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1 Basic

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1.1 default code

```
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
using namespace std;
ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
```

1.2 .vimrc

```
set nu rnu ts=4 sw=4 bs=2 ai hls cin mouse=a
color default
sy on
inoremap {<CR> {<CR>}<C-o>0
inoremap jk <Esc>
nnoremap J 5j
nnoremap K 5k
nnoremap run :w<bar>!g++ -std=c++14 -DLOCAL -Wfatal-
    errors -o test "%" && echo "done." && time ./test<
```

1.3 Increase Stack Size (linux)

```
#include <sys/resource.h>
void increase_stack_size() {
  const rlim_t ks = 64*1024*1024;
   struct rlimit rl;
   int res=getrlimit(RLIMIT_STACK, &rl);
   if(res==0){
     if(rl.rlim_cur<ks){</pre>
       rl.rlim_cur=ks;
       res=setrlimit(RLIMIT_STACK, &rl);
} } }
```

1.4 Misc

```
編譯參數: -std=c++14 -Wall -Wshadow (-fsanitize=
    undefined)
mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }
#define SECs ((double)clock() / CLOCKS_PER_SEC)
double startTime;
bool TIME() { // 比最大可執行時間小一點
    return SECs - startTime > 0.8;
int main() {
    startTime = SECs;
struct KeyHasher {
  size_t operator()(const Key& k) const {
    return k.first + k.second * 100000;
typedef unordered_map<Key,int,KeyHasher> map_t;
// builtin function 可以代的值為int32
```

```
|__builtin_popcountll // 二進位有幾個1
|__builtin_clzll // 左起第一個1之前0的個數
|__builtin_parityll // 1的個數的奇偶性
|__builtin_mul_overflow(a,b,&h) // a*b是否溢位
```

1.5 check

```
for ((i=0;;i++))
do
    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done
```

1.6 python-related

```
parser:
int(eval(num.replace("/","//")))
from fractions import Fraction
from decimal import Decimal, getcontext, ROUND_HALF_UP,
     ROUND_CEILING, ROUND_FLOOR
getcontext().prec = 250 # set precision
getcontext().rounding = ROUND_HALF_UP
itwo = Decimal(0.5)
two = Decimal(2)
format(x, '0.10f') # set precision
N = 200
def angle(cosT):
  """given cos(theta) in decimal return theta"""
  for i in range(N):
    cosT = ((cosT + 1) / two) ** itwo
 sinT = (1 - cosT * cosT) ** itwo
return sinT * (2 ** N)
pi = angle(Decimal(-1))
"""round to 2 decimal places"""
sum = Decimal(input())
sum.quantize(Decimal('.00'), ROUND_HALF_UP)
"""Fraction"""
x = Fraction(1, 3) # 1/3
x.as_integer_ratio() # (1, 3)
"""input list of integers"""
arr = list(map(int, input().split()))
```

2 flow

2.1 ISAP $O(V^3)$

```
struct Maxflow {
    static const int MAXV = 20010;
    static const int INF = 1000000;
    struct Edge {
        int v, c, r;
        Edge(int _v, int _c, int _r):
            v(_v), c(_c), r(_r) {}
    };
    int s, t;
    vector<Edge> G[MAXV*2];
    int iter[MAXV*2], d[MAXV*2], gap[MAXV*2], tot;
    void init(int x) {
        tot = x+2;
        s = x+1, t = x+2;
        for(int i = 0; i <= tot; i++) {
        G[i].clear();
        iter[i] = d[i] = gap[i] = 0;
    }
}</pre>
```

```
void addEdge(int u, int v, int c) {
    G[u].push_back(Edge(v, c, SZ(G[v]) ));
G[v].push_back(Edge(u, 0, SZ(G[u]) - 1));
  int dfs(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
      Edge &e = G[p][i];
       if(e.c > 0 \& d[p] == d[e.v]+1) {
         int f = dfs(e.v, min(flow, e.c));
         if(f) {
           e.c -= f;
           G[e.v][e.r].c += f;
           return f:
    if((--gap[d[p]]) == 0) d[s] = tot;
    else {
    d[p]++;
      iter[p] = 0;
      ++gap[d[p]];
    return 0;
  int solve() {
    int res = 0;
    gap[0] = tot;
    for(res = 0; d[s] < tot; res += dfs(s, INF));</pre>
    return res;
  void reset() {
    for(int i=0; i <= tot; i++) {
      iter[i]=d[i]=gap[i]=0;
```

2.2 MinCostFlow

```
struct zkwflow{
  static const int maxN=10000;
struct Edge{ int v,f,re; ll w;};
int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
  vector<Edge> E[maxN];
  void init(int _n,int _s,int _t){
    n=_n,s=_s,t=_t;
    for(int i=0;i<n;i++) E[i].clear();</pre>
  void addEdge(int u,int v,int f,ll w){
    E[u].push_back(\{v,f,(int)E[v].size(),w\});
    E[v].push\_back({u,0,(int)}E[u].size()-1,-w});
  bool SPFA(){
    fill_n(dis,n,LLONG_MAX); fill_n(vis,n,false);
    queue<int> q; q.push(s); dis[s]=0;
    while (!q.empty()){
      int u=q.front(); q.pop(); vis[u]=false;
for(auto &it:E[u]){
         if(it.f>0&&dis[it.v]>dis[u]+it.w){
           dis[it.v]=dis[u]+it.w;
           if(!vis[it.v]){
             vis[it.v]=true; q.push(it.v);
    return dis[t]!=LLONG_MAX;
  int DFS(int u,int nf){
    if(u==t) return nf;
    int res=0; vis[u]=true;
    for(int &i=ptr[u];i<(int)E[u].size();i++){</pre>
       auto &it=E[u][i]
       if(it.f>0&&dis[it.v]==dis[u]+it.w&&!vis[it.v]){
         int tf=DFS(it.v,min(nf,it.f));
         res+=tf,nf-=tf,it.f-=tf;
         E[it.v][it.re].f+=tf;
         if(nf==0){ vis[u]=false; break; }
    }
    return res;
  pair<int,ll> flow(){
    int flow=0; ll cost=0;
    while (SPFA()){
      fill_n(ptr,n,0);
```

```
int f=DFS(s,INT_MAX)
       flow+=f; cost+=dis[t]*f;
     return{ flow,cost };
  } // reset: do nothing
} flow;
2.3 Dinic O(V^2E)
#define SZ(x) (int)x.size()
#define PB push_back
struct Dinic{
  struct Edge{ int v,f,re; };
int n,s,t,level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
  n = _n;  s = _s;  t = _t;
  for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
    E[u].PB({v,f,SZ(E[v])});
     E[v].PB({u,0,SZ(E[u])-1});
  bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;
queue<int> que;
     que.push(s);
     level[s] = 0;
    while (!que.empty()){
       int u = que.front(); que.pop();
       for (auto it : E[u]){
  if (it.f > 0 && level[it.v] == -1){
            level[it.v] = level[u]+1;
            que.push(it.v);
     return level[t] != -1;
  int DFS(int u, int nf){
    if (u == t) return nf;
     int res = 0;
     for (auto &it : E[u]){
       if (it.f > 0 && level[it.v] == level[u]+1){
         int tf = DFS(it.v, min(nf,it.f));
         res += tf; nf -= tf; it.f -= tf;
E[it.v][it.re].f += tf;
         if (nf == 0) return res;
     } }
     if (!res) level[u] = -1;
     return res;
  int flow(int res=0){
     while (BFS())
       res += DFS(s,2147483647);
     return res;
} }flow;
```

2.4 Kuhn Munkres 最大完美二分匹配 $O(n^3)$

```
ll t = lx[x]+ly[y]-g[x][y];
            if(t==0){
              pa[y]=x
              if(!my[y]){augment(y);return;}
              vy[y]=1, q.push(my[y]);
            }else if(sy[y]>t) pa[y]=x,sy[y]=t;
         _ }
       11 cut = INF;
       for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;</pre>
          if(vy[j]) ly[j] += cut;
         else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
          if(!my[y]){augment(y); return;}
          vy[y]=1, q.push(my[y]);
   } } }
   ll solve(){
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
     fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
    lx[x] = max(lx[x], g[x][y]);</pre>
     for(int x=1; x<=n; ++x) bfs(x);</pre>
     11 \text{ ans} = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
     return ans;
} }graph;
```

2.5 SW min-cut (不限 S-T 的 min-cut) $O(V^3)$

```
// global min cut
struct SW{ // O(V^3)
  int n,vst[MXN],del[MXN];
  int edge[MXN][MXN],wei[MXN];
  void init(int _n){
    n = _n; FZ(edge); FZ(del);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  void search(int &s, int &t){
    FZ(vst); FZ(wei);
s = t = -1;
    while (true){
       int mx=-1, cur=0;
for (int i=0; i<n; i++)</pre>
         if (!del[i] && !vst[i] && mx<wei[i])</pre>
       cur = i, mx = wei[i];
if (mx == -1) break;
       vst[cur] = 1;
       s = t; t = cur;
for (int i=0; i<n; i++)
         if (!vst[i] && !del[i]) wei[i] += edge[cur][i];
    }
  int solve(){
    int res = 2147483647;
     for (int i=0,x,y; i<n-1; i++){</pre>
       search(x,y);
       res = min(res,wei[y]);
       del[y] = 1;
       for (int j=0; j<n; j++)
         edge[x][j] = (edge[j][x] += edge[y][j]);
    return res;
} }graph;
```

2.6 Max flow with lower/upper bound

```
// flow use ISAP
// Max flow with lower/upper bound on edges
// source = 1 , sink = n
int in[ N ] , out[ N ];
int l[ M ] , r[ M ] , a[ M ] , b[ M ];//0-base,a下界,b
上界
int solve(){
```

```
flow.init( n );
                     //n為點的數量,m為邊的數量,點是1-
for( int i = 0 ; i < m ; i ++ ){
  in[ r[ i ] ] += a[ i ];
out[ l[ i ] ] += a[ i ];
flow.addEdge( l[ i ] , r[ i ] , b[ i ] - a[ i ] );
// flow from l[i] to r[i] must in [a[ i ], b[ i ]]
int nd = 0;
for( int i = 1 ; i <= n ; i ++ ){
  if( in[ i ] < out[ i ] ){</pre>
     flow.addEdge( i , flow.t , out[ i ] - in[ i ] );
nd += out[ i ] - in[ i ];
  if( out[ i ] < in[ i ] )
     flow.addEdge( flow.s , i , in[ i ] - out[ i ] );
// original sink to source
flow.addEdge( n , 1 , INF );
if( flow.maxflow() != nd )
  return -1; // no solution
int ans = flow.G[ 1 ].back().c; // source to sink
flow.G[1].back().c = flow.G[n].back().c = 0;
// take out super source and super sink
for( size_t i = 0 ; i < flow.G[ flow.s ].size() ; i</pre>
  ++ ){
flow.G[ flow.s ][ i ].c = 0;
  Edge \&\bar{e} = flow.\bar{G}[\bar{f}low.s][i];
  flow.G[ e.v ][ e.r ].c = 0;
for( size_t i = 0 ; i < flow.G[ flow.t ].size() ; i</pre>
  flow.G[ flow.t ][ i ].c = 0;
Edge &e = flow.G[ flow.t ][ i ];
  flow.G[e.v][e.r].c = 0;
flow.addEdge( flow.s , 1 , INF
flow.addEdge( n , flow.t , INF );
flow.reset();
return ans + flow.maxflow();
```

2.7 Flow Method

source with

```
Maximize c^T x subject to Ax \le b, x \ge 0;
with the corresponding symmetric dual problem,
Minimize b^T y subject to A^T y \geq c, y \geq 0.
Maximize c^T x subject to Ax \le b;
with the corresponding asymmetric dual problem,
Minimize b^T y subject to A^T y = c, y \ge 0.
Minimum vertex cover on bipartite graph =
Maximum matching on bipartite graph
Minimum edge cover on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
Independent set on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
找出最小點覆蓋,做完dinic之後,從源點dfs只走還有流量的
    邊, 左邊沒被走到的點跟右邊被走到的點就是答案, 其他
    點為最大獨立集
Maximum density subgraph (\sum W_e + \sum W_v) / |V|
Binary search on answer:
For a fixed D, construct a Max flow model as follow:
Let S be Sum of all weight( or inf)
1. from source to each node with cap = S
2. For each (u,v,w) in E, (u->v,cap=w), (v->u,cap=w)
3. For each node v, from v to sink with cap = S + 2 * D - deg[v] - 2 * (W of v)
where deg[V] = \sum weight of edge associated with v If maxflow < S * |V|, D is an answer.
Requiring subgraph: all vertex can be reached from
```

```
edge whose cap > 0.
2-SAT: (a or b) and (not a or c) => (a->c) and (b->c)
If a and b are in the same SCC, then it is impossible.
If a and not a are in the same SCC, then it is
   impossible.
```

