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1 Basic

1.1 default code

```
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
using namespace std;
ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
```

1.2 .vimrc

```
set nu rnu ts=4 sw=4 bs=2 ai hls cin mouse=a
color default
sy on
inoremap {<CR> {<CR>}<C-o>0
inoremap jk <Esc>
nnoremap J 5j
nnoremap K 5k
nnoremap run :w<bar>!g++ -std=c++14 -DLOCAL -Wfatal-
errors -o test "%" && echo "done." && time ./test<
CR>
```

1.3 Increase Stack Size (linux)

```
#include <sys/resource.h>
void increase_stack_size() {
  const rlim_t ks = 64*1024*1024;
  struct rlimit rl;
  int res=getrlimit(RLIMIT_STACK, &rl);
  if(res==0){
    if(rl.rlim_cur<ks){
      rl.rlim_cur=ks;
      res=setrlimit(RLIMIT_STACK, &rl);
} } }</pre>
```

1.4 Misc

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```
編譯參數: -std=c++14 -Wall -Wshadow (-fsanitize= undefined)

mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }

#define SECs ((double)clock() / CLOCKS_PER_SEC)
double startTime;
bool TIME() { // 比最大可執行時間小一點
    return SECs - startTime > 0.8;
}
int main() {
    startTime = SECs;
}
```

1.5 check

```
for ((i=0;;i++))
do

    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done
```

1.6 python-related

```
parser:
int(eval(num.replace("/","//")))
from fractions import Fraction
from decimal import Decimal, getcontext, ROUND_HALF_UP,
      ROUND_CEILING, ROUND_FLOOR
getcontext().prec = 250 # set precision
getcontext().rounding = ROUND_HALF_UP
itwo = Decimal(0.5)
two = Decimal(2)
format(x, '0.10f') # set precision
N = 200
def angle(cosT):
    """given cos(theta) in decimal return theta"""
  for i in range(N):
  cosT = ((cosT + 1) / two) ** itwo 
 sinT = (1 - cosT * cosT) ** itwo 
 return sinT * (2 ** N)
pi = angle(Decimal(-1))
"""round to 2 decimal places"""
sum = Decimal(input())
sum.quantize(Decimal('.00'), ROUND_HALF_UP)
"""Fraction"""
x = Fraction(1, 3) # 1/3
x.as_integer_ratio() # (1, 3)
"""input list of integers"""
arr = list(map(int, input().split()))
```

2 flow

2.1 ISAP $O(V^3)$

```
struct Maxflow {
    static const int MAXV = 20010;
    static const int INF = 1000000;
    struct Edge {
        int v, c, r;
        Edge(int _v, int _c, int _r):
            v(_v), c(_c), r(_r) {}
    };
    int s, t;
    vector<Edge> G[MAXV*2];
    int iter[MAXV*2], d[MAXV*2], gap[MAXV*2], tot;
```

```
void init(int x) {
    tot = x+2;
    s = x+1, t = x+2;
    for(int i = 0; i <= tot; i++) {</pre>
      G[i].clear();
       iter[i] = d[i] = gap[i] = 0;
  void addEdge(int u, int v, int c) {
    G[u].push_back(Edge(v, ć, SZ(G[v]) ));
G[v].push_back(Edge(u, 0, SZ(G[u]) - 1));
  int dfs(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
       Edge &e = G[p][i];
       if(e.c > 0 \& d[p] == d[e.v]+1) {
         int f = dfs(e.v, min(flow, e.c));
         if(f) {
           e.c -= f;
           G[e.v][e.r].c += f;
           return f;
    } } }
    if((--gap[d[p]]) == 0) d[s] = tot;
    else {
      d[p]++;
      iter[p] = 0;
      ++gap[d[p]];
    return 0;
  int solve() {
    int res = 0;
    gap[0] = tot;
    for(res = 0; d[s] < tot; res += dfs(s, INF));</pre>
    return res;
  void reset() {
    for(int i=0;i<=tot;i++) {</pre>
      iter[i]=d[i]=gap[i]=0;
} } flow;
```

2.2 MinCostFlow

```
struct zkwflow{
  static const int maxN=10000;
  struct Edge{ int v,f,re; ll w;};
int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
vector<Edge> E[maxN];
  void init(int _n,int _s,int _t){
    n=_n,s=_s,t=_t;
    for(int i=0;i<n;i++) E[i].clear();</pre>
  void addEdge(int u,int v,int f,ll w){
    E[u].push_back({v,f,(int)E[v].size(),w});
    E[v].push_back(\{u,0,(int)E[u].size()-1,-w\});
  bool SPFA(){
    fill_n(dis,n,LLONG_MAX); fill_n(vis,n,false);
    queue<int> q; q.push(s); dis[s]=0;
    while (!q.empty()){
  int u=q.front(); q.pop(); vis[u]=false;
  for(auto &it:E[u]){
         if(it.f>0&&dis[it.v]>dis[u]+it.w){
           dis[it.v]=dis[u]+it.w;
           if(!vis[it.v]){
             vis[it.v]=true; q.push(it.v);
    return dis[t]!=LLONG_MAX;
  int DFS(int u,int nf){
    if(u==t) return nf;
    int res=0; vis[u]=true;
    for(int &i=ptr[u];i<(int)E[u].size();i++){</pre>
      auto &it=E[u][i]
       if(it.f>0&&dis[it.v]==dis[u]+it.w&&!vis[it.v]){
         int tf=DFS(it.v,min(nf,it.f));
         res+=tf,nf-=tf,it.f-=tf;
         E[it.v][it.re].f+=tf;
         if(nf==0){ vis[u]=false; break; }
```

```
return res;
}
pair<int,ll> flow(){
  int flow=0; ll cost=0;
  while (SPFA()){
    fill_n(ptr,n,0);
    int f=DFS(s,INT_MAX);
    flow+=f; cost+=dis[t]*f;
}
  return{ flow,cost };
} // reset: do nothing
} flow;
```

2.3 Dinic $O(V^2E)$

```
#define SZ(x) (int)x.size()
#define PB push_back
struct Dinic{
  struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
    n = _n;    s = _s;    t = _t;

    for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
    E[u].PB({v,f,SZ(E[v])});
    E[v].PB({u,0,SZ(E[u])-1});
  bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;</pre>
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()){
      int u = que.front(); que.pop();
      for (auto it : E[u]){
  if (it.f > 0 && level[it.v] == -1){
           level[it.v] = level[u]+1;
           que.push(it.v);
    } } }
    return level[t] != -1;
  int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
      if (it.f > 0 && level[it.v] == level[u]+1){
        int tf = DFS(it.v, min(nf,it.f));
         res += tf; nf -= tf; it.f -= tf;
         E[it.v][it.re].f += tf;
         if (nf == 0) return res;
    if (!res) level[u] = -1;
    return res;
  int flow(int res=0){
    while ( BFS() )
      res += DFS(s,2147483647);
    return res;
} }flow;
```

2.4 Kuhn Munkres 最大完美二分匹配 $O(n^3)$

```
struct KM{ // max weight, for min negate the weights
  int n, mx[MXN], my[MXN], pa[MXN];
  ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void init(int _n) { // 1-based
    n = _n;
    for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
  }
  void addEdge(int x, int y, ll w) {g[x][y] = w;}
  void augment(int y) {
    for(int x, z; y; y = z)
        x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
  }</pre>
```

```
void bfs(int st) {
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
     queue<int> q; q.push(st);
     for(;;) {
       while(q.size()) {
         int x=q.front(); q.pop(); vx[x]=1;
          for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
            ll t = lx[x]+ly[y]-g[x][y];
            if(t==0){
              pa[y]=x
              if(!my[y]){augment(y);return;}
              vy[y]=1, q.push(my[y]);
            }else if(sy[y]>t) pa[y]=x,sy[y]=t;
       } }
       il cut = INF;
       for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;</pre>
          if(vy[j]) ly[j] += cut;
          else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
  if(!my[y]){augment(y);return;}</pre>
         vy[y]=1, q.push(my[y]);
    } }
  ll solve(){
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
     fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
    lx[x] = max(lx[x], g[x][y]);</pre>
     for(int x=1; x<=n; ++x) bfs(x);
     ll ans = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
     return ans;
} }graph;
```

```
2.5 Flow Method
Maximize c^T x subject to Ax \le b, x \ge 0;
with the corresponding symmetric dual problem,
Minimize b^T y subject to A^T y \ge c, y \ge 0.
Maximize c^Tx subject to Ax \le b;
with the corresponding asymmetric dual problem,
Minimize b^Ty subject to A^Ty=c,y\geq 0.
Minimum vertex cover on bipartite graph =
Maximum matching on bipartite graph
Minimum edge cover on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
König Theorem
最小點覆蓋:選出最少的點,滿足每條邊至少有一個端點被選
 「分圖中・最小點覆蓋 = 最大匹配
Independent set on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
__分圖中·最大獨立集 =n- 最小點覆蓋
找出最小點覆蓋,做完dinic之後
從源點dfs只走還有流量的邊
左邊沒被走到的點跟右邊被走到的點就是答案,
其他點為最大獨立集
最大閉包(最大權閉合子圖)
源點連到所有正權點流量為點權
所有負權點連到匯點流量為點權(絕對值)
所有圖上的邊權重為 INF
路徑覆蓋數量
把每個點拆成 入點 和 出點,轉化為二分圖
原圖頂點數 - 二分圖最大匹配數
Maximum density subgraph (\sum W_e + \sum W_v) / IVI
Binary search on answer:
```

For a fixed D, construct a Max flow model as follow:

3 Math

3.1 FFT

```
// const int MXN = 262144 (MXN must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx; //real() ,imag()
const ld PI = acosl(-1);
const cplx I(0, 1);
struct FFT{
  cplx omega[MXN+1];
  FFT(){ //pre_fft
    for(int i=0; i<=MXN; i++)
  omega[i] = exp(i * 2 * PI / MXN * I);</pre>
  // n must be 2^k
  void fft(int n, cplx a[], bool inv=false){
    int basic = MXN / n;
    int theta = basic;
    for (int m = n; m >= 2; m >>= 1) {
      int mh = m >> 1;
       for (int i = 0; i < mh; i++) {
      cplx w = omega[inv ? MXN-(i*theta%MXN) : i*theta%
           MXN];
       for (int j = i; j < n; j += m) {
        int k = j + mh;
cplx x = a[j] - a[k];
         a[j] += a[k];
        a[k] = w * x;
      theta = (theta * 2) % MXN;
    for (int j = 1; j < n - 1; j++) {
       for (int k = n \gg 1; k \gg (i ^= k); k \gg = 1);
       if (j < i) swap(a[i], a[j]);</pre>
     if(inv) for (i = 0; i < n; i++) a[i] /= n;
  cplx arr[MXN+1];
  inline void mul(int _n,ll a[],int _m,ll b[],ll ans[])
       {
    int n=1,sum=_n+_m-1;
    while(n<sum)</pre>
      n <<=1;
     for(int i=0;i<n;i++) {</pre>
      double x=(i<_n?a[i]:0), y=(i<_m?b[i]:0);
      arr[i]=complex<double>(x+y,x-y);
    fft(n,arr);
    for(int i=0;i<n;i++)</pre>
       arr[i]=arr[i]*arr[i];
     fft(n,arr,true);
     for(int i=0;i<sum;i++)</pre>
      ans[i]=(long long)(arr[i].real()/4+0.5);
}fft;
```

3.2 Faulhaber $(\sum_{i=1}^{n} i^p)$

```
/* faulhaber's formula -
* cal power sum formula of all p=1~k in O(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinactories
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
  while(b) {
    int q,t;
    q=a/b; t=b; b=a-b*q; a=t;
     t=b0; b0=a0-b0*q; a0=t;
    t=b1; b1=a1-b1*q; a1=t;
  return a0<0?a0+mod:a0;
inline void pre() {
  /* combinational
  for(int i=0;i<=MAXK;i++) {</pre>
    cm[i][0]=cm[i][i]=1;
    for(int_j=1;j<i;j++)</pre>
       cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);
  /* inverse */
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
  /* bernoulli */
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
  for(int i=2;i<MAXK;i++) {</pre>
    if(i&1) { b[i]=0; continue; }
    b[i]=1;
    for(int j=0;j<i;j++)</pre>
       b[i]=sub(b[i]
                 mul(cm[i][j],mul(b[j], inv[i-j+1])));
  /* faulhaber */
  // sigma_x=1~n {x^p} =
// 1/(p+1) * sigma_j=0~p {C(p+1,j)*Bj*n^(p-j+1)}
  for(int i=1;i<MAXK;i++) {</pre>
    co[i][0]=0;
    for(int j=0; j<=i; j++)</pre>
       co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]))
  }
}
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
inline int solve(int n,int p) {
  int sol=0, m=n;
  for(int i=1;i<=p+1;i++) {</pre>
    sol=add(sol,mul(co[p][i],m));
    m = mul(m, n);
  return sol;
```

3.3 Chinese Remainder

3.4 Miller Rabin

