Thesis

Tim Lichtnau

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1 Saturated Topologies

Consider a topology \mathbb{T} finer than the Zariski topology.

Definition 1.1. A smooth atlas of X is some $\hat{X} \in \mathbb{T}, \hat{X} \to X$ T-cover

Definition 1.2. \mathbb{T} is saturated if Beeing in the topology descents along \mathbb{T} -covers between affines, i.e. every affine schemes that has a smooth atlas lies itself in \mathbb{T} .

The saturated closure of a topology \mathbb{T} is the topology \mathbb{T}' defined by (todo finite sums of?)

$$X \in \mathbb{T}'$$
 iff X is affine $\wedge \exists$ smooth at as of X

Lemma 1.3. Using ZLC, this is the smallest saturated topology containing \mathbb{T} .

Proof. Obviously $1 \in \mathbb{T}'$. Types which have a smooth atlas are stable by dependent sums by the proof of $\ref{thm:proof}$. For the saturatedness consider some \mathbb{T}' -cover $\mathbb{T}' \ni X' \to X$. By replacing X' with some smooth atlas, we may assume that $X' \in \mathbb{T}$. As every fiber $X'_x \in \mathbb{T}'$, we merely find a smooth atlas $\tilde{X}'_x \to X'_x$. Then by Zariski local choice there exists a Zariski atlas $\hat{X} \to X$ and a commutative diagram

$$Y \equiv \sum_{x:\hat{X}} \tilde{X}'_x \longrightarrow \sum_{x:X} X'_x = X'$$

$$\downarrow \qquad \qquad \downarrow$$

$$\hat{X} \xrightarrow{Zax} X$$

As $X' \in \mathbb{T}$ and $Y \to X'$ is fibered in \mathbb{T} (4.3) we have $Y \in \mathbb{T}$. But $Y \to \hat{X}$ is a \mathbb{T} -cover and $\hat{X} \to X$ is a \mathbb{T} -cover, $Y \to X$ is a \mathbb{T} -cover. Hence $X \in \mathbb{T}'$.

Lemma 1.4. A type T is a sheaf wrt to \mathbb{T}' iff it is a sheaf wrt to \mathbb{T}

Proof. As $\mathbb{T} \subset \mathbb{T}'$ the \to direction is clear. Now, let $X \in \mathbb{T}'$. We have to show that $T \to T^{\|X\|}$ is an equivalence. Choose $\mathbb{T} \ni Y \to X$. Then we have a commutative diagram

$$T \xrightarrow{\simeq} T^{\|X\|} \downarrow \\ T^{\|Y\|}$$

So $T \to T^{\|X\|}$ has a left-inverse. Thus it suffices to show that any $f: T^{\|X\|}$ has a preimage. Choose t: T, s.th. cnst_t^Y is the composite $\|Y\| \to \|X\| \stackrel{f}{\to} T$. We have $\|Y\| \to (\operatorname{cnst}_t^X = f)$. But as $Y \in \mathbb{T}$ and $\Delta_t = f$ is a sheaf (as an identity type in the sheaf $T^{\|X\|}$) we are done. \square

Remark 1. We never used that we only talk about T-covers.

Lemma 1.5. Every saturated affine (i.e. Spec $A \in \mathbb{T}'$) it \mathbb{T} -merely inhabited.

Proof. We have $||X|| \to ||\operatorname{Spec} A||$ for some smooth atlas $\mathbb{T} \ni X \to \operatorname{Spec} A$.

Question 1. Does the converse hold, i.e. is every T-merely inhabited affine saturated?

2 Lex Modalities

Lemma 2.1 (Stability resuls). Lex Modalities are stable under

- 1. Conjunction
- 2. Composition

Lemma 2.2. Let \bigcirc be a lex-modality. Let X be \bigcirc -modal and $B: X \to \mathcal{U}_{\bigcirc}$ be a family of modal types. Then $\sum_{x:X} B_x$ is \bigcirc -modal

Lemma 2.3. Let $B: \bullet X \to \mathcal{U}$. Then $\bullet(\sum_{x:X} B(\eta x)) = \sum_{x:\bullet X} \bullet Bx$

Proof. Observe that

$$\sum_{x:X} B\eta x \to \sum_{x:\bullet X} Bx$$

is a \bullet -equivalence, because for all modal types T, the type $Bx \to T$ is modal for any $x : \bullet X$. Then it follows by [ref?].

3 Atlas

Definition 3.1. A \mathbb{T} -atlas of X is a \mathbb{T} -cover Spec $A \to X$ out of an affine scheme.

