fib(n):  
    if n <= 1:  
        return n;  
    return fib(n-1) + fib(n-2);
```

How many times is the `fib` function called if you call `fib(3)` (including the original call of `fib(3)`)?

- first call: `fib(3)`-recurses;
- first recurse: second call, `fib(2)`-recurses; third call, `fib(1)`-returns;
- second recurse: fourth call, `fib(1)`-returns; fifth call, `fib(0)`-returns;
- five total calls

Is the code efficient? Why or why not?

- Looking for answers stating no, this approach involves repeated computation of the same value in different branches of the recursion.
- Actual complexity is  $\mathcal{O}(2^N)$  (not expected to know or derive)
- To compute the  $N$ th Fibonacci number we can do an iterative approach:

```

fib = 1;
f(int N) {
    for i in (1..N) {
        fib += fib;
    }
    return fib;
}

```

## Sorting

1. What is the worst-case runtime of merge-sort on an array of size  $n$ ? Explain your reasoning.

Answer:  $O(n \log(n))$ , see for example a reverse-sorted array: will have to build full binary tree and merge all individual elements.

2. True or false:
  - If you are sorting an array where the key is unique for each element, you will get the same result whether your sort algorithm is stable or not.
    - True
  - Quick sort has the same average case time complexity as its worst-case time complexity.
    - False: Worst is  $O(n^2)$  (always pick largest/smallest as pivot), average is  $O(n \log n)$  (similar to merge sort).
3. Consider insertion sort, selection sort, merge sort, and quick sort. Which sort(s) never compare the same two elements twice?
  - Quicksort: elements are divided according to pivot, then sorted within branches
  - Mergesort: once we compare two elements and merge, these sub-lists are sorted
  - Insertion sort: when we insert element, we are comparing to sorted elements, compare once
  - Selection sort: DOES potentially compare twice, because we *select* the next smallest element from unsorted list and swap with the

## Hashing

- What is the time complexity of adding an element to a hash map?
  - $O(1)$
- Multiple choice: what is a *collision* in a hash table?
  - Two key-value pairs that have equal keys but hash to different indices.
  - Two key-value pairs that have different keys but hash to the same index. (Correct)
  - Two key-value pairs that have equal keys but different values.
  - Two key-value pairs that have different keys and hash to different indices.

## Trees

- What is the asymptotic (big-O) height of a balanced binary tree with  $N$  elements?
  - $O(\log N)$
- What is the big-O height of a worst-case unbalanced binary tree with  $N$  elements?
  - $O(N)$
- Given tree, write out order of elements in (post | pre | in) order.