```
// n < 4,759,123,141
                             3: 2, 7, 61
// n < 1,122,004,669,633
                             4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383
                                    6 : pirmes <= 13
// n < 2^{^{\circ}}64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL mul(LL x,LL y,LL mod){
  LL ret=x*y-(LL)((long double)x/mod*y)*mod;
  // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
  return ret<0?ret+mod:ret;</pre>
LL magic[]={}
bool witness(LL a, LL n, LL u, int t){
  if(!a) return 0;
  LL x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    LL nx=mul(x,x,n);
    if(nx==1&&x!=1&&x!=n-1) return 1;
    x=nx;
  }
  return x!=1;
bool miller_rabin(LL n) {
  int s=(magic number size)
  // iterate s times of witness on n
  if(n<2) return 0;</pre>
  if(!(n\&1)) return n == 2;
  ll u=n-1; int t=0;
  // n-1 = u*2^t
  while(!(u&1)) u>>=1, t++;
  while(s--){
    LL a=magic[s]%n;
    if(witness(a,n,u,t)) return 0;
  return 1;
}
```

3.5 Pollard Rho

```
// does not work when n is prime O(n^(1/4))
LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
LL pollard_rho(LL n) {
   if(!(n&1)) return 2;
   while(true){
      LL y=2, x=rand()%(n-1)+1, res=1;
      for(int sz=2; res==1; sz*=2) {
        for(int i=0; i<sz && res<=1; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
        }
        y = x;
      }
      if (res!=0 && res!=n) return res;
}</pre>
```

3.6 Josephus Problem

```
int josephus(int n, int m){ //n人每m次
  int ans = 0;
  for (int i=1; i<=n; ++i)
     ans = (ans + m) % i;
  return ans;
}</pre>
```

3.7 Gaussian Elimination

```
const int GAUSS_MOD = 100000007LL;
struct GAUSS{
  int n;
  vector<vector<int>> v;
  int ppow(int a , int k){
```

```
if(k == 0) return 1;
    if(k % 2 == 0) return ppow(a * a % GAUSS_MOD , k >>
    1); if(k % 2 == 1) return ppow(a * a % GAUSS_MOD , k >>
          1) * a % GAUSS_MOD;
  vector<int> solve(){
    vector<int> ans(n);
    REP(now , 0 , n){
      REP(i , now , n) if(v[now][now] == 0 && v[i][now]
            != 0)
      swap(v[i] , v[now]); // det = -det;
if(v[now][now] == 0) return ans;
      int inv = ppow(v[now][now] , GAUSS_MOD - 2);
REP(i , 0 , n) if(i != now){
         int tmp = v[i][now] * inv % GAUSS_MOD;
         }
    REP(i , 0 , n) ans[i] = v[i][n + 1] * ppow(v[i][i] , GAUSS_MOD - 2) % GAUSS_MOD;
    return ans;
  // gs.v.clear() , gs.v.resize(n , vector<int>(n + 1 ,
        0));
} gs;
```

3.8 Inverse Matrix

0));

```
int GAUSS_MOD;
struct GAUSS{
  int n;
  vector<vector<int> > v;
  vector<vector<int> > rev;
  int mul(int x,int y,int mod){
    int ret=x*y-(int)((long double)x/mod*y)*mod;
    return ret<0?ret+mod:ret;</pre>
  int ppow(int a, int b){//res=(a^b)%m
    int res=1, k=a;
    while(b){
      if((b&1)) res=mul(res,k,GAUSS_MOD)%GAUSS_MOD;
      k=mul(k,k,GAUSS_MOD)%GAUSS_MOD;
      b>>=1:
    return res%GAUSS_MOD;
  bool solve(){
    for(int now = 0; now < n; now++){
      int ch;
      for(ch = now; ch < n && !v[ch][now]; ch++);</pre>
      if(ch >= n) return 0;
      for(int i = now; i < n; i++) if(v[now][now] == 0
           && v[i][now] != 0){
          swap(v[i] , v[now]); // det = -det;
swap(rev[i], rev[now]);
      if(v[now][now] == 0) return 0;
      GAUSS_MOD) %= GAUSS_MOD;
(rev[i][j] += GAUSS_MOD - tmp * rev[now][j] %
GAUSS_MOD) %= GAUSS_MOD;
        }
      }
    }
    return 1;
}} gs;
signed main(){
  int n, p; //n*n matrix, MOD=p
  cin>>n>>p; //if(!n && !p) return 0;
  GAUSS\_MOD = p; gs.n = n;
  gs.v.clear() , gs.v.resize(n + 1, vector<int>(n + 2 ,
```

3.9 模反元素

```
|long long inv(long long a,long long m){
    long long x,y;
    long long d=exgcd(a,m,x,y);
    if(d==1) return (x+m)%m;
    else return -1; //-1為無解
}
```

3.10 ax+by=gcd

```
PII gcd(int a, int b){
   if(b == 0) return {1, 0};
   PII q = gcd(b, a % b);
   return {q.second, q.first - q.second * (a / b)};
}
int exgcd(int a,int b,long long &x,long long &y) {
   if(b == 0){x=1,y=0;return a;}
   int now=exgcd(b,a%b,y,x);
   y-=a/b*x;
   return now;
}
```

3.11 Discrete sqrt

```
void calcH(LL &t, LL &h, const LL p) {
   LL tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
}

// solve equation x^2 mod p = a
bool solve(LL a, LL p, LL &x, LL &y) {
   if(p == 2) { x = y = 1; return true; }
   int p2 = p / 2, tmp = mypow(a, p2, p);
   if (tmp == p - 1) return false;
   if ((p + 1) % 4 == 0) {
        x = mypow(a, (p+1)/4, p); y=p-x; return true;
   } else {
        LL t, h, b, pb; calcH(t, h, p);
        if (t >= 2) {
            do {b = rand() % (p - 2) + 2;
            } while (mypow(b, p / 2, p) != p - 1);
            pb = mypow(a, h / 2, p);
        for (int step = 2; step <= t; step++) {
            int ss = (((LL)(s * s) % p) * a) % p;
            for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);
            if (ss + 1 == p) s = (s * pb) % p;
            pb = ((LL)pb * pb) % p;
            } x = ((LL)s * a) % p; y = p - x;
      } return true;
}</pre>
```

3.12 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
  for( int i = 2 ; i < m ; i ++ )
    inv[ i ] = ((LL)(m - m / i) * inv[m % i]) % m;
}</pre>
```

3.13 Roots of Polynomial 找多項式的根

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ]; // a[0..n](coef) must be
    filled
int n; // degree of polynomial must be filled
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a[], int n, double x){
  double tmp=1,sum=0;
  for(int i=0;i<=n;i++)</pre>
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double l,double r,double a[],int n){
  int sl=sign(f(a,n,l)), sr=sign(f(a,n,r));
if(sl==0) return l; if(sr==0) return r;
  if(sl*sr>0) return inf;
  while(r-l>eps){
     double mid=(l+r)/2;
     int ss=sign(f(a,n,mid));
     if(ss==0) return mid;
     if(ss*sl>0) l=mid; else r=mid;
  return 1:
void solve(int n,double a[],double x[],int &nx){
   if(n==1){ x[1]=-a[0]/a[1]; nx=1; return; }
  double da[10], dx[10]; int ndx;
  for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0;
  if(ndx==0){
     double tmp=binary(-inf,inf,a,n);
     if (tmp<inf) x[++nx]=tmp;</pre>
    return;
  double tmp;
tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;
for(int i=1;i<=ndx-1;i++){</pre>
     tmp=binary(dx[i],dx[i+1],a,n);
     if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
} // roots are stored in x[1..nx]
```

3.14 Combination thearom

```
const ll mod = 1e9 + 7;
ll fac[(int)2e6 + 1], inv[(int)2e6 + 1];
ll getinv(ll a){ return qpow(a, mod-2); }
void init(int n){
  fac[0] = 1;
  for(int i = 1; i <= n; i++){
    fac[i] = fac[i-1] * i % mod;
  }
  inv[n] = getinv(fac[n]);
  for(int i = n - 1; i >= 0; i--){
    inv[i] = inv[i + 1] * (i + 1) % mod;
  }
}
ll C(int n, int m){
  if(m > n) return 0;
  return fac[n] * inv[m] % mod * inv[n-m] % mod;
}
```

3.15 Primes

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771

* 1001010013, 1000512343, 987654361, 999991231

* 999888733, 98789101, 987777733, 999991921, 1010101333
* 1010102101, 1000000000039, 100000000000037
* 2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[ N ] , p_tbl[ N ];
vector<int> primes;
void sieve() {
   mu[ 1 ] = p_tbl[ 1 ] = 1;
for( int i = 2 ; i < N ; i ++ ){
   if( !p_tbl[ i ] ){</pre>
         p_tbl[ i ] = i;
        primes.push_back( i );
mu[ i ] = -1;
      for( int p : primes ){
  int x = i * p;
  if( x >= M ) break;
         p_{tbl}[x] = p;
         mu[x] = -mu[i];
if(i%p == 0){
            mu[x] = 0;
            break;
vector<int> factor( int x ){
   vector<int> fac{ 1 };
   while(x > 1){
      int fn = SZ(fac), p = p_tbl[x], pos = 0;
      while( x \% p == 0){
         x \neq p;
         for( int i = 0 ; i < fn ; i ++ )
  fac.PB( fac[ pos ++ ] * p );</pre>
   } }
   return fac;
```

3.16 Phi

3.17 Int Sqrt

```
LL intSqrt(LL S) { //return origin val when S <= 0
    if (S <= 0) return S;
    LL x = S;
    for (LL nx;;x = nx){
        nx = (x+S/x)>>1LL;
        if(nx >= x) break;
    }
    return x;
}
```

3.18 Result

```
• Lucas' Theorem : For n,m\in\mathbb{Z}^* and prime P, C(m,n)\mod P=\Pi(C(m_i,n_i)) where m_i is the i-th digit of m in base P.
```

- Stirling approximation : $n! \approx \sqrt{2\pi n} \left(\frac{n}{2}\right)^n e^{\frac{1}{12n}}$
- Stirling Numbers(permutation |P|=n with k cycles): S(n,k)= coefficient of x^k in $\Pi_{i=0}^{n-1}(x+i)$

```
• Stirling Numbers(Partition n elements into k non-empty set): S(n,k)=\frac{1}{k!}\sum_{j=0}^k(-1)^{k-j}{k\choose j}j^n
```

```
    Pick's Theorem: A = i + b/2 - 1
在二維座標平面中畫上網格・對於任何簡單多邊形
A: 面積、i: 內部的格點數、b: 邊上的格點數
```

```
 \begin{split} \bullet & \text{ Catalan number } : \ C_n = \binom{2n}{n}/(n+1) \\ & C_n^{n+m} - C_{n+1}^{n+m} = (m+n)! \frac{n-m+1}{n+1} \quad for \quad n \geq m \\ & C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)!n!} \\ & C_0 = 1 \quad and \quad C_{n+1} = 2(\frac{2n+1}{n+2})C_n \\ & C_0 = 1 \quad and \quad C_{n+1} = \sum_{i=0}^n C_i C_{n-i} \quad for \quad n \geq 0 \end{split}
```

• Euler Characteristic: planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2 V,E,F,C: number of vertices, edges, faces(regions), and components

• Kirchhoff's theorem : $A_{ii}=deg(i), A_{ij}=(i,j)\in E$?-1:0, Deleting any one row, one column, and cal the det(A)

• Polya' theorem (c is number of color \cdot m is the number of cycle size): $(\sum_{i=1}^m c^{gcd(i,m)})/m$

• Burnside lemma: $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

・ 錯排公式: (n 個人中·每個人皆不再原來位置的組合數): dp[0] = 1; dp[1] = 0; dp[i] = (i-1)*(dp[i-1]+dp[i-2]);

• Bell 數 (有 n 個人,把他們拆組的方法總數): $B_0=1$ $B_n=\sum_{k=0}^n s(n,k)$ (second-stirling) $B_{n+1}=\sum_{k=0}^n \binom{n}{k} B_k$