3 Math

3.1 FFT

```
// const int MAXN = 262144;
// (must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx; //real() ,imag()
const ld PI = acosl(-1);
const cplx I(0, 1);
cplx omega[MAXN+1];
void pre_fft(){
  for(int i=0; i<=MAXN; i++)
  omega[i] = exp(i * 2 * PI / MAXN * I);</pre>
// n must be 2^k
void fft(int n, cplx a[], bool inv=false){
  int basic = MAXN / n;
  int theta = basic;
  for (int m = n; m >= 2; m >>= 1) {
    int mh = m >> 1;
    for (int i = 0; i < mh; i++) {
  cplx w = omega[inv ? MAXN-(i*theta%MAXN)</pre>
                            : i*theta%MAXN];
       for (int j = i; j < n; j += m) {
         int k = j + mh;
         cplx x = a[j] - a[k];
        a[j] += a[k];
        a[k] = w * x;
    } }
    theta = (theta * 2) % MAXN;
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
    for (int k = n \gg 1; k \gg (i ^= k); k \gg 1);
    if (j < i) swap(a[i], a[j]);</pre>
  if(inv) for (i = 0; i < n; i++) a[i] /= n;
cplx arr[MAXN+1];
inline void mul(int _n,ll a[],int _m,ll b[],ll ans[]){
  int n=1,sum=_n+_m-1;
  while(n<sum)</pre>
    n << =1;
  for(int i=0;i<n;i++) {</pre>
    double x=(i<_n?a[i]:0), y=(i<_m?b[i]:0);
    arr[i]=complex<double>(x+y,x-y);
  fft(n,arr);
  for(int i=0;i<n;i++)</pre>
    arr[i]=arr[i]*arr[i];
  fft(n,arr,true);
  for(int i=0;i<sum;i++)</pre>
    ans[i]=(long long int)(arr[i].real()/4+0.5);
```

3.2 O(1)mul

```
LL mul(LL x,LL y,LL mod){
  LL ret=x*y-(LL)((long double)x/mod*y)*mod;
  // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
  return ret<0?ret+mod:ret;
}</pre>
```

3.3 Faulhaber $(\sum_{i=1}^{n} i^{p})$

```
/* faulhaber's formula - 
 * cal power sum formula of all p=1~k in 0(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinactories
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
  while(b) {
    int q,t;
    q=a/b; t=b; b=a-b*q; a=t;
    t=b0; b0=a0-b0*q; a0=t;
    t=b1; b1=a1-b1*q; a1=t;
  return a0<0?a0+mod:a0;</pre>
inline void pre() {
  /* combinational
  for(int i=0;i<=MAXK;i++) {</pre>
    cm[i][0]=cm[i][i]=1;
    for(int j=1; j<i; j++)</pre>
       cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);
  /* inverse */
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
   /* bernoulli */
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
  for(int i=2;i<MAXK;i++) {</pre>
    if(i&1) { b[i]=0; continue; }
    b[i]=1;
    for(int j=0;j<i;j++)</pre>
       b[i]=sub(b[i],
                 mul(cm[i][j],mul(b[j], inv[i-j+1])));
  /* faulhaber */
  // sigma_x=1~n {x^p} = 
// 1/(p+1) * sigma_j=0~p {C(p+1,j)*Bj*n^(p-j+1)}
  for(int i=1;i<MAXK;i++) {</pre>
    co[i][0]=0;
    for(int j=0;j<=i;j++)</pre>
       co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]))
  }
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
inline int solve(int n,int p) {
  int sol=0,m=n;
  for(int i=1;i<=p+1;i++) {</pre>
    sol=add(sol,mul(co[p][i],m));
    m = mul(m, n);
  return sol;
```

3.4 Chinese Remainder

3.5 Miller Rabin

```
// n < 4,759,123,141
                              3: 2, 7, 61
                                   2, 13, 23, 1662803
6: pirmes <= 13
// n < 1,122,004,669,633
// n < 3,474,749,660,383
// n < 2^{64}
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL magic[]={}
bool witness(LL a, LL n, LL u, int t){
  if(!a) return 0;
  LL x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    LL nx=mul(x,x,n);
    if(nx==1\&&x!=1\&&x!=n-1) return 1;
    x=nx;
  return x!=1;
bool miller_rabin(LL n) {
  int s=(magic number size)
  // iterate s times of witness on n
  if(n<2) return 0;</pre>
  if(!(n&1)) return n == 2;
  ll u=n-1; int t=0;
  // n-1 = u*2^t
  while(!(u&1)) u>>=1, t++;
  while(s--){
    LL a=magic[s]%n;
    if(witness(a,n,u,t)) return 0;
  return 1;
}
```

3.6 Pollard Rho

```
// does not work when n is prime 0(n^(1/4))
LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
LL pollard_rho(LL n) {
   if(!(n&1)) return 2;
   while(true){
      LL y=2, x=rand()%(n-1)+1, res=1;
      for(int sz=2; res==1; sz*=2) {
        for(int i=0; i<sz && res<=1; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
        }
        y = x;
    }
   if (res!=0 && res!=n) return res;
}</pre>
```

3.7 Josephus Problem

```
int josephus(int n, int m){ //n人每m次
  int ans = 0;
  for (int i=1; i<=n; ++i)
      ans = (ans + m) % i;
  return ans;
}</pre>
```

3.8 Matrix

```
\{1, 0\}
};
int mypow(int y){
 while(y){
   if( y&1 ) { ans = mul(ans, base); } //實作矩陣乘法
   base = mul(base, base);//實作矩陣乘法
   y >>= 1;
 return ans[0][0];
```

Gaussian Elimination

```
const int GAUSS_MOD = 100000007LL;
struct GAUSS{
  int n;
  vector<vector<int>> v;
  int ppow(int a , int k){
  if(k == 0) return 1;
     if(k % 2 == 0) return ppow(a * a % GAUSS_MOD , k >>
          1);
     if(k \% 2 == 1) return ppow(a * a % GAUSS_MOD , k >>
           1) * a % GAUSS_MOD;
  vector<int> solve(){
     vector<int> ans(n);
     REP(now , 0 , n){
       REP(i , now , n) if(v[now][now] == 0 \& v[i][now]
             != 0)
       swap(v[i] , v[now]); // det = -det;
if(v[now][now] == 0) return ans;
       int inv = ppow(v[now][now] , GAUSS_MOD - 2);
       REP(i , 0 , n) if(i != now){
          int tmp = v[i][now] * inv % GAUSS_MOD;
         REP(j , now , n + 1) (v[i][j] += GAUSS_MOD -
tmp * v[now][j] % GAUSS_MOD) %= GAUSS_MOD;
       }
    REP(i , 0 , n) ans[i] = v[i][n + 1] * ppow(v[i][i]
      , GAUSS_MOD - 2) % GAUSS_MOD;
     return ans;
  // gs.v.clear() , gs.v.resize(n , vector<int>(n + 1 ,
         0));
} gs;
```

Inverse Matrix

```
int GAUSS_MOD;
struct GAUSS{
 int n;
 vector<vector<int> > v;
 vector<vector<int> > rev;
 int mul(int x,int y,int mod){
   int ret=x*y-(int)((long double)x/mod*y)*mod;
   return ret<0?ret+mod:ret;</pre>
 int ppow(int a, int b){//res=(a^b)%m
   int res=1, k=a;
   while(b){
     if((b&1)) res=mul(res,k,GAUSS_MOD)%GAUSS_MOD;
     k=mul(k,k,GAUSS_MOD)%GAUSS_MOD;
     b>>=1;
   return res%GAUSS_MOD;
 bool solve(){
   for(int now = 0; now < n; now++){
     for(ch = now; ch < n && !v[ch][now]; ch++);</pre>
     if(ch >= n) return 0;
     swap(v[i], v[now]); // det = -det;
         swap(rev[i], rev[now]);
     }
```

```
if(v[now][now] == 0) return 0;
     GAUSS_MOD) %= GAUSS_MOD;
          (rev[i][j] += GAUSS_MOD - tmp * rev[now][j] %
               GAUSS_MOD) %= GAUSS_MOD;
       }
     }
    }
    return 1;
}} gs;
signed main(){
  int n, p; //n*n matrix, MOD=p
  cin>>n>>p; //if(!n && !p) return 0;
  GAUSS\_MOD = p; gs.n = n;
  gs.v.clear(), gs.v.resize(n + 1, vector<int>(n + 2,
       0));
  gs.rev.clear() , gs.rev.resize(n + 1, vector<int>(n +
  2 , 0));
for(int i = 0; i < n; i++){
    for(int j = 0; j < n; j++){
      cin>>gs.v[i][j];
      if(i == j) gs.rev[i][j] = 1;
  if(!gs.solve()) cout << "singular\n";</pre>
  else{
    for(int i = 0; i < n; i++){
      int inv = gs.ppow(gs.v[i][i] , p - 2);
      for(int j = 0; j < n; j++)
          cout << (gs.rev[i][j] * inv % p) <<" ";</pre>
      cout<<"\n";
   }
  }
  cout << "\n";
3.11
       模反元素
long long inv(long long a,long long m){
    long long x,y;
    long long d=exgcd(a,m,x,y);
```

```
if(d==1) return (x+m)%m;
    else return -1; //-1為無解
}
```

3.12 ax+by=gcd

```
PII gcd(int a, int b){
  if(b == \emptyset) return {1, \emptyset};
  PII q = gcd(b, a \% b);
  return {q.second, q.first - q.second * (a / b)};
int exgcd(int a,int b,long long &x,long long &y) {
    if(b == 0)\{x=1,y=0; return a;\}
    int now=exgcd(b,a%b,y,x);
    y=a/b*x;
    return now;
}
```

3.13 Discrete sqrt

```
void calcH(LL &t, LL &h, const LL p) {
  LL tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
\frac{1}{r} solve equation x^2 mod p = a
bool solve(LL a, LL p, LL &x, LL &y) {
  if(p == 2) { x = y = 1; return true; }
  int p2 = p / 2, tmp = mypow(a, p2, p);
  if (tmp == p - 1) return false;
  if ((p + 1) % 4 == 0) {
```

```
x=mypow(a,(p+1)/4,p); y=p-x; return true;
} else {
  LL t, h, b, pb; calcH(t, h, p);
  if (t >= 2) {
     do \{b = rand() \% (p - 2) + 2;
     while (mypow(b, p / 2, p) != p - 1);
  pb = mypow(b, h, p);
} int s = mypow(a, h / 2, p);
for (int step = 2; step <= t; step++) {
  int ss = (((LL)(s * s) % p) * a) % p;
}</pre>
     for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);</pre>
     if (ss + 1 == p) s = (s * pb) % p;
     pb = ((LL)pb * pb) % p;
  x = ((LL)s * a) % p; y = p - x;
} return true;
```

3.14 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
for( int i = 2 ; i < m ; i ++ )</pre>
    inv[i] = ((LL)(m - m / i) * inv[m % i]) % m;
```

Roots of Polynomial 找多項式的根 3.15

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ]; // a[0..n](coef) must be
    filled
int n; // degree of polynomial must be filled
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a[], int n, double x){
  double tmp=1,sum=0;
  for(int i=0;i<=n;i++)</pre>
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double l,double r,double a[],int n){
  int sl=sign(f(a,n,l)),sr=sign(f(a,n,r));
if(sl==0) return l; if(sr==0) return r;
  if(sl*sr>0) return inf;
  while(r-l>eps){
    double mid=(l+r)/2;
    int ss=sign(f(a,n,mid));
    if(ss==0) return mid;
    if(ss*sl>0) l=mid; else r=mid;
  }
  return 1:
void solve(int n,double a[],double x[],int &nx){
  if(n==1){ x[1]=-a[0]/a[1]; nx=1; return; }
double da[10], dx[10]; int ndx;
for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0;
  if(ndx==0){
    double tmp=binary(-inf,inf,a,n);
    if (tmp<inf) x[++nx]=tmp;</pre>
    return;
  double tmp;
tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
  for(int i=1;i<=ndx-1;i++){</pre>
    tmp=binary(dx[i],dx[i+1],a,n);
    if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
} // roots are stored in x[1..nx]
```

Combination thearom

3.16

```
const ll mod = 1e9 + 7;
ll fac[(int)2e6 + 1], inv[(int)2e6 + 1];
ll getinv(ll a){ return qpow(a, mod-2); }
void init(int n){
  fac[0] = 1;
  for(int i = 1; i <= n; i++){
    fac[i] = fac[i-1] * i % mod;
  inv[n] = getinv(fac[n]);
for(int i = n - 1; i >= 0; i--){
    inv[i] = inv[i + 1] * (i + 1) % mod;
il C(int n, int m){
  if(m > n) return 0;
  return fac[n] * inv[m] % mod * inv[n-m] % mod;
```

3.17 Primes

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231

* 999888733, 98789101, 987777733, 999991921, 1010101333

* 1010102101, 1000000000039, 10000000000037
* 2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[N], p_tbl[N];
vector<int> primes;
void sieve() {
   mu[1] = p_tbl[1] = 1;
   for( int i = 2 ; i < N ; i ++ ){
  if( !p_tbl[ i ] ){</pre>
        p_tbl[ i ] = i;
        primes.push_back( i );
        mu[ i ] = -1;
     for( int p : primes ){
  int x = i * p;
  if( x >= M ) break;
        p_tbl[ x ] = p;
mu[ x ] = -mu[ i ];
        if(i \% p == 0){
           mu[x] = 0;
           break:
} } } }
vector<int> factor( int x ){
   vector<int> fac{ 1 };
   while (x > 1)
     int fn = SZ(fac), p = p_tbl[ x ], pos = 0;
      while( x \% p == 0 ){
        for( int i = 0 ; i < fn ; i ++ )
fac.PB( fac[ pos ++ ] * p );</pre>
   } }
   return fac;
}
```

3.18 Phi

```
ll_phi(ll n){ // 計算小於n的數中與n互質的有幾個
  ll res = n, a=n;
for(ll i=2;i*i<=a;i++){ // O(sqrtN)</pre>
     if(a\%i==0){
       res = res/i*(i-1);
       while(a\%i==0) a/=i;
  if(a>1) res=res/a*(a-1);
  return res;
}
```

3.19 Result

For $n,m\in\mathbb{Z}^*$ and prime P, C(m,n) mod $P=\Pi(C(m_i,n_i))$ where m_i is the i-th digit of m in base P.

