

# Operating Systems

## Chapter 4 Threads, SMP, and Microkernels

# Agenda

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- 4.1 Processes and Threads
- 4.2 Types of Threads

# 4.1 Processes and Threads

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- 4.1.1 Introduction
- 4.1.2 Multithreading
- 4.1.3 Thread Functionality

## 4.1.1 Introduction(1/2)

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- These two characteristics are treated independently by the operating system OS 区别对待以下问题
  - Resource ownership( 资源所有权 ) - process includes a virtual address space to hold the process image
  - Scheduling/execution( 调度 / 执行 )- follows an execution path that may be interleaved with other processes

## 4.1.1 Introduction(2/2)

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- Scheduling/execution
  - Dispatching is referred to as a thread or lightweight process( 调度的单位称为线程或轻量进程 )
- Resource
  - Resource of ownership is referred to as a process or task( 资源所有权的单位称为进程或者任务 )

# 4.1 Processes and Threads

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- 4.1.1 Introduction
- 4.1.2 Multithreading
- 4.1.3 Thread Functionality

## 4.1.2 Multithreading(1/9)

- Single-Thread( 单线程 )
  - *Single-threaded approach* refers to the traditional approach of a single thread of execution per process, in which the concept of a thread is not recognized ( 单线程指一个进程中只有一个线程在执行的传统方法，线程的概念并不明确 )
    - MS-DOS supports a single thread( 单进程、单线程 )
    - Some UNIX supports multiple user processes but only supports one thread per process ( 多进程，每个进程单线程 )

## 4.1.2 Multithreading(2/9)

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- Multithreading( 多线程 )
  - *Multithreading* refers to the ability of an OS to support multiple threads of execution within a single process( 指操作系统支持在一个进程中执行多个线程的能力 )
    - Windows, OpenEuler, Solaris, Linux, Mach, and OS/2 support multiple threads



