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1 Basic

1.1 default code

```
#pragma GCC optimize("O3,unroll-loops")
#pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
#include <bits/stdc++.h>
using namespace std;
ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
```

1.2 .vimrc

```
set nu rnu ts=4 sw=4 bs=2 ai hls cin mouse=a
color default
sy on
inoremap {<CR> {<CR>}<C-o>0
inoremap jk <Esc>
nnoremap J 5j
nnoremap K 5k
nnoremap run :w<bar>!g++ -std=c++14 -DLOCAL -Wfatal-
errors -o test "%" && echo "done." && time ./test<
CR>
```

1.3 Increase Stack Size (linux)

```
#include <sys/resource.h>
void increase_stack_size() {
    const rlim_t ks = 64*1024*1024;
    struct rlimit rl;
    int res=getrlimit(RLIMIT_STACK, &rl);
    if(res==0){
        if(rl.rlim_cur<ks){
            rl.rlim_cur=ks;
            res=setrlimit(RLIMIT_STACK, &rl);
        } } }
```

1.4 Misc

```
編譯參數: -std=c++14 -Wall -Wshadow (-fsanitize=
undefined)

mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }

#define SECs ((double)clock() / CLOCKS_PER_SEC)
double startTime;
bool TIME() { // 比最大可執行時間小一點
    return SECs - startTime > 0.8;
}
int main() {
    startTime = SECs;
}
```

```

struct KeyHasher {
    size_t operator()(const Key& k) const {
        return k.first + k.second * 100000;
    } };
typedef unordered_map<Key,int,KeyHasher> map_t;
// builtin function 可以代的值為int32
__builtin_popcountll    // 二進位有幾個1
__builtin_clzll         // 左起第一個1之前0的個數
__builtin_parityll      // 1的個數的奇偶性
__builtin_mul_overflow(a,b,&h) // a*b是否溢位

```

1.5 check

```

for ((i=0;;i++))
do
    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done

```

1.6 python-related

```

parser:
int(eval(num.replace("/", "")))

from fractions import Fraction
from decimal import Decimal, getcontext, ROUND_HALF_UP,
    ROUND_CEILING, ROUND_FLOOR
getcontext().prec = 250 # set precision
getcontext().rounding = ROUND_HALF_UP

itwo = Decimal(0.5)
two = Decimal(2)

format(x, '0.10f') # set precision

N = 200
def angle(cosT):
    """given cos(theta) in decimal return theta"""
    for i in range(N):
        cosT = ((cosT + 1) / two) ** itwo
        sinT = (1 - cosT * cosT) ** itwo
        return sinT * (2 ** N)
pi = angle(Decimal(-1))

"""round to 2 decimal places"""
sum = Decimal(input())
sum.quantize(Decimal('.00'), ROUND_HALF_UP)

"""Fraction"""
x = Fraction(1, 3) # 1/3
x.as_integer_ratio() # (1, 3)

"""input list of integers"""
arr = list(map(int, input().split()))

"""把字元轉成ascii再轉回字串"""
chr(ord('a'))

```

2 flow

2.1 ISAP $O(V^3)$

```

struct Maxflow {
    static const int MAXV = 20010;
    static const int INF = 1000000;
    struct Edge {
        int v, c, r;
        Edge(int _v, int _c, int _r):
            v(_v), c(_c), r(_r) {}
    };
};

```

```

int s, t;
vector<Edge> G[MAXV*2];
int iter[MAXV*2], d[MAXV*2], gap[MAXV*2], tot;
void init(int x) {
    tot = x+2;
    s = x+1, t = x+2;
    for(int i = 0; i <= tot; i++) {
        G[i].clear();
        iter[i] = d[i] = gap[i] = 0;
    }
}
void addEdge(int u, int v, int c) {
    G[u].push_back(Edge(v, c, SZ(G[v])));
    G[v].push_back(Edge(u, 0, SZ(G[u]) - 1));
}
int dfs(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {
        Edge &e = G[p][i];
        if(e.c > 0 && d[p] == d[e.v]+1) {
            int f = dfs(e.v, min(flow, e.c));
            if(f) {
                e.c -= f;
                G[e.v][e.r].c += f;
                return f;
            }
        }
    }
    if( (--gap[d[p]]) == 0) d[s] = tot;
    else {
        d[p]++;
        iter[p] = 0;
        ++gap[d[p]];
    }
    return 0;
}
int solve() {
    int res = 0;
    gap[0] = tot;
    for(res = 0; d[s] < tot; res += dfs(s, INF));
    return res;
}
void reset() {
    for(int i=0;i<=tot;i++) {
        iter[i]=d[i]=gap[i]=0;
    }
}
} } }flow;

```

2.2 MinCostFlow

```

struct zkwflow{
    static const int maxN=10000;
    struct Edge{ int v,f,re; ll w;};
    int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
    vector<Edge> E[maxN];
    void init(int _n,int _s,int _t){
        n=_n,s=_s,t=_t;
        for(int i=0;i<n;i++) E[i].clear();
    }
    void addEdge(int u,int v,int f,ll w){
        E[u].push_back({v,f,(int)E[v].size(),w});
        E[v].push_back({u,0,(int)E[u].size()-1,-w});
    }
    bool SPFA(){
        fill_n(dis,n,LLONG_MAX); fill_n(vis,n,false);
        queue<int> q; q.push(s); dis[s]=0;
        while (!q.empty()){
            int u=q.front(); q.pop(); vis[u]=false;
            for(auto &it:E[u]){
                if(it.f>0&&dis[it.v]>dis[u]+it.w){
                    dis[it.v]=dis[u]+it.w;
                    if(!vis[it.v]){
                        vis[it.v]=true; q.push(it.v);
                    }
                }
            }
        }
        return dis[t]!=LLONG_MAX;
    }
    int DFS(int u,int nf){
        if(u==t) return nf;
        int res=0; vis[u]=true;
        for(int &i=ptr[u];i<(int)E[u].size();i++){
            auto &it=E[u][i];
            if(it.f>0&&dis[it.v]==dis[u]+it.w&&!vis[it.v]){
                int tf=DFS(it.v,min(nf,it.f));
                res+=tf,nf-=tf,it.f-=tf;
            }
        }
    }
};

```

```

        E[it.v][it.re].f+=tf;
        if(nf==0){ vis[u]=false; break; }
    }
    return res;
}
pair<int,ll> flow(){
    int flow=0; ll cost=0;
    while (SPFA()){
        fill_n(ptr,n,0);
        int f=DFS(s,INT_MAX);
        flow+=f; cost+=dis[t]*f;
    }
    return{ flow,cost };
} // reset: do nothing
} flow;

```

2.3 Dinic $O(V^2E)$

```

#define SZ(x) (int)x.size()
#define PB push_back
struct Dinic{
    struct Edge{ int v,f,re; };
    int n,s,t,level[MXN];
    vector<Edge> E[MXN];
    void init(int _n, int _s, int _t){
        n = _n; s = _s; t = _t;
        for (int i=0; i<n; i++) E[i].clear();
    }
    void add_edge(int u, int v, int f){
        E[u].PB({v,f,SZ(E[v])});
        E[v].PB({u,0,SZ(E[u])-1});
    }
    bool BFS(){
        for (int i=0; i<n; i++) level[i] = -1;
        queue<int> que;
        que.push(s);
        level[s] = 0;
        while (!que.empty()){
            int u = que.front(); que.pop();
            for (auto it : E[u]){
                if (it.f > 0 && level[it.v] == -1){
                    level[it.v] = level[u]+1;
                    que.push(it.v);
                }
            }
            return level[t] != -1;
        }
    }
    int DFS(int u, int nf){
        if (u == t) return nf;
        int res = 0;
        for (auto &it : E[u]){
            if (it.f > 0 && level[it.v] == level[u]+1){
                int tf = DFS(it.v, min(nf,it.f));
                res += tf; nf -= tf; it.f -= tf;
                E[it.v][it.re].f += tf;
                if (nf == 0) return res;
            }
        }
        if (!res) level[u] = -1;
        return res;
    }
    int flow(int res=0){
        while (BFS())
            res += DFS(s,2147483647);
        return res;
    }
} flow;

```

2.4 Kuhn Munkres 最大完美二分匹配 $O(n^3)$

```

struct KM{ // max weight, for min negate the weights
    int n, mx[MXN], my[MXN], pa[MXN];
    ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
    bool vx[MXN], vy[MXN];
    void init(int _n) { // 1-based
        n = _n;
        for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
    }
    void addEdge(int x, int y, ll w) {g[x][y] = w;}
    void augment(int y) {

```

```

        for(int x, z; y; y = z)
            x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
    }
    void bfs(int st) {
        for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;
        queue<int> q; q.push(st);
        for(;;) {
            while(q.size()) {
                int x=q.front(); q.pop(); vx[x]=1;
                for(int y=1; y<=n; ++y) if(!vy[y]){
                    ll t = lx[x]+ly[y]-g[x][y];
                    if(t==0){
                        pa[y]=x;
                        if(!my[y]){augment(y);return;}
                        vy[y]=1, q.push(my[y]);
                    }else if(sy[y]>t) pa[y]=x,sy[y]=t;
                }
            }
            ll cut = INF;
            for(int y=1; y<=n; ++y)
                if(!vy[y]&&cut>sy[y]) cut=sy[y];
            for(int j=1; j<=n; ++j){
                if(vx[j]) lx[j] -= cut;
                if(vy[j]) ly[j] += cut;
                else sy[j] -= cut;
            }
            for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
                if(!my[y]){augment(y);return;}
                vy[y]=1, q.push(my[y]);
            }
        }
    }
    ll solve(){
        fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
        fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
        for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
            lx[x] = max(lx[x], g[x][y]);
        for(int x=1; x<=n; ++x) bfs(x);
        ll ans = 0;
        for(int y=1; y<=n; ++y) ans += g[my[y]][y];
        return ans;
    }
} graph;

```

2.5 Flow Method

Maximize $c^T x$ subject to $Ax \leq b, x \geq 0$;
with the corresponding symmetric dual problem,
Minimize $b^T y$ subject to $A^T y \geq c, y \geq 0$.

Maximize $c^T x$ subject to $Ax \leq b$;
with the corresponding asymmetric dual problem,
Minimize $b^T y$ subject to $A^T y = c, y \geq 0$.

Minimum vertex cover on bipartite graph =
Maximum matching on bipartite graph

Minimum edge cover on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)

König Theorem

最小點覆蓋：選出最少的點，滿足每條邊至少有一個端點被選
二分圖中，最小點覆蓋 = 最大匹配

Independent set on bipartite graph =
vertex number - Minimum vertex cover(Maximum matching)
二分圖中，最大獨立集 = n - 最小點覆蓋
找出最小點覆蓋，做完dinic之後，
從源點dfs只走還有流量的邊，
左邊沒被走到的點跟右邊被走到的點就是答案，
其他點為最大獨立集

最大閉包(最大權閉合子圖)

源點連到所有正權點流量為點權
所有負權點連到匯點流量為點權(絕對值)
所有圖上的邊權重為 INF

路徑覆蓋數量

把每個點拆成 入點 和 出點，轉化為二分圖
原圖頂點數 - 二分圖最大匹配數

Maximum density subgraph $(\sum W_e + \sum W_v) / |V|$

Binary search on answer:

For a fixed D, construct a Max flow model as follow:

Let S be Sum of all weight(or inf)

1. from source to each node with cap = S
2. For each (u,v,w) in E, (u→v, cap=w), (v→u, cap=w)
3. For each node v, from v to sink with cap = S + 2 * D - deg[v] - 2 * (W of v)

where $\deg[v] = \sum \text{weight of edge associated with } v$

If maxflow < S * |V|, D is an answer.

Requiring subgraph: all vertex can be reached from source with edge whose cap > 0.

3 Math

3.1 FFT

```
// const int MXN = 262144 (MXN must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx; //real() ,imag()
const ld PI = acos(-1);
const cplx I(0, 1);

struct FFT{
    cplx omega[MXN+1];
    FFT(){ //pre_fft
        for(int i=0; i<=MXN; i++)
            omega[i] = exp(i * 2 * PI / MXN * I);
    }
    // n must be 2^k
    void fft(int n, cplx a[], bool inv=false){
        int basic = MXN / n;
        int theta = basic;
        for (int m = n; m >= 2; m >>= 1) {
            int mh = m >> 1;
            for (int i = 0; i < mh; i++) {
                cplx w = omega[inv ? MXN-(i*theta%MXN) : i*theta%MXN];
                for (int j = i; j < n; j += m) {
                    int k = j + mh;
                    cplx x = a[j] - a[k];
                    a[j] += a[k];
                    a[k] = w * x;
                }
            }
            theta = (theta * 2) % MXN;
        }
        int i = 0;
        for (int j = 1; j < n - 1; j++) {
            for (int k = n >> 1; k > (i ^ k); k >>= 1);
            if (j < i) swap(a[i], a[j]);
        }
        if(inv) for (i = 0; i < n; i++) a[i] /= n;
    }
    cplx arr[MXN+1];
    inline void mul(int _n, ll a[], int _m, ll b[], ll ans[])
    {
        int n=1, sum=_n+_m-1;
        while(n<sum)
            n<<=1;
        for(int i=0; i<n; i++) {
            double x=(i<_n?a[i]:0), y=(i<_m?b[i]:0);
            arr[i]=complex<double>(x+y, x-y);
        }
        fft(n, arr);
        for(int i=0; i<n; i++)
            arr[i]=arr[i]*arr[i];
        fft(n, arr, true);
        for(int i=0; i<sum; i++)
            ans[i]=(long long)(arr[i].real()/4+0.5);
    }
}fft;
```

3.2 Faulhaber $(\sum_{i=1}^n i^p)$

```
/* faulhaber' s formula -
 * cal power sum formula of all p=1~k in O(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinatorics
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
    int a=x, b=mod, a0=1, a1=0, b0=0, b1=1;
    while(b) {
        int q, t;
        q=a/b; t=b; b=a-b*q; a=t;
        t=b0; b0=a0-b0*q; a0=t;
        t=b1; b1=a1-b1*q; a1=t;
    }
    return a0<0?a0+mod:a0;
}
inline void pre() {
    /* combinational */
    for(int i=0; i<=MAXK; i++) {
        cm[i][0]=cm[i][i]=1;
        for(int j=1; j<i; j++)
            cm[i][j]=add(cm[i-1][j-1], cm[i-1][j]);
    }
    /* inverse */
    for(int i=1; i<=MAXK; i++) inv[i]=getinv(i);
    /* bernoulli */
    b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
    for(int i=2; i<=MAXK; i++) {
        if(i&1) { b[i]=0; continue; }
        b[i]=1;
        for(int j=0; j<i; j++)
            b[i]=sub(b[i], mul(cm[i][j], mul(b[j], inv[i-j+1])));
    }
    /* faulhaber */
    // sigma_x=1~n {x^p} =
    // 1/(p+1) * sigma_j=0~p {C(p+1, j)*Bj*n^(p-j+1)}
    for(int i=1; i<=MAXK; i++) {
        co[i][0]=0;
        for(int j=0; j<=i; j++)
            co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]));
    }
}
/* sample usage: return f(n,p) = sigma_x=1~n (x^p) */
inline int solve(int n, int p) {
    int sol=0, m=n;
    for(int i=1; i<=p+1; i++) {
        sol=add(sol, mul(co[p][i], m));
        m = mul(m, n);
    }
    return sol;
}
```

