# Thesis

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### 1 Saturated Topologies

Consider a topology  $\mathbb{T}$  finer than the Zariski topology.

**Definition 1.1.** A smooth atlas of X is some  $\hat{X} \in \mathbb{T}, \hat{X} \to X$  T-cover

**Definition 1.2.**  $\mathbb{T}$  is saturated if Beeing in the topology descents along  $\mathbb{T}$ -covers between affines, i.e. every affine schemes that has a smooth atlas lies itself in  $\mathbb{T}$ .

The saturated closure of a topology  $\mathbb{T}$  is the topology  $\mathbb{T}'$  defined by (todo finite sums of?)

$$X \in \mathbb{T}'$$
 iff X is affine  $\wedge \exists$  smooth at as of X

**Lemma 1.3.** Using ZLC, this is the smallest saturated topology containing  $\mathbb{T}$ .

*Proof.* Obviously  $1 \in \mathbb{T}'$ . Types which have a smooth atlas are stable by dependent sums by the proof of  $\ref{thm:proof}$ . For the saturatedness consider some  $\mathbb{T}'$ -cover  $\mathbb{T}' \ni X' \to X$ . By replacing X' with some smooth atlas, we may assume that  $X' \in \mathbb{T}$ . As every fiber  $X'_x \in \mathbb{T}'$ , we merely find a smooth atlas  $\tilde{X}'_x \to X'_x$ . Then by Zariski local choice there exists a Zariski atlas  $\hat{X} \to X$  and a commutative diagram

$$Y \equiv \sum_{x:\hat{X}} \tilde{X}'_x \longrightarrow \sum_{x:X} X'_x = X'$$

$$\downarrow \qquad \qquad \downarrow$$

$$\hat{X} \xrightarrow{Zax} X$$

As  $X' \in \mathbb{T}$  and  $Y \to X'$  is fibered in  $\mathbb{T}$  (4.3) we have  $Y \in \mathbb{T}$ . But  $Y \to \hat{X}$  is a  $\mathbb{T}$ -cover and  $\hat{X} \to X$  is a  $\mathbb{T}$ -cover,  $Y \to X$  is a  $\mathbb{T}$ -cover. Hence  $X \in \mathbb{T}'$ .

**Lemma 1.4.** A type T is a sheaf wrt to  $\mathbb{T}'$  iff it is a sheaf wrt to  $\mathbb{T}$ 

*Proof.* As  $\mathbb{T} \subset \mathbb{T}'$  the  $\to$  direction is clear. Now, let  $X \in \mathbb{T}'$ . We have to show that  $T \to T^{\|X\|}$  is an equivalence. Choose  $\mathbb{T} \ni Y \to X$ . Then we have a commutative diagram

$$T \xrightarrow{\simeq} T^{\|X\|} \downarrow \\ T^{\|Y\|}$$

So  $T \to T^{\|X\|}$  has a left-inverse. Thus it suffices to show that any  $f: T^{\|X\|}$  has a preimage. Choose t: T, s.th.  $\operatorname{cnst}_t^Y$  is the composite  $\|Y\| \to \|X\| \stackrel{f}{\to} T$ . We have  $\|Y\| \to (\operatorname{cnst}_t^X = f)$ . But as  $Y \in \mathbb{T}$  and  $\Delta_t = f$  is a sheaf (as an identity type in the sheaf  $T^{\|X\|}$ ) we are done.  $\square$ 

Remark 1. We never used that we only talk about T-covers.

**Lemma 1.5.** Every saturated affine (i.e. Spec  $A \in \mathbb{T}'$ ) it  $\mathbb{T}$ -merely inhabited.

*Proof.* We have  $||X|| \to ||\operatorname{Spec} A||$  for some smooth atlas  $\mathbb{T} \ni X \to \operatorname{Spec} A$ .

Question 1. Does the converse hold, i.e. is every T-merely inhabited affine saturated?

