definition

Thesis

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1 Does not fit yet

Algebraic space = Classical algebraic space. Let \mathbb{T} be a saturated topology. Let U be an affine in \mathbb{T} , $R:U^2\to \operatorname{Prop}$ be a smooth equivalence relation, meaning that the fibers R_x are smooth algebraic spaces for x:U. I have shown previously that the identity types then are algebraic spaces. Let U/R denote the sheafifaction of the set truncation of the homotopy quotient. We want to show, that As U is projective we can choose $\tilde{R}:U^2\to \mathbb{T}$ such that $\|\tilde{R}xy\|_{\mathbb{T}}=Rxy$. Consider the sheafification of the homotopy quotient $U//\tilde{R}$, this will be a 1-stack whose identity types are in \mathbb{T} . Hence it suffices to show that the map $f:U//\tilde{R}\to U/R$ is fibered in smooth 0-stacks. Consider a term in U/R. By descent we may assume its of the form [x] for some x:U. I claim that the map

$$\tilde{R}_x \to \mathrm{fib}_f[x] = \sum_{t:U//\tilde{R}} ft =_{U/R} [x]$$

is an equivalence, this is en

2 A minimal approach

Fix \mathbb{T} a merely inhabited class of types.

Definition 2.1. Let $C \supset \mathbb{T}$ be a superclass stable under Σ . The class of smooth C-stacks is the smallest intermediate class $\mathbb{T} \subset \tilde{\mathbb{T}} \subset C$ that is closed under covers between types in C: If X,Y:C and $X\to Y$ is fibered in \tilde{T} , then $X\in \tilde{\mathbb{T}}$ iff $Y\in \tilde{\mathbb{T}}$.

The closedness under covers assumption is the conjunction of closed under \sum (as C \sum -stable) and closed under quotients.

Lemma 2.2. smooth C-stacks contain 1 and are closed under \sum .

Example 2.3. smooth Aff-stacks = saturation of \mathbb{T} . Indeed: By definition, the saturation of \mathbb{T} is is obtained by quotients of \mathbb{T} by \mathbb{T} -covers. We have shown, that its closed under covers between affines.

Definition 2.4. We call X a C'-stack, iff there merely exists some affine Spec $A \to X$ fibered in smooth C-stacks.

We call X a C-stack, iff its a C'-stack and $X \in C$.

Definition 2.5. The (smooth ∞ -stacks are the (smooth) \mathcal{U} -stacks.

Lemma 2.6. X is a n' stack iff its an n + 1-stack

Proof. If its

Lemma 2.7. C-stacks are closed under id-types.

Proof. This is similar to 14.2.

Warning. The previous lemma does not hold for smooth stacks: Identity types of things in \mathbb{T} could be empty.

THIS IS UNUSUAL, but surprisingly useful. Let $n \geq 0$.

Example 2.8. Affine smooth 0-stacks are the saturation of \mathbb{T} .

Definition 2.9. X is a (smooth) 0-stack, if its a (smooth) 0-type-stack.

Theorem 2.10 (TODO). Let X be a type. TFAE for all n:

- 1. X is a smooth n-type-stack.
- 2. Inductively, There merely exists some $U \in \mathbb{T}$ with a map $U \to X$ fibered in smooth n-1-stacks.
- 2' Inductively, as the previous one but additionally the id-types of X are n-1-stacks.
- 3. Inductively, There merely exists some smooth n-1-stack U with a map $U \to X$ fibered in smooth n-1-stacks.
- 3' Inductively, as the previous one but additionally the id-types of X are n-1-stacks.

If one of the conditions is satisfied we call X a smooth n-stack.

Proof. Induction $n-1\mapsto n$, $n\geq 1$.

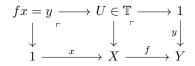
- 1. \Rightarrow 2 We have to show, that the class in 2. is closed under \sum and closed under quotients between *n*-types. This was already done.
- $2. \Rightarrow 3$ Clear
- $3. \Rightarrow 3'$ By 2.7 and independence of the truncation level (TODO).
- 3'. \Rightarrow 3 , 2' \Rightarrow 2 Clear
 - 3'. \Rightarrow 1. by induction, smooth n-1-stacks = smooth n-1-type-stacks \subset smooth n-type-stacks. Now use stability under covers between n-types.
 - $3' \Rightarrow 2'$ Use $3 \Rightarrow 1 \Rightarrow 2$.

2.1 Descent

For this subsection lets assume St a class of sheaves, such that \mathbb{T} is contained in it and for any map $X \to Y$ fibered in \mathbb{T} , $X \in \mathsf{St}$ iff $Y \in \mathsf{St}$. We call types in this class stacky.

Lemma 2.11. Let \mathbb{T} satisfy descent, i.e. beeing affine in the topology has descent. Let $X \in \mathsf{St}$ and Y a type. Let $f: X \twoheadrightarrow Y$ be fibered in \mathbb{T} and surjective. Then $\bullet Y$ is stacky.

Proof. Claim: Y is separated. Proof: By surjectivity of f we may only show that for any x: X, y: Y, the type $fx =_Y y$ is a sheaf. If we define U to be the fiber over y, it is in \mathbb{T} by assumption. But then $fx =_Y y$ is the outer pullback



os.

of stacky types, in particular sheaves.

 $\square(Claim)$

Consider $X \xrightarrow{f} Y \xrightarrow{\eta} \bullet Y$. As X is stacky, it suffices to show, that the fibers are in \mathbb{T} . As beeing affine in \mathbb{T} is a sheaf, we may just show that for all y:Y, the fibers over $\eta y:\bullet Y$ are in \mathbb{T} . As η is a monomorphism by 4.4, η restricts to an equivalence

$$\operatorname{fib}_f y \to \operatorname{fib}_{\eta f}(\eta y)$$

But the left hand side is in \mathbb{T} by assumption.

Theorem 2.12. Assume \mathbb{T} have descent. Then St is a sheaf.