3.3 Chinese Remainder

```
LL x[N], m[N];
LL CRT(LL x1, LL m1, LL x2, LL m2) {
    LL g = __gcd(m1, m2);
    if((x2 - x1) % g) return -1; // no sol
    m1 /= g; m2 /= g;
    pair<LL, LL> p = gcd(m1, m2);
    LL lcm = m1 * m2 * g;
    LL res = p.first * (x2 - x1) * m1 + x1;
    return (res % lcm + lcm) % lcm;
}
LL solve(int n) { // n>=2, be careful with no solution
    LL res=CRT(x[0], m[0], x[1], m[1]), p=m[0]/__gcd(m[0], m[1])*m[1];
    for(int i=2; i<n; i++){
        res=CRT(res, p, x[i], m[i]);
        p=p/__gcd(p, m[i])*m[i];
    }
    return res;
}
```

```
} }
```

3.4 Miller Rabin

```
// n < 4,759,123,141      3 : 2, 7, 61
// n < 1,122,004,669,633 4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383 6 : pimes <= 13
// n < 2^64              7 :
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL mul(LL x, LL y, LL mod){
    LL ret=x*y-(LL)((long double)x/mod*y)*mod;
    // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
    return ret<0?ret+mod:ret;
}

LL magic[]={}
bool witness(LL a, LL n, LL u, int t){
    if(!a) return 0;
    LL x=mypow(a,u,n);
    for(int i=0; i<t; i++){
        LL nx=mul(x,x,n);
        if(nx==1&&x!=1&&x!=n-1) return 1;
        x=nx;
    }
    return x!=1;
}

bool miller_rabin(LL n) {
    int s=(magic number size)
    // iterate s times of witness on n
    if(n<2) return 0;
    if(!(n&1)) return n == 2;
    ll u=n-1; int t=0;
    // n-1 = u*2^t
    while(!(u&1)) u>>=1, t++;
    while(s--){
        LL a=magic[s]%n;
        if(witness(a,n,u,t)) return 0;
    }
    return 1;
}
```

3.5 Pollard Rho

```
// does not work when n is prime 0(n^(1/4))
LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
LL pollard_rho(LL n) {
    if(!(n&1)) return 2;
    while(true){
        LL y=2, x=rand()%(n-1)+1, res=1;
        for(int sz=2; res==1; sz*=2) {
            for(int i=0; i<sz && res<=1; i++) {
                x = f(x, n);
                res = __gcd(abs(x-y), n);
            }
            y = x;
        }
        if (res!=0 && res!=n) return res;
    }
}
```

3.6 Josephus Problem

```
int josephus(int n, int m){ //n人每m次
    int ans = 0;
    for (int i=1; i<=n; ++i)
        ans = (ans + m) % i;
    return ans;
}
```

3.7 Gaussian Elimination

```
const int GAUSS_MOD = 100000007LL;
struct GAUSS{
    int n;
    vector<vector<int>> v;
    int ppow(int a, int k){
        if(k == 0) return 1;
        if(k % 2 == 0) return ppow(a * a % GAUSS_MOD, k >> 1);
        if(k % 2 == 1) return ppow(a * a % GAUSS_MOD, k >> 1) * a % GAUSS_MOD;
    }
    vector<int> solve(){
        vector<int> ans(n);
        REP(now, 0, n){
            REP(i, now, n) if(v[now][now] == 0 && v[i][now] != 0)
                swap(v[i], v[now]); // det = -det;
            if(v[now][now] == 0) return ans;
            int inv = ppow(v[now][now], GAUSS_MOD - 2);
            REP(i, 0, n) if(i != now){
                int tmp = v[i][now] * inv % GAUSS_MOD;
                REP(j, now, n+1) (v[i][j] += GAUSS_MOD - tmp * v[now][j] % GAUSS_MOD) %= GAUSS_MOD;
            }
        }
        REP(i, 0, n) ans[i] = v[i][n+1] * ppow(v[i][i], GAUSS_MOD - 2) % GAUSS_MOD;
        return ans;
    }
    // gs.v.clear(), gs.v.resize(n, vector<int>(n+1, 0));
} gs;
```

3.8 Inverse Matrix

```
int GAUSS_MOD;
struct GAUSS{
    int n;
    vector<vector<int>> v;
    vector<vector<int>> rev;
    int mul(int x, int y, int mod){
        int ret=x*y-(int)((long double)x/mod*y)*mod;
        return ret<0?ret+mod:ret;
    }
    int ppow(int a, int b){//res=(a^b)%m
        int res=1, k=a;
        while(b){
            if((b&1)) res=mul(res,k,GAUSS_MOD)%GAUSS_MOD;
            k=mul(k,k,GAUSS_MOD)%GAUSS_MOD;
            b>>=1;
        }
        return res%GAUSS_MOD;
    }
    bool solve(){
        for(int now = 0; now < n; now++){
            int ch;
            for(ch = now; ch < n && !v[ch][now]; ch++);
            if(ch >= n) return 0;
            for(int i = now; i < n; i++) if(v[now][now] == 0 && v[i][now] != 0){
                swap(v[i], v[now]); // det = -det;
                swap(rev[i], rev[now]);
            }
            if(v[now][now] == 0) return 0;
            int inv = ppow(v[now][now], GAUSS_MOD - 2);
            for(int i = 0; i < n; i++) if(i != now){
                int tmp = v[i][now] * inv % GAUSS_MOD;
                for(int j = 0; j < n; j++) {
                    (v[i][j] += GAUSS_MOD - tmp * v[now][j] % GAUSS_MOD) %= GAUSS_MOD;
                    (rev[i][j] += GAUSS_MOD - tmp * rev[now][j] % GAUSS_MOD) %= GAUSS_MOD;
                }
            }
        }
        return 1;
    }
} gs;

signed main(){
    int n, p; //n*n matrix, MOD=p
```

```

cin>>n>>p; //if(!n && !p) return 0;
GAUSS_MOD = p; gs.n = n;
gs.v.clear() , gs.v.resize(n + 1, vector<int>(n + 2 ,
0));
gs.rev.clear() , gs.rev.resize(n + 1, vector<int>(n +
2 , 0));
for(int i = 0; i < n; i++){
    for(int j = 0; j < n; j++){
        cin>>gs.v[i][j];
        if(i == j) gs.rev[i][j] = 1;
    }
}
if(!gs.solve()) cout << "singular\n";
else{
    for(int i = 0; i < n; i++){
        int inv = gs.ppow(gs.v[i][i] , p - 2);
        for(int j = 0; j < n; j++){
            cout << (gs.rev[i][j] * inv % p) << " ";
            cout<<"\n";
        }
    }
    cout << "\n";
}

```

3.9 模反元素

```

long long inv(long long a, long long m){
    long long x, y;
    long long d = exgcd(a, m, x, y);
    if(d == 1) return (x + m) % m;
    else return -1; // -1 為無解
}

```

3.10 $ax+by=gcd$

```

PII gcd(int a, int b){
    if(b == 0) return {1, 0};
    PII q = gcd(b, a % b);
    return {q.second, q.first - q.second * (a / b)};
}

int exgcd(int a, int b, long long &x, long long &y) {
    if(b == 0){x = 1, y = 0; return a;}
    int now = exgcd(b, a % b, y, x);
    y -= a / b * x;
    return now;
}

```

3.11 Discrete sqrt

```

void calcH(LL &t, LL &h, const LL p) {
    LL tmp = p - 1; for(t = 0; (tmp & 1) == 0; tmp /= 2) t++; h = tmp;
}
// solve equation  $x^2 \bmod p = a$ 
bool solve(LL a, LL p, LL &x, LL &y) {
    if(p == 2) {x = y = 1; return true;}
    int p2 = p / 2, tmp = mypow(a, p2, p);
    if(tmp == p - 1) return false;
    if((p + 1) % 4 == 0) {
        x = mypow(a, (p + 1) / 4, p); y = p - x; return true;
    } else {
        LL t, h, b, pb; calcH(t, h, p);
        if(t >= 2) {
            do {b = rand() % (p - 2) + 2;} while (mypow(b, p / 2, p) != p - 1);
            pb = mypow(b, h, p);
        } int s = mypow(a, h / 2, p);
        for(int step = 2; step <= t; step++) {
            int ss = (((LL)(s * s) % p) * a) % p;
            for(int i = 0; i < t - step; i++) ss = mul(ss, ss, p);
            if(ss + 1 == p) s = (s * pb) % p;
            pb = ((LL)pb * pb) % p;
        } x = ((LL)s * a) % p; y = p - x;
    } return true;
}

```

3.12 Prefix Inverse

```

void solve(int m){
    inv[1] = 1;
    for(int i = 2; i < m; i++){
        inv[i] = ((LL)(m - m / i) * inv[m % i]) % m;
    }
}

```

3.13 Roots of Polynomial 找多項式的根

```

const double eps = 1e-12;
const double inf = 1e+12;
double a[10], x[10]; // a[0..n](coef) must be filled
int n; // degree of polynomial must be filled
int sign(double x){return (x < -eps)?(-1):(x > eps);}
double f(double a[], int n, double x){
    double tmp = 1, sum = 0;
    for(int i = 0; i <= n; i++){
        sum = sum + a[i] * tmp; tmp = tmp * x;
    } return sum;
}
double binary(double l, double r, double a[], int n){
    int sl = sign(f(a, n, l)), sr = sign(f(a, n, r));
    if(sl == 0) return l; if(sr == 0) return r;
    if(sl * sr > 0) return inf;
    while(r - l > eps){
        double mid = (l + r) / 2;
        int ss = sign(f(a, n, mid));
        if(ss == 0) return mid;
        if(ss * sl > 0) l = mid; else r = mid;
    }
    return l;
}
void solve(int n, double a[], double x[], int &nx){
    if(n == 1){x[1] = -a[0] / a[1]; nx = 1; return;}
    double da[10], dx[10]; int ndx;
    for(int i = n; i >= 1; i--) da[i - 1] = a[i] * i;
    solve(n - 1, da, dx, ndx);
    nx = 0;
    if(ndx == 0){
        double tmp = binary(-inf, inf, a, n);
        if(tmp < inf) x[++nx] = tmp;
        return;
    }
    double tmp;
    tmp = binary(-inf, dx[1], a, n);
    if(tmp < inf) x[++nx] = tmp;
    for(int i = 1; i <= ndx - 1; i++){
        tmp = binary(dx[i], dx[i + 1], a, n);
        if(tmp < inf) x[++nx] = tmp;
    }
    tmp = binary(dx[ndx], inf, a, n);
    if(tmp < inf) x[++nx] = tmp;
} // roots are stored in x[1..nx]

```

3.14 Combination theorem

```

const ll mod = 1e9 + 7;
ll fac[(int)2e6 + 1], inv[(int)2e6 + 1];
ll getinv(ll a){return qpow(a, mod - 2);}
void init(int n){
    fac[0] = 1;
    for(int i = 1; i <= n; i++){
        fac[i] = fac[i - 1] * i % mod;
    }
    inv[n] = getinv(fac[n]);
    for(int i = n - 1; i >= 0; i--){
        inv[i] = inv[i + 1] * (i + 1) % mod;
    }
}
ll C(int n, int m){
    if(m > n) return 0;
    return fac[n] * inv[m] % mod * inv[n - m] % mod;
}

```


3.15 Primes

```

/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
 * 999983, 1097774749, 1076767633, 100102021, 999997771
 * 1001010013, 1000512343, 987654361, 999991231
 * 999888733, 98789101, 987777733, 999991921, 1010101333
 * 1010102101, 10000000000039, 1000000000000037
 * 2305843009213693951, 4611686018427387847
 * 9223372036854775783, 18446744073709551557 */
int mu[ N ], p_tbl[ N ];
vector<int> primes;
void sieve() {
    mu[ 1 ] = p_tbl[ 1 ] = 1;
    for( int i = 2 ; i < N ; i ++ ){
        if( !p_tbl[ i ] ){
            p_tbl[ i ] = i;
            primes.push_back( i );
            mu[ i ] = -1;
        }
        for( int p : primes ){
            int x = i * p;
            if( x >= M ) break;
            p_tbl[ x ] = p;
            mu[ x ] = -mu[ i ];
            if( i % p == 0 ){
                mu[ x ] = 0;
                break;
            }
        }
    }
}
vector<int> factor( int x ){
    vector<int> fac{ 1 };
    while( x > 1 ){
        int fn = SZ(fac), p = p_tbl[ x ], pos = 0;
        while( x % p == 0 ){
            x /= p;
            for( int i = 0 ; i < fn ; i ++ )
                fac.PB( fac[ pos ++ ] * p );
        }
    }
    return fac;
}

```

3.16 Phi

```

ll phi(ll n){ // 計算小於n的數中與n互質的有幾個
    ll res = n, a=n;
    for(ll i=2;i*i<=a;i++){ // O(sqrtN)
        if(a%i==0){
            res = res/i*(i-1);
            while(a%i==0) a/=i;
        }
    }
    if(a>1) res=res/a*(a-1);
    return res;
}

```

3.17 Int Sqrt

```

LL intSqrt(LL S) { //return origin val when S <= 0
    if (S <= 0) return S;
    LL x = S;
    for (LL nx;;x = nx){
        nx = (x+S/x)>>1LL;
        if(nx >= x) break;
    }
    return x;
}

```

3.18 Result

- Lucas' Theorem :
For $n, m \in \mathbb{Z}^*$ and prime P , $C(m, n) \bmod P = \prod (C(m_i, n_i))$ where m_i is the i -th digit of m in base P .
- Stirling approximation :
$$n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\frac{1}{12n}}$$
- Stirling Numbers(permutation $|P| = n$ with k cycles):
$$S(n, k) = \text{coefficient of } x^k \text{ in } \prod_{i=0}^{n-1} (x+i)$$

- Stirling Numbers(Partition n elements into k non-empty set):

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

- Pick' s Theorem : $A = i + b/2 - 1$
在二維座標平面中畫上網格，對於任何簡單多邊形
 A : 面積、 i : 內部的格點數、 b : 邊上的格點數
- Catalan number : $C_n = \binom{2n}{n} / (n+1)$
$$C_n^{n+m} - C_{n+1}^{n+m} = (m+n)! \frac{n-m+1}{n+1} \quad \text{for } n \geq m$$

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)n!}$$

$$C_0 = 1 \quad \text{and} \quad C_{n+1} = 2 \left(\frac{2n+1}{n+2} \right) C_n$$

$$C_0 = 1 \quad \text{and} \quad C_{n+1} = \sum_{i=0}^n C_i C_{n-i} \quad \text{for } n \geq 0$$

