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Basic

1.1 default code

```
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
using namespace std;
ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
```

1.2 .vimrc

```
set nu rnu ts=4 sw=4 bs=2 ai hls cin mouse=a
color default
sv on
inoremap {<CR>} {<CR>} {<C>>}
inoremap jk <Esc>
nnoremap J 5j
nnoremap K 5k
nnoremap run :w<br/>-bar>!g++ -std=c++14 -DLOCAL -Wfatal-errors -o test "%" && echo "done." && time ./test<
```

1.3 Increase Stack Size (linux)

```
#include <sys/resource.h>
void increase_stack_size() {
  const rlim_t ks = 64*1024*1024;
  struct rlimit rl;
  int res=getrlimit(RLIMIT_STACK, &rl);
  if(res==0){
     if(rl.rlim_cur<ks){</pre>
       rl.rlim_cur=ks;
       res=setrlimit(RLIMIT_STACK, &rl);
} } }
```

1.4 Misc

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```
編譯參數:-std=c++14 -Wall -Wshadow (-fsanitize=
    undefined)
mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }
#define SECs ((double)clock() / CLOCKS_PER_SEC)
double startTime;
bool TIME() { // 比最大可執行時間小一點
    return SECs - startTime > 0.8;
int main() {
    startTime = SECs;
```

1.5 check

```
for ((i=0;;i++))
do

    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done
```

1.6 python-related

```
parser:
int(eval(num.replace("/","//")))
from fractions import Fraction
from decimal import Decimal, getcontext, ROUND_HALF_UP,
      ROUND_CEILING, ROUND_FLOOR
getcontext().prec = 250 # set precision
getcontext().rounding = ROUND_HALF_UP
itwo = Decimal(0.5)
two = Decimal(2)
format(x, '0.10f') # set precision
N = 200
def angle(cosT):
    """given cos(theta) in decimal return theta"""
  for i in range(N):
  cosT = ((cosT + 1) / two) ** itwo 
 sinT = (1 - cosT * cosT) ** itwo 
 return sinT * (2 ** N)
pi = angle(Decimal(-1))
"""round to 2 decimal places"""
sum = Decimal(input())
sum.quantize(Decimal('.00'), ROUND_HALF_UP)
"""Fraction"""
x = Fraction(1, 3) # 1/3
x.as_integer_ratio() # (1, 3)
"""input list of integers"""
arr = list(map(int, input().split()))
```

2 flow

2.1 ISAP $O(V^3)$

```
struct Maxflow {
    static const int MAXV = 20010;
    static const int INF = 1000000;
    struct Edge {
        int v, c, r;
        Edge(int _v, int _c, int _r):
            v(_v), c(_c), r(_r) {}
    };
    int s, t;
    vector<Edge> G[MAXV*2];
    int iter[MAXV*2], d[MAXV*2], gap[MAXV*2], tot;
```

```
void init(int x) {
    tot = x+2;
    s = x+1, t = x+2;
    for(int i = 0; i <= tot; i++) {</pre>
      G[i].clear();
       iter[i] = d[i] = gap[i] = 0;
  void addEdge(int u, int v, int c) {
    G[u].push_back(Edge(v, ć, SZ(G[v]) ));
G[v].push_back(Edge(u, 0, SZ(G[u]) - 1));
  int dfs(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
       Edge &e = G[p][i];
       if(e.c > 0 \& d[p] == d[e.v]+1) {
         int f = dfs(e.v, min(flow, e.c));
         if(f) {
           e.c -= f;
           G[e.v][e.r].c += f;
           return f;
    } } }
    if((--gap[d[p]]) == 0) d[s] = tot;
    else {
      d[p]++;
      iter[p] = 0;
      ++gap[d[p]];
    return 0;
  int solve() {
    int res = 0;
    gap[0] = tot;
    for(res = 0; d[s] < tot; res += dfs(s, INF));</pre>
    return res;
  void reset() {
    for(int i=0;i<=tot;i++) {</pre>
      iter[i]=d[i]=gap[i]=0;
} } flow;
```

2.2 MinCostFlow

```
struct zkwflow{
  static const int maxN=10000;
  struct Edge{ int v,f,re; ll w;};
int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
vector<Edge> E[maxN];
  void init(int _n,int _s,int _t){
    n=_n,s=_s,t=_t;
    for(int i=0;i<n;i++) E[i].clear();</pre>
  void addEdge(int u,int v,int f,ll w){
    E[u].push_back({v,f,(int)E[v].size(),w});
    E[v].push_back(\{u,0,(int)E[u].size()-1,-w\});
  bool SPFA(){
    fill_n(dis,n,LLONG_MAX); fill_n(vis,n,false);
    queue<int> q; q.push(s); dis[s]=0;
    while (!q.empty()){
  int u=q.front(); q.pop(); vis[u]=false;
  for(auto &it:E[u]){
         if(it.f>0&&dis[it.v]>dis[u]+it.w){
           dis[it.v]=dis[u]+it.w;
           if(!vis[it.v]){
             vis[it.v]=true; q.push(it.v);
    return dis[t]!=LLONG_MAX;
  int DFS(int u,int nf){
    if(u==t) return nf;
    int res=0; vis[u]=true;
    for(int &i=ptr[u];i<(int)E[u].size();i++){</pre>
      auto &it=E[u][i]
       if(it.f>0&&dis[it.v]==dis[u]+it.w&&!vis[it.v]){
         int tf=DFS(it.v,min(nf,it.f));
         res+=tf,nf-=tf,it.f-=tf;
         E[it.v][it.re].f+=tf;
         if(nf==0){ vis[u]=false; break; }
```

```
freturn res;

pair<int,ll> flow(){
  int flow=0; ll cost=0;
  while (SPFA()){
    fill_n(ptr,n,0);
    int f=DFS(s,INT_MAX);
    flow+=f; cost+=dis[t]*f;
}
  return{ flow,cost };
} // reset: do nothing
} flow;
```

2.3 Dinic $O(V^2E)$

```
#define SZ(x) (int)x.size()
#define PB push_back
struct Dinic{
  struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
    n = _n;    s = _s;    t = _t;

    for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
    E[u].PB({v,f,SZ(E[v])});
    E[v].PB({u,0,SZ(E[u])-1});
  bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;</pre>
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()){
      int u = que.front(); que.pop();
       for (auto it : E[u]){
         if (it.f > 0 && level[it.v] == -1){
           level[it.v] = level[u]+1;
           que.push(it.v);
    } } }
    return level[t] != -1;
  int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
      if (it.f > 0 && level[it.v] == level[u]+1){
        int tf = DFS(it.v, min(nf,it.f));
         res += tf; nf -= tf; it.f -= tf;
         E[it.v][it.re].f += tf;
         if (nf == 0) return res;
    if (!res) level[u] = -1;
    return res;
  int flow(int res=0){
    while ( BFS() )
      res += DFS(s,2147483647);
    return res;
} }flow;
```

2.4 Kuhn Munkres 最大完美二分匹配 $O(n^3)$

```
struct KM{ // max weight, for min negate the weights
  int n, mx[MXN], my[MXN], pa[MXN];
  ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void init(int _n) { // 1-based
    n = _n;
    for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
  }
  void addEdge(int x, int y, ll w) {g[x][y] = w;}
  void augment(int y) {
    for(int x, z; y; y = z)
        x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
}</pre>
```

```
void bfs(int st) {
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
     queue<int> q; q.push(st);
     for(;;) {
       while(q.size()) {
         int x=q.front(); q.pop(); vx[x]=1;
          for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
            11 t = 1x[x]+1y[y]-g[x][y];
            if(t==0){
              pa[y]=x
              if(!my[y]){augment(y);return;}
              vy[y]=1, q.push(my[y]);
            }else if(sy[y]>t) pa[y]=x,sy[y]=t;
         }
       11 cut = INF;
       for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;</pre>
          if(vy[j]) ly[j] += cut;
          else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
  if(!my[y]){augment(y);return;}</pre>
         vy[y]=1, q.push(my[y]);
    } }
  ll solve(){
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
     fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
    lx[x] = max(lx[x], g[x][y]);</pre>
     for(int x=1; x<=n; ++x) bfs(x);
     ll\ ans = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
     return ans;
} }graph;
```

2.5 SW min-cut (不限 S-T 的 min-cut) $O(V^3)$

```
// global min cut
struct SW{ // O(V^3)
   int n,vst[MXN],del[MXN];
   int edge[MXN][MXN],wei[MXN];
   void init(int _n){
    n = _n; FZ(edge); FZ(del);
   void addEdge(int u, int v, int w){
     edge[u][v] += w; edge[v][u] += w;
   void search(int &s, int &t){
     FZ(vst); FZ(wei);
     s = t = -1;
     while (true){
       int mx=-1, cur=0;
for (int i=0; i<n; i++)
  if (!del[i] && !vst[i] && mx<wei[i])</pre>
            cur = i, mx = wei[i];
       if (mx == -1) break;
       vst[cur] = 1;
       s = t; t = cur;
       for (int i=0; i<n; i++)
  if (!vst[i] && !del[i]) wei[i] += edge[cur][i];</pre>
   int solve(){
     int res = 2147483647;
     for (int i=0,x,y; i<n-1; i++){
       search(x,y);
       res = min(res,wei[y]);
       del[y] = 1;
       for (int j=0; j<n; j++)</pre>
          edge[x][j] = (edge[j][x] += edge[y][j]);
     return res;
} }graph;
```

2.6 Max flow with lower/upper bound

```
// flow use ISAP
// Max flow with lower/upper bound on edges
// source = 1 , sink = n
int in[N], out[N];
int l[M], r[M], a[M], b[M];//0-base,a下界,b
     上界
int solve(){
  flow.init(n); //n為點的數量,m為邊的數量,點是1-
       base
  for( int i = 0 ; i < m ; i ++ ){
  in[ r[ i ] ] += a[ i ];
  out[ l[ i ] ] += a[ i ];</pre>
    flow.addEdge( l[ i ] , r[ i ] , b[ i ] - a[ i ] );
// flow from l[i] to r[i] must in [a[ i ], b[ i ]]
  int nd = 0;
  for( int i = 1 ; i <= n ; i ++ ){
  if( in[ i ] < out[ i ] ){</pre>
      flow.addEdge( i , flow.t , out[ i ] - in[ i ] );
nd += out[ i ] - in[ i ];
    if( out[ i ] < in[ i ] )</pre>
       flow.addEdge( flow.s , i , in[ i ] - out[ i ] );
  // original sink to source
  flow.addEdge( n , 1 , INF );
if( flow.maxflow() != nd )
    return -1; // no solution
  int ans = flow.G[ 1 ].back().c; // source to sink
  flow.G[1].back().c = flow.G[n].back().c = 0;
  // take out super source and super sink
  for( size_t i = 0 ; i < flow.G[ flow.s ].size() ; i</pre>
       ++ ){
     flow.G[flow.s][i].c = 0;
    Edge &e = flow.G[ flow.s ][ i ];
    flow.G[ e.v ][ e.r ].c = 0;
  for( size_t i = 0 ; i < flow.G[ flow.t ].size() ; i</pre>
       ++ ){
     flow.\hat{G}[flow.t][i].c = 0;
    Edge &e = flow.G[ flow.t ]['i ];
    flow.G[ e.v ][ e.r ].c = 0;
  flow.addEdge( flow.s , 1 , INF );
  flow.addEdge( n , flow.t , INF );
  flow.reset();
  return ans + flow.maxflow();
```

2.7 Flow Method

Maximize $c^T x$ subject to $Ax \le b$, $x \ge 0$;

Let S be Sum of all weight(or inf)

1. from source to each node with cap = S

2. For each (u,v,w) in E, (u->v,cap=w), (v->u,cap=w)

```
with the corresponding symmetric dual problem,
Minimize b^T y subject to A^T y \geq c, y \geq 0.
Maximize c^T x subject to Ax \le b;
with the corresponding asymmetric dual problem,
Minimize b^T y subject to A^T y = c, y \ge 0.
Minimum vertex cover on bipartite graph =
Maximum matching on bipartite graph
Minimum edge cover on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
Independent set on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
找出最小點覆蓋、做完dinic之後、從源點dfs只走還有流量的
    邊 · 左邊 沒 被 走 到 的 點 跟 右 邊 被 走 到 的 點 就 是 答 案 · 其 他
    點為最大獨立集
Maximum density subgraph (\sum W_e + \sum W_v) / |V|
Binary search on answer:
For a fixed D, construct a Max flow model as follow:
```

3 Math

3.1 FFT

```
// const int MXN = 262144 (MXN must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx; //real() ,imag()
const ld PI = acosl(-1);
const cplx I(0, 1);
struct FFT{
  cplx omega[MXN+1];
  FFT(){ //pre_fft
    for(int i=0; i<=MXN; i++)
  omega[i] = exp(i * 2 * PI / MXN * I);</pre>
  // n must be 2^k
  void fft(int n, cplx a[], bool inv=false){
     int basic = MXN / n;
     int theta = basic;
     for (int m = n; m >= 2; m >>= 1) {
       int mh = m >> 1;
       for (int i = 0; i < mh; i++) {
       cplx w = omega[inv ? MXN-(i*theta%MXN) : i*theta%
           MXN];
       for (int j = i; j < n; j += m) {
         int k = j + mh;
         cplx x = a[j] - a[k];
         a[j] += a[\bar{k}];
         a[k] = w * x;
       } }
       theta = (theta * 2) % MXN;
    int i = 0;
for (int j = 1; j < n - 1; j++) {
       for (int k = n >> 1; k > (i ^= k); k >>= 1);
       if (j < i) swap(a[i], a[j]);
     if(inv) for (i = 0; i < n; i++) a[i] /= n;
  cplx arr[MXN+1];
  inline void mul(int _n,ll a[],int _m,ll b[],ll ans[])
     int n=1, sum=_n+_m-1;
     while(n<sum)</pre>
      n < < =1;
     for(int i=0;i<n;i++) {</pre>
       double x=(i<_n?a[i]:0),y=(i<_m?b[i]:0);</pre>
       arr[i]=complex<double>(x+y,x-y);
     fft(n,arr);
     for(int i=0;i<n;i++)</pre>
       arr[i]=arr[i]*arr[i];
     fft(n,arr,true);
     for(int i=0;i<sum;i++)</pre>
       ans[i]=(long long)(arr[i].real()/4+0.5);
}fft;
```

$3.2 \ O(1)$ mul