Proof. St is seperated: This follows from the embedding St into the seperated (TODO) type of sheaves

Let $U \in \mathbb{T}$ and $P : ||U|| \to \mathsf{St}$. We want to construct a filler



Claim: $\bullet(\sum_{x:||U||} Px)$ is stacky.

Proof. of the claim. We want to apply the previous lemma to the map

$$\sum_{x:U} P|x| \to \sum_{x:||U||} Px$$

The domain is in St by stability under Σ . The fibers are equivalent to $U \in \mathbb{T} \subset \mathsf{St}$.

The claim provides the map $1 \to St$. The diagram commutes: Assuming $x : \|\operatorname{Spec} A\|$ we wish to show $Px = \sum_{x:\|U\|} Px$. Using univalence, we may show that the maps

$$Px \to \sum_{x: \|U\|} Px \overset{\eta}{\to} \bullet \sum_{x: \|U\|} Px$$

are both equivalences. The first one is an equivalence as ||U|| is contractible. Hence the middle term is a sheaf, thus the unit map is an equivalence as well.

Corollary. For all $n : \mathbb{N} \cup \{\infty\}$, the class of (smooth) (n-)stacks satisfy descent.

3 Saturated Topologies

Consider a topology \mathbb{T} finer than the Zariski topology.

Definition 3.1. A smooth atlas of X is some $\hat{X} \in \mathbb{T}, \hat{X} \to X$ T-cover

Definition 3.2. \mathbb{T} is saturated if Beeing in the topology descents along \mathbb{T} -covers between affines, i.e. every affine schemes that has a smooth atlas lies itself in \mathbb{T} . The saturated closure of a topology \mathbb{T} is the topology \mathbb{T}' defined by (todo finite sums of?)

 $X \in \mathbb{T}'$ iff X is affine $\wedge \exists$ smooth at as of X

Lemma 3.3. Using ZLC, this is the smallest saturated topology containing \mathbb{T} .

Proof. Obviously $1 \in \mathbb{T}'$. Types which have a smooth atlas are stable by dependent sums by the proof of $\ref{thm:proof:equation}$. For the saturatedness consider some \mathbb{T}' -cover $\mathbb{T}' \ni X' \to X$. By replacing X' with some smooth atlas, we may assume that $X' \in \mathbb{T}$. As every fiber $X'_x \in \mathbb{T}'$, we merely find a smooth atlas $\tilde{X}'_x \to X'_x$. Then by Zariski local choice there exists a Zariski atlas $\hat{X} \to X$ and a commutative diagram

$$Y \equiv \sum_{x:\hat{X}} \tilde{X}'_x \longrightarrow \sum_{x:X} X'_x = X'$$

$$\downarrow \qquad \qquad \downarrow$$

$$\hat{X} \xrightarrow{Zax} X$$

As $X' \in \mathbb{T}$ and $Y \to X'$ is fibered in \mathbb{T} (6.3) we have $Y \in \mathbb{T}$. But $Y \to \hat{X}$ is a \mathbb{T} -cover and $\hat{X} \to X$ is a \mathbb{T} -cover, $Y \to X$ is a \mathbb{T} -cover. Hence $X \in \mathbb{T}'$.

Lemma 3.4. A type T is a sheaf wrt to \mathbb{T}' iff it is a sheaf wrt to \mathbb{T}

Proof. As $\mathbb{T} \subset \mathbb{T}'$ the \to direction is clear. Now, let $X \in \mathbb{T}'$. We have to show that $T \to T^{\|X\|}$ is an equivalence. Choose $\mathbb{T} \ni Y \to X$. Then we have a commutative diagram

$$T \xrightarrow{\simeq} T^{\|X\|} \downarrow \\ T^{\|Y\|}$$

So $T \to T^{\|X\|}$ has a left-inverse. Thus it suffices to show that any $f: T^{\|X\|}$ has a preimage. Choose t: T, s.th. cnst_t^Y is the composite $\|Y\| \to \|X\| \stackrel{f}{\to} T$. We have $\|Y\| \to (\operatorname{cnst}_t^X = f)$. But as $Y \in \mathbb{T}$ and $\Delta_t = f$ is a sheaf (as an identity type in the sheaf $T^{\|X\|}$) we are done. \square

Remark 1. We never used that we only talk about T-covers.

Lemma 3.5. Every saturated affine (i.e. Spec $A \in \mathbb{T}'$) it \mathbb{T} -merely inhabited.

Proof. We have $||X|| \to ||\operatorname{Spec} A||$ for some smooth atlas $\mathbb{T} \ni X \to \operatorname{Spec} A$.

Question 1. Does the converse hold, i.e. is every T-merely inhabited affine saturated?

4 Lex Modalities

Lemma 4.1 (Stability resuls). Lex Modalities are stable under

- 1. Conjunction
- 2. Composition

Lemma 4.2. Let \bigcirc be a lex-modality. Let X be \bigcirc -modal and $B: X \to \mathcal{U}_{\bigcirc}$ be a family of modal types. Then $\sum_{x:X} B_x$ is \bigcirc -modal

Lemma 4.3. Let $B: \bullet X \to \mathcal{U}$. Then $\bullet(\sum_{x:X} B(\eta x)) = \sum_{x:\bullet X} \bullet Bx$

Proof. Observe that

$$\sum_{x:X} B\eta x \to \sum_{x:\bullet X} Bx$$

is a \bullet -equivalence, because for all modal types T, the type $Bx \to T$ is modal for any $x : \bullet X$. Then it follows by [ref?].

Lemma 4.4. For a type X the following are equivalent:

- \bullet the identity types of X are sheaves
- the unit $X \to \bullet X$ is a monomorphism

In this case we call X seperated

5 Atlas

Definition 5.1. A \mathbb{T} -atlas of X is a \mathbb{T} -cover Spec $A \to X$ out of an affine scheme.