Remark 2. Any good enough TODO scheme has a Zariski atlas. If \mathbb{T} is finer than the Zariski-topology then in the definition we may replace affine scheme by good enough scheme, if its just about the question whether a type admits an atlas.

Example 3.2. Let X be a (1-)type. X has a Zariski-atlas, iff there exists some $f : \operatorname{Spec} A \to X$ fibered in types of the form $\operatorname{Spec}(R_{f_1} \times \ldots \times R_{f_n})$ for $(f_1, \ldots, f_n) \in Um(R)$.

Remark 3. If one applies ZLC to an affine scheme Spec A the resulting principal open cover $D(f_i), f_i \in A$ will induce indeed a zariski atlas $\bigsqcup D(f_i) \to \operatorname{Spec} A$, because the fiber over $x : \operatorname{Spec} A$ is $\bigsqcup D(f_i(x))$.

Question: Does every zariski atlas of Spec A have this form? Weird Zariski Atlasses

Example 3.3. \mathbb{P}^n has a zariski atlas given by the standart homogeneous principal opens $\sum_{i=0}^{n} D_{+}(x_i)$. The fiber over a point $[y_0:\ldots y_n]$ is $D(y_0)+\ldots D(y_n)$ where $(y_1,\ldots,y_n)\in Um(R)$.

4 Local Choice

In this section let \mathbb{T} denote a topology finer than the zariski topology.

Definition 4.1. Let Cov be a class of morphisms (which we think of n-atlasses of some n), containing \mathbb{T} -atlas, (stable under pullback NECESSARY TODO?) A type S has $local \ choice$ wrt Cov if for any \mathbb{T} -surjective map $X \to Y$ and any map $f: S \to Y$ there exists a map $p': S' \to S$ in Cov and a commutative diagram

$$\begin{array}{ccc}
S' & ---- & X \\
\downarrow & & \downarrow & p \\
S & \xrightarrow{f} & Y
\end{array}$$

Proposition 4.2. Assume that Cov is stable under composition and that Zariski-covers are in Cov. S has \mathbb{T} -local choice wrt Cov if it has a projective Cover, i.e. there exists a projective (or, assuming ZLC, affine scheme resp.) \hat{S} with a map $g: \hat{S} \to S$ in Cov.

Proof. By stability under composition of Cov, We may assume that $g: \hat{S} \to S$ is the identity. As p is \mathbb{T} -surjective, for any x: S there merely is a $\operatorname{Spec} B_x \in T$ and a map $\operatorname{Spec} B_x \to \|\operatorname{fib}_p(x)\|$. Claim: No matter on the assumptions (on $S=\hat{S}$), there exists a Zariski cover $S' \xrightarrow{p'} S$ with S' projective (affine resp.) and a term in

$$\prod_{x:S'} \sum_{\operatorname{Spec} B_x \in T} \operatorname{Spec} B_x \to \|\operatorname{fib}_p(fp'x)\|$$

Proof: In the case of projectivity, just use $p' = \mathrm{id}_S$ and in the case of having ZLC and S beeing affine, use ZLC (3). \square (Claim) By setting

$$(S'' := \sum_{x \in S'} \operatorname{Spec} B_x) \xrightarrow{\pi} S'$$

the projection, we are now in the situation that for any t:S'' we merely have a point in $\mathrm{fib}_p((p''(t)))$ and $S''\to S'$ is a \mathbb{T} -cover, thus it is in Cov. Moreover, S'' is a projective type (affine), as it is a dependent sum of projectives (affines). Hence again we now can find a lift $S''\to X$. making

$$S'' \longrightarrow X$$

$$\downarrow \qquad \qquad \downarrow$$

$$S' \qquad \qquad p$$

$$\downarrow^{p'} \qquad \qquad \downarrow$$

$$S \stackrel{\text{id}}{\longrightarrow} S$$

commute. Now $S'' \to S' \to S$ as the composition of Zariski-covers and Cover is a Cover [...]as desired.

The next lemma shows, that the class of types which have a \mathbb{T} -cover is stable under dependent sums.

Lemma 4.3. Let $\mathcal{U}' \subset \mathcal{U}$ be stable under dependent sums (e.g. \mathbb{T} -inhabited types) Let X be a type with a map $p: U \to X$ fibered in \mathcal{U}' . For any x: X, let Y_x be a type and moreover for any u: U, we are given a map $q_u: V_u \to Y_{p(u)}$ fibered in \mathcal{U}' . Then the induced map

$$p: \sum_{u:U} V_u \to \sum_{x:X} Y_x$$

is fibered in \mathcal{U}'

Proof. The fiber of p over some $(x,y) \in \sum_{x:X} Y_x$ is given by

$$\sum_{u: \operatorname{fib}_p x} \operatorname{fib}_{q_u}(y')$$

where $y': Y_{p(u)}$ (depending on u) is the transport of $y: Y_x$ along x = p(u). As \mathcal{U}' is stable under dependent sum those fibers are again in \mathcal{U}' . This shows the result.