- Euler Characteristic:
planar graph: $V - E + F - C = 1$
convex polyhedron: $V - E + F = 2$
 V, E, F, C : number of vertices, edges, faces(regions), and components

- Kirchhoff's theorem :
 $A_{ii} = \deg(i), A_{ij} = (i, j) \in E ? -1 : 0$, Deleting any one row, one column, and cal the $\det(A)$

- Polya' theorem (c is number of color, m is the number of cycle size):
$$\left(\sum_{i=1}^m c^{gcd(i, m)} \right) / m$$

- Burnside lemma:
$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

- 錯排公式: (n 個人中，每個人皆不再原來位置的組合數):
 $dp[0] = 1; dp[1] = 0;$
 $dp[i] = (i-1) * (dp[i-1] + dp[i-2]);$

- Bell 數 (有 n 個人，把他們拆組的方法總數) :
 $B_0 = 1$
 $B_n = \sum_{k=0}^n s(n, k) \quad (\text{second - stirling})$
 $B_{n+1} = \sum_{k=0}^n \binom{n}{k} B_k$

- Wilson's theorem :
 $(p-1)! \equiv -1 \pmod{p}$

- Fermat's little theorem :
 $a^p \equiv a \pmod{p}$

- Euler's totient function:
 $A^{B^C} \bmod p = \text{pow}(A, \text{pow}(B, C, p-1)) \bmod p$

- 歐拉函數降冪公式:
 $A^B \bmod C = A^{B \bmod \phi(C) + \phi(C)} \bmod C$

- 用歐拉函數求模反元素:
如果 a 和 n 互質，則 a 對 n 的模反元素
 $a^{-1} \equiv a^{\phi(n)-1} \pmod{n}$

- 6 的倍數:
 $(a-1)^3 + (a+1)^3 + (-a)^3 + (-a)^3 = 6a$

- 上高斯 (向上取整) :
 $\lceil \frac{a}{b} \rceil = \frac{a+b-1}{b}$

- 點到直線距離公式:
$$d = \frac{|ax_0 + by_0 + c|}{\sqrt{a^2 + b^2}}$$

4 Geometry

4.1 definition

```

#define all(a) a.begin(),a.end()
ostream& operator<<(ostream& os, const Pt& pt) {
    return os << "(" << pt.x << ", " << pt.y << ")";
}
typedef long double ld;
const ld eps = 1e-8;
const ld pi = acos(-1);
int dcmp(ld x) {
    if(abs(x) < eps) return 0;
    else return x < 0 ? -1 : 1;
}
struct Pt {
    ld x, y;
    Pt(ld _x=0, ld _y=0):x(_x), y(_y) {}
    Pt operator+(const Pt &a) const {
        return Pt(x+a.x, y+a.y);
    }
    Pt operator-(const Pt &a) const {

```

```

    return Pt(x-a.x, y-a.y); }
Pt operator*(const ld &a) const {
    return Pt(x*a, y*a); }
Pt operator/(const ld &a) const {
    return Pt(x/a, y/a); }
ld operator*(const Pt &a) const {
    return x*a.x + y*a.y; }
ld operator^(const Pt &a) const {
    return x*a.y - y*a.x; }
bool operator<(const Pt &a) const {
    return x < a.x || (x == a.x && y < a.y); }
//return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0 &&
    dcmp(y-a.y) < 0); }
bool operator==(const Pt &a) const {
    return dcmp(x-a.x) == 0 && dcmp(y-a.y) == 0; }
};
ld norm2(const Pt &a) {
    return a*a; }
ld norm(const Pt &a) {
    return sqrt(norm2(a)); }
Pt perp(const Pt &a) {
    return Pt(-a.y, a.x); }
Pt rotate(const Pt &a, ld ang) {
    return Pt(a.x*cos(ang)-a.y*sin(ang), a.x*sin(ang)+a.y
        *cos(ang)); }
bool collinear(Pt a, Pt b, Pt c) { return ((b - a) ^ (c
    - a)) == 0; }
struct Circle {
    Pt o; ld r;
    Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
};

```

4.2 Intersection of 2 lines

```

// NAN(parallel), INF(overlapping)
Pt LLIntersect(Line a, Line b) {
    Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
    ld f1 = (p2-p1)^(q1-p1), f2 = (p2-p1)^(p1-q2), f;
    if(dcmp(f=f1+f2) == 0)
        return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
    return q1*(f2/f) + q2*(f1/f);
}

```

4.3 halfPlaneIntersection

```

// O(nlogn)
// 傳入 vector<Line>
// (半平面為點 st 往 ed 的逆時針方向)
// 回傳值為形成的凸多邊形的頂點 vector
// assume that Lines intersect
vector<Pt> HPI(vector<Line> P) {
    sort(P.begin(), P.end(), [&](Line l, Line m) {
        if (argcmp(l.v, m.v)) return true;
        if (argcmp(m.v, l.v)) return false;
        return PtSide(l.s, m) > 0;
    });
    int n = P.size(), l = 0, r = -1;
    for (int i = 0; i < n; i++) {
        if (i and !argcmp(P[i-1].v, P[i].v)) continue;
        while (l < r and PtSide(LLIntersect(P[r-1], P[r]),
            P[i]) <= 0) r--;
        while (l < r and PtSide(LLIntersect(P[l], P[l+1]),
            P[i]) <= 0) l++;
        P[++r] = P[i];
    }
    while (l < r and PtSide(LLIntersect(P[r-1], P[r]),
        P[l]) <= 0) r--;
    while (l < r and PtSide(LLIntersect(P[l], P[l+1]),
        P[r]) <= 0) l++;
    if (r - l <= 1 or !argcmp(P[l].v, P[r].v))
        return {}; // empty
    if (PtSide(LLIntersect(P[l], P[r]), P[l+1]) <= 0) {
        assert(0);
        return {}; // infinity
    }
    vector<Line> lns = vector(P.begin() + l, P.begin()
        + r + 1);
}

```

```

lns.push_back(lns[0]);
vector<Pt> hpi;
for(int i = 1; i < lns.size(); i++) hpi.push_back(
    LLIntersect(lns[i-1], lns[i]));
return hpi;
}

```

4.4 Convex Hull

```

double cross(Pt o, Pt a, Pt b){
    return (a-o) ^ (b-o);
}
vector<Pt> convex_hull(vector<Pt> pt){
    sort(pt.begin(),pt.end());
    int top=0;
    vector<Pt> stk(2*pt.size());
    for (int i=0; i<(int)pt.size(); i++){
        while (top >= 2 && cross(stk[top-2],stk[top-1],pt[i]
            ]) <= 0) // 如果想要有點共線的點 · 把 <= 改成 <
            top--;
        stk[top++] = pt[i];
    }
    for (int i=pt.size()-2, t=top+1; i>=0; i--){
        while (top >= t && cross(stk[top-2],stk[top-1],pt[i]
            ]) <= 0)
            top--;
        stk[top++] = pt[i];
    }
    stk.resize(top-1);
    return stk;
}

```

4.5 Convex Hull trick

```

struct Convex {
    int n;
    vector<Pt> A, V, L, U;
    Convex(const vector<Pt> &A) : A(_A), n(_A.size()) {
        // n >= 3
        auto it = max_element(all(A));
        L.assign(A.begin(), it + 1);
        U.assign(it, A.end()), U.push_back(A[0]);
        for (int i = 0; i < n; i++) {
            V.push_back(A[(i + 1) % n] - A[i]);
        }
    }
    int PtSide(Pt p, Line L) {
        return dcmp((L.b - L.a)^(p - L.a));
    }
    int inside(Pt p, const vector<Pt> &h, auto f) {
        auto it = lower_bound(all(h), p, f);
        if (it == h.end()) return 0;
        if (it == h.begin()) return p == *it;
        return 1 - dcmp((p - *prev(it))^( *it - *prev(it)))
            ;
    }
    // 1. whether a given point is inside the CH
    // ret 0: out, 1: on, 2: in
    int inside(Pt p) {
        return min(inside(p, L, less{}), inside(p, U,
            greater{}));
    }
    static bool cmp(Pt a, Pt b) { return dcmp(a ^ b) > 0; }
}
// 2. Find tangent points of a given vector
// ret the idx of far/closer tangent point
int tangent(Pt v, bool close = true) {
    assert(v != Pt{});
    auto l = V.begin(), r = V.begin() + L.size() - 1;
    if (v < Pt{}) l = r, r = V.end();
    if (close) return (lower_bound(l, r, v, cmp) - V.
        begin()) % n;
    return (upper_bound(l, r, v, cmp) - V.begin()) % n;
}
// 3. Find 2 tang pts on CH of a given outside point
// return index of tangent points
// return {-1, -1} if inside CH
array<int, 2> tangent2(Pt p) {
}

```



```

array<int, 2> t{-1, -1};
if (inside(p) == 2) return t;
if (auto it = lower_bound(all(L), p); it != L.end()
    and p == *it) {
    int s = it - L.begin();
    return {(s + 1) % n, (s - 1 + n) % n};
}
if (auto it = lower_bound(all(U), p, greater{}); it
    != U.end() and p == *it) {
    int s = it - U.begin() + L.size() - 1;
    return {(s + 1) % n, (s - 1 + n) % n};
}
for (int i = 0; i != t[0]; i = tangent((A[t[0] = i]
    - p), 0));
for (int i = 0; i != t[1]; i = tangent((p - A[t[1]
    = i]), 1));
return t;
}
int find(int l, int r, Line L) {
    if (r < l) r += n;
    int s = PtSide(A[l % n], L);
    return *ranges::partition_point(views::iota(l, r),
        [&](int m) {
            return PtSide(A[m % n], L) == s;
        }) - 1;
};
// 4. Find intersection point of a given line
// intersection is on edge (i, next(i))
vector<int> intersect(Line L) {
    int l = tangent(L.a - L.b), r = tangent(L.b - L.a);
    if(PtSide(A[l], L) == 0) return {l};
    if(PtSide(A[r], L) == 0) return {r};
    if (PtSide(A[l], L) * PtSide(A[r], L) > 0) return
        {};
    return {find(l, r, L) % n, find(r, l, L) % n};
}
};

```

4.6 Intersection of 2 segments

```

int ori( const Pt& o , const Pt& a , const Pt& b ){
    LL ret = ( a - o ) ^ ( b - o );
    return (ret > 0) - (ret < 0);
}
// p1 == p2 || q1 == q2 need to be handled
bool banana( const Pt& p1 , const Pt& p2 ,
    const Pt& q1 , const Pt& q2 ){
    if( ( ( p2 - p1 ) ^ ( q2 - q1 ) ) == 0 ){ // parallel
        if( ori( p1 , p2 , q1 ) ) return false;
        return ( ( p1 - q1 ) * ( p2 - q1 ) ) <= 0 ||
            ( ( p1 - q2 ) * ( p2 - q2 ) ) <= 0 ||
            ( ( q1 - p1 ) * ( q2 - p1 ) ) <= 0 ||
            ( ( q1 - p2 ) * ( q2 - p2 ) ) <= 0;
    }
    return (ori( p1, p2, q1 ) * ori( p1, p2, q2 ) <= 0) &&
        (ori( q1, q2, p1 ) * ori( q1, q2, p2 ) <= 0);
}

```

4.7 Intersection of Polygon and Circle

```

ld PCIntersect(vector<Pt> v, Circle cir) {
    for(int i = 0 ; i < (int)v.size() ; ++i) v[i] = v[i]
        - cir.o;
    ld ans = 0, r = cir.r;
    int n = v.size();
    for(int i = 0 ; i < n ; ++i) {
        Pt pa = v[i], pb = v[(i+1)%n];
        if(norm(pa) < norm(pb)) swap(pa, pb);
        if(dcmp(norm(pb)) == 0) continue;
        ld s, h, theta;
        ld a = norm(pb), b = norm(pa), c = norm(pb-pa);
        ld cosB = (pb*(pb-pa))/a/c, B = acos(cosB);
        if(cosB > 1) B = 0;
        else if(cosB < -1) B = PI;
        ld cosC = (pa*pb)/a/b, C = acos(cosC);
        if(cosC > 1) C = 0;
        else if(cosC < -1) C = PI;
        if(a > r) {

```

```

            s = (C/2)*r*r;
            h = a*b*sin(C)/c;
            if(h < r && B < PI/2) s -= (acos(h/r)*r*r - h*
                sqrt(r*r-h*h));
        }
        else if(b > r) {
            theta = PI - B - asin(sin(B)/r*a);
            s = 0.5*a*r*sin(theta) + (C-theta)/2*r*r;
        }
        else s = 0.5*sin(C)*a*b;
        ans += abs(s)*dcmp(v[i]^v[(i+1)%n]);
    }
    return abs(ans);
}

```

4.8 Circle cover

```

#define N 1021
#define D long double
struct CircleCover{
    int C; Circle c[ N ]; //填入C(圖數量),c(圖陣列)
    bool g[ N ][ N ], overlap[ N ][ N ];
    // Area[i] : area covered by at least i circles
    D Area[ N ];
    void init( int _C ){ C = _C; }
    bool Cinter( Circle& a , Circle& b , Pt& p1 , Pt& p2
        ){
        Pt o1 = a.o , o2 = b.o;
        D r1 = a.r , r2 = b.r;
        if( norm( o1 - o2 ) > r1 + r2 ) return {};
        if( norm( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
            return {};
        D d2 = ( o1 - o2 ) * ( o1 - o2 );
        D d = sqrt(d2);
        if( d > r1 + r2 ) return false;
        Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
        D A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
        Pt v=Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
        p1 = u + v; p2 = u - v;
        return true;
    }
}
struct Teve {
    Pt p; D ang; int add;
    Teve() {}
    Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
    bool operator<(const Teve &a)const {
        return ang < a.ang;
    }
}teve[ N * 2 ];
// strict: x = 0, otherwise x = -1
bool disjunct( Circle& a, Circle& b, int x )
{return dcmp( norm( a.o - b.o ) - a.r - b.r ) > x;}
bool contain( Circle& a, Circle& b, int x )
{return dcmp( a.r - b.r - norm( a.o - b.o ) ) > x;}
bool contain(int i, int j){
    /* c[j] is non-strictly in c[i]. */
    return (dcmp(c[i].r - c[j].r) > 0 ||
        (dcmp(c[i].r - c[j].r) == 0 && i < j) ) &&
        contain(c[i], c[j], -1);
}
void solve(){
    for( int i = 0 ; i <= C + 1 ; i ++ )
        Area[ i ] = 0;
    for( int i = 0 ; i < C ; i ++ )
        for( int j = 0 ; j < C ; j ++ )
            overlap[i][j] = contain(i, j);
    for( int i = 0 ; i < C ; i ++ )
        for( int j = 0 ; j < C ; j ++ )
            g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                disjunct(c[i], c[j], -1));
    for( int i = 0 ; i < C ; i ++ ){
        int E = 0, cnt = 1;
        for( int j = 0 ; j < C ; j ++ )
            if( j != i && overlap[j][i] )
                cnt ++;
        for( int j = 0 ; j < C ; j ++ )
            if( i != j && g[i][j] ){
                Pt aa, bb;
                Cinter(c[i], c[j], aa, bb);
                D A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
                D B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);