# 2 Lex Modalities

Lemma 2.1 (Stability resuls). Lex Modalities are stable under

- 1. Conjunction
- 2. Composition

**Lemma 2.2.** Let  $\bigcirc$  be a lex-modality. Let X be  $\bigcirc$ -modal and  $B: X \to \mathcal{U}_{\bigcirc}$  be a family of modal types. Then  $\sum_{x:X} B_x$  is  $\bigcirc$ -modal

**Lemma 2.3.** Let  $B: \bullet X \to \mathcal{U}$ . Then  $\bullet(\sum_{x:X} B(\eta x)) = \sum_{x:\bullet X} \bullet Bx$ 

*Proof.* Observe that

$$\sum_{x:X} B\eta x \to \sum_{x:\bullet X} Bx$$

is a  $\bullet$ -equivalence, because for all modal types T, the type  $Bx \to T$  is modal for any  $x : \bullet X$ . Then it follows by [ref?].

### 3 Atlas

**Definition 3.1.** A  $\mathbb{T}$ -atlas of X is a  $\mathbb{T}$ -cover Spec  $A \to X$  out of an affine scheme.

**Remark 2.** Any good enough TODO scheme has a Zariski atlas. If  $\mathbb{T}$  is finer than the Zariski-topology then in the definition we may replace affine scheme by good enough scheme, if its just about the question whether a type admits an atlas.

**Example 3.2.** Let X be a (1-)type. X has a Zariski-atlas, iff there exists some  $f : \operatorname{Spec} A \to X$  fibered in types of the form  $\operatorname{Spec}(R_{f_1} \times \ldots \times R_{f_n})$  for  $(f_1, \ldots, f_n) \in Um(R)$ .

**Remark 3.** If one applies ZLC to an affine scheme Spec A the resulting principal open cover  $D(f_i), f_i \in A$  will induce indeed a zariski atlas  $\bigsqcup D(f_i) \to \operatorname{Spec} A$ , because the fiber over  $x : \operatorname{Spec} A$  is  $\bigsqcup D(f_i(x))$ .

Question: Does every zariski atlas of Spec A have this form? Weird Zariski Atlasses

**Example 3.3.**  $\mathbb{P}^n$  has a zariski atlas given by the standart homogeneous principal opens  $\sum_{i=0}^{n} D_+(x_i)$ . The fiber over a point  $[y_0:\ldots y_n]$  is  $D(y_0)+\ldots D(y_n)$  where  $(y_1,\ldots,y_n)\in Um(R)$ .

#### 4 Local Choice

In this section let  $\mathbb{T}$  denote a topology finer than the zariski topology.

**Definition 4.1.** Let Cov be a class of morphisms (which we think of n-atlasses of some n), containing  $\mathbb{T}$ -atlas, (stable under pullback NECESSARY TODO?) A type S has  $local \ choice$  wrt Cov if for any  $\mathbb{T}$ -surjective map  $X \to Y$  and any map  $f: S \to Y$  there exists a map  $p': S' \to S$  in Cov and a commutative diagram

$$S' \xrightarrow{p} X$$

$$\downarrow \qquad \qquad p \downarrow$$

$$S \xrightarrow{f} Y$$

**Proposition 4.2.** Assume that Cov is stable under composition.

- If  $\hat{S} \to S$  is a Cover and  $\hat{S}$  has  $\mathbb{T}$ -local choice, then S has  $\mathbb{T}$ -local choice.
- Affine schemes have  $\mathbb{T}$ -local choice.
- Any type admitting a Cov Atlas Spec  $A \to S$  has  $\mathbb{T}$ -local choice.