## 4.1.2 Multithreading(3/9)

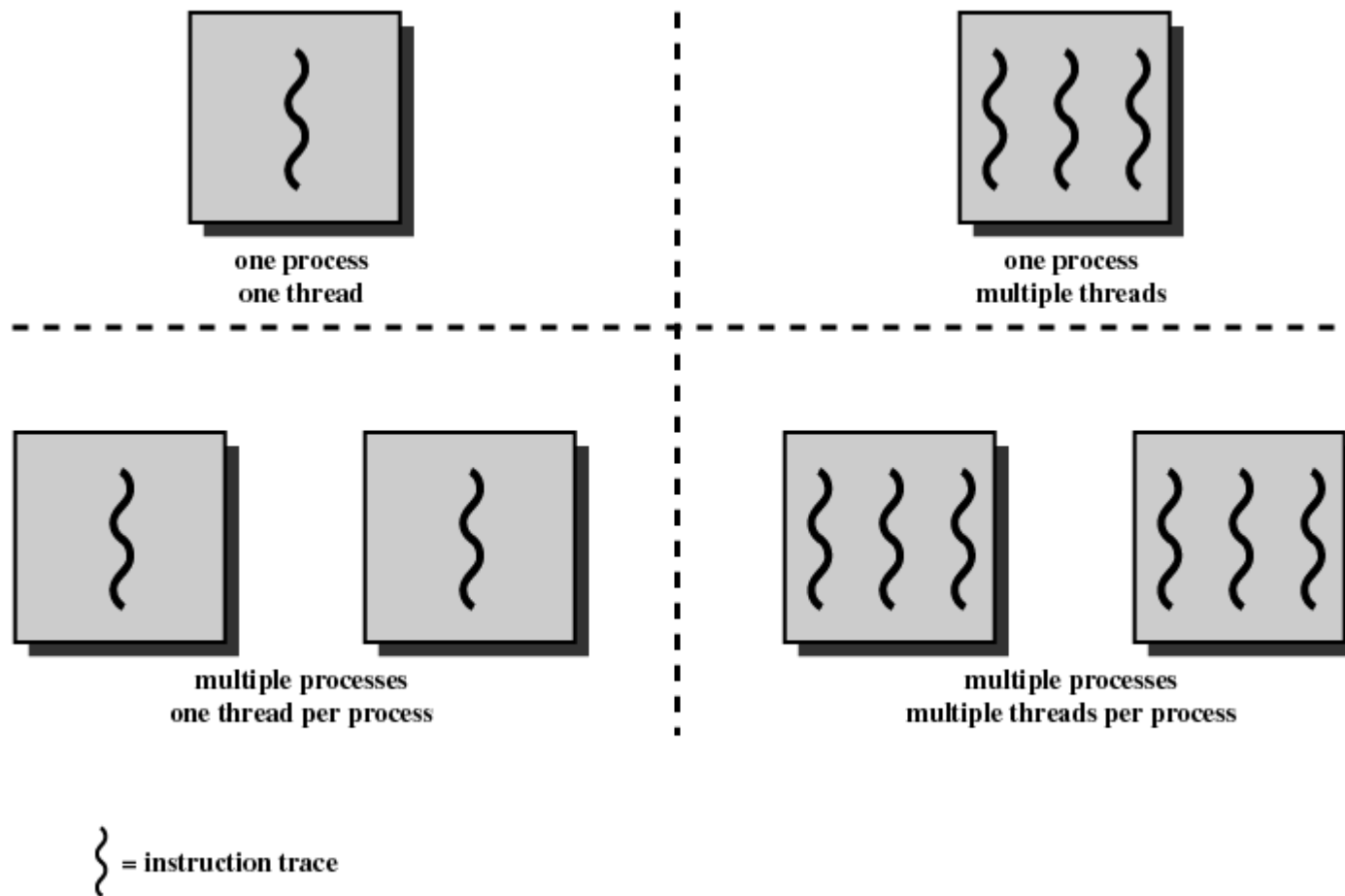


Figure 4.1 Threads and Processes [ANDE97]

## 4.1.2 Multithreading(4/9)

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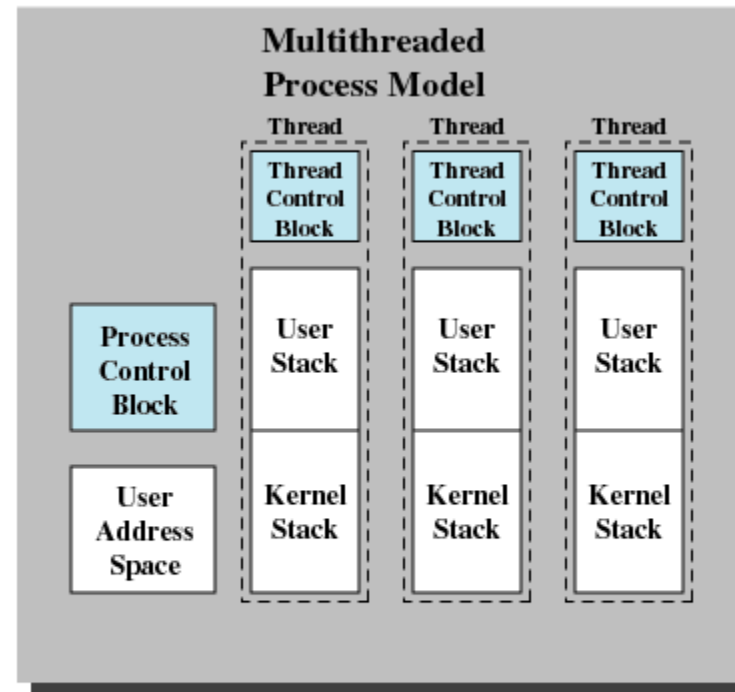
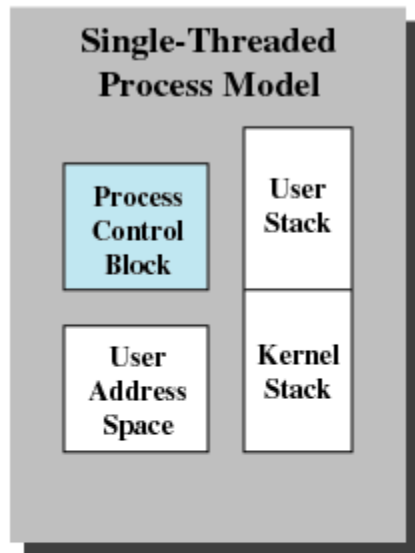
- **Process** : Unit of Resource allocation and Unit of Protection( 资源分配和保护的单位 )
  - Have a virtual address space( 虚拟地址空间 ) which holds the process image( code, data, stack and PCB)
  - Protected access to processors, memory, other processes( 与其他进程通信 ), files, and I/O resources
  - Contains one or more threads

