```
堯哥 cheat
#include <bits/stdc++.h>
using namespace std;
using 11 = long long;
using vi = vector<int>:
using pii = pair<int, int>;
using pll = pair<ll, 11>;
template<typename T> using vec = vector<T>;
template<typename T> using Prior = priority_queue<T>;
template<typename T> using prior = priority queue<T, vector<T>.
#define yccc ios_base::sync_with_stdio(false), cin.tie(0)
#define al(a) a.begin(),a.end()
#define F first
#define S second
#define REP(i, n) for(int i = 0; i < n; i++)
#define REP1(i, n) for(int i = 1; i <= n; i++)
#define eb emplace_back
#define pb push back
#define mp(a, b) make_pair(a, b)
#define debug(x) cout << " > " << #x << ": " << x << endl:
#define devec(v) cout << " > " << #v << ": "; for (auto i : v) cout << i
<< ' ': cout << endl:
#define devec2(v) cout << " > " << #v << ":\n"; for (auto i : v) { for
(auto k : i) cout << ' ' << k: cout << endl: }
const int INF = 1e9:
const int nINF = -1e9:
const 11 11INF = 4*1e18:
const int MOD = 1e9+7;
11& pmod(11& a, 11 b) { a = (a+b) % MOD; return a; }
11& pmod(11& a, 11 b, 11 c) { a = (a+b) % c; return a; }
11& mmod(11& a, 11 b) { a = (a-b+MOD) % MOD; return a; }
11& mmod(11& a, 11 b, 11 c) { a = (a-b+c) % c; return a; }
11& tmod(11& a, 11 b) { a = (a*b) % MOD; return a; }
11 mul(11 a, 11 b) { return (a*b) % MOD; }
11 mul(11 a, 11 b, 11 c) { return (a*b) % c; }
vim setting :
syntax enable
syntax off
set number
set nonumber
set tabstop=4
colo default/darkblue/koehler/desert/ron/torte
data structure
Seament Tree
```

#define IR(X) ((X * 2) + 2) #define MAXN 500005 struct Node{ long long sum; long long lazy_tag; int size; }; short dataseq[MAXN]; Node seq[MAXN * 4 + 5]; void build(int l, int r, int L, int R, int index); void modify(int l, int r, int L, int R, int index, long long Query(int l, int r, int L, int R, int index); void pull(int index); void pull(int index);

seq[index].sum = seq[IL(index)].sum + seq[IR(index)].sum;

seq[IL(index)].lazy_tag += seq[index].lazy_tag;

seq[index].size = seq[IL(index)].size + seq[IR(index)].size;

seq[IL(index)].sum += seq[index].lazy_tag * seq[IL(index)].size;

#define IL(X) ((X * 2) + 1)

seq[index].lazy_tag = 0;