```
LL mul(LL x,LL y,LL mod){
  LL ret=x*y-(LL)((long double)x/mod*y)*mod;
  // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
  return ret<0?ret+mod:ret;
}</pre>
```

3.3 Faulhaber $(\sum_{i=1}^{n} i^p)$

```
/* faulhaber' s formula -
 * cal power sum formula of all p=1\sim k in O(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinactories
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
  while(b) {
    int q,t;
    q=a/b; t=b; b=a-b*q; a=t;
t=b0; b0=a0-b0*q; a0=t;
    t=b1; b1=a1-b1*q; a1=t;
  return a0<0?a0+mod:a0;</pre>
/* combinational
  for(int i=0;i<=MAXK;i++) {</pre>
    cm[i][0]=cm[i][i]=1;
    for(int j=1;j<i;j++)
  cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);</pre>
  /* inverse */
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
   * bernoulli */
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
for(int i=2;i<MAXK;i++) {</pre>
    if(i&1) { b[i]=0; continue; }
    b[i]=1;
    for(int j=0;j<i;j++)</pre>
      b[i]=sub(b[i],
                 mul(cm[i][j],mul(b[j], inv[i-j+1])));
  /* faulhaber */
  // sigma_x=1~n \{x^p\} = 
// 1/(p+1) * sigma_j=0~p \{C(p+1,j)*Bj*n^(p-j+1)\}
  for(int i=1;i<MAXK;i++) {</pre>
    co[i][0]=0;
    for(int j=0;j<=i;j++)</pre>
      co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]))
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
inline int solve(int n,int p) {
  int sol=0,m=n;
  for(int i=1;i<=p+1;i++) {
    sol=add(sol,mul(co[p][i],m));
    m = mul(m, n);
  return sol;
```

3.4 Chinese Remainder

3.5 Miller Rabin

|}

```
// n < 4,759,123,141
                              3 : 2, 7, 61
4 : 2, 13, 23, 1662803
// n < 1,122,004,669,633
// n < 3,474,749,660,383
                                    6 : pirmes <= 13
// n < 2^{64}
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL magic[]={}
bool witness(LL a,LL n,LL u,int t){
  if(!a) return 0;
  LL x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    LL nx=mul(x,x,n);
    if(nx==1&&x!=1&&x!=n-1) return 1;
    x=nx;
  }
  return x!=1;
bool miller_rabin(LL n) {
  int s=(magic number size)
  // iterate s times of witness on n
  if(n<2) return 0;</pre>
  if(!(n\&1)) return n == 2;
  ll u=n-1; int t=0;
  // n-1 = u*2^t
  while(!(u&1)) u>>=1, t++;
  while(s--){
    LL a=magic[s]%n;
    if(witness(a,n,u,t)) return 0;
  return 1;
}
```

3.6 Pollard Rho

```
// does not work when n is prime 0(n^(1/4))
LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
LL pollard_rho(LL n) {
   if(!(n&1)) return 2;
   while(true){
      LL y=2, x=rand()%(n-1)+1, res=1;
      for(int sz=2; res==1; sz*=2) {
        for(int i=0; i<sz && res<=1; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
        }
        y = x;
   }
   if (res!=0 && res!=n) return res;
} </pre>
```

3.7 Josephus Problem

```
int josephus(int n, int m){ //n人每m次
   int ans = 0;
   for (int i=1; i<=n; ++i)
        ans = (ans + m) % i;
   return ans;
}</pre>
```

3.8 Matrix

```
//矩陣乘法
for(int i = 0; i < n; i++){
   for(int j = 0; j < n; j++){
      for(int k = 0; k < n; k++){
        ret[i][j] += a[i][k] * b[k][j];
      }
}</pre>
```

```
//矩陣快速冪
int base[2][2] = {
                    int ans[2][2] = {
 {1, 1},
{1, 0}
                      {1, 0},
                       \{0, 1\}
int mypow(int y){
 while(v){
   if( y&1 ) { ans = mul(ans, base); } //實作矩陣乘法
   base = mul(base, base);//實作矩陣乘法
  return ans[0][0];
```

3.9 **Gaussian Elimination**

```
const int GAUSS_MOD = 100000007LL;
struct GAUSS{
  int n;
  vector<vector<int>> v;
  int ppow(int a , int k){
    if(k == 0) return 1;
    if(k % 2 == 0) return ppow(a * a % GAUSS_MOD , k >>
          1);
     if(k \% 2 == 1) return ppow(a * a % GAUSS_MOD , k >>
          1) * a % GAUSS_MOD;
  vector<int> solve(){
    vector<int> ans(n);
    REP(now , 0 , n){
       REP(i , now , n) if(v[now][now] == 0 \& v[i][now]
             != 0)
       swap(v[i] , v[now]); // det = -det;
if(v[now][now] == 0) return ans;
       int inv = ppow(v[now][now] , GAUSS_MOD - 2);
       REP(i , 0 , n) if(i != now){
  int tmp = v[i][now] * inv % GAUSS_MOD;
         REP(j, now, n + 1) (v[i][j] += GAUSS\_MOD -
              tmp * v[now][j] % GAUSS_MOD) %= GAUSS_MOD;
    ŘEP(i , 0 , n) ans[i] = v[i][n + 1] * ppow(v[i][i]
      , GAUSS_MOD - 2) % GAUSS_MOD;
    return ans;
  // gs.v.clear() , gs.v.resize(n , vector<int>(n + 1 ,
        0));
} gs;
```

3.10 Inverse Matrix

```
int GAUSS_MOD;
struct GAUSS{
 int n;
 vector<vector<int> > v;
  vector<vector<int> > rev;
  int mul(int x,int y,int mod){
    int ret=x*y-(int)((long double)x/mod*y)*mod;
    return ret<0?ret+mod:ret;</pre>
  int ppow(int a, int b){//res=(a^b)%m
    int res=1, k=a;
    while(b){
      if((b&1)) res=mul(res,k,GAUSS_MOD)%GAUSS_MOD;
      k=mul(k,k,GAUSS_MOD)%GAUSS_MOD;
      b>>=1;
   }
    return res%GAUSS_MOD;
  bool solve(){
    for(int now = 0; now < n; now++){
      int ch;
      for(ch = now; ch < n && !v[ch][now]; ch++);</pre>
      if(ch >= n) return 0;
      for(int i = now; i < n; i++) if(v[now][now] == 0
    && v[i][now] != 0){</pre>
```

```
swap(v[i] , v[now]); // det = -det;
swap(rev[i], rev[now]);
                      if(v[now][now] == 0) return 0;
                      int inv = ppow(v[now][now] , GAUSS_MOD - 2);
for(int i = 0; i < n; i++) if(i != now){
  int tmp = v[i][now] * inv % GAUSS_MOD;
  for(int i = 0; i)</pre>
                             for(int j = 0; j < n; j++) {
  (v[i][j] += GAUSS_MOD - tmp * v[now][j] %</pre>
                                    GAUSS_MOD) %= GAUSS_MOD;
(rev[i][j] += GAUSS_MOD - tmp * rev[now][j] %
                                                       GAUSS_MOD) %= GAUSS_MOD;
                     }
              }
              return 1;
}} gs;
signed main(){
       int n, p; //n*n matrix, MOD=p
       cin>>n>>p; //if(!n && !p) return 0;
       GAUSS\_MOD = p; gs.n = n;
       gs.v.clear() , gs.v.resize(n + 1, vector < int > (n + 2, vector 
                          0));
       gs.rev.clear() , gs.rev.resize(n + 1, vector<int>(n +
       2 , 0));
for(int i = 0; i < n; i++){
for(int j = 0; j < n; j++){
                      cin>>gs.v[i][j];
                      if(i == j) gs.rev[i][j] = 1;
              }
       if(!gs.solve()) cout << "singular\n";</pre>
       else{
               for(int i = 0; i < n; i++){
                      int inv = gs.ppow(gs.v[i][i] , p - 2);
                      for(int j = 0; j < n; j++)
                                   cout << (gs.rev[i][j] * inv % p) <<" ";
                      cout<<"\n";</pre>
       }
       cout << "\n";
                            模反元素
 3.11
 long long inv(long long a,long long m){
               long long x,y;
               long long d=exgcd(a,m,x,y);
               if(d==1) return (x+m)%m;
```

```
else return -1; //-1為無解
}
```

3.12 ax+by=gcd

```
PII gcd(int a, int b){
  if(b == \emptyset) return {1, \emptyset};
  PII q = gcd(b, a \% b);
  return {q.second, q.first - q.second * (a / b)};
}
int exgcd(int a,int b,long long &x,long long &y) {
    if(b == 0){x=1,y=0;return a;}
    int now=exgcd(b,a%b,y,x);
    y=a/b*x;
    return now;
}
```

3.13 Discrete sqrt

```
void calcH(LL &t, LL &h, const LL p) {
 LL tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
// solve equation x^2 \mod p = a
bool solve(LL a, LL p, LL &x, LL &y) {
 if(p == 2) { x = y = 1; return true; }
```

```
int p2 = p / 2, tmp = mypow(a, p2, p);
if (tmp == p - 1) return false;
if ((p + 1) % 4 == 0) {
    x=mypow(a,(p+1)/4,p); y=p-x; return true;
} else {
    LL t, h, b, pb; calcH(t, h, p);
    if (t >= 2) {
        do {b = rand() % (p - 2) + 2;
        } while (mypow(b, p / 2, p) != p - 1);
        pb = mypow(b, h, p);
    } int s = mypow(a, h / 2, p);
    for (int step = 2; step <= t; step++) {
        int ss = (((LL)(s * s) % p) * a) % p;
        for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);
        if (ss + 1 == p) s = (s * pb) % p;
        pb = ((LL)pb * pb) % p;
        y = p - x;
} return true;
}</pre>
```

3.14 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
  for( int i = 2 ; i < m ; i ++ )
    inv[ i ] = ((LL)(m - m / i) * inv[m % i]) % m;
}</pre>
```

3.15 Roots of Polynomial 找多項式的根

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ]; // a[0..n](coef) must be
    filled
int n; // degree of polynomial must be filled
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a[], int n, double x){
  double tmp=1,sum=0;
  for(int i=0;i<=n;i++)
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double l,double r,double a[],int n){
  int sl=sign(f(a,n,l)),sr=sign(f(a,n,r));
  if(sl==0) return l; if(sr==0) return r;
  if(sl*sr>0) return inf;
  while(r-l>eps){
    double mid=(l+r)/2;
    int ss=sign(f(a,n,mid));
    if(ss==0) return mid;
    if(ss*sl>0) l=mid; else r=mid;
  }
  return 1;
void solve(int n,double a[],double x[],int &nx){
  if(n==1){ x[1]=-a[0]/a[1]; nx=1; return; }
double da[10], dx[10]; int ndx;
  for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0;
  if(ndx==0){
    double tmp=binary(-inf,inf,a,n);
    if (tmp<inf) x[++nx]=tmp;</pre>
  double tmp;
  tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;
for(int i=1;i<=ndx-1;i++){</pre>
    tmp=binary(dx[i],dx[i+1],a,n);
    if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
} // roots are stored in x[1..nx]
```

3.16 Combination thearom

```
const ll mod = 1e9 + 7;
ll fac[(int)2e6 + 1], inv[(int)2e6 + 1];
ll getinv(ll a){ return qpow(a, mod-2); }
void init(int n){
  fac[0] = 1;
  for(int i = 1; i <= n; i++){
    fac[i] = fac[i-1] * i % mod;
  }
  inv[n] = getinv(fac[n]);
  for(int i = n - 1; i >= 0; i--){
    inv[i] = inv[i + 1] * (i + 1) % mod;
  }
}
ll C(int n, int m){
  if(m > n) return 0;
  return fac[n] * inv[m] % mod * inv[n-m] % mod;
}
```

3.17 Primes

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231
* 999888733, 98789101, 987777733, 999991921, 1010101333
 1010102101, 10000000000039, 1000000000000037
2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[ N ] , p_tbl[ N ];
vector<int> primes;
void sieve() {
  mu[1] = p_tbl[1] = 1;
  for( int i = 2 ; i < N ; i ++ ){
  if( !p_tbl[ i ] ){</pre>
       p_tbl[ i ] = i;
       primes.push_back( i );
mu[ i ] = -1;
     for( int p : primes ){
       int x = i * p;
       if( x >= M ) break;
       p_tbl[ x ] = p;
mu[ x ] = -mu[ i ];
       if(i \% p == 0){
          mu[x] = 0;
          break:
1 1 1 1
vector<int> factor( int x ){
  vector<int> fac{ 1 };
  while(x > 1){
     int fn = SZ(fac), p = p_tbl[x], pos = 0;
     while( x \% p == 0 ){
       x \neq p;
       for( int i = 0 ; i < fn ; i ++ )
fac.PB( fac[ pos ++ ] * p );</pre>
  } }
  return fac;
}
```