5 AlgebraicSpace

We first need to define a notion of algebraic space and smooth algebraic space, which should be the smallest class of types that satisfies the following:

- Stability under finite limits 6.1
- has Descent
- (nice) Schemes are contained in it
- affines in \mathbb{T} are smooth algebraic spaces. (there are probably more).
- stable under smooth quotients: If X is an algebraic space, Y modal 0-type and $X \to Y$ is \mathbb{T} -surjective and fibered in smooth algebraic spaces, then Y is an algebraic space. Additionally, if X is smooth, then Y is smooth.

Definition 5.1. An affine Scheme U is called flat, if

$$||U||_{\mathbb{T}} \to (U \in \mathbb{T})$$

We call a proposition algebraic, if its merely of the form $||U||_{\mathbb{T}}$ for some flat affine U.

Lemma 5.2. Algebraic propositions are algebraic spaces.

Proof. We have $U \to ||U||_{\mathbb{T}}$ where U is affine, hence an algebraic space and the fibers are in \mathbb{T} by flatness of U, hence they are smooth algebraic spaces. By stability under quotients, our algebraic proposition is an algebraic space.

Definition 5.3. An smooth equivalence relation on a set U is some equivalence relation $R: U^2 \to \text{Prop}$, whose fibers are in \mathbb{T}

Lemma 5.4. let U be an algebraic space (e.g. affine scheme) and $R:U^2\to \operatorname{Prop}$ be a smooth equivalence relation Then U/R is an algebraic space

Proof. The map $U \to U/R$ is fibered in \mathbb{T} , in particular fibered in smooth algebraic spaces. By stability under smooth quotients, U/R is an algebraic space.

Recall the definition of T-atlas??

Lemma 5.5. The class of types admitting a T-atlas is closed under taking identity types.

Proof. This is an example of 11.2.

Question 2. Is it stable under dependent sums?

Proposition 5.6. every proposition P having a \mathbb{T} -atlas is an algebraic proposition.

Proof. Let $V \to P$ be a \mathbb{T} -atlas. have to show TFAE $||V||_{\mathbb{T}} \to P \to (V \in \mathbb{T}) \to ||V||_{\mathbb{T}}$. Proof: $||V||_{\mathbb{T}} \to P$ as P is modal prop. Secondly, because $V \to P$ is a \mathbb{T} -cover. Thirdly, because things in \mathbb{T} are automatically \mathbb{T} -merely inhabited.

Hence P is an algebraic proposition.

Corollary. The identity types of U/R are algebraic propositions.

Definition 5.7. A modal set X is a classical algebraic space iff it is merely of the form U/R for some afine U and $R:U^2\to \operatorname{Prop}$ an algebraic equivalence relation. Equivalently there exists some \mathbb{T} -atlas $U\to X$ (i.e. out of an affine). We call X smooth if U can be choosen to be in \mathbb{T} .

Corollary. Classical Algebraic spaces that are propositions are algebraic propositions.

Lemma 5.8. smooth classical Algebraic spaces which are affine are in \mathbb{T} .

Proof. Saturatedness of the topology.

6 n-stacks

Definition 6.1. Let \mathbb{T} be a subcanonical topology finer than the Zariski topology. Let $n \geq -2$. A type X

- is a (smooth) -2-stack if it is contractible
- is A (n+1)-stack, if
 - -X is a \mathbb{T} -sheaf
 - For any $x, y : X \ x =_X y$ is a *n*-stack
 - There exists an n-atlas, i.e. a T-surjective map $\operatorname{Spec} A \to X$ fibered in
 - * \mathbb{T} , if $n \leq 0$
 - * smooth n-stacks, if n > 0.
- X is a smooth n+1-stack if
 - -X is a (n+1)-stack
 - There exists a *n*-atlas Spec $A \to X$ with Spec $A \in \mathbb{T}$

Lemma 6.2. One could only alternatively talk about (smooth) n-stacks for $n \geq 1$, define them by induction as above. Then later define:

- A (smooth) -1-stack is a (smooth) 1- stack is a proposition.
- A (smooth) 0-stack is a (smooth) 1- that is a 0-type.

 \square Proof.

Lemma 6.3. A (smooth) n-stack is a (smooth) n + 1-stack.