Remark 2. Any good enough TODO scheme has a Zariski atlas. If \mathbb{T} is finer than the Zariski-topology then in the definition we may replace affine scheme by good enough scheme, if its just about the question whether a type admits an atlas.

Example 5.2. Let X be a (1-)type. X has a Zariski-atlas, iff there exists some $f : \operatorname{Spec} A \to X$ fibered in types of the form $\operatorname{Spec}(R_{f_1} \times \ldots \times R_{f_n})$ for $(f_1, \ldots, f_n) \in Um(R)$.

Remark 3. If one applies ZLC to an affine scheme Spec A the resulting principal open cover $D(f_i), f_i \in A$ will induce indeed a zariski atlas $\bigsqcup D(f_i) \to \operatorname{Spec} A$, because the fiber over $x : \operatorname{Spec} A$ is $\bigsqcup D(f_i(x))$.

Question: Does every zariski atlas of Spec A have this form? Weird Zariski Atlasses

Example 5.3. \mathbb{P}^n has a zariski atlas given by the standart homogeneous principal opens $\sum_{i=0}^{n} D_{+}(x_i)$. The fiber over a point $[y_0 : \ldots : y_n]$ is $D(y_0) + \ldots D(y_n)$ where $(y_1, \ldots, y_n) \in Um(R)$.

6 Local Choice

In this section let \mathbb{T} denote a topology finer than the zariski topology.

Definition 6.1. Let Cov be a class of morphisms (which we think of n-atlasses of some n), containing \mathbb{T} -atlas, (stable under pullback NECESSARY TODO?) A type S has $local \ choice$ wrt Cov if for any \mathbb{T} -surjective map $X \to Y$ and any map $f: S \to Y$ there exists a map $p': S' \to S$ in Cov and a commutative diagram

$$S' \xrightarrow{p} X$$

$$\downarrow \qquad \qquad p \downarrow$$

$$S \xrightarrow{f} Y$$

Proposition 6.2. Assume that Cov is stable under composition.

- If $\hat{S} \to S$ is a Cover and \hat{S} has \mathbb{T} -local choice, then S has \mathbb{T} -local choice.
- Affine schemes have \mathbb{T} -local choice.
- Any type admitting a Cov Atlas Spec $A \to S$ has \mathbb{T} -local choice.

Proof. The first point follows from stability under composition of Cov. the third point follows from the second. By the first point, we may assume that S is affine. As p is \mathbb{T} -surjective, for any x:S there merely is a $\operatorname{Spec} B_x \in T$ and a map $\operatorname{Spec} B_x \to \|\operatorname{fib}_p(x)\|$. As S is projective, we have a term in

$$\prod_{x:S} \sum_{\operatorname{Spec} B_x \in T} \operatorname{Spec} B_x \to \|\operatorname{fib}_p(fx)\|$$

By setting

$$(S' := \sum_{x:S} \operatorname{Spec} B_x) \xrightarrow{\pi} S$$

the projection, we are now in the situation that for any t:S' we merely have a point in $\mathrm{fib}_p((p'(t)))$ and $S'\to S$ is a \mathbb{T} -cover, thus it is in Cov. Moreover, S' is affine, as it is a dependent sum of affines. Hence again we now can find a lift $S'\to X$ making

$$S' \longrightarrow Y$$

$$\downarrow p' \downarrow p \downarrow$$

$$S \longrightarrow X$$

commute.

The next lemma shows, that the class of types equipped with a T-atlas is stable under dependent sums.

Lemma 6.3. Let $\mathcal{U}' \subset \mathcal{U}$ be stable under dependent sums (e.g. \mathbb{T} -inhabited types) Let X be a type with a map $p: U \to X$ fibered in \mathcal{U}' . For any x: X, let Y_x be a type and moreover for any u: U, we are given a map $q_u: V_u \to Y_{p(u)}$ fibered in \mathcal{U}' . Then the induced map

$$p: \sum_{u:U} V_u \to \sum_{x:X} Y_x$$

is fibered in \mathcal{U}'

Proof. The fiber of p over some $(x,y) \in \sum_{x:X} Y_x$ is given by

$$\sum_{u: \mathrm{fib}_p} \mathrm{fib}_{q_u}(y')$$

where $y': Y_{p(u)}$ (depending on u) is the transport of $y: Y_x$ along x = p(u). As \mathcal{U}' is stable under dependent sum those fibers are again in \mathcal{U}' . This shows the result.

Theorem 6.4. Let \mathcal{U}' be a class stable under dependent sums. The class of types admitting a \mathcal{U}' -atlas is closed under dependent sums. If \mathbb{T} is a topology, the same holds for smooth \mathcal{U}' -atlasses (i.e. with domain in \mathbb{T}).

Proof. Let us construct some atlas Spec $A \to \sum_{x:X} B_x$ For any x:X we merely have an atlas $V_x \to B_x$, i.e. with V_x affine. X has local choice wrt atlasses by (6.2) using \mathcal{U}' is \sum -stable (we use the trivial topology).

If additionally, all the B_x and X are smooth n-stacks, just observe that we can choose the affine V_{pu} to lie in \mathbb{T} , Accordingly $\sum_{u:U} V_{pu} \in T$ as \mathbb{T} is stable under Σ . By Local choice for X, we merely find U affine, an atlas $p:U\to X$ with

$$\prod_{u:U} \sum_{V_{p(u)} \in T} (q: V_{p(u)} \to B_{p(u)}) \times (q \text{ fibered in smooth } n \text{ stacks })$$

Now the desired map is $\sum_{u:U} V_{pu} \to \sum_{x:X} B_x$, because it is an atlas by 6.3

9

7 Algebraic Space

We first need to define a notion of algebraic space and smooth algebraic space, which should be the smallest class of types that satisfies the following:

- Stability under finite limits 9.1
- has Descent
- (nice) Schemes are contained in it
- \bullet affines in \mathbb{T} are smooth algebraic spaces. (there are probably more).
- stable under smooth quotients: If X is an algebraic space, Y modal 0-type and $X \to Y$ is \mathbb{T} -surjective and fibered in smooth algebraic spaces, then Y is an algebraic space. Additionally, if X is smooth, then Y is smooth.

