Intercalação

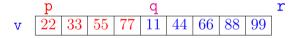


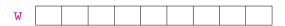
Fonte: http://csumplugged.org/sorting-algorithms PF 9

http://www.ime.usp.br/~pf/algoritmos/aulas/mrgsrt.html

(□) (∰) (E) (E) (O)

Intercalação





k v [11]

Intercalação



Interca**l**ação

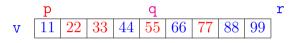
Problema: Dados $\mathbf{v}[\mathbf{p}:\mathbf{q}]$ e $\mathbf{v}[\mathbf{q}:\mathbf{r}]$ crescentes, rearranjar $\mathbf{v}[\mathbf{p}:\mathbf{r}]$ de modo que ele fique em ordem crescente.

Para que valores de ${\color{red} {\bf q}}$ o problema faz sentido?

Entra:

p q 1 7 22 33 55 77 11 44 66 88 99

Sai:



Intercalação



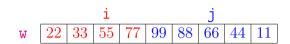
Intercalação

Intercalação

Interca**l**ação

				k			
٧	11	22	33				

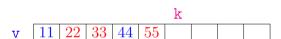


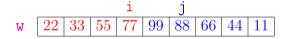


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Intercalação

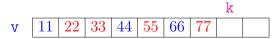
Intercalação





Intercalação

Intercalação



Intercalação

v 11 22 33 44 55 66 77 88 99

Consumo de tempo

n := r - p

linha	proporcional a
1–2	?
3	?
4	?
5 - 6	?
7	?
8–13	?
total	7

Conclusão

A função intercale consome 5n + 2 unidades de tempo.

O algoritmo intercale consome O(n) unidades de tempo.

Também escreve-se

O algoritmo intercale consome tempo O(n).

Intercalação

```
def intercale(p, q, r, v):
     e = v[p:q] # clone
 1
 2
     d = v[q:r] # clone
     d.reverse() # método mutador
     w = e + d
     i = 0
 5
     j = r-p-1
 6
 7
     for k in range(p,r):
        if w[i] <= w[j]):
 8
           v[k] = w[i]
 9
            i += 1
10
11
        else:
           v[k] = w[j]
12
            j -= 1
13
```

Consumo de tempo

$$n := r - p$$

linha	pro	porcional a		
1–2	=	n		
3				
4	=	n		
5–6	=	1		
7	=	r-p+1	=	n + 1
8–13	=	r-p	=	n
total	\approx	5n + 2		

Ordenação: algoritmo Mergesort



Fonte: https://www.youtube.com/watch?v=XaqR3G_NVoo

PF 9

http://www.ime.usp.br/~pf/algoritmos/aulas/mrgsrt.html

Ordenação

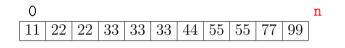
 $v[0:\underline{n}] \text{ \'e crescente se } v[0] \leq \dots \leq v[\underline{n}{-}1].$

Problema: Rearranjar um vetor v[0:n] de modo que ele fique crescente.

Entra:

```
0
33 | 55 | 33 | 44 | 33 | 22 | 11 | 99 | 22 | 55 | 77
```

Sai:



merge_sort

Rearranja v[p : r] em ordem crescente.

<□> <□> < □> < ₹> < ₹> ₹ 9<€

merge_sort

Rearranja v[p : r] em ordem crescente.

merge_sort

Rearranja v[p:r] em ordem crescente.

merge_sort

4 m > 4 m >

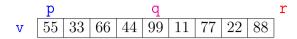
Rearranja v[p : r] em ordem crescente.

merge_sort

Rearranja v[p : r] em ordem crescente.

p q v | 55 | 33 | 66 | 44 | 99 | 11 | 77 | 22 | 88 |

Mergesort

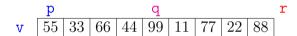


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Mergesort



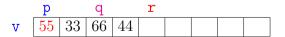
Mergesort

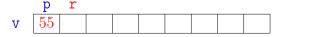




Mergesort

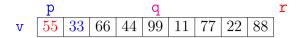


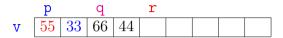


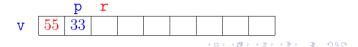




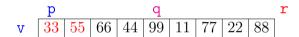






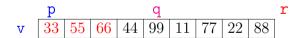


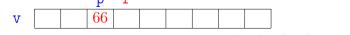
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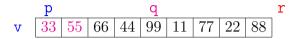


Mergesort





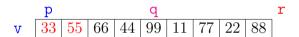
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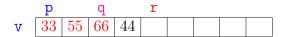
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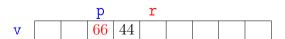