void push(int index){

```
seq[IR(index)].lazy_tag += seq[index].lazy_tag;
    seq[IR(index)].sum += seq[index].lazy_tag * seq[IR(index)].size;
    seq[index].lazy_tag = 0;
void build(int 1, int r, int L, int R, int index){
        seq[index].sum = dataseq[1];
        seq[index].size = 1;
        seq[index].lazy_tag = 0;
    int M = (L + R) / 2;
    build(1, M, L, M, IL(index));
          build(M + 1, r, M + 1, R, IR(index));
    pull(index);
void modify(int 1, int r, int L, int R, int index, long long Add){
    if(1 == L \&\& r == R){
        seq[index].lazy_tag += Add;
        seq[index].sum += Add * seq[index].size;
        return:
    push(index);
    int M = (L + R) / 2;
    if(r \ll M)
        modify(1, r, L, M, IL(index), Add);
    else if(1 > M){
        modify(1, r, M + 1, R, IR(index), Add);
    }else{
        modify(1, M, L, M, IL(index), Add);
        modify(M + 1, r, M + 1, R, IR(index), Add);
    pull(index):
long long Query(int 1, int r, int L, int R, int index){
    if(1 == L \&\& r == R){
        return sealindexl.sum:
    int M = (L + R) / 2;
    push(index);
    pull(index):
    if(r <= M){
        return Query(1, r, L, M, IL(index));
    }else if(1 > M){
        return Query(1, r, M + 1, R, IR(index));
        return Query(1, M, L, M, IL(index)) +
        Query(M + 1, r, M + 1, R, IR(index));
樹狀樹組:
#define lowbit(x) (x &(-x))
int bit [MAX_N+1]={0};
int i=1
while(i<=n){
           //陣列a存放原始資料
           int ans=0;
           for(int ii=i ; ii>=(i-lowbit(i)+1); ){
                      if ( bit[ii] != 0){
                                 ans+=bit[ii];
                                 ii-=lowbit(ii);
                      }
                      else{
                                 ans+=a[ii];
                                 ii--
                      }
           bit[i]=ans;
}
```

```
long long query_sum(int x){
          long long ans=0;
           for(;x;x-=lowbit(x)) ans+=bit[x];
           return ans:
void add(int x.int v){
           for(:x<=N:x+=lowbit(x)) bit[x]+=v: //N應該是元素個數
Sparse Table
int n:
int v[1000009];
int sparse[22][1000009]:
// O(nlogn) preprocess O(1)Query
// sp[x][y] is the answer from (v[x], v[x+2^y-1])
inline void init()
    for (int i = 0; i < n; ++i)
        sparse[0][i] = v[i];
    for (int j = 1; (1 << j) <= n; ++j)
        for (int i = 0; i + (1 << j) <= n; ++i)
        sparse[j][i] = min(
        sparse[j - 1][i],
        sparse[j - 1][i + (1 << (j - 1)])
inline int query(int 1, int r)
    int k = _lg(r - 1 + 1);
    return min(sparse[k][1], sparse[k][r - (1 << k) + 1]);
Graph
Minimum Spanning Tree
           Kruskal
struct DSU // disjoint set no rank-comp-merge
    vector(int) fa:
    DSU(int n) : fa(n) { iota(fa.begin(), fa.end(), 0); } // auto fill fa
    int find(int x) { return fa[x] == x ? x : fa[x] = find(fa[x]); }
    void merge(int x, int y) { fa[find(x)] = find(y); }
int kruskal(int V, vector<tuple<int, int, int>> E) // save all edges into
E, instead of saving graph via adjacency list
    sort(E.begin(), E.end());
   DSU dsu(V);
    int mcnt = 0:
    int ans = 0;
    for (auto e : E)
        int w, u, v; // w for start, u for des, v for val
        tie(w, u, v) = E;
        if (dsu.find(u) == dsu.find(v))
            continue:
        dsu.merge(u, v);
        ans += w;
        if (++mcnt == V - 1)
            break:
    return ans;
}
    Prim
#define enp pair<int, int> // pair<edge_val, node>
int prim pg(vector<vector<enp>> E){
    vector<bool> vis;
    vis.resize(E.size(), false);
    vis[0] = true;
    priority queue<enp> pq;
    for(auto e: E[0]){
        pq.emplace(-e.first, e.second);
```

```
Shortest Path Bellman-Ford
#include <iostream>
#include <vector>
#include <list>
#include <utility:
                            // for std::pair<>
#include <iomanip>