3.18 Phi

3.19 Result

```
• Lucas' Theorem : For n,m\in\mathbb{Z}^* and prime P, C(m,n) mod P=\Pi(C(m_i,n_i)) where
   m_i is the i\text{-th} digit of m in base P.
• Stirling approximation :
   n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\frac{1}{12n}}
• Stirling Numbers(permutation |P|=n with k cycles):
   S(n,k) = \text{coefficient of } x^k \text{ in } \Pi_{i=0}^{n-1}(x+i)
• Stirling Numbers(Partition n elements into k non-empty set):
   S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n
• Pick's Theorem : A=i+b/2-1 在二維座標平面中畫上網格·對於任何簡單多邊形
   A: 面積、i: 內部的格點數、b: 邊上的格點數
• Catalan number : C_n={2n\choose n}/(n+1) C_n^{n+m}-C_{n+1}^{n+m}=(m+n)!\frac{n-m+1}{n+1}\quad for\quad n\geq m
   C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{(n+1)!n!}
   C_0 = 1 and C_{n+1} = 2(\frac{2n+1}{n+2})C_n

C_0 = 1 and C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i} for n \ge 0
• Euler Characteristic:
   planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2
   V, E, F, C: number of vertices, edges, faces(regions), and compo-
• Kirchhoff's theorem : A_{ii}=deg(i), A_{ij}=(i,j)\in E ?-1:0, Deleting any one row, one column, and cal the det(A)
ullet Polya' theorem (c is number of color \cdot m is the number of cycle
   (\sum_{i=1}^{\stackrel{.}{m}}c^{\gcd(i,m)})/m
• Burnside lemma: |X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|
• 錯排公式: (n 個人中·每個人皆不再原來位置的組合數):
   dp[0] = 1; \dot{dp}[1] = 0;
   dp[i] = (i-1) * (dp[i-1] + dp[i-2]);
• Bell 數 (有 n 個人, 把他們拆組的方法總數):
   B_n = \sum_{k=0}^n s(n,k) (second – stirling)
   B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k
• Wilson's theorem :
   (p-1)! \equiv -1 \pmod{p}
• Fermat's little theorem :
   a^p \equiv a \pmod{p}
• Euler's totient function:
   A^{B^C} mod p = pow(A, pow(B, C, p - 1)) mod p
• 歐拉函數降冪公式: A^B \mod C = A^B \mod \phi(c) + \phi(c) \mod C
• 用歐拉函數求模反元素:
   如果 a 和 n 互質, 則 a 對 n 的模反元素 a^{-1} \equiv a^{\phi(n)-1} (mod\ n)
• 6 的倍數: (a-1)^3 + (a+1)^3 + (-a)^3 + (-a)^3 = 6a
• 上高斯 (向上取整):
   \lceil \frac{a}{b} \rceil = \frac{a+b-1}{b}
```

4 Geometry

4.1 definition

```
#define all(a) a.begin(),a.end()
ostream& operator<<(ostream& os, const Pt& pt) {
   return os << "(" << pt.x << ", " << pt.y << ")";}
typedef long double ld;
const ld eps = 1e-8;
const ld pi = acos(-1);
int dcmp(ld x) {
   if(abs(x) < eps) return 0;
   else return x < 0 ? -1 : 1;</pre>
```

```
struct Pt {
  ld x, y;
Pt(ld _x=0, ld _y=0):x(_x), y(_y) {}
   Pt operator+(const Pt &a) const {
  return Pt(x+a.x, y+a.y); }
Pt operator-(const Pt &a) const {
   return Pt(x-a.x, y-a.y); }
Pt operator*(const ld &a) const {
  return Pt(x*a, y*a); }
Pt operator/(const ld &a) const {
     return Pt(x/a, y/a); }
   ld operator*(const Pt &a) const {
     return x*a.x + y*a.y;
   ld operator^(const Pt &a) const {
     return x*a.y - y*a.x;
   bool operator<(const Pt &a) const {</pre>
     return x < a.x || (x == a.x && y < a.y); }
     //return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0 \&\&
          dcmp(y-a.y) < 0); }
   bool operator==(const Pt &a) const {
     return dcmp(x-a.x) == 0 &\& dcmp(y-a.y) == 0; }
};
ld norm2(const_Pt &a) {
   return a*a; }
ld norm(const Pt &a) {
   return sqrt(norm2(a)); }
Pt perp(const Pt &a) {
return Pt(-a.y, a.x); }
Pt rotate(const Pt &a, ld ang) {
   return Pt(a.x*cos(ang)-a.y*sin(ang), a.x*sin(ang)+a.y
       *cos(ang)); }
struct Circle {
  Pt o; ld r;
  Circle(Pt _{o}=Pt(0, 0), ld _{r}=0):o(_{o}), r(_{r}) {}
};
```

4.2 Intersection of 2 lines

```
Pt LLIntersect(Line a, Line b) {
   Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
   ld f1 = (p2-p1)^(q1-p1),f2 = (p2-p1)^(p1-q2),f;
   if(dcmp(f=f1+f2) == 0)
     return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
   return q1*(f2/f) + q2*(f1/f);
}
```

4.3 halfPlaneIntersection

```
|// 0(nlogn)
// 傳入 vector<Line>
// (半平面為點 st 往 ed 的逆時針方向)
// 回傳值為形成的凸多邊形的頂點 vector
// assume that Lines intersect
vector<Pt> HPI(vector<Line> P) {
     sort(P.begin(), P.end(), [&](Line l, Line m) {
          if (argcmp(l.v, m.v)) return true;
if (argcmp(m.v, l.v)) return false;
return PtSide(l.s, m) > 0;
     });
     int n = P.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {
    if (i and !argcmp(P[i - 1].v, P[i].v)) continue</pre>
          while (l < r and PtSide(LLIntersect(P[r-1], P[r</pre>
          ]), P[i]) <= 0) r--;
while (l < r and PtSide(LLIntersect(P[l], P[l
+1]), P[i]) <= 0) l++;
           P[++r] = P[i];
     while (l < r and PtSide(LLIntersect(P[r-1], P[r]),</pre>
           P[l]) <= 0) r--
     while (l < r and PtSide(LLIntersect(P[l], P[l+1]),</pre>
          P[r]) <= 0) l++;
     if (r - 1 <= 1 or !argcmp(P[l].v, P[r].v))</pre>
           return {}; // empty
      if (PtSide(LLIntersect(P[l], P[r]), P[l+1]) <= 0) {</pre>
          assert(0);
```

4.4 Convex Hull

```
double cross(Pt o, Pt a, Pt b){
 return (a-o) ^ (b-o);
vector<Pt> convex_hull(vector<Pt> pt){
 sort(pt.begin(),pt.end());
  int top=0;
 vector<Pt> stk(2*pt.size());
 for (int i=0; i<(int)pt.size(); i++){</pre>
   while (top >= 2 && cross(stk[top-2],stk[top-1],pt[i
       ]) <= 0) // 如果想要有點共線的點,把 <= 改成 <
     ton--
   stk[top++] = pt[i];
 for (int i=pt.size()-2, t=top+1; i>=0; i--){
   while (top >= t && cross(stk[top-2],stk[top-1],pt[i
       ]) <= 0)
     top--;
   stk[top++] = pt[i];
 stk.resize(top-1);
 return stk;
```

4.5 Convex Hull trick

```
struct Convex {
  vector<Pt> A, V, L, U;
 Convex(const vector<Pt> &_A) : A(_A), n(_A.size()) {
      // n >= 3
    auto it = max_element(all(A));
   L.assign(A.begin(), it + 1);
U.assign(it, A.end()), U.push_back(A[0]);
    for (int i = 0; i < n; i++) {
      V.push\_back(A[(i + 1) % n] - A[i]);
   }
  int PtSide(Pt p, Line L) {
   return dcmp((L.b - L.a)^{p - L.a);
  int inside(Pt p, const vector<Pt> &h, auto f) {
   auto it = lower_bound(all(h), p, f);
    if (it == h.end()) return 0;
    if (it == h.begin()) return p == *it;
    return 1 - dcmp((p - *prev(it))^(*it - *prev(it)))
 // 1. whether a given point is inside the CH
 // ret 0: out, 1: on, 2: in
  int inside(Pt p) {
    return min(inside(p, L, less{}), inside(p, U,
        greater{}));
 static bool cmp(Pt a, Pt b) { return dcmp(a \land b) > 0;
  // 2. Find tangent points of a given vector
  // ret the idx of far/closer tangent point
  int tangent(Pt v, bool close = true) {
    assert(v != Pt{});
    auto l = V.begin(), r = V.begin() + L.size() - 1;
    if (v < Pt{}) l = r, r = V.end();</pre>
    if (close) return (lower_bound(l, r, v, cmp) - V.
        begin()) % n;
    return (upper_bound(1, r, v, cmp) - V.begin()) % n;
```

```
// 3. Find 2 tang pts on CH of a given outside point
   // return index of tangent points
   // return {-1, -1} if inside CH
   array<int, 2> tangent2(Pt p) {
     array<int, 2> t{-1, -1};
if (inside(p) == 2) return t;
     if (auto it = lower_bound(all(L), p); it != L.end()
           and p == *it) {
       int s = it - L.begin();
       return {(s + 1) % n, (s - 1 + n) % n};
     if (auto it = lower_bound(all(U), p, greater{}); it
           != U.end() and p == *it) {
       int s = it - U.begin() + L.size() - 1;
       return \{(s + 1) \% n, (s - 1 + n) \% n\};
     for (int i = 0; i != t[0]; i = tangent((A[t[0] = i]
           - p), 0));
     for (int i = 0; i != t[1]; i = tangent((p - A[t[1]
          = i]), 1));
     return t;
   int find(int l, int r, Line L) {
  if (r < l) r += n;</pre>
     int s = PtSide(A[l % n], L);
     return *ranges::partition_point(views::iota(l, r),
       [&](int m) {
         return PtSide(A[m % n], L) == s;
       }) - 1;
   };
// 4. Find intersection point of a given line
   // intersection is on edge (i, next(i))
   vector<int> intersect(Line L) {
     int l = tangent(L.a - L.b), r = tangent(L.b - L.a);
     if(PtSide(A[1], L) == 0) return {1};
if(PtSide(A[r], L) == 0) return {r};
     if (PtSide(\bar{A}[\bar{1}], \bar{L}) * PtSide(\bar{A}[r], \bar{L}) > 0) return
          {};
     return {find(1, r, L) % n, find(r, l, L) % n};
};
```

4.6 Intersection of 2 segments

4.7 Point In Polygon

4.8 Circle cover

```
#define N 1021
#define D long double
struct CircleCover{
   int C; Circle c[ N ]; //填入C(圓數量),c(圓陣列)
  bool g[ N ][ N ], overlap[ N ][ N ];
// Area[i] : area covered by at least i circles
  D Area[ N ];
void init( int _C ){ C = _C; }
bool CCinter( Circle& a , Circle& b , Pt& p1 , Pt& p2
     Pt o1 = a.o , o2 = b.o;
D r1 = a.r , r2 = b.r;
     if( norm( o1 - o2 ) > r1 + r2 ) return {};
if( norm( o1 - o2 ) < max(r1, r2) - min(r1, r2) )</pre>
           return {};
     D d2 = (o1 - o2) * (o1 - o2);
     D d = sqrt(d2);
if( d > r1 + r2 ) return false;
     Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
     D A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
     Pt v=Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
p1 = u + v; p2 = u - v;
     return true;
   struct Teve {
     Pt p; D ang; int add;
     Teve() {}
     Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
     bool operator<(const Teve &a)const
  {return ang < a.ang;}
}eve[ N * 2 ];
   // strict: x = 0, otherwise x = -1
  bool disjuct( Circle& a, Circle &b, int x )
{return dcmp( norm( a.o - b.o ) - a.r - b.r ) > x;}
bool contain( Circle& a, Circle &b, int x )
  return (dcmp(c[i].r - c[j].r) > 0 \mid \mid
                (dcmp(c[i].r - c[j].r) == 0 \&\& i < j)) \&\&
                      contain(c[i], c[j], -1);
   void solve(){
     for( int i = 0 ; i <= C + 1 ; i ++ )
        Area[ i ] = 0;
     for( int i = 0 ; i < C ; i ++ )
for( int j = 0 ; j < C ; j ++
           overlap[i][j] = contain(i, j);
     for( int i = 0 ; i < C ; i ++ )
for( int j = 0 ; j < C ; j ++
           g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                           disjuct(c[i], c[j], -1));
     for( int i = 0 ; i < C ; i ++ ){
        int E = 0, cnt = 1;
        for( int j = 0 ; j < C ; j ++ )
  if( j != i && overlap[j][i] )</pre>
             cnt ++;
        for( int j = 0 ; j < C ; j ++ )
  if( i != j && g[i][j] ){
    Pt aa, bb;</pre>
             CCinter(c[i], c[j], aa, bb);
D A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
D B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);
             eve[E ++] = Teve(bb, B, 1);
             eve[E ++] = Teve(aa, A, -1);
             if(B > A) cnt ++;
        if( E == 0 ) Area[ cnt ] += pi * c[i].r * c[i].r;
        else{
           sort( eve , eve + E );
           eve[E] = eve[0];
for( int j = 0 ; j < E ; j ++ ){
  cnt += eve[j].add;</pre>
             Area[cnt] += (eve[j].p ^{\circ} eve[j + 1].p) * 0.5;
             D theta = eve[j + 1].ang - eve[j].ang;
if (theta < 0) theta += 2.0 * pi;
             Area[cnt] +=
                (theta - sin(theta)) * c[i].r*c[i].r * 0.5;
}}}};;
```