• Wilson's theorem : $(p-1)! \equiv -1 (mod \ p)$

• Fermat's little theorem : $a^p \equiv a (mod\ p)$

• Euler's totient function: $A^{B^C} \mod p = pow(A, pow(B, C, p-1)) \mod p$

• 歐拉函數降幂公式: $A^B \mod C = A^B \mod \phi(c) + \phi(c) \mod C$

• 用歐拉函數求模反元素: 如果 a 和 n 互質,則 a 對 n 的模反元素 $a^{-1} \equiv a^{\phi(n)-1} (mod\ n)$

• 6 的倍數: $(a-1)^3 + (a+1)^3 + (-a)^3 + (-a)^3 = 6a$

• 點到直線距離公式: $d = \frac{|ax_0 + by_0 + c|}{\sqrt{a^2 + b^2}}$

4 Geometry

4.1 definition

```
#define all(a) a.begin(),a.end()
ostream& operator<<(ostream& os, const Pt& pt) {
    return os << "(" << pt.x << ", " << pt.y << ")";}
typedef long double ld;
const ld eps = 1e-8;
const ld pi = acos(-1);
int dcmp(ld x) {
    if(abs(x) < eps) return 0;
    else return x < 0 ? -1 : 1;
}
struct Pt {
    ld x, y;
    Pt(ld _x=0, ld _y=0):x(_x), y(_y) {}
    Pt operator+(const Pt &a) const {
        return Pt(x+a.x, y+a.y);
    }
    Pt operator-(const Pt &a) const {</pre>
```

```
return Pt(x-a.x, y-a.y); }
Pt operator*(const ld &a) const {
    return Pt(x*a, y*a);
  Pt operator/(const ld &a) const {
    return Pt(x/a, y/a);
  ld operator*(const Pt &a) const {
  return x*a.x + y*a.y; }
ld operator^(const Pt &a) const {
    return x*a.y - y*a.x;
  bool operator<(const Pt &a) const {</pre>
    return x < a.x | | (x == a.x && y < a.y); }
    //return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0 \&\&
         dcmp(y-a.y) < 0); }
  bool operator==(const Pt &a) const {
    return dcmp(x-a.x) == 0 && dcmp(y-a.y) == 0; }
ld norm2(const Pt &a) {
  return a*a; ]
ld norm(const Pt &a) {
  return sqrt(norm2(a)); }
Pt perp(const Pt &a) {
return Pt(-a.y, a.x); }
Pt rotate(const Pt &a, ld ang) {
  return Pt(a.x*cos(ang)-a.y*sin(ang), a.x*sin(ang)+a.y
       *cos(ang)); }
bool collinear(Pt a, Pt b, Pt c) { return ((b - a) ^ (c
       (a)) == 0; 
struct Circle {
  Pt o; ld r;
  Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
```

4.2 Intersection of 2 lines

```
// NAN(parallel), INF(overlapping)
Pt LLIntersect(Line a, Line b) {
  Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
  ld f1 = (p2-p1)^(q1-p1),f2 = (p2-p1)^(p1-q2),f;
  if(dcmp(f=f1+f2) == 0)
    return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
  return q1*(f2/f) + q2*(f1/f);
}
```

4.3 halfPlaneIntersection

// 0(nlogn)

```
// 傳入 vector<Line>
// (半平面為點 st 往 ed 的逆時針方向)
// 回傳值為形成的凸多邊形的頂點 vector
// assume that Lines intersect
vector<Pt> HPI(vector<Line> P) {
     sort(P.begin(), P.end(), [&](Line l, Line m) {
   if (argcmp(l.v, m.v)) return true;
          if (argcmp(m.v, l.v)) return false;
          return PtSide(l.s, m) > 0;
    });
    int n = P.size(), l = 0, r = -1;
for (int i = 0; i < n; i++) {
    if (i and !argcmp(P[i - 1].v, P[i].v)) continue</pre>
          while (l < r and PtSide(LLIntersect(P[r-1], P[r</pre>
               ]), P[i]) <= 0) r-
          while (l < r and PtSide(LLIntersect(P[l], P[l</pre>
               +1]), P[i]) <= 0) l++;
          P[++r] = P[i];
     while (l < r and PtSide(LLIntersect(P[r-1], P[r]),</pre>
          P[1]) <= 0) r-
     while (l < r and PtSide(LLIntersect(P[l], P[l+1]),</pre>
          P[r]) <= 0) l++;
     if (r - 1 <= 1 or !argcmp(P[l].v, P[r].v))
    return {}; // empty</pre>
     if (PtSide(LLIntersect(P[l], P[r]), P[l+1]) <= 0) {</pre>
          assert(0);
          return {}; // infinity
     vector<Line> lns = vector(P.begin() + 1, P.begin()
          + r + 1);
```

```
lns.push_back(lns[0]);
vector<Pt> hpi;
for(int i = 1; i < lns.size(); i++) hpi.push_back(
    LLIntersect(lns[i-1], lns[i]));
return hpi;
}</pre>
```

4.4 Convex Hull

```
double cross(Pt o, Pt a, Pt b){
  return (a-o) ^ (b-o);
vector<Pt> convex_hull(vector<Pt> pt){
  sort(pt.begin(),pt.end());
  int top=0;
  vector<Pt> stk(2*pt.size());
  for (int i=0; i<(int)pt.size(); i++){</pre>
    while (top >= 2 && cross(stk[top-2],stk[top-1],pt[i
]) <= 0) // 如果想要有點共線的點·把 <= 改成 <
      top--;
    stk[top++] = pt[i];
  for (int i=pt.size()-2, t=top+1; i>=0; i--){
    while (top >= t && cross(stk[top-2],stk[top-1],pt[i
        ]) <= 0)
      top--;
    stk[top++] = pt[i];
  stk.resize(top-1);
  return stk;
```

4.5 Convex Hull trick

```
struct Convex {
  int n;
  vector<Pt> A, V, L, U;
  Convex(const vector<Pt> &_A) : A(_A), n(_A.size()) {
      // n >= 3
    auto it = max_element(all(A));
    L.assign(A.begin(), it + 1);
    U.assign(it, A.end()), U.push_back(A[0]);
for (int i = 0; i < n; i++) {</pre>
      V.push\_back(A[(i + 1) % n] - A[i]);
  int PtSide(Pt p, Line L) {
    return dcmp((L.b - L.a)^{p - L.a);
  int inside(Pt p, const vector<Pt> &h, auto f) {
    auto it = lower_bound(all(h), p, f);
    if (it == h.end()) return 0;
    if (it == h.begin()) return p == *it;
    return 1 - dcmp((p - *prev(it))^(*it - *prev(it)))
  ^{\prime}// 1. whether a given point is inside the CH
  // ret 0: out, 1: on, 2: in
  int inside(Pt p) {
    return min(inside(p, L, less{}), inside(p, U,
        greater{}));
  static bool cmp(Pt a, Pt b) { return dcmp(a \land b) > 0;
  // 2. Find tangent points of a given vector
  // ret the idx of far/closer tangent point
  int tangent(Pt v, bool close = true) {
    assert(v != Pt{});
    auto l = V.begin(), r = V.begin() + L.size() - 1;
    if (v < Pt{}) l = r, r = V.end();
    if (close) return (lower_bound(l, r, v, cmp) - V.
        begin()) % n;
    return (upper_bound(l, r, v, cmp) - V.begin()) % n;
  // 3. Find 2 tang pts on CH of a given outside point
  // return index of tangent points
  // return {-1, -1} if inside CH
  array<int, 2> tangent2(Pt p) {
```

```
array<int, 2> t{-1, -1};
if (inside(p) == 2) return t;
     if (auto it = lower_bound(all(L), p); it != L.end()
          and p == *it) {
       int s = it - L.begin();
       return {(s + 1) % n, (s - 1 + n) % n};
    if (auto it = lower_bound(all(U), p, greater{}); it
          != U.end() and p == *it) {
       int s = it - U.begin() + L.size() - 1;
       return \{(s + 1) \% n, (s - 1 + n) \% n\};
    for (int i = 0; i != t[0]; i = tangent((A[t[0] = i]
    - p), 0));
for (int i = 0; i != t[1]; i = tangent((p - A[t[1]
         = i]), 1));
    return t;
  int find(int l, int r, Line L) {
  if (r < l) r += n;</pre>
    int s = PtSide(A[1 % n], L);
    return *ranges::partition_point(views::iota(l, r),
       [&](int m) {
         return PtSide(A[m % n], L) == s;
       }) - 1;
  };
// 4. Find intersection point of a given line
  // intersection is on edge (i, next(i))
  vector<int> intersect(Line L) {
    int l = tangent(L.a - L.b), r = tangent(L.b - L.a);
    if(PtSide(A[1], L) == 0) return {l};
     if(PtSide(A[r], L) == 0) return {r};
    if (PtSide(A[l], L) * PtSide(A[r], L) > 0) return
    return {find(l, r, L) % n, find(r, l, L) % n};
  }
|};
```

4.6 Intersection of 2 segments

4.7 Intersection of Polygon and Circle

```
ld PCIntersect(vector<Pt> v, Circle cir) {
  for(int i = 0 ; i < (int)v.size() ; ++i) v[i] = v[i]</pre>
       - cir.o;
  ld ans = 0, r = cir.r;
  int n = v.size();
for(int i = 0; i < n; ++i) {
  Pt pa = v[i], pb = v[(i+1)%n];</pre>
    if(norm(pa) < norm(pb)) swap(pa, pb);</pre>
    if(dcmp(norm(pb)) == 0) continue;
    ld s, h, theta;
    ld a = norm(pb), b = norm(pa), c = norm(pb-pa);
    ld cosB = (pb*(pb-pa))/a/c, B = acos(cosB);
    if(cosB > 1) B = 0;
    else if(cosB < -1) B = PI;</pre>
    ld cosC = (pa*pb)/a/b, C = acos(cosC);
    if(cosC > 1) C = 0;
    else if(cosC < -1) C = PI;</pre>
    if(a > r) {
```

4.8 Circle cover

```
#define N 1021
#define D long double
struct CircleCover{
  int C; Circle c[ N ]; //填入C(圓數量),c(圓陣列)
  bool g[ N ][ N ], overlap[ N ][ N ];
// Area[i] : area covered by at least i circles
  D Area[ N ];
  void init( int _C ){ C = _C; }
bool CCinter( Circle& a , Circle& b , Pt& p1 , Pt& p2
     Pt o1 = a.o, o2 = b.o;
     D r1 = a.r , r2 = b.r;
     if( norm( o1 - o2 ) > r1 + r2 ) return {};
if( norm( o1 - o2 ) < max(r1, r2) - min(r1, r2) )</pre>
           return {};
     D d2 = (o1 - o2) * (o1 - o2);
     D d = sqrt(d2);
     if( d > r1 + r2 ) return false;
     Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
     D A=sqrt((r_1+r_2+d)*(r_1-r_2+d)*(r_1+r_2-d)*(-r_1+r_2+d));
     Pt v=Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
p1 = u + v; p2 = u - v;
     return true;
  struct Teve {
     Pt p; D ang; int add;
     Teve() {}
     Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
     bool operator<(const Teve &a)const
  {return ang < a.ang;}
}eve[ N * 2 ];
  // strict: x = 0, otherwise x = -1
  bool disjuct( Circle& a, Circle &b, int x )
{return dcmp( norm( a.o - b.o ) - a.r - b.r ) > x;}
bool contain( Circle& a, Circle &b, int x )
   {return dcmp( a.r - b.r - norm( a.o - b.o ) ) > x;}
  bool contain(int i, int j){
     /* c[j] is non-strictly in c[i]. */
     return (dcmp(c[i].r - c[j].r) > 0 ||
(dcmp(c[i].r - c[j].r) == 0 && i < j) ) &&
                     contain(c[i], c[j], -1);
  void solve(){
     for( int i = 0 ; i \leftarrow C + 1 ; i + + )
        Area[i] = 0;
     for( int i = 0 ; i < C ; i ++ )
for( int j = 0 ; j < C ; j ++ )</pre>
          overlap[i][j] = contain(i, j);
     for( int i = 0; i < C; i ++ )
for( int j = 0; j < C; j ++ )
g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                          disjuct(c[i], c[j], -1));
     for( int i = 0 ; i < C ; i ++ ){
        int E = 0, cnt = 1;
        for( int j = 0; j < C; j ++)
           if( j != i && overlap[j][i] )
             cnt ++;
        for( int j = 0 ; j < C ; j ++ )
  if( i != j && g[i][j] ){</pre>
             Pt aa, bb;
             CCinter(c[i], c[j], aa, bb);
             D A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
D B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);
```

4.9 Tangent line of two circles

```
vector<Line> go( const Cir& c1 , const Cir& c2 , int
      sign1 ){
   // sign1 = 1 for outer tang, -1 for inter tang
   vector<Line> ret;
   double d_{sq} = norm2(c1.0 - c2.0);
   if( d_sq < eps ) return ret;</pre>
   double d = sqrt( d_sq );
  Pt v = ( c2.0 - c1.0 ) / d;
double c = ( c1.R - sign1 * c2.R ) / d;
if( c * c > 1 ) return ret;
   double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
                   v.Y * c + sign2 * h * v.X };
      Pt p1 = c1.0 + n * c1.R;
      Pt p2 = c2.0 + n * (c2.R * sign1);
     if( fabs( p1.X - p2.X ) < eps and
   fabs( p1.Y - p2.Y ) < eps )
  p2 = p1 + perp( c2.0 - c1.0 );</pre>
      ret.push_back( { p1 , p2 } );
   return ret;
}
```

