

## Problem.1

You are given an integer array **nums** and an integer **k**. Determine whether there exists a contiguous subarray of length exactly **k** that contains at least two equal elements. If such a subarray exists, return true; otherwise, return false.

Example 1:

**Input:** nums = [1,2,3,2,4,5], k = 3

**Output:** true

**Explanation:** The subarray [2,3,2] (indices 1–3) contains the duplicate element 2

Example 2:

**Input:** nums = [1,2,3,4,5,6], k = 3

**Output:** false

**Explanation:** All contiguous subarrays of length 3 contain distinct elements.

## Problem.2

Given an integer array **nums** and two integers **k** and **t**, determine whether there exist two distinct indices **i** and **j** such that:

$$\text{abs}(i - j) \leq k$$

$$\text{abs}(\text{nums}[i] - \text{nums}[j]) \leq t$$

Example 1:

**Input:** nums = [1,5,9,3,1], k = 2, t = 2

**Output:** true

**Explanation:** Indices 1 and 3:  $\text{abs}(1-3) = 2 \leq 2$ ,  $\text{abs}(5-3) = 2 \leq 2$

Example 2:

**Input:** nums = [1,10,20,30,40], k = 2, t = 5

**Output:** false

**Explanation:** Within any sliding window of length 3, the difference between any two elements is greater than 5.

**Tips:**

$1 \leq \text{nums.length} \leq 10^4$

$-10^9 \leq \text{nums}[i], \text{nums}[j] \leq 10^9$

$0 \leq k \leq 10000$

$0 \leq t \leq 10^9$

### Problem.3

You are given an integer array **nums** and two integers **k** and **f**. Determine whether there exists a contiguous subarray of length at most **k** in which some element appears at least **f** times.

Example 1:

**Input:** `nums = [1,2,2,3,2,4]`, `k = 4`, `f = 3`

**Output:** true

**Explanation:** The contiguous subarray `[2,2,3,2]` (indices 1–4, length 4) contains the element 2 appearing 3 times.

Example 2:

**Input:** `nums = [1,1,1,2,2,2]`, `k = 2`, `f = 3`

**Output:** false

**Explanation:** In any contiguous subarray of length  $\leq 2$ , the maximum frequency of any element is 2.

**Tips:**

$1 \leq \text{nums.length} \leq 10^4$

$-10^9 \leq \text{nums}[i] \leq 10^9$

$1 \leq k \leq \text{nums.length}$

$2 \leq f \leq k$

## Problem 4

The Catalan numbers are a sequence of natural numbers that appear in various counting problems in combinatorics.

They can be defined recursively as:

$$C_0 = 1, \quad C_{n+1} = \sum_{i=0}^n C_i \times C_{n-i}$$

### Tasks

For each implementation below:

1. **Correctness** – Is the implementation correct?

If not, what needs to be fixed?

2. **Characteristics and Use Case** –

Describe the characteristics of this implementation, its advantages, and the types of situations it is most suitable for.

3. **New Knowledge** –

Identify any C/C++ programming syntax or algorithmic details you were not familiar with before, and briefly explain them.

4. **Code Style** –

Add comments to each function following the **Google C++ Coding Style**.

5. **Your Version** –

After analyzing all ten versions, write your own implementation of the Catalan number function that you consider the **most elegant and effective**.

## 10 Implementations of the Catalan Number

1	<pre>// Version 1: Simple recursive implementation int catalan_recursive(int n) {     if (n &lt;= 1) return 1;     int res = 0;     for (int i = 0; i &lt; n; ++i) {         res += catalan_recursive(i) * catalan_recursive(n - 1 - i);     }     return res; }</pre>
2	<pre>// Version 2: Iterative dynamic programming #include &lt;vector&gt; int catalan_iterative(int n) {     if (n &lt;= 1) return 1;     std::vector&lt;int&gt; dp(n + 1, 0);     dp[0] = dp[1] = 1;     for (int i = 2; i &lt;= n; ++i) {         for (int j = 0; j &lt; i; ++j)             dp[i] += dp[j] * dp[i - 1 - j];     }     return dp[n]; }</pre>
3	<pre>// Version 3: Recursive implementation with memoization #include &lt;unordered_map&gt;  long long catalan_memoization(int n, std::unordered_map&lt;int, long long&gt;&amp; memo) {     if (n &lt;= 1) return 1;     if (memo.find(n) != memo.end()) return memo[n];     long long res = 0;     for (int i = 0; i &lt; n; ++i)         res += catalan_memoization(i, memo) * catalan_memoization(n - 1 - i, memo);     return memo[n] = res; }</pre>
4	<pre>// Version 4: Using Boost multiprecision for large n #include &lt;boost/multiprecision/cpp_int.hpp&gt; #include &lt;vector&gt;  boost::multiprecision::cpp_int catalan_bigint(int n) {     using boost::multiprecision::cpp_int;     if (n &lt;= 1) return 1;     std::vector&lt;cpp_int&gt; dp(n + 1);     dp[0] = dp[1] = 1;</pre>