Definition 7.1. An affine Scheme U is called flat, if

$$||U||_{\mathbb{T}} \to (U \in \mathbb{T})$$

Lemma 7.2. The converse holds always

Proof. because things in \mathbb{T} are automatically \mathbb{T} -merely inhabited

Recall the definition of T-atlas 5.1

Definition 7.3. We call a modal proposition algebraic, if one of the equivalent conditions is satisfied:

- 1. its merely of the form $||U||_{\mathbb{T}}$ for some flat affine U.
- 2. There is a \mathbb{T} -surjective map out of a flat affine U.
- 3. It has a T-atlas.

Proof.

 $1 \Leftrightarrow 2$ Clear.

- $1 \Rightarrow 3$ we show that $U \to ||U||_{\mathbb{T}}$ is a T-atlas. Every fiber is in T, because U is flat.
- $3 \Rightarrow 1$ Let $V \to P$ be a \mathbb{T} -atlas. have to show TFAE $||V||_{\mathbb{T}} \to P \to (V \in \mathbb{T}) \stackrel{7.2}{\to} ||V||_{\mathbb{T}}$. Proof: $||V||_{\mathbb{T}} \to P$ as P is modal prop. Secondly, because $V \to P$ is a \mathbb{T} -cover. Hence P is an algebraic proposition.

Lemma 7.4. Algebraic propositions are algebraic spaces.

Proof. We have $U \to ||U||_{\mathbb{T}}$ where U is affine, hence an algebraic space and the fibers are in \mathbb{T} by flatness of U, hence they are smooth algebraic spaces. By stability under quotients, our algebraic proposition is an algebraic space.

Definition 7.5. An smooth equivalence relation on a set U is some equivalence relation $R: U^2 \to \text{Prop}$, whose fibers are in \mathbb{T}

Lemma 7.6. let U be an algebraic space (e.g. affine scheme) and $R: U^2 \to \text{Prop be a}$ smooth equivalence relation Then U/R is an algebraic space

Proof. The map $U \to U/R$ is fibered in \mathbb{T} , in particular fibered in smooth algebraic spaces. By stability under smooth quotients, U/R is an algebraic space.

Corollary. Let U be affine and R a smooth equivalence relation. The identity types of U/R, i.e. the propositions R(x,y), are algebraic propositions.

Proof. By 14.2, the class of types admitting a \mathbb{T} -atlas is closed under taking identity types. U/R is a type admitting a \mathbb{T} -atlas, hence its identity types admit them as well.

Definition 7.7. A modal set X is a classical algebraic space iff it is merely of the form U/R for some affine U and $R:U^2\to \operatorname{Prop}$ a smooth equivalence relation. Equivalently there exists some \mathbb{T} -atlas $U\to X$ (i.e. out of an affine). We call X smooth if U can be choosen to be in \mathbb{T} .

Corollary (of 7.3). Classical Algebraic spaces that are propositions are algebraic propositions.

Remark 4. Assume Saturatedness of the topology. smooth classical Algebraic spaces which are affine are in \mathbb{T} .

Question 2. Is the class of classical algebraic spaces stable under smooth quotients? If its not, how should we enlarge it?

Try: Assume R is fibered in smooth algebraic spaces. Choose $U \to T$ a \mathbb{T} -atlas. For any x:U the fiber R_x merely has an atlas $\tilde{R}_x \to R_x$. As U has choice (its affine), we find some \mathbb{T} -cover $\tilde{U} = \sum_x \tilde{R}_x \to \sum_x R_x$. Goal: Find for all t:U/R a \mathbb{T} -atlas $V_t \to \mathrm{fib}_{\parallel}(t)$. Then $\sum_t V_t$ will be affine, because its the total space of a \mathbb{T} -cover of an affine. Moreover, $\sum_t V_t \to \sum_t \mathrm{fib}_{\parallel}(t) \to U/R$ will be a \mathbb{T} -cover, as $V_t \in \mathbb{T}$. This is what we wanted to show.

8 *n*-stacks

Definition 8.1. Let \mathbb{T} be a subcanonical topology finer than the Zariski topology. Let $n \geq -2$. A type X

- is a (smooth) -2-stack if it is contractible
- is A (n+1)-stack, if
 - -X is a \mathbb{T} -sheaf
 - For any $x, y : X \ x =_X y$ is a *n*-stack
 - There exists an n-atlas, i.e. a T-surjective map $\operatorname{Spec} A \to X$ fibered in
 - * \mathbb{T} , if n < 0
 - * smooth n-stacks, if n > 0.
- X is a smooth n+1-stack if
 - -X is a (n+1)-stack
 - There exists a *n*-atlas Spec $A \to X$ with Spec $A \in \mathbb{T}$

Lemma 8.2. One could only alternatively talk about (smooth) n-stacks for $n \ge 1$, define them by induction as above. Then later define:

- A (smooth) -1-stack is a (smooth) 1- stack is a proposition.
- A (smooth) 0-stack is a (smooth) 1- that is a 0-type.

Proof.

Lemma 8.3. A (smooth) n-stack is a (smooth) n + 1-stack.

Proof. Induction. Be aware of the induction start, where maybe no atlas is assumed! We need, that \mathbb{T} is subcanonical to conclude that affines are \mathbb{T} -sheaves.

Remark 5. If one changes the definition of atlas to be a map out of a scheme, then smooth -1 atlas will be scheme in T. Otherwise propositional -1-stack are not 0-stacks.

9 Stability results

Theorem 9.1. Let $n \ge -2$. Smooth / n-stacks are stable by dependent sums.