*Proof.* The first point follows from stability under composition of Cov. the third point follows from the second. By the first point, we may assume that S is affine. As p is  $\mathbb{T}$ -surjective, for any x:S there merely is a  $\operatorname{Spec} B_x \in T$  and a map  $\operatorname{Spec} B_x \to \|\operatorname{fib}_p(x)\|$ . As S is projective, we have a term in

$$\prod_{x:S} \sum_{\operatorname{Spec} B_x \in T} \operatorname{Spec} B_x \to \|\operatorname{fib}_p(fx)\|$$

By setting

$$(S' := \sum_{x:S} \operatorname{Spec} B_x) \xrightarrow{\pi} S$$

the projection, we are now in the situation that for any t:S' we merely have a point in  $\mathrm{fib}_p((p'(t)))$  and  $S'\to S$  is a  $\mathbb{T}$ -cover, thus it is in Cov. Moreover, S' is affine, as it is a dependent sum of affines. Hence again we now can find a lift  $S'\to X$  making

$$S' \longrightarrow Y$$

$$\downarrow p' \downarrow p \downarrow$$

$$S \longrightarrow X$$

commute.

The next lemma shows, that the class of types equipped with a T-atlas is stable under dependent sums.

**Lemma 4.3.** Let  $\mathcal{U}' \subset \mathcal{U}$  be stable under dependent sums (e.g.  $\mathbb{T}$ -inhabited types) Let X be a type with a map  $p: U \to X$  fibered in  $\mathcal{U}'$ . For any x: X, let  $Y_x$  be a type and moreover for any u: U, we are given a map  $q_u: V_u \to Y_{p(u)}$  fibered in  $\mathcal{U}'$ . Then the induced map

$$p: \sum_{u:U} V_u \to \sum_{x:X} Y_x$$

is fibered in  $\mathcal{U}'$ 

*Proof.* The fiber of p over some  $(x,y) \in \sum_{x:X} Y_x$  is given by

$$\sum_{u: \mathrm{fib}_p} \mathrm{fib}_{q_u}(y')$$

where  $y': Y_{p(u)}$  (depending on u) is the transport of  $y: Y_x$  along x = p(u). As  $\mathcal{U}'$  is stable under dependent sum those fibers are again in  $\mathcal{U}'$ . This shows the result.

**Theorem 4.4.** Let  $\mathcal{U}'$  be a class stable under dependent sums. The class of types admitting a  $\mathcal{U}'$ -atlas is closed under dependent sums. If  $\mathbb{T}$  is a topology, the same holds for smooth  $\mathcal{U}'$ -atlasses (i.e. with domain in  $\mathbb{T}$ ).

*Proof.* Let us construct some atlas Spec  $A \to \sum_{x:X} B_x$  For any x:X we merely have an atlas  $V_x \to B_x$ , i.e. with  $V_x$  affine. X has local choice wrt atlasses by (4.2) using  $\mathcal{U}'$  is  $\sum$ -stable (we use the trivial topology).

If additionally, all the  $B_x$  and X are smooth n-stacks, just observe that we can choose the affine  $V_{pu}$  to lie in  $\mathbb{T}$ , Accordingly  $\sum_{u:U} V_{pu} \in T$  as  $\mathbb{T}$  is stable under  $\Sigma$ . By Local choice for X, we merely find U affine, an atlas  $p:U\to X$  with

$$\prod_{u:U} \sum_{V_{p(u)} \in T} (q: V_{p(u)} \to B_{p(u)}) \times (q \text{ fibered in smooth } n \text{ stacks })$$

Now the desired map is  $\sum_{u:U} V_{pu} \to \sum_{x:X} B_x$ , because it is an atlas by 4.3

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### 5 Algebraic Space

We first need to define a notion of algebraic space and smooth algebraic space, which should be the smallest class of types that satisfies the following:

- Stability under finite limits 7.1
- has Descent
- (nice) Schemes are contained in it
- $\bullet$  affines in  $\mathbb{T}$  are smooth algebraic spaces. (there are probably more).
- stable under smooth quotients: If X is an algebraic space, Y modal 0-type and  $X \to Y$  is  $\mathbb{T}$ -surjective and fibered in smooth algebraic spaces, then Y is an algebraic space. Additionally, if X is smooth, then Y is smooth.