```

4.9 Tangent line of two circles

4.10 Minimum distance of two convex

4.11 Poly Union

```

return (p.x-p1.x)/(p2.x-p1.x);
}
double polyUnion(int n){ //py[0~n-1] must be filled
    int i,j,ii,jj,ta,tb,r,d; double z,w,s,sum=0,tc,td;
    for(i=0;i<n;i++) py[i][py[i].n]=py[i][0];
    for(i=0;i<n;i++){
        for(ii=0;ii<py[i].n;ii++){
            r=0;
            c[r++]=make_pair(0.0,0); c[r++]=make_pair(1.0,0);
            for(j=0;j<n;j++){
                if(i==j) continue;
                for(jj=0;jj<py[j].n;jj++){
                    ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))
                    ;
                    tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj
                    +1]));
                    if(ta==0 && tb==0){
                        if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[i
                        ][ii])>0&&j<i){
                            c[r++]=make_pair(segP(py[j][jj],py[i][ii
                            ],py[i][ii+1]),1);
                            c[r++]=make_pair(segP(py[j][jj+1],py[i][
                            ii],py[i][ii+1]),-1);
                        }
                    }else if(ta>=0 && tb<0){
                        tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
                        td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
                        c[r++]=make_pair(tc/(tc-td),1);
                    }else if(ta<0 && tb>=0){
                        tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
                        td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
                        c[r++]=make_pair(tc/(tc-td),-1);
                    } } }
            sort(c,c+r);
            z=min(max(c[0].first,0.0),1.0); d=c[0].second; s
            =0;
            for(j=1;j<r;j++){
                w=min(max(c[j].first,0.0),1.0);
                if(!d) s+=w-z;
                d+=c[j].second; z=w;
            }
            sum+=(py[i][ii]^py[i][ii+1])*s;
        } }
    return sum/2;
}

```

4.12 Minkowski sum

```
// P, Q, R(return) are counterclockwise order convex
    polygon
#define all(a) a.begin(),a.end()
vector<Pt> Minkowski(vector<Pt> P, vector<Pt> Q) {
    auto cmp = [&](Pt a, Pt b) {
        return Pt{a.y, a.x} < Pt{b.y, b.x};
    };
    auto reorder = [&](auto &R) {
        rotate(R.begin(), min_element(all(R), cmp), R.
            end());
        R.push_back(R[0]), R.push_back(R[1]);
    };
    const int n = P.size(), m = Q.size();
    reorder(P), reorder(Q);
    vector<Pt> R;
    for (int i = 0, j = 0, s; i < n or j < m; ) {
        R.push_back(P[i] + Q[j]);
        s = dcmp((P[i + 1] - P[i]) ^ (Q[j + 1] - Q[j]));
        if (s >= 0) i++;
        if (s <= 0) j++;
    }
    rotate(R.begin(), min_element(all(R)), R.end());
    return R;
}
```

4.13 Area of Rectangles

```
struct AreaofRectangles{
#define cl(x) (x<<1)
```

```

#define cr(x) (x<<11)
ll n, id, sid;
pair<ll,ll> tree[MXN<<3]; // count, area
vector<ll> ind;
tuple<ll,ll,ll> scan[MXN<<1];
void pull(int i, int l, int r){
    if(tree[i].first) tree[i].second = ind[r+1] -
        ind[l];
    else if(l != r){
        int mid = (l+r)>>1;
        tree[i].second = tree[cl(i)].second + tree[
            cr(i)].second;
    }
    else tree[i].second = 0;
}
void upd(int i, int l, int r, int ql, int qr, int v
){
    if(ql <= l && r <= qr){
        tree[i].first += v;
        pull(i, l, r); return;
    }
    int mid = (l+r) >> 1;
    if(ql <= mid) upd(cl(i), l, mid, ql, qr, v);
    if(qr > mid) upd(cr(i), mid+1, r, ql, qr, v);
    pull(i, l, r);
}
void init(int _n){
    n = _n; id = sid = 0;
    ind.clear(); ind.resize(n<<1);
    fill(tree, tree+(n<<2), make_pair(0, 0));
}
void addRectangle(int lx, int ly, int rx, int ry){
    ind[id++] = lx; ind[id++] = rx;
    scan[sid++] = make_tuple(ly, 1, lx, rx);
    scan[sid++] = make_tuple(ry, -1, lx, rx);
}
ll solve(){
    sort(ind.begin(), ind.end());
    ind.resize(unique(ind.begin(), ind.end()) - ind
        .begin());
    sort(scan, scan + sid);
    ll area = 0, pre = get<0>(scan[0]);
    for(int i = 0; i < sid; i++){
        auto [x, v, l, r] = scan[i];
        area += tree[1].second * (x-pre);
        upd(1, 0, ind.size()-1, lower_bound(ind
            .begin(), ind.end(), l)-ind.begin(),
            lower_bound(ind.begin(), ind.end(), r)-
            ind.begin()-1, v);
        pre = x;
    }
    return area;
}
} rect;

```

4.14 Min dist on Cuboid

```

typedef LL T;
T r;
void turn(T i, T j, T x, T y, T z,
    T x0, T y0, T L, T W, T H) {
    if (z==0) { T R = x*x+y*y; if (R<r) r=R; return; }
    if(i>=0 && i<2) turn(i+1, j, x0+L+z, y, x0+L-x,
        x0+L, y0, H, W, L);
    if(j>=0 && j<2) turn(i, j+1, x, y0+W+z, y0+W-y,
        x0, y0+W, L, H, W);
    if(i<=0 && i>-2) turn(i-1, j, x0-z, y, x-x0,
        x0-H, y0, H, W, L);
    if(j<=0 && j>-2) turn(i, j-1, x, y0-z, y-y0,
        x0, y0-H, L, H, W);
}
T solve(T L, T W, T H,
    T x1, T y1, T z1, T x2, T y2, T z2){
    if( z1!=0 && z1!=H ){
        if( y1==0 || y1==W )
            swap(y1,z1), swap(y2,z2), swap(W,H);
        else swap(x1,z1), swap(x2,z2), swap(L,H);
    }
    if (z1==H) z1=0, z2=H-z2;
    r=INF; turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H);
    return r;
}

```

```

}

```

4.15 Distance of Line and Point

```

ld Distance_of_Line_and_Point(Line l, Pt p) {
    ld cross_product = abs((p - l.s) ^ l.v);
    ld line_length = sqrtl(l.v * l.v);
    return cross_product / line_length;
}

```

4.16 Angle of two vector

```

// radian of OA and OB (directed angle)
ld Angle_of_two_vector(Pt A, Pt B, Pt O) {
    ld a = (A - O) * (B - O);
    ld b = (A - O) ^ (B - O);
    ld theta = atan2(b, a);
    return theta;
}

```

4.17 極角排序

```

//極角排序
//atan2(y, x) version
// p is reference point
// 180 度開始，逆時針排序，剛好在 180 度會排最後
bool cmp(Pt &lhs, Pt rhs) {
    return atan2((lhs - p).y, (lhs - p).x) < atan2((rhs
        - p).y, (rhs - p).x);
}

//cross product version
// p is reference point
// 270 度開始，逆時針排序，剛好在 270 度會排最後
bool cmp(const Pt& lhs, const Pt& rhs) {
    if ((lhs < p) ^ (rhs < p)) return (lhs < p) < (rhs
        < p);
    return ((lhs - p) ^ (rhs - p)) > 0;
}

```

4.18 Heart of Triangle

```

Pt inCenter( Pt &A, Pt &B, Pt &C) { // 內心
    double a = norm(B-C), b = norm(C-A), c = norm(A-B);
    return (A * a + B * b + C * c) / (a + b + c);
}
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心
    Pt bb = b - a, cc = c - a;
    double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc);
    return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d;
}
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心
    Pt ba = b - a, ca = c - a, bc = b - c;
    double Y = ba.Y * ca.Y * bc.Y,
        A = ca.X * ba.Y - ba.X * ca.Y,
        x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A,
        y0= -ba.X * (x0 - c.X) / ba.Y + ca.Y;
    return Pt(x0, y0);
}

```

5 Graph

5.1 Lowest Common Ancestor $O(\lg n)$

```

struct LCA {
    int n, ti, lgN;
    int anc[MXN + 5][__lg(MXN) + 1] = {0};
    int MaxLength[MXN][__lg(MXN) + 1] = {0};
    int time_in[MXN] = {0};
    int time_out[MXN] = {0};
}

```

```

LCA(int _n, int f):n(_n), ti(0), lgN(__lg(n)) {
    dfs(f, f, 0);
    build();
}
void dfs(int now, int f, int len_to_father) { // dfs
    for anc, time, Lenth
    ti++;
    anc[now][0] = f;
    time_in[now] = ti;
    MaxLength[now][0] = len_to_father;
    for (auto i : graph[now]) {
        if (i.first == f) continue;
        dfs(i.first, now, i.second);
    }
    time_out[now] = ti;
}
void build() { // build anc[], MaxLength[]
    for (int i = 1; i <= lgN; ++i) {
        for (int u = 1; u <= n; ++u) {
            anc[u][i] = anc[anc[u][i - 1]][i - 1];
            MaxLength[u][i] = max(MaxLength[u][i - 1],
                                   MaxLength[anc[u][i - 1]][i - 1]);
            // dis[u][i] += dis[anc[u][i - 1]][i - 1]
            // + dis[u][i - 1];
        }
    }
}
bool isAncestor(int x, int y) {
    return time_in[x] <= time_in[y] && time_out[x] >=
        time_out[y];
}
int getLCA(int u, int v) {
    if (isAncestor(u, v)) return u;
    if (isAncestor(v, u)) return v;
    for (int i = lgN; i >= 0; --i) {
        if (!isAncestor(anc[u][i], v)) {
            u = anc[u][i];
        }
    }
    return anc[u][0];
}
int getMax(int u, int v) { //獲得路徑上最大邊權
    int lca = getLCA(u, v);
    int maxx = -1;
    for (int i = lgN; i >= 0; --i) {
        // u to lca
        if (!isAncestor(anc[u][i], lca)) {
            maxx = max(maxx, MaxLength[u][i]);
            u = anc[u][i];
        }

        // v to lca
        if (!isAncestor(anc[v][i], lca)) {
            maxx = max(maxx, MaxLength[v][i]);
            v = anc[v][i];
        }
    }
    if (u != lca) maxx = max(maxx, MaxLength[u][0]);
    if (v != lca) maxx = max(maxx, MaxLength[v][0]);
    return maxx;
}
};

```

5.2 Hamiltonian path $O(n^2 2^n)$

```

//dp[i][j] = 目前在j節點走過{i}節點的最短路徑
for(int i=1; i < (1 << n); i++) {
    for(int j = 1; j < n; j++) {
        if(!((1 << j) & i) && (i & 1)) {
            for(int k = 0; k < n; k++) {
                if(j == k) continue;
                if((1 << k) & i) dp[j][i | (1 << j)] =
                    min(dp[j][i | (1 << j)], dp[k][i] + dis[k][j]);
            }
        }
    }
}
}

```

5.3 MaximumClique 最大團

```

#define N 111
struct MaxClique { // 0-base
    typedef bitset<N> Int;
    Int linkto[N], v[N];
    int n;
    void init(int _n) {
        n = _n;
        for(int i = 0; i < n; i++) {
            linkto[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a, int b) {
        v[a][b] = v[b][a] = 1;
    }
    int popcount(const Int& val) {
        return val.count();
    }
    int lowbit(const Int& val) {
        return val._Find_first();
    }
    int ans, stk[N];
    int id[N], di[N], deg[N];
    Int cans;
    void maxclique(int elem_num, Int candi) {
        if(elem_num > ans) {
            ans = elem_num; cans.reset();
            for(int i = 0; i < elem_num; i++)
                cans[id[stk[i]]] = 1;
        }
        int potential = elem_num + popcount(candi);
        if(potential <= ans) return;
        int pivot = lowbit(candi);
        Int smaller_candi = candi & (~linkto[pivot]);
        while(smaller_candi.count() && potential > ans) {
            int next = lowbit(smaller_candi);
            candi[next] = !candi[next];
            smaller_candi[next] = !smaller_candi[next];
            potential--;
            if(next == pivot || (smaller_candi & linkto[next]).count()) {
                stk[elem_num] = next;
                maxclique(elem_num + 1, candi & linkto[next]);
            }
        }
    }
    int solve() {
        for(int i = 0; i < n; i++) {
            id[i] = i; deg[i] = v[i].count();
        }
        sort(id, id + n, [&](int id1, int id2) {
            return deg[id1] > deg[id2];
        });
        for(int i = 0; i < n; i++) di[id[i]] = i;
        for(int i = 0; i < n; i++)
            for(int j = 0; j < n; j++)
                if(v[i][j]) linkto[di[i]][di[j]] = 1;
        Int cand; cand.reset();
        for(int i = 0; i < n; i++) cand[i] = 1;
        ans = 1;
        cans.reset(); cans[0] = 1;
        maxclique(0, cand);
        return ans;
    }
} solver;

```

5.4 MaximalClique 極大團

```

#define N 80
struct MaxClique { // 0-base
    typedef bitset<N> Int;
    Int lnk[N], v[N];
    int n;
    void init(int _n) {
        n = _n;
        for(int i = 0; i < n; i++) {
            lnk[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a, int b) {
        v[a][b] = v[b][a] = 1;
    }
    int ans, stk[N], id[N], di[N], deg[N];
    Int cans;
    void dfs(int elem_num, Int candi, Int ex) {
        if(candi.none() && ex.none()) {
            cans.reset();
            for(int i = 0; i < elem_num; i++)

```

```

    cans[id[stk[i]]] = 1;
    ans = elem_num; // cans is a maximal clique
    return;
}
int pivot = (candilex)._Find_first();
Int smaller_candi = candi & (~lnk[pivot]);
while(smaller_candi.count()){
    int nxt = smaller_candi._Find_first();
    candi[nxt] = smaller_candi[nxt] = 0;
    ex[nxt] = 1;
    stk[elem_num] = nxt;
    dfs(elem_num+1, candi&lnk[nxt], ex&lnk[nxt]);
}
}
int solve(){
    for(int i = 0; i < n; i++){
        id[i] = i; deg[i] = v[i].count();
    }
    sort(id, id + n, [&](int id1, int id2){
        return deg[id1] > deg[id2]; });
    for(int i = 0; i < n; i++) di[id[i]] = i;
    for(int i = 0; i < n; i++)
        for(int j = 0; j < n; j++)
            if(v[i][j]) lnk[di[i]][di[j]] = 1;
    ans = 1; cans.reset(); cans[0] = 1;
    dfs(0, Int(string(n, '1')), 0);
    return ans;
}
} solver;