## 4.1.2 Multithreading(5/9)

- **Thread** : Unit of Scheduling/Execution( 调度执行的单位 )
  - Each thread has:
    - An execution state (running, ready, etc.)
    - Saved thread context when not running (pc)
    - Has an execution stack
    - Some per-thread static storage for local variables 局部变量
    - Access to the memory and resources of its process 比如全局变量

## 4.1.2 Multithreading(6/9)

Distinction Between Threads and Processes From the Point of View of Process Management ( 从进程管理的角度看线程和进程 )

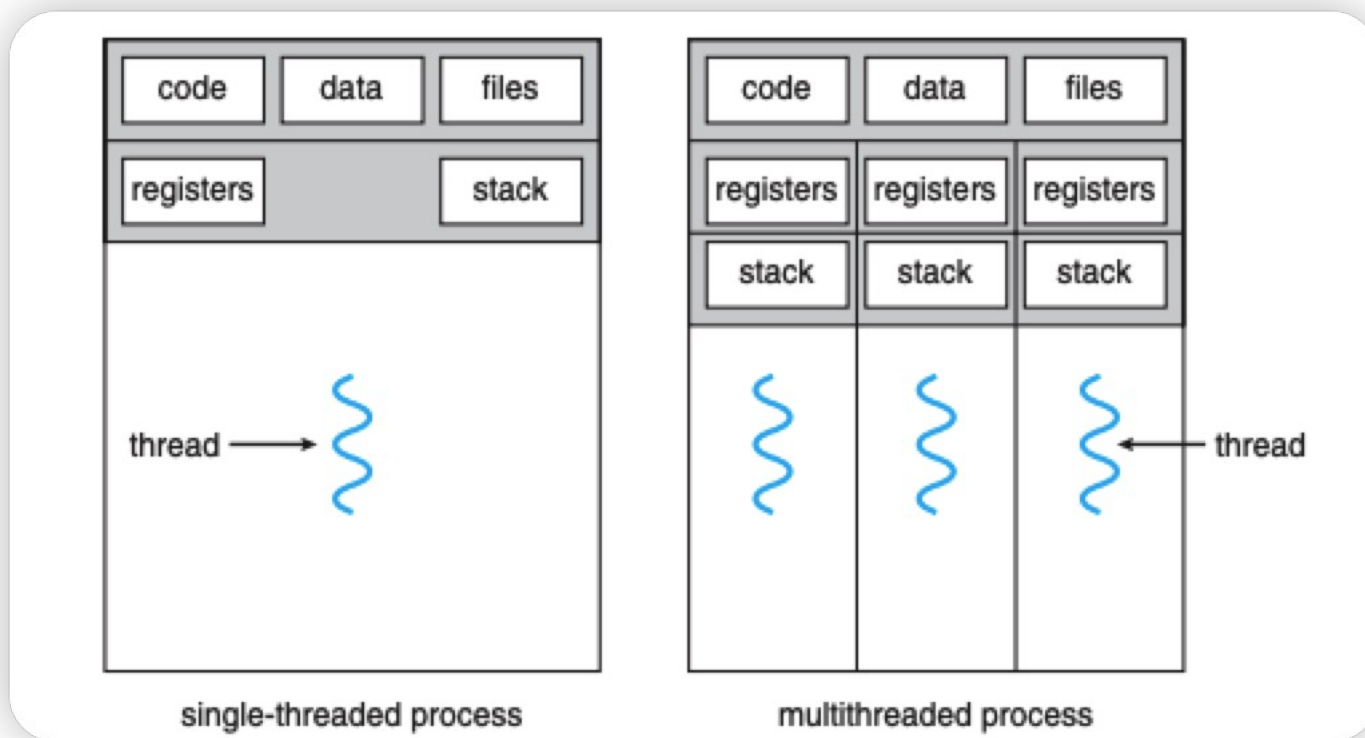


线程 = 进程 - 共享资源

Figure 4.2 Single Threaded and Multithreaded Process Models

## 4.1.2 Multithreading(7/9)

多线程共享进程的代码段和数据段（全局变量），每个线程有自己的程序计数值和局部变量 stack



## 4.1.2 Multithreading(8/9)

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### Benefits( 优点 ) of Threads