4.10 Minimum distance of two convex

4.11 Poly Union

```
struct PY{
  int n; Pt pt[5]; double area;
  Pt& operator[](const int x){ return pt[x]; }
  void init(){ //n,pt[0~n-1] must be filled
    area=pt[n-1]^pt[0];
    for(int i=0;i<n-1;i++) area+=pt[i]^pt[i+1];
    if((area/=2)<0)reverse(pt,pt+n),area=-area;
} };
PY py[500]; pair<double,int> c[5000];
inline double segP(Pt &p,Pt &p1,Pt &p2){
  if(dcmp(p1.x-p2.x)==0) return (p.y-p1.y)/(p2.y-p1.y);
```

```
return (p.x-p1.x)/(p2.x-p1.x);
double polyUnion(int n){ //py[0~n-1] must be filled
   int i,j,ii,jj,ta,tb,r,d; double z,w,s,sum=0,tc,td;
   for(i=0;i<n;i++) py[i][py[i].n]=py[i][0];
for(i=0;i<n;i++){</pre>
     for(ii=0;ii<py[i].n;ii++){</pre>
       r=0:
       c[r++]=make\_pair(0.0,0); c[r++]=make\_pair(1.0,0);
        for(j=0;j<n;j++){</pre>
          if(i==j) continue;
          for(jj=0;jj<py[j].n;jj++){</pre>
            ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))
            tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj
                 +1]));
            if(ta==0 \&\& tb==0){
               if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[
                    i][ii])>0&&j<i){
                 c[r++]=make_pair(segP(py[j][jj],py[i][ii
                      ],py[i][ii+1]),1);
                 c[r++]=make_pair(segP(py[j][jj+1],py[i][
                      ii],py[i][ii+1]),-1);
            }else if(ta>=0 && tb<0){
              tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
c[r++]=make_pair(tc/(tc-td),1);
            }else if(ta<0 && tb>=0){
              tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
               c[r++]=make_pair(tc/(tc-td),-1);
       } } }
       sort(c,c+r)
       z=min(max(c[0].first,0.0),1.0); d=c[0].second; s
            =0;
        for(j=1;j<r;j++){</pre>
          w=min(max(c[j].first,0.0),1.0);
          if(!d) s+=w-z;
          d+=c[j].second; z=w;
       sum+=(py[i][ii]^py[i][ii+1])*s;
   } }
   return sum/2;
}
```

4.12 Minkowski sum

```
// P, Q, R(return) are counterclockwise order convex
    polygon
#define all(a) a.begin(),a.end()
vector<Pt> Minkowski(vector<Pt> P, vector<Pt> Q) {
    auto cmp = [&](Pt a, Pt b) {
        return Pt{a.y, a.x} < Pt{b.y, b.x};</pre>
    auto reorder = [&](auto &R) {
        rotate(R.begin(), min_element(all(R), cmp), R.
            end());
        R.push_back(R[0]), R.push_back(R[1]);
    const int n = P.size(), m = Q.size();
    reorder(P), reorder(Q);
    vector<Pt> R;
    for (int i = 0, j = 0, s; i < n or j < m; ) {
        R.push_back(P[i] + Q[j]);
        s = dcmp((P[i + 1] - P[i]) \wedge (Q[j + 1] - Q[j]))
        if (s >= 0) i++;
        if (s <= 0) j++;
  rotate(R.begin(), min_element(all(R)), R.end());
    return R;
```

4.13 Area of Rectangles

```
struct AreaofRectangles{
#define cl(x) (x<<1)</pre>
```

```
pair<ll,ll> tree[MXN<<3]; // count, area</pre>
    vector<ll> ind;
    tuple<ll, ll, ll, ll> scan[MXN<<1];</pre>
    void pull(int i, int l, int r){
         if(tree[i].first) tree[i].second = ind[r+1] -
             ind[l];
         else if(l != r){
             int mid = (l+r)>>1;
             tree[i].second = tree[cl(i)].second + tree[
                  cr(i)].second;
         else
                 tree[i].second = 0;
    void upd(int i, int l, int r, int ql, int qr, int v
         if(ql \ll l \& r \ll qr){
             tree[i].first += v;
             pull(i, l, r); return;
         int mid = (l+r) \gg 1;
         if(ql <= mid) upd(cl(i), l, mid, ql, qr, v);</pre>
         if(qr > mid) upd(cr(i), mid+1, r, ql, qr, v);
         pull(i, l, r);
    void init(int _n){
    n = _n; id = sid = 0;
         ind.clear(); ind.resize(n<<1);</pre>
         fill(tree, tree+(n<<2), make_pair(0, 0));</pre>
    void addRectangle(int lx, int ly, int rx, int ry){
         ind[id++] = lx; ind[id++] = rx;
scan[sid++] = make_tuple(ly, 1, lx, rx);
         scan[sid++] = make\_tuple(ry, -1, lx, rx);
    ll solve(){
         sort(ind.begin(), ind.end());
         ind.resize(unique(ind.begin(), ind.end()) - ind
              .begin());
         sort(scan, scan + sid);
         11 area = 0, pre = get<0>(scan[0]);
         for(int i = 0; i < sid; i++){
             auto [x, v, l, r] = scan[i];
             area += tree[1].second * (x-pre);
upd(1, 0, ind.size()-1, lower_bound(ind.
    begin(), ind.end(), l)-ind.begin(),
                  lower_bound(ind.begin(),ind.end(),r)-
                  ind.begin()-1, v);
             pre = x;
         return area;
    }rect;
```

4.14 Min dist on Cuboid

4.15 Distance of Line and Point

```
ld Distance_of_Line_and_Point(Line 1, Pt p) {
    ld cross_product = abs((p - l.s) ^ l.v);
    ld line_length = sqrtl(l.v * l.v);
    return cross_product / line_length;
}
```

4.16 Angle of two vector

```
// radian of OA and OB (directed angle)
ld Angle_of_two_vector(Pt A, Pt B, Pt 0) {
    ld a = (A - 0) * (B - 0);
    ld b = (A - 0) ^ (B - 0);
    ld theta = atan2(b, a);
    return theta;
}
```

4.17 極角排序

|}

```
//極角排序
//atan2(y, x) version
// p is reference point
// 180 度開始, 逆時針排序, 剛好在 180 度會排最後
bool cmp(Pt &lhs, Pt rhs) {
    return atan2((lhs - p).y, (lhs - p).x) < atan2((rhs - p).y, (rhs - p).x);
}

//cross product version
// p is reference point
// 270 度開始, 逆時針排序, 剛好在 270 度會排最後
bool cmp(const Pt& lhs, const Pt& rhs) {
    if ((lhs < p) ^ (rhs < p)) return (lhs < p) < (rhs < p);
    return ((lhs - p) ^ (rhs - p)) > 0;
}
```

4.18 Heart of Triangle

```
Pt inCenter( Pt &A, Pt &B, Pt &C) { // 內心 double a = norm(B-C), b = norm(C-A), c = norm(A-B); return (A * a + B * b + C * c) / (a + b + c); }
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心 Pt bb = b - a, cc = c - a; double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc); return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d; }
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心 Pt ba = b - a, ca = c - a, bc = b - c; double Y = ba.Y * ca.Y * bc.Y, A = ca.X * ba.Y - ba.X * ca.Y, x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A, y0= -ba.X * (x0 - c.X) / ba.Y + ca.Y; return Pt(x0, y0); }
```

5 Graph

5.1 Lowest Common Ancestor O(lqn)

```
struct LCA {
  int n, ti, lgN;
  int anc[MXN + 5][__lg(MXN) + 1] = {0};
  int MaxLength[MXN][__lg(MXN) + 1] = {0};
  int time_in[MXN] = {0};
  int time_out[MXN] = {0};
```

```
LCA(int _n, int f):n(_n), ti(0), lgN(__lg(n)) {
   dfs(f, f, 0);
   build();
  void dfs(int now, int f, int len_to_father) { // dfs
        for anc, time, Lenth
    ti++;
    anc[now][0] = f;
    time_in[now] = ti;
    MaxLength[now][0] = len_to_father;
    for (auto i : graph[now]) {
         if (i.first == f) continue;
         dfs(i.first, now, i.second);
    time_out[now] = ti;
  void build() {      // build anc[][], MaxLength[][]
for (int i = 1; i <= lgN; ++i) {</pre>
       for (int u = 1; u <= n; ++u) {
         anc[u][i] = anc[anc[u][i - 1]][i - 1];
         MaxLength[u][i] = max(MaxLength[u][i - 1],
        MaxLength[anc[u][i - 1]][i - 1]);

// dis[u][i] += dis[anc[u][i - 1]][i - 1]

// + dis[u][i - 1];
    }
  bool isAncestor(int x, int y) {
    return time_in[x] <= time_in[y] && time_out[x] >=
         time_out[y];
  int getLCA(int u, int v) {
    if (isAncestor(u, v)) return u;
    if (isAncestor(v, u)) return v;
    for (int i = lgN; i >= 0; --i) {
      if (!isAncestor(anc[u][i], v)) {
        u = anc[u][i];
      }
    }
    return anc[u][0];
  int getMAX(int u, int v) { //獲得路徑上最大邊權
    int lca = getLCA(u, v);
    int maxx = -1;
    for (int i = lgN; i >= 0; --i) {
      // u to lca
      if (!isAncestor(anc[u][i], lca))
        maxx = max(maxx, MaxLength[u][i]);
         u = anc[u][i];
      // v to lca
      if (!isAncestor(anc[v][i], lca))
        maxx = max(maxx, MaxLength[v][i]);
         v = anc[v][i];
    if (u != lca) maxx = max(maxx, MaxLength[u][0]);
    if (v != lca) maxx = max(maxx, MaxLength[v][0]);
    return maxx;
};
```

5.2 Hamiltonian path $O(n^22^n)$

5.3 MaximumClique 最大團

```
#define N 111
struct MaxClique{ // 0-base
   typedef bitset<N> Int;
   Int linkto[N] , v[N];
   int n:
   void init(int _n){
     n = _n;
     for(int i = 0; i < n; i ++){
       linkto[i].reset(); v[i].reset();
  void addEdge(int a , int b)
{ v[a][b] = v[b][a] = 1; }
int popcount(const Int& val)
   { return val.count(); }
   int lowbit(const Int& val)
   { return val._Find_first(); }
   int ans , stk[N];
   int id[N] , di[N] , deg[N];
   Int cans;
   void maxclique(int elem_num, Int candi){
     if(elem_num > ans){
       ans = elem_num; cans.reset();
for(int i = 0; i < elem_num; i ++)</pre>
          cans[id[stk[i]]] = 1;
     int potential = elem_num + popcount(candi);
     if(potential <= ans) return;</pre>
     int pivot = lowbit(candi);
Int smaller_candi = candi & (~linkto[pivot]);
     while(smaller_candi.count() && potential > ans){
       int next = lowbit(smaller_candi);
candi[next] = !candi[next];
        smaller_candi[next] = !smaller_candi[next];
        potential --:
        if(next == pivot || (smaller_candi & linkto[next
             ]).count()){
          stk[elem_num] = next;
          maxclique(elem_num + 1, candi & linkto[next]);
   } } }
   int solve(){
     for(int i = 0; i < n; i ++){
       id[i] = i; deg[i] = v[i].count();
     sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
     for(int i = 0; i < n; i ++) di[id[i]] = i;
for(int i = 0; i < n; i ++)
  for(int j = 0; j < n; j ++)</pre>
          if(v[i][j]) linkto[di[i]][di[j]] = 1;
     Int cand; cand.reset();
     for(int i = 0 ; i < n ; i ++) cand[i] = 1;
     ans = 1;
     cans.reset(); cans[0] = 1;
     maxclique(0, cand);
     return ans;
} }solver;
```

5.4 MaximalClique 極大團

```
#define N 80
struct MaxClique{ // 0-base
    typedef bitset<N> Int;
    Int lnk[N] , v[N];
    int n;
    void init(int _n){
        n = _n;
        for(int i = 0 ; i < n ; i ++){
             lnk[i].reset(); v[i].reset();
        }
    void addEdge(int a , int b)
    { v[a][b] = v[b][a] = 1; }
    int ans , stk[N], id[N] , di[N] , deg[N];
    Int cans;
    void dfs(int elem_num, Int candi, Int ex){
        if(candi.none()&&ex.none()){
            cans.reset();
            for(int i = 0 ; i < elem_num ; i ++)</pre>
```

```
cans[id[stk[i]]] = 1;
ans = elem_num; // cans is a maximal clique
    int pivot = (candilex)._Find_first();
    Int smaller_candi = candi & (~lnk[pivot]);
    while(smaller_candi.count()){
       int nxt = smaller_candi._Find_first();
       candi[nxt] = smaller_candi[nxt] = 0;
       ex[nxt] = 1;
       stk[elem_num] = nxt;
       dfs(elem_num+1,candi&lnk[nxt],ex&lnk[nxt]);
  int solve(){
    for(int i = 0 ; i < n ; i ++){
       id[i] = i; deg[i] = v[i].count();
    sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
    for(int i = 0; i < n; i ++) di[id[i]] = i;
for(int i = 0; i < n; i ++)
       for(int j = 0; j < n; j ++)
         if(v[i][j]) lnk[di[i]][di[j]] = 1;
    ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n,'1')), 0);
    return ans;
} }solver;
```