4.9 Tangent line of two circles

```
vector<Line> go( const Cir& c1 , const Cir& c2 , int
     sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = norm2(c1.0 - c2.0);
  if( d_sq < eps ) return ret;
double d = sqrt( d_sq );</pre>
  Pt v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if( c * c > 1 ) return ret;
  double h = sqrt( max( 0.0 , 1.0 - c * c ) );
  for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
                v.Y * c + sign2 * h * v.X };
     Pt p1 = c1.0 + n * c1.R;
     Pt p2 = c2.0 + n * (c2.R * sign1);
     if( fabs( p1.X - p2.X ) < eps and
    fabs( p1.Y - p2.Y ) < eps )
    p2 = p1 + perp( c2.0 - c1.0 );</pre>
     ret.push_back( { p1 , p2 } );
  return ret;
```

4.10 Minimum distance of two convex

```
double TwoConvexHullMinDis(Pt P[],Pt Q[],int n,int m){
  int mn=0,mx=0; double tmp,ans=1e9;
for(int i=0;i<n;++i) if(P[i].y<P[mn].y) mn=i;</pre>
  for(int i=0;i<m;++i) if(Q[i].y>Q[mx].y) mx=i;
  P[n]=P[0]; Q[m]=Q[0];
  for (int i=0;i<n;++i)</pre>
    while(tmp=((Q[mx+1]-P[mn+1])^(P[mn]-P[mn+1]))>((Q[
         mx]-P[mn+1])^(P[mn]-P[mn+1])) mx=(mx+1)%m;
    if(tmp<0) // pt to segment distance</pre>
      ans=min(ans,dis(Line(P[mn],P[mn+1]),Q[mx]));
    else // segment to segment distance
      ans=min(ans,dis(Line(P[mn],P[mn+1]),Line(Q[mx],Q[
           mx+17)));
    mn=(mn+1)%n;
  }
  return ans;
}
```

4.11 Poly Union

```
struct PY{
  int n; Pt pt[5]; double area;
  Pt& operator[](const int x){ return pt[x]; }
  void init(){ //n,pt[0~n-1] must be filled
    area=pt[n-1]^pt[0];
    for(int i=0;i<n-1;i++) area+=pt[i]^pt[i+1];</pre>
    if((area/=2)<0)reverse(pt,pt+n),area=-area;</pre>
PY py[500]; pair<double,int> c[5000];
inline double segP(Pt &p,Pt &p1,Pt &p2){
  if(dcmp(p1.x-p2.x)==0) return (p.y-p1.y)/(p2.y-p1.y);
  return (p.x-p1.x)/(p2.x-p1.x);
double polyUnion(int n){ //py[0~n-1] must be filled
  int i,j,ii,jj,ta,tb,r,d; double z,w,s,sum=0,tc,td;
  for(i=0;i<n;i++) py[i][py[i].n]=py[i][0];</pre>
  for(i=0;i<n;i++){</pre>
    for(ii=0;ii<py[i].n;ii++){</pre>
      r=0:
       c[r++]=make\_pair(0.0,0); c[r++]=make\_pair(1.0,0);
       for(j=0;j<n;j++){</pre>
         if(i==j) continue
         for(jj=0;jj<py[j].n;jj++){</pre>
           ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))
           tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj
                +17));
           if(ta==0 \&\& tb==0){
```

```
if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[
                 i][ii])>0&&j<i){
               c[r++]=make_pair(segP(py[j][jj],py[i][ii
                    ],py[i][ii+1]),1);
               c[r++]=make_pair(segP(py[j][jj+1],py[i][
                    ii],py[i][ii+1]),-1);
         }else if(ta>=0 && tb<0){</pre>
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
c[r++]=make_pair(tc/(tc-td),1);
         }else if(ta<0 && tb>=0){
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
            td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
c[r++]=make_pair(tc/(tc-td),-1);
    } } }
    sort(c,c+r);
    z=min(max(c[0].first,0.0),1.0); d=c[0].second; s
          =0;
     for(j=1;j<r;j++){</pre>
       w=min(max(c[j].first,0.0),1.0);
       if(!d) s+=w-z;
       d+=c[j].second; z=w;
    sum+=(py[i][ii]^py[i][ii+1])*s;
} }
return sum/2;
```

4.12 Minkowski sum

```
// P, Q, R(return) are counterclockwise order convex
    polygon
#define all(a) a.begin(),a.end()
vector<Pt> Minkowski(vector<Pt> P, vector<Pt> Q) {
    auto cmp = [\&](Pt a, Pt b) {
        return Pt{a.y, a.x} < Pt{b.y, b.x};
    auto reorder = [&](auto &R) {
        rotate(R.begin(), min_element(all(R), cmp), R.
             end());
        R.push\_back(R[0]), R.push\_back(R[1]);
    const int n = P.size(), m = Q.size();
    reorder(P), reorder(Q);
    vector<Pt> R;
    for (int i = 0, j = 0, s; i < n or j < m; ) {
    R.push_back(P[i] + Q[j]);
        s = dcmp((P[i + 1] - P[i]) \land (Q[j + 1] - Q[j]))
        if (s >= 0) i++;
        if (s <= 0) j++;
  rotate(R.begin(), min_element(all(R)), R.end());
    return R;
}
```

4.13 Area of Rectangles

```
struct AreaofRectangles{
#define cl(x) (x<<1)
#define cr(x) (x<<1/1)
    ll n, id, sid;
    pair<ll,ll> tree[MXN<<3];</pre>
                               // count, area
    vector<ll> ind;
    tuple<ll, ll, ll, ll> scan[MXN<<1];</pre>
    void pull(int i, int l, int_r){
        if(tree[i].first) tree[i].second = ind[r+1] -
            ind[l];
        else if(l != r){
            int mid = (l+r)>>1;
            tree[i].second = tree[cl(i)].second + tree[
                 cr(i)].second;
        else
                tree[i].second = 0;
    void upd(int i, int l, int r, int ql, int qr, int v
```

```
if(ql \ll l \& r \ll qr)
         tree[i].first += v;
         pull(i, l, r); return;
    int mid = (l+r) >> 1;
    if(ql <= mid) upd(cl(i), l, mid, ql, qr, v);</pre>
    if(qr > mid) upd(cr(i), mid+1, r, ql, qr, v);
    pull(i, l, r);
void init(int _n){
    n = _n; id = sid = 0;
    ind.clear(); ind.resize(n<<1);</pre>
    fill(tree, tree+(n<<2), make_pair(0, 0));</pre>
void addRectangle(int lx, int ly, int rx, int ry){
    ind[id++] = lx; ind[id++] = rx;
    scan[sid++] = make\_tuple(ly, 1, lx, rx);
    scan[sid++] = make_tuple(ry, -1, lx, rx);
ll solve(){
    sort(ind.begin(), ind.end());
    ind.resize(unique(ind.begin(), ind.end()) - ind
         .begin());
    sort(scan, scan + sid);
    11 area = 0, pre = get<0>(scan[0]);
    for(int i = 0; i < sid; i++){
        auto [x, v, l, r] = scan[i];
         area += tree[1].second * (x-pre);
        upd(1, 0, ind.size()-1, lower_bound(ind.
begin(), ind.end(), l)-ind.begin(),
             lower_bound(ind.begin(),ind.end(),r)-
             ind.begin()-1, v);
        pre = x;
    return area;
}rect;
```

4.14 Min dist on Cuboid

```
typedef LL T;
Tr;
x0+L, y0, H, W, L);
if(j>=0 && j< 2) turn(i, j+1, x, y0+W+z, y0+W-y,
x0, y0+W, L, H, W);
  x0, y0-H, L, H, W);
T solve(T L, T W, T H,
       T \times 1, T \times 1, T \times 2, T \times 2, T \times 2){
  if( z1!=0 && z1!=H ){
    if( y1==0 || y1==W )
     swap(y1,z1), swap(y2,z2), swap(W,H);
  else swap(x1,z1), swap(x2,z2), swap(L,H);
  if (z1==H) z1=0, z2=H-z2;
  r=INF; turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H);
  return r;
}
```

4.15 Heart of Triangle

```
Pt inCenter( Pt &A, Pt &B, Pt &C) { // 內心 double a = norm(B-C), b = norm(C-A), c = norm(A-B); return (A * a + B * b + C * c) / (a + b + c); }
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心 Pt bb = b - a, cc = c - a; double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc); return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d; }
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心 Pt ba = b - a, ca = c - a, bc = b - c;
```

```
double Y = ba.Y * ca.Y * bc.Y,
   A = ca.X * ba.Y - ba.X * ca.Y,
   x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A,
   y0= -ba.X * (x0 - c.X) / ba.Y + ca.Y;
   return Pt(x0, y0);
}
```

5 Graph

5.1 DSU 並查集 & MST

```
struct DSU {// 並查集
     vector<int> fa, sz;
DSU(int n = 0) : fa(n), sz(n, 1) {
         iota(fa.begin(), fa.end(), 0);
     int Find(int x) { // 路徑壓縮
while (x != fa[x])
            x = fa[x] = fa[fa[x]];
         return x:
     }
    bool Merge(int x, int y) { //合併
         x = Find(x), y = Find(y);
         if (x == y) return false; // 是否為連通
         if (sz[x] > sz[y]) swap(x, y);
         fa[x] = y;
         sz[y] += sz[x];
         return true;
    }
|};
```

5.2 Lowest Common Ancestor O(lqn)

```
struct LCA {
  int n, ti,
                lgN;
  int anc[MXN + 5][__lg(MXN) + 1] = {0};
int MaxLength[MXN][__lg(MXN) + 1] = {0};
  int time_in[MXN] = {0};
int time_out[MXN] = {0};
  LCA(int _n, int f):n(_n), ti(0), lgN(__lg(n)) {
  dfs(f, f, 0);
    build();
  void dfs(int now, int f, int len_to_father) { // dfs
         for anc, time, Lenth
     ti++;
     anc[now][0] = f;
     time_in[now] = ti;
    MaxLength[now][0] = len_to_father;
     for (auto i : graph[now]) {
          if (i.first == f) continue;
          dfs(i.first, now, i.second);
     time_out[now] = ti;
  void build() {  // build anc[][], MaxLength[][]
for (int i = 1; i <= lgN; ++i) {</pre>
        for (int u = 1; u <= n; ++u) {
          anc[u][i] = anc[anc[u][i - 1]][i - 1];
MaxLength[u][i] = max(MaxLength[u][i - 1],
                       MaxLength[anc[u][i - 1]\bar{[}i - 1]);
          // dis[u][i] += dis[anc[u][i - 1]][i - 1]
// + dis[u][i - 1];
    }
  bool isAncestor(int x, int y) {
     return time_in[x] <= time_in[y] && time_out[x] >=
          time_out[y];
  int getLCA(int u, int v) {
  if (isAncestor(u, v)) return u;
  if (isAncestor(v, u)) return v;
     for (int i = lgN; i >= 0; --i) {
       if (!isAncestor(anc[u][i], v)) {
          u = anc[u][i];
```

```
}
    }
    return anc[u][0];
  }
  int getMAX(int u, int v) { //獲得路徑上最大邊權
    int lca = getLCA(u, v);
    int maxx = -1;
    for (int i = lgN; i >= 0; --i) {
// u to lca
      if (!isAncestor(anc[u][i], lca)) {
        maxx = max(maxx, MaxLength[u][i]);
        u = anc[u][i];
      // v to lca
      if (!isAncestor(anc[v][i], lca)) {
        maxx = max(maxx, MaxLength[v][i]);
        v = anc[v][i];
      }
    if (u != lca) maxx = max(maxx, MaxLength[u][0]);
    if (v != lca) maxx = max(maxx, MaxLength[v][0]);
    return maxx;
  }
};
```

5.3 Hamiltonian path $O(n^22^n)$

5.4 MaximumClique 最大團

```
#define N 111
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int linkto[N] , v[N];
  void init(int _n){
    n = _n;
    for(int i = 0; i < n; i ++){
      linkto[i].reset(); v[i].reset();
  void addEdge(int a , int b)
  \{ v[a][b] = v[b][a] = 1; \}
  int popcount(const Int& val)
  { return val.count(); } int lowbit(const Int& val)
  { return val._Find_first(); }
  int ans , stk[N];
int id[N] , di[N] , deg[N];
  Int cans;
  void maxclique(int elem_num, Int candi){
    if(elem_num > ans){
      ans = elem_num; cans.reset();
for(int i = 0; i < elem_num; i ++)</pre>
         cans[id[stk[i]]] = 1;
    int potential = elem_num + popcount(candi);
    if(potential <= ans) return;</pre>
    int pivot = lowbit(candi);
    Int smaller_candi = candi & (~linkto[pivot]);
    while(smaller_candi.count() && potential > ans){
      int next = lowbit(smaller_candi);
      candi[next] = !candi[next];
      smaller_candi[next] = !smaller_candi[next];
      potential --;
```

```
if(next == pivot || (smaller_candi & linkto[next
             ]).count()){
          stk[elem_num] = next;
          maxclique(elem_num + 1, candi & linkto[next]);
  } } }
   int solve(){
     for(int i = 0; i < n; i + +){
       id[i] = i; deg[i] = v[i].count();
     sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
for(int i = 0 ; i < n ; i ++) di[id[i]] = i;</pre>
     for(int i = 0 ; i < n ; i ++)</pre>
       for(int j = 0; j < n; j ++)
  if(v[i][j]) linkto[di[i]][di[j]] = 1;</pre>
     Int cand; cand.reset();
     for(int i = 0; i < n; i ++) cand[i] = 1;
     ans = 1:
     cans.reset(); cans[0] = 1;
     maxclique(0, cand);
     return ans;
} }solver;
```