1. Takes less time to create a new thread than a process ( 创建快 ) 不重头分配内存等资源
2. Less time to terminate a thread than a process ( 结束快 )
3. Less time to switch between two threads within the same process ( 切换快 ) 相同镜像，不需要切换页表
4. Since threads within the same process share memory and files, they can communicate with each other without invoking the kernel ( 通信快 ) 不经过内核通信

## 4.1.2 Multithreading(9/9)

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- Threads are Affected by Many Process Action
  - **Suspending** a process involves suspending all threads of the process since all threads share the same address space ( 挂起进程会挂起该进程的所有线程 )
  - **Termination** of a process, terminates all threads within the process ( 终止进程会终止该进程的所有线程 )

# 4.1 Processes and Threads

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- 4.1.1 Introduction
- 4.1.2 Multithreading
- 4.1.3 Thread Functionality ( 线程功能特性 )



## 4.1.3 Thread Functionality(1/5)

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- key states for a thread
  - Running, Ready,Blocked.
- Operations associated with a change in thread state
  - Spawn( 派生 )
    - Spawn another thread
  - Block( 阻塞 )
  - Unblock( 解除阻塞 )
  - Finish
    - Deallocate register context and stacks

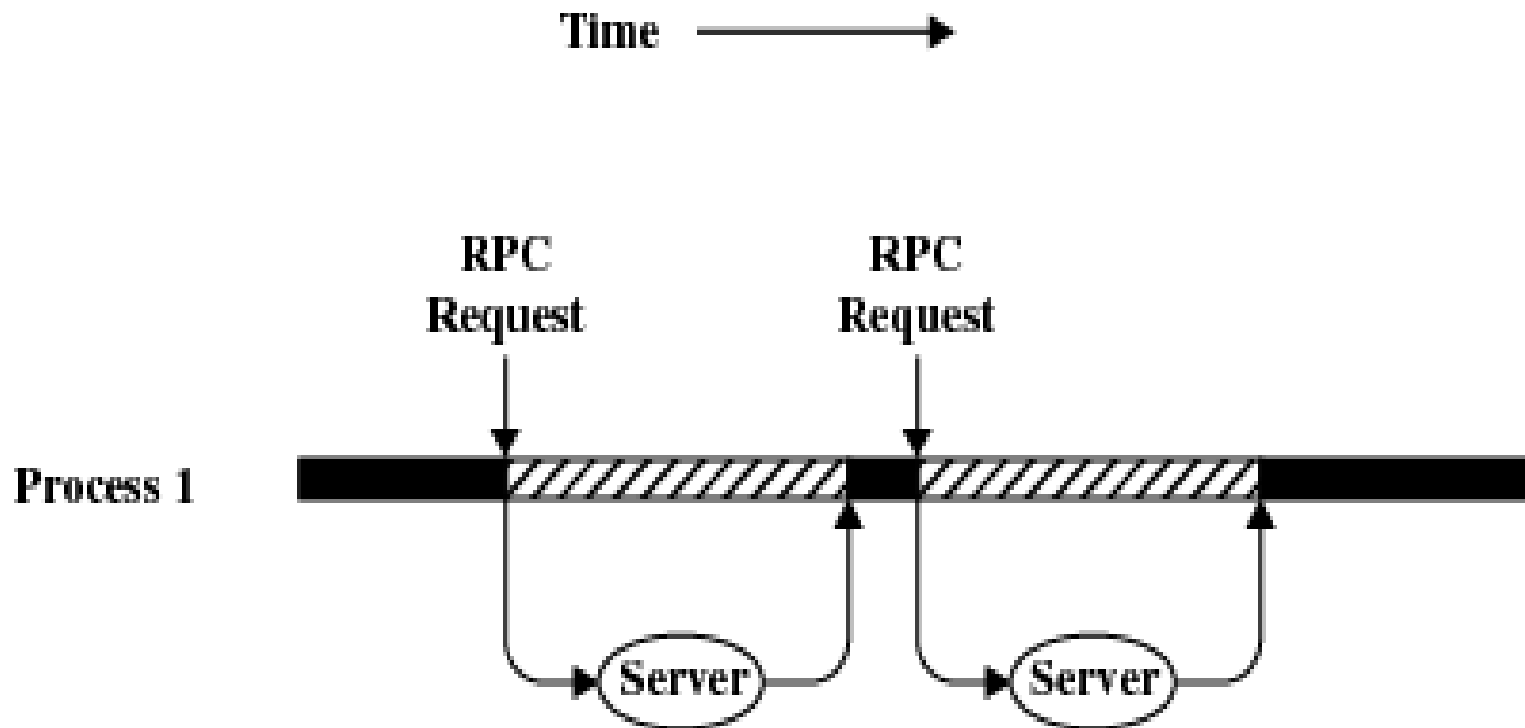
## 4.1.3 Thread Functionality(2/5)