```
• Stirling approximation :
   n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\frac{1}{12n}}
• Stirling Numbers(permutation |P|=n with k cycles):
   S(n,k) = \text{coefficient of } x^k \text{ in } \prod_{i=0}^{n-1} (x+i)
• Stirling Numbers(Partition n elements into k non-empty set):
   S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n
• Pick's Theorem : A=i+b/2-1 在二維座標平面中畫上網格·對於任何簡單多邊形
   A: 面積、i: 內部的格點數、b: 邊上的格點數
• Catalan number : C_n = {2n \choose n}/(n+1)
   C_n^{n+m} - C_{n+1}^{n+m} = (m+n)! \frac{n-m+1}{n+1} for n \ge m
   C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{(n+1)!n!}
   C_0 = 1 and C_{n+1} = 2(\frac{2n+1}{n+2})C_n

C_0 = 1 and C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i} for n \ge 0
• Euler Characteristic:
   planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2
   V,E,F,C\colon number of vertices, edges, faces(regions), and compo-
• Kirchhoff's theorem :
   A_{ii}=deg(i), A_{ij}=(i,j)\in E ?-1:0, Deleting any one row, one column, and call the det(A)
\bullet Polya' theorem (c is number of \operatorname{color}\cdot m is the number of \operatorname{cycle}
   size):
   (\sum_{i=1}^{m} c^{\gcd(i,m)})/m
• Burnside lemma: |X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|
• 錯排公式: (n 個人中·每個人皆不再原來位置的組合數):
   dp[0] = 1; \dot{dp}[1] = 0;

dp[i] = (i-1)*(dp[i-1] + dp[i-2]);
• Bell 數 (有 n 個人, 把他們拆組的方法總數):

\overline{B_0} = 1

B_n = \sum_{k=0}^{n} s(n, k) \text{ (second - stirling)}

   B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k

    Wilson's theorem

   (p-1)! \equiv -1 \pmod{p}
• Fermat's little theorem :
   a^p \equiv a \pmod{p}
• Euler's totient function:
        \mod p = pow(A, pow(B, C, p-1)) mod \ p
• 歐拉函數降冪公式: A^B \mod C = A^B \mod \phi(c) + \phi(c) \mod C
• 用歐拉函數求模反元素:
   如果 a 和 n 互質,則 a 對 n 的模反元素 a^{-1} \equiv a^{\phi(n)-1} (mod \ n)
• 6 的倍數:  (a-1)^3 + (a+1)^3 + (-a)^3 + (-a)^3 = 6a 
• 上高斯 (向上取整):
   \lceil \frac{a}{b} \rceil = \frac{a+b-1}{b}
```

4 Geometry

4.1 definition

```
typedef long double ld;
const ld eps = 1e-8;
int dcmp(ld x) {
   if(abs(x) < eps) return 0;
   else return x < 0 ? -1 : 1;
}
struct Pt {
   ld x, y;
   Pt(ld _x=0, ld _y=0):x(_x), y(_y) {}
   Pt operator+(const Pt &a) const {
      return Pt(x+a.x, y+a.y); }
   Pt operator-(const Pt &a) const {
      return Pt(x-a.x, y-a.y); }
   Pt operator*(const ld &a) const {</pre>
```

```
return Pt(x*a, y*a); }
Pt operator/(const ld &a) const {
    return Pt(x/a, y/a);
  ld operator*(const Pt &a) const {
    return x*a.x + y*a.y;
  ld operator^(const Pt &a) const {
    return x*a.y - y*a.x;
  bool operator<(const Pt &a) const {</pre>
    return x < a.x | | (x == a.x && y < a.y); }
    //return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0 \&\&
         dcmp(y-a.y) < 0); }
  bool operator==(const Pt &a) const {
    return dcmp(x-a.x) == 0 && dcmp(y-a.y) == 0; }
ld norm2(const Pt &a) {
  return a*a; }
ld norm(const Pt &a) {
  return sqrt(norm2(a)); }
Pt perp(const Pt &a) {
return Pt(-a.y, a.x); }
Pt rotate(const Pt &a, ld ang) {
  return Pt(a.x*cos(ang)-a.y*sin(ang), a.x*sin(ang)+a.y
       *cos(ang)); }
struct Line {
  Pt s, e, v; // start, end, end-start
  ld ang;
  Line(Pt _s=Pt(0, 0), Pt _e=Pt(0, 0)):s(_s), e(_e) { v }
       = e-s; ang = atan2(v.y, v.x); }
  bool operator<(const Line &L) const {</pre>
    return ang < L.ang;</pre>
} };
struct Circle {
  Pt o; ld r;
  Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
```

4.2 Intersection of 2 lines

```
Pt LLIntersect(Line a, Line b) {
   Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
   ld f1 = (p2-p1)^(q1-p1),f2 = (p2-p1)^(p1-q2),f;
   if(dcmp(f=f1+f2) == 0)
     return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
   return q1*(f2/f) + q2*(f1/f);
}
```

4.3 halfPlaneIntersection

```
|// 0(nlogn)
// 傳入 vector<Line>
// (半平面為點 st 往 ed 的逆時針方向)
// 回傳值為形成的凸多邊形的頂點 vector
 // 對於點或線的解,將 '>' 改為 '>='
bool onleft(Line L, Pt p) { return dcmp(L.v ^ (p - L.s)
     ) > 0;  }
// 假設線段是有交點的
vector<Pt> HPI(vector<Line> &L) {
   sort(L.begin(), L.end()); // 按角度排序
   int n = L.size(), fir, las;
   Pt *p = new Pt[n];
   Line *q = new Line[n];
  q[fir = las = 0] = L[0];
for (int i = 1; i < n; i++) {
  while (fir < las && !onleft(L[i], p[las - 1])) las</pre>
     while (fir < las && !onleft(L[i], p[fir])) fir++;</pre>
     q[++las] = L[i];
     if (dcmp(q[las].v \land q[las - 1].v) == 0) {
       if (onleft(q[las], L[i].s)) q[las] = L[i];
     if (fir < las)</pre>
       p[las - 1] = LLIntersect(q[las - 1], q[las]);
   while (fir < las && !onleft(q[fir], p[las - 1])) las</pre>
```

```
if (las - fir <= 1) return {};</pre>
p[las] = LLIntersect(q[las], q[fir]);
int m = 0;
vector<Pt> ans(las - fir + 1);
for (int i = fir; i <= las; i++) ans[m++] = p[i];
return ans;
```

4.4 Convex Hull

```
double cross(Pt o, Pt a, Pt b){
 return (a-o) ^ (b-o);
vector<Pt> convex_hull(vector<Pt> pt){
  sort(pt.begin(),pt.end());
  int top=0;
 vector<Pt> stk(2*pt.size());
 for (int i=0; i<(int)pt.size(); i++){</pre>
   while (top >= 2 && cross(stk[top-2],stk[top-1],pt[i
       ]) <= 0) // 如果想要有點共線的點,把 <= 改成 <
     ton--:
   stk[top++] = pt[i];
 for (int i=pt.size()-2, t=top+1; i>=0; i--){
   while (top >= t && cross(stk[top-2],stk[top-1],pt[i
       ]) <= 0)
     top--;
   stk[top++] = pt[i];
 stk.resize(top-1);
 return stk;
```

Convex Hull trick 4.5

```
/* Given a convexhull, answer querys in O(lgN)
CH should not contain identical points, the area should
be > 0, min pair(x, y) should be listed first */
double det( const Pt& p1 , const Pt& p2 )
{ return p1.X * p2.Y - p1.Y * p2.X; }
struct Conv{
  int n;
  vector<Pt> a;
  vector<Pt> upper, lower;
  Conv(vector < Pt > _a) : a(_a){}
    n = a.size();
    int ptr = 0;
    for(int i=1; i<n; ++i) if (a[ptr] < a[i]) ptr = i;
for(int i=0; i<=ptr; ++i) lower.push_back(a[i]);
for(int i=ptr; i<n; ++i) upper.push_back(a[i]);</pre>
    upper.push_back(a[0]);
  int sign( LL x ){ // fixed when changed to double
  return x < 0 ? -1 : x > 0; }
  pair<LL,int> get_tang(vector<Pt> &conv, Pt vec){
    int l = 0, r = (int)conv.size() - 2;
    for(; l + 1 < r; ){
int mid = (l + r) / 2;
       if(sign(det(conv[mid+1]-conv[mid],vec))>0)r=mid;
       else l = mid;
    return max(make_pair(det(vec, conv[r]), r)
                 make_pair(det(vec, conv[0]), 0));
  void upd_tang(const Pt &p, int id, int &i0, int &i1){
    if(det(a[i0] - p, a[id] - p) > 0) i0 = id;
    if(det(a[i1] - p, a[id] - p) < 0) i1 = id;
  void bi_search(int l, int r, Pt p, int &i0, int &i1){
    if(l == r) return;
upd_tang(p, l % n, i0, i1);
    int sl=sign(det(a[l % n] - p, a[(l + 1) % n] - p));
    for(; l + 1 < r; ) {
       int mid = (l + r) / 2;
       int smid=sign(det(a[mid%n]-p, a[(mid+1)%n]-p));
       if (smid == sl) l = mid;
       else r = mid;
```

```
upd_tang(p, r % n, i0, i1);
int bi_search(Pt u, Pt v, int l, int r) {
  int sl = sign(det(v - u, a[l % n] - u));
  for( ; l + 1 < r; )</pre>
    int mid = (l + r) / 2;
    int smid = sign(det(v - u, a[mid % n] - u));
    if (smid == sl) l = mid;
    else r = mid:
  return 1 % n;
}
^{\prime}// 1. whether a given point is inside the CH
bool contain(Pt p) {
  if (p.X < lower[0].X || p.X > lower.back().X)
       return 0;
  int id = lower_bound(lower.begin(), lower.end(), Pt
       (p.X, -INF)) - lower.begin();
  if (lower[id].X == p.X) {
  if (lower[id].Y > p.Y) return 0;
}else if(det(lower[id-1]-p,lower[id]-p)<0)return 0;</pre>
  id = lower_bound(upper.begin(), upper.end(), Pt(p.X
  , INF), greater<Pt>()) - upper.begin();
if (upper[id].X == p.X) {
    if (upper[id].Y < p.Y) return 0;</pre>
  }else if(det(upper[id-1]-p,upper[id]-p)<0)return 0;</pre>
  return 1;
// 2. Find 2 tang pts on CH of a given outside point
// return true with i0, i1 as index of tangent points
// return false if inside CH
bool get_tang(Pt p, int &i0, int &i1) {
  if (contain(p)) return false;
  i0 = i1 = 0;
  int id = lower_bound(lower.begin(), lower.end(), p)
        lower.begin();
  bi_search(0, id, p, i0, i1);
bi_search(id, (int)lower.size(), p, i0, i1);
  id = lower_bound(upper.begin(), upper.end(), p,
      greater<Pt>()) - upper.begin();
  bi_search((int)lower.size() - 1, (int)lower.size()
      -1 + id, p, i0, i1);
  bi_search((int)lower.size() - 1 + id, (int)lower.
      size() - 1 + (int)upper.size(), p, i0, i1);
  return true;
// 3. Find tangent points of a given vector
// ret the idx of vertex has max cross value with vec
int get_tang(Pt vec){
  pair<LL, int> ret = get_tang(upper, vec);
  ret.second = (ret.second+(int)lower.size()-1)%n;
  ret = max(ret, get_tang(lower, vec));
  return ret.second;
// 4. Find intersection point of a given line
// return 1 and intersection is on edge (i, next(i))
// return 0 if no strictly intersection
bool get_intersection(Pt u, Pt v, int &i0, int &i1){
 int p0 = get_tang(u - v), p1 = get_tang(v - u);
if(sign(det(v-u,a[p0]-u))*sign(det(v-u,a[p1]-u))<0){</pre>
   if (p0 > p1) swap(p0, p1);
   i0 = bi_search(u, v, p0, p1);
   i1 = bi_search(u, v, p1, p0 + n);
   return 1;
 }
 return 0;
};
```

4.6 Intersection of 2 segments