Proof. Induction. We need, that \mathbb{T} is subcanonical to conclude that affines are \mathbb{T} -sheaves. \square

Remark 4. If one changes the definition of atlas to be a map out of a scheme, then smooth -1 atlas will be scheme in \mathbb{T} . Otherwise propositional -1-stack are not 0-stacks.

7 Stability results

Theorem 7.1. Let $n \geq -2$. Smooth / n-stacks are stable by dependent sums.

Proof. Induction. For n=-2 its okay. Let $B:X\to \mathcal{U}$ be a family of n+1-stacks indexed over a n+1-stack X, then surely the total space $\sum_{x:X}Bx$ is a \mathbb{T} -sheaf as \mathbb{T} -sheaves are stable under dependent sum. The identity types in a \sum type are \sum of identity types. It remains to construct some n-atlas $\operatorname{Spec} A\to \sum_{x:X}B_x$ For any x:X we merely have an n-atlas $V_x\to B_x$, i.e. with V_x affine. Claim: X has local choice for X wrt n-atlasses. Proof: n-atlasses contain zariski-atlasses, because \mathbb{T} is finer than the Zariski topology. n-stacks are stable under dependent sum by induction, thus n-atlasses are stable under composition. $\square(\operatorname{Claim})$

By (4.2) for X, we merely find U affine, an n-atlas $p: U \to X$ with

$$\prod_{u:U} \sum_{V_{p(u)} \in T} (q: V_{p(u)} \to B_{p(u)}) \times (q \text{ fibered in smooth } n \text{ stacks })$$

Now the desired map is $\sum_{u:U} V_{pu} \to \sum_{x:X} B_x$, because it is an *n*-atlas by 4.3 If additionally, all the B_x and X are smooth *n*-stacks, just observe that we can choose the affine V_{pu} to lie in \mathbb{T} , Accordingly $\sum_{u:U} V_{pu} \in T$ as \mathbb{T} is stable under Σ .

Corollary. n-atlasses are stable under composition.

Lemma 7.2. n + 1-stacks are closed under taking closed (open) subtypes.

Proof. First we show:if X has an n-atlas and Y is a closed (open) subtype of X, then Y has an n-atlas. Choose an n-atlas Spec $A \to X$. The pullback to Y has have the same fibers. If Y is closed, and the total space is a closed subtype of Spec A, hence it will be affine. if Y is an open subtype of X, then the pullback is an open subtype of Spec A, hence by zariski local choice merely of the form $\bigcup_{i=1}^n D(a_i) \subset A$. As n-atlasses are stable under composition 6, it suffices to show, that the map $f: \bigsqcup_i D(a_i) \to \bigcup_{i=1}^n D(a_i)$ is a Zariski-atlas, because then it will be an n-atlas as well. Let $x: \bigcup_{i=1}^n D(a_i)$, i.e. there merely exists an i, such that $a_i(x)$ is invertible. The fiber is exactly $D(a_1(x)) + \ldots + D(a_n(x))$. thus we are done. (MAYBE OUTSOURCE THIS and say open subschemes of affines have zariski atlas)

Corollary. Let X be a quasi-projective scheme that is a sheaf. Then X is a 0-stack.

Proof. It suffices to see that X has a zariski atlas. Use .

Definition 7.3. A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

Lemma 7.4. Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

Lemma 7.5. Given a local property P of morphisms of modal n-types, a morphism $f: X \to Y$ has P if there exists an n-atlas of f having P.

The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n = 0, 1.

8 Descent

Theorem 8.1. Let T be a modal n-type. The Proposition, that P is a (smooth) n-stack, is modal.

9 Fundamental Theorem of algebraic spaces

9.1 For groupoids

Lemma 9.1. If R woheadrightarrow X o X is a \mathbb{T} -httpy-coequalizer diagram of two \mathbb{T} -covers between affines, then X is a 1-stack.

9.2 For sets

Lemma 9.2. Denote $\mathbb{T}Set$ for the sets that are \mathbb{T} -sheaves. Assume given a $\mathbb{T}set\ X$ then the following maps are mutually inverse

$$\begin{split} \sum_{R:X\to X\to\mathbb{T}\,\text{Prop}} R \text{ equivalence relation} &\simeq \sum_{Y:\mathbb{T}\mathsf{Set}} \sum_{p:X\to Y} p \text{ } \mathbb{T} \text{surjective} \\ R &\mapsto (X/R,[_]) \\ \lambda x, y.(p(x) = p(y)) & \hookleftarrow (Y,p) \end{split}$$

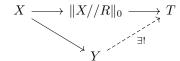
where X/R is defined by applying $L_T\|_{-}\|_0$ at the higher inductive type X//R.