5.5 MaximalClique 極大團

```
#define N 80
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int lnk[N] , v[N];
  void init(int _n){
    n = _n;
    for(int i = 0; i < n; i ++){
      lnk[i].reset(); v[i].reset();
  void addEdge(int a , int b)
  \{ v[a][b] = v[b][a] = 1; \}
  int ans , stk[N], id[N] , di[N] , deg[N];
  Int cans;
  void dfs(int elem_num, Int candi, Int ex){
    if(candi.none()&ex.none()){
      cans.reset();
       for(int i = 0 ; i < elem_num ; i ++)</pre>
        cans[id[stk[i]]] = 1;
      ans = elem_num; // cans is a maximal clique
    int pivot = (candilex)._Find_first();
    Int smaller_candi = candi & (~lnk[pivot]);
    while(smaller_candi.count()){
      int nxt = smaller_candi._Find_first();
      candi[nxt] = smaller_candi[nxt] = 0;
      ex[nxt] = 1;
      stk[elem_num] = nxt;
      dfs(elem_num+1,candi&lnk[nxt],ex&lnk[nxt]);
  } }
  int solve(){
    for(int i = 0; i < n; i ++){
      id[i] = i; deg[i] = v[i].count();
    sort(id , id + n , [&](int id1, int id2){
           return deg[id1] > deg[id2]; });
    for(int i = 0 ; i < n ; i ++) di[id[i]] = i;
for(int i = 0 ; i < n ; i ++)</pre>
      for(int j = 0; j < n; j ++)
        if(v[i][j]) lnk[di[i]][di[j]] = 1;
    ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n,'1')), 0);
    return ans;
} }solver;
```

5.6 BCC based on vertex 點雙聯通分量

```
#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
  int n,nScc,step,dfn[MXN],low[MXN];
  vector<int> E[MXN],sccv[MXN];
```

```
int top,stk[MXN];
  void init(int _n) { // 初始化n點
    n = _n; nScc = step = 0;
    for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v) // 無向邊
{ E[u].PB(v); E[v].PB(u); }
  void DFS(int u, int f) {
    dfn[u] = low[u] = step++;
    stk[top++] = u
     for (auto v:E[u]) {
       if (v == f) continue;
       if (dfn[v] == -1) {
        DFS(v,u);
         low[u] = min(low[u], low[v]);
         if (low[v] >= dfn[u]) {
           int z
           sccv[nScc].clear();
           do {
             z = stk[--top]
             sccv[nScc].PB(z);
           } while (z != v);
           sccv[nScc++].PB(u);
      }else
        low[u] = min(low[u],dfn[v]);
  } }
  vector<vector<int>>> solve() { // 回傳(size=2 橋, size
      >2 點雙連通分量)
     vector<vector<int>> res;
     for (int i=0; i<n; i++)</pre>
      dfn[i] = low[i] = -1;
     for (int i=0; i<n; i++)
      if (dfn[i] == -1) {
        top = 0;
        DFS(i,i);
    REP(i,nScc) res.PB(sccv[i]);
    return res;
}graph;
```

5.7 Strongly Connected Component 強連通分 量

```
#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x)) //fill zero
struct Scc{
  int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
  void init(int _n){
    n = _n;
for (int i=0; i<MXN; i++)</pre>
       E[i].clear(), rE[i].clear();
  void addEdge(int u, int v){
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u){
    vst[u]=1;
    for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u){
    vst[u] = 1; bln[u] = nScc;
     for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve(){
    nScc = 0;
    vec.clear();
    FZ(vst);
    for (int i=0; i<n; i++)
  if (!vst[i]) DFS(i);</pre>
    reverse(vec.begin(),vec.end());
    FZ(vst);
    for (auto v : vec)
       if (!vst[v]){
         rDFS(v); nScc++;
```

```
|<sub>};</sub>
```

5.8 Maximum General graph Matching

```
// should shuffle vertices and edges
const int N=100005, E=(2e5)*2+40;
struct Graph{ // 1-based; match:
  int to[E],bro[E],head[N],e,lnk[N],vis[N],stp,n;
  void init(int _n){
    stp=0; e=1; n=_n;
    for(int i=1;i<=n;i++) head[i]=lnk[i]=vis[i]=0;</pre>
  void add_edge(int u,int v){
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i]
      if(!lnk[v]){ lnk[x]=v,lnk[v]=x; return true; }
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(vis[lnk[v]]<stp){</pre>
        int w=lnk[v]; lnk[x]=v,lnk[v]=x,lnk[w]=0;
        if(dfs(w)) return true;
        lnk[w]=v, lnk[v]=w, lnk[x]=0;
    return false;
  int solve(){
    int ans=0;
    for(int i=1;i<=n;i++) if(!lnk[i]) stp++,ans+=dfs(i)</pre>
    return ans;
 }
}graph;
```

5.9 ManhattanMST

```
//return {{u,v},w}: u <-> v (w), 需要再手動去重
//need Point definition
vector<pair<pair<int,int>, int>> ManhattanMST(vector<Pt</pre>
  vector<int> id(P.size());
 iota(id.begin(),id.end(), 0);
  vector<pair<int,int>, int>> edg;
  for (int k = 0; k < 4; k++) {
    sort(id.begin(),id.end(), [&](int i, int j) {
      return (P[i] - P[j]).x < (P[j] - P[i]).y;</pre>
   });
    map<int, int> sweep;
    for (int i : id) {
      auto it = sweep.lower_bound(-P[i].y);
      while (it != sweep.end()) {
        int j = it->second;
        Pt d = P[i] - P[j];
        if (d.y > d.x) break;
        edg.push_back(\{\{i, j\}, d.x + d.y\});
        it = sweep.erase(it);
      }
      sweep[-P[i].y] = i;
    for (Pt &p : P) {
      if (k \% 2) p.x = -p.x;
      else swap(p.x, p.y);
  return edg;
```

5.10 Min Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
  struct Edge { int v,u; double c; };
  int n, m, prv[V][V], prve[V][V], vst[V];
  Edge e[E];
  vector<int> edgeID, cycle, rho;
  double d[V][V];
  void init( int _n )
  \{ n = _n; m = 0; \}
  // WARNING: TYPE matters
  void addEdge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
  void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;</pre>
     for(int i=0; i<n; i++) {</pre>
       fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;</pre>
         if(d[i][v]<inf<sup>'</sup>&& d[i+1][u]>d[i][v]+e[j].c) {
           d[i+1][u] = d[i][v]+e[j].c;
           prv[i+1][u] = v;
           prve[i+1][u] = j;
  } } } }
  double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1;
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
       for(int k=0; k<n; k++) {</pre>
         if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i</pre>
              ])/(n-k));
         else avg=max(avg,inf);
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
     fill(vst,0); edgeID.clear(); cycle.clear(); rho.
          clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
       vst[st]++
       edgeID.PB(prve[i][st]);
       rho.PB(st);
     while (vst[st] != 2) {
       if(rho.empty()) return inf;
       int v = rho.back(); rho.pop_back();
       cycle.PB(v);
       vst[v]++;
    reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
    return mmc;
} }mmc;
```

5.11 Directed Graph Min Cost Cycle

```
// works in O(N M)
#define INF 1000000000000000LL
#define N 5010
#define M 200010
struct edge{
  int to; LL w;
  edge(int a=0, LL b=0): to(a), w(b){}
};
struct node{
  LL d; int u, next;
  node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
struct DirectedGraphMinCycle{
  vector<edge> g[N], grev[N];
LL dp[N][N], p[N], d[N], mu;
  bool inq[N];
  int n, bn, bsz, hd[N];
  void b_insert(LL d, int u){
    int i = d/mu;
```

if(i >= bn) return; b[++bsz] = node(d, u, hd[i]);hd[i] = bsz;void init(int _n){ n = _n; for(int i = 1 ; i <= n ; i ++) g[i].clear(); void addEdge(int ai , int bi , LL ci) { g[ai].push_back(edge(bi,ci)); } LL solve(){ fill(dp[0], dp[0]+n+1, 0);for(int i=1; i<=n; i++){ fill(dp[i]+1, dp[i]+n+1, INF); for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){ for(int k=0; k<(int)g[j].size(); k++) dp[i][g[j][k].to] =min(dp[i][g][k].to],</pre> dp[i-1][j]+g[j][k].w);} } mu=INF; LL bunbo=1; for(int i=1; i<=n; i++) if(dp[n][i] < INF){</pre> LL a=-INF, b=1; for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){ if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre> a = dp[n][i]-dp[j][i];b = n-j;if(mu*b > bunbo*a) mu = a, bunbo = b; if(mu < 0) return -1; // negative cycle</pre> if(mu == INF) return INF; // no cycle if(mu == 0) return 0; for(int i=1; i<=n; i++)</pre> for(int j=0; j<(int)g[i].size(); j++) g[i][j].w *= bunbo;</pre> memset(p, 0, sizeof(p)); queue<int> q; for(int i=1; i<=n; i++){</pre> q.push(i); inq[i] = true; while(!q.empty()){ int i=q.front(); q.pop(); inq[i]=false; for(int j=0; j<(int)g[i].size(); j++){ if(p[g[i][j].to] > p[i]+g[i][j].w-mu){ p[g[i][j].to] = p[i]+g[i][j].w-mu;if(!inq[g[i][j].to]){ q.push(g[i][j].to); inq[g[i][j].to] = true;} } } for(int i=1; i<=n; i++) grev[i].clear();</pre> for(int i=1; i<=n; i++)</pre> for(int j=0; j<(int)g[i].size(); j++){ g[i][j].w += p[i]-p[g[i][j].to];</pre> grev[g[i][j].to].push_back(edge(i, g[i][j].w)); LL mldc = n*mu; for(int i=1; i<=n; i++){</pre> bn=mldc/mu, bsz=0; memset(hd, 0, sizeof(hd)); fill(d+i+1, d+n+1, INF); b_insert(d[i]=0, i); for(int j=0; j<=bn-1; j++) for(int k=hd[j]; k; k=</pre> b[k].next){ int u = b[k].u; LL du = b[k].d;if(du > d[u]) continue; for(int l=0; l<(int)g[u].size(); l++) if(g[u][l</pre>].to > i){ if(d[g[u][ĺ].to] > du + g[u][l].w){ d[g[u][l].to] = du + g[u][l].w;b_insert(d[g[u][l].to], g[u][l].to); for(int j=0; j<(int)grev[i].size(); j++) if(grev[</pre> i][j].to > i) mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w); return mldc / bunbo; } }graph;

5.12 K-th Shortest Path

```
// time: O(|E| \setminus |g| |E| + |V| \setminus |g| |V| + K)
// memory: 0(|E| \lg |E| + |V|)
struct KSP{ // 1-base
  struct nd{
     int u, v; ll d;
     nd(int ui = 0, int vi = 0, ll di = INF)
     \{ u = ui; v = vi; d = di; \}
  struct heap{
     nd* edge; int dep; heap* chd[4];
  static int cmp(heap* a,heap* b)
  { return a->edge->d > b->edge->d; }
  struct node{
     int v; ll d; heap* H; nd* E;
     node(){}
    node(ll _d, int _v, nd* _E)
{ d =_d; v = _v; E = _E; }
node(heap* _H, ll _d)
     \{ H = _H; d = _d; \}
     friend bool operator<(node a, node b)
     { return a.d > b.d; }
  int n, k, s, t;
ll dst[N];
  nd *nxt[ N ];
  vector<nd*> g[ N ], rg[ N ];
heap *nullNd, *head[ N ];
  void init( int _n , int _k , int _s , int _t ){
    n = _n;    k = _k;    s = _s;    t = _t;
    for( int i = 1 ; i <= n ; i ++ ){
        g[ i ].clear();
        rg[ i ].clear();
    }
}</pre>
       nxt[ i ] = NULL; head[ i ] = NULL;
dst[ i ] = -1;
  } }
  void addEdge( int ui , int vi , ll di ){
     nd* e = new nd(ui, vi, di);
g[_ui ].push_back( e );
     rg[ vi ].push_back( e );
  queue<int> dfsQ;
  void dijkstra(){
     while(dfsQ.size()) dfsQ.pop();
     priority_queue<node> Q;
     Q.push(node(0, t, NULL));
     while (!Q.empty()){
       node p = Q.top(); Q.pop();
        if(dst[p.v] != -1) continue;
       dst[ p.v ] = p.d;
nxt[ p.v ] = p.E;
       dfsQ.push( p.v_);
       for(auto e: rg[ p.v ])
          Q.push(node(p.d + e->d, e->u, e));
  } }
  heap* merge(heap* curNd, heap* newNd){
     if(curNd == nullNd) return newNd;
     heap* root = new heap;
     memcpy(root, curNd, sizeof(heap));
if(newNd->edge->d < curNd->edge->d){
       root->edge = newNd->edge
       root->chd[2] = newNd->chd[2];
       root->chd[3] = newNd->chd[3];
       newNd->edge = curNd->edge;
       newNd - > chd[2] = curNd - > chd[2];
       newNd - chd[3] = curNd - chd[3];
     if(root->chd[0]->dep < root->chd[1]->dep)
       root->chd[0] = merge(root->chd[0],newNd);
       root->chd[1] = merge(root->chd[1],newNd);
     root->dep = max(root->chd[0]->dep, root->chd[1]->
          dep) + 1;
     return root;
  vector<heap*> V;
  void build(){
     nullNd = new heap;
     nullNd->dep = 0;
     nullNd->edge = new nd;
```