```
int ori( const Pt& o , const Pt& a , const Pt& b ){
  LL ret = ( a - o ) ^ ( b - o );
   return (ret > 0) - (ret < 0);</pre>
// p1 == p2 || q1 == q2 need to be handled
bool banana( const Pt& p1 , const Pt& p2
   const Pt& q1 , const Pt& q2 ){

if( ( p2 - p1 ) ^ ( q2 - q1 ) ) == 0 ){ // parallel
      if( ori( p1 , p2 , q1 ) ) return false;
return ( ( p1 - q1 ) * ( p2 - q1 ) ) <= 0 ||
```

4.7 Point In Polygon

4.8 Tangent line of two circles

```
vector<Line> go( const Cir& c1 , const Cir& c2 , int
     sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_{sq} = norm2(c1.0 - c2.0);
  if( d_sq < eps ) return ret;
double d = sqrt( d_sq );</pre>
  Pt v = (c2.0 - c1.0) / d;
  double c = (c1.R - sign1 * c2.R) / d;
  if( c * c > 1 ) return ret;
double h = sqrt( max( 0.0 , 1.0 - c * c ) );
  for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
                v.Y * c + sign2 * h * v.X $;
     Pt p1 = c1.0 + n * c1.R;
     Pt p2 = c2.0 + n * (c2.R * sign1);
     if( fabs( p1.X - p2.X ) < eps and
       fabs( p1.Y - p2.Y ) < eps )
p2 = p1 + perp( c2.0 - c1.0 );
     ret.push_back( { p1 , p2 } );
  return ret;
}
```

4.9 Minimum distance of two convex

4.10 Area of Rectangles

```
struct AreaofRectangles{
#define cl(x) (x<<1)
#define cr(x) (x<<1|1)
     ll n, id, sid;
     pair<ll,ll> tree[MXN<<3]; // count, area</pre>
     vector<ll> ind;
     tuple<ll, ll, ll, ll> scan[MXN<<1];</pre>
     void puli(int i, int l, int r){
   if(tree[i].first) tree[i].second = ind[r+1] -
              ind[l];
          else if(l != r){
              int mid = (l+r)>>1;
              tree[i].second = tree[cl(i)].second + tree[
                   cr(i)].second;
          else
                   tree[i].second = 0;
     void upd(int i, int l, int r, int ql, int qr, int v
          if(ql <= l \&\& r <= qr){}
              tree[i].first += v;
              pull(i, l, r); return;
          int mid = (l+r) >> 1;
          if(ql <= mid) upd(cl(i), l, mid, ql, qr, v);</pre>
         if(qr > mid) upd(cr(i), mid+1, r, ql, qr, v);
pull(i, l, r);
     void init(int _n){
         n = _n; id = sid = 0;
ind.clear(); ind.resize(n<<1);</pre>
          fill(tree, tree+(n<<2), make_pair(0, 0));
     void addRectangle(int lx, int ly, int rx, int ry){
          ind[id++] = lx; ind[id++] = rx;
          scan[sid++] = make\_tuple(ly, 1, lx, rx);
          scan[sid++] = make\_tuple(ry, -1, lx, rx);
     ll solve(){
          sort(ind.begin(), ind.end());
          ind.resize(unique(ind.begin(), ind.end()) - ind
               .begin());
         sort(scan, scan + sid);
ll area = 0, pre = get<0>(scan[0]);
          for(int i = 0; i < sid; i++){
              auto [x, v, l, r] = scan[i];
area += tree[1].second * (x-pre);
              upd(1, 0, ind.size()-1, lower_bound(ind.
begin(), ind.end(), l)-ind.begin(),
                   lower_bound(ind.begin(),ind.end(),r)-
                   ind.begin()-1, v);
              pre = x;
          return area;
     }rect;
}
```

4.11 Min dist on Cuboid

```
r=INF; turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H); return r; }
```

4.12 Heart of Triangle

```
Pt inCenter( Pt &A, Pt &B, Pt &C) { // 内心 double a = norm(B-C), b = norm(C-A), c = norm(A-B); return (A * a + B * b + C * c) / (a + b + c); }

Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心 Pt bb = b - a, cc = c - a; double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc); return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d; }

Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心 Pt ba = b - a, ca = c - a, bc = b - c; double Y = ba.Y * ca.Y * bc.Y, A = ca.X * ba.Y - ba.X * ca.Y, x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A, y0= -ba.X * (x0 - c.X) / ba.Y + ca.Y; return Pt(x0, y0); }
```

5 Graph

5.1 DSU 並查集 & MST

```
struct DSU {// 並查集
    vector<int> fa, sz;
DSU(int n = 0) : fa(n), sz(n, 1) {
         iota(fa.begin(), fa.end(), 0);
     int Find(int x) { // 路徑壓縮
while (x != fa[x])
             x = fa[x] = fa[fa[x]];
         return x;
    bool Merge(int x, int y) { //合併
         x = Find(x), y = Find(y);
if (x == y) return false; // 是否為連通
         if (sz[x] > sz[y]) swap(x, y);
         fa[x] = y;
         sz[y] += sz[x];
         return true;
int MST(int n, int m, vector<tuple<int, int, int>> &edge
     ){ //0 base
     sort(edge.begin(), edge.end());
     DSU dsu(n+1); // 初始化並查集
     int res = 0, flag=1; // 最小生成樹邊權和
for (auto &[w, u, v] : edge)
         if(dsu.Merge(u, v)) {
              res += w; //合併並統計答案
//graph[u].push_back({v,w});
              //graph[v].push_back({u,w});
          //else edges.push_back({w,u,v});
    return res;
int main(){
    int n, m; //點數,邊數
     cin >> n >> m;
    vector<tuple<int, int, int>> edge(m);
for (auto &[w, u, v] : edge) cin >> u >> v >> w;
     cout << MST(n, m, edge);</pre>
}
```

5.2 Lowest Common Ancestor O(lgn)

```
struct LCA {
  int n, ti, lgN;
  int anc[MXN + 5][__lg(MXN) + 1] = {0};
```

```
int MaxLength[MXN][\_lg(MXN) + 1] = {0};
   int time_in[MXN] = {0};
   int time_out[MXN] = \{0\};
  LCA(int _n, int f):n(_n), ti(0), lgN(__lg(n)) {
  dfs(f, f, 0);
     build();
   void dfs(int now, int f, int len_to_father) { // dfs
        for anc, time, Lenth
     ti++;
     anc[now][0] = f;
     time_in[now] = ti;
     MaxLength[now][0] = len_to_father;
     for (auto i : graph[now]) {
         if (i.first == f) continue
         dfs(i.first, now, i.second);
     time_out[now] = ti;
   void build() { // build anc[][], MaxLength[][]
     for (int i = 1; i <= lgN; ++i) {
       for (int u = 1; u <= n; ++u) {
         anc[u][i] = anc[anc[u][i - 1]][i - 1];
         MaxLength[u][i] = max(MaxLength[u][i - 1]
                    MaxLength[anc[u][i - 1]][i - 1]);
         // dis[u][i] += dis[anc[u][i - 1]][i - 1]
// + dis[u][i - 1];
    }
   bool isAncestor(int x, int y) {
     return time_in[x] <= time_in[y] && time_out[x] >=
         time_out[y];
   int getLCA(int u, int v) {
     if (isAncestor(u, v)) return u;
if (isAncestor(v, u)) return v;
     for (int i = lgN; i >= 0; --i) {
       if (!isAncestor(anc[u][i], v)) {
         u = anc[u][i];
       }
     return anc[u][0];
   int getMAX(int u, int v) { //獲得路徑上最大邊權
     int lca = getLCA(u, v);
     int maxx = -1;
     for (int i = lgN; i >= 0; --i) {
       // u to lca
       if (!isAncestor(anc[u][i], lca)) {
         maxx = max(maxx, MaxLength[u][i]);
         u = anc[u][i];
       // v to lca
       if (!isAncestor(anc[v][i], lca))
         maxx = max(maxx, MaxLength[v][i]);
         v = anc[v][i];
       }
     if (u != lca) maxx = max(maxx, MaxLength[u][0]);
if (v != lca) maxx = max(maxx, MaxLength[v][0]);
     return maxx;
};
```

5.3 Hamiltonian path $O(n^22^n)$

```
//dp[i][j] = 目前在j節點走過{i}節點的最短路徑
for(int i=1; i < (1 << n); i++ ) {
    for(int j = 1; j < n; j++ ) {
        if(!(1 << j) & i)&&(i&1)) {
            for( int k = 0; k < n; k++ ) {
                if(j == k) continue;
                if( (1<<k)&i ) dp[j][i|(1<<j)]=
                      min(dp[j][i|(1<<j)],dp[k][i]+dis[k][j]);
        }
    }
    }
}
```

5.4 MaximumClique 最大團

```
#define N 111
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int linkto[N] , v[N];
  int n:
  void init(int _n){
     n = _n;
     for(int i = 0 ; i < n ; i ++){
       linkto[i].reset(); v[i].reset();
  void addEdge(int a , int b)
{ v[a][b] = v[b][a] = 1; }
int popcount(const Int& val)
  { return val.count(); }
  int lowbit(const Int& val)
  { return val._Find_first(); }
  int ans , stk[N];
int id[N] , di[N] , deg[N];
  Int cans;
  void maxclique(int elem_num, Int candi){
     if(elem_num > ans){
       ans = elem_num; cans.reset();
       for(int i = 0 ; i < elem_num ; i ++)</pre>
          cans[id[stk[i]]] = 1;
     int potential = elem_num + popcount(candi);
     if(potential <= ans) return;</pre>
     int pivot = lowbit(candi);
Int smaller_candi = candi & (~linkto[pivot]);
     while(smaller_candi.count() && potential > ans){
       int next = lowbit(smaller_candi);
       candi[next] = !candi[next];
       smaller_candi[next] = !smaller_candi[next];
       potential --:
       if(next == pivot || (smaller_candi & linkto[next
            ]).count()){
          stk[elem_num] = next;
          maxclique(elem_num + 1, candi & linkto[next]);
  } } }
  int solve(){
    for(int i = 0 ; i < n ; i ++){
  id[i] = i; deg[i] = v[i].count();</pre>
     sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
     for(int i = 0; i < n; i ++) di[id[i]] = i;</pre>
     for(int i = 0; i < n; i ++)
for(int j = 0; j < n; j +-
          if(v[i][j]) linkto[di[i]][di[j]] = 1;
     Int cand; cand.reset();
for(int i = 0 ; i < n ; i ++) cand[i] = 1;</pre>
     ans = 1;
     cans.reset(); cans[0] = 1;
     maxclique(0, cand);
     return ans;
} }solver;
```

5.5 MaximalClique 極大團

```
#define N 80
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int lnk[N] , v[N];
  int n;
  void init(int _n){
    n = _n;
    for(int i = 0 ; i < n ; i ++){
        lnk[i].reset(); v[i].reset();
    }}
  void addEdge(int a , int b)
  { v[a][b] = v[b][a] = 1; }
  int ans , stk[N], id[N] , di[N] , deg[N];
  Int cans;
  void dfs(int elem_num, Int candi, Int ex){
    if(candi.none()&&ex.none()){
        cans.reset();
        for(int i = 0 ; i < elem_num ; i ++)</pre>
```

```
cans[id[stk[i]]] = 1;
ans = elem_num; // cans is a maximal clique
    int pivot = (candilex)._Find_first();
    Int smaller_candi = candi & (~lnk[pivot]);
    while(smaller_candi.count()){
       int nxt = smaller_candi._Find_first();
       candi[nxt] = smaller_candi[nxt] = 0;
       ex[nxt] = 1;
       stk[elem_num] = nxt;
       dfs(elem_num+1,candi&lnk[nxt],ex&lnk[nxt]);
  int solve(){
    for(int i = 0; i < n; i ++){
       id[i] = i; deg[i] = v[i].count();
    sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
    for(int i = 0; i < n; i ++) di[id[i]] = i;
for(int i = 0; i < n; i ++)</pre>
       for(int j = 0; j < n; j ++)
         if(v[i][j]) ink[di[i]][di[j]] = 1;
    ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n,'1')), 0);
     return ans;
} }solver;
```

5.6 BCC based on vertex 點雙聯通分量