	<pre> for (int i = 2; i &lt;= n; ++i) {     dp[i] = 0;     for (int j = 0; j &lt; i; ++j)         dp[i] += dp[j] * dp[i - 1 - j];     } return dp[n]; } </pre>
5	<pre> // Version 5: Using factorial and combinatorial formula #include &lt;cmath&gt;  long long factorial(int n) {     long long res = 1;     for (int i = 2; i &lt;= n; ++i) res *= i;     return res; }  long long catalan_formula(int n) {     if (n &lt;= 1) return 1;     return factorial(2 * n) / (factorial(n + 1) * factorial(n)); } </pre>
6	<pre> // Version 6: Compile-time computation using constexpr constexpr long long catalan_constexpr(int n) {     if (n &lt;= 1) return 1;     long long dp[64] = {0}; // supports n &lt;= 63     dp[0] = dp[1] = 1;     for (int i = 2; i &lt;= n; ++i) {         dp[i] = 0;         for (int j = 0; j &lt; i; ++j)             dp[i] += dp[j] * dp[i - 1 - j];         }     return dp[n]; } </pre>
7	<pre> // Version 7: Parallel recursive computation using async #include &lt;future&gt;  long long catalan_parallel(int n) {     if (n &lt;= 1) return 1;     std::vector&lt;std::future&lt;long long&gt;&gt; futures;     long long res = 0;     for (int i = 0; i &lt; n; ++i) {         futures.push_back(std::async(std::launch::async, catalan_parallel, i));     } } </pre>

	<pre>     }     for (int i = 0; i &lt; n; ++i) {         res += futures[i].get() * catalan_parallel(n - 1 - i);     }     return res; } </pre>
8	<pre> // Version 8: Safe implementation using std::optional #include &lt;optional&gt; #include &lt;vector&gt;  std::optional&lt;long long&gt; catalan_safe(int n) {     if (n &lt; 0) return std::nullopt;     if (n &lt;= 1) return 1;     std::vector&lt;long long&gt; dp(n + 1);     dp[0] = dp[1] = 1;     for (int i = 2; i &lt;= n; ++i)         for (int j = 0; j &lt; i; ++j)             dp[i] += dp[j] * dp[i - 1 - j];     return dp[n]; } </pre>
9	<pre> // Version 9: Template metaprogramming version template&lt;int N&gt; struct Catalan {     static constexpr long long value = ([]() constexpr {         long long sum = 0;         for (int i = 0; i &lt; N; ++i)             sum += Catalan&lt;i&gt;::value * Catalan&lt;N - 1 - i&gt;::value;         return sum;     })(); };  template&lt;&gt; struct Catalan&lt;0&gt; { static constexpr long long value = 1; }; template&lt;&gt; struct Catalan&lt;1&gt; { static constexpr long long value = 1; }; </pre>
10	<pre> // Version 10: Conceptual matrix-based recurrence demonstration #include &lt;array&gt;  using Matrix2x2 = std::array&lt;std::array&lt;long long, 2&gt;, 2&gt;;  Matrix2x2 matrix_multiply(const Matrix2x2&amp; a, const Matrix2x2&amp; b) {     Matrix2x2 result = {{ {0, 0}, {0, 0} }};     for (int i = 0; i &lt; 2; ++i) </pre>

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        for (int j = 0; j < 2; ++j)
            for (int k = 0; k < 2; ++k)
                result[i][j] += a[i][k] * b[k][j];
    return result;
}

long long catalan_matrix(int n) {
    if (n <= 1) return 1;
    // This version illustrates the concept of matrix recurrence
    Matrix2x2 base = {{ {1, 1}, {1, 0} }};
    Matrix2x2 result = {{ {1, 0}, {0, 1} }};
    for (int i = 0; i < n; ++i)
        result = matrix_multiply(result, base);
    return result[0][0];
}
```