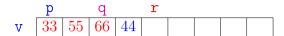


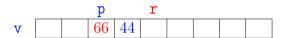






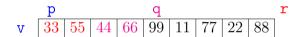


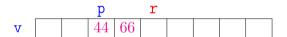




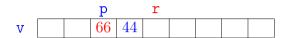


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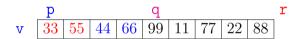




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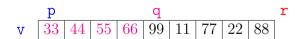


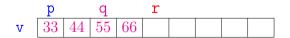
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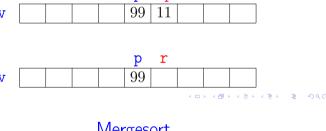
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Mergesort

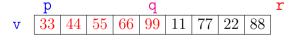


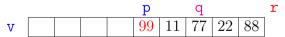


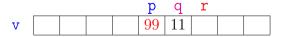
Mergesort v 33 44 55 66 99 11 77 22 88 p q r v 99 11 77 22 88 Mergesort p q r v 33 44 55 66 99 11 77 22 88 r v 99 11 77 22 88

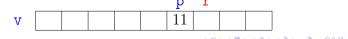


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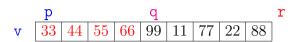


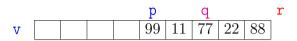






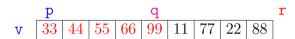
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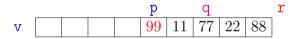


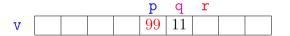


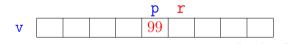
			p	q	r	
v			99	11		

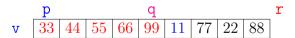
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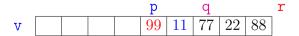


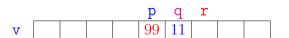






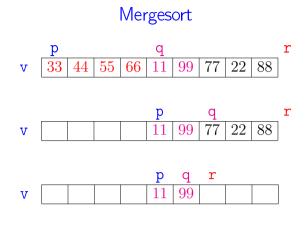


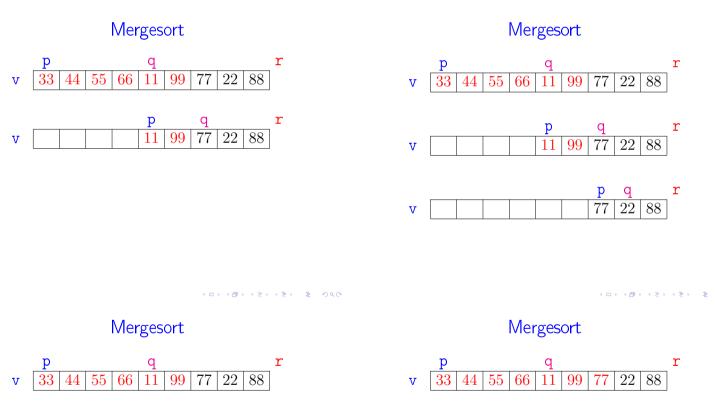


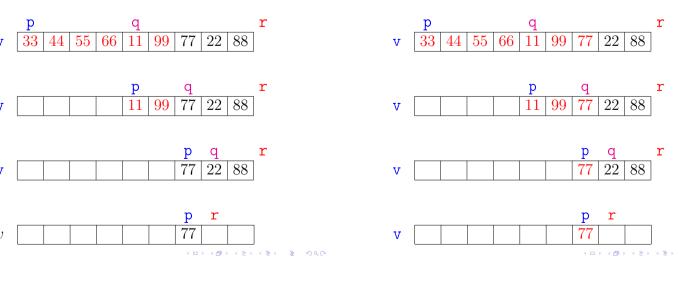




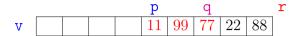
Mergesort v 33 44 55 66 99 11 77 22 88 p q r v 99 11 77 22 88 p q r v 99 11 77 22 88 Mergesort v 99 7 22 88





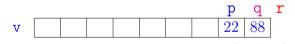


p q r v 33 44 55 66 11 99 77 22 88



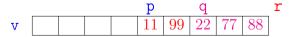


Mergesort



Mergesort

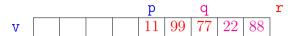






Mergesort



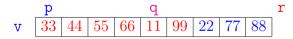






Mergesort







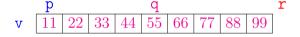
v 33 44 55 66 11 22 77 88 99 r

			p		q			r
V			11	22	77	88	99	

Mergesort



Mergesort



Mergesort

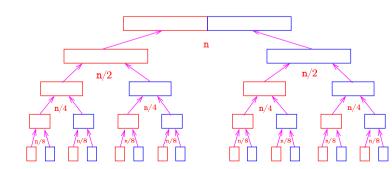
Correção

```
def merge_sort (p, r, v):
1    if p < r-1:
2        q = (p + r) // 2
3        merge_sort(p, q, v)
4        merge_sort(q, r, v)
5    intercale(p, q, r, v)</pre>
```

A função está correta?