                           // for std::setw()
const int Max_Distance = 100;
class Graph_SP{
                            // SP serves as Shortest Path
private:
    std::vector<std::list<std::pair<int.int>>> AdiList:
    std::vector<int> predecessor, distance;
public:
    Graph_SP():num_vertex(0){};
   Graph SP(int n):num vertex(n){
       AdiList.resize(num vertex):
   void AddEdge(int from, int to, int weight);
   void PrintDataArray(std::vector<int> array);
    void InitializeSingleSource(int Start):
                                               // 以Start作為起點
                                               // 對edge(X,Y)進行Relax
   void Relax(int X, int Y, int weight);
   bool BellmanFord(int Start = 0);
                                               // 以Start作為起點
                                               // if there is negative
cycle, return false
bool Graph_SP::BellmanFord(int Start){
   InitializeSingleSource(Start):
   for (int i = 0; i < num_vertex-1; i++) {</pre>
                                                           // |V-1|次的
iteration
        // for each edge belonging to E(G)
        for (int j = 0 ; j < num_vertex; j++) {</pre>
                                                           // 把AdjList最
外層的vector走一遍
            for (std::list<std::pair<int,int> >::iterator itr =
AdiList[i].begin():
                itr != AdjList[j].end(); itr++) {
                                                           // 各個vector
中、所有edge走一遍
               Relax(j, (*itr).first, (*itr).second);
   // check if there is negative cycle
   for (int i = 0; i < num vertex; i++) {
        for (std::list<std::pair<int,int> >::iterator itr =
AdjList[i].begin();
            itr != AdjList[i].end(); itr++) {
            if (distance[(*itr).first] > distance[i]+(*itr).second) { //
i是from, *itr是to
               return false;
   return true;
```

```
void Graph SP::InitializeSingleSource(int Start){
    distance.resize(num vertex);
    predecessor.resize(num_vertex);
    for (int i = 0; i < num \ vertex; i++) {
        distance[i] = Max_Distance;
        predecessor[i] = -1;
    distance[Start] = 0:
void Graph_SP::Relax(int from, int to, int weight){
    if (distance[to] > distance[from] + weight) {
        distance[to] = distance[from] + weight;
        predecessor[to] = from;
void Graph SP::AddEdge(int from, int to, int weight){
    AdjList[from].push_back(std::make_pair(to,weight));
Diikstra
vector<vector<con>> Graph: //
vector<int> dis; // distance;
priority_queue<con> pq;
    pq.emplace(con(dis[0] = 0, 0));
    while(pq.size()){
        con cur = pq.top();
        pq.pop();
        for(auto it: Graph[cur.second]){
            if(dis[it.second] != it.first) continue;
            dis[it.second] = min(dis[it.second], cur.first + it.first);
            pq.push(con(dis[it.second], it.second));
Flovd-Warshall
// all pairs shortest path
#include <iostream>
#include <vector>
#include <iomanip>
                        // for setw()
const int MaxDistance = 1000:
class Graph_SP_AllPairs{
private:
    int num_vertex;
    std::vector< std::vector<int> > AdiMatrix, Distance, Predecessor;
public:
    Graph_SP_AllPairs():num_vertex(0){};
    Graph_SP_AllPairs(int n);
    void AddEdge(int from, int to, int weight);
    void PrintData(std::vector< std::vector<int> > array);
    void InitializeData():
    void FloydWarshall();
};
Graph_SP_AllPairs::Graph_SP_AllPairs(int n):num_vertex(n){
    // Constructor, initialize AdjMatrix with 0 or MaxDistance
    AdjMatrix.resize(num_vertex);
    for (int i = 0; i < num vertex; i++) {
        AdjMatrix[i].resize(num_vertex, MaxDistance);
        for (int j = 0; j < num\_vertex; j++) {
            if (i == j){}
                AdjMatrix[i][j] = 0;
void Graph SP AllPairs::InitializeData(){
```