```

5.5 BCC based on vertex 點雙聯通分量

```

#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
    int n, nScc, step, dfn[MXN], low[MXN];
    vector<int> E[MXN], sccv[MXN];
    int top, stk[MXN];
    void init(int _n) { // 初始化n點
        n = _n; nScc = step = 0;
        for (int i=0; i<n; i++) E[i].clear();
    }
    void addEdge(int u, int v) // 無向邊
    { E[u].PB(v); E[v].PB(u); }
    void DFS(int u, int f) {
        dfn[u] = low[u] = step++;
        stk[top++] = u;
        for (auto v:E[u]) {
            if (v == f) continue;
            if (dfn[v] == -1) {
                DFS(v, u);
                low[u] = min(low[u], low[v]);
                if (low[v] >= dfn[u]) {
                    int z;
                    sccv[nScc].clear();
                    do {
                        z = stk[--top];
                        sccv[nScc].PB(z);
                    } while (z != v);
                    sccv[nScc++].PB(u);
                }
            } else
                low[u] = min(low[u], dfn[v]);
        }
    }
    vector<vector<int>> solve() { // 回傳(size=2 橋, size
        >2 點雙聯通分量)
        vector<vector<int>> res;
        for (int i=0; i<n; i++)
            dfn[i] = low[i] = -1;
        for (int i=0; i<n; i++)
            if (dfn[i] == -1) {
                top = 0;
                DFS(i, i);
            }
        REP(i, nScc) res.PB(sccv[i]);
        return res;
    }
} graph;

```

5.6 Strongly Connected Component 強連通分量

```

#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x)) //fill zero
struct Scc{
    int n, nScc, vst[MXN], bln[MXN];
    vector<int> E[MXN], rE[MXN], vec;
    void init(int _n){
        n = _n;
        for (int i=0; i<MXN; i++)
            E[i].clear(), rE[i].clear();
    }
    void addEdge(int u, int v){
        E[u].PB(v); rE[v].PB(u);
    }
    void DFS(int u){
        vst[u]=1;
        for (auto v : E[u]) if (!vst[v]) DFS(v);
        vec.PB(u);
    }
    void rDFS(int u){
        vst[u] = 1; bln[u] = nScc;
        for (auto v : rE[u]) if (!vst[v]) rDFS(v);
    }
    void solve(){
        nScc = 0;
        vec.clear();
        FZ(vst);
        for (int i=0; i<n; i++)
            if (!vst[i]) DFS(i);
        reverse(vec.begin(), vec.end());
        FZ(vst);
        for (auto v : vec)
            if (!vst[v]){
                rDFS(v); nScc++;
            }
    }
};

```

5.7 ManhattanMST

```

//return {{u,v},w}: u <-> v (w), 需要再手動去重
//need Point definition
vector<pair<pair<int,int>, int>> ManhattanMST(vector<Pt
> P) {
    vector<int> id(P.size());
    iota(id.begin(), id.end(), 0);
    vector<pair<pair<int,int>, int>> edg;
    for (int k = 0; k < 4; k++) {
        sort(id.begin(), id.end(), [&](int i, int j) {
            return (P[i] - P[j]).x < (P[j] - P[i]).y;
        });
        map<int, int> sweep;
        for (int i : id) {
            auto it = sweep.lower_bound(-P[i].y);
            while (it != sweep.end()) {
                int j = it->second;
                Pt d = P[i] - P[j];
                if (d.y > d.x) break;
                edg.push_back({{i, j}, d.x + d.y});
                it = sweep.erase(it);
            }
            sweep[-P[i].y] = i;
        }
        for (Pt &p : P) {
            if (k % 2) p.x = -p.x;
            else swap(p.x, p.y);
        }
    }
    return edg;
}

```

5.8 Min Mean Cycle

```

/* minimum mean cycle O(VE) */
struct MMC{

```



```

#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
struct Edge { int v,u; double c; };
int n, m, prv[V][V], prve[V][V], vst[V];
Edge e[E];
vector<int> edgeID, cycle, rho;
double d[V][V];
void init( int _n )
{ n = _n; m = 0; }
// WARNING: TYPE matters
void addEdge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
    for(int i=0; i<n; i++) d[0][i]=0;
    for(int i=0; i<n; i++) {
        fill(d[i+1], d[i+1]+n, inf);
        for(int j=0; j<m; j++) {
            int v = e[j].v, u = e[j].u;
            if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                d[i+1][u] = d[i][v]+e[j].c;
                prv[i+1][u] = v;
                prve[i+1][u] = j;
            }
        }
    }
}
double solve(){
    // returns inf if no cycle, mmc otherwise
    double mmc=inf;
    int st = -1;
    bellman_ford();
    for(int i=0; i<n; i++) {
        double avg=-inf;
        for(int k=0; k<n; k++) {
            if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])/(n-k));
            else avg=max(avg,inf);
        }
        if (avg < mmc) tie(mmc, st) = tie(avg, i);
    }
    fill(vst,0); edgeID.clear(); cycle.clear(); rho.clear();
    for (int i=n; !vst[st]; st=prv[i--][st]) {
        vst[st]++;
        edgeID.PB(prve[i][st]);
        rho.PB(st);
    }
    while (vst[st] != 2) {
        if(rho.empty()) return inf;
        int v = rho.back(); rho.pop_back();
        cycle.PB(v);
        vst[v]++;
    }
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
    return mmc;
} }mmc;

```

5.9 Directed Graph Min Cost Cycle

```

// works in O(N M)
#define INF 100000000000000LL
#define N 5010
#define M 200010
struct edge{
    int to; LL w;
    edge(int a=0, LL b=0): to(a), w(b){}
};
struct node{
    LL d; int u, next;
    node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
}b[M];
struct DirectedGraphMinCycle{
    vector<edge> g[N], grev[N];
    LL dp[N][N], p[N], d[N], mu;
    bool inq[N];
    int n, bn, bsz, hd[N];
    void b_insert(LL d, int u){
        int i = d/mu;
        if(i >= bn)return;
        b[++bsz] = node(d, u, hd[i]);
    }

```

```

        hd[i] = bsz;
    }
    void init( int _n ){
        n = _n;
        for( int i = 1 ; i <= n ; i ++ )
            g[ i ].clear();
    }
    void addEdge( int ai , int bi , LL ci )
    { g[ai].push_back(edge(bi,ci)); }
    LL solve(){
        fill(dp[0], dp[0]+n+1, 0);
        for(int i=1; i<=n; i++){
            fill(dp[i+1], dp[i+1]+n+1, INF);
            for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){
                for(int k=0; k<(int)g[j].size(); k++){
                    dp[i][g[j][k].to] = min(dp[i][g[j][k].to],
                        dp[i-1][j]+g[j][k].w);
                }
            }
            mu=INF; LL bunbo=1;
            for(int i=1; i<=n; i++) if(dp[n][i] < INF){
                LL a=-INF, b=1;
                for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
                    if(a*(n-j) < b*(dp[n][i]-dp[j][i])){
                        a = dp[n][i]-dp[j][i];
                        b = n-j;
                    }
                }
                if(mu*b > bunbo*a)
                    mu = a, bunbo = b;
            }
            if(mu < 0) return -1; // negative cycle
            if(mu == INF) return INF; // no cycle
            if(mu == 0) return 0;
            for(int i=1; i<=n; i++){
                for(int j=0; j<(int)g[i].size(); j++){
                    g[i][j].w *= bunbo;
                }
            }
            memset(p, 0, sizeof(p));
            queue<int> q;
            for(int i=1; i<=n; i++){
                q.push(i);
                inq[i] = true;
            }
            while(!q.empty()){
                int i=q.front(); q.pop(); inq[i]=false;
                for(int j=0; j<(int)g[i].size(); j++){
                    if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
                        p[g[i][j].to] = p[i]+g[i][j].w-mu;
                        if(!inq[g[i][j].to]){
                            q.push(g[i][j].to);
                            inq[g[i][j].to] = true;
                        }
                    }
                }
            }
            for(int i=1; i<=n; i++) grev[i].clear();
            for(int i=1; i<=n; i++){
                for(int j=0; j<(int)g[i].size(); j++){
                    g[i][j].w += p[i]-p[g[i][j].to];
                    grev[g[i][j].to].push_back(edge(i, g[i][j].w));
                }
            }
            LL mldc = n*mu;
            for(int i=1; i<=n; i++){
                bn=mldc/mu, bsz=0;
                memset(hd, 0, sizeof(hd));
                fill(d+i+1, d+i+1+n+1, INF);
                b_insert(d[i]=0, i);
                for(int j=0; j<=bn-1; j++) for(int k=hd[j]; k; k=
                    b[k].next){
                    int u = b[k].u;
                    LL du = b[k].d;
                    if(du > d[u]) continue;
                    for(int l=0; l<(int)g[u].size(); l++) if(g[u][l].to > i){
                        if(d[g[u][l].to] > du + g[u][l].w){
                            d[g[u][l].to] = du + g[u][l].w;
                            b_insert(d[g[u][l].to], g[u][l].to);
                        }
                    }
                }
                for(int j=0; j<(int)grev[i].size(); j++) if(grev[i][j].to > i)
                    mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
            }
            return mldc / bunbo;
        }
    } }graph;

```

5.10 DominatorTree

```

struct DominatorTree{ // O(N)
#define REP(i,s,e) for(int i=(s);i<=(e);i++)
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
    int n, m, s;
    vector<int> g[ MAXN ], pred[ MAXN ];
    vector<int> cov[ MAXN ];
    int dfn[ MAXN ], nfd[ MAXN ], ts;
    int par[ MAXN ]; //idom[u] s到u的最後一個必經點
    int sdom[ MAXN ], idom[ MAXN ];
    int mom[ MAXN ], mn[ MAXN ];
    inline bool cmp( int u, int v )
    { return dfn[ u ] < dfn[ v ]; }
    int eval( int u ){
        if( mom[ u ] == u ) return u;
        int res = eval( mom[ u ] );
        if( cmp( sdom[ mn[ mom[ u ] ] ], sdom[ mn[ u ] ] ) )
            mn[ u ] = mn[ mom[ u ] ];
        return mom[ u ] = res;
    }
    void init( int _n, int _m, int _s ){
        ts = 0; n = _n; m = _m; s = _s;
        REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
    }
    void addEdge( int u, int v ){
        g[ u ].push_back( v );
        pred[ v ].push_back( u );
    }
    void dfs( int u ){
        ts++;
        dfn[ u ] = ts;
        nfd[ ts ] = u;
        for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
            par[ v ] = u;
            dfs( v );
        }
    }
    void build(){
        REP( i, 1, n ){
            dfn[ i ] = nfd[ i ] = 0;
            cov[ i ].clear();
            mom[ i ] = mn[ i ] = sdom[ i ] = i;
        }
        dfs( s );
        REPD( i, n, 2 ){
            int u = nfd[ i ];
            if( u == 0 ) continue;
            for( int v : pred[ u ] ) if( dfn[ v ] ){
                eval( v );
                if( cmp( sdom[ mn[ v ] ], sdom[ u ] ) )
                    sdom[ u ] = sdom[ mn[ v ] ];
            }
            cov[ sdom[ u ] ].push_back( u );
            mom[ u ] = par[ u ];
            for( int w : cov[ par[ u ] ] ){
                eval( w );
                if( cmp( sdom[ mn[ w ] ], par[ u ] ) )
                    idom[ w ] = mn[ w ];
                else idom[ w ] = par[ u ];
            }
            cov[ par[ u ] ].clear();
        }
        REP( i, 2, n ){
            int u = nfd[ i ];
            if( u == 0 ) continue;
            if( idom[ u ] != sdom[ u ] )
                idom[ u ] = idom[ idom[ u ] ];
        }
    } } domT;

```

5.11 K-th Shortest Path

```

// time: O(|E| \lg |E| + |V| \lg |V| + K)
// memory: O(|E| \lg |E| + |V|)
struct KSP{ // 1-base
    struct nd{
        int u, v; ll d;
        nd(int ui = 0, int vi = 0, ll di = INF)
            : u(ui), v(vi), d(di) {}
    };
};

```

```

struct heap{
    nd* edge; int dep; heap* chd[4];
};
static int cmp(heap* a, heap* b)
{ return a->edge->d > b->edge->d; }
struct node{
    int v; ll d; heap* H; nd* E;
    node(){
        node(ll _d, int _v, nd* _E)
            : d(_d), v(_v), E(_E) {}
        node(heap* _H, ll _d)
            : H(_H), d(_d) {}
        friend bool operator<(node a, node b)
            { return a.d > b.d; }
    };
    int n, k, s, t;
    ll dst[ N ];
    nd *nxt[ N ];
    vector<nd*> g[ N ], rg[ N ];
    heap *nullNd, *head[ N ];
    void init( int _n, int _k, int _s, int _t ){
        n = _n; k = _k; s = _s; t = _t;
        for( int i = 1; i <= n; i ++ ){
            g[ i ].clear(); rg[ i ].clear();
            nxt[ i ] = NULL; head[ i ] = NULL;
            dst[ i ] = -1;
        }
    }
    void addEdge( int ui, int vi, ll di ){
        nd* e = new nd(ui, vi, di);
        g[ ui ].push_back( e );
        rg[ vi ].push_back( e );
    }
    queue<int> dfsQ;
    void dijkstra(){
        while(dfsQ.size()) dfsQ.pop();
        priority_queue<node> Q;
        Q.push(node(0, t, NULL));
        while (!Q.empty()){
            node p = Q.top(); Q.pop();
            if(dst[p.v] != -1) continue;
            dst[ p.v ] = p.d;
            nxt[ p.v ] = p.E;
            dfsQ.push( p.v );
            for(auto e: rg[ p.v ] )
                Q.push(node(p.d + e->d, e->u, e));
        }
    }
    heap* merge(heap* curNd, heap* newNd){
        if(curNd == nullNd) return newNd;
        heap* root = new heap;
        memcpy(root, curNd, sizeof(heap));
        if(newNd->edge->d < curNd->edge->d){
            root->edge = newNd->edge;
            root->chd[2] = newNd->chd[2];
            root->chd[3] = newNd->chd[3];
            newNd->edge = curNd->edge;
            newNd->chd[2] = curNd->chd[2];
            newNd->chd[3] = curNd->chd[3];
        }
        if(root->chd[0]->dep < root->chd[1]->dep)
            root->chd[0] = merge(root->chd[0], newNd);
        else
            root->chd[1] = merge(root->chd[1], newNd);
        root->dep = max(root->chd[0]->dep, root->chd[1]->
            dep) + 1;
        return root;
    }
    vector<heap*> V;
    void build(){
        nullNd = new heap;
        nullNd->dep = 0;
        nullNd->edge = new nd;
        fill(nullNd->chd, nullNd->chd+4, nullNd);
        while(not dfsQ.empty()){
            int u = dfsQ.front(); dfsQ.pop();
            if(!nxt[ u ]) head[ u ] = nullNd;
            else head[ u ] = head[nxt[ u ]->v];
            V.clear();
            for( auto&& e : g[ u ] ){
                int v = e->v;
                if( dst[ v ] == -1 ) continue;
                e->d += dst[ v ] - dst[ u ];
                if( nxt[ u ] != e ){