- *Proof.* Well-definedness: The map $[_]: X \to ||X//R||_0 \to L_T ||X//R||_0$ is the composition of a surjective with a \mathbb{T} -surjective map [ref?], hence its \mathbb{T} -surjective. Conversely given (Y,p) as Y is a sheaf, we have for all x,y:X that $p(x)=_Y p(y)$ is a sheaf.
 - If x, y : X then we have a chain of equivalences

$$R(x,y) \simeq (\bar{x} = \|X//R\|_0 \bar{y}) \to ([x] = L_T \|X//R\|_0 [y])$$

where the first map is plain HoTT and the second map is ap, i.e. the unit of the modality [ref?], but as the $\bar{x} = \|X//R\|_0$ \bar{y} is already a sheaf, it is an isomorphism as well.

• Let (Y,p) be in the RHS. Let $R(x,y)=(p(x)=p(y)):\mathbb{T}$ Prop. By plain HoTT, There is a map $\eta:\|X//R\|_0\to Y$ (defined by the universal property of the set truncation and by induction on the higher inductive type X//R on canonical terms through the map $p:X\to Y$). I claim η exhibits Y as the sheafification of $\|X//R\|_0$. Let T be another Tsheaf equipped with a map $\|X//R\|_0\to T$. By precomposition we obtain a map $X\to T$. Claim: it factors uniquely through $p:X\to Y$.



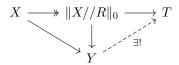
Proof:

Existence: We want to define a map $Y \to T$. Let y:Y. As p is \mathbb{T} -surjective and T is a sheaf, we may assume we merely have some element in the fiber of p over y. Now push this element through

$$\|\operatorname{fib}_p y\| \to \|X//R\|_0 \to T$$

where the first map is by Plain HoTT and the second one was given by assumption. One can easily check this makes the diagram commute. Uniqueness follows as $X \to Y$ is \mathbb{T} -surjective and \mathbb{T} is a sheaf.

By surjectivity of $X \to ||X//R||_0$ there exists a unique factorization



making the right triangle commute. This is what we wanted to show.

Definition 9.3. An equivalence relation R on a type X is called:

- if for all x, y : X the proposition R(x, y) is a -1-stack.
- smooth if its and for any y: X its fibers:

$$R_y := \sum_{x:X} R(x,y)$$

are affine in \mathbb{T} .

Lemma 9.4. Assume that \mathbb{T} satisfies descent for propositions and for sets 7.1, i.e. that a modal proposition being a (-1)-stack is a sheaf. Assume that a modal set beeing affine in \mathbb{T} is a sheaf. Assume given a \mathbb{T} set X, then the following types are equivalent:

- The type of smooth equivalence relations over X.
- The type of Tsets Y with identity types beeing stacks and an -1-atlas X to Y (in V2 a T-cover).

Proof. By the equivalence in 8.2, it is enough to check that:

• The identity types in X/R are (-1)-stacks if and only if the relation R is . For any x,y:X we know that:

$$R(x,y) \simeq [x] =_{X/R} [y]$$

so the direct direction is immediate. For the converse we use the assumption that a modal proposition being a (-1)-stack is a sheaf and that the map [_]: $X \to X/R$ is \mathbb{T} -surjective.

• The fibers of:

$$[_]: X \to X/R$$

are affine in \mathbb{T} if and only if the relation R is smooth. For any y:X we have that:

$$\sum_{x:X} R(x,y) \simeq \mathrm{fib}_{[.]}([y])$$

so the direct direction is immediate. Here as well the converse follows from \mathbb{T} -surjectivity of $[\]$ and that the topology has descent.

Corollary. Assume \mathbb{T} satisfies descent for propositions and for sets. A type is a 0-stack iff its merely the \mathbb{T} -quotient of an affine scheme by a smooth equivalence relation.

Theorem 9.5. Assume \mathbb{T} satisfies descent for propositions. The quotient of a 0-stack $X \in \mathbb{T}$ Set by an 0-smooth equivalence relation R is a 0-stack. TODO

Proof. The identity types in X/R are propositional 0-stacks, hence (-1)-Truncations of -1-stacks by 10.2 as desired.

How to find an atlas: todo. How to proceed, if we could choose all atlasses we want at the same time?

Remark 5. This is equivalent to saying that 1-stacks that are 0-types are geomeric 0-stacks: One direction we prove later. If R is a 0-smooth equivalence relation on a 0-stack X, then X/R is a 1-stack by observing that any -1-atlas $X' \to X$ gives a 0-atlas $X' \to X \to X/R$. Moreover, X/R is a 0-type, hence by assumption a 0-stack.