Distance.resize(num vertex);

Predecessor.resize(num vertex);

```
for (int i = 0; i < num_vertex; i++) {</pre>
        Distance[i].resize(num vertex):
        Predecessor[i].resize(num_vertex, -1);
        for (int j = 0; j < num_vertex; j++) {
            Distance[i][j] = AdjMatrix[i][j];
              if (Distance[i][j] != 0 && Distance[i][j] != MaxDistance) {
                Predecessor[i][i] = i:
    }
void Graph_SP_AllPairs::FloydWarshall(){
    InitializeData();
    std::cout << "initial Distance[]:\n";</pre>
    PrintData(Distance):
    std::cout << "\ninitial Predecessor[]:\n";</pre>
    PrintData(Predecessor):
    for (int k = 0; k < num_vertex; k++) {</pre>
        std::cout << "\nincluding vertex(" << k << "):\n";</pre>
        for (int i = 0; i < num_vertex; i++) {</pre>
            for (int j = 0; j < num_vertex; j++) {</pre>
                if ((Distance[i][j] > Distance[i][k]+Distance[k][j])
                      && (Distance[i][k] != MaxDistance)) {
                     Distance[i][j] = Distance[i][k]+Distance[k][j];
                     Predecessor[i][j] = Predecessor[k][j];
        // print data after including new vertex and updating the shortest
paths
        std::cout << "Distance[]:\n":</pre>
        PrintData(Distance);
        std::cout << "\nPredecessor[]:\n";
        PrintData(Predecessor):
void Graph_SP_AllPairs::PrintData(std::vector< std::vector<int> > array){
    for (int i = 0; i < num_vertex; i++){</pre>
        for (int j = 0; j < num_vertex; j++) {</pre>
            std::cout << std::setw(5) << array[i][j];
        std::cout << std::endl:
void Graph_SP_AllPairs::AddEdge(int from, int to, int weight){
    AdjMatrix[from][to] = weight;
Lowest Common Ancestor
#define MAXN 200005
#define MAXLOG 200
int D[MAXN]:
int P[MAXLOG][MAXLOG];
#include <cmath>
#include <algorithm>
using namespace std:
#define MAXN 200005
#define MAXLOG 200
int N = MAXN;
int lgN = log(N) / log(2);
int D[MAXN];
int P[MAXLOG][MAXLOG];
int LCA(int u, int v)
    if (D[u] \rightarrow D[v])
        swap(u, v):
    int s = D[v] - D[u]; // adjust D until D[v] = D[u]
    for (int i = 0; i <= lgN; ++i)
        if (s & (1 << i))
            v = P[v][i];
    if (u == v)
        return v;
```

```
// because they are at same depth
    // jump up if they are different
    // think about that if P[u][i] == P[v][i]
   // then that point must be the ancestor of LCA or LCA itself
   // by this, we will stop at LCA's child
    for (int i = lgN; i >= 0; --i)
       if (P[u][i] != P[v][i])
           u = P[u][i];
           v = P[v][i];
    return P[u][0];
void ComputeP()
   int n = N:
   for (int i = 0; i < lgN; ++i) // to lgN enough
       for (int x = 0; x < n; ++x)
           if (P[x][i] == -1)
               P[x][i + 1] = -1;
            else
               P[x][i+1] = P[P[x][i]][i]; // equal to move on the
parent direction
               // And P[x][i] move 2 ^ n steps to a parent we call it v
               // P[y][i] means continue move 2 ^ n step from y to a
parent we call z
                // so the total equal to move 2 ^ n * 2 ^ n steps from x
to z
                // which is move 2 ^{\circ} (n + 1) steps to z
TREE-Centroid
  int subTsize[200005]:
vector<int> adi[200005]:
int n; // n for node num ??
pair<int, int> Tree_Centroid(int v, int pa)
    // return (最大子樹節點數,節點ID)
   subTsize[v] = 1:
    pair<int, int> res(INT_MAX, -1); // ans: tree cnetroid
    int max_subT = 0; // 最大子樹節點數
    for (size_t i = 0; i < adj[v].size(); ++i)
       int x = adj[v][i];
       if (x == pa)
           continué:
       res = min(res, Tree_Centroid(x, v));
       subTsize[v] += subTsize[x];
       max_subT = max(max_subT, subTsize[x]);
   res = min(res, make_pair(max(max_subT, n - subTsize[v]), v)); // (n -
subTsize[v]) for maybe parent tree is the biggest
   // min because all res will be greater than n/2;
```

Kosaraju_for_SCC class Kosaraju_for_SCC(int NodeNum; vector<vector<int>> G; vector<vector<int>> GT; stack<int> st; vector<bool> visited; vector<int>> c; int sccID; public: void init(int N){

return res:

// the min one is the tree centroid

```
NodeNum = N:
        G.clear();
        G.resize(N + 5);
        GT.clear():
        GT.resize(N + 5);
        while(!st.emptv())
            st.pop():
        visited.clear():
        visited.resize(N + 5, false);
        scc.clear();
        scc.resize(N + 5);
        sccID = 1:
    void addEdge(int w, int v){
        G[w].emplace back(v):
        GT[v].emplace_back(w);
    void DFS(bool isG, int v, int k = -1){
        visited[v] = true;
        scc[v] = k;
        vector<vector<int>> &dG = (isG ? G : GT);
        for(int w: dG[v])
            if(!visited[w]){
                DFS(isG, w, k);
        if(isG){
            st.push(v);
    void Kosaraju(int N){
        visited.clear();
        visited.resize(N + 5, false):
        for (int i = 1; i <= N; i++){
            if(!visited[i])
                DFS(true, i):
        visited.clear();
        visited.resize(N + 5, false);
        while(!st.empty()){
            if(!visited[st.top()])
                DFS(false, st.top(), sccID++);
            st.pop();
    vector<vector<int>> generateReG(){
        vector<vector<int>> reG;
        reG.resize(sccID);
        for (int i = 1; i <= NodeNum; i++){
            for(int w: G[i]){
               if(scc[i] == scc[w])
                   continue:
               reG[scc[i]].emplace_back(scc[w]);
        return reG;
};
```

Tarjan for SCC class tarjan for SCC{ private: vector<vector<int>> G; // adjacency list vector<int> D: vector<int> L; vector<int> sccID: stack<int> st; // for SccID vector<bool> inSt; vector<vector<int>> reG; int timeStamp, sccIDstamp; public: void init(int size = 1){ G.clear(); G.resize(size + 3); D.clear();

```
D.resize(size + 3, 0);
        L.clear();
       L.resize(size + 3, 0);
        sccID.clear():
        sccID.resize(size + 3, 0);
        while(!st.empty())
            st.pop():
        inSt.clear():
        inSt.resize(size + 3, false);
        reG.clear();
        sccIDstamp = timeStamp = 1:
    void addEdge(int from, int to){
        G[from].emplace_back(to);
    void DFS(int v, int pa){ //call DFS(v,v) at first
        D[v] = L[v] = timeStamp++; //timestamp > 0
        st.push(v);
        inSt[v] = true;
        for(int w: G[v]){ // directed graph don't need w == pa
            if(!D[w]){ // D[w] = 0 if not visited}
                DFS(w, v);
                L[v] = min(L[v], L[w]);
            }else if(inSt[w])
            { /* w has been visited.
                if we don't add this, the L[v] will think that v can back
to node whose index less to v.
                !inSt[w] is true that v -> w is a forward edge
                opposite it's a cross edge
                L[v] = min(L[v], D[w]); // why D[w] instead of L[w]??
        if(D[v] == L[v]){
           int w:
            do{
                w = st.top();
                st.pop();
                sccID[w] = sccIDstamp; // scc ID for this pooint at which
SCC
                inSt[w] = false;
           } while (w != v);
            sccIDstamp++;
    void generateReG(int N = 1){
        reG.clear();
        reG.resize(sccIDstamp):
        for (int i = 1; i <= N; i++){
            for(int w: G[i]){
                if(sccID[i] == sccID[w])
                    continue:
                reG[sccID[i]].emplace_back(sccID[w]);
    bool visited(int v){
        return D[v];
};
           Tarjan for ArticulationPointBridge
```

Tarjan for ArticulationPointBridge class tarjan{ vector<vector<int>> G; // adjacency List vector(int> D; // visit or visited and D-value vector(int> L; // for L-value vector(one edgeBridge; vector(int> APnode; int timestamp; tarjan(int size = 1){ timestamp = 0; G.resize(size); D.resize(size, 0); ledgeBridge.clear();

```
APnode.clear();
    void init(int size = 1){
       tarjan(size);
   void addedge(int u, int v)
    { // undirected graph
       G[u].push_back(v);
       G[v].push_back(u);
    void DFS(int v. int pa)
   { // 使用 DFS(v,v) 來呼叫函數
       D[v] = L[v] = timestamp++;
       int Childcount = 0;
       bool isAP = false:
       for(int w: G[v]){
           if(w == pa)
               continue:
           if(!D[w])
           { // 用 D[w] = 0 當作沒走過
               DFS(w, v);
               Childcount++:
               if(D[v] <= L[w])
                   isAP = true; // 結 論 2 對於除了 r 點以外的所有點 v.v 點
在 G 上為 AP 的充要條件為其在 T 中至少有一個子節點 w 滿足 D(v) ≤ L(w)
               if(D[v] < L[w])
                   edgeBridge.emplace_back(v,w);// 結 論 3 對於包含 r 在內
的所有點 v 和 v 在 T 中的子節點 w, 邊 e(v, w) 在圖 G 中為bridge 的充要條件為
D(v) < L(w)_{\circ}
               L[v] = min(L[v], L[w]);
           L[v] = min(L[v], D[w]);
       if(v == pa && Childcount < 2)
           isAP = false:
       if(isAP)
           APnode.emplace back(v):
};
          Tarjan for BridgeCC
int timestamp = 1;
int bccid = 1:
int D[MAX_N];
int L[MAX N];
int bcc[MAX_N];
stack<int> st:
vector<int> adi[MAX N]:
bool inSt[MAX_N];
void DFS(int v, int fa) { //call DFS(v,v) at first
   D[v] = L[v] = timestamp++; //timestamp > 0
   st.emplace(v):
    for (int w:adj[v]) {
       if( w==fa ) continue;
       if (!D[w]) \{ // D[w] = 0 \text{ if not visited}
           L[v] = min(L[v], L[w]);
       L[v] = min(L[v], D[w]);
    if (L[v]==D[v]) {
       bccid++;
       int x;
           x = st.top(); st.pop();
           bcc[x] = bccid;
       } while (x!=v);
   return ;
FLOW
          Ford Fulkerson
```