```
#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
  int n,nScc,step,dfn[MXN],low[MXN];
vector<int> E[MXN],sccv[MXN];
  int top,stk[MXN];
  void init(int _n) { // 初始化n點
     n = _n; nScc = step = 0;
for (int i=0; i<n; i++) E[i].clear();</pre>
  void addEdge(int u, int v) // 無向邊
  { E[u].PB(v); E[v].PB(u); } void DFS(int u, int f) {
     dfn[u] = low[u] = step++;
     stk[top++] = u;
     for (auto v:E[u]) {
       if (v == f) continue;
if (dfn[v] == -1) {
         DFS(v,u);
          low[u] = min(low[u], low[v]);
          if (low[v] >= dfn[u]) {
            int z;
            sccv[nScc].clear();
            do {
              z = stk[--top];
              sccv[nScc].PB(z);
            } while (z != v);
            sccv[nScc++].PB(u);
       }else
         low[u] = min(low[u],dfn[v]);
  vector<vector<int>> solve() { // 回傳(size=2 橋, size
       >2 點雙連通分量)
     vector<vector<int>> res;
     for (int i=0; i<n; i++)</pre>
     dfn[i] = low[i] = -1;
for (int i=0; i<n; i++)</pre>
       if (dfn[i] == -1) {
         top = 0;
         DFS(i,i);
     REP(i,nScc) res.PB(sccv[i]);
     return res;
}graph;
```

Strongly Connected Component 強連通分 5.9 Min Mean Cycle

```
#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x)) //fill zero
struct Scc{
  int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
  void init(int _n){
    n = _n;
for (int i=0; i<MXN; i++)</pre>
      E[i].clear(), rE[i].clear();
  void addEdge(int u, int v){
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u){
    vst[u]=1;
    for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u){
    vst[u] = 1; bln[u] = nScc;
    for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve(){
    nScc = 0;
    vec.clear();
    FZ(vst);
    for (int i=0; i<n; i++)</pre>
      if (!vst[i]) DFS(i);
    reverse(vec.begin(),vec.end());
    FZ(vst);
    for (auto v : vec)
      if (!vst[v]){
        rDFS(v); nScc++;
  }
};
```

Maximum General graph Matching

```
// should shuffle vertices and edges
const int N=100005, E=(2e5)*2+40;
struct Graph{ // 1-based; match: i <-> lnk[i]
  int to[E],bro[E],head[N],e,lnk[N],vis[N],stp,n;
  void init(int _n){
    stp=0; e=1; n=_n;
for(int i=1;i<=n;i++) head[i]=lnk[i]=vis[i]=0;</pre>
  void add_edge(int u,int v){
  to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
       int v=to[i];
       if(!lnk[v]){ lnk[x]=v,lnk[v]=x; return true; }
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
       if(vis[lnk[v]]<stp){</pre>
         int w=lnk[v]; lnk[x]=v,lnk[v]=x,lnk[w]=0;
if(dfs(w)) return true;
         lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans=0;
    for(int i=1;i<=n;i++) if(!lnk[i]) stp++,ans+=dfs(i)</pre>
    return ans;
}graph;
```

```
/* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
   struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
   Edge e[E];
   vector<int> edgeID, cycle, rho;
   double d[V][V];
   void init( int _n )
   \{ n = _n; m = 0; \}
   // WARNING: TYPE matters
  void addEdge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);
  for(int i=0; i=1); inf</pre>
       for(int j=0; j<m; j++) {
  int v = e[j].v, u = e[j].u;
  if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
            d[i+1][u] = d[i][v]+e[j].c;
            prv[i+1][u] = v;
            prve[i+1][u] = j;
   double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1;
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
        for(int k=0; k<n; k++) {</pre>
          if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i</pre>
               ])/(n-k)):
          else avg=max(avg,inf);
       if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
     fill(vst,0); edgeID.clear(); cycle.clear(); rho.
          clear();
     for (int i=n; !vst[st]; st=prv[i--][st]) {
       vst[st]++;
        edgeID.PB(prve[i][st]);
       rho.PB(st);
     while (vst[st] != 2) {
       if(rho.empty()) return inf;
        int v = rho.back(); rho.pop_back();
       cycle.PB(v);
       vst[v]++;
     reverse(ALL(edgeID));
     edgeID.resize(SZ(cycle));
     return mmc;
} }mmc;
```

5.10 Directed Graph Min Cost Cycle

```
// works in O(N M)
#define INF 1000000000000000LL
#define N 5010
#define M 200010
struct edge{
  int to; LL w;
  edge(int a=0, LL b=0): to(a), w(b){}
};
struct node{
  LL d; int u, next;
  node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
}b[M];
struct DirectedGraphMinCycle{
  vector<edge> g[N], grev[N];
  LL dp[N][N], p[N], d[N], mu;
  bool inq[N];
```

```
int n, bn, bsz, hd[N];
void b_insert(LL d, int u){
  int i = d/mu;
  if(i >= bn) return;
  b[++bsz] = node(d, u, hd[i]);
  hd[i] = bsz;
void init( int _n ){
  n = _n;
  for( int i = 1; i <= n; i ++)
     g[ i ].clear();
void addEdge( int ai , int bi , LL ci )
{ g[ai].push_back(edge(bi,ci)); }
LL solve(){
  fill(dp[0], dp[0]+n+1, 0);
  dp[i][g[j][k].to] =min(dp[i][g[j][k].to]
                                      dp[i-1][j]+g[j][k].w);
  mu=INF; LL bunbo=1;
  for(int i=1; i<=n; i++) if(dp[n][i] < INF){</pre>
     LL a=-INF, b=1;
     for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
  if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre>
          a = dp[n][i]-dp[j][i];
          b = n-j;
     if(mu*b > bunbo*a)
       mu = a, bunbo = b;
  if(mu < 0) return -1; // negative cycle
  if(mu == INF) return INF; // no cycle
   if(mu == 0) return 0;
  for(int i=1; i<=n; i++)</pre>
     for(int j=0; j<(int)g[i].size(); j++)
g[i][j].w *= bunbo;</pre>
  memset(p, 0, sizeof(p));
  queue<int> q;
  for(int i=1; i<=n; i++){</pre>
     q.push(i);
     inq[i] = true;
  while(!q.empty()){
     int i=q.front(); q.pop(); inq[i]=false;
     for(int j=0; j<(int)g[i].size(); j++){
  if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
          p[g[i][j].to] = p[i]+g[i][j].w-mu;
          if(!inq[g[i][j].to]){
            q.push(g[i][j].to);
            inq[g[i][j].to] = true;
  } } }
for(int i=1; i<=n; i++) grev[i].clear();</pre>
  for(int i=1; i<=n; i++)</pre>
    for(int j=0; j<(int)g[i].size(); j++){
  g[i][j].w += p[i]-p[g[i][j].to];
  grev[g[i][j].to].push_back(edge(i, g[i][j].w));</pre>
  LL \ mldc = n*mu;
  for(int i=1; i<=n; i++){</pre>
    bn=mldc/mu, bsz=0;
memset(hd, 0, sizeof(hd));
fill(d+i+1, d+n+1, INF);
     b_insert(d[i]=0, i);
     for(int j=\bar{0}; j<=bn-1; j++) for(int k=hd[j]; k; k=b[k].next){
       int u = b[k].u;
       LL du = b[k].d;
       if(du > d[u]) continue;
       for(int l=0; l<(int)g[u].size(); l++) if(g[u][l</pre>
         ].to > i){
    if(d[g[u][l].to] > du + g[u][l].w){
        d[g[u][l].to] = du + g[u][l].w;
            b_insert(d[g[u][l].to], g[u][l].to);
     } } }
     for(int j=0; j<(int)grev[i].size(); j++) if(grev[</pre>
          i][j].to > i)
       mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
  }
```

```
5.11 K-th Shortest Path
```

return mldc / bunbo;

} }graph;

```
// time: 0(|E| \lg |E| + |V| \lg |V| + K)
// memory: 0(|E| \lg |E| + |V|)
struct KSP{ // 1-base
  struct nd{
     int u, v; ll d;
     nd(int ui = 0, int vi = 0, ll di = INF)
{ u = ui; v = vi; d = di; }
  };
  struct heap{
     nd* edge; int dep; heap* chd[4];
  static int cmp(heap* a,heap* b)
   { return a->edge->d > b->edge->d; }
   struct node{
     int v; ll d; heap* H; nd* E;
     node(){}
     node(ll _d, int _v, nd* _E)
{ d =_d; v = _v; E = _E; }
node(heap* _H, ll _d)
     {H = _H; d = _d; }
     friend bool operator<(node a, node b)
     { return a.d > b.d; }
  int n, k, s, t;
  11 dst[ N ];
  nd *nxt[ N ];
  vector<nd*> g[ N ], rg[ N ];
  heap *nullNd, *head[ N ];
  void init( int _n , int _k , int _s , int _t ){
     n = _n; k = _k; s = _s; t = _t;
for( int i = 1 ; i <= n ; i ++ ){
    g[ i ].clear(); rg[ i ].clear();
    nxt[ i ] = NULL; head[ i ] = NULL;
    dst[ i ] = -1;
}</pre>
  void addEdge( int ui , int vi , ll di ){
  nd* e = new nd(ui, vi, di);
     g[ ui ].push_back( e );
     rg[ vi ].push_back( e );
  queue<int> dfsQ;
  void dijkstra(){
     while(dfsQ.size()) dfsQ.pop();
     priority_queue<node> Q;
     Q.push(node(0, t, NULL));
     while (!Q.empty()){
        node p = Q.top(); Q.pop();
if(dst[p.v] != -1) continue;
        dst[ p.v ] = p.d;
nxt[ p.v ] = p.E;
        dfsQ.push( p.v );
        for(auto e: rg[ p.v ])
          Q.push(node(p.d + e->d, e->u, e));
  heap* merge(heap* curNd, heap* newNd){
   if(curNd == nullNd) return newNd;
     heap* root = new heap;
     memcpy(root, curNd, sizeof(heap));
if(newNd->edge->d < curNd->edge->d){
        root->edge = newNd->edge;
        root->chd[2] = newNd->chd[2];
root->chd[3] = newNd->chd[3];
        newNd->edge = curNd->edge;
        newNd - chd[2] = curNd - chd[2];
        newNd - > chd[3] = curNd - > chd[3];
     if(root->chd[0]->dep < root->chd[1]->dep)
        root->chd[0] = merge(root->chd[0],newNd);
        root->chd[1] = merge(root->chd[1],newNd);
     root->dep = max(root->chd[0]->dep, root->chd[1]->
          dep) + 1;
     return root;
  vector<heap*> V;
```

```
void build(){
     nullNd = new heap;
     nullNd->dep = 0;
     nullNd->edge = new nd;
     fill(nullNd->chd, nullNd->chd+4, nullNd);
     while(not dfsQ.empty()){
       int u = dfsQ.front(); dfsQ.pop();
if(!nxt[ u ]) head[ u ] = nullNd;
       else head[ u ] = head[nxt[ u ]->v];
       V.clear();
        for( auto&& e : g[ u ] ){
          int v = e \rightarrow v;
         if( dst[ v ] == -1 ) continue;
e->d += dst[ v ] - dst[ u ];
if( nxt[ u ] != e ){
            heap* p = new heap;
            fill(p->chd, p->chd+4, nullNd);
            p->dep = 1;
            p->edge = e;
            V.push_back(p);
       if(V.empty()) continue;
make_heap(V.begin(), V.end(), cmp);
#define L(X) ((X<<1)+1)</pre>
#define R(X) ((X<<1)+2)
       for( size_t i = 0 ; i < V.size() ; i ++ ){
  if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
          else V[i]->chd[2]=nullNd;
          if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
          else V[i]->chd[3]=nullNd;
       head[u] = merge(head[u], V.front());
  } }
  vector<ll> ans
  void first_K(){
     ans.clear();
     priority_queue<node> Q;
     if( dst[ s ] == -1 ) return;
     ans.push_back( dst['s ] );
if( head[s] != nullNd )
     Q.push(node(head[s], dst[s]+head[s]->edge->d));
for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();</pre>
       ans.push_back( p.d );
       if(head[ p.H->edge->v ] != nullNd){
          q.H = head[ p.H->edge->v ];
          q.d = p.d + q.H->edge->d;
          Q.push(q);
       for( int i = 0; i < 4; i ++ )
          if( p.H->chd[ i ] != nullNd ){
            q.H = p.H->chd[i];
            q.d = p.d - p.H->edge->d + p.H->chd[i]->
                 edge->d;
            Q.push( q );
  } }
  void solve(){ // ans[i] stores the i-th shortest path
     dijkstra();
     build();
     first_K(); // ans.size() might less than k
} }solver;
```

5.12 Floryd Warshall

5.13 SPFA

```
#define MXN 200005
struct SPFA{
  int n;
  LL inq[MXN], len[MXN];
  vector < LL> dis;
  vector<pair<int, LL>> edge[MXN];
  void init(int _n){
    n = _n;
    dis.clear(); dis.resize(n, 1e18);
    for(int i = 0; i < n; i++){
       edge[i].clear();
       inq[i] = len[i] = 0;
  } }
  void addEdge(int u, int v, LL w){
    edge[u].push_back({v, w});
  vector<LL> solve(int st = 0){
    deque<int> dq; //return {-1} if has negative cycle
dq.push_back(st); //otherwise return dis from st
    inq[st] = 1; dis[st] = 0;
    while(!dq.empty()){
       int u = dq.front(); dq.pop_front();
       inq[u] = 0;
       for(auto [to, d] : edge[u]){
         if(dis[to] > d+dis[u]){
           dis[to] = d+dis[u];
len[to] = len[u]+1;
           if(len[to] > n) return {-1};
           if(inq[to]) continue;
           (!dq.empty()&&dis[dq.front()] > dis[to]?
                dq.push_front(to) : dq.push_back(to));
           inq[to] = 1;
    } } }
    return dis;
} }spfa;
```

5.14 Tree Hash