```
fill(nullNd->chd, nullNd->chd+4, nullNd);
    while(not dfsQ.empty()){
       int u = dfsQ.front(); dfsQ.pop();
       if(!nxt[ u ]) head[ u ] = nullNd;
else head[ u ] = head[nxt[ u ]->v];
       V.clear();
       for( auto\&\& e : g[u]){
         int v = e->v;
         if( dst[ v ] == -1 ) continue;
         e->d += dst[ v ] - dst[ u ];
if( nxt[ u ] != e ){
            heap* p = new heap;
            fill(p->chd, p->chd+4, nullNd);
            p->dep = 1;
            p->edge = \acute{e};
            V.push_back(p);
       if(V.empty()) continue;
       make_heap(V.begin(), V.end(), cmp);
#define L(X) ((X<<1)+1)
#define R(X) ((X<<1)+2)
       for( size_t i = 0 ; i < V.size() ; i ++ ){</pre>
         if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
         else V[i]->chd[2]=nullNd;
         if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
         else V[i]->chd[3]=nullNd;
       head[u] = merge(head[u], V.front());
  } }
  vector<ll> ans;
  void first_K(){
    ans.clear();
    priority_queue<node> Q;
if( dst[ s ] == -1 ) return;
    ans.push_back( dst[ s ] );
    if( head[s] != nullNd )
    Q.push(node(head[s], dst[s]+head[s]->edge->d));
for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();</pre>
       ans.push_back( p.d );
       if(head[ p.H->edge->v ] != nullNd){
         q.H = head[p.H->edge->v];
         q.d = p.d + q.H->edge->d;
         Q.push(q);
       for( int i = 0 ; i < 4 ; i ++ )
  if( p.H->chd[ i ] != nullNd ){
    q.H = p.H->chd[ i ];
            q.d = p.d - p.H->edge->d + p.H->chd[i]->
                edge->d;
            Q.push( q );
  } }
         }
  void solve(){ // ans[i] stores the i-th shortest path
    dijkstra();
    build();
    first_K(); // ans.size() might less than k
} }solver;
```

5.13 Floryd Warshall

5.14 SPFA

```
#define MXN 200005
struct SPFA{
```

```
int n;
  LL inq[MXN], len[MXN];
  vector<LL> dis;
  vector<pair<int, LL>> edge[MXN];
  void init(int _n){
    n = _n;
    dis.clear(); dis.resize(n, 1e18);
for(int i = 0; i < n; i++){</pre>
      edge[i].clear();
       inq[i] = len[i] = 0;
  void addEdge(int u, int v, LL w){
    edge[u].push_back({v, w});
  vector<LL> solve(int st = 0){
    deque<int> dq; //return {-1} if has negative cycle
    dq.push_back(st); //otherwise return dis from st
inq[st] = 1; dis[st] = 0;
    while(!dq.empty()){
       int u = dq.front(); dq.pop_front();
       inq[u] = 0;
       for(auto [to, d] : edge[u]){
         if(dis[to] > d+dis[u]){
           dis[to] = d+dis[u];
           len[to] = len[u]+1;
           if(len[to] > n) return {-1};
           if(inq[to]) continue;
           (!dq.empty()&&dis[dq.front()] > dis[to]?
               dq.push_front(to) : dq.push_back(to));
           inq[to] = 1;
    } } }
    return dis;
} }spfa;
```

5.15 Tree Hash

```
//限定root = 1
//從 dfs(1,1) 開始
int subtree_sz[MXN];
vector<int> edge[MXN];
int dfs(int u, int f) {
  vector<pair<int, int>> h;
  subtree\_sz[u] = 1;
  for (int child : edge[u]) {
    if (child == f) continue;
    int tmp = dfs(child, u);
    h.push_back(make_pair(tmp, subtree_sz[child]));
    subtree_sz[u] += subtree_sz[child];
  sort(h.begin(), h.end());
  int ret = subtree_sz[u];
  for (auto v : h) {
    ret = ((ret * p) % MOD + v.first) % MOD;
    ret = ret * v.second % MOD;
  return ret;
}
```

5.16 HeavyLightDecomposition

```
}
// 第二次 dfs
int top[MXN]; // 每個節點所在的鏈的頂端節點
int dfn[MXN]; // 節點編號,編號為在線段樹上的位置
int rnk[MXN]; // 編號為哪個節點
int bottom[MXN]; // 維護每個節點的子樹中最大 dfn 編號
int cnt = 0:
int dfs_hld(int x, int f){
    top[x] = (son[fa[x]] == x ? top[fa[x]] : x);
    rnk[cnt] = x;
    bottom[x] = dfn[x] = cnt++;
                 bottom[x] = max(bottom[x], dfs_hld(
    if(son[x])
        son[x], x)); // 更新子樹最大編號
    for(int i : edge[x]){
       if(i == f || i == son[x])
                                    continue;
       bottom[x] = max(bottom[x], dfs_hld(i, x)); //
            更新子樹最大編號
    return bottom[x];
}
// 求出 lca
// 不斷跳鏈·直到 u,v 跳到同一條鏈上為止
// 每次跳鏈選所在的鏈頂端深度較深的一端往上跳
int getLca(int u, int v) {
   while(top[u] != top[v]){
     if(dep[top[u]] > dep[top[v]])
         u = fa[top[u]];
      else
         v = fa[top[v]];
    return dep[u] > dep[v] ? v : u;
}
// 路徑權重總和
int query(int u, int v) {
    int ret = 0
    while(top[u] != top[v]){
        if (dep[top[u]] > dep[top[v]]){
           ret += segtree.query(dfn[top[u]], dfn[u]);
           u = fa[top[u]];
       }
       else{
           ret += segtree.query(dfn[top[v]], dfn[v]);
           v = fa[top[v]];
   }
    // 最後到同一條鏈上
    ret += segtree.query(min(dfn[u], dfn[v]), max(dfn[u
        ], dfn[v]));
    return ret;
}
```

5.17 差分約束

約束條件 $V_j - V_i \leq W$ addEdge(V_i, V_j, W) and run bellman-ford or spfa

6 String

6.1 PalTree O(n)

```
diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++:
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0;
    newNode(0,1), newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

6.2 Longest Increasing Subsequence

```
vector<int> getLIS(vector<int> a){
  vector<int> lis;
  for(int i : a){
    if(lis.empty() || lis.back() < i) | lis.push_Back(
        i);
    else      *lower_bound(lis.begin(), lis.end(), i) =
        i;
  }
  return lis;
}</pre>
```

6.3 Longest Common Subsequence O(nlgn)

6.4 KMP

```
| /* len-failure[k]:
在k結尾的情況下‧這個子字串可以由開頭長度為(len-failure[k])的部分重複出現來表達
| failure[k]為次長相同前綴後綴| 如果我們不只想求最多‧而且以0-base做為考量
```

```
那可能的長度由大到小會是
failuer[k] · failure[failuer[k]-1]
 failure[failure[failuer[k]-1]-1]..
直到有值為0為止 */
int failure[MXN];
vector<int> KMP(string& t, string& p) {
    vector<int> ret;
    if(p.size() > t.size()) return ret;
    for(int i = 1, j = failure[0] = -1; i < p.size(); i</pre>
        while(j \ge 0 \&\& p[j + 1] != p[i]) j = failure[j]
        if(p[j + 1] == p[i]) j++;
        failure[i] = j;
    for(int i = 0, j = -1; i < t.size(); i++) {
        while (j >= 0 && p[j + 1] != t[i]) j = failure[
        j];
if(p[j + 1] == t[i]) j++;
        if(j == p.size() - 1) {
            ret.push_back(i - p.size() + 1);
            j = failure[j];
        }
    return ret;
}
```

6.5 SAIS O(n)

```
/*** SA· 將字串的所有後綴排序後的數組 ***/
/* SA[i]儲存排序後第i小的後綴從哪裡開始 */
/**** H[i] 為第i小的字串跟第i-1小的LCP ***/
/**** 註:LCP(Longest Common Prefix) ****/
/**** ex:S = "babd", SA[0] = 1("abd") ****/
/** SA[1] = 0("babd"), SA[2] = 2("bd") **/
/*** H[0] = 0, H[1] = 0, H[2] = 1("b") ***/
/* 傳入參數:ip 陣列放字串,len為字串長度 */
/* 需保證ip[len]為0, 且字串裡的元素不為0 */
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
  hei[N], r[N];
int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
    memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
    mkhei(n);
  void mkhei(int n){
    REP(i,n) r[\_sa[i]] = i;
    hei[0] = 0;
    REP(i,n) if(r[i]) {
       int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
      hei[r[i]] = ans;
  void sais(int *s, int *sa, int *p, int *q, bool *t,
      int *c, int_n, int z){
    bool uniq = t[n-1] = true, neq;
    int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
         lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
    memcpy(x, c, sizeof(int) * z); \
    memcpy(x + 1, c, sizeof(int) * (z - 1)); \
    REP(i,n) if(sa[i] \& !t[sa[i]-1]) sa[x[s[sa[i]-1]] 
    ]-1]]++] = sa[i]-1; \
memcpy(x, c, sizeof(int) * z); \
for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i
         ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
    MSO(c, z);
    REP(i,n) uniq \&= ++c[s[i]] < 2;
    REP(i,z-1) c[i+1] += c[i];
    if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
```

```
MAGIC(REP1(i,1,n-1) if(t[i] \&\& !t[i-1]) sa[--x[s[i] k])
        ]]]=p[q[i]=nn++]=i);
    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
      neq=lst<0||memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa|)
           [i])*sizeof(int));
      ns[q[lst=sa[i]]]=nmxz+=neq;
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
    MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[s[p[
        nsa[i]]]] = p[nsa[i]]);
}sa;
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
  // should padding a zero in the back
  // ip is int array, len is array length
// ip[0..n-1] != 0, and ip[len] = 0
  ip[len++] = 0;
  sa.build(ip, len, 128);
  for (int i=0; i<len; i++) {
    H[i] = sa.hei[i + 1];
    SA[i] = sa.\_sa[i + 1];
  // resulting height, sa array \in [0,len)
```

6.6 Z Value O(n)

```
//z[i] = lcp(s[1...n-1],s[i...n-1])
int z[MAXN];
void Z_value(const string& s) {
   int i, j, left, right, len = s.size();
   left=right=0; z[0]=len;
   for(i=1;i<len;i++) {
      j=max(min(z[i-left],right-i),0);
      for(;i+j<len&s[i+j]==s[j];j++);
      z[i]=j;
      if(i+z[i]>right) {
       right=i+z[i];
      left=i;
   } }
}
```

6.7 Manacher Algorithm O(n)

```
|// 求以每個字元為中心的最長回文半徑
// 頭尾以及每個字元間都加入一個
// 沒出現過的字元,這邊以'@'為例
// s為傳入的字串·len為字串長度
// z為儲存答案的陣列 (有包含'@'要小心)
                       "@a@b@a@a@c@"
// ex: s = "abaac" ->
                      [12141232121]
void z_value_pal(char *s,int len,int *z){
  len=(len<<1)+1;
  for(int i=len-1;i>=0;i--)
    s[i]=i&1?s[i>>1]:'@';
  z[0]=1;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i < r?min(z[l+l-i],r-i):1;
    \label{eq:while} \begin{aligned} & \text{while}(i-z[i]>=0\&\&i+z[i]<len\&\&s[i-z[i]]==s[i+z[i]]) \end{aligned}
        ++z[i];
    if(i+z[i]>r) l=i,r=i+z[i];
} }
```

6.8 Smallest Rotation

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
   if(a+k == b || s[a+k] < s[b+k])
     {b += max(0, k-1); break;}
  if(s[a+k] > s[b+k]) {a = b; break;}
```

```
} return a;
}
```

6.9 Cyclic LCS

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]=\{0,-1, -1,-1, -1,0\};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al, j=bl, l=0;
  while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++;
    i+=mov[dir][0];
    j+=mov[dir][1];
  return 1;
inline void reroot(int r) { // r = new base row
  int i=r,j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
  if(j>bl) return;
  pred[i][j]=L;
while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
       pred[i][j]=L;
    } else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
       i++;
       pred[i][j]=L;
    } else {
      j++;
int cyclic_lcs() {
   // a, b, al, bl should be properly filled
   // note: a WILL be altered in process
              -- concatenated after itself
  char tmp[MAXL];
  if(al>bl) {
    swap(al,bl);
    strcpy(tmp,a);
    strcpy(a,b);
    strcpy(b,tmp);
  }
  strcpy(tmp,a);
  strcat(a,tmp);
  // basic lcs
  for(int i=0;i<=2*al;i++) {
  dp[i][0]=0;</pre>
    pred[i][0]=U;
  for(int j=0;j<=bl;j++) {
  dp[0][j]=0;</pre>
    pred[0][j]=L;
  for(int i=1;i<=2*al;i++) {</pre>
    for(int j=1;j<=bl;j++) {
  if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;</pre>
       else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
       if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
       else if(a[i-1]==b[j-1]) pred[i][j]=LU;
       else pred[i][j]=U;
  } }
// do_cyclic lcs
  int clcs=0;
  for(int i=0;i<al;i++) {</pre>
    clcs=max(clcs,lcs_length(i));
    reroot(i+1);
  // recover a
  a[al]='\0':
  return clcs;
```

