STOPPING TIMES

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Definition 1. A *binary string of length* n is a tuple $(x_1, x_2, \dots, x_n)$ where each x_k is 0 or 1. Such a string is said to be *stopped* at index k if every index of the tuple in $(k/2, k]$ is zero. The *stopping time* of a binary string is the smallest k such that the string is stopped at index k, or ∞ if no such k exists. Note that, except 1 and ∞, every stopping time is even.

We are primarily interested in binary strings of length n which have stopping time n. Indeed, let $g(n)$ be the number of such strings. Our first goal is to show that $g(n)$ satifies the recurrence

$$\begin{aligned} g(1) = g(2) = g(4) &= 1 \\ g(4n) &= 2g(4n-2) \\ g(4n+2) &= 2g(4n) - g(2n). \end{aligned}$$

This will identify the sequence $a(n) = g(2n)$ as the *Narayana–Zidek–Capell numbers*, A2083 in the OEIS, which begin

$$1, 1, 1, 2, 3, 6, 11, 22, 42, 84, 165, 330, 654, \dots$$

This recurrence was discovered thanks to the OEIS.

Our first definition discusses *finite* binary strings, but it is more convenient to work with *infinite* binary strings which are nonzero in only finitely many places. This is equivalent, and has the benefit that every such infinite string has a finite stopping time when considered as a sufficiently long, finite binary string.

Definition 2. Let

$$V = \bigoplus_{k=1}^{\infty} (\mathbb{Z}/2\mathbb{Z})$$

be the direct sum of infinitely many copies of $\mathbb{Z}/2\mathbb{Z}$. That is, the set of all infinite tuples $(x_1, x_2, x_3, \dots)$ where each x_k is 0 or 1, and only finitely many are 1. For each positive integer k, let $\bigcirc_k$ be the set of elements of V which are zero beyond position k and, when the first k entries are regarded as a finite binary string, they have stopping time k. That is, $\bigcirc_1 = \{0\} \subset V$,

$$\bigcirc_{2k} = \{v \in V : \forall j > k, v_j = 0, \text{ and } v \text{ has stopping time } 2k\}$$

and $\bigcirc_{2k+1} = \emptyset$ for every integer $k \geq 1$.

It is clear that $g(n) = |\bigcirc_n|$ for every positive integer n.

Theorem 1. *The sequence $g(n)$ satisfies the recurrence*

$$\begin{aligned} g(1) = g(2) &= g(4) = 1 \\ g(4n) &= 2g(4n-2) \\ g(4n+2) &= 2g(4n) - g(2n). \end{aligned}$$

Proof. Each nonzero element in V has a final nonzero entry. Let $e_0 \colon V \to V$ be a map which inserts a 0 in the position of this final entry, shifting the previous entry to the right. Similarly, let e_1 insert a 1 into this position. Symbolically,

$$e_0(b) = \begin{cases} 0 \text{ if } b = 0 \\ (b_1, ..., b_j, 0, 1, 0, 0, ...) \text{ if } b = (b_1, ..., b_j, 1, 0, 0, ...) \end{cases}$$

and

$$e_1(b) = \begin{cases} (1, 0, 0, \dots) \text{ if } b = 0 \\ (b_1, ..., b_j, 1, 1, 0, 0, ...) \text{ if } b = (b_1, ..., b_j, 1, 0, 0, ...) \end{cases} .$$

Note that e_0 and e_1 are injective, $e_0(V) \cap e_1(V) = \emptyset$, and $e_0(V) \cup e_1(V) = V$.

For any positive integer k,

$$\bigcirc_{2k+2} \subseteq e_0(\bigcirc_{2k}) \cup e_1(\bigcirc_{2k}).$$

Indeed, the final nonzero entry of every string in $\bigcirc_{2k+2}$ occurs at position $2k+1$. Removing the immediately preceding entry produces a string in $\bigcirc_{2k}$—e_0 and e_1 simply add the entry back.

In the other direction, for $k \geq 2$ we have

$$e_1(\bigcirc_{2k}) \subseteq \bigcirc_{2k+2},$$

because inserting a 1 into position k will never decrease the stopping time. For *odd* $k \geq 2$ we have

$$e_0(\bigcirc_{2k}) \subseteq \bigcirc_{2k+2},$$

because inserting a 0 at position k will only decrease the stopping time if the new stopping time *is* k; this is impossible if k is odd. This shows that

$$S_{2k+2} = e_0(S_{2k}) \cup e_1(S_{2k}),$$

and thus

$$g(2k+2) = 2g(2k)$$

for k odd.

To handle the even case, let

$$\hexagon^1_{2k} = \{(b_1, ..., b_{2k}, 1, 0, 0, ...) \in V \text{ s.t. } (b_1, ..., b_{2k}, 0, 0, ...) \in \hexagon_{2k}\}.$$

(Note that $|\hexagon^1_{2k}| = |\hexagon_{2k}| = g(2k)$.) Then

$$e_0(\hexagon_{4k}) \subseteq \hexagon_{4k+2} \cup \hexagon^1_{2k}.$$

This implies

$$\hexagon_{4k+2} \cup \hexagon^1_{2k} = e_0(\hexagon_{4k}) \cup e_1(\hexagon_{4k}).$$

(That the left contains the right is clear given the preceding inclusion. For the other direction, the only problem is $\hexagon^1_{2k} \subseteq e_0(\hexagon_{4k}) \cup e_1(\hexagon_{4k})$. This is easily handled by considering the element immediately preceding the final nonzero element.) Therefore

$$g(4k+2) + |\hexagon^1_{2k}| = 2g(4k),$$

or

$$g(4k+2) = 2g(4k) - g(2k).$$

This establishes the recurrence. $\square$