```
//限定root = 1
//從 dfs(1,1) 開始
int subtree_sz[MXN];
vector<int> edge[MXN];
int dfs(int u, int f) {
  vector<pair<int, int>> h;
  subtree\_sz[u] = 1;
  for (int child : edge[u]) {
    if (child == f) continue;
    int tmp = dfs(child, u);
    h.push_back(make_pair(tmp, subtree_sz[child]));
    subtree_sz[u] += subtree_sz[child];
  sort(h.begin(), h.end());
  int ret = subtree_sz[u];
  for (auto v : h) {
  ret = ((ret * p) % MOD + v.first) % MOD;
    ret = ret * v.second % MOD;
  return ret;
}
```

5.15 HeavyLightDecomposition

```
// 詢問,修改複雜度 O(log^2 n)
// 1-base

int sz[MXN], dep[MXN], son[MXN], fa[MXN];

// 第一次 dfs
// 找重兒子 需要紀錄當前節點的子樹大小(sz)、深度(dep)、
重兒子(son)、父節點(fa)
// 沒有子節點 son[x] = 0
void dfs_sz(int x, int f, int d) { //當前節點 x · 父節點 f·深度 d
sz[x] = 1; dep[x] = d; fa[x] = f;
```

```
for(int i : edge[x]) {
        if(i == f)
                     continue;
       dfs_sz(i, x, d+
sz[x] += sz[i];
                    d+1);
        if(sz[son[x]] < sz[i])</pre>
                                 son[x] = i;
    }
}
// 第二次 dfs
int top[MXN]; // 每個節點所在的鏈的頂端節點
int dfn[MXN]; // 節點編號,編號為在線段樹上的位置
int rnk[MXN]; // 編號為哪個節點
int bottom[MXN]; // 維護每個節點的子樹中最大 dfn 編號
int cnt = 0;
int dfs_hld(int x, int f){
   top[x] = (son[fa[x]] == x ? top[fa[x]] : x);
    rnk[cnt] = x
    bottom[\bar{x}] = dfn[x] = cnt++;
        if(son[x])
    for(int i : edge[x]){
        if(i == f || i == son[x])
                                    continue;
        bottom[x] = max(bottom[x], dfs_hld(i, x)); //
            更新子樹最大編號
    return bottom[x];
}
// 求出 lca
// 不斷跳鏈·直到 u,v 跳到同一條鏈上為止
// 每次跳鏈選所在的鏈頂端深度較深的一端往上跳
int getLca(int u, int v) {
   while(top[u] != top[v]){
      if(dep[top[u]] > dep[top[v]])
         u = fa[top[u]];
      else
          v = fa[top[v]];
    return dep[u] > dep[v] ? v : u;
}
// 路徑權重總和
int query(int u, int v) {
    int ret = 0;
    while(top[u] != top[v]){
        if (dep[top[u]] > dep[top[v]]){
            ret += segtree.query(dfn[top[u]], dfn[u]);
            u = fa[top[u]];
        }
        else{
            ret += segtree.query(dfn[top[v]], dfn[v]);
            v = fa[top[v]];
    }
    // 最後到同一條鏈上
    ret += segtree.query(min(dfn[u], dfn[v]), max(dfn[u
        ], dfn[v]));
    return ret;
}
```

5.16 差分約束

約束條件 $V_j - V_i \leq W$ addEdge(V_i, V_j, W) and run bellman-ford or spfa

6 String

6.1 PalTree O(n)

```
// state[i]代表第i個字元為結尾的最長回文編號
// len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文子字串在整個字串中的出現次數
// fail[s]是他長度次長的回文後綴·aba的fail是a
const int MXN = 1000010;
struct PalT{
  int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
```

```
int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN] = \{-1\};
  int newNode(int l,int f){
    len[tot]=l,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
    diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++;
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int v=lst;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0;
    newNode(0,1), newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

6.2 Longest Increasing Subsequence

```
vector<int> getLIS(vector<int> a){
  vector<int> lis;
  for(int i : a){
    if(lis.empty() || lis.back() < i) lis.push_Back(
        i);
    else *lower_bound(lis.begin(), lis.end(), i) =
        i;
  }
  return lis;
}</pre>
```

6.3 Longest Common Subsequence O(nlgn)

6.4 KMP

```
/* len-failure[k]:
在k結尾的情況下‧這個子字串可以由開頭
長度為(len-failure[k])的部分重複出現來表達
failure[k] 為次長相同前綴後綴
如果我們不只想求最多,而且以0-base做為考量
 ・那可能的長度由大到小會是
failuer[k] · failure[failuer[k]-1]
 failure[failure[failuer[k]-1]-1]..
直到有值為0為止 *.
int failure[MXN];
vector<int> KMP(string& t, string& p) {
    vector<int> ret;
    if(p.size() > t.size()) return ret;
    for(int i = 1, j = failure[0] = -1; i < p.size(); i
       while(j \ge 0 \& p[j + 1] != p[i]) j = failure[j]
       if(p[j + 1] == p[i]) j++;
       failure[i] = j;
    for(int i = 0, j = -1; i < t.size(); i++) {
       while (j \ge 0 \& p[j + 1] != t[i]) j = failure[
       if(p[j + 1] == t[i]) j++;
       if(j == p.size() - 1) {
           ret.push_back(i - p.size() + 1);
           j = failure[j];
       }
    return ret;
}
```

6.5 SAIS O(n)

```
/*** SA· 將字串的所有後綴排序後的數組 ***/
/* SA[i]儲存排序後第i小的後綴從哪裡開始 */
/**** H[i] 為第i小的字串跟第i-1小的LCP ***/
/**** 註: LCP(Longest Common Prefix) ****/
/*** ex:S = "babd", SA[0] = 1("abd") ****/
/** SA[1] = 0("babd"), SA[2] = 2("bd") **/
/** H[0] = 0, H[1] = 0, H[2] = 1("b") ***/
/* 傳入參數:ip 陣列放字串·len為字串長度 */
/* 需保證ip[len]為0, 且字串裡的元素不為0 */
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
  hei[N], r[N];
int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
  memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
    mkhei(n);
  void mkhei(int n){
    REP(i,n) r[\_sa[i]] = i;
    hei[0] = 0;
    REP(i,n) if(r[i]) {
       int ans = i>0? max(hei[r[i-1]] - 1, 0) : 0;
       hei[r[i]] = ans;
  void sais(int *s, int *sa, int *p, int *q, bool *t,
       int *c, int n, int z){
    bool uniq = t[n-1] = true, neq;
    int n = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
         lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
    memcpy(x, c, sizeof(int) * z); \
    memcpy(x + 1, c, sizeof(int) * (z - 1)); \
    REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i
    ]-1]]++] = sa[i]-1; \
memcpy(x, c, sizeof(int) * z); \
```

```
for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i
                            ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
             MSO(c, z);
              REP(i,n) uniq \&= ++c[s[i]] < 2;
              REP(i,z-1) c[i+1] += c[i];
              if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
             for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i
+1] ? t[i+1] : s[i]<s[i+1]);
             MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i
             ]]]=p[q[i]=nn++]=i);
REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
                    \label{lem:neq} \begin{tabular}{ll} neq=lst<0 | lmemcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa(i))) & lem: lmemcmp(s+sa[i],s+lst) &
                                   [i])*sizeof(int));
                    ns[q[lst=sa[i]]]=nmxz+=neq;
              sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
             + 1);
MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[s[p[
                           nsa[i]]]] = p[nsa[i]]);
      }
}sa;
 int H[ N ], SA[ N ];
 void suffix_array(int* ip, int len) {
       // should padding a zero in the back
       // ip is int array, len is array length
       // ip[0..n-1] != 0, and ip[len] = 0
       ip[len++] = 0;
       sa.build(ip, len, 128);
       for (int i=0; i<len; i++) {
             H[i] = sa.hei[i + 1];
              SA[i] = sa.\_sa[i + 1];
        // resulting height, sa array \in [0,len)
}
```

6.6 Z Value O(n)

```
//z[i] = lcp(s[1...n-1],s[i...n-1])
int z[MAXN];
void Z_value(const string& s) {
   int i, j, left, right, len = s.size();
   left=right=0; z[0]=len;
   for(i=1;i<len;i++) {
      j=max(min(z[i-left],right-i),0);
      for(;i+j<len&&s[i+j]==s[j];j++);
      z[i]=j;
      if(i+z[i]>right) {
        right=i+z[i];
        left=i;
   } }
}
```

6.7 Manacher Algorithm O(n)

```
|// 求以每個字元為中心的最長回文半徑
// 頭尾以及每個字元間都加入一個
// 沒出現過的字元‧這邊以'@'為例
// s為傳入的字串·len為字串長度
// z為儲存答案的陣列 (有包含'@'要小心)
                   // ex: s = "abaac" ->
// z =
                   [12141232121]
void z_value_pal(char *s,int len,int *z){
  len=(len<<1)+1
  for(int i=len-1;i>=0;i--)
    s[i]=i&1?s[i>>1]:'@';
  z[0]=1;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i < r?min(z[l+l-i],r-i):1;
    while(i-z[i]>=0&&i+z[i]<len&&s[i-z[i]]==s[i+z[i]])</pre>
       ++z[i];
   if(i+z[i]>r) l=i,r=i+z[i];
} }
```

6.8 Smallest Rotation

```
//rotate(begin(s), begin(s)+minRotation(s), end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
    if(a+k == b | | s[a+k] < s[b+k])
      {b += max(0, k-1); break;}
  if(s[a+k] > s[b+k]) {a = b; break;}
  } return a;
}
```

6.9 Cyclic LCS

#define L 0

#define LU 1
#define U 2

```
const int mov[3][2]=\{0,-1,-1,-1,-1,0\};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al,j=bl,l=0;
  while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++;
    i+=mov[dir][0];
    j+=mov[dir][1];
  return 1:
inline void reroot(int r) { // r = new base row
  int i=r, j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
  if(j>bl) return;
  pred[i][j]=L;
while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
      pred[i][j]=L;
    } else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
       i++;
      pred[i][j]=L;
    } else {
      j++;
} } }
// note: a WILL be altered in process
              -- concatenated after itself
  char tmp[MAXL];
  if(al>bl) {
    swap(al,bl);
    strcpy(tmp,a);
    strcpy(a,b);
    strcpy(b,tmp);
  strcpy(tmp,a);
  strcat(a,tmp);
  // basic lcs
  for(int i=0;i<=2*al;i++) {</pre>
    dp[i][0]=0;
    pred[i][0]=U;
  for(int j=0;j<=bl;j++) {
  dp[0][j]=0;</pre>
    pred[0][j]=L;
  for(int i=1;i<=2*al;i++) {
    for(int j=1;j<=bl;j++) {</pre>
      if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
      else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
else if(a[i-1]==b[j-1]) pred[i][j]=LU;
       else pred[i][j]=U;
  } }
// do cyclic lcs
  int clcs=0;
  for(int i=0;i<al;i++) {</pre>
    clcs=max(clcs,lcs_length(i));
```

```
reroot(i+1);
}
// recover a
a[al]='\0';
return clcs;
}
```

6.10 Hash

```
//字串雜湊前的idx是0-base,雜湊後為1-base
//即區間為 [0,n-1] -> [1,n]
//若要取得區間[L,R]的值則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為len的字串和
//(h[i+len-1] - h[i-1] * qpow(p, len) % modl + modl)
//從j開始長度為len的字串是否相同
#define x first
#define y second
pair<int,int> Hash[MXN];
void build(const string& s){
  pair<int, int> val = make_pair(0,0);
  Hash[0]=val;
for(int i=1; i<=s.size(); i++){
val.x = (val.x * P1 + s[i-1]) % MOD;
  val.y = (val.y * P2 + s[i-1]) % MOD;
  Hash[i] = val;
  }
bool cmp( int i, int j, int len ) {
    return ((Hash[i+len-1].x-Hash[i-1].x*qpow(P1,len)%
MOD+MOD)%MOD == (Hash[j+len-1].x-Hash[j-1].x*
          qpow(P1,len)%MOD+MOD)%MOD)
    && ((Hash[i+len-1].y-Hash[i-1].y*qpow(P2,len)%MOD+
          MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow(
          P2, len)%MOD+MOD)%MOD);
}
```