Proof. Induction. For n=-2 its okay. Let $B:X\to \mathcal{U}$ be a family of n+1-stacks indexed over a n+1-stack X, then surely the total space $\sum_{x:X}Bx$ is a \mathbb{T} -sheaf as \mathbb{T} -sheaves are stable under dependent sum. The identity types in a \sum type are \sum of identity types. Admitting an n-atlas is stable under dependent sum: We apply 6.4 to the class of (smooth) n-atlasses, which is stable under depent sum by induction.

Corollary. n-atlasses are stable under composition.

Lemma 9.2. n+1-stacks are closed under taking closed (open) subtypes.

Proof. First we show:if X has an n-atlas and Y is a closed (open) subtype of X, then Y has an n-atlas. Choose an n-atlas Spec $A \to X$. The pullback to Y has have the same fibers. If Y is closed, and the total space is a closed subtype of Spec A, hence it will be affine. if Y is an open subtype of X, then the pullback is an open subtype of Spec A, hence by zariski local choice merely of the form $\bigcup_{i=1}^n D(a_i) \subset A$. As n-atlasses are stable under composition 9, it suffices to show, that the map $f: \bigsqcup_i D(a_i) \to \bigcup_{i=1}^n D(a_i)$ is a Zariski-atlas, because then it will be an n-atlas as well. Let $x: \bigcup_{i=1}^n D(a_i)$, i.e. there merely exists an i, such that $a_i(x)$ is invertible. The fiber is exactly $D(a_1(x)) + \ldots + D(a_n(x))$. thus we are done. (MAYBE OUTSOURCE THIS and say open subschemes of affines have zariski atlas)

Corollary. Let X be a quasi-projective scheme that is a sheaf. Then X is a 0-stack.

Proof. It suffices to see that X has a zariski atlas. Use .

Definition 9.3. A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

Lemma 9.4. Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

Lemma 9.5. Given a local property P of morphisms of modal n-types, a morphism $f: X \to Y$ has P if there exists an n-atlas of f having P.

The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n = 0, 1.

10 Descent

Theorem 10.1. Let T be a modal n-type. The Proposition, that P is a (smooth) n-stack, is modal.

11 Fundamental Theorem of algebraic spaces

11.1 For groupoids

Lemma 11.1. If R woheadrightarrow X o X is a \mathbb{T} -htpy-coequalizer diagram of two \mathbb{T} -covers between affines, then X is a 1-stack.

11.2 For sets

Lemma 11.2. Denote $\mathbb{T}Set$ for the sets that are \mathbb{T} -sheaves. Assume given a $\mathbb{T}set\ X$ then the following maps are mutually inverse

$$\begin{split} \sum_{R:X\to X\to\mathbb{T}\,\text{Prop}} R \text{ equivalence relation} &\simeq \sum_{Y:\mathbb{T}\mathsf{Set}} \sum_{p:X\to Y} p \text{ } \mathbb{T} \text{surjective} \\ R &\mapsto (X/R,[\ \]) \\ \lambda x,y.(p(x)=p(y)) &\leftarrow (Y,p) \end{split}$$

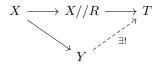
where X/R is defined by applying $L_T\|_{-}\|_0$ at the higher inductive type X//R.

- *Proof.* Well-definedness: The map $[\cdot]: X \to ||X//R||_0 \to L_T ||X//R||_0$ is the composition of a surjective with a \mathbb{T} -surjective map [ref?], hence its \mathbb{T} -surjective. Conversely given (Y,p) as Y is a sheaf, we have for all x,y:X that $p(x)=_Y p(y)$ is a sheaf.
 - If x, y : X then we have a chain of equivalences

$$R(x,y) \simeq (\bar{x} =_{\|X//R\|_0} \bar{y}) \to ([x] =_{L_T\|X//R\|_0} [y])$$

where the first map is plain HoTT and the second map is ap, i.e. the unit of the modality [ref?], but as the $\bar{x} = \|X//R\|_0$ \bar{y} is already a sheaf, it is an isomorphism as well.

• Let (Y,p) be in the RHS. Let $R(x,y)=(p(x)=p(y)):\mathbb{T}$ Prop. By plain HoTT, There is a map $\eta:X//R\to Y$ (defined by the universal property of the set truncation and by induction on the higher inductive type X//R on canonical terms through the map $p:X\to Y$). I claim η exhibits Y as the localization for \mathbb{T} Set-modality of X//R. Let T be another \mathbb{T} Set equipped with a map $X//R\to T$. By precomposition we obtain a map $X\to T$. Claim: it factors uniquely through $p:X\to Y$.



Proof:

Existence: We want to define a map $Y \to T$. Let y: Y. As p is \mathbb{T} -surjective and T is a sheaf, we may assume we merely have some element in the fiber of p over y. Now push this element through

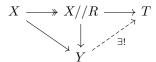
$$\|\operatorname{fib}_p y\| \to \|X//R\|_0 \to T$$

where the first map is by Plain HoTT and the second one is induced from $X//R \to T$ by assumption and the fact that T is a set.. One can easily check this makes the diagram commute. Uniqueness follows from $X \to Y$ beeing \mathbb{T} -surjective and the following Fact: Two parellel maps $Y \rightrightarrows T$ into a \mathbb{T} Set T are already equal if the become equal after

precomposition with a \mathbb{T} -surjection $X \to Y$.

Proof of the fact: Let y:Y. The goal is an identity type of a \mathbb{T} Set, hence a \mathbb{T} Prop. Hence As the fiber over y in X is \mathbb{T} -merely inhabited, we may assume an actual term in the fiber. As $X \to Y$ equalizes the arrows, this term allows us to conclude. \Box (fact) \Box (Claim)

We apply the fact to the (T-)surjectivity of $X \to X//R$ to get a unique factorization



П

making the right triangle commute. This is what we wanted to show.

Definition 11.3. An equivalence relation R on a type X is called:

- redundant if for all x, y : X the proposition R(x, y) is a -1-stack.
- smooth if its and for any y: X its fibers:

$$R_y :\equiv \sum_{x:X} R(x,y)$$

are affine in \mathbb{T} .