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- Suspending a process involves suspending all threads of the process
  - swap address space to disk
- Blocking a thread involves blocking the process , right ?
  - Better not

## 4.1.3 Thread Functionality(3/5)

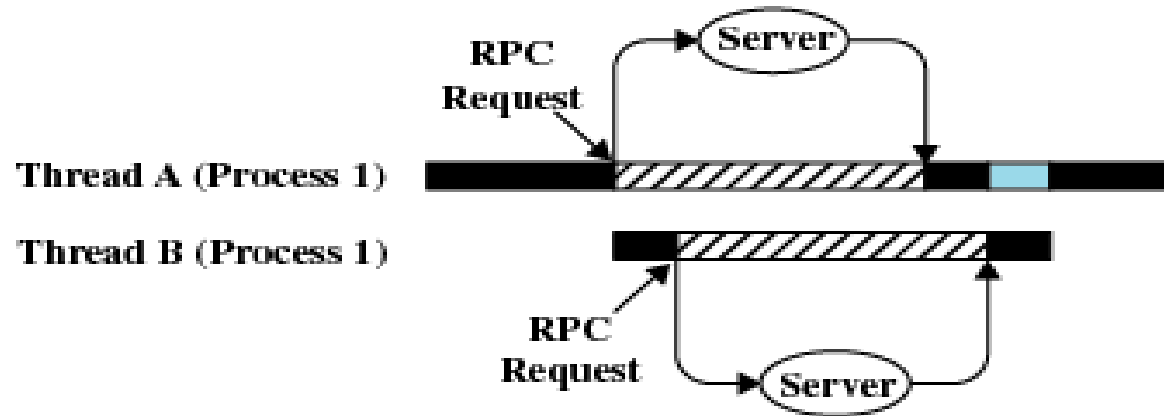
### Remote Procedure Call Using Single Thread






(a) RPC Using Single Thread

## 4.1.3 Thread Functionality(4/5)

### Remote Procedure Call Using Two Threads



(b) RPC Using One Thread per Server (on a uniprocessor)

-  Blocked, waiting for response to RPC
-  Blocked, waiting for processor, which is in use by Thread B
-  Running

**Figure 4.3 Remote Procedure Call (RPC) Using Threads**

## 4.1.3 Thread Functionality(5/5)

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- THREAD Synchronization 线程同步
- Why?
- In a process:
  - All threads share the same address space and other resources(open files).
  - Any alteration of a resource by one thread affects the others.

