

# Compilers-1 - CS3510

## Toy Cool Programs

### Correct COOL Programs

Please see below some of the MIPS assembly language instructions

- 1) **li** : Load immediate value instruction
- 2) **sw** : store word
- 3) **move** : move the value stored in the address to registers
- 4) **addiu** : add the immediate unsigned without any overflow
- 5) **bne** : branch on equal
- 6) **str\_const(l)** it contains all the string literals in our code sections
- 7) Function calls use branch instructions which are in assembly language.
- 8) **la** : stand for local address
- 9) **\$sp** is stack pointer
- 10) **\$fp** is frame pointer
- 11) Labels are set of instructions to be executed which are in assembly language when we were directed to go there

### Program1:

#### Average of N numbers rounded off

Main.main:

```
addiu $sp $sp -32    #space on stack have been reserved and moving stack pointer
sw $fp 32($sp)       #save the frame pointer onto stack
sw $s0 28($sp)        #preserve $s0
sw $ra 24($sp)        #preserve the Returned address pointer
addiu $fp $sp 4
move $s0 $a0
la $a0 str_const0     #str_const0 is string represented by StringSymbol
sw $a0 0($sp)         #store word in register a0 to 0($sp) address
addiu $sp $sp -4
move $a0 $s0
bne $a0 $zero label0  #branch to label0 if a0 not equal to zero
la $a0 str_const5     #these statements get executed after label0 returns
li $t1 1
```

```
jal _dispatch_abort    #end program execution
```

## **Program 2:**

### **Factorial of a number N**

Main.fact:

addiu \$sp \$sp -20	#for the factorial function assembly code
sw \$fp 20(\$sp)	#reserving space on stack
sw \$s0 16(\$sp)	#preserve the frame pointer onto stack
sw \$ra 12(\$sp)	#preserve \$s0
addiu \$fp \$sp 4	#preserve the Returned address pointer
move \$s0 \$a0	#moving frame pointer by 4 bits
lw \$s1 20(\$fp)	#restore the value
la \$t2 int_const0	#the integer constant in local address
move \$t1 \$s1	#move the number read into \$t1
la \$a0 bool_const1	#checking with boolean constant
beq \$t1 \$t2 label2	#branch on equal to label2 if t1 = t2
la \$a1 bool_const0	
jal equality_test	#jump to label

## **Program 3**

### **GCD of 2 numbers:**

Main.gcd:

addiu \$sp \$sp -16	#this method was called in the main method
sw \$fp 16(\$sp)	#preserve the frame pointer onto stack
sw \$s0 12(\$sp)	
sw \$ra 8(\$sp)	
addiu \$fp \$sp 4	#moving frame pointer by 4 bits
move \$s0 \$a0	

Main.main:

addiu \$sp \$sp -12	#reserving space on stack
sw \$fp 12(\$sp)	#preserve the frame pointer onto stack
sw \$s0 8(\$sp)	
sw \$ra 4(\$sp)	
addiu \$fp \$sp 4	

```

move $s0 $a0
la    $a0 str_const0    #string constant in the local address
sw    $a0 0($sp)
addiu $sp $sp -4        #reserving space on stack
move  $a0 $s0
bne   $a0 $zero label7  #branch on not equal to label7 and execute instructions
la    $a0 str_const4
li    $t1 1
jal   _dispatch_abort   #jump to label and exit execution and abort is given

```

## **Program 4**

### **Pattern:**

Main.pattern:

```

addiu $sp $sp -28        #reserving space on stack
sw    $fp 28($sp)
sw    $s0 24($sp)
sw    $ra 20($sp)
addiu $fp $sp 4          #moving frame pointer by 4 bits
move  $s0 $a0
la    $s1 int_const0     #initialising two integers
la    $s2 int_const0     # labels follow after this which execute the instructions

```

Main.main:

```

addiu $sp $sp -16        #reserving space on stack
sw    $fp 16($sp)
sw    $s0 12($sp)
sw    $ra 8($sp)
addiu $fp $sp 4          #moving frame pointer by 4 bits
move  $s0 $a0
la    $a0 str_const2     #string constant in the local address
sw    $a0 0($sp)
addiu $sp $sp -4
move  $a0 $s0            #this main method calls pattern method
bne   $a0 $zero label8
la    $a0 str_const4
li    $t1 1
jal   _dispatch_abort

```

## **Program 5**

### **String Reverse:**

Main.main:

```
    addiu  $sp $sp -28          #reserving space on stack
    sw     $fp 28($sp)
    sw     $s0 24($sp)
    sw     $ra 20($sp)
    addiu  $fp $sp 4
    move   $s0 $a0
    la     $a0 str_const0      #string constant in the local address
    sw     $a0 0($sp)
    addiu  $sp $sp -4
    move   $a0 $s0             #labels are present in the code which tell
    bne    $a0 $zero label0    #which operation to perform according to the code
    la     $a0 str_const2
    li     $t1 1
    jal    _dispatch_abort     #end program execution
```

The instructions have different labels to be executed in their program flow the labels and the instructions there will be evaluated accordingly. There will be different branches happening due to conditional statements depending on their boolean values we need to execute certain statements so labels represent these different branches which happens and certain statements in the code will be executed on the branch we go on due to the conditional statement boolean value. Please see that the explanations of some of the instructions have been mentioned in the points above.

## **Incorrect Programs**

### **Program 1**

The identifier for Int starts with a capital letter X, where the identifiers should start with small letters.

**Error:**

"incorrect1.cl", line 3: syntax error at or near TYPEID = X

"incorrect1.cl", line 5: syntax error at or near TYPEID = X

Compilation halted due to lex and parse errors

### **Program 2**

Unterminated string constant

**Error:**

"incorrect2.cl", line 3: syntax error at or near ERROR = Unterminated string constant

Compilation halted due to lex and parse errors

**Program 3**

Error due to dash, improper comment.

**Error:**

"incorrect3.cl", line 8: syntax error at or near '-'

Compilation halted due to lex and parse errors

**Program 4**

Here **'fi'** is a key word in the language

**Error:**

"incorrect4.cl", line 3: syntax error at or near FI

"incorrect4.cl", line 6: syntax error at or near FI

Compilation halted due to lex and parse errors

**Program 5**

White space zero length is not allowed by the COOL language

**Error:**

"incorrect5.cl", line 4: syntax error at or near ERROR = \342

Compilation halted due to lex and parse errors