Lemma 11.4. Assume that \mathbb{T} satisfies descent for propositions and for sets 10.1, i.e. that a modal proposition being a (-1)-stack is a sheaf. Assume that a modal set beeing affine in \mathbb{T} is a sheaf. Assume given a \mathbb{T} set X, then the following types are equivalent:

- ullet The type of redundant smooth equivalence relations over X.
- The type of Tsets Y with identity types beeing stacks and an -1-atlas X to Y (in V2 a T-cover).

Proof. By the equivalence in 11.2, it is enough to check that:

• The identity types in X/R are (-1)-stacks if and only if the relation R is redundant . For any x,y:X we know that:

$$R(x,y) \simeq [x] =_{X/R} [y]$$

so the direct direction is immediate. For the converse we use the assumption that a modal proposition being a (-1)-stack is a sheaf and that the map [$_$]: $X \to X/R$ is \mathbb{T} -surjective.

• The fibers of:

$$[_]: X \to X/R$$

are affine in \mathbb{T} if and only if the relation R is smooth. For any y:X we have that:

$$\sum_{x \in X} R(x, y) \simeq \mathrm{fib}_{[-]}([y])$$

so the direct direction is immediate. Here as well the converse follows from \mathbb{T} -surjectivity of $[_]$ and that the topology has descent.

Corollary. Assume \mathbb{T} satisfies descent for propositions and for sets. A type is a 0-stack iff its merely the \mathbb{T} -quotient of an affine scheme by a smooth equivalence relation.

Theorem 11.5. Assume \mathbb{T} satisfies descent for propositions. The quotient of a 0-stack $X \in \mathbb{T}$ Set by an 0-smooth equivalence relation R is a 0-stack. TODO

Proof. The identity types in X/R are propositional 0-stacks, hence (-1)-Truncations of -1-stacks by 13.2 as desired.

How to find an atlas: todo. How to proceed, if we could choose all atlasses we want at the same time?

Remark 6. This is equivalent to saying that 1-stacks that are 0-types are geomeric 0-stacks: One direction we prove later. If R is a 0-smooth equivalence relation on a 0-stack X, then X/R is a 1-stack by observing that any -1-atlas $X' \to X$ gives a 0-atlas $X' \to X \to X/R$. Moreover, X/R is a 0-type, hence by assumption a 0-stack.

Example 11.6. There are open affine subschemes U of affine schemes $\operatorname{Spec} A$, which are not (disjoints unions of) principal open

Proof. Consider $A = R[x, y, u, v]/(xy + ux^2 + vy^2), X = \operatorname{Spec} A$ and consider the open U = D(x, y).

We cant expect U to be a disjoint union of principal opens (todo). However, D(x,y) is affine: We have maps $U \to R$ given by $f = -v/x = (y + ux)/y^2, g = -u/y = (x + vy)/x^2$. Then $D(f) \cup D(g) = \operatorname{Spec} R^X$, as yf + xg = 1 in R^U . Taking preimages under the affinization map, $U_f \cup U_g = X$ and one checks this defines an open affine cover (for example : $U_f \simeq \operatorname{Spec} R[x, u, f^{\pm 1}, g]/(xy + ux^2 + uy^2)$ with y := (1 - gx)/f.) But on both of this open subsets the affinization map is an isomorphism hence the affinization of X is an isomorphism. compare (Hartshorne II.2.17)

Lemma 11.7. Let $f: X \to Y$ be surjective. There exists a Zariski Cover $X' \to X$ such that $X' \to Y$ is a Zariski cover iff there exists a Zariski Cover $X' \to X$, some $n: \mathbb{N}$ and an open affine embedding $X' \hookrightarrow Y^n$ over Y.

12 Saturated Topologies revisited

Lemma 12.1 (1.1). We want that every n-1-atlas of a smooth n-atlas has the additional requirement in the definition of smooth n-atlas. It turns out, that for this topology needs to be saturated: The following are equivalent

- 1. Beeing in the topology descents along \mathbb{T} -covers between affines, i.e. \mathbb{T} is saturated.
- 2. A smooth n -stack X that is an affine scheme lies in the Topology \mathbb{T} .
- 3. Let $n \geq 0$. If T is a smooth n-stack, then any n-1-atlas $U \to T$ satisfies $U \in \mathbb{T}$.
- 4. If $U \xrightarrow{f} V \xrightarrow{g} W$ are maps between affines and f and gf are \mathbb{T} covers, then g is a \mathbb{T} Cover

Proof. $1 \Rightarrow 2$

Induction. This holds for n=-1. Assume it holds for n-1. Choose a n-1-atlas with T source, i.e. $T \ni \operatorname{Spec} A \to X$ fibered in smooth n-1-stacks. As it is affine, all the fibers of the atlas are affine smooth n-1-stacks, hence by induction they lie in \mathbb{T} , thus the atlas is a \mathbb{T} -cover between affines, hence $X \in \mathbb{T}$.

 $2 \Rightarrow 3$

As $U \to T$ is fibered in smooth n-1 stacks, all the fibers are in particular smooth n-stacks by 8.3. By stability under dependent sum $U = \sum_{t:T} U_t$ is a smooth n-stack that is affine, hence by assumption (2) it lies in the topology.

 $3 \Rightarrow 1$

Let $X \to Y$ be a \mathbb{T} -cover with X affine in \mathbb{T} and Y affine. Then Y is a smooth 0-stack, But $Y \to Y$ is a -1-atlas, hence by assumption $Y \in T$.

 $4 \Rightarrow 1$

Obvious

 $1 \Rightarrow 4$

Check fiberwise \Box

If $n \geq$, replacing \mathbb{T} by its saturation \mathbb{T}' does change the notion of (smooth) n-stack, but we have the following statement, that tells us, that if we start with 0- \mathbb{T} -stacks then the notion of smoothness does not see the difference between \mathbb{T} and its saturation.