```

```

    heap* p = new heap;
    fill(p->chd, p->chd+4, nullNd);
    p->dep = 1;
    p->edge = e;
    V.push_back(p);
} }
if(V.empty()) continue;
make_heap(V.begin(), V.end(), cmp);
#define L(X) ((X<<1)+1)
#define R(X) ((X<<1)+2)
for( size_t i = 0 ; i < V.size() ; i ++ ){
    if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
    else V[i]->chd[2]=nullNd;
    if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
    else V[i]->chd[3]=nullNd;
}
head[u] = merge(head[u], V.front());
} }
vector<ll> ans;
void first_K(){
    ans.clear();
    priority_queue<node> Q;
    if( dst[ s ] == -1 ) return;
    ans.push_back( dst[ s ] );
    if( head[s] != nullNd )
        Q.push(node(head[s], dst[s]+head[s]->edge->d));
    for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
        node p = Q.top(), q; Q.pop();
        ans.push_back( p.d );
        if(head[ p.H->edge->v ] != nullNd){
            q.H = head[ p.H->edge->v ];
            q.d = p.d + q.H->edge->d;
            Q.push(q);
        }
        for( int i = 0 ; i < 4 ; i ++ )
            if( p.H->chd[ i ] != nullNd ){
                q.H = p.H->chd[ i ];
                q.d = p.d - p.H->edge->d + p.H->chd[ i ]->
                    edge->d;
                Q.push( q );
            }
    }
}
void solve(){ // ans[i] stores the i-th shortest path
    dijkstra();
    build();
    first_K(); // ans.size() might less than k
} } solver;

```

5.12 Floryd Warshall

```

for( int k=0 ; k < n ; k++ )
    for( int i=0 ; i < n ; i++ )
        for( int j=0 ; j < n ; j++ )
            if( dis[i][j] > dis[i][k]+dis[k][j] && dis[i][k]
                < INF && dis[k][j] < INF )
                dis[i][j]=dis[i][k]+dis[k][j];
for( int i=0 ; i < n ; i++ )
    for( int j=0 ; j < n ; j++ )
        for( int k=0 ; k < n && dis[i][j] != negINF ; k++ )
            if( dis[k][k] < 0 && dis[i][k] != INF && dis[k]
                ][j] != INF )
                dis[i][j]=negINF;

```

5.13 Minimum Steiner Tree

```

// Minimum Steiner Tree 重要點的mst
//  $O(V^3 \log V + V^2 \log V)$ 
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
    int n , dst[V][V] , dp[1 << T][V] , tdst[V];
    void init( int _n ){
        n = _n;
        for( int i = 0 ; i < n ; i ++ ){
            for( int j = 0 ; j < n ; j ++ )
                dst[ i ][ j ] = INF;
            dst[ i ][ i ] = 0;

```

```

} }
void add_edge( int ui , int vi , int wi ){
    dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
}
void shortest_path(){ // using spfa may faster
    for( int k = 0 ; k < n ; k ++ )
        for( int i = 0 ; i < n ; i ++ )
            for( int j = 0 ; j < n ; j ++ )
                dst[ i ][ j ] = min( dst[ i ][ j ] ,
                    dst[ i ][ k ] + dst[ k ][ j ] );
} // call shorest_path before solve
int solve( const vector<int>& ter ){
    int t = (int)ter.size();
    for( int i = 0 ; i < ( 1 << t ) ; i ++ )
        for( int j = 0 ; j < n ; j ++ )
            dp[ i ][ j ] = INF;
    for( int i = 0 ; i < n ; i ++ )
        dp[ 0 ][ i ] = 0;
    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
        if( msk == ( msk & (-msk) ) ){
            int who = __lg( msk );
            for( int i = 0 ; i < n ; i ++ )
                dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
            continue;
        }
        for( int i = 0 ; i < n ; i ++ )
            for( int submsk = ( msk - 1 ) & msk ; submsk ;
                submsk = ( submsk - 1 ) & msk )
                dp[ msk ][ i ] = min( dp[ msk ][ i ] ,
                    dp[ submsk ][ i ] +
                    dp[ msk ^ submsk ][ i ] );
        for( int i = 0 ; i < n ; i ++ ){
            tdst[ i ] = INF;
            for( int j = 0 ; j < n ; j ++ )
                tdst[ i ] = min( tdst[ i ] ,
                    dp[ msk ][ j ] + dst[ j ][ i ] );
        }
        for( int i = 0 ; i < n ; i ++ )
            dp[ msk ][ i ] = tdst[ i ];
    }
    int ans = INF;
    for( int i = 0 ; i < n ; i ++ )
        ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
    return ans;
} } solver;

```

5.14 虛樹

```

vector<int> virTree(vector<int> ver, LCA &lca) {
    auto cmp = [&](int u, int v){return time_in[u] <
        time_in[v];};
    sort(ver.begin(), ver.end(), cmp); //用dfn排序
    vector<int> res(ver.begin(), ver.end());
    for(int i = 1; i < ver.size(); i++){
        res.push_back(lca.getLCA(ver[i-1], ver[i])); //把
        LCA丟進虛樹內
    }
    sort(res.begin(), res.end(), cmp); //再用dfn排序
    res.erase(unique(res.begin(), res.end()), res.end());
    //丟掉重複的點
    return res;
}

```

5.15 Tree Hash

```

map<vector<int>, int> id;
int dfs(int x, int f){
    vector<int> sub;
    for( int v : edge[x] ){
        if (v != f)
            sub.push_back(dfs(v, x));
    }
    sort(sub.begin(), sub.end());
    if (!id.count(sub))
        id[sub] = id.size();
    return id[sub];
}

```

5.16 HeavyLightDecomposition

```
// 詢問,修改複雜度  $O(\log^2 n)$ 
// 1-base

int sz[MXN], dep[MXN], son[MXN], fa[MXN];

// 第一次 dfs
// 找重兒子 需要紀錄當前節點的子樹大小(sz)、深度(dep)、
// 重兒子(son)、父節點(fa)
// 沒有子節點 son[x] = 0
void dfs_sz(int x, int f, int d) { //當前節點 x · 父節點 f · 深度 d
    sz[x] = 1; dep[x] = d; fa[x] = f;
    for(int i : edge[x]) {
        if(i == f) continue;
        dfs_sz(i, x, d+1);
        sz[x] += sz[i];
        if(sz[son[x]] < sz[i]) son[x] = i;
    }
}

// 第二次 dfs
int top[MXN]; // 每個節點所在的鏈的頂端節點
int dfn[MXN]; // 節點編號,編號為在線段樹上的位置
int rnk[MXN]; // 編號為哪個節點
int bottom[MXN]; // 維護每個節點的子樹中最大 dfn 編號
int cnt = 0;
int dfs_hld(int x, int f){
    top[x] = (son[fa[x]] == x ? top[fa[x]] : x);
    rnk[cnt] = x;
    bottom[x] = dfn[x] = cnt++;
    if(son[x]) bottom[x] = max(bottom[x], dfs_hld(son[x], x)); // 更新子樹最大編號
    for(int i : edge[x]){
        if(i == f || i == son[x]) continue;
        bottom[x] = max(bottom[x], dfs_hld(i, x)); // 更新子樹最大編號
    }
    return bottom[x];
}

// 求出 lca
// 不斷跳鏈 · 直到 u,v 跳到同一條鏈上為止
// 每次跳鏈選所在的鏈頂端深度較深的一端往上跳
int getLca(int u, int v) {
    while(top[u] != top[v]){
        if(dep[top[u]] > dep[top[v]])
            u = fa[top[u]];
        else
            v = fa[top[v]];
    }
    return dep[u] > dep[v] ? v : u;
}

// 路徑權重總和
int query(int u, int v) {
    int ret = 0;
    while(top[u] != top[v]){
        if(dep[top[u]] > dep[top[v]]){
            ret += segtree.query(dfn[top[u]], dfn[u]);
            u = fa[top[u]];
        }
        else{
            ret += segtree.query(dfn[top[v]], dfn[v]);
            v = fa[top[v]];
        }
    }
    // 最後到同一條鏈上
    ret += segtree.query(min(dfn[u], dfn[v]), max(dfn[u], dfn[v]));
    return ret;
}
```

5.17 Graph Thearom

- 差分約束條件:
約束條件 $V_j - V_i \leq W$ addEdge(V_i, V_j, W) and run bellman-ford or spfa

- 龜兔賽跑演算法:
開始賽跑, 兔子一次走兩格、烏龜一次走一格直到他們相遇停止
此時讓兔子返回起始點, 兩者以相同走一格的速度繼續前進, 他們就會在環入口會合
- 2-SAT 條件:
滿足 $(x_1 \text{ or } y_1) \text{ and } (x_2 \text{ or } y_2) \text{ and } \dots$ 對於一個限制 $(x \text{ or } y)$ · 則加兩條邊 $x \rightarrow \neg y, y \rightarrow \neg x$

6 String

6.1 PalTree $O(n)$

```
// state[i]代表第i個字元為結尾的最長回文編號
// len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文字字串在整個字串中的出現次數
// fail[s]是他長度次長的回文後綴 · aba的fail是a
const int MXN = 1000010;
struct PalT{
    int nxt[MXN][26], fail[MXN], len[MXN];
    int tot, lst, n, state[MXN], cnt[MXN], num[MXN];
    int diff[MXN], sfail[MXN], fac[MXN], dp[MXN];
    char s[MXN] = {"-1"};
    int newNode(int l, int f){
        len[tot] = l, fail[tot] = f, cnt[tot] = num[tot] = 0;
        memset(nxt[tot], 0, sizeof(nxt[tot]));
        diff[tot] = (l > 0 ? l - len[f] : 0);
        sfail[tot] = (l > 0 && diff[tot] == diff[f] ? sfail[f] : f);
        return tot++;
    }
    int getfail(int x){
        while(s[n-len[x]-1] != s[n]) x = fail[x];
        return x;
    }
    int getmin(int v){
        dp[v] = fac[n-len[sfail[v]]-diff[v]];
        if(diff[v] == diff[fail[v]])
            dp[v] = min(dp[v], dp[fail[v]]);
        return dp[v]+1;
    }
    int push(){
        int c = s[n] - 'a', np = getfail(lst);
        if(!(lst = nxt[np][c])){
            lst = newNode(len[np]+2, getfail(fail[np]));
            nxt[np][c] = lst; num[lst] = num[fail[lst]]+1;
        }
        fac[n] = n;
        for(int v = lst; len[v] > 0; v = sfail[v])
            fac[n] = min(fac[n], getmin(v));
        return ++cnt[lst], lst;
    }
    void init(const char *_s){
        tot = lst = n = 0;
        newNode(0, 1), newNode(-1, 1);
        for(; s[n];) s[n+1] = _s[n], ++n, state[n-1] = push();
        for(int i = tot-1; i > 1; i--) cnt[fail[i]] += cnt[i];
    }
}palt;
```

6.2 Longest Increasing Subsequence

```
vector<int> getLIS(vector<int> a){
    vector<int> lis;
    for(int i : a){
        if(lis.empty() || lis.back() < i) lis.push_Back(i);
        else *lower_bound(lis.begin(), lis.end(), i) = i;
    }
    return lis;
}
```

6.3 Longest Common Subsequence $O(n \lg n)$

```
int LCS(string& s1, string& s2) {
    vector<int> p[128]; // 假設字元範圍為 0 ~ 127
    for (int i = 0; i < s2.size(); ++i) p[s2[i]].
        push_back(i);
    vector<int> v;
    v.push_back(-1);

    for (int i = 0; i < s1.size(); ++i)
        for (int j = p[s1[i]].size() - 1; j >= 0; --j) {
            int n = p[s1[i]][j];

            if (n > v.back())
                v.push_back(n);
            else
                *lower_bound(v.begin(), v.end(), n) = n;
        }
    return v.size() - 1;
};
```

6.4 KMP

/* len-failure[k]:
在k結尾的情況下，這個子字串可以由開頭
長度為(len-failure[k])的部分重複出現來表達

failure[k]為次長相同前綴後綴
如果我們不只想求最多，而且以0-base做為考量
，那可能的長度由大到小會是
failure[k]、failure[failuer[k]-1]
、failure[failure[failuer[k]-1]-1]..
直到有值為0為止 */
int failure[MXN];
vector<int> KMP(string& t, string& p) {
 vector<int> ret;
 if(p.size() > t.size()) return ret;
 for(int i = 1, j = failure[0] = -1; i < p.size(); i
 ++){
 while(j >= 0 && p[j + 1] != p[i]) j = failure[j
];
 if(p[j + 1] == p[i]) j++;
 failure[i] = j;
 }
 for(int i = 0, j = -1; i < t.size(); i++) {
 while (j >= 0 && p[j + 1] != t[i]) j = failure[j
];
 if(p[j + 1] == t[i]) j++;
 if(j == p.size() - 1) {
 ret.push_back(i - p.size() + 1);
 j = failure[j];
 }
 }
 return ret;
}

6.5 SAIS $O(n)$

/* ** SA · 將字串的所有後綴排序後的數組 ** */
/* SA[i]儲存排序後第i小的後綴從哪裡開始 */
/* ** H[i]為第i小的字串跟第i-1小的LCP ** */
/* ** 註：LCP(Longest Common Prefix) ** */
/* ** ex:S = "babd", SA[0] = 1("abd") ** */
/* ** SA[1] = 0("babd"), SA[2] = 2("bd") ** */
/* ** H[0] = 0, H[1] = 0, H[2] = 1("b") ** */
/* 傳入參數:ip 陣列放字串，len為字串長度 */
/* 需保證ip[len]為0，且字串裡的元素不為0 */
const int N = 300010;
struct SA{
#define REP(i,n) for (int i=0; i<int(n); i++)
#define REP1(i,a,b) for (int i=(a); i<=int(b); i++)
 bool _t[N*2];
 int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
 hei[N], r[N];
 int operator [] (int i){ return _sa[i]; }
 void build(int *s, int n, int m){
 memcpy(_s, s, sizeof(int) * n);
 sais(_s, _sa, _p, _q, _t, _c, n, m);
 mkhei(n);
 }

```
}  
void mkhei(int n){  
    REP(i,n) r[_sa[i]] = i;  
    hei[0] = 0;  
    REP(i,n) if(r[i]) {  
        int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;  
        while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;  
        hei[r[i]] = ans;  
    }  
}  
void sais(int *s, int *sa, int *p, int *q, bool *t,  
    int *c, int n, int z){  
    bool uniq = t[n-1] = true, neq;  
    int nn = 0, nmzx = -1, *nsa = sa + n, *ns = s + n,  
        lst = -1;  
#define MS0(x,n) memset((x),0,n*sizeof(*(x)))  
#define MAGIC(XD) MS0(sa, n); \  
    memcpy(x, c, sizeof(int) * z); \  
    XD; \  
    memcpy(x + 1, c, sizeof(int) * (z - 1)); \  
    REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i]  
        ]-1]]++ = sa[i]-1; \  
    memcpy(x, c, sizeof(int) * z); \  
    for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i]  
        ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;  
    MS0(c, z);  
    REP(i,n) uniq &= ++c[s[i]] < 2;  
    REP(i,z-1) c[i+1] += c[i];  
    if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }  
    for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i  
        +1] ? t[i+1] : s[i]<s[i+1]);  
    MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i  
        ]]] = p[q[i]=nn++] = i);  
    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {  
        neq=lst<0||memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa  
            [i])*sizeof(int));  
        ns[q[lst=sa[i]]]=nmzx+=neq;  
    }  
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmzx  
        + 1);  
    MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[s[p[  
        nsa[i]]]]] = p[nsa[i]]);  
    }  
}sa;  
int H[ N ], SA[ N ];  
void suffix_array(int* ip, int len) {  
    // should padding a zero in the back  
    // ip is int array, len is array length  
    // ip[0..n-1] != 0, and ip[len] = 0  
    ip[len++] = 0;  
    sa.build(ip, len, 128);  
    for (int i=0; i<len; i++) {  
        H[i] = sa.hei[i + 1];  
        SA[i] = sa._sa[i + 1];  
    }  
    // resulting height, sa array \in [0,len)  
}
```