5.5 BCC based on vertex 點雙聯通分量

```
#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
  int n,nScc,step,dfn[MXN],low[MXN];
  vector<int> E[MXN],sccv[MXN];
  int top,stk[MXN];
  void init(int _n) { // 初始化n點
    n = _n; nScc = step = 0;
for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v) // 無向邊
  { E[u].PB(v); E[v].PB(u); } void DFS(int u, int f) {
    dfn[u] = low[u] = step++;
    stk[top++] = u;
    for (auto v:E[u]) {
      if (v == f) continue;
if (dfn[v] == -1) {
         DFS(v,u);
         low[u] = min(low[u], low[v]);
         if (low[v] >= dfn[u]) {
           int z;
           sccv[nScc].clear();
           do {
             z = stk[--top];
             sccv[nScc].PB(z);
           } while (z != v);
           sccv[nScc++].PB(u);
      }else
         low[u] = min(low[u],dfn[v]);
  vector<vector<int>> solve() { // 回傳(size=2 橋, size
      >2 點雙連通分量)
    vector<vector<int>> res;
    for (int i=0; i<n; i++)</pre>
    dfn[i] = low[i] = -1;
for (int i=0; i<n; i++)</pre>
      if (dfn[i] == -1) {
         top = 0;
        DFS(i,i);
    REP(i,nScc) res.PB(sccv[i]);
    return res;
}graph;
```

5.6 Strongly Connected Component 強連通分量

```
#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x)) //fill zero
struct Scc{
  int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
  void init(int _n){
    n = _n;
for (int i=0; i<MXN; i++)
       E[i].clear(), rE[i].clear();
  void addEdge(int u, int v){
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u){
    vst[u]=1;
     for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u){
  vst[u] = 1; bln[u] = nScc;
     for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve(){
    nScc = 0;
     vec.clear();
     FZ(vst);
     for (int i=0; i<n; i++)
       if (!vst[i]) DFS(i);
     reverse(vec.begin(),vec.end());
     FZ(vst);
     for (auto v : vec)
       if (!vst[v]){
         rDFS(v); nScc++;
  }
};
```

5.7 ManhattanMST

```
//return {{u,v},w}: u <-> v (w), 需要再手動去重
//need Point definition
vector<pair<int,int>, int>> ManhattanMST(vector<Pt</pre>
    > P) {
  vector<int> id(P.size());
  iota(id.begin(),id.end(), 0);
  vector<pair<pair<int,int>, int>> edg;
for (int k = 0; k < 4; k++) {</pre>
    sort(id.begin(),id.end(), [&](int i, int j) {
       return (P[i] - P[j]).x < (P[j] - P[i]).y;</pre>
    map<int, int> sweep;
for (int i : id) {
       auto it = sweep.lower_bound(-P[i].y);
       while (it != sweep.end()) {
         int j = it->second;
Pt d = P[i] - P[j];
         if (d.y > d.x) break;
         edg.push_back(\{\{i, j\}, d.x + d.y\});
         it = sweep.erase(it);
      }
       sweep[-P[i].y] = i;
    for (Pt &p : P) {
       if (k \% 2) p.x = -p.x;
       else swap(p.x, p.y);
  }
  return edg;
```

5.8 Min Mean Cycle

```
/* minimum mean cycle O(VE) */
struct MMC{
```

```
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
  struct Edge { int v,u; double c; };
  int n, m, prv[V][V], prve[V][V], vst[V];
  Edge e[E];
  vector<int> edgeID, cycle, rho;
  double d[V][V];
  void init( int _n )
  { n = _n; m = 0; }
// WARNING: TYPE matters
  void addEdge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , úi , ci }; } void bellman_ford() {
    d[i+1][u] = d[i][v]+e[j].c;
          prv[i+1][u] = v;
          prve[i+1][u] = j;
  double solve(){
    // returns inf if no cycle, mmc otherwise
    double mmc=inf;
    int st = -1;
    bellman_ford();
    for(int i=0; i<n; i++) {</pre>
      double avg=-inf;
      for(int k=0; k<n; k++) {
  if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])</pre>
             ])/(n-k));
        else avg=max(avg,inf);
      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    fill(vst,0); edgeID.clear(); cycle.clear(); rho.
        clear();
    for (int i=n; !vst[st]; st=prv[i--][st]) {
      vst[st]++;
      edgeID.PB(prve[i][st]);
      rho.PB(st);
    while (vst[st] != 2) {
      if(rho.empty()) return inf;
      int v = rho.back(); rho.pop_back();
      cycle.PB(v);
      vst[v]++;
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
    return mmc;
} }mmc;
```

5.9 Directed Graph Min Cost Cycle

```
// works in O(N M)
#define INF 10000000000000000LL
#define N 5010
#define M 200010
struct edge{
  int to; LL w;
  edge(int a=0, LL b=0): to(a), w(b){}
struct node{
  LL d; int u, next;
node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
}b[M];
struct DirectedGraphMinCycle{
  vector<edge> g[N], grev[N];
  LL dp[N][N], p[N], d[N], mu;
  bool inq[N];
  int n, bn, bsz, hd[N];
  void b_insert(LL d, int u){
    int i = d/mu;
    if(i >= bn) return;
    b[++bsz] = node(d, u, hd[i]);
```

```
hd[i] = bsz;
   void init( int _n ){
     n = _n;
for( int i = 1 ; i <= n ; i ++ )</pre>
         g[ i ].clear();
   void addEdge( int ai , int bi , LL ci )
   { g[ai].push_back(edge(bi,ci)); }
   LL solve(){
      fill(dp[0], dp[0]+n+1, 0);
      for(int i=1; i<=n; i++){</pre>
         fill(dp[i]+1, dp[i]+n+1, INF);
for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){
    for(int k=0; k<(int)g[j].size(); k++)</pre>
              dp[i][g[j][k].to] =min(dp[i][g[j][k].to],
                                              dp[i-1][j]+g[j][k].w);
      mu=INF; LL bunbo=1;
      for(int i=1; i<=n; i++) if(dp[n][i] < INF){
  LL a=-INF, b=1;</pre>
         for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
   if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre>
              a = dp[n][i]-dp[j][i];
              b = n-j;
         if(mu*b > bunbo*a)
           mu = a, bunbo = b;
      if(mu < 0) return -1; // negative cycle
if(mu == INF) return INF; // no cycle</pre>
      if(mu == 0) return 0;
      for(int i=1; i<=n; i++)
  for(int j=0; j<(int)g[i].size(); j++)
  g[i][j].w *= bunbo;</pre>
      memset(p, 0, sizeof(p));
      queue<int> q;
      for(int i=1; i<=n; i++){</pre>
         q.push(i);
         inq[i] = true;
      while(!q.empty()){
         int i=q.front(); q.pop(); inq[i]=false;
for(int j=0; j<(int)g[i].size(); j++){</pre>
            if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
              p[g[i][j].to] = p[i]+g[i][j].w-mu;
if(!inq[g[i][j].to]){
                 q.push(g[i][j].to);
                 inq[g[i][j].to] = true;
      for(int i=1; i<=n; i++) grev[i].clear();
for(int i=1; i<=n; i++)
    for(int j=0; j<(int)g[i].size(); j++){</pre>
           g[i][j].w += p[i]-p[g[i][j].to];
           grev[g[i][j].to].push_back(edge(i, g[i][j].w));
      LL mldc = n*mu;
      for(int i=1; i<=n; i++){</pre>
         bn=mldc/mu, bsz=0;
        memset(hd, 0, sizeof(hd));
fill(d+i+1, d+n+1, INF);
         b_insert(d[i]=0, i);
         for(int j=\bar{0}; j \leftarrow bn-1; j++) for(int k=hd[j]; k; k=
              b[k].next){
            int u = b[k].u;
           LL du = b[k].d;
            if(du > d[u]) continue;
            for(int l=0; l<(int)g[u].size(); l++) if(g[u][l
    ].to > i){
              if(d[g[u][l].to] > du + g[u][l].w){
  d[g[u][l].to] = du + g[u][l].w;
                 b_insert(d[g[u][l].to], g[u][l].to);
         } } }
         for(int j=0; j<(int)grev[i].size(); j++) if(grev[
    i][j].to > i)
            mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
      return mldc / bunbo;
} }graph;
```

5.10 DominatorTree

```
struct DominatorTree{ // O(N)
#define REP(i,s,e) for(int i=(s);i<=(e);i++)
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
  int n , m , s;
  vector< int > g[ MAXN ] , pred[ MAXN ];
vector< int > cov[ MAXN ];
  int dfn[ MAXN ] , nfd[ MAXN ] , ts;
  int par[ MAXN ]; //idom[u] s到u的最後一個必經點 int sdom[ MAXN ], idom[ MAXN ]; int mom[ MAXN ], mn[ MAXN ]; inline bool cmp( int u , int v ) { return dfn[ u ] < dfn[ v ]; } int eval( int u ){
     if( mom[ u ] == u ) return u;
     int res = eval( mom[ u ] );
if(cmp( sdom[ mn[ mom[ u ] ] ] , sdom[ mn[ u ] ] ))
        mn[u] = mn[mom[u]];
     return mom[ u ] = res;
   void init( int _n , int _m , int _s ){
     ts = 0; n = _n; m = _m; s = _s;

REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
  void addEdge( int u , int v ){
  g[ u ].push_back( v );
  pred[ v ].push_back( u );
   void dfs( int u ){
     ts++;
     dfn[u] = ts;
     nfd[ ts ] = u;
for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
        par[ v ] = u;
        dfs(v);
   void build(){
     mom[ i ] = mn[ i ] = sdom[ i ] = i;
     dfs( s );
     REPD( i , n , 2 ){
  int u = nfd[ i ];
  if( u == 0 ) continue ;
        for( int v : pred[ u ] ) if( dfn[ v ] ){
          eval( v );
if( cmp( sdom[ mn[ v ] ] , sdom[ u ] ) )
   sdom[ u ] = sdom[ mn[ v ] ];
        cov[ sdom[ u ] ].push_back( u );
        mom[ u ] = par[ u ];
        for( int w : cov[ par[ u ] ] ){
           eval( w );
           if( cmp( sdom[ mn[ w ] ] , par[ u ] ) )
          idom[w] = mn[w];
else idom[w] = par[u];
        cov[ par[ u ] ].clear();
     REP( i , 2 , n ){
  int u = nfd[ i ];
        if( u == 0 ) continue ;
if( idom[ u ] != sdom[ u ] )
           idom[ u ] = idom[ idom[ u ] ];
```