Ford_Fulkerso
// O((V+E)F)

```
#define mayn 101
// remember to change used into the maxNode size -- kattis elementary math
bool used[maxn];
int End:
vector<int> V[maxn];
vector<tuple<int. int>> E:
// x=>y 可以流 C
// if undirected or 2-direc edge, bakcward Capacity become C;
// Graph build by edge array
// 反向邊的編號只要把自己的編號 xor 1 就能取得
void add_edge(int x, int y,int c)
    V[x].emplace_back( E.size() );
   E.emplace_back(y,c);
V[y].emplace_back( E.size() );
    E.emplace_back(x,0);
int dfs(int v. int f)
    if( v==End ) return f;
    used[v] = true;
    int e,w;
    for( int eid : V[v] )
        tie(e,w) = E[eid];
        if( used[e] || w==0 ) continue;
        w = dfs(e, min(w,f));
        if( w>0 )
            // 更新流量
            get<1>(E[eid ]) -= w;
            get<1>(E[eid^1]) += w;
            return w;
    return 0;// Fail!
int ffa(int s,int e)
    int ans = 0, f;
    End = e:
    while(true)
        memset(used, 0, sizeof(used));
        f = dfs(s, INT_MAX);
        if( f<=0 ) break;
        ans += f:
    return ans;
Dinic's Algorithm
// O(V^2E) O(VE) finding argument path
// if unit capacity networe then O(min(V^3/2, E^1/2) E)
// solving bipartite matching O(E V^1/2)
#define maxn 101
#define INT MAX 10000000
int End, dist[maxn];
vector<tuple<int, int, int>> V[maxn];
// 1st for node-index, 2nd for cap, 3nd for the index of the edge and the
vector[u];
void addEdge(int u, int v, int c){
    V[u].emplace_back(v, c, V[v].size());
    V[v].emplace_back(u, 0, V[u].size() - 1);
bool bfs(int s) {
    memset(dist, -1, sizeof(dist));
    queue<int> qu;
    qu.emplace(s);
    dist[s]=0;
    while( !qu.empty() ) {
        int S = qu.front(); qu.pop();
        for(auto &p : V[S]) {
```

```
int E, C;
            tie(E, C, ignore) = p;
            if( dist[E]==-1 && C!=0 ) {
    dist[E]=dist[S]+1;
                qu.emplace(E);
    return dist[End] != -1:
int dfs(int v, int f) {
    int e,w,rev;
    if( v==End || f==0 ) return f:
    for( auto &t : V[v] )
        tie(e,w,rev) = t;
        if( dist[e]!=dist[v]+1 || w==0 )
            continue:
        w = dfs(e, min(w,f));
        if( w>0 )
            get<1>(t) -= w:
            get<1>(V[e][rev]) += w;
            return w;
    dist[v] = -1; //優化. 這個點沒用了
    return 0:// Fail!
int dinic(int s,int e)
    int ans = 0, f;
    End = e:
    while(bfs(s))
        while( f = dfs(s, INT MAX) )
            ans += f;
    return ans;
```