6.6 Z Value $O(n)$

```
//z[i] = lcp(s[1...n-1],s[i...n-1])  
int z[MXN];  
void Z_value(const string& s) {  
    int i, j, left, right, len = s.size();  
    left=right=0; z[0]=len;  
    for(i=1;i<len;i++) {  
        j=max(min(z[i-left],right-i),0);  
        for(;i+j<len&&s[i+j]==s[j];j++);  
        z[i]=j;  
        if(i+z[i]>right) {  
            right=i+z[i];  
            left=i;  
        }  
    }  
}
```

6.7 Manacher Algorithm $O(n)$

```
// 求以每個字元為中心的最長回文半徑  
// 頭尾以及每個字元間都加入一個
```



```
// 沒出現過的字元 · 這邊以' @' 為例
// s為傳入的字串 · len為字串長度
// z為儲存答案的陣列 (有包含'@'要小心)
// ex: s = "abaac" -> "@a@b@a@a@c@"
// z = [12141232121]
void z_value_pal(char *s,int len,int *z){
    len=(len<<1)+1;
    for(int i=len-1;i>=0;i--){
        s[i]=i&1?s[i>>1]:'@';
        z[0]=1;
        for(int i=1,l=0,r=0;i<len;i++){
            z[i]=i<r?min(z[l+l-i],r-i):1;
            while(i-z[i]>=0&&i+z[i]<len&&s[i-z[i]]==s[i+z[i]])
                ++z[i];
            if(i+z[i]>r) l=i,r=i+z[i];
        }
    }
}
```

6.8 Smallest Rotation

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
    int a = 0, N = s.size(); s += s;
    rep(b,0,N) rep(k,0,N) {
        if(a+k == b || s[a+k] < s[b+k])
            {b += max(0, k-1); break;}
        if(s[a+k] > s[b+k]) {a = b; break;}
    } return a;
}
```

6.9 Cyclic LCS

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]={0,-1, -1,-1, -1,0};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
    int i=r+al,j=bl,l=0;
    while(i>r) {
        char dir=pred[i][j];
        if(dir==LU) l++;
        i+=mov[dir][0];
        j+=mov[dir][1];
    }
    return l;
}
inline void reroot(int r) { // r = new base row
    int i=r,j=1;
    while(j<=bl&&pred[i][j]!=LU) j++;
    if(j>bl) return;
    pred[i][j]=L;
    while(i<2*al&&j<=bl) {
        if(pred[i+1][j]==U) {
            i++;
            pred[i][j]=L;
        } else if(j<bl&&pred[i+1][j+1]==LU) {
            i++;
            j++;
            pred[i][j]=L;
        } else {
            j++;
        }
    }
}
int cyclic_lcs() {
    // a, b, al, bl should be properly filled
    // note: a WILL be altered in process
    // -- concatenated after itself
    char tmp[MAXL];
    if(al>bl) {
        swap(al,bl);
        strcpy(tmp,a);
        strcpy(a,b);
        strcpy(b,tmp);
    }
    strcpy(tmp,a);
    strcat(a,tmp);
}
```

```
// basic lcs
for(int i=0;i<=2*al;i++) {
    dp[i][0]=0;
    pred[i][0]=U;
}
for(int j=0;j<=bl;j++) {
    dp[0][j]=0;
    pred[0][j]=L;
}
for(int i=1;i<=2*al;i++) {
    for(int j=1;j<=bl;j++) {
        if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
        else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
        if(dp[i][j-1]==dp[i-1][j]) pred[i][j]=L;
        else if(a[i-1]==b[j-1]) pred[i][j]=LU;
        else pred[i][j]=U;
    }
}
// do cyclic lcs
int clcs=0;
for(int i=0;i<al;i++) {
    clcs=max(clcs,lcs_length(i));
    reroot(i+1);
}
// recover a
a[al]='\0';
return clcs;
}
```

6.10 Hash

```
//字串雜湊前的idx是0-base · 雜湊後為1-base
//即區間為 [0,n-1] -> [1,n]
//若要取得區間[L,R]的值得則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為len的字串和
//(h[i+len-1] - h[i-1] * qpow(p, len) % mod1 + mod1)
//從j開始長度為len的字串是否相同
#define x first
#define y second
pair<int,int> Hash[MXN];
void build(const string& s){
    pair<int,int> val = make_pair(0,0);
    Hash[0]=val;
    for(int i=1; i<=s.size(); i++){
        val.x = (val.x * P1 + s[i-1]) % MOD;
        val.y = (val.y * P2 + s[i-1]) % MOD;
        Hash[i] = val;
    }
}
bool cmp( int i, int j, int len ) {
    return ((Hash[i+len-1].x-Hash[i-1].x*qpow(P1,len)%MOD+MOD)%MOD == (Hash[j+len-1].x-Hash[j-1].x*qpow(P1,len)%MOD+MOD)%MOD)
    && ((Hash[i+len-1].y-Hash[i-1].y*qpow(P2,len)%MOD+MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow(P2,len)%MOD+MOD)%MOD);
}
```

7 Data Structure

7.1 Segment tree

```
// !!!注意build()時初始化用的陣列也是1-base
// !!!query(0, 0) 會報錯
#define cl(x) (x*2)
#define cr(x) (x*2+1)
struct segmentTree {
    int n;
    vector<int> seg, tag, cov;
    segmentTree(int _n): n(_n) {
        seg = tag = cov = vector<int>(n * 4, 0);
    }
    void push(int i, int L, int R) {
        if(cov[i]) {

```

```

        seg[i] = cov[i] * (R - L + 1);
        if(L < R) {
            cov[cl(i)] = cov[cr(i)] = cov[i];
            tag[cl(i)] = tag[cr(i)] = 0;
        }
        cov[i] = 0;
    }
    if(tag[i]) {
        seg[i] += tag[i] * (R - L + 1);
        if(L < R) {
            tag[cl(i)] += tag[i];
            tag[cr(i)] += tag[i];
        }
        tag[i] = 0;
    }
}

void pull(int i, int L, int R) {
    if(L >= R) return;
    int mid = L + R >> 1;
    push(cl(i), L, mid);
    push(cr(i), mid + 1, R);
    seg[i] = seg[cl(i)] + seg[cr(i)];
}

void build(vector<int>& arr, int i = 1, int L = 1,
int R = -1) {
    if(R == -1) R = n;
    if(L == R) return void(seg[i] = arr[L]);
    int mid = L + R >> 1;
    build(arr, cl(i), L, mid);
    build(arr, cr(i), mid + 1, R);
    pull(i, L, R);
}

int query(int rL, int rR, int i = 1, int L = 1, int
R = -1) {
    if(R == -1) R = n;
    push(i, L, R);
    if(rL <= L && R <= rR) return seg[i];
    int mid = L + R >> 1, ret = 0;
    if(rL <= mid) ret += query(rL, rR, cl(i), L,
mid);
    if(mid < rR) ret += query(rL, rR, cr(i), mid +
1, R);
    return ret;
}

void update(int rL, int rR, int val, int i = 1, int
L = 1, int R = -1) {
    if(R == -1) R = n;
    push(i, L, R);
    if(rL <= L && R <= rR) return void(tag[i] = val
);
    int mid = L + R >> 1;
    if(rL <= mid) update(rL, rR, val, cl(i), L, mid
);
    if(mid < rR) update(rL, rR, val, cr(i), mid +
1, R);
    pull(i, L, R);
}

void cover(int rL, int rR, int val, int i = 1, int
L = 1, int R = -1) {
    if(R == -1) R = n;
    push(i, L, R);
    if(rL <= L && R <= rR) return void(cov[i] = val
);
    int mid = L + R >> 1;
    if(rL <= mid) cover(rL, rR, val, cl(i), L, mid
);
    if(mid < rR) cover(rL, rR, val, cr(i), mid +
1, R);
    pull(i, L, R);
}
};

/* Test Case:
4
1 2 3 4
5
2 1 3
1 1 3 1
2 1 3
1 1 4 1
2 1 4
*/

```

7.2 持久化 SMT

```

struct node{
    node *l, *r;
    int val;
};

vector<node*> ver;
int arr[MXN] = {0};

//0-base
struct SegmentTree{
    int n;
    node *root;
    void build(int _n){
        n = _n;
        root = build(0, n-1);
    }
    node* build(int L, int R){
        node *x = new node();
        if(L == R){ x->val = arr[L]; return x;}
        int mid = (L+R)/2;
        x->l = build(L, mid);
        x->r = build(mid + 1, R);
        x->val = x->l->val + x->r->val;
        return x;
    }
    int query(node *ro, int L, int R){return query(ro, 0,
n-1, L, R);}
    int query(int L, int R){return query(root, 0, n-1, L,
R);}
    int query(node *x, int L, int R, int recL, int recR){
        if(recL <= L && R <= recR) return x->val;
        int mid = (L+R)/2, res = 0;
        if(recL <= mid) res += query(x->l, L, mid, recL,
recR);
        if(mid < recR) res += query(x->r, mid+1, R, recL,
recR);
        return res;
    }
    void update(int pos, int v){update(root, 0, n-1, pos,
v);}
    void update(node *x, int L, int R, int pos, int v){
        if(L == R){ x->val = v; arr[L] = v; return;}
        int mid = (L+R)/2;
        if(pos <= mid) update(x->l, L, mid, pos, v);
        else update(x->r, mid+1, R, pos, v);
        x->val = x->l->val + x->r->val;
    }
    node *update_ver(node *pre, int l, int r, int pos,
int v){
        node *x = new node(); //當前位置建立新節點
        if(l == r){
            x->val = v;
            return x;
        }
        int mid = (l+r)>>1;
        if(pos <= mid){ //更新左邊
            x->l = update_ver(pre->l, l, mid, pos, v); //左邊
            節點連向新節點
            x->r = pre->r; //右邊連到原本的右邊
        }
        else{ //更新右邊
            x->l = pre->l; //左邊連到原本的左邊
            x->r = update_ver(pre->r, mid+1, r, pos, v); //
            右邊節點連向新節點
        }
        x->val = x->l->val + x->r->val;
        return x;
    }
} seg;

void add_ver(int x, int v){ //修改位置 x 的值為 v
    ver.push_back(seg.update_ver(ver.back(), 0, seg.n
-1, x, v));
}

```

7.3 持久化並查集

```

struct DSU {

```

```

int n;
vector<int> fa, sz;
vector<tuple<int, int, int, int>> ver;
DSU(int _n): n(_n), fa(n), sz(n, 1) {
    iota(fa.begin(), fa.end(), 0);
}
int find(int x) {
    return fa[x] == x ? x : find(fa[x]);
}
void merge(int x, int y) {
    x = find(x), y = find(y);
    if(sz[x] < sz[y]) swap(x, y);
    ver.push_back({x, sz[x], y, fa[y]});
    if(x == y) return;
    sz[x] += sz[y];
    fa[y] = x;
}
void undo() {
    if(ver.empty()) return;
    auto [x, szx, y, fy] = ver.back();
    ver.pop_back();
    sz[x] = szx;
    fa[y] = fy;
}
};

```

7.4 Trie

```

struct trie{
    trie *nxt[26];
    int cnt; //紀錄有多少個字串以此節點結尾
    int sz; //有多少字串的前綴包括此節點
    trie():cnt(0),sz(0){
        memset(nxt,0,sizeof(nxt));
    }
};

trie *root = new trie(); //創建新的字典樹

void insert(string& s){
    trie *now = root; //每次從根結點出發
    for(auto i:s){
        now->sz++;
        if(now->nxt[i-'a'] == NULL){
            now->nxt[i-'a'] = new trie();
        }
        now = now->nxt[i-'a']; //走到下一個字母
    }
    now->cnt++; now->sz++;
}

int query_prefix(string& s){ //查詢有多少前綴為 s
    trie *now = root; //每次從根結點出發
    for(auto i:s){
        if(now->nxt[i-'a'] == NULL){
            return 0;
        }
        now = now->nxt[i-'a'];
    }
    return now->sz;
}

int query_count(string& s){ //查詢字串 s 出現次數
    trie *now = root; //每次從根結點出發
    for(auto i:s){
        if(now->nxt[i-'a'] == NULL){
            return 0;
        }
        now = now->nxt[i-'a'];
    }
    return now->cnt;
}

```

7.5 Treap (interval reverse)

//拆出[a,b]區間就如同下面所展示先使用splitByTh()拆出
//左右,再把左區間拆成l, m最後merge()回去

```

//反轉區間時又記得使用^=可以直接反轉01

//treap拆區間時從後面拆是因為這樣[a,b]的關係
//不用重新考慮·要是先拆前面b的位置會變成b-a+1
//0-base
//splitByTh(root,a-1,l,m);
//splitByTh(m,b-a+1,m,r);

mt19937 gen(chrono::steady_clock::now().
    time_since_epoch().count());
struct Treap {
    int key, pri, sz, tag, sum;
    Treap *L, *R;
    Treap( int val ) {
        sum=key=val, pri=gen(), sz=1, tag=0;
        L=R=NULL;
    };
};
int Size( Treap *a ) { return !a?0:a->sz; }
void pull( Treap *a ) {
    a->sz=Size(a->L)+Size(a->R)+1;
    a->sum=a->key;
    if( a->L ) a->sum+=a->L->sum;
    if( a->R ) a->sum+=a->R->sum;
}
void push( Treap *a ) {
    if( a && a->tag ) {
        swap(a->L,a->R);
        if( a->L ) a->L->tag^=1;
        if( a->R ) a->R->tag^=1;
        a->tag=0;
    }
}
Treap *merge(Treap *a, Treap *b) {
    if( !a || !b ) return a?b;
    push(a), push(b);
    if( a->pri > b->pri ) {
        a->R=merge(a->R,b);
        pull(a); return a;
    }
    b->L=merge(a,b->L);
    pull(b); return b;
}
void print(Treap *a) {
    if( !a ) return;
    push(a);
    print(a->L);
    cout.put(a->key);
    print(a->R);
}
Treap *buildTreap( int n, string& str ) {
    Treap *root=NULL;
    for( int i=0; i < n; i++ )
        root=merge(root,new Treap(str[i]));
    return root;
}
void splitbyk( Treap *x, int k, Treap *&a, Treap *&b )
{
    if(!x) a=b=NULL;
    else if( x->key <= k ) {
        a=x;
        splitbyk(x->R,k,a->R,b);
        pull(a);
    }
    else {
        b=x;
        splitbyk(x->L,k,a,b->L);
        pull(b);
    }
}
void splitByTh( Treap *x, int k, Treap *&a, Treap *&b )
{
    if( !x ) { a=b=NULL; return; }
    push(x);
    if( Size(x->L)+1 <= k ) {
        a=x;
        splitByTh(x->R,k-Size(x->L)-1,a->R,b);
        pull(a);
    }
    else {
        b=x;
        splitByTh(x->L,k,a,b->L);
        pull(b);
    }
}

