

Finding n_0 .

$$f(n) = n^2 + n + 3 \text{ where } g(n) = n^2$$

So we have:

Equation from Big Oh Definition:

$$0 \leq f(n) \leq cg(n) \text{ for all } n \geq n_0.$$

Find N_0 and c . They both depend on each other so we can simply choose c in this case.

$$n^2 + n + 3 \leq c(n^2)$$

Let us pick $c = 2$.

$f(n)$:

$$f(1): 1 + 1 + 3 = 5$$

$$f(2): 4 + 2 + 3 = 9$$

$$f(3): 9 + 3 + 3 = 15$$

$cg(n)$:

$$cg(1): 2(1) = 2$$

$$cg(2): 2(4) = 8$$

$$cg(3): 2(9) = 18$$

$$n_0 = 3.$$

So in this case, there is a constant c (2) where $f(n) \leq cg(n)$ for all $n \geq n_0(3)$.

If we were to plot both equations we could see that when $n = 1$ and 2 , function $f(n)$ is higher than function $cg(n)$. However, when $n \geq 3$, we can see that $cg(n)$ is higher than $f(n)$. The definition of Big Oh states that this inequality will hold true for every $n \geq n_0$. Since $n_0 = 3$, we can expect the behavior to continue for any size n when c is 2 .

We could of chosen a different c , for example $c = 4$.

$cg(n)$:

$$cg(1): 4(1) = 4$$

$$cg(2): 4(4) = 16$$

$$cg(3): 4(9) = 36$$

In this case, $n_0 = 2$ when $c = 4$.

STEPS:

1. Pick a c .
2. Make a table for $f(n)$ and $cg(n)$.
3. Find where $cg(n) \geq f(n)$. This is n_0 .