

Recitation Problems - Com S 311

Week of Feb 5th - Feb 10th

1. Derive the worst-case runtime of the following loop structure as a function of n and determine its Big-O upper bound. You must show the derivation of the end result. Assume the runtime is expressed as follows:

$$\begin{aligned}\text{Runtime} &= \sum_{i=1}^n \sum_{j=i}^n \sum_{k=i}^j 1 \\ \text{Runtime} &= \sum_{i=1}^n \sum_{j=i}^n \sum_{k=i}^j 1 \\ &= \sum_{i=1}^n \sum_{j=i}^n (j - i + 1) \\ &= \sum_{i=1}^n \frac{(n - i + 1)(n - i + 2)}{2} \quad [\text{See equation (1) below}] \\ &= \frac{1}{2} \sum_{i=1}^n (n^2 - ni + 2n - ni + i^2 - 2i + n - i + 2) \\ &= \frac{1}{2} \sum_{i=1}^n (n^2 - 2ni - 3i + 3n + 2 + i^2) \\ &= \frac{1}{2} \sum_{i=1}^n (n^2 - i(2n + 3) + 3n + 2 + i^2) \\ &= \frac{1}{2} \left[n^3 - \frac{n(n+1)}{2}(2n+3) + 3n^2 + 2n + \frac{n(n+1)(2n+1)}{6} \right] \\ &= \frac{1}{2} \left[n^3 - \frac{2n^3 + 5n^2 + 3n}{2} + 3n^2 + 2n + \frac{2n^3 + 3n^2 + n}{6} \right] \\ &= \frac{1}{2} \left[\frac{6n^3 - 6n^3 - 15n^2 - 9n + 18n^2 + 12n + 2n^3 + 3n^2 + n}{6} \right] \\ &= \frac{1}{12} [2n^3 + 6n^2 + 4n] \\ &\in O(n^3)\end{aligned}$$

$$\begin{aligned}\sum_{j=i}^n (j - i + 1) &= 1 + 2 + 3 + \dots + (n - i + 1) \\ &= \frac{(n - i + 1)(n - i + 2)}{2}\end{aligned} \tag{1}$$

2. You are given an array A of integers, sorted in increasing order, and a target integer T . Write an algorithm to verify whether there exist **two** integers x and y in A such that $x + y = T$. Discuss the derivation of the runtime of your algorithm.

Example: $A = [3, 6, 9, 10, 11]$

- $T = 16 \rightarrow \text{Answer} = \text{Yes}$
- $T = 14 \rightarrow \text{Answer} = \text{Yes}$
- $T = 11 \rightarrow \text{Answer} = \text{No}$
- $T = 22 \rightarrow \text{Answer} = \text{No}$

Algorithm 1 Solution

```
1:  $i = 0$ ;    // leftmost index of array
2:  $j = n - 1$ ; // rightmost index of array
3:  $found = false$ ; // boolean flag
4: while  $i < j$  &&  $!found$  do
5:   if  $A[i] + A[j] < T$  then // implies that we need to look for bigger number on left
6:      $i++$ ;
7:   else if  $A[i] + A[j] > T$  then // implies that we need to look for smaller number on right
8:      $j--$ ;
9:   else
10:      $found = true$ ;
11: return  $found$ ;
```

Runtime Analysis:

The runtime depends on the number of times the while loop iterates in the worst case. The loop iterates as long as $j - i > 0$. Initially, $j - i$ is $n - 1$. In each iteration, this value decrements by 1. Therefore, the number of times the loop iterates is bounded by n . In the body of the loop, we do constant time operations. Hence, the runtime $\in O(n)$.