# Agenda

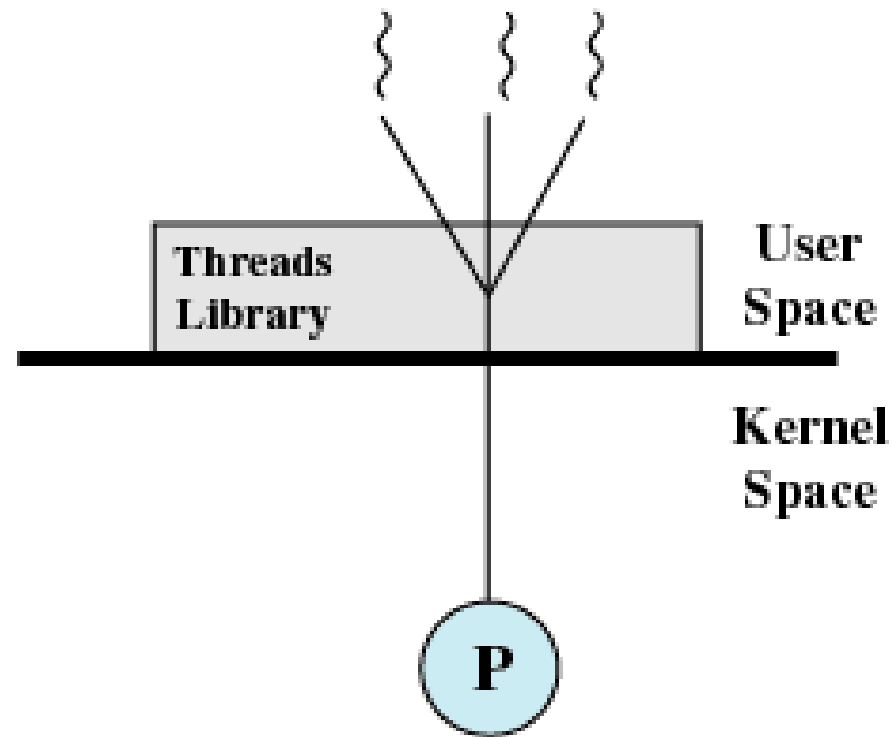
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- 4.1 Processes and Threads
- 4.2 Types of Threads
  - 4.2.1 User-Level Introduction
  - 4.2.2 Kernel-Level Thread
  - 4.2.3 Combined Approaches
  - 4.2.4 Other Arrangements

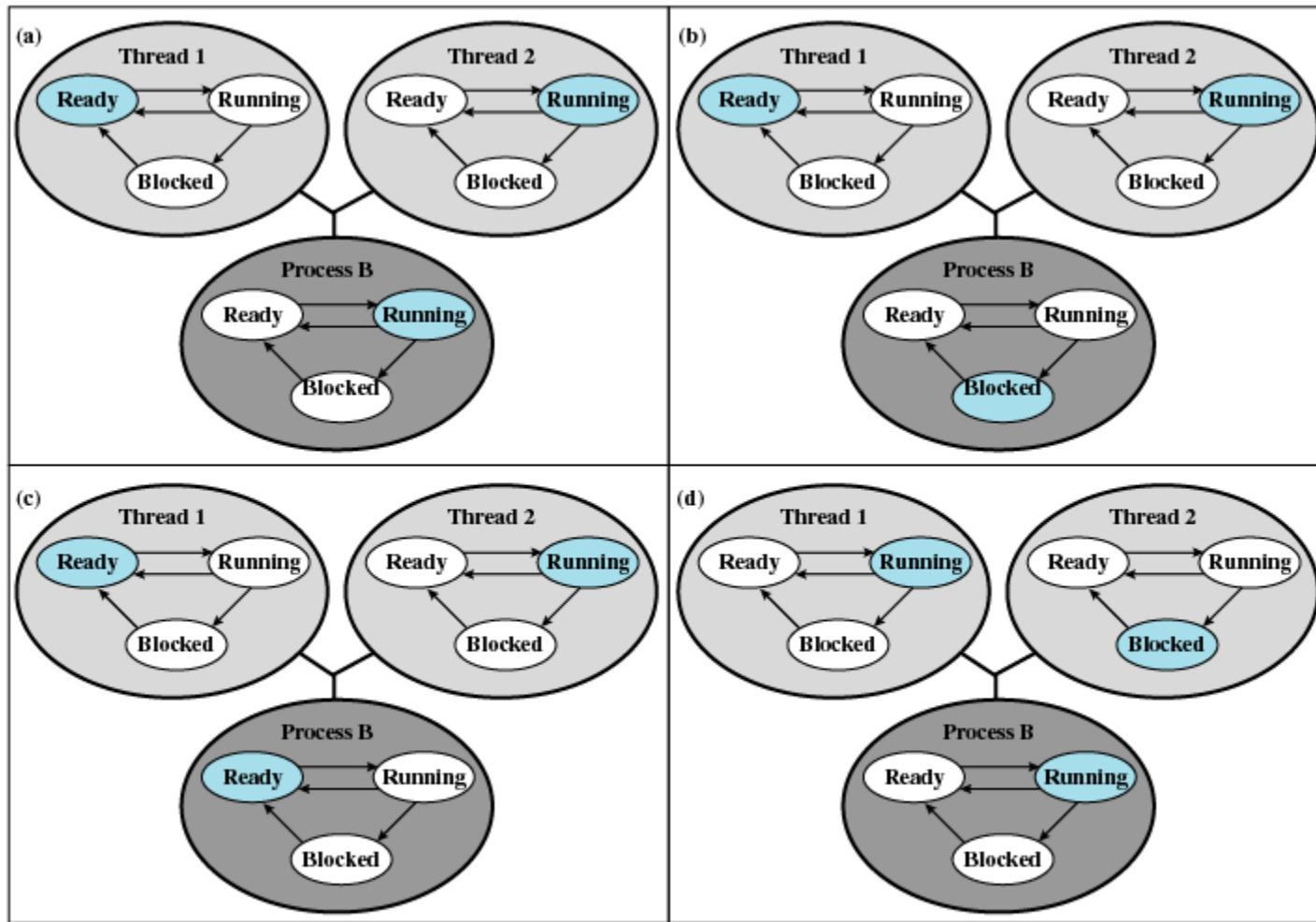
## 4.2 Types of Threads(1/10)