Proposition 12.2. Let X be a 0-stack that is a weak smooth 0-stack, i.e. there exists a \mathbb{T}' -atlas $\mathbb{T}' \ni X' \to X$ (i.e. fibered in \mathbb{T}'). Then X is a smooth 0-stack.

Proof. Wlog $X' \in \mathbb{T}$. Choose a -1-atlas Spec $A \to X$ (i.e. fibered in \mathbb{T}). As the fibers of $X' \to X$ merely have smooth atlasses $\tilde{X}'_x \to X'_x$, we can use Local choice to obtain a commutative diagram $Y = \sum_{x':X'} \tilde{X}'_x$

$$\tilde{X} \xrightarrow{\mathbb{T}} \operatorname{Spec} A$$

$$\mathbb{T} \downarrow \qquad \qquad \downarrow$$

$$X' \xrightarrow{\mathbb{T}'} X$$

As $Y \to X'$ is a \mathbb{T} -cover and $X' \in \mathbb{T}$ we conclude $Y \in \mathbb{T}$. Hence we found a smooth \mathbb{T} -atlas of X.

12.1 Zariski Topology is not saturated

Example 12.3 (Weird Zariski Atlasses). Assume those equivalent conditions on the Zariski topology. There exist Zariski atlasses of affines $\operatorname{Spec} A = X$ which are not of the form $D(a_1) + \ldots + D(a_n) \to \operatorname{Spec} A$ for $(a_1, \ldots, a_n) \in Um(A)$

Proof. Indeed, using the first example, choose $U \subset \operatorname{Spec} A$ affine not principal open, then choosing a Zariski atlas $V \to U$ gives $V + X \to U + X \to X$ where $V + X \to X$ is a Zariski cover and $V + X \to U + X$ is a Zariski cover. From (4), we deduce that $U + X \to X$ is a Zariski cover, but U is not a disjoint union of principal opens in $\operatorname{Spec} A$.

Example 12.4. Assume those equivalent conditions on the Zariski topology. Every affine open proposition U is principal open!

Proof. Let $V \to U$ be a Zariski atlas. Then $V+1 \to U+1$ is a Zariski atlas with $V+1 \in \mathbb{T}$ and U+1 affine, hence by (1) $U+1 \in \mathbb{T}$, hence U is a disjoint union of principal opens hence, as it is a proposition, its principal open.

13 beeing a stack is indepent of the truncation level

Lemma 13.1. Let $n \ge 0$. A n-stack is an modal n-type.

Proof. The n-Truncation an n-type. Now conclude by induction.

We want to show that the notion of stack makes sense, i.e. beeing a stack should not depend on the truncation level.

Lemma 13.2. Assume \mathbb{T} is saturated and satisfies descent for propositions. Let P be a modal proposition. Then TFAE

- 1. For some m > 0, P is a m stack
- 2. There exists some fp algebra A such that $\operatorname{Spec} A \to P$ and P is logically equivalent to $(\operatorname{Spec} A \in \mathbb{T})$.
- 3. P is equivalent to $\|\operatorname{Spec} A\|_{\mathbb{T}}$ for some fp A, i.e. P is a -1-stack.

Proof.

- 1. ⇒ 2. Let Spec $A \to P$ be a m-1 atlas. Assume Spec $A \in \mathbb{T}$. Then $\|\operatorname{Spec} A\| \to P$ so as P is a sheaf, we have P. Conversely, if x:P, then the fiber over x is Spec A and a smooth m-1 stack, hence belongs to the topology by 12.1.
- 2. \Rightarrow 3. We have to show: There exists some flat algebra such that P is logically equivalent to $\|\operatorname{Spec} A\|_{\mathbb{T}}$. By assumption we have $\operatorname{Spec} A \to P \to (\operatorname{Spec} A \in \mathbb{T})$, so we deduce $\|\operatorname{Spec} A\|_{\mathbb{T}} \to P \to (\operatorname{Spec} A \in \mathbb{T})$, as P is a modal proposition. In particular A is flat. Conversely $P \to (\operatorname{Spec} A \in \mathbb{T}) \to \|\operatorname{Spec} A\|_{\mathbb{T}}$, where the first arrow is by assumption.

 $3. \Rightarrow 1. 8.3$

Lemma 13.3. A smooth -1-stack P is contractible.

Proof. Choose a \mathbb{T} -cover $\mathbb{T} \ni \operatorname{Spec} A \to P$. As P is a proposition we have $\|\operatorname{Spec} A\| \to P$. As P is a sheaf we have P.

Example 13.4. A 0-stack is a \mathbb{T} -sheaf whose identity types are (-1)- \mathbb{T} runcations of ((affine ?)) schemes and there exists a \mathbb{T} -atlas Spec $A \to X$.

Why are schemes 0-stacks? This holds in special case, for example if the scheme is quasi projective.

Theorem 13.5. Let \mathbb{T} be saturated. Assume the topology satisfies descent Let $m, n \geq -2$. Given an n-type T that is a (smooth) m-stack then T is a (smooth) n-stack.

Proof. By 8.3 we may assume $m \ge n \ge -2$.

If $m \le 1$ this is clear. Now assume $m \ge 2$. Induction. Inductionstart m = 2. Let us prove the case of m = 2, n = 1, the cases $-2 \le n < 1$ are immediate from this.

Choose a 1-atlas $X' \to T$, i.e. its fibered in smooth 1-stacks. As T is a groupoid and X' is a set, the fibers are actually sets, i.e. smooth 0-stacks.