7 Data Structure

7.1 Segment tree

```
// !!!注意build()時初始化用的陣列也是1-base
//!!!query(0,0) 會報錯
#define cl(x)(x*2)
#define cr(x) (x*2+1)
struct segmentTree {
    int n:
    vector<int> seg, tag, cov;
segmentTree(int _n): n(_n) {
         seg=tag=cov=vector<int>(n*4,0);
    void push(int i, int L, int R) {
         if(cov[i]) {
             seg[i]=cov[i]*(R-L+1);
             if(L < R)
                 cov[cl(i)]=cov[cr(i)]=cov[i];
                 tag[cl(i)]=tag[cr(i)]=0;
             cov[i]=0;
         if(tag[i]) {
             seg[i]+=tag[i]*(R-L+1);
             if(L < R) {
                 tag[cl(i)]+=tag[i];
                 tag[cr(i)]+=tag[i];
             tag[i]=0;
        }
    void pull(int i, int L, int R) {
   if(L >= R) return;
         int mid=(L+R)>>1;
        push(cl(i),L,mid)
        push(cr(i),mid+1,R);
```

```
seg[i]=seg[cl(i)]+seg[cr(i)];
void build(vector<int>& arr, int i=1, int L=1, int
    R=-1) {
    if(R == -1) R=n;
    if(L == R) {
         seg[i]=arr[L];
         return;
    int mid=(L+R)>>1;
    build(arr,cl(i),L,mid);
    build(arr,cr(i),mid+1,R);
    pull(i,L,R);
int query(int rL, int rR, int i=1, int L=1, int R
    =-1) {
    if(R == -1) R=n;
    push(i,L,R);
    if(rL <= L && R <= rR) return seg[i];</pre>
    int mid=(L+R)>>1, ret=0;
    if(rL <= mid) ret+=query(rL,rR,cl(i),L,mid);</pre>
    if(mid < rR) ret+=query(rL,rR,cr(i),mid+1,R);</pre>
    return ret;
void update(int rL, int rR, int val, int i=1, int L
    =1, int R=-1) {
    if(R == -1) R=n;
    push(i,L,R);
    if(rL <= L && R <= rR) {
         tag[i]=val;
        return;
    int mid=(L+R)>>1;
    if(rL <= mid) update(rL,rR,val,cl(i),L,mid);</pre>
    if(mid < rR) update(rL,rR,val,cr(i),mid+1,R);</pre>
    pull(i,L,R);
void cover(int rL, int rR, int val, int i=1, int L
    =1, int R=-1) {
if(R == -1) R=n;
    push(i,L,R);
    if(rL <= L && R <= rR) {
        cov[i]=val;
         return;
    int mid=(L+R)>>1;
    if(rL <= mid) cover(rL,rR,val,cl(i),L,mid);</pre>
    if(mid < rR) cover(rL,rR,val,cr(i),mid+1,R);</pre>
    pull(i,L,R);
}
Test Case:
1 2 3 4
2 1 3
1 1 3 1
2 1 3
1 1 4 1
2 1 4
```

7.2 持久化 SMT

```
struct node{
  node *1, *r;
  int val;
};

vector<node *> ver;
int arr[MXN] = {0};

//0-base
struct SegmentTree{
  int n;
  node *root;
  void build(int _n){
    n = _n;
    root = build(0, n-1);
}
```

```
node* build(int L, int R){
    node *x = new node();
    if(L == R){x->val = arr[L]; return x;}
    int mid = (L+R)/2;
    x->l = build(L, mid);
    x->r = build(mid + 1, R);
    x->val = x->l->val + x->r->val;
    return x;
  int query(node *ro, int L, int R){return query(ro, 0,
       n-1, L, R);}
  int query(int L, int R){return query(root, 0, n-1, L,
       R);}
  int query(node *x, int L, int R, int recL, int recR){
    if(recL <= L && R <= recR) return x->val;
    int mid = (L+R)/2, res = 0;
    if(recL <= mid) res += query(x->1, L, mid, recL,
        recR);
    if(mid < recR) res += query(x->r, mid+1, R, recL,
        recR);
    return res;
  }
  void update(int pos, int v){update(root, 0, n-1, pos,
  void update(node *x, int L, int R, int pos, int v){
  if(L == R){ x->val = v; arr[L] = v; return;}
    int mid = (L+R)/2;
    if(pos <= mid) update(x->1, L, mid, pos, v);
    else
                   update(x->r, mid+1, R, pos, v);
    x->val = x->l->val + x->r->val;
  node *update_ver(node *pre, int 1, int r, int pos,
      int v){
    node *x = new node(); //當前位置建立新節點
    if(l == r){
      x->val = v;
      return x;
    int mid = (l+r)>>1;
if(pos <= mid){ //更新左邊
      x->l = update_ver(pre->l, l, mid, pos, v); //左邊
          節點連向新節點
      x->r = pre->r; //右邊連到原本的右邊
    else{ //更新右邊
      x->l = pre->l; //左邊連到原本的左邊
      x->r = update_ver(pre->r, mid+1, r, pos, v); //
          右邊節點連向新節點
    x->val = x->l->val + x->r->val;
    return x;
}} seg;
void add_ver(int x,int v){
                              //修改位置 x 的值為 v
    ver.push_back(seg.update_ver(ver.back(), 0, seg.n
         -1, x, v));
}
```

7.3 Trie

```
struct trie{
 trie *nxt[26];
            //紀錄有多少個字串以此節點結尾
 int cnt;
 int sz;
            //有多少字串的前綴包括此節點
 trie():cnt(0),sz(0){
     memset(nxt,0,sizeof(nxt));
 }
};
trie *root = new trie(); //創建新的字典樹
void insert(string& s){
 trie *now = root; // 每次從根結點出發
 for(auto i:s){
   now->sz++;
   if(now->nxt[i-'a'] == NULL){
     now->nxt[i-'a'] = new trie();
   now = now->nxt[i-'a']; //走到下一個字母
```

```
now->cnt++; now->sz++;
int query_prefix(string& s){ //查詢有多少前綴為 s
 trie *now = root;
                     // 每次從根結點出發
  for(auto i:s){
   if(now->nxt[i-'a'] == NULL){
     return 0;
   now = now->nxt[i-'a'];
 return now->sz;
}
int query_count(string& s){ //查詢字串 s 出現次數
 trie *now = root;
                    // 每次從根結點出發
  for(auto i:s){
   if(now->nxt[i-'a'] == NULL){
     return 0;
   now = now->nxt[i-'a'];
 }
  return now->cnt;
```

7.4 Treap (interval reverse)

```
//拆出[a,b]區間就如同下面所展示先使用splitByTh()拆出
//左右,再把左區間拆成1, m最後merge()回去
//反轉區間時又記得使用^=可以直接反轉01
//treap 拆區間時從後面拆是因為這樣[a,b]的關係
//不用重新考慮·要是先拆前面b的位置會變成b-a+1
//0-base
//splitByTh(root,a-1,l,m);
//splitByTh(m,b-a+1,m,r);
mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
struct Treap {
  int key, pri, sz, tag, sum;
Treap *L, *R;
  Treap( int val ) {
    sum=key=val, pri=gen(), sz=1, tag=0;
    L=R=NULL;
};};
int Size( Treap *a ) { return !a?0:a->sz;}
void pull( Treap *a ) {
  a \rightarrow sz = Size(a \rightarrow L) + Size(a \rightarrow R) + 1;
  a->sum=a->key;
  if( a \rightarrow L ) a \rightarrow sum += a \rightarrow L \rightarrow sum;
  if( a \rightarrow R ) a \rightarrow sum += a \rightarrow R \rightarrow sum;
void push( Treap *a ) {
  if( a && a->tag ) {
    swap(a->L,a->R);
    if( a->L ) a->L->tag^=1;
if( a->R ) a->R->tag^=1;
    a \rightarrow tag=0:
Treap *merge(Treap *a, Treap *b) {
  if(!a || !b ) return a?a:b;
  push(a), push(b);
  if( a->pri > b->pri ) {
    a \rightarrow R = merge(a \rightarrow R, b);
    pull(a); return a;
  b \rightarrow L = merge(a, b \rightarrow L);
  pull(b); return b;
void print(Treap *a) {
  if( !a ) return;
  push(a);
  print(a->L);
  cout.put(a->key);
  print(a->R);
Treap *buildTreap( int n, string& str ) {
```

```
Treap *root=NULL;
  for( int i=0 ; i < n ; i++ )</pre>
    root=merge(root, new Treap(str[i]));
  return root:
void splitbyk( Treap *x, int k, Treap *&a, Treap *&b )
  if(!x) a=b=NULL;
  else if( x->key <= k ) {
    splitbyk(x->R,k,a->R,b);
    pull(a);
  else {
    b=x
    splitbyk(x->L,k,a,b->L);
    pull(b);
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
  if( !x ) { a=b=NULL; return; }
  push(x);
  if( Size(x->L)+1 \le k ) {
    splitByTh(x->R,k-Size(x->L)-1,a->R,b);
    pull(a);
  }
  else {
    b=x;
    splitByTh(x->L,k,a,b->L);
    pull(b);
  }
}
signed main() {
  string str;
  int n, m;
  cin>>n>>m>>str;
  Treap *root;
  root=buildTreap(n,str);
  for( int i=0 ; i < m ; i++ ) {</pre>
    int a, b;
    cin>>a>>b;
    Treap *1, *m, *r;
    splitByTh(root,b,l,r);
    splitByTh(l,a-1,l,m);
    m->tag^=1;
    root=merge(l,merge(m,r));
  print(root);
}
```

7.5 Treap (interval erase)

```
//區間移除使用bitset維護區間值
mt19937 gen(chrono::steady_clock::now().
      time_since_epoch().count());
struct Treap {
char key;
int pri, sz;
bitset<128> tag;
   Treap *L, *R;
   Treap( char val ) {
     key=val, pri=gen(), sz=1;
     L=R=NULL;
     tag.set(key);
}; };
int Size( Treap *a ) { return !a?0:a->sz;}
void pull( Treap *a ) {
  if(!a) return;
  a \rightarrow sz = Size(a \rightarrow L) + Size(a \rightarrow R) + 1;
   a->tag=a->tag.reset();
   a->tag=a->tag.set(a->key);
  if( a \rightarrow L ) a \rightarrow tagl = a \rightarrow L \rightarrow tag;
   if( a \rightarrow R ) a \rightarrow tag = a \rightarrow R \rightarrow tag;
Treap *merge( Treap *a, Treap *b ) {
   if( !a || !b ) return a?a:b;
   if( a->pri > b->pri ) {
     a \rightarrow R = merge(a \rightarrow R, b);
```

```
pull(a);
    return a;
 b \rightarrow L = merge(a, b \rightarrow L);
 pull(b);
  return b;
Treap *buildTreap( int n, string& str ) {
 Treap *root=NULL;
  for( int i=0 ; i < n ; i++ )</pre>
    root=merge(root, new Treap(str[i]));
  return root;
void print( Treap *a ) {
 if( !a ) return;
 print(a->L);
 cout.put(a->key);
 print(a->R);
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
  if( !x ) { a=b=NULL; return; }
  if( Size(x->L)+1 \le k ) {
    splitByTh(x->R,k-Size(x->L)-1,a->R,b);
    pull(a);
 else {
   b=x;
    splitByTh(x->L,k,a,b->L);
   pull(b);
void erase( Treap *&x, char ch ) {
 if( !x || !x->tag.test(ch) ) return;
 erase(x->L,ch);
  erase(x->R,ch);
  if(x->key == ch) {
   Treap *l=x->L, *r=x->R;
    x=NULL;
   x=merge(l,r);
 pull(x);
signed main() {
 string str;
  int n, m;
 cin>>n>>m>>str;
 Treap *root;
  root=buildTreap(n,str);
  for( int i=0 ; i < m ; i++ ) {</pre>
    char c;
    int a, b;
   cin>>á>>b>>c;
Treap *l, *m, *r;
if( !root || !root->tag.test(c) ) continue;
    splitByTh(root,b,l,r);
    splitByTh(l,a-1,l,m);
    if( m || !m->tag.test(c) ) erase(m,c);
    root=merge(l,merge(m,r));
 print(root);
```