MinCost MaxFlow // by finkela

```
Bipartite Matching
konig' algorithm
#include <vector>
#include <cstring>
using namespace std:
vector<int> V[205]:
// V[i]記錄了左半邊可以配到右邊的那些點
int match[205]; // A<=B
// match[i] 記錄了右半邊配對到左半邊的哪個點
bool used[205]:
int n:
bool dfs(int v)
   for(int e:V[v])
       if( used[e] ) continue;
       used[e] = true;
       if( match[e] == -1 || dfs( match[e] ) )
           match[e] = v;
           return true;
    return false;
int konig()
   memset(match,-1,sizeof(match));
```

```
int ans=0;
for(int i=1;i<=n;++i)
{
    memset(used, 0, sizeof(used));
    if( dfs(i) )
        ans++;
}
return ans;
}</pre>
```

KM algorithm

// Max weight perfect bipartite matching // $O(V^3)$ // by jinkela

Graph Matching(untest!!) blossom algorithm // by jinkela // 最大圖匹配// O(V²(V+E))

```
最近點對
template <typename T>
T ClosestPairSquareDistance(typename vector<Point<T>>::iterator 1,
                           typename vector<Point<T>>::iterator r)
   auto delta = numeric_limits<T>::max();
   if (r - 1 > 1)
       auto m = 1 + (r - 1 >> 1):
       nth_element(1, m, r); // Lexicographical order in default
       delta = min(ClosestPairSquareDistance<T>(1, m),
                   ClosestPairSquareDistance<T>(m, r));
       auto square = [&](T y) { return y * y; };
       auto sgn = [=](T a, T b) {
           return square(a - b) <= delta ? 0 : a < b ? -1 : 1;
       vector<Point<T>> x_near[2];
       copy_if(1, m, back_inserter(x_near[0]), [=](Point<T> a) {
           return sgn(a.x, x) == 0;
       copy_if(m, r, back_inserter(x_near[1]), [=](Point<T> a) {
            return sgn(a.x, x) == 0;
       for (int i = 0, j = 0; i < x_near[0].size(); ++i)
            while (j < x_near[1].size() and
                  sgn(x_near[1][j].y, x_near[0][i].y) == -1)
            for (int k = j; k < x_near[1].size() and
                           sgn(x_near[1][k].y, x_near[0][i].y) == 0;
                ++k)
               delta = min(delta, (x_near[0][i] - x_near[1][k]).norm());
       inplace_merge(1, m, r, [](Point<T> a, Point<T> b) {
           return a.y < b.y;
       });
   return delta:
```

```
if (ch[i])
                     ch[i]->~KDNode();
         T dfs(const Point<T> a.
               T dis = numeric limits<T>::max().
               bool parity = 0)
             dis = min(dis, (v - q).norm());
             bool isRight = parity ? v.x < q.x : v.y < q.y;</pre>
             if (ch[isRight])
                dis = min(dis, ch[isRight]->dfs(q, dis, parity ^ 1));
             if (ch[isRight ^ 1] and [](T x) {
                     return x * x;
                (parity ? v.x - q.x : v.y - q.y) < dis)
                dis = min(dis, ch[isRight ^ 1]->dfs(q, dis, parity ^ 1));
             return dis:
     } * root:
    KDNode *buildKDTree(typename vector<Point<T>>::iterator 1,
                         typename vector<Point<T>>::iterator r,
                         bool parity = 0)
         if (r == 1)
             return nullptr;
         auto m = 1 + (r - 1 >> 1);
         nth_element(1, m, r, [=](Point<T> a, Point<T> b) {
            return parity ? a.x < b.x : a.y < b.y;
         });
         return new KDNode(*m.
                           buildKDTree(l, m, parity ^ 1),
                           buildKDTree(m + 1, r, parity ^ 1));
public:
    KDTree(vector<Point<T>> A) : root{buildKDTree(A.begin(), A.end())} {}
    T nearestNeighborSquareDistance(Point<T> q)
         return root->dfs(q);
};
nCr
using i64 = unsigned long long;
#define maxn 300005
i64 fact[maxn], tcaf[maxn]:
#define P 998244353
#define REP1(i, n) for (int i = 1; i \leftarrow (int)(n); ++i)
#define REP(i, n) for (int i = (int)(n) - 1; i >= 0; --i)
void init(int n){
    fact[0] = 1;
    for (int i = 1; i <= n; i++)
         fact[i] = i * fact[i - 1] % P;
    for (int i = n; i >= 0; --i)
         tcaf[i] = deg(fact[i], -1);
}
i64 deg(i64 x, i64 d) {
    if (d < 0) d += P - 1;
    i64 y = 1;
    while (d) {
        if (d & 1) (y *= x) %= P;
         d /= 2;
         (x *= x) %= P;
     return y;
}
i64 cnk(int n, int k) {
    if (k < 0 \mid | k > n) return 0;
```

```
return fact[n] * tcaf[k] % P * tcaf[n - k] % P;
}
some theorem
Maximum Independent Set
General: [NPC] maximum clique of complement of G
Tree: [P] Greedy
```