A correção da função, que se apóia na correção do intercale, pode ser demonstrada por indução em $\mathbf{n} := \mathbf{r} - \mathbf{p}$.

Consumo de tempo: versão MAC0122



Consumo de tempo: versão MAC0122

O consumo de tempo em cada nível da recursão é proporcional a n.

Há cerca de lg n níveis de recursão.

nível	consumo de tempo (proporcional a)	
1	$pprox \mathtt{n}$	= n
2	$\approx \frac{\mathrm{n}}{2} + \frac{\mathrm{n}}{2}$	= n
3	$\approx \frac{n}{4} + \frac{n}{4} + \frac{n}{4} + \frac{n}{4}$	= n
	•••	
lg n	$\approx 1+1+1+1\cdots+1+1$	= n
Total	$\sim n \log n = O(n \log n)$	

 $| \text{lotal} \approx n \lg n = O(n \lg n)$

4D> 4B> 4B> B 9Q0

Consumo de tempo: outra versão

```
def merge_sort (p, r, v):
     if p < r-1:
         q = (p + r) // 2
3
        merge_sort(p, q, v)
        merge_sort(q, r, v)
4
5
         intercale(p, q, r, v)
               consumo na linha (proporcional a)
        1
        2
               ?
               ?
        3
               ?
         T(n) = ?
```

Consumo de tempo: outra versão

T(n) := consumo de tempo quando <math>n = r - p

$$\begin{split} &T(1)=1\\ &T(\textbf{n})=T(\lceil \textbf{n}/2\rceil)+T(\lceil \textbf{n}/2\rceil)+\textbf{n} \ \ \text{para} \ \textbf{n}=2,3,4,\ldots \end{split}$$

Solução: $T(n) \in O(n \log n)$.

Demonstração: ...

Consumo de tempo: outra versão

```
def merge_sort (p, r, v) {
    if p < r-1:
1
2
       q = (p + r) // 2
       merge_sort(p, q, v)
3
4
       merge_sort(q, r, v)
5
       intercale(p, q, r, v)
```

Consumo de tempo?

def merge_sort (p, r, v):

T(n) := consumo de tempo quando <math>n = r - p

10 + 40 + 45 + 45 + 5 99 C

Consumo de tempo: outra versão

```
if p < r-1:
          q = (p + r) // 2
3
          merge_sort(p, q, v)
4
          merge_sort(q, r, v)
5
          intercale(p, q, r, v)
                consumo na linha (proporcional a)
                  = 1
          2
                  = 1
          3
                  = T(|\mathbf{n}/2|)
          4
                  = T(\lceil \frac{n}{2} \rceil)
```

Conclusão

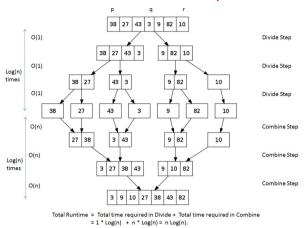
 $T(n) = T(\lceil n/2 \rceil) + T(\lfloor n/2 \rfloor) + n + 2$

O consumo de tempo da função merge_sort é proporcional a $n \lg n$.

O consumo de tempo da função merge_sort é $O(n \lg n)$.

4 D > 4 B > 4 B > 4 B > 9 Q @

Consumo de tempo



Fonte: http://images.1233.tw/in-place-quicksort-algorithm/

merge_sort: versão iterativa

```
def merge_sort (n, v):
    b = 1
    while b < n:
        p = 0
        while p + b < n:
        r = p + 2*b
        if r > n:        r = n
            intercale(p, p+b, r, v)
        p = p + 2*b
        b = 2*b
```

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Divisão e conquista

Algoritmos por divisão-e-conquista têm três passos em cada nível da recursão:

Dividir: o problema é dividido em subproblemas de tamanho menor;

Conquistar: os subproblemas são resolvidos recursivamente e subproblemas "pequenos" são resolvidos diretamente;

Combinar: as soluções dos subproblemas são combinadas para obter uma solução do problema original.

Exemplo: ordenação por intercalação (merge_sort).

←□ → ←□ → ←□ → ←□ → □