Example 9.6. There are open affine subschemes U of affine schemes $\operatorname{Spec} A$, which are not (disjoints unions of) principal open

Proof. Consider $A=R[x,y,u,v]/(xy+ux^2+vy^2), X=\operatorname{Spec} A$ and consider the open U=D(x,y).

We cant expect U to be a disjoint union of principal opens (todo). However, D(x,y) is affine: We have maps $U \to R$ given by $f = -v/x = (y + ux)/y^2$, $g = -u/y = (x + vy)/x^2$. Then $D(f) \cup D(g) = \operatorname{Spec} R^X$, as yf + xg = 1 in R^U . Taking preimages under the affinization map, $U_f \cup U_g = X$ and one checks this defines an open affine cover (for example : $U_f \simeq \operatorname{Spec} R[x, u, f^{\pm 1}, g]/(xy + ux^2 + uy^2)$ with y := (1 - gx)/f.) But on both of this open subsets the affinization map is an isomorphism hence the affinization of X is an isomorphism. compare (Hartshorne II.2.17)

Lemma 9.7. Let $f: X \to Y$ be surjective. There exists a Zariski Cover $X' \to X$ such that $X' \to Y$ is a Zariski cover iff there exists a Zariski Cover $X' \to X$, some $n: \mathbb{N}$ and an open affine embedding $X' \hookrightarrow Y^n$ over Y.

10 Saturated Topologies revisited

Lemma 10.1 (1.1). We want that every n-1-atlas of a smooth n-atlas has the additional requirement in the definition of smooth n-atlas. It turns out, that for this topology needs to be saturated: The following are equivalent

- 1. Beeing in the topology descents along \mathbb{T} -covers between affines, i.e. \mathbb{T} is saturated.
- 2. A smooth n -stack X that is an affine scheme lies in the Topology \mathbb{T} .
- 3. Let $n \geq 0$. If T is a smooth n-stack, then any n-1-atlas $U \to T$ satisfies $U \in \mathbb{T}$.
- 4. If $U \xrightarrow{f} V \xrightarrow{g} W$ are maps between affines and f and gf are \mathbb{T} covers, then g is a \mathbb{T} Cover

Proof. $1 \Rightarrow 2$

Induction. This holds for n=-1. Assume it holds for n-1. Choose a n-1-atlas with T source, i.e. $T\ni \operatorname{Spec} A\to X$ fibered in smooth n-1-stacks. As it is affine, all the fibers of the atlas are affine smooth n-1-stacks, hence by induction they lie in \mathbb{T} , thus the atlas is a \mathbb{T} -cover between affines, hence $X\in\mathbb{T}$.

 $2 \Rightarrow 3$

As $U \to T$ is fibered in smooth n-1 stacks, all the fibers are in particular smooth n-stacks by 5.11. By stability under dependent sum $U = \sum_{t:T} U_t$ is a smooth n-stack that is affine, hence by assumption (2) it lies in the topology.

 $3 \Rightarrow 1$

Let $X \to Y$ be a \mathbb{T} -cover with X affine in \mathbb{T} and Y affine. Then Y is a smooth 0-stack, But $Y \to Y$ is a -1-atlas, hence by assumption $Y \in T$.

 $4 \Rightarrow 1$

Obvious

 $1 \Rightarrow 4$

Check fiberwise

If $n \geq$, replacing \mathbb{T} by its saturation \mathbb{T}' does change the notion of (smooth) n-stack, but we have the following statement, that tells us, that if we start with 0- \mathbb{T} -stacks then the notion of smoothness does not see the difference between \mathbb{T} and its saturation.

Proposition 10.2. Let X be a 0-stack that is a weak smooth 0-stack, i.e. there exists a \mathbb{T}' -atlas $\mathbb{T}' \ni X' \to X$ (i.e. fibered in \mathbb{T}'). Then X is a smooth 0-stack.

Proof. Wlog $X' \in \mathbb{T}$. Choose a -1-atlas Spec $A \to X$ (i.e. fibered in \mathbb{T}). As the fibers of $X' \to X$ merely have smooth atlasses $\tilde{X}'_x \to X'_x$, we can use Local choice to obtain a commutative diagram $Y = \sum_{x':X'} \tilde{X}'_x$

$$\tilde{X} \xrightarrow{\mathbb{T}} \operatorname{Spec} A$$

$$\mathbb{T} \downarrow \qquad \qquad \downarrow$$

$$X' \xrightarrow{\mathbb{T}'} X$$

As $Y \to X'$ is a \mathbb{T} -cover and $X' \in \mathbb{T}$ we conclude $Y \in \mathbb{T}$. Hence we found a smooth \mathbb{T} -atlas of X.