### 1. User-Level Threads (ULT , 用户级线程 )

- Multithread implemented by a threads library 线程库
- All thread management is done by the application
- The kernel is not aware of the existence of threads & scheduling is done on a process basis



## 4.2 Types of Threads(2/10)



Colored state  
is current state

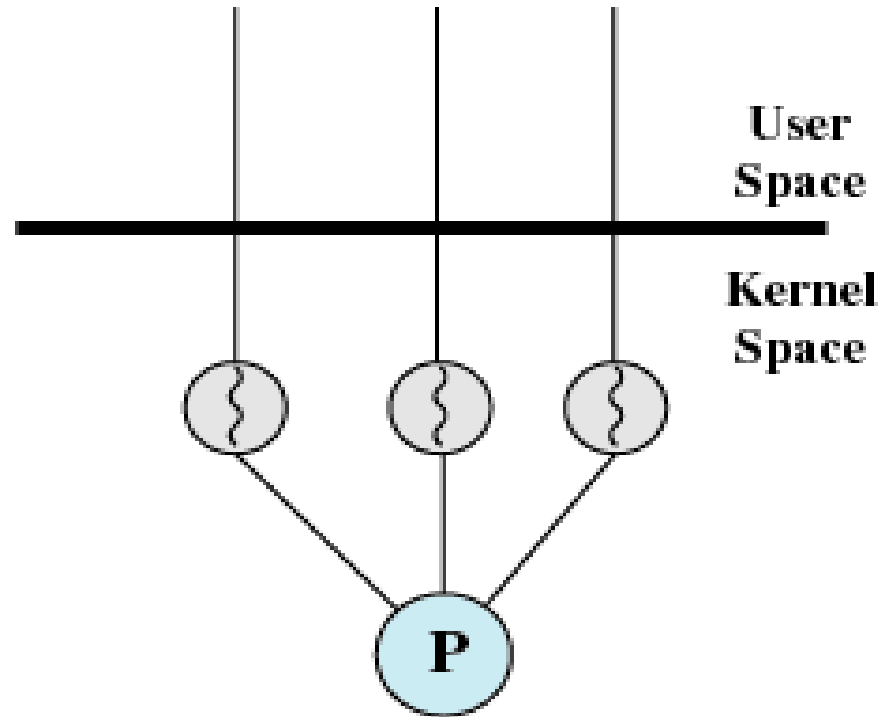
Figure 4.7 Examples of the Relationships Between User-Level Thread States and Process States