Bipartite Graph: [P] Maximum Cardinality Bipartite Matching

Minimum Dominating Set

General: [NPC] Tree: [P] DP

Bipartite Graph: [NPC]

Minimum Vertex Cover General: [NPC] V - MIS

Tree: [P] Greedy, from leaf to root

Bipartite Graph: [P] Maximum Cardinality Bipartite

Matching

Minimum Edge Cover

General: [P] V - Maximum Matching

Bipartite Graph: [P] Greedy, strategy: cover small

degree node first.

(Min/Max)Weighted: [P]: Minimum/Minimum Weight Matching

Pick's Theorem A = i + b/2 - 1

我们称Y中所有边的两个端点为被Y所匹配。我们遍历所有R中没有被Y匹配的顶点,寻找所有长度为偶数的路径,路径中匹配边和未匹配边交替出现。事实上不存在长度为奇数的路径,不然我们就找到了一个增广路径。我们将L中被标记过的顶点和R中未被标记的顶点合成一个新的点集合X,我们接下来证明X是最小顶点覆盖。

9.8 Formulas or Theorems

We have encountered some rarely used formulas or theorems in some programming contest problems before. Knowing them will give you an *unfair advantage* over other contestants if one of these rare formulas or theorems is used in the programming contest that you join.

- 1. Cayley's Formula: There are n^{n-2} spanning trees of a complete graph with n labeled vertices. Example: UVa 10843 Anne's game.
- 2. Derangement: A permutation of the elements of a set such that none of the elements appear in their original position. The number of derangements der(n) can be computed as follow: $der(n) = (n-1) \times (der(n-1) + der(n-2))$ where der(0) = 1 and der(1) = 0. A basic problem involving derangement is UVa 12024 Hats (see Section 5.6).
- 3. Erdős Gallai's Theorem gives a necessary and sufficient condition for a finite sequence of natural numbers to be the degree sequence of a simple graph. A sequence of nonnegative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be the degree sequence of a simple graph on n vertices iff $\sum_{i=1}^n d_i$ is even and $\sum_{i=1}^k d_i \leq k \times (k-1) + \sum_{i=k+1}^n \min(d_i,k)$ holds for $1 \leq k \leq n$. Example: UVa 10720 Graph Construction.
- Euler's Formula for Planar Graph⁶: V E + F = 2, where F is the number of faces⁷ of the Planar Graph. Example: UVa 10178 Count the Faces.
- 5. Moser's Circle: Determine the number of pieces into which a circle is divided if n points on its circumference are joined by chords with no three internally concurrent. Solution: q(n) = n C₄ + n C₂ + 1. Example: UVa 10213 How Many Pieces of Land?
- 6. Pick's Theorem⁸: Let I be the number of integer points in the polygon, A be the area of the polygon, and b be the number of integer points on the boundary, then $A = i + \frac{b}{2} 1$. Example: UVa 10088 Trees on My Island.
- 7. The number of spanning tree of a complete bipartite graph $K_{n,m}$ is $m^{n-1} \times n^{m-1}$. Example: UVa 11719 Gridlands Airport.