10.1 Zariski Topology is not saturated

Example 10.3 (Weird Zariski Atlasses). Assume those equivalent conditions on the Zariski topology. There exist Zariski atlasses of affines $\operatorname{Spec} A = X$ which are not of the form $D(a_1) + \ldots + D(a_n) \to \operatorname{Spec} A$ for $(a_1, \ldots, a_n) \in Um(A)$

Proof. Indeed, using the first example, choose $U \subset \operatorname{Spec} A$ affine not principal open, then choosing a Zariski atlas $V \to U$ gives $V + X \to U + X \to X$ where $V + X \to X$ is a Zariski cover and $V + X \to U + X$ is a Zariski cover. From (4), we deduce that $U + X \to X$ is a Zariski cover, but U is not a disjoint union of principal opens in $\operatorname{Spec} A$.

Example 10.4. Assume those equivalent conditions on the Zariski topology. Every affine open proposition U is principal open!

Proof. Let $V \to U$ be a Zariski atlas. Then $V+1 \to U+1$ is a Zariski atlas with $V+1 \in \mathbb{T}$ and U+1 affine, hence by (1) $U+1 \in \mathbb{T}$, hence U is a disjoint union of principal opens hence, as it is a proposition, its principal open.

11 beeing a stack is indepent of the truncation level

Lemma 11.1. Let $n \ge 0$. A n-stack is an modal n-type.

Proof. The n-Truncationis an n-type. Now conclude by induction.

We want to show that the notion of stack makes sense, i.e. beeing should not depend on the truncation level.

Lemma 11.2. Assume \mathbb{T} is saturated and satisfies descent for propositions. Let P be a modal proposition. Then TFAE

- 1. For some $m \geq 0$, P is a m stack
- 2. There exists some fp algebra A such that Spec $A \to P$ and P is logically equivalent to (Spec $A \in \mathbb{T}$).
- 3. P is equivalent to $\|\operatorname{Spec} A\|_{\mathbb{T}}$ for some $fp\ A$, i.e. P is a -1-stack.

Proof.

- 1. \Rightarrow 2. Let Spec $A \to P$ be a m-1 atlas. Assume Spec $A \in \mathbb{T}$. Then $\|\operatorname{Spec} A\| \to P$ so as P is a sheaf, we have P. Conversely, if x : P, then the fiber over x is Spec A and a smooth m-1 stack, hence belongs to the topology by 9.1.
- 2. \Rightarrow 3. We have to show: There exists some fp algebra such that P is logically equivalent to $\|\operatorname{Spec} A\|_{\mathbb{T}}$. By assumption we have $\operatorname{Spec} A \to P$, so we deduce $\|\operatorname{Spec} A\|_{\mathbb{T}} \to P$ as P is a modal proposition. Conversely $P \to (\operatorname{Spec} A \in \mathbb{T}) \to \|\operatorname{Spec} A\|_{\mathbb{T}}$, where the first arrow is by assumption.
- $3. \Rightarrow 1. \ 5.11$

Lemma 11.3. A smooth -1-stack P is contractible.

Proof. Choose a \mathbb{T} -cover $\mathbb{T} \ni \operatorname{Spec} A \to P$. As P is a proposition we have $\|\operatorname{Spec} A\| \to P$. As P is a sheaf we have P.

Example 11.4. A 0-stack is a \mathbb{T} -sheaf whose identity types are (-1)- \mathbb{T} runcations of ((affine ?)) schemes and there exists a \mathbb{T} -atlas Spec $A \to X$.

Why are schemes 0-stacks? This holds in special case, for example if the scheme is quasi projective.

Theorem 11.5. Let \mathbb{T} be saturated. Assume the topology satisfies descent Let $m, n \geq -2$. Given an n-type T that is a (smooth) m-stack then T is a (smooth) n-stack.

Proof. By 5.11 we may assume $m \ge n \ge -2$.

If $m \le 1$ this is clear. Now assume $m \ge 2$. Induction. Inductionstart m = 2. Let us prove the case of m = 2, n = 1, the cases $-2 \le n < 1$ are immediate from this.

Choose a 1-atlas $X' \to T$, i.e. its fibered in smooth 1-stacks. As T is a groupoid and X' is a set, the fibers are actually sets, i.e. smooth 0-stacks.