```

```

}
}
signed main() {
    string str;
    int n, m;
    cin >> n >> m >> str;
    Treap *root;
    root = buildTreap(n, str);
    for (int i = 0; i < m; i++) {
        int a, b;
        cin >> a >> b;
        Treap *l, *m, *r;
        splitByTh(root, b, l, r);
        splitByTh(l, a - 1, l, m);
        m->tag ^= 1;
        root = merge(l, merge(m, r));
    }
    print(root);
}

```

7.6 BIT

```

#define lowbit(x) (x & -x)
struct BIT {
    int n;
    vector<int> bit;
    BIT(int _n): n(_n), bit(_n + 1), C(_n + 1) {}
    void update(int x, int val) {
        for (; x <= n; x += lowbit(x)) bit[x] += val;
    }
    void update(int L, int R, int val) {
        update(L, val), update(R + 1, -val);
    }
    int query(int x) {
        int res = 0;
        for (; x; x -= lowbit(x)) res += bit[x];
        return res;
    }
    int query(int L, int R) {
        return query(R) - query(L - 1);
    }
    int getMax(int l, int r) {
        int ans = 0;
        while (l <= r) {
            ans = max(ans, bit[r--]);
            for (; l <= r - lowbit(r); r -= lowbit(r))
                ans = max(ans, C[r]);
        }
        return ans;
    }
    int kth(int k) {
        int sum = 0, x = 0;
        for (int i = __lg(n); ~i; i--) {
            x += 1 << i;
            if (x >= n || sum + bit[x] >= k) x -= 1 << i;
            else sum += bit[x];
        }
        return x + 1;
    }
};

```

7.7 Black Magic

```

#include <bits/extc++.h>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update> set_t;
tree<int, null_type, less_equal<int>, rb_tree_tag,
    tree_order_statistics_node_update> mt_t;
#include <ext/pb_ds/assoc_container.hpp>
typedef cc_hash_table<int, int> umap_t;
// gp_hash_table<int, int>
typedef priority_queue<int> heap;
#include <ext/rope>
using namespace __gnu_cxx;
int main() {
    // Insert some entries into s.

```

```

set_t s; s.insert(12); s.insert(505);
// The order of the keys should be: 12, 505.
assert(*s.find_by_order(0) == 12);
assert(*s.find_by_order(3) == 505);
// The order of the keys should be: 12, 505.
assert(s.order_of_key(12) == 0);
assert(s.order_of_key(505) == 1);
// Erase an entry.
s.erase(12);
// The order of the keys should be: 505.
assert(*s.find_by_order(0) == 505);
// The order of the keys should be: 505.
assert(s.order_of_key(505) == 0);
// if we want to delete less_equal tag tree
mt_t.erase(mt_t.find_by_order(mt_t.order_of_key(val))
    );

heap h1, h2; h1.join(h2);

rope<char> r[2];
r[1] = r[0]; // persistenet
string t = "abc";
r[1].insert(0, t.c_str());
r[1].erase(1, 1);
cout << r[1].substr(0, 2);
}

```

8 Others

8.1 SOS dp

```

for (int i = 0; i < (1 << N); ++i)
    F[i] = A[i];
for (int i = 0; i < N; ++i) for (int mask = 0; mask < (1 << N); ++mask) {
    if (mask & (1 << i))
        F[mask] += F[mask ^ (1 << i)];
}

```

8.2 De Bruijn sequence

```

// return cyclic array of length k^n such that every
// array of length n using 0~k-1 appears as a subarray.
vector<int> DeBruijn(int k, int n) {
    if (k == 1) return {0};
    vector<int> aux(k * n), res;
    function<void(int, int)> f = [&](int t, int p) -> void {
        if (t > n) { if (n % p == 0)
            for (int i = 1; i <= p; ++i) res.push_back(aux[i]);
        } else {
            aux[t] = aux[t - p]; f(t + 1, p);
            for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t]) f(t + 1, t);
        }
    };
    f(1, 1); return res;
}

```

8.3 CDQ 分治

```

// cdq分治使用的結構u, v, w為排序物的三個維度
// ans記錄了有幾項三維都小於等於自己
// cnt記錄了相同物有幾個。在使用cdq之前必先去重。
// 並且將相同元素紀錄至cnt中。可使用map來做到這步
// cdq使用的BIT就是普通求和的BIT。大小就開維度的
// 值域範圍。若值域大於2e6則要先進行離散化
struct triple {int u, v, w, ans, cnt;};
BIT *bt;
void cdq(int L, int R, vector<triple>& arr) {
    if (R - L <= 1) return;
    int mid = L + R >> 1;
    vector<triple> temp;
    cdq(L, mid, arr), cdq(mid, R, arr);
    for (int i = L, j = mid; i < mid || j < R;) {

```

```

    for(; i < mid && (j >= R || arr[i].v <= arr[j].v); i++) {
        bt->update(arr[i].w, arr[i].cnt);
        temp.push_back(arr[i]);
    }
    if(j < R) {
        arr[j].ans += bt->query(arr[j].w);
        temp.push_back(arr[j]);
        j++;
    }
}
for(int i = L; i < mid; i++)
    bt->update(arr[i].w, -arr[i].cnt);
copy(temp.begin(), temp.end(), arr.begin() + L);
}
signed main()
{
    // n 個數 k 值域範圍
    int n, k;
    cin >> n >> k;
    map<tuple<int, int, int>, int> mp;
    vector<int> res(n, 0);
    vector<triple> arr;
    bt = new BIT(k + 1);
    for(int i = 0; i < n; i++) {
        int x, y, z;
        cin >> x >> y >> z;
        mp[{x, y, z}]++;
    }
    for(auto t : mp)
        arr.push_back({get<0>(t.first), get<1>(t.first),
            get<2>(t.first), 0, t.second});
    cdq(0, arr.size(), arr);
    for(auto &[x,y,z,a,b] : arr) res[a + b - 1] += b;
    for(int i : res) cout << i << '\n';
}

```

8.4 3D LIS

```

#define lowbit(x) (x&-x)
const int MAXN=1e5+5;
struct BIT {
    int n;
    vector<int> bit;
    BIT(int _n):n(_n), bit(_n+1,0) {}
    int query(int x) {
        int res=0;
        for(; x > 0; x-=lowbit(x)) res=max(res,bit[x]);
        return res;
    }
    void update(int x, int val) {
        for(; x <= n; x+=lowbit(x)) {
            if(val < 0) bit[x]=0;
            else bit[x]=max(bit[x],val);
        }
    }
}bt(MAXN);
struct triple {
    int u, v, w, ans, cnt;
    bool operator<(triple b) { return u<b.u; }
};
bool cmp(triple a, triple b) {return a.v<b.v;}
void cdq(int L, int R, vector<triple>& arr) {
    if(R-L <= 1) return;
    int mid=L+R>>1;
    cdq(L,mid,arr);
    sort(arr.begin()+L,arr.begin()+mid,cmp);
    sort(arr.begin()+mid,arr.begin()+R,cmp);
    for(int i=L, j=mid; i < mid || j < R; ) {
        for(; i < mid && (j >= R || arr[i].v < arr[j].v) ; i++) bt.update(arr[i].w, arr[i].ans);
        if(j < R) {
            arr[j].ans=max(bt.query(arr[j].w-1)+1, arr[j].ans);
            j++;
        }
    }
    for(int i=L; i < mid; i++) bt.update(arr[i].w, -1);
    sort(arr.begin()+L,arr.begin()+R);
}

```

```

    cdq(mid,R,arr);
}
signed main()
{
    ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    ;
    int n, res=0;
    cin>>n;
    vector<int> ls;
    vector<triple> arr;
    for(int i=0; i < n; i++) {
        int a, b;
        cin>>a>>b;
        arr.push_back({i,a,b,1,1}); //{第一維,第二維,第三維,
            答案,數量}
        ls.push_back(b);
    }
    sort(ls.begin(),ls.end());
    ls.resize(unique(ls.begin(),ls.end())-ls.begin());
    for(auto &t : arr) t.w=lower_bound(ls.begin(),ls.
        end(),t.w)-ls.begin()+1;
    n=arr.size();
    cdq(0,n,arr);
    for(int i=0; i < n; i++) res=max(res,arr[i].ans);
    cout<<res<<'\n';
}

```

8.5 Ternary Search

```

while(L <= R) {
    int ml = L + (R - L) / 3, mr = R - (R - L) / 3;
    if(L == R) return L;
    else if(checker(ml) < checker(mr)) L = ml + 1;
    else R = mr - 1;
}

```

8.6 Max Subrectangle

```

const int N = 1e5+5;
int n, a[N], l[N], r[N];
long long ans;
int main() {
    while(cin>>n) {
        ans = 0;
        for(int i = 1; i <= n; i++) cin>>a[i], l[i] = r[i] = i;
        for(int i = 1; i <= n; i++)
            while(l[i] > 1 && a[i] <= a[l[i] - 1]) l[i] = l[l[i] - 1];
        for(int i = n; i >= 1; i--)
            while(r[i] < n && a[i] <= a[r[i] + 1]) r[i] = r[r[i] + 1];
        for(int i = 1; i <= n; i++)
            ans = max(ans, (long long)(r[i] - l[i] + 1) * a[i]);
        cout<<ans<<'\n';
    }
}

```

8.7 Maximal Rectangle

```

const int MXN = 300;
int maximalRectangle(vector<vector<char>>& matrix) {
    int a[MXN][], l[MXN][], r[MXN][];
    int n = matrix.size(), m = matrix[0].size(), ans = 0;
    for(int i = 1; i <= n; i++) {
        for(int j = 1; j <= m; j++) l[j] = r[j] = j;
        char c;
        for(int j = 1; j <= m; j++) { //對每一個直行做
            統計·若是上一個a[j]也是1則會變成2
            c = matrix[i - 1][j - 1];
            if(c == '1') a[j]++;
            else if(c == '0') a[j] = 0;
        }
    }
}

```



```

    for(int j = 1; j <= m; j++) while(l[j] != 1 &&
        a[l[j] - 1] >= a[j]) l[j] = l[l[j] - 1];
    for(int j = m; j >= 1; j--) while(r[j] != m &&
        a[r[j] + 1] >= a[j]) r[j] = r[r[j] + 1];
    for(int j = 1; j <= m; j++) ans = max(ans, (r[j]
        ] - l[j] + 1) * a[j]);
}
return ans;
}

```

8.8 p-Median

```

for(int i = 1; i <= n; i++) {
    for(int j = i; j <= n; j++) {
        dis[i][j] = 0;
        for(int k = i; k <= j; k++) dis[i][j] += abs(arr[k]
            - arr[i + j >> 1]);
        if(i == 1) dp[i][j] = dis[i][j];
    }
}
for(int i = 2; i <= p; i++) {
    for(int j = i; j <= n; j++) {
        dp[i][j] = INF;
        for(int k = i; k <= j; k++) {
            if(dp[i][j] > dp[i - 1][k - 1] + dis[k][j]) {
                dp[i][j] = dp[i - 1][k - 1] + dis[k][j];
                fa[i][j] = k - 1;
            }
        }
    }
}

```

8.9 Tree Knapsack

```

int dfs(int u) {
    int p = 1;
    dp[u][1] = s[u];
    for (int v : edge[u]) {
        int siz = dfs(v);
        for (int i = min(p, m + 1); i; i--)
            for (int j = 1; j <= siz && i + j <= m + 1; j++)
                dp[u][i + j] = max(dp[u][i + j], dp[u][i] + dp[
                    v][j]);
        p += siz;
    }
    return p;
}

```

8.10 質數個數

- 10^2 內有 25 個質數
- 10^3 內有 168 個質數
- 10^4 內有 1229 個質數
- 10^5 內有 9592 個質數
- 10^6 內有 78498 個質數
- 10^7 內有 664579 個質數
- 10^8 內有 5761455 個質數
- 10^9 內有 50847534 個質數
- 10^{12} 內有 37607912018 個質數
- 10^{18} 內有 24739954287740860 個質數

8.11 AC-Automaton

```

// 1-based
// n is the number of patterns
struct Automaton {
    static const int MXN = 1e6;
    int n, cnt, vis[MXN], rev[MXN], indeg[MXN], ans[MXN]
};

```

```

queue<int> q;
struct trie_node {
    vector<int> son;
    int fail, flag, ans;
    trie_node(): son(27), fail(0), flag(0) {}
} trie[MXN];
void init(int _n) {
    n = _n, cnt = 1;
    for (int i = 1; i <= n; i++) vis[i] = 0;
}
// insert a string s with number num
// num is the index of the pattern
void insert(string s, int num) {
    int u = 1, len = s.size();
    for (int i = 0; i < len; i++) {
        int v = s[i] - 'a';
        if (!trie[u].son[v]) trie[u].son[v] = ++cnt
            ;
        u = trie[u].son[v];
    }
    if (!trie[u].flag) trie[u].flag = num;
    rev[num] = trie[u].flag;
}
void getfail() {
    for (int i = 0; i < 26; i++) trie[0].son[i] =
        1;
    q.push(1);
    trie[1].fail = 0;
    while (q.size()) {
        int u = q.front(); q.pop();
        int Fail = trie[u].fail;
        for (int i = 0; i < 26; i++) {
            int v = trie[u].son[i];
            if (!v) {
                trie[u].son[i] = trie[Fail].son[i];
                continue;
            }
            trie[v].fail = trie[Fail].son[i];
            indeg[trie[Fail].son[i]]++;
            q.push(v);
        }
    }
}
void topu() {
    for (int i = 1; i <= cnt; i++)
        if (!indeg[i]) q.push(i);
    while (q.size()) {
        int fr = q.front(); q.pop();
        vis[trie[fr].flag] = trie[fr].ans;
        int u = trie[fr].fail;
        trie[u].ans += trie[fr].ans;
        if (!--indeg[u]) q.push(u);
    }
}
void query(string &s) {
    int u = 1, len = s.size();
    for (int i = 0; i < len; i++) u = trie[u].son[s
        [i] - 'a'], trie[u].ans++;
}
void solve(string &s) {
    getfail();
    query(s);
    topu();
    for (int i = 1; i <= n; i++) ans[i] = vis[rev[i]
        ]];
}
} AC;

```