5.11 K-th Shortest Path

```
// time: O(|E| \lg |E| + |V| \lg |V| + K)
// memory: O(|E| \lg |E| + |V|)
struct KSP{ // 1-base
    struct nd{
      int u, v; ll d;
      nd(int ui = 0, int vi = 0, ll di = INF)
      { u = ui; v = vi; d = di; }
};
```

```
struct heap{
  nd* edge; int dep; heap* chd[4];
static int cmp(heap* a,heap* b)
{ return a->edge->d > b->edge->d; }
struct node{
   int v; ll d; heap* H; nd* E;
   node(){}
  node(ll _d, int _v, nd* _E)
{ d =_d; v = _v; E = _E; }
node(heap* _H, ll _d)
{ H = _H; d = _d; }
   friend bool operator<(node a, node b)
   { return a.d > b.d; }
int n, k, s, t;
ll dst[ N ];
nd *nxt[ N ];
vector<nd*> g[ N ], rg[ N ];
heap *nullNd, *head[ N ];
void init( int _n , int _k , int _s , int _t ){
  n = _n; k = _k; s = _s; t = _t;
for( int i = 1 ; i <= n ; i ++ ){
    g[ i ].clear(); rg[ i ].clear();
    nxt[ i ] = NULL; head[ i ] = NULL;
    dst[ i ] = -1;
}</pre>
void addEdge( int ui , int vi , ll di ){
  nd* e = new nd(ui, vi, di);
g[ ui ].push_back( e );
rg[ vi ].push_back( e );
queue<int> dfsQ:
void dijkstra(){
   while(dfsQ.size()) dfsQ.pop();
   priority_queue<node> Q;
   Q.push(node(0, t, NULL));
   while (!Q.empty()){
     node p = Q.top(); Q.pop();
if(dst[p.v] != -1) continue;
     dst[ p.v ] = p.d;
     nxt[ p.v ] = p.E;
      dfsQ.push( p.v )
     for(auto e: rg[ p.v ])
        Q.push(node(p.d + e->d, e->u, e));
heap* merge(heap* curNd, heap* newNd){
   if(curNd == nullNd) return newNd;
   heap* root = new heap;
   memcpy(root, curNd, sizeof(heap));
   if(newNd->edge->d < curNd->edge->d){
     root->edge = newNd->edge;
root->chd[2] = newNd->chd[2]
     root->chd[3] = newNd->chd[3];
     newNd->edge = curNd->edge;
newNd->chd[2] = curNd->chd[2];
     newNd - > chd[3] = curNd - > chd[3];
   if(root->chd[0]->dep < root->chd[1]->dep)
     root->chd[0] = merge(root->chd[0],newNd);
     root->chd[1] = merge(root->chd[1],newNd);
   root->dep = max(root->chd[0]->dep, root->chd[1]->
        dep) + 1;
  return root;
vector<heap*> V;
void build(){
   nullNd = new heap;
   nullNd->dep = 0;
   nullNd->edge = new nd;
fill(nullNd->chd, nullNd->chd+4, nullNd);
   while(not dfsQ.empty()){
     int u = dfsQ.front(); dfsQ.pop();
if(!nxt[ u ]) head[ u ] = nullNd;
else head[ u ] = head[nxt[ u ]->v];
     V.clear();
      for( auto&& e : g[ u ] ){
        int v = e->v;
        if( dst[ v ] == -1 ) continue;
e->d += dst[ v ] - dst[ u ];
if( nxt[ u ] != e ){
```

```
heap* p = new heap;
            fill(p->chd, p->chd+4, nullNd);
            p->dep = 1;
            p->edge = e;
            V.push_back(p);
       } }
       if(V.empty()) continue;
make_heap(V.begin(), V.end(), cmp);
#define L(X) ((X<<1)+1)
#define R(X) ((X<<1)+2)
       for( size_t i = 0 ; i < V.size() ; i ++ ){
  if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
          else V[i]->chd[2]=nullNd;
         if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
         else V[i]->chd[3]=nullNd;
       head[u] = merge(head[u], V.front());
  } }
  vector<ll> ans;
  void first_K(){
     ans.clear();
     priority_queue<node> Q;
     if( dst[ s ] == -1 ) return;
ans.push_back( dst[ s ] );
     if( head[s] != nullNd )
       Q.push(node(head[s], dst[s]+head[s]->edge->d));
     for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();</pre>
       ans.push_back( p.d );
       if(head[ p.H->edge->v ] != nullNd){
         q.H = head[p.H->edge->v];
         q.d = p.d + q.H->edge->d;
         Q.push(q);
       for( int i = 0 ; i < 4 ; i ++ )
  if( p.H->chd[ i ] != nullNd ){
    q.H = p.H->chd[ i ];
            q.d = p.d - p.H->edge->d + p.H->chd[i]->
                 edge->d;
            Q.push( q );
  } }
         }
  void solve(){ // ans[i] stores the i-th shortest path
     dijkstra();
     build():
     first_K(); // ans.size() might less than k
} }solver;
```

5.12 Floryd Warshall

5.13 Minimum Steiner Tree

```
} }
   void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
   void shortest_path(){ // using spfa may faster
      for( int k = 0 ; k < n ; k ++ )
  for( int i = 0 ; i < n ; i ++ )</pre>
            for( int j = 0; j < n; j ++ )

dst[ i ][ j ] = min( dst[ i ][ j ],

dst[ i ][ k ] + dst[ k ][ j ]);
   }// call shorest_path before solve
   int solve( const vector<int>& ter ){
      int t = (int)ter.size();
for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
      for( int j = 0 ; j < n ; j ++ )

dp[ i ][ j ] = INF;

for( int i = 0 ; i < n ; i ++ )
         dp[0][i] = 0;
      for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
  if( msk == ( msk & (-msk) ) ){</pre>
            int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
             continue;
          for( int i = 0 ; i < n ; i ++ )</pre>
             for( int submsk = ( msk - 1 ) & msk ; submsk ;
    submsk = ( submsk - 1 ) & msk )
                   for( int i = 0 ; i < n ; i ++ ){</pre>
            tdst[ i ] = INF;
             for( int j = 0 ; j < n ; j ++ )
  tdst[ i ] = min( tdst[ i ],</pre>
                                  dp[ msk ][ j ] + dst[ j ][ i ] );
         for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = tdst[ i ];</pre>
      int ans = INF;
      for( int i = 0 ; i < n ; i ++ )
ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
      return ans;
} }solver;
```

5.14 虚樹

```
vector<int> virTree(vector<int> ver, LCA &lca) {
    auto cmp = [&](int u, int v){return time_in[u] <
        time_in[v];};
    sort(ver.begin(),ver.end(),cmp); //用dfn排序
    vector<int>res(ver.begin(),ver.end());
    for(int i = 1; i < ver.size(); i++){
        res.push_back(lca.getLCA(ver[i-1],ver[i]));//把
        LCA丟進虚樹內
    }
    sort(res.begin(),res.end(),cmp); //再用dfn排序
    res.erase(unique(res.begin(),res.end()), res.end())
    ; //去掉重複的點
    return res;
}</pre>
```

5.15 Tree Hash

```
map<vector<int>, int> id;
int dfs(int x, int f){
  vector<int> sub;
  for (int v : edge[x]){
    if (v != f)
       sub.push_back(dfs(v, x));
  }
  sort(sub.begin(), sub.end());
  if (!id.count(sub))
    id[sub] = id.size();
  return id[sub];
}
```

5.16 HeavyLightDecomposition

```
// 詢問,修改複雜度 0(log^2 n)
// 1-base
int sz[MXN], dep[MXN], son[MXN], fa[MXN];
// 第一次 dfs
// 找重兒子 需要紀錄當前節點的子樹大小(sz)、深度(dep)、
    重兒子(son)、父節點(fa)
// 沒有子節點 son[x] = 0
void dfs_sz(int x, int f, int d) { //當前節點 x · 父節
    點 f・深度 d
    sz[x] = 1; dep[x] = d; fa[x] = f;
    for(int i : edge[x]) {
       if(i == f)
                     continue:
       dfs_sz(i, x, d+1);
sz[x] += sz[i];
       if(sz[son[x]] < sz[i])</pre>
                                 son[x] = i;
   }
}
// 第二次 dfs
int top[MXN]; // 每個節點所在的鏈的頂端節點
int dfn[MXN]; // 節點編號,編號為在線段樹上的位置
int rnk[MXN]; // 編號為哪個節點
int bottom[MXN]; // 維護每個節點的子樹中最大 dfn 編號
int cnt = 0;
int dfs_hld(int x, int f){
  top[x] = (son[fa[x]] == x ? top[fa[x]] : x);
    rnk[cnt] = x
    bottom[x] = dfn[x] = cnt++;
        on[x]) bottom[x] = max(bottom[x], dfs_hld(son[x], x)); // 更新子樹最大編號
    if(son[x])
    for(int i : edge[x]){
        if(i == f | i == son[x])
                                    continue;
       bottom[x] = max(bottom[x], dfs_hld(i, x)); //
            更新子樹最大編號
    return bottom[x];
}
// 求出 lca
// 不斷跳鏈·直到 u,v 跳到同一條鏈上為止
// 每次跳鏈選所在的鏈頂端深度較深的一端往上跳
int getLca(int u, int v) {
   while(top[u] != top[v]){
     if(dep[top[u]] > dep[top[v]])
          u = fa[top[u]];
          v = fa[top[v]];
    return dep[u] > dep[v] ? v : u;
}
// 路徑權重總和
int query(int u, int v) {
    int ret = 0;
    while(top[u] != top[v]){
        if (dep[top[u]] > dep[top[v]]){
            ret += segtree.query(dfn[top[u]], dfn[u]);
            u = fa[top[u]];
       }
       else{
            ret += segtree.query(dfn[top[v]], dfn[v]);
            v = fa[top[v]];
       }
    // 最後到同一條鏈上
    ret += segtree.query(min(dfn[u], dfn[v]), max(dfn[u
        ], dfn[v]));
    return ret;
}
```

- 龜兔賽跑演算法: 開始賽跑,兔子一次走兩格、烏龜一次走一格直到他們相遇停止 此時讓兔子返回起始點,兩者以相同走一格的速度繼續前進,他們就會在環入口
- 2-SAT 條件: 滿足 (x_1ory_1) $and(x_2ory_2)$ and ... 對於一個限制 (xory) , 則加兩條邊

6 String

6.1 PalTree O(n)

```
// state[i]代表第i個字元為結尾的最長回文編號
// len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文子字串在整個字串中的出現次數
// fail[s]是他長度次長的回文後綴·aba的fail是a
const int MXN = 1000010;
struct PalT{
  int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
  int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN]={-1};
int newNode(int l,int f){
    len[tot]=1,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++;
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0;
    newNode(0,1), newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

Longest Increasing Subsequence

```
vector<int> getLIS(vector<int> a){
  vector<int> lis:
  for(int i : a){
    if(lis.empty() || lis.back() < i)</pre>
                                           lis.push Back(
         i);
               *lower_bound(lis.begin(), lis.end(), i) =
         i;
  return lis;
}
```

5.17 Graph Thearom

預集條件 $V_j - V_i \leq W$ addEdge (V_i, V_j, W) and run bellman-ford or 6.3 Longest Common Subsequence O(nlgn)spfa

6.4 KMP

```
/* len-failure[k]:
在k結尾的情況下,這個子字串可以由開頭
長度為(len-failure[k])的部分重複出現來表達
failure[k] 為次長相同前綴後綴
如果我們不只想求最多,而且以0-base做為考量
 · 那可能的長度由大到小會是
failuer[k] · failure[failuer[k]-1]
、failure[failure[failuer[k]-1]-1]..
直到有值為0為止 */
int failure[MXN];
vector<int> KMP(string& t, string& p) {
    vector<int> ret;
    if(p.size() > t.size()) return ret;
    for(int i = 1, j = failure[0] = -1; i < p.size(); i
        ++) {
       while(j \ge 0 \& p[j + 1] != p[i]) j = failure[j]
        if(p[j + 1] == p[i]) j++;
       failure[i] = j;
    for(int i = 0, j = -1; i < t.size(); i++) {
       while (j \ge 0 \& p[j + 1] != t[i]) j = failure[
           j];
       if(p[j + 1] == t[i]) j++;
       if(j == p.size() - 1) {
    ret.push_back(i - p.size() + 1);
           j = failure[j];
    return ret;
}
```

6.5 SAIS O(n)

```
/*** SA· 將字串的所有後綴排序後的數組 ***/
/* SA[i]儲存排序後第i小的後綴從哪裡開始 */
/**** H[i] 為第i小的字串跟第i-1小的LCP ***/
/**** 註:LCP(Longest Common Prefix) ****/
/**** ex:S = "babd", SA[0] = 1("abd") ****/
/** SA[1] = 0("babd"), SA[2] = 2("bd") **/
/*** H[0] = 0, H[1] = 0, H[2] = 1("b") ***/
/* 傳入參數:ip 陣列放字串·len為字串長度 */
/* 需保證ip[len]為0, 且字串裡的元素不為0 */
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
 hei[N], r[N];
int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
    memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
mkhei(n);
```

```
void mkhei(int n){
     REP(i,n) r[\_sa[i]] = i;
     hei[0] = 0;
     REP(i,n) if(r[i]) {
       int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
       while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
       hei[r[i]] = ans;
  void sais(int *s, int *sa, int *p, int *q, bool *t,
       int *c, int n, int z){
     bool uniq = t[n-1] = true, neq;
     int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
         lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
   memcpy(x, c, sizeof(int) * z); \
    \label{eq:memcpy} \begin{array}{lll} \text{memcpy}(x + 1, \ c, \ sizeof(int) * (z - 1)); \\ \text{REP}(i,n) \ if(sa[i] \&\& \ !t[sa[i]-1]) \ sa[x[s[sa[i]-1]]) \end{array}
         ]-1]]++] = sa[i]-1;
    memcpy(x, c, sizeof(int) * z); \
for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i
          ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
     MS0(c, z);
     REP(i,n) uniq \&= ++c[s[i]] < 2;
     REP(i,z-1) c[i+1] += c[i];
     if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
    ]]]=p[q[i]=nn++]=i);
REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
       neq=lst<0||memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa]
            [i])*sizeof(int));
       ns[q[lst=sa[i]]]=nmxz+=neq;
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
    MAGIC(for(int i = nn - 1; i \ge 0; i--) sa[--x[s[p[
         nsa[i]]]] = p[nsa[i]]);
}sa;
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
  // should padding a zero in the back
  // ip is int array, len is array length
  // ip[0..n-1] != 0, and ip[len] = 0
  ip[len++] = 0;
  sa.build(ip, len, 128);
for (int i=0; i<len; i++) {</pre>
    H[i] = sa.hei[i + 1];
     SA[i] = sa.\_sa[i + 1];
   // resulting height, sa array \in [0,len)
}
```

6.6 Z Value O(n)