Now consider $R:=X'\times_T X'$. As X' is in particular a 0-stack and 0-stacks are stable under dependent sums, R will be a 0-stack. Choose a a \mathbb{T} -cover $R'\to R$ with R' affine. Now $R'\to R\to X'$ is a map between affine schemes i.e it is fibered in smooth 0-stacks that are affine. As \mathbb{T} is saturated, the fibers of $R'\to X'$ are in \mathbb{T} . As X'//R' is a 1-stack by ??, it suffices to show that $X'//R'\to X'//R$ is a \mathbb{T} -cover. Pick a term in X'//R. As the fiber beeing in \mathbb{T} is sheaf If additionally T is assumed to be a smooth 2-stack, then we can assume X' to be in the topology. This will force R to be a smooth 0-stack, so we may choose R' Assume m>2 and the statement is proven for all (n',m')<(n,m) in lexicographical ordering. As the identity types of T are n-1-types and m-1 stacks by induction they are n-1 stacks. Let $X\to T$ be an m-1-atlas, i.e. fibered in smooth m-1-stacks with X affine. The fibers are in particular n-1-types, so by induction they are smooth n-1-stacks. Hence $X\to T$ is an n-1-atlas. If, additionally T is assumed to be a smooth m-stack, we can choose $X\in \mathbb{T}$, hence $X\to T$ witnesses that T is a smooth n-stack.

12 Stability under Quotients

Definition 12.1. A morphism between *n*-stacks is smooth if it is fibered in

- \mathbb{T} if $n \leq 0$
- smooth n-stacks if n > 0.

Lemma 12.2. Let C be a class of affine schemes. The class of types Y which admit a map $\operatorname{Spec} A \to Y$ fibered in C is closed under taking identity types.

Proof. By assumption we can choose a map $p:V\to Y$ out of an affine fibered in C. Let y,y':Y. Then we have the map

$$(\operatorname{fib}_p y) \times_V (\operatorname{fib}_p y') \to y = y'$$

$$(v, q: y = pv), (v', q': y' = pv'), (h: v = v') \mapsto q \cdot h \cdot q'^{-1}$$

The domain is a fiber product of affines, hence affine. The fiber over j: y = y' looks like

$$\sum_{v} (\underbrace{\sum_{v'} (h:v=v')}) \times (q:y=pv) \times (q':y'=pv') \times (q\cdot h\cdot q'^{-1}=j) \simeq \sum_{v} (v=py) \simeq \operatorname{fib}_{p} y$$

Lemma 12.3 (TODO). If a modal proposition P admits a -2-atlas Spec $A \to P$ then P is a -1-stack.

Theorem 12.4. Let $f: X \to Y$ be a \mathbb{T} -surjective smooth morphism between modal n-types. If X is a (smooth) stack, then Y a (smooth) stack.

Proof. Induction. For n=-2 its clear. Let X be a n-stack. Lets first construct the n-1-atlas of Y. We merely find a $V \twoheadrightarrow X$ which is an n-1-atlas. Then $V \to X \to Y$ is an n-atlas because it is \mathbb{T} -surjective and is fibered in the correct \sum -stable class of types, i.e. \mathbb{T}

if $n \le 1$ and smooth n-1-stacks for n > 1. Hence Y is an n+1-stack. As Y is an n-type, Y is an n-stack 10.5.

If additionally X is assumed to be smooth, then V can be assumed to lie in \mathbb{T} which directly gives us that Y has a smooth atlas.

It remains to show that the identity types of Y are n-1-stacks. By 11.2 we find some n-1-atlas $p:W\to y=y'$. The map is smooth, because the fibers of p are smooth. If $n=0,\ y=y'$ is a -1-stack by 11.3. If n>0, W is an n-1-stack and p is smooth, so by induction y=y' is an n-1-stack.

13 Local properties

Definition 13.1. Let Cov be the property of morphisms of n-stacks defined by asking that the morphism is \mathbb{T} -surjective and fibered in smooth n-stacks. Its stable under basechange. A property of n-stacks is local if P(1) holds, P is stable by dependent sums and given a $Cover\ X \to Y$ we have PX iff PY.

Example 13.2. beeing smooth n-stack is a local property of stacks.

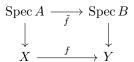
Proof. We have to show: If $f: X \to Y$ is a T-surjective map fibered in smooth n-stacks between n-stacks, then X is a smooth n-stack iff Y is a smooth n-stack. The only if is clear by stability under dependent sums. The other direction is 11.4.

Definition 13.3. A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

Lemma 13.4. Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

Lemma 13.5 ([ref?]). Given a local property P of morphisms of n-stacks, a morphism $f: X \to Y$ has P if there exists an n-atlas of f having P.

Example 13.6. A morphism of n-stacks is smooth iff there exists an n-atlas of f



such that \tilde{f} is a \mathbb{T} -cover.

r The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n = 0, 1.