Now consider $R := X' \times_T X'$. As X' is in particular a 0-stack and 0-stacks are stable under dependent sums, R will be a 0-stack. Choose a a \mathbb{T} -cover $R' \to R$ with R' affine. Now $R' \to R \to X'$ is a map between affine schemes i.e it is fibered in smooth 0-stacks that are affine. As \mathbb{T} is saturated, the fibers of $R' \to X'$ are in \mathbb{T} . As X'//R' is a 1-stack by ??, it suffices to show that $X'//R' \to X'//R$ is a \mathbb{T} -cover. Pick a term in X'//R. As the fiber beeing in \mathbb{T} is sheaf If additionally T is assumed to be a smooth 2-stack, then we can assume X' to be in the topology. This will force R to be a smooth 0-stack, so we may choose R' Assume m > 2 and the statement is proven for all (n', m') < (n, m) in lexicographical ordering. As the identity types of T are n-1-types and m-1 stacks by induction they are n-1 stacks. Let $X \to T$ be an m-1-atlas, i.e. fibered in smooth m-1-stacks with X affine. The fibers are in particular n-1-types, so by induction they are smooth n-1-stacks. Hence $X \to T$ is an n-1-atlas. If, additionally T is assumed to be a smooth m-stack, we can choose $X \in \mathbb{T}$, hence $X \to T$ witnesses that T is a smooth n-stack.

14 Stability under Quotients

Definition 14.1. A morphism between n-stacks is smooth if it is fibered in

- \mathbb{T} if $n \leq 0$
- smooth n-stacks if n > 0.

Lemma 14.2. Let C be a class of types stable under \sum . The class $\mathsf{HasAtlas}_C$ of types Y which admit a map $\mathsf{Spec}\,A \to Y$ fibered in C is stable under finite limits, i

Proof. Obviously 1 has an atlas, and the class of types admitting an atlas is stable by \sum by 6.4. It remains to show, that identity types in Y have an atlas provided that Y has an atlas.

By assumption we can choose a map $p:V\to Y$ out of an affine fibered in C. Let y,y':Y. Then we have the map

$$(\operatorname{fib}_p y) \times_V (\operatorname{fib}_p y') \to y = y'$$
$$(v, q : y = pv), (v', q' : y' = pv'), (h : v = v') \mapsto q \cdot h \cdot q'^{-1}$$

The fiber over j: y = y' looks like

$$\sum_{v} (\underbrace{\sum_{v'} (h:v=v'))}_{\mathsf{isContr}} \times (q:y=pv) \times (q':y'=pv') \times (q\cdot h\cdot q'^{-1}=j) \simeq \sum_{v} (v=py) \simeq \mathrm{fib}_p \, y$$

Hence the map is fibered in C. It suffices to show, that $(\operatorname{fib}_p y) \times_V (\operatorname{fib}_p y')$ has an atlas, because then we can compose such an atlas with the above map to obtain an atlas of y = y'. By assumption the fibers of p have an atlas, so we can choose $q:W\to\operatorname{fib}_p y,q':W'\to\operatorname{fib}_p y'$ atlasses. Then $W\times_V W'\to(\operatorname{fib}_p y)\times_V (\operatorname{fib}_p y')$ is an atlas: The domain is a fiber product of affines, hence affine. The fiber over (x,x') is equivalent to the product of fibers $(\operatorname{fib}_q x)\times(\operatorname{fib}_{q'} x')$ which is in C by stability under dependent sums (so in particular under finite products).

Theorem 14.3. Let $f: X \to Y$ be a \mathbb{T} -surjective smooth morphism between modal n-types. If X is a (smooth) stack, then Y a (smooth) stack.

(*) This can only hold if we define -1-stacks to be modal propositions with a -2-atlas Spec $A \to P$, i.e. algebraic propositions 7.3

Proof. Induction. For n=-2 its clear. Let X be a n-stack. Lets first construct the n-1-atlas of Y. We merely find a $V \twoheadrightarrow X$ which is an n-1-atlas. Then $V \to X \to Y$ is an n-atlas because it is \mathbb{T} -surjective and is fibered in the correct Σ -stable class of types, i.e. \mathbb{T} if $n \le 1$ and smooth n-1-stacks for n > 1. Hence Y is an n+1-stack. As Y is an n-type, Y is an n-stack 13.5.

If additionally X is assumed to be smooth, then V can be assumed to lie in \mathbb{T} which directly gives us that Y has a smooth atlas.

It remains to show that the identity types of Y are n-1-stacks. As Y has an n-1-atlas, by 14.2 we find some n-1-atlas $p:W\to y=y'$. The map is smooth. If n=0, y=y' is a -1-stack by (*). If n>0, W is an n-1-stack and p is smooth, so by induction y=y' is an n-1-stack.

Remark 7 (Using descent but not induction). Hugo suggested an alternative argument proving that the identity types of Y are n-1-stacks, which presumable avoids 13.5 but uses descent for n-stacks: For x:X,y:Y we have that

$$(f(x) = y) \simeq (1 \times_X \operatorname{fib}_f y)$$

is an n-stack by stability under \sum . Because it is an n-1-type, it is a n-1-stack by 13.5. Now conclude that every identity type of Y is an n-1-stack by using descent for n-1-stacks and \mathbb{T} -surjectivity of f.

15 Local properties

Definition 15.1. Let Cov be the property of morphisms of n-stacks defined by asking that the morphism is \mathbb{T} -surjective and fibered in smooth n-stacks. Its stable under basechange. A property of n-stacks is local if P(1) holds, P is stable by dependent sums and given a $Cover\ X \to Y$ we have PX iff PY.

Example 15.2. beeing smooth n-stack is a local property of stacks.

Proof. We have to show: If $f: X \to Y$ is a T-surjective map fibered in smooth n-stacks between n-stacks, then X is a smooth n-stack iff Y is a smooth n-stack. The only if is clear by stability under dependent sums. The other direction is 14.3.

Definition 15.3. A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

Lemma 15.4. Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

Lemma 15.5 ([ref?]). Given a local property P of morphisms of n-stacks, a morphism $f: X \to Y$ has P if there exists an n-atlas of f having P.

Example 15.6. A morphism of n-stacks is smooth iff there exists an n-atlas of f

such that \tilde{f} is a \mathbb{T} -cover.

The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n = 0, 1.