```
//z[i] = lcp(s[1...n-1],s[i...n-1])
int z[MAXN];
void Z_value(const string& s) {
  int i, j, left, right, len = s.size();
  left=right=0; z[0]=len;
  for(i=1;i<len;i++) {
    j=max(min(z[i-left],right-i),0);
    for(;i+j<len&&s[i+j]==s[j];j++);
    z[i]=j;
    if(i+z[i]>right) {
      right=i+z[i];
      left=i;
    }
}
```

6.7 Manacher Algorithm O(n)

```
|// 求以每個字元為中心的最長回文半徑
|// 頭尾以及每個字元間都加入一個
```

```
// 沒出現過的字元,這邊以'@'為例
// s為傳入的字串·len為字串長度
// z為儲存答案的陣列 (有包含'@'要小心)
// ex: s = "abaac" -> "@a@b@a@a@c@"
                      [12141232121]
void z_value_pal(char *s,int len,int *z){
  len=(len<<1)+1;
  for(int i=len-1;i>=0;i--)
    s[i]=i&1?s[i>>1]:'@';
  z[0]=1;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i < r?min(z[l+l-i],r-i):1;
    while(i-z[i] >= 0 \& i+z[i] < len \& s[i-z[i]] == s[i+z[i]])
        ++z[i];
    if(i+z[i]>r) l=i,r=i+z[i];
| } }
```

6.8 Smallest Rotation

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
    if(a+k == b || s[a+k] < s[b+k])
      {b += max(0, k-1); break;}
  if(s[a+k] > s[b+k]) {a = b; break;}
  } return a;
}
```

6.9 Cyclic LCS

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]=\{0,-1,-1,-1,-1,0\};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al,j=bl,l=0;
  while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++;
    i+=mov[dir][0];
    j+=mov[dir][1];
  return 1;
inline void reroot(int r) { // r = new base row
  int i=r, j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
  if(j>bl) return;
  pred[i][j]=L;
while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
    pred[i][j]=L;
} else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
      i++;
      j++:
      pred[i][j]=L;
    } else {
      j++;
} } }
int cyclic_lcs() {
  // a, b, al, bl should be properly filled
  // note: a WILL be altered in process
               - concatenated after itself
  char tmp[MAXL];
  if(al>bl) {
    swap(al,bl);
    strcpy(tmp,a);
    strcpy(a,b);
    strcpy(b,tmp);
  strcpy(tmp,a);
  strcat(a,tmp);
```

```
// basic lcs
for(int i=0;i<=2*al;i++) {</pre>
   dp[i][0]=0;
   pred[i][0]=U;
for(int j=0;j<=bl;j++) {
  dp[0][j]=0;</pre>
   pred[0][j]=L;
for(int i=1;i<=2*al;i++) {
  for(int j=1;j<=bl;j++) {
    if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
    if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;</pre>
      else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
      if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
else if(a[i-1]==b[j-1]) pred[i][j]=LU;
      else pred[i][j]=U;
} }
// do cyclic lcs
int clcs=0;
for(int i=0;i<al;i++) {</pre>
   clcs=max(clcs,lcs_length(i));
   reroot(i+1);
// recover a
a[al]='\0':
return clcs;
```

6.10 Hash

```
//字串雜湊前的idx是0-base · 雜湊後為1-base
//即區間為 [0,n-1] -> [1,n]
//若要取得區間[L,R]的值則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為len的字串和
//(h[i+len-1] - h[i-1] * qpow(p, len) % modl + modl)
//從j開始長度為len的字串是否相同
#define x first
#define y second
pair<int,int> Hash[MXN];
void build(const string& s){
  pair<int,int> val = make_pair(0,0);
  Hash[0]=val;
  for(int i=1; i<=s.size(); i++){
val.x = (val.x * P1 + s[i-1]) % MOD;
val.y = (val.y * P2 + s[i-1]) % MOD;</pre>
  Hash[i] = val;
}
bool cmp( int i, int j, int len ) {
     return ((Hash[i+len-1].x-Hash[i-1].x*qpow(P1,len)%
         MOD+MOD)%MOD == (Hash[j+len-1].x-Hash[j-1].x*
          qpow(P1,len)%MOD+MOD)%MOD)
     && ((Hash[i+len-1].y-Hash[i-1].y*qpow(P2,len)%MOD+
         MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow(
         P2,len)%MOD+MOD)%MOD);
}
```

7 Data Structure

7.1 Segment tree

```
//!!!注意build()時初始化用的陣列也是1-base
//!!!query(0, 0) 會報錯
#define cl(x) (x*2)
#define cr(x) (x*2+1)

struct segmentTree {
    int n;
    vector<int> seg, tag, cov;
    segmentTree(int _n): n(_n) {
        seg = tag = cov = vector<int>(n * 4, 0);
    }
    void push(int i, int L, int R) {
        if(cov[i]) {
```

```
seg[i] = cov[i] * (R - L + 1);
             if(\bar{L} < R) {
                 cov[cl(i)] = cov[cr(i)] = cov[i];
                 tag[cl(i)] = tag[cr(i)] = 0;
             cov[i] = 0;
        if(tag[i]) {
             seg[i] += tag[i] * (R - L + 1);
             if(L < R) {
                 tag[cl(i)] += tag[i];
                 tag[cr(i)] += tag[i];
             tag[i] = 0;
        }
    void pull(int i, int L, int R) {
        if(L >= R) return;
        int mid = L + R \gg 1;
        push(cl(i), L, mid);
        push(cr(i), mid + 1, R);
        seg[i] = seg[cl(i)] + seg[cr(i)];
    void build(vector<int>& arr, int i = 1, int L = 1,
        int R = -1) {
        if(R == -1) \bar{R} = n;
        if(L == R) return void(seg[i] = arr[L]);
        int mid = L + R \gg 1;
        build(arr, cl(i), L, mid);
build(arr, cr(i), mid + 1, R);
pull(i, L, R);
    int query(int rL, int rR, int i = 1, int L = 1, int
         R = -1) {
        if(R == -1) R = n;
        push(i, L, R);
        if(rL <= L && R <= rR) return seg[i];</pre>
        int mid = L + R \gg 1, ret = 0;
        if(rL <= mid) ret += query(rL, rR, cl(i), L,</pre>
             mid);
        if(mid < rR ) ret += query(rL, rR, cr(i), mid +</pre>
              1, R);
        return ret;
    void update(int rL, int rR, int val, int i = 1, int
          L = 1, int R = -1) {
        if(R == -1) R = n;
        push(i, L, R);
        if(rL <= L && R <= rR) return void(tag[i] = val
        int mid = L + R \gg 1;
        if(rL <= mid) update(rL, rR, val, cl(i), L, mid</pre>
        if(mid < rR ) update(rL, rR, val, cr(i), mid +</pre>
             1, R);
        pull(i, L, R);
    void cover(int rL, int rR, int val, int i = 1, int
        L = 1, int R = -1) {
        if(R = -1) R = n;
        push(i, L, R);
        if(rL <= L && R <= rR) return void(cov[i] = val</pre>
        int mid = L + R \gg 1;
        if(rL <= mid) cover(rL, rR, val, cl(i), L, mid)</pre>
        if(mid < rR ) cover(rL, rR, val, cr(i), mid +</pre>
        pull(i, L, R);
    }
   Test Case:
    1 2 3 4
    2 1 3
    1 1 3 1
    2 1 3
    1 1 4 1
    2 1 4
*/
```

7.2 持久化 SMT

```
struct node{
  node *l, *r;
  int val;
vector<node *> ver;
int arr[MXN] = \{0\};
//0-base
struct SegmentTree{
  int n;
  node *root;
  void build(int _n){
    n = _n;
   root = build(0, n-1);
  node* build(int L, int R){
    node *x = new node();
    if(L == R){ x->val = arr[L]; return x;}
    int mid = (L+R)/2;
    x->l = build(L, mid)
    x->r = build(mid + 1, R);
    x->val = x->l->val + x->r->val;
    return x;
  int query(node *ro, int L, int R){return query(ro, 0,
       n-1, L, R);}
  int query(int L, int R){return query(root, 0, n-1, L,
       R);}
  int query(node *x, int L, int R, int recL, int recR){
    if(recL <= L && R <= recR) return x->val;
int mid = (L+R)/2, res = 0;
    if(recL <= mid) res += query(x->1, L, mid, recL,
        recR);
    if(mid < recR) res += query(x->r, mid+1, R, recL,
        recR);
    return res;
  void update(int pos, int v){update(root, 0, n-1, pos,
       v);}
  void update(node *x, int L, int R, int pos, int v){
    if(L == R){x-val = v; arr[L] = v; return;}
    int mid = (L+R)/2;
    if(pos <= mid) update(x->1, L, mid, pos, v);
                  update(x->r, mid+1, R, pos, v);
    else
    x->val = x->l->val + x->r->val;
  node *update_ver(node *pre, int l, int r, int pos,
      int v){
    node *x = new node();
                            //當前位置建立新節點
    if(l == r){
     x->val = v;
      return x;
    int mid = (l+r)>>1;
    if(pos <= mid){ //更新左邊
     x->l = update_ver(pre->l, l, mid, pos, v); //左邊
          節點連向新節點
     x->r = pre->r; //  右邊連到原本的右邊
    else{ //更新右邊
     x->l = pre->l; // 左邊連到原本的左邊
      x->r = update_ver(pre->r, mid+1, r, pos, v); //
          右邊節點連向新節點
    x->val = x->l->val + x->r->val;
    return x;
}} seg;
                             //修改位置 x 的值為 v
void add_ver(int x,int v){
    ver.push_back(seg.update_ver(ver.back(), 0, seg.n
        -1, x, v));
}
```

7.3 持久化並查集

```
int n:
     vector<int> fa, sz;
     vector<tuple<int, int, int, int>> ver;
     DSU(int _n): n(_n), fa(n), sz(n, 1) {
          iota(fa.begin(), fa.end(), 0);
     int find(int x) {
         return fa[x] == x ? x : find(fa[x]);
     void merge(int x, int y) {
    x = find(x), y = find(y);
    if(sz[x] < sz[y]) swap(x, y);</pre>
         ver.push\_back({x, sz[x], y, fa[y]});
         if(x == y) return;
         sz[x] += sz[y];
         fa[y] = x;
     void undo() {
         if(ver.empty()) return;
         auto [x, szx, y, fy] = ver.back();
         ver.pop_back();
         sz[x] = szx;
         fa[y] = fy;
};
```

7.4 Trie

```
struct trie{
 trie *nxt[26];
            //紀錄有多少個字串以此節點結尾
  int cnt;
             //有多少字串的前綴包括此節點
  int sz;
  trie():cnt(0),sz(0){
     memset(nxt,0,sizeof(nxt));
 }
};
trie *root = new trie(); //創建新的字典樹
void insert(string& s){
 trie *now = root; // 每次從根結點出發
  for(auto i:s){
   now->sz++:
   if(now->nxt[i-'a'] == NULL){
     now->nxt[i-'a'] = new trie();
   now = now->nxt[i-'a']; //走到下一個字母
  now->cnt++; now->sz++;
int query_prefix(string& s){ //查詢有多少前綴為 s
 trie *now = root;
                    // 每次從根結點出發
  for(auto i:s){
   if(now->nxt[i-'a'] == NULL){
     return 0;
   now = now->nxt[i-'a'];
  return now->sz;
int query_count(string& s){ //查詢字串 s 出現欠數
  trie *now = root;
                     // 每次從根結點出發
  for(auto i:s){
   if(now->nxt[i-'a'] == NULL){
     return 0;
   now = now->nxt[i-'a'];
  }
  return now->cnt;
}
```

7.5 Treap (interval reverse)

//拆出[a,b]區間就如同下面所展示先使用splitByTh()拆出//左右,再把左區間拆成1,m最後merge()回去