6.10 Hash

```
//字串雜湊前的idx是0-base,雜湊後為1-base
//即區間為 [0,n-1] -> [1,n]
//若要取得區間[L,R]的值則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為len的字串和
//(h[i+len-1] - h[i-1] * qpow(p, len) % modl + modl)
//從j開始長度為len的字串是否相同
#define x first
#define y second
pair<int,int> Hash[MXN];
void build(const string& s){
  pair<int,int> val = make_pair(0,0);
  Hash[0]=val;
  for(int i=1; i<=s.size(); i++){
val.x = (val.x * P1 + s[i-1]) % MOD;
val.y = (val.y * P2 + s[i-1]) % MOD;
  Hash[i] = val;
bool cmp( int i, int j, int len ) {
    return ((Hash[i+len-1].x-Hash[i-1].x*qpow(P1,len)%)
          MOD+MOD)\%\overline{M}OD == (\overline{Hash[j+len-1]}.x-\overline{Hash[j-1]}.x*
           qpow(P1,len)%MOD+MOD)%MOD)
     && ((Hash[i+len-1].y-Hash[i-1].y*qpow(P2,len)%MOD+
          MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow(
          P2, len)%MOD+MOD)%MOD);
}
```

7 Data Structure

7.1 Segment tree

```
//!!!注意build()時初始化用的陣列也是1-base
//!!!query(0, 0) 會報錯
#define cl(x) (x*2)
#define cr(x) (x*2+1)
struct segmentTree {
    int n;
    vector<int> seg, tag, cov;
    segmentTree(int _n): n(_n) {
         seg = tag = cov = vector<int>(n * 4, 0);
    void push(int i, int L, int R) {
         if(cov[i]) {
             seg[i] = cov[i] * (R - L + 1);
             if(L < R) {
                 cov[cl(i)] = cov[cr(i)] = cov[i];
                 tag[cl(i)] = tag[cr(i)] = 0;
             cov[i] = 0;
         if(tag[i]) {
             seg[i] += tag[i] * (R - L + 1);
             if(L < R) {
                 tag[cl(i)] += tag[i];
                 tag[cr(i)] += tag[i];
             tag[i] = 0;
    void pull(int i, int L, int R) {
         if(L >= R) return;
        int mid = L + R \gg 1;
        push(cl(i), L, mid);
        push(cr(i), mid + 1, R);
        seg[i] = seg[cl(i)] + seg[cr(i)];
    void build(vector<int>& arr, int i = 1, int L = 1,
         int R = -1) {
if(R == -1) R = n;
         if(L == R) return void(seg[i] = arr[L]);
         int mid = L + R \gg 1;
        build(arr, cl(i), L, mid);
build(arr, cr(i), mid + 1, R);
```

```
pull(i, L, R);
    int query(int rL, int rR, int i = 1, int L = 1, int
          R = -1) \{
         if(R == -1) R = n;
         push(i, L, R);
         if(rL <= L && R <= rR) return seg[i];</pre>
         int mid = L + R \gg 1, ret = 0;
         if(rL <= mid) ret += query(rL, rR, cl(i), L,</pre>
             mid);
         if(mid < rR ) ret += query(rL, rR, cr(i), mid +</pre>
              1, R);
         return ret;
    void update(int rL, int rR, int val, int i = 1, int
          L = 1, int R = -1) {
         if(R == -1) R = n;
         push(i, L, R);
         if(rL <= L && R <= rR) return void(tag[i] = val</pre>
             );
         int mid = L + R \gg 1;
         if(rL <= mid) update(rL, rR, val, cl(i), L, mid</pre>
         if(mid < rR ) update(rL, rR, val, cr(i), mid +</pre>
             1, R);
         pull(i, L, R);
    void cover(int rL, int rR, int val, int i = 1, int
         L = 1, int R = -1) {
        if(R == -1) R = n;
push(i, L, R);
         if(rL <= L && R <= rR) return void(cov[i] = val</pre>
         int mid = L + R \gg 1;
         if(rL <= mid) cover(rL, rR, val, cl(i), L, mid)</pre>
         if(mid < rR ) cover(rL, rR, val, cr(i), mid +</pre>
        1, R);
pull(i, L, R);
    }
};
/*
    Test Case:
    1 2 3 4
    2 1 3
    1 1 3 1
    2 1 3
    1 1 4 1
    2 1 4
```

7.2 持久化 SMT

```
struct node{
 node *l,
  int val;
vector<node *> ver;
int arr[MXN] = \{0\};
//0-base
struct SegmentTree{
 int n;
node *root;
  void build(int _n){
   n = _n;
    root = build(0, n-1);
 node* build(int L, int R){
    node *x = new node();
    if(L == R){x->val = arr[L]; return x;}
    int mid = (L+R)/2;
    x->l = build(L, mid);
   x->r = build(mid + 1, R);
    x->val = x->l->val + x->r->val;
    return x;
```

```
int query(node *ro, int L, int R){return query(ro, 0,
  n-1, L, R);}
int query(int L, int R){return query(root, 0, n-1, L,
       R);}
  int query(node *x, int L, int R, int recL, int recR){
    if(recL <= L && R <= recR) return x->val;
    int mid = (L+R)/2, res = 0;
    if(recL <= mid) res += query(x->1, L, mid, recL,
        recR);
    if(mid < recR) res += query(x->r, mid+1, R, recL,
       recR);
    return res;
  void update(int pos, int v){update(root, 0, n-1, pos,
  void update(node *x, int L, int R, int pos, int v){
    if(L == R){x->val = v; arr[L] = v; return;}
    int mid = (L+R)/2;
    if(pos <= mid) update(x->1, L, mid, pos, v);
                  update(x->r, mid+1, R, pos, v);
   x->val = x->l->val + x->r->val;
  node *update_ver(node *pre, int l, int r, int pos,
     int v){
    node *x = new node();
                           //當前位置建立新節點
    if(l == r){
     x->val = v;
     return x;
    int mid = (l+r)>>1;
    if(pos <= mid){ //更新左邊
     x->l = update_ver(pre->l, l, mid, pos, v); //左邊
          節點連向新節點
     x->r = pre->r; //右邊連到原本的右邊
    else{ //更新右邊
     x->l = pre->l; //左邊連到原本的左邊
     x->r = update_ver(pre->r, mid+1, r, pos, v); //
         右邊節點連向新節點
   x->val = x->l->val + x->r->val;
   return x;
}} seg;
                            //修改位置 x 的值為 v
void add_ver(int x,int v){
   ver.push_back(seg.update_ver(ver.back(), 0, seg.n
        -1, x, v));
}
```

7.3 Trie

```
struct trie{
 trie *nxt[26];
            ·//紀錄有多少個字串以此節點結尾
 int cnt;
            //有多少字串的前綴包括此節點
 int sz;
 trie():cnt(0),sz(0){
     memset(nxt,0,sizeof(nxt));
};
trie *root = new trie(); //創建新的字典樹
void insert(string& s){
 trie *now = root; // 每次從根結點出發
 for(auto i:s){
   now->sz++;
   if(now->nxt[i-'a'] == NULL){
     now->nxt[i-'a'] = new trie();
   now = now->nxt[i-'a']; //走到下一個字母
 now->cnt++; now->sz++;
}
int query_prefix(string& s){ //查詢有多少前綴為 s
 trie *now = root;
                    // 每次從根結點出發
 for(auto i:s){
   if(now->nxt[i-'a'] == NULL){
     return 0;
```

7.4 Treap (interval reverse)

```
// 拆 出 [a,b] 區 間 就 如 同 下 面 所 展 示 先 使 用 splitByTh() 拆 出
//左右,再把左區間拆成1, m最後merge()回去
//反轉區間時又記得使用^=可以直接反轉01
//treap 拆區間時從後面拆是因為這樣[a,b]的關係
//不用重新考慮,要是先拆前面b的位置會變成b-a+1
//0-base
//splitByTh(root,a-1,l,m);
//splitByTh(m,b-a+1,m,r);
mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
struct Treap {
  int key, pri, sz, tag, sum;
Treap *L, *R;
  Treap( int val ) {
    sum=key=val, pri=gen(), sz=1, tag=0;
    L=R=NULL;
};};
int Size( Treap *a ) { return !a?0:a->sz;}
void pull( Treap *a ) {
  a \rightarrow sz = Size(a \rightarrow L) + Size(a \rightarrow R) + 1;
  a->sum=a->key;
  if( a\rightarrow L ) a\rightarrow sum+=a\rightarrow L\rightarrow sum;
  if( a\rightarrow R ) a\rightarrow sum+=a\rightarrow R->sum;
void push( Treap *a ) {
  if( a && a->tag ) {
    swap(a->L,a->R);
    if( a\rightarrow L ) a\rightarrow L\rightarrow tag^{=1};
    if( a->R ) a->R->tag^=1;
    a \rightarrow tag=0;
Treap *merge(Treap *a, Treap *b) {
  if( !a || !b ) return a?a:b;
  push(a), push(b);
  if( a->pri > b->pri ) {
    a \rightarrow R = merge(a \rightarrow R, b);
    pull(a); return a;
  b \rightarrow L = merge(a, b \rightarrow L);
  pull(b); return b;
void print(Treap *a) {
  if( !a ) return;
  push(a);
  print(a->L);
  cout.put(a->key);
  print(a->R);
Treap *buildTreap( int n, string& str ) {
  Treap *root=NULL;
  for( int i=0 ; i < n ; i++ )</pre>
    root=merge(root,new Treap(str[i]));
  return root;
void splitbyk( Treap *x, int k, Treap *&a, Treap *&b )
  if(!x) a=b=NULL;
  else if( x->key <= k ) {
    a=x;
```

```
splitbyk(x->R,k,a->R,b);
    pull(a);
  else {
    splitbyk(x->L,k,a,b->L);
    pull(b);
  }
}
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
  if( !x ) { a=b=NULL; return; }
  push(x);
  if( Size(x->L)+1 <= k ) {
    splitByTh(x->R,k-Size(x->L)-1,a->R,b);
    pull(a);
  else {
    b=x:
    splitByTh(x->L,k,a,b->L);
    pull(b);
}
signed main() {
  string str;
  int n, m;
  cin>>n>>m>>str;
  Treap *root;
  root=buildTreap(n,str);
  for( int i=0 ; i < m ; i++ ) {
    int a, b;
    cin>>a>>b;
Treap *l, *m, *r;
    splitByTh(root,b,l,r);
    splitByTh(l,a-1,l,m);
    m->tag^=1;
    root=merge(l,merge(m,r));
  print(root);
}
```

7.5 Treap (interval erase)

```
//區間移除使用bitset維護區間值
mt19937 gen(chrono::steady_clock::now().
     time_since_epoch().count());
struct Treap {
char key;
int pri, sz;
bitset<128> tag;
  Treap *L, *R;
  Treap( char val ) {
     key=val, pri=gen(), sz=1;
     L=R=NULL:
    tag.set(key);
}; };
int Size( Treap *a ) { return !a?0:a->sz;}
void pull( Treap *a ) {
  if(!a) return;
  a \rightarrow sz = Size(a \rightarrow L) + Size(a \rightarrow R) + 1;
  a->tag=a->tag.reset();
  a->tag=a->tag.set(a->key);
  if( a->L ) a->tagl=a->L->tag;
  if( a \rightarrow R ) a \rightarrow tag = a \rightarrow R \rightarrow tag;
Treap *merge( Treap *a, Treap *b ) {
  if( !a || !b ) return a?a:b;
  if( a->pri > b->pri ) {
    a \rightarrow R = merge(a \rightarrow R, b);
    pull(a);
     return a;
  b->L=merge(a,b->L);
  pull(b);
  return b;
Treap *buildTreap( int n, string& str ) {
  Treap *root=NULL;
  for( int i=0 ; i < n ; i++ )</pre>
```

```
root=merge(root, new Treap(str[i]));
  return root;
void print( Treap *a ) {
  if(!a) return;
 print(a->L);
  cout.put(a->key);
 print(a->R);
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
  if( !x ) { a=b=NULL; return; }
  if( Size(x->L)+1 \le k ) {
    splitByTh(x->R,k-Size(x->L)-1,a->R,b);
   pull(a);
  else {
    b=x:
    splitByTh(x->L,k,a,b->L);
    pull(b);
void erase( Treap *&x, char ch ) {
 if( !x || !x->tag.test(ch) ) return;
  erase(x->L,ch);
  erase(x->R,ch);
  if( x \rightarrow key == ch ) {
    Treap *l=x->L, *r=x->R;
    x=NULL;
   x=merge(l,r);
 pull(x);
signed main() {
 string str;
  int n, m;
  cin>>n>>m>>str;
 Treap *root;
  root=buildTreap(n,str);
  for( int i=0 ; i < m ; i++ ) {</pre>
    char c;
    int a, b;
    cin>>a>>b>>c;
   Treap *1, *m, *r;
if( !root || !root->tag.test(c) ) continue;
    splitByTh(root,b,l,r);
    splitByTh(l,a-1,l,m);
    if( m || !m->tag.test(c) ) erase(m,c);
    root=merge(l,merge(m,r));
 print(root);
```