7.6 Link-Cut Tree

```
struct Splay {
    static Splay nil, mem[MEM], *pmem;
    Splay *ch[2], *f;
    int val, rev, size;
    Splay (int _val=-1) : val(_val), rev(0), size(1)
    { f = ch[0] = ch[1] = &nil; }
    bool isr()
    { return f->ch[0] != this && f->ch[1] != this; }
    int dir()
    { return f->ch[0] == this ? 0 : 1; }
    void setCh(Splay *c, int d){
        ch[d] = c;
        if (c != &nil) c->f = this;
        pull();
```

```
void push(){
    if( !rev ) return;
    swap(ch[0], ch[1]);
if (ch[0] != &nil) ch[0]->rev ^= 1;
    if (ch[1] != &nil) ch[1]->rev ^= 1;
    rev=0;
  void pull(){
    size = ch[0] -> size + ch[1] -> size + 1;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::
    mem;
Splay *nil = &Splay::nil;
void rotate(Splay *x){
  Splay *p = x->f;
int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d)
  p->pull(); x->pull();
vector<Splay*> splayVec;
void splay(Splay *x){
  splayVec.clear();
  for (Splay *q=x;; q=q->f){
    splayVec.push_back(q);
    if (q->isr()) break;
  reverse(begin(splayVec), end(splayVec));
for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir())
       rotate(x->f),rotate(x);
    else rotate(x), rotate(x);
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x){
  Splay *q = nil;
  for (;x!=nil;x=x->f){
    splay(x)
    x->setCh(q, 1);
    q = x;
  return q;
}
void chroot(Splay *x){
  access(x);
  splay(x);
  x \rightarrow rev \stackrel{\wedge}{=} 1;
  x->push(); x->pull();
void link(Splay *x, Splay *y){
  access(x);
  splay(x);
  chroot(y)
  x \rightarrow setCh(y, 1);
void cut_p(Splay *y) {
  access(y);
  splay(y)
  y->push();
  y->ch[0] = y->ch[0]->f = nil;
void cut(Splay *x, Splay *y){
  chroot(x);
  cut_p(y);
Splay* get_root(Splay *x) {
  access(x);
  splay(x);
  for(; x - ch[0] != nil; x = x - ch[0])
    x -> push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
```

```
x = get_root(x);
y = get_root(y);
return x == y;
}
Splay* lca(Splay *x, Splay *y) {
  access(x);
  access(y);
  splay(x);
  if (x->f == nil) return x;
  else return x->f;
}
```

7.7 BIT

```
#define lowbit(x) (x&-x)
struct BIT {
    int n;
    vector<int> bit;
    BIT(int _n):n(_n), bit(n+1) {}
    void update( int x, int val ) {
        for(;x <= n; x += lowbit(x)) bit[x] += val;</pre>
    void range_update( int L, int R, int val ) {
        update(L,val), update(R+1,-val);
    int query( int x ) {
        int res = 0;
        for(;x; x -= lowbit(x)) res += bit[x];
        return res;
    int range_query( int L, int R ) {
        return query(R)-query(L-1);
};
```

7.8 Black Magic

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
    tree_order_statistics_node_update> set_t;
tree<int,null_type,less_equal<int>,rb_tree_tag,
    tree_order_statistics_node_update> mt_t;
#include <ext/pb_ds/assoc_container.hpp>
typedef cc_hash_table<int,int> umap_t;
// gp_hash_table<int, int>
typedef priority_queue<int> heap;
#include<ext/rope>
using namespace __gnu_cxx;
int main(){
  // Insert some entries into s.
  set_t s; s.insert(12); s.insert(505);
  // The order of the keys should be: 12, 505.
  assert(*s.find_by_order(0) == 12);
  assert(*s.find_by_order(3) == 505);
  // The order of the keys should be: 12, 505.
  assert(s.order_of_key(12) == 0);
  assert(s.order_of_key(505) == 1);
  // Erase an entry.
  s.erase(12);
  // The order of the keys should be: 505.
  assert(*s.find_by_order(0) == 505);
  // The order of the keys should be: 505.
  assert(s.order_of_key(505) == 0);
  // if we want to delete less_equal tag tree
  mt_t.erase(mt_t.find_by_order(mt_t.order_of_key(val))
       );
  heap h1 , h2; h1.join( h2 );
  rope<char> r[ 2 ];
r[ 1 ] = r[ 0 ]; // persistenet
string t = "abc";
 r[1].insert(0, t.c_str());
r[1].erase(1,1);
cout << r[1].substr(0,2);
```

8 Others

8.1 SOS dp

```
for(int i = 0; i<(1<<N); ++i)
  F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1<<
      N); ++mask){
    if(mask & (1<<i))
      F[mask] += F[mask^(1<<i)];
}</pre>
```

8.2 Max subrectangle

```
const int N = 1e5+5;
int n, a[N], l[N], r[N];
long long ans;
int main() {
  while (cin>>n) {
    ans = 0;
    for (int i = 1; i \le n; i++) cin>>a[i], l[i] = r[i]
    for (int i = 1; i <= n; i++)
      while (l[i] > 1 \& a[i] <= a[l[i] - 1]) l[i] = l[
          l[i] - 1];
    for (int i = n; i >= 1; i--)
      while (r[i] < n \& a[i] <= a[r[i] + 1]) r[i] = r[
          r[i] + 1];
    for (int i = 1; i <= n; i++)
      ans = max(ans, (long long)(r[i] - l[i] + 1) * a[i]
          ]);
    cout<<ans<<"\n";
}
```

8.3 De Brujin sequence

```
// return cyclic array of length k^n such that every
// array of length n using 0~k-1 appears as a subarray.
vector<int> DeBruijn(int k,int n){
   if(k=1) return {0};
   vector<int> aux(k*n),res;
   function<void(int,int)> f=[&](int t,int p)->void{
      if(t>n){   if(n%p==0)
         for(int i=1;i<=p;++i) res.push_back(aux[i]);
   }else{
      aux[t]=aux[t-p]; f(t+1,p);
      for(aux[t]=aux[t-p]+1;aux[t]<k;++aux[t]) f(t+1,t)
      ;
   }
   };
   f(1,1); return res;
}</pre>
```

8.4 CDQ 分治

```
if(j < R) {
      arr[j].ans += bt->query(arr[j].w);
      temp.push_back(arr[j]);
      j++;
  for(int i = L; i < mid; i++)</pre>
   bt->update(arr[i].w, -arr[i].cnt);
  copy(temp.begin(), temp.end(), arr.begin() + L);
signed main()
{
 // n 個數 k 值域範圍
 int n, k;
 cin >> n >> k;
 map<tuple<int, int, int>, int> mp;
 vector<int> res(n, 0);
 vector<triple> arr
 bt = new BIT(k + 1);
 for(int i = 0; i < n; i++) {
      int x, y, z;
      cin >> x >> y >> z;
      mp[{x, y, z}]++;
  for(auto t : mp)
   arr.push_back({get<0>(t.first), get<1>(t.first),
        get<2>(t.first), 0, t.second});
  cdq(0, arr.size(), arr);
 for(auto \&[x,y,z,a,b] : arr) res[a + b - 1] += b;
  for(int i : res) cout << i << '\n';</pre>
```

8.5 3D LIS

```
#define lowbit(x) (x&-x)
const int MAXN=1e5+5;
struct BIT {
  int n:
  vector<int> bit;
  BIT( int _n ):n(_n), bit(_n+1,0) {}
  int query( int x ) {
    int res=0;
    for(; x > 0 ; x-=lowbit(x) ) res=max(res,bit[x]);
    return res;
  void update( int x, int val )
    for(; x <= n ; x+=lowbit(x) ) {
  if( val < 0 ) bit[x]=0;</pre>
       else bit[x]=max(bit[x],val);
    }
}bt(MAXN);
struct triple {
  int u, v, w, ans, cnt;
  bool operator<( triple b ) { return u<b.u; }</pre>
bool cmp( triple a, triple b ) {return a.v<b.v;}</pre>
void cdq( int L, int R, vector<triple>& arr ) {
  if( R-L <= 1 ) return;</pre>
  int mid=L+R>>1:
  cdq(L,mid,arr);
  sort(arr.begin()+L,arr.begin()+mid,cmp);
  sort(arr.begin()+mid,arr.begin()+R,cmp);
  for( int i=L, j=mid ; i < mid || j < R ; ) {
  for(; i < mid && ( j >= R || arr[i].v < arr[j].v )</pre>
    ; i++ ) bt.update(arr[i].w,arr[i].ans); if( j < R ) {
       arr[j].ans=max(bt.query(arr[j].w-1)+1,arr[j].ans)
       j++;
    }
  for( int i=L ; i < mid ; i++ ) bt.update(arr[i].w,-1)</pre>
  sort(arr.begin()+L,arr.begin()+R);
  cdq(mid,R,arr);
signed main()
```

```
ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0)
int n, res=0;
cin>>n;
vector<int> ls;
vector<triple> arr;
for( int i=0 ; i < n ; i++ ) {
 int a, b;
 cin>>a>>b:
  arr.push_back({i,a,b,1,1});//{第一維,第二維,第三維,
      答案,數量}
 ls.push_back(b);
sort(ls.begin(),ls.end());
ls.resize(unique(ls.begin(),ls.end())-ls.begin());
for( auto &t : arr ) t.w=lower_bound(ls.begin(),ls.
    end(),t.w)-ls.begin()+1;
n=arr.size();
cdq(0,n,arr);
for( int i=0 ; i < n ; i++ ) res=max(res,arr[i].ans);</pre>
cout<<res<<'\n';
```

8.6 Ternary Search

```
while(L <= R) {
   int ml = L + (R - L) / 3, mr = R - (R - L) / 3;
   if(L == R) return L;
   else if( checker(ml) < checker(mr) ) L = ml + 1;
   else R = mr - 1;
}</pre>
```

8.7 Max Subrectangle

```
const int N = 1e5+5;
int n, a[N], l[N], r[N];
long long ans;
int main() {
  while (cin>>n) {
    ans = 0;
    for (int i = 1; i <= n; i++) cin>>a[i], l[i] = r[i]
    for (int i = 1; i <= n; i++)
      while (l[i] > 1 \&\& a[i] <= a[l[i] - 1]) l[i] = l[
           l[i] - 1];
    for (int i = n; i >= 1; i--)
      while (r[i] < n \&\& a[i] <= a[r[i] + 1]) r[i] = r[
    r[i] + 1];
for (int i = 1; i <= n; i++)
      ans = max(ans, (long long)(r[i] - l[i] + 1) * a[i]
           ]);
    cout<<ans<<"\n";
  }
}
```

8.8 Maximal Rectangle

8.9 p-Median

8.10 Tree Knapsack

8.11 AC-Automaton

```
// 1-based
// n is the number of patterns
struct Automaton {
    static const int MXN = 1e6;
    int n, cnt, vis[MXN], rev[MXN], indeg[MXN], ans[MXN
        ];
    queue<int> q;
    struct trie_node {
        vector<int> son;
        int fail, flag, ans;
        trie_node(): son(27), fail(0), flag(0) {}
    } trie[MXN];
    void init(int _n) {
        n = _n, cnt = 1;
        for (int i = 1; i <= n; i++) vis[i] = 0;
    // insert a string s with number num
    // num is the index of the pattern
    void insert(string s, int num) {
        int u = 1, len = s.size();
        for (int i = 0; i < len; i++) {
    int v = s[i] - 'a';
            if (!trie[u].son[v]) trie[u].son[v] = ++cnt
            u = trie[u].son[v];
        if (!trie[u].flag) trie[u].flag = num;
        rev[num] = trie[u].flag;
    void getfail() {
```

```
for (int i = 0; i < 26; i++) trie[0].son[i] =</pre>
         q.push(1);
         trie[1].fail = 0;
         while (q.size()) {
             int u = q.front(); q.pop();
int Fail = trie[u].fail;
             for (int i = 0; i < 26; i++) {
                  int v = trie[u].son[i];
                  if (!v) {
                      trie[u].son[i] = trie[Fail].son[i];
                      continue;
                  trie[v].fail = trie[Fail].son[i];
                  indeg[trie[Fail].son[i]]++;
                  q.push(v);
             }
         }
    void topu() {
         for (int i = 1; i <= cnt; i++)</pre>
             if (!indeg[i]) q.push(i);
         while (q.size()) {
             int fr = q.front(); q.pop();
             vis[trie[fr].flag] = trie[fr].ans;
             int u = trie[fr].fail;
             trie[u].ans += trie[fr].ans;
             if (!--indeg[u]) q.push(u);
         }
    void query(string &s) {
         int u = 1, len = s.size();
         for (int i = 0; i < len; i++) u = trie[u].son[s
    [i] - 'a'], trie[u].ans++;</pre>
    void solve(string &s) {
         getfail();
         query(s);
         topu();
         for (int i = 1; i <= n; i++) ans[i] = vis[rev[i
             ]];
} AC;
```

8.12 Aho-Corasick

```
struct ACautomata{
  struct Node{
    int cnt,i;
    Node *go[26], *fail, *dic;
    Node (){
      cnt = 0; fail = 0; dic = 0; i = 0;
      memset(go,0,sizeof(go));
  }pool[1048576],*root;
  int nMem,n_pattern;
  Node* new_Node(){
    pool[nMem] = Node();
    return &pool[nMem++];
  void init() {
    nMem=0;root=new_Node();n_pattern=0;
    add("");
  void add(const string &str) { insert(root,str,0); }
  void insert(Node *cur, const string &str, int pos){
  for(int i=pos;i<str.size();i++){</pre>
      if(!cur->go[str[i]-'a'])
         cur->go[str[i]-'a'] = new_Node();
      cur=cur->go[str[i]-'a'];
    cur->cnt++; cur->i=n_pattern++;
  void make_fail(){
    queue<Node*> que;
    que.push(root);
    while (!que.empty()){
      Node* fr=que.front(); que.pop();
      for (int i=0; i<26; i++){
        if (fr->go[i]){
```