```
//反轉區間時又記得使用^=可以直接反轉01
//treap 拆區間時從後面拆是因為這樣 [a,b] 的關係
//不用重新考慮·要是先拆前面b的位置會變成b-a+1
//0-base
//splitByTh(root,a-1,l,m);
//splitByTh(m,b-a+1,m,r);
mt19937 gen(chrono::steady_clock::now().
     time_since_epoch().count());
struct Treap {
  int key, pri, sz, tag, sum;
  Treap *L, *R;
  Treap( int val ) {
     sum=key=val, pri=gen(), sz=1, tag=0;
    L=R=NULL;
};};
int Size( Treap *a ) { return !a?0:a->sz;}
void pull( Treap *a ) {
  a \rightarrow sz = Size(a \rightarrow L) + Size(a \rightarrow R) + 1;
  a->sum=a->key;
  if( a \rightarrow L ) a \rightarrow sum += a \rightarrow L \rightarrow sum;
  if( a \rightarrow R ) a \rightarrow sum + = a \rightarrow R \rightarrow sum;
void push( Treap *a ) {
  if( a && a->tag ) {
    swap(a->L,a->R);
if( a->L ) a->L->tag^=1;
if( a->R ) a->R->tag^=1;
    a \rightarrow tag=0;
Treap *merge(Treap *a, Treap *b) {
  if(!a || !b ) return a?a:b;
  push(a), push(b);
  if( a->pri > b->pri ) {
    a \rightarrow R = merge(a \rightarrow R, b);
    pull(a); return a;
  b \rightarrow L = merge(a, b \rightarrow L);
  pull(b); return b;
}
void print(Treap *a) {
  if( !a ) return;
  push(a);
  print(a->L);
  cout.put(a->key);
  print(a->R);
Treap *buildTreap( int n, string& str ) {
  Treap *root=NULL;
  for( int i=0 ; i < n ; i++ )</pre>
    root=merge(root, new Treap(str[i]));
  return root;
}
void splitbyk( Treap *x, int k, Treap *&a, Treap *&b )
  if(!x) a=b=NULL;
  else if( x->key <= k ) {
    a=x:
    splitbyk(x->R,k,a->R,b);
    pull(a);
  else {
    splitbyk(x->L,k,a,b->L);
    pull(b);
  }
}
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
  if( !x ) { a=b=NULL; return; }
  push(x);
  if( Size(x->L)+1 \le k ) {
    splitByTh(x->R,k-Size(x->L)-1,a->R,b);
    pull(a);
  else {
    b=x
    splitByTh(x->L,k,a,b->L);
    pull(b);
```

```
}
signed main() {
    string str;
    int n, m;
    cin>n>>m>>str;
    Treap *root;
    root=buildTreap(n,str);
    for( int i=0 ; i < m ; i++ ) {
        int a, b;
        cin>>a>>b;
        Treap *l, *m, *r;
        splitByTh(root,b,l,r);
        splitByTh(l,a-1,l,m);
        m->tag^=1;
        root=merge(l,merge(m,r));
    }
    print(root);
}
```

7.6 BIT

```
#define lowbit(x) (x&-x)
struct BIT {
    int n;
    vector<int> bit;
    BIT(int _n):n(_n), bit(n + 1) {}
    void update(int x, int val) {
        for(; x <= n; x += lowbit(x)) bit[x] += val;
    }
    void update(int L, int R, int val) {
            update(L, val), update(R + 1, -val);
    }
    int query(int x) {
        int res = 0;
        for(; x; x -= lowbit(x)) res += bit[x];
        return res;
    }
    int query(int L, int R) {
        return query(R) - query(L - 1);
    }
};</pre>
```

7.7 Black Magic

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
    tree_order_statistics_node_update> set_t;
tree<int, null_type, less_equal<int>, rb_tree_tag,
    tree_order_statistics_node_update> mt_t;
#include <ext/pb_ds/assoc_container.hpp>
typedef cc_hash_table<int,int> umap_t;
// gp_hash_table<int, int>
typedef priority_queue<int> heap;
#include<ext/rope>
using namespace __gnu_cxx;
int main(){
  // Insert some entries into s.
  set_t s; s.insert(12); s.insert(505);
  // The order of the keys should be: 12, 505.
 assert(*s.find_by_order(0) == 12)
 assert(*s.find_by_order(3) == 505);
 // The order of the keys should be: 12, 505. assert(s.order_of_key(12) == 0);
 assert(s.order_of_key(505) == 1);
  // Erase an entry.
  s.erase(12);
 // The order of the keys should be: 505.
 assert(*s.find_by_order(0) == 505);
  // The order of the keys should be: 505.
 assert(s.order_of_key(505) == 0);
  // if we want to delete less_equal tag tree
  mt_t.erase(mt_t.find_by_order(mt_t.order_of_key(val))
      );
  heap h1 , h2; h1.join( h2 );
```

```
rope<char> r[ 2 ];
r[ 1 ] = r[ 0 ]; // persistenet
string t = "abc";
r[ 1 ].insert( 0 , t.c_str() );
r[ 1 ].erase( 1 , 1 );
cout << r[ 1 ].substr( 0 , 2 );
}</pre>
```

8 Others

8.1 SOS dp

8.2 De Brujin sequence

```
// return cyclic array of length k^n such that every
// array of length n using 0~k-1 appears as a subarray.
vector<int> DeBruijn(int k,int n){
   if(k=1) return {0};
   vector<int> aux(k*n),res;
   function<void(int,int)> f=[&](int t,int p)->void{
      if(t>n){   if(n%p==0)
          for(int i=1;i<=p;++i) res.push_back(aux[i]);
   }else{
      aux[t]=aux[t-p]; f(t+1,p);
      for(aux[t]=aux[t-p]+1;aux[t]<k;++aux[t]) f(t+1,t)
      ;
   }
   };
   f(1,1); return res;
}</pre>
```

8.3 CDQ 分治

```
//cdq分治使用的結構u, v, w為排序物的三個維度
//ans記錄了有幾項三維都小於等於自己
//cnt記錄了相同物有幾個·在使用cdq之前必先去重。
//並且將相同元素紀錄至cnt中,可使用map來做到這步
//cdq使用的BIT就是普通求和的BIT,大小就開維度的
//值域範圍·若值域大於2e6則要先進行離散化
struct triple {int u, v, w, ans, cnt;};
BIT *bt;
void cdq(int L, int R, vector<triple>& arr) {
  if(R - L <= 1) return;</pre>
  int mid = L + R \gg 1;
  vector<triple> temp;
 cdq(L, mid, arr), cdq(mid, R, arr);
for(int i = L, j = mid; i < mid || j < R;) {
    for(; i < mid && (j >= R || arr[i].v <= arr[j].v);</pre>
        i++) {
      bt->update(arr[i].w, arr[i].cnt);
      temp.push_back(arr[i]);
    if(j < R) {
      arr[j].ans += bt->query(arr[j].w);
      temp.push_back(arr[j]);
      j++;
    }
  for(int i = L; i < mid; i++)</pre>
    bt->update(arr[i].w, -arr[i].cnt);
  copy(temp.begin(), temp.end(), arr.begin() + L);
signed main()
  // n 個數 k 值域範圍
  int n, k;
```

8.4 3D LIS

```
#define lowbit(x) (x&-x)
const int MAXN=1e5+5;
struct BIT {
  int n;
  vector<int> bit;
  BIT( int _n ):n(_n), bit(_n+1,0) {} int query( int x ) {
    int res=0;
    for(; x > 0; x-=lowbit(x)) res=max(res,bit[x]);
    return res:
  }
  void update( int x, int val ) {
    for(; x <= n ; x+=lowbit(x) ) {
  if( val < 0 ) bit[x]=0;</pre>
      else bit[x]=max(bit[x],val);
    }
}bt(MAXN);
struct triple {
  int u, v, w, ans, cnt;
  bool operator<( triple b ) { return u<b.u; }</pre>
bool cmp( triple a, triple b ) {return a.v<b.v;}</pre>
void cdq( int L, int R, vector<triple>& arr ) {
  if( R-L <= 1 ) return;</pre>
  int mid=L+R>>1;
  cdq(L,mid,arr);
  sort(arr.begin()+L,arr.begin()+mid,cmp);
  sort(arr.begin()+mid,arr.begin()+R,cmp);
  for( int i=L, j=mid ; i < mid || j < R ; ) {
  for(; i < mid && ( j >= R || arr[i].v < arr[j].v )</pre>
    ; i++ ) bt.update(arr[i].w,arr[i].ans); if( j < R ) {
      arr[j].ans=max(bt.query(arr[j].w-1)+1,arr[j].ans)
    }
  for( int i=L ; i < mid ; i++ ) bt.update(arr[i].w,-1)</pre>
  sort(arr.begin()+L,arr.begin()+R);
  cdq(mid,R,arr);
signed main()
  ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0)
  int n, res=0;
  cin>>n;
  vector<int> ls;
  vector<triple> arr;
  for( int i=0 ; i < n ; i++ ) {
    int a, b;
    cin>>a>>b;
    arr.push_back({i,a,b,1,1});//{第一維,第二維,第三維,
         答案,數量}
    ls.push_back(b);
  sort(ls.begin(),ls.end());
  ls.resize(unique(ls.begin(),ls.end())-ls.begin());
```

8.5 Ternary Search

```
while(L <= R) {
   int ml = L + (R - L) / 3, mr = R - (R - L) / 3;
   if(L == R) return L;
   else if( checker(ml) < checker(mr) ) L = ml + 1;
   else R = mr - 1;
}</pre>
```

8.6 Max Subrectangle

```
const int N = 1e5+5;
int n, a[N], l[N], r[N];
long long ans;
int main() {
  while (cin>>n) {
    ans = 0;
    for (int i = 1; i \le n; i++) cin>>a[i], l[i] = r[i]
    for (int i = 1; i <= n; i++)
      while (l[i] > 1 \&\& a[i] <= a[l[i] - 1]) l[i] = l[
           l[i] - 1];
    for (int i = n; i >= 1; i--)
      while (r[i] < n \& a[i] <= a[r[i] + 1]) r[i] = r[
           r[i] + 1];
    for (int i = 1; i <= n; i++)
      ans = max(ans, (long long)(r[i] - l[i] + 1) * a[i]
           ]);
    cout<<ans<<"\n";
  }
}
```

8.7 Maximal Rectangle

8.8 p-Median

8.9 Tree Knapsack

8.10 質數個數

```
• 10 ^ 2 內有 25 個質數
```

- 10 ^ 3 內有 168 個質數
- 10 ^ 4 內有 1229 個質數
- 10 ^ 5 內有 9592 個質數
- 10 ^ 6 內有 78498 個質數
- 10 ^ 7 內有 664579 個質數
- 10 ^ 8 內有 5761455 個質數
- 10 ^ 9 內有 50847534 個質數
- 10 ^ 12 內有 37607912018 個質數
- 10 ^ 18 內有 24739954287740860 個質數

8.11 AC-Automaton

```
// 1-based
// n is the number of patterns
struct Automaton {
    static const int MXN = 1e6;
    int n, cnt, vis[MXN], rev[MXN], indeg[MXN], ans[MXN
         ];
    queue<int> q;
    struct trie_node {
         vector<int> son;
         int fail, flag, ans;
         trie_node(): son(27), fail(0), flag(0) {}
    } trie[MXN];
    void init(int _n) {
        n = _n, cnt = 1;
         for (int i = 1; i <= n; i++) vis[i] = 0;
    // insert a string s with number num
    // num is the index of the pattern
    void insert(string s, int num) {
         int u = 1, len = s.size();
for (int i = 0; i < len; i++) {
   int v = s[i] - 'a';</pre>
             if (!trie[u].son[v]) trie[u].son[v] = ++cnt
             u = trie[u].son[v];
```

```
if (!trie[u].flag) trie[u].flag = num;
         rev[num] = trie[u].flag;
    void getfail() {
         for (int i = 0; i < 26; i++) trie[0].son[i] =
             1;
         q.push(1);
         trie[1].fail = 0;
         while (q.size()) {
             int u = q.front(); q.pop();
int Fail = trie[u].fail;
             for (int i = 0; i < 26; i++) {
                  int v = trie[u].son[i];
                  if (!v) {
                      trie[u].son[i] = trie[Fail].son[i];
                      continue;
                  trie[v].fail = trie[Fail].son[i];
                  indeg[trie[Fail].son[i]]++;
                  q.push(v);
        }
    void topu() {
         for (int i = 1; i <= cnt; i++)
    if (!indeg[i]) q.push(i);</pre>
         while (q.size()) {
             int fr = q.front(); q.pop();
             vis[trie[fr].flag] = trie[fr].ans;
             int u = trie[fr].fail;
             trie[u].ans += trie[fr].ans;
             if (!--indeg[u]) q.push(u);
        }
    void query(string &s) {
         int u = 1, len = s.size();
         for (int i = 0; i < len; i++) u = trie[u].son[s
             [i] - 'a'], trie[u].ans++;
    void solve(string &s) {
         getfail();
         query(s);
         topu();
         for (int i = 1; i <= n; i++) ans[i] = vis[rev[i
} AC;
```

