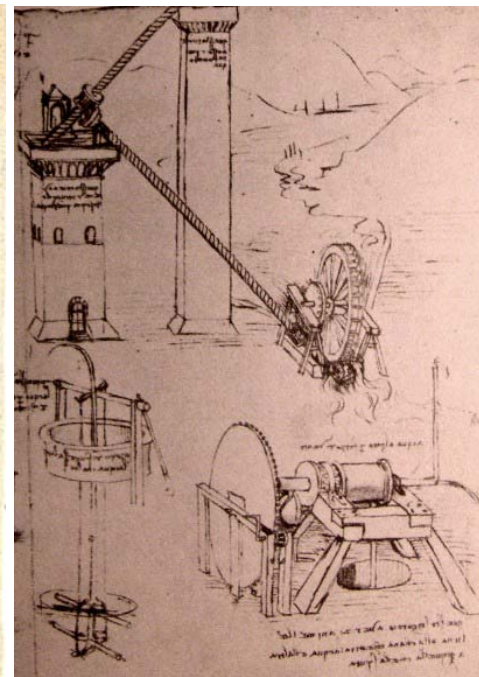
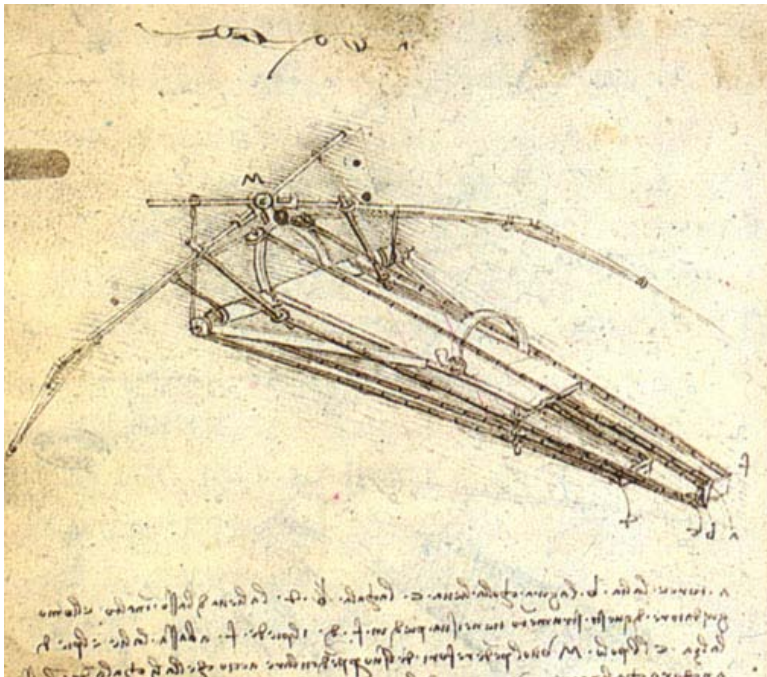


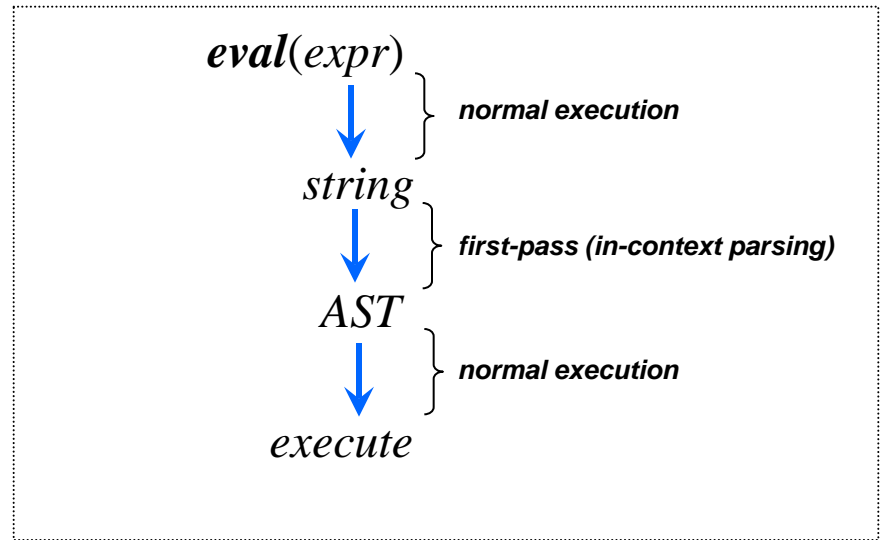
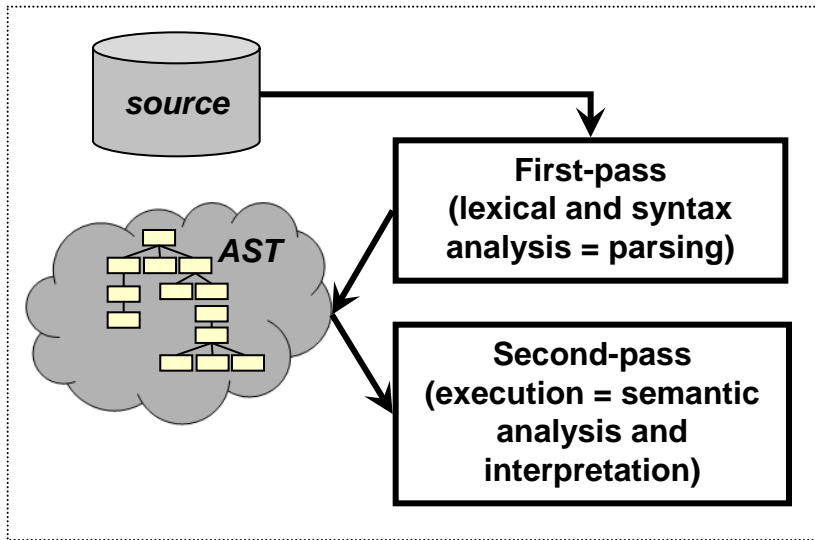
HY540 – Advanced Topics in Programming Language Development