7.6 Link-Cut Tree

```
struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
 int val, rev, size;
Splay (int _val=-1) : val(_val), rev(0), size(1)
  \{ f = ch[0] = ch[1] = &nil; \}
 bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir()
  { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d){
    ch[d] = c;
    if (c != &nil) c->f = this;
   pull();
  void push(){
    if( !rev ) return;
    swap(ch[0], ch[1]);
    if (ch[0] != &nil) ch[0]->rev ^= 1;
    if (ch[1] != &nil) ch[1]->rev ^= 1;
    rev=0;
  void pull(){
    size = ch[0] -> size + ch[1] -> size + 1;
```

```
if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::
Splay *nil = &Splay::nil;
void rotate(Splay *x){
  Splay *p = x->f
  int d = x - sdir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(); x->pull();
}
vector<Splay*> splayVec;
void splay(Splay *x){
  splayVec.clear();
  for (Splay *q=x;; q=q->f){
    splayVec.push_back(q);
    if (q->isr()) break;
  reverse(begin(splayVec), end(splayVec));
for (auto it : splayVec) it->push();
while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir())
      rotate(x->f),rotate(x);
    else rotate(x), rotate(x);
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x){
  Splay *q = nil;
  for (;x!=nil;x=x->f){
    splay(x)
    x \rightarrow setCh(q, 1);
    q = x;
  }
  return q;
}
void chroot(Splay *x){
  access(x);
  splay(x);
  x->rev ^= 1;
  x->push(); x->pull();
void link(Splay *x, Splay *y){
  access(x);
  splay(x)
  chroot(y);
  x->setCh(y, 1);
void cut_p(Splay *y) {
  access(y);
  splay(y)
  y->push();
  y - ch[0] = y - ch[0] - f = nil;
}
void cut(Splay *x, Splay *y){
  chroot(x);
  cut_p(y);
Splay* get_root(Splay *x) {
  access(x);
  splay(x);
  for(; x - ch[0] != nil; x = x - ch[0])
    x->push();
  splay(x);
  return x:
bool conn(Splay *x, Splay *y) {
  x = get_root(x);
  y = get_root(y);
  return x == y;
Splay* lca(Splay *x, Splay *y) {
  access(x);
  access(y);
  splay(x);
  if (x->f == nil) return x;
  else return x->f;
```

7.7 BIT

|}

```
#define lowbit(x) (x&-x)
struct BIT {
    int n;
    vector<int> bit;
    BIT(int _n):n(_n), bit(n + 1) {}
    void update(int x, int val) {
        for(; x <= n; x += lowbit(x)) bit[x] += val;</pre>
    void update(int L, int R, int val) {
        update(L, val), update(R + 1, -val);
    int query(int x) {
        int res = 0;
        for(; x; x -= lowbit(x)) res += bit[x];
        return res;
    int query(int L, int R) {
        return query(R) - query(L - 1);
};
```

7.8 Black Magic

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
    tree_order_statistics_node_update> set_t;
tree<int,null_type,less_equal<int>,rb_tree_tag,
    tree_order_statistics_node_update> mt_t;
#include <ext/pb_ds/assoc_container.hpp>
typedef cc_hash_table<int,int> umap_t;
// gp_hash_table<int, int>
typedef priority_queue<int> heap;
#include<ext/rope>
using namespace __gnu_cxx;
int main(){
  // Insert some entries into s.
  set_t s; s.insert(12); s.insert(505);
  // The order of the keys should be: 12, 505.
 assert(*s.find_by_order(0) == 12);
assert(*s.find_by_order(3) == 505);
 // The order of the keys should be: 12, 505.
 assert(s.order_of_key(12) == 0)
 assert(s.order_of_key(505) == 1);
 // Erase an entry.
  s.erase(12);
  // The order of the keys should be: 505.
 assert(*s.find_by_order(0) == 505);
  // The order of the keys should be: 505.
 assert(s.order_of_key(505) == 0);
 // if we want to delete less_equal tag tree
 mt_t.erase(mt_t.find_by_order(mt_t.order_of_key(val))
      );
 heap h1 , h2; h1.join( h2 );
  rope<char> r[ 2 ];
 r[1] = r[0]; // persistenet
string t = "abc";
 r[1].insert(0, t.c_str());
r[1].erase(1,1);
  cout << r[ 1 ].substr( 0 , 2 );</pre>
```

8 Others

8.1 SOS dp

```
| for(int i = 0; i<(1<<N); ++i)
| F[i] = A[i];
```

```
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1<< N); ++mask){
  if(mask & (1<<i))
    F[mask] += F[mask^(1<<i)];
}</pre>
```

8.2 De Brujin sequence

```
// return cyclic array of length k^n such that every
// array of length n using 0~k-1 appears as a subarray.
vector<int> DeBruijn(int k,int n){
   if(k==1) return {0};
   vector<int> aux(k*n),res;
   function<void(int,int)> f=[&](int t,int p)->void{
      if(t>n){   if(n%p==0)
          for(int i=1;i<=p;++i) res.push_back(aux[i]);
   }else{
      aux[t]=aux[t-p]; f(t+1,p);
      for(aux[t]=aux[t-p]+1;aux[t]<k;++aux[t]) f(t+1,t)
      ;
   }
   };
   f(1,1); return res;
}</pre>
```

8.3 CDQ 分治

```
//cdq分治使用的結構u, v, w為排序物的三個維度
//ans記錄了有幾項三維都小於等於自己
//cnt記錄了相同物有幾個·在使用cdq之前必先去重
//並且將相同元素紀錄至cnt中,可使用map來做到這步
//cdq使用的BIT就是普通求和的BIT·大小就開維度的
//值域範圍·若值域大於2e6則要先進行離散化
struct triple {int u, v, w, ans, cnt;};
BIT *bt:
void cdq(int L, int R, vector<triple>& arr) {
  if(R - L <= 1) return;</pre>
  int mid = L + R \gg 1;
  vector<triple> temp;
  cdq(L, mid, arr), cdq(mid, R, arr);
for(int i = L, j = mid; i < mid || j < R;) {</pre>
    for(; i < mid && (j >= R || arr[i].v <= arr[j].v);</pre>
        i++) {
      bt->update(arr[i].w, arr[i].cnt);
      temp.push_back(arr[i]);
    if(j < R) {
      arr[j].ans += bt->query(arr[j].w);
      temp.push_back(arr[j]);
      j++;
    }
  for(int i = L; i < mid; i++)</pre>
    bt->update(arr[i].w, -arr[i].cnt);
  copy(temp.begin(), temp.end(), arr.begin() + L);
signed main()
  // n 個數 k 值域範圍
  int n, k;
  cin >> n >> k;
  map<tuple<int, int, int>, int> mp;
  vector<int> res(n, 0);
  vector<triple> arr;
  bt = new BIT(k + 1);
for(int i = 0; i < n; i++) {
      int x, y, z;
cin >> x >> y >> z;
      mp[{x, y, z}]++;
  for(auto t : mp)
    arr.push_back({get<0>(t.first), get<1>(t.first),
        get<2>(t.first), 0, t.second});
  cdq(0, arr.size(), arr);
  for(auto &[x,y,z,a,b] : arr) res[a + b - 1] += b;
for(int i : res) cout << i << '\n';</pre>
```

8.4 3D LIS

```
#define lowbit(x) (x&-x)
const int MAXN=1e5+5;
struct BIT {
  int n;
  vector<int> bit;
BIT( int _n ):n(_n), bit(_n+1,0) {}
  int query( int x ) {
    int res=0;
    for(; x > 0; x-=lowbit(x)) res=max(res,bit[x]);
    return res:
  void update( int x, int val ) {
  for(; x <= n ; x+=lowbit(x) ) {
    if( val < 0 ) bit[x]=0;</pre>
       else bit[x]=max(bit[x],val);
    }
}bt(MAXN);
struct triple {
  int u, v, w, ans, cnt;
  bool operator<( triple b ) { return u<b.u; }</pre>
bool cmp( triple a, triple b ) {return a.v<b.v;}</pre>
void cdq( int L, int R, vector<triple>& arr ) {
  if( R-L <= 1 ) return;</pre>
  int mid=L+R>>1;
  cdq(L,mid,arr);
  sort(arr.begin()+L,arr.begin()+mid,cmp);
  sort(arr.begin()+mid,arr.begin()+R,cmp);
  for( int i=L, j=mid ; i < mid || j < R ; ) {
  for(; i < mid && ( j >= R || arr[i].v < arr[j].v )</pre>
           i++ ) bt.update(arr[i].w,arr[i].ans);
    if(j < R) {
      arr[j].ans=max(bt.query(arr[j].w-1)+1,arr[j].ans)
      j++;
    }
  for( int i=L ; i < mid ; i++ ) bt.update(arr[i].w,-1)</pre>
  sort(arr.begin()+L,arr.begin()+R);
  cdq(mid,R,arr);
signed main()
  ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0)
  int n, res=0;
  cin>>n;
  vector<int> ls;
  vector<triple> arr;
  for( int i=0 ; i < n ; i++ ) {</pre>
    int a, b;
    cin>>a>>b;
    arr.push\_back({i,a,b,1,1});//{第一維,第二維,第三維,
         答案,數量}
    ls.push_back(b);
  sort(ls.begin(),ls.end());
  ls.resize(unique(ls.begin(),ls.end())-ls.begin());
  for( auto &t : arr ) t.w=lower_bound(ls.begin(),ls.
       end(),t.w)-ls.begin()+1;
  n=arr.size();
  cdq(0,n,arr);
  for( int i=0 ; i < n ; i++ ) res=max(res,arr[i].ans);</pre>
  cout<<res<<'\n';
```

8.5 Ternary Search

```
while(L <= R) {
   int ml = L + (R - L) / 3, mr = R - (R - L) / 3;
   if(L == R) return L;
   else if( checker(ml) < checker(mr) ) L = ml + 1;
   else R = mr - 1;
}</pre>
```

8.6 Max Subrectangle

```
const int N = 1e5+5;
int n, a[N], l[N], r[N];
long long ans;
int main() {
  while (cin>>n) {
    ans = 0;
    for (int i = 1; i \le n; i++) cin>>a[i], l[i] = r[i]
    for (int i = 1; i <= n; i++)
      while (l[i] > 1 \& a[i] <= a[l[i] - 1]) l[i] = l[
          l[i] - 1];
    for (int i = n; i >= 1; i--)
      while (r[i] < n \&\& a[i] <= a[r[i] + 1]) r[i] = r[
          r[i] + 1];
    for (int i = 1; i <= n; i++)
      ans = max(ans, (long long)(r[i] - l[i] + 1) * a[i]
          ]);
    cout<<ans<<"\n";
  }
}
```

8.7 Maximal Rectangle

```
const int MXN = 300;
int maximalRectangle(vector<vector<char>>& matrix) {
     int a[MXN]{}, l[MXN]{}, r[MXN]{};
     int n = matrix.size(), m = matrix[0].size(), ans =
     for(int i = 1; i <= n; i++) {
         for(int j = 1; j \leftarrow m; j++) l[j] = r[j] = j;
         char c;
         for(int j = 1; j <= m; j++) { //對每一個直行做
統計·若是上一個a[j]也是1則會變成2
              c = matrix[i - 1][j - 1];
              if (c == '\bar{1}') a[j]++;
              else if (c == '0') a[j] = 0;
         for(int j = 1; j <= m; j++) while(l[j] != 1 &&
    a[l[j] - 1] >= a[j]) l[j] = l[l[j] - 1];
for(int j = m; j >= 1; j--) while(r[j] != m &&
              a[r[j] + 1] >= a[j] r[j] = r[r[j] + 1];
         return ans;
}
```

8.8 p-Median

8.9 Tree Knapsack

```
int dfs(int u) {
   int p = 1;
                                                                                       void solve(string &s) {
   dp[u][1] = s[u];
                                                                                            getfail();
                                                                                            query(s);
   for (int v : edge[u]) {
     int siz = dfs(v);
                                                                                            topu();
     for (int i = min(p, m + 1); i; i--)
  for (int j = 1; j <= siz && i + j <= m + 1; j++)
    dp[u][i + j] = max(dp[u][i + j], dp[u][i] + dp[</pre>
                                                                                            for (int i = 1; i <= n; i++) ans[i] = vis[rev[i</pre>
                                                                                                  ]];
                                                                                } AC;
                v][j]);
     p += siz;
   return p;
}
```

8.10 AC-Automaton

```
// 1-based
// n is the number of patterns
struct Automaton {
    static const int MXN = 1e6;
    int n, cnt, vis[MXN], rev[MXN], indeg[MXN], ans[MXN
        ];
    queue<int> q;
    struct trie_node {
        vector<int> son;
        int fail, flag, ans;
        trie_node(): son(27), fail(0), flag(0) {}
    } trie[MXN];
    void init(int _n) {
        n = n, cnt = 1;
        for (int i = 1; i <= n; i++) vis[i] = 0;
    // insert a string s with number num
    // num is the index of the pattern
    void insert(string s, int num) {
        int u = 1, len = s.size();
        for (int i = 0; i < len; i++) {
   int v = s[i] - 'a';
             if (!trie[u].son[v]) trie[u].son[v] = ++cnt
             u = trie[u].son[v];
        if (!trie[u].flag) trie[u].flag = num;
        rev[num] = trie[u].flag;
    void getfail() {
        for (int i = 0; i < 26; i++) trie[0].son[i] =
        q.push(1);
        trie[1].fail = 0;
        while (q.size()) {
             int u = q.front(); q.pop();
             int Fail = trie[u].fail;
             for (int i = 0; i < 26; i++) {
                 int v = trie[u].son[i];
                 if (!v) {
                      trie[u].son[i] = trie[Fail].son[i];
                      continue;
                 trie[v].fail = trie[Fail].son[i];
                 indeg[trie[Fail].son[i]]++;
                 q.push(v);
            }
        }
    void topu() {
        for (int i = 1; i <= cnt; i++)
    if (!indeg[i]) q.push(i);</pre>
        while (q.size()) {
             int fr = q.front(); q.pop();
             vis[trie[fr].flag] = trie[fr].ans;
             int u = trie[fr].fail;
             trie[u].ans += trie[fr].ans;
             if (!--indeg[u]) q.push(u);
        }
    void query(string &s) {
        int u = 1, len = s.size();
        for (int i = 0; i < len; i++) u = trie[u].son[s
    [i] - 'a'], trie[u].ans++;</pre>
```