## 4.2 Types of Threads(3/10)

### 2. Kernel-Level Threads( 内核级线程 )

- Kernel maintains context information for the process and the threads
- Scheduling is done on a thread basis



(b) Pure kernel-level

## 4.2 Types of Threads(4/10)

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- Advantages of ULT to KLT
  - Less overhead of thread switches(mode switches do not required) 切换开销小
  - Scheduling can be application specific 调度策略根据应用可以不同
  - ULTs can run on any operating system without modify the underlying kernel 无需底层内核修改

## 4.2 Types of Threads(5/10)

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- Disadvantages of ULT to KLT
  - One thread is blocked, all other threads of the process are blocked(ULT 按进程调度 )
  - A multithreaded application cannot take advantage of multiprocessing ( 线程不能分配到多核 )
  - Ways to work around these drawbacks:
    - Multiple processes 用多进程代替多线程
    - Jacketing 套管 针对阻塞问题

## 4.2 Types of Threads(6/10)

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- KLT
  - Kernel manage the threads
    - Windows provides API
  - Kernel routines themselves can be multithreaded
- Advantages of KLT to ULT
  - Overcomes the two principal drawbacks of the ULT
    - Multiple threads in one process can simultaneously run on multiple processors 多 CPU 执行
    - One threads blocked cannot make the other threads within the same process blocked 仅阻塞单个线程

## 4.2 Types of Threads(7/10)

- Disadvantages of KLT to ULT
  - The principal disadvantage is that thread switch requires mode switches( 模式切换 ) to the kernel

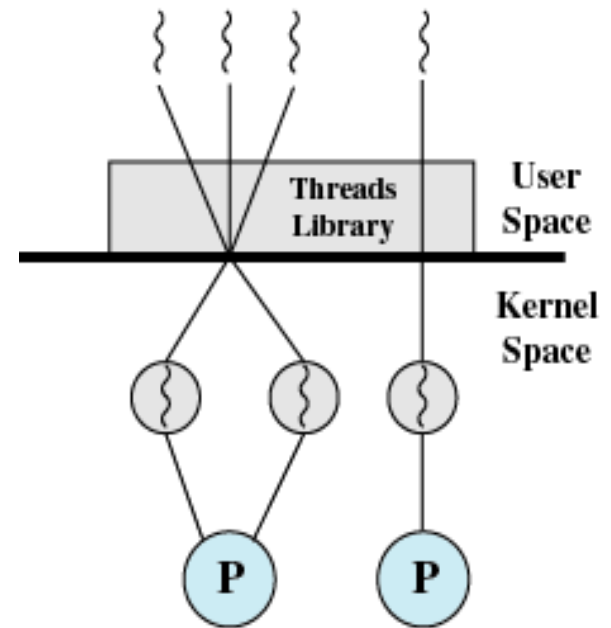
**Table 4.1** Thread and Process Operation Latencies ( $\mu s$ )

Operation	User-Level Threads	Kernel-Level Threads	Processes
Null Fork	34	948	11,300
Signal Wait	37	441	1,840

## 4.2 Types of Threads(8/10)

### 3. Combined Approaches 混合方式

- thread creation/scheduling/synchronization is done completely in user space 分配调度同步可以在用户空间
- The multiple ULTs from a single application are mapped onto some (smaller or equal) number (adjustable) of KLTs. 允许映射到内核



(c) Combined

## 4.2 Types of Threads(9/10)

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- Advantages of Combined Approaches
  - multiple threads within the same application run in parallel on multiple processors 多 cpu 运行
  - A blocking system call only blocks the thread. 仅阻塞单个线程

## 4.2 Types of Threads(10/10)

### Relationship Between Threads and Processes

**Table 4.2 Relationship Between Threads and Processes**

Threads:Processes	Description	Example Systems
1:1	Each thread of execution is a unique process with its own address space and resources.	Traditional UNIX implementations
M:1	A process defines an address space and dynamic resource ownership. Multiple threads may be created and executed within that process.	Windows NT, Solaris, Linux OS/2, OS/390, MACH
1:M	A thread may migrate from one process environment to another. This allows a thread to be easily moved among distinct systems.	Ra (Clouds), Emerald
M:N	Combines attributes of M:1 and 1:M cases.	TRIX