Project

Project (1/11)

- *Small interpreted meta language named **sin***
 - ❑ compile to AST and run
 - ❑ no *eval()* is required (*optional*)



Project (2/11)

- *Untyped object-based* with dynamically typed variables (like *JavaScript*)
- Objects *ex nihilo* as field dictionaries
 - Object construction expressions are optional
 - $\{ [\langle \textit{key} \rangle \langle \textit{value} \rangle (, \langle \textit{key} \rangle \langle \textit{value} \rangle)^*] \}$
 - $a = \{ x \ 1, \ y \ 2, \ z \ f() \};$
 - $a = \{ \}; \ a.x = 1; \ a.y = 2; \ a.z = f();$
 - If object construction expressions are not supported, you need basically the expression $\{ \}$ to imply “*make a new object*”

Project (3/11)

- Basic reference counting for objects
 - very primitive garbage collection not working for cyclic references
- All key types are allowed but are treated as strings just before hashing
 - `t[0] ≡ t["0"]`
 - `t[t] ≡ t["object(0x12)"]`
 - assuming unique serial number for objects, else you may use the native object address or reference
 - Reserve a *keys* and *size* key that can't collide with user-defined keys
 - *keys* = object with all keys, *size* = number of fields
 - e.g. `a.#keys` and `a.#size`

Project (4/11)

- Simple syntax but always *of your choice*
 - *You may want to justify that*
 - Don't be biased or affected by the syntax of my previous or later examples
- Statement based
 - Notion of a global program as a series of statements with the following simplifications
 - Nested functions but no access to outer locals (only globals)
 - No closures
 - No lambda (anonymous functions) - *optional*
 - No true methods (only syntactic sugar)
 - $a \rightarrow f(b) \equiv a.f(a, b)$ assuming \rightarrow means method invocation
 - No functors - *optional*
 - No function expressions - *optional*

Project (5/11)

- Function environment *fenv* is a hash table *Environment* (internal object) carrying
 - Arguments
 - Locals
- Global variables are stored in a global *genv* that is used for variables outside functions
- Inside functions *genv* and *fenv* are normally accessed
 - with every function call getting its reserved *fenv* copy
 - created at *call site (caller)* with the actual arguments and then expanded at the *callee site (called)* with the locals

Project (6/11)

- Extensible library functions programmed in native code
 - the language in which you make the *sin*
 - *Environment* as parameter to library functions
 - to extract actual arguments
 - You need also a way to return a value from library functions
 - you can reserve a *retval* field in the *Environment* object
 - the environment is also to be destroyed at the *call site* after the invocation returns and the return value is (optionally) taken

Project (7/11)

- You will need to support native objects (pointers or references) as types in *sin*
 - those are returned by library functions and they are created and managed exclusively by native code
- You will need this feature since
 - ASTs will be native objects
 - handled by the set of tree manipulation library functions you will create and install

Project (8/11)

- Basic metaprogramming features of *sin*
 - *.< expr >.* *shift to meta level ($expr \rightarrow AST$)*
 - *.~ var* *assume var already carries an AST*
 - *.! expr* *compile (execute) an AST (meta expression)*
- Optional features
 - *.@ expr* *compile a string (runtime) expr to an AST*
 - *eval(expr)* is *.!.@expr*
 - *.# expr* *unparse a meta expression ($AST \rightarrow text$)*
- *There is no separation of metafunctions from normal functions*
 - *In principle, any program expression may be meta code*
 - *You can use .! to compile a manually produced AST via lib functions*

Project (9/11)

- Metaprograms you may develop for testing (1/2)
 - simple local optimizations
 - optimized versions of functions like power
 - loop unrolling for statically known number of iterations
 - adding diagnostics into functions
 - wrapping invocations to specific functions with *before* and *after* messages
 - aspectual transformations
 - simple *advice* for simple code patterns like adding *Design by Contract* calls in methods
 - in every function *f* with a *self* first argument add the following code:

```
precond = self.pref_f; if (precond) precondition(self);
```

Project (10/11)

■ Metaprograms you may develop for testing (2/2)

□ generating object factories

- Meta functions accepting pairs of slot identifiers and initial values, producing a respective object factory function

```
st_factory = .!factory("name", "", "address", "", "semester", 1);  
student = st_factory.new();
```

□ static function analysis

- exit paths *compile warnings*
- simple dead code elimination *if (false), return; <code>*
- assignment in condition *compile warnings*

□ static function style checker

- function size in statements
- expression complexity

Project (11/11)

■ *Rules and terms*

- ❑ Up to four persons
- ❑ Any implementation language allowed
- ❑ No report needed
- ❑ Examination is a formal presentation
- ❑ Examples from three categories at least required
- ❑ Presentation date (*final*) is Wednesday 09/09/09
- ❑ Plenary discussions with all teams biweekly possible