**Definition 5.1.** An affine Scheme U is called flat, if

$$||U||_{\mathbb{T}} \to (U \in \mathbb{T})$$

Lemma 5.2. The converse holds always

*Proof.* because things in  $\mathbb{T}$  are automatically  $\mathbb{T}$ -merely inhabited

Recall the definition of T-atlas 3.1

**Definition 5.3.** We call a modal proposition algebraic, if one of the equivalent conditions is satisfied:

- 1. its merely of the form  $||U||_{\mathbb{T}}$  for some flat affine U.
- 2. There is a  $\mathbb{T}$ -surjective map out of a flat affine U.
- 3. It has a  $\mathbb{T}$ -atlas.

Proof.

 $1 \Leftrightarrow 2$  Clear.

- $1 \Rightarrow 3$  we show that  $U \to ||U||_{\mathbb{T}}$  is a T-atlas. Every fiber is in T, because U is flat.
- $3 \Rightarrow 1$  Let  $V \to P$  be a  $\mathbb{T}$ -atlas. have to show TFAE  $||V||_{\mathbb{T}} \to P \to (V \in \mathbb{T}) \stackrel{5.2}{\to} ||V||_{\mathbb{T}}$ . Proof:  $||V||_{\mathbb{T}} \to P$  as P is modal prop. Secondly, because  $V \to P$  is a  $\mathbb{T}$ -cover. Hence P is an algebraic proposition.

**Lemma 5.4.** Algebraic propositions are algebraic spaces.

*Proof.* We have  $U \to ||U||_{\mathbb{T}}$  where U is affine, hence an algebraic space and the fibers are in  $\mathbb{T}$  by flatness of U, hence they are smooth algebraic spaces. By stability under quotients, our algebraic proposition is an algebraic space.

**Definition 5.5.** An smooth equivalence relation on a set U is some equivalence relation  $R: U^2 \to \text{Prop}$ , whose fibers are in  $\mathbb{T}$ 

**Lemma 5.6.** let U be an algebraic space (e.g. affine scheme) and  $R: U^2 \to \text{Prop be a}$  smooth equivalence relation Then U/R is an algebraic space

*Proof.* The map  $U \to U/R$  is fibered in  $\mathbb{T}$ , in particular fibered in smooth algebraic spaces. By stability under smooth quotients, U/R is an algebraic space.

**Corollary.** Let U be affine and R a smooth equivalence relation. The identity types of U/R, i.e. the propositions R(x,y), are algebraic propositions.

*Proof.* By 12.2, the class of types admitting a  $\mathbb{T}$ -atlas is closed under taking identity types. U/R is a type admitting a  $\mathbb{T}$ -atlas, hence its identity types admit them as well.

**Definition 5.7.** A modal set X is a classical algebraic space iff it is merely of the form U/R for some affine U and  $R:U^2\to \operatorname{Prop}$  a smooth equivalence relation. Equivalently there exists some  $\mathbb{T}$ -atlas  $U\to X$  (i.e. out of an affine). We call X smooth if U can be choosen to be in  $\mathbb{T}$ .

**Corollary** (of 5.3). Classical Algebraic spaces that are propositions are algebraic propositions.

**Remark 4.** Assume Saturatedness of the topology. smooth classical Algebraic spaces which are affine are in  $\mathbb{T}$ .

**Question 2.** Is the class of classical algebraic spaces stable under smooth quotients? If its not, how should we enlarge it?

Try: Assume R is fibered in smooth algebraic spaces. Choose  $U \to T$  a  $\mathbb{T}$ -atlas. For any x:U the fiber  $R_x$  merely has an atlas  $\tilde{R}_x \to R_x$ . As U has choice (its affine), we find some  $\mathbb{T}$ -cover  $\tilde{U} = \sum_x \tilde{R}_x \to \sum_x R_x$ . Goal: Find for all t:U/R a  $\mathbb{T}$ -atlas  $V_t \to \mathrm{fib}_{\parallel}(t)$ . Then  $\sum_t V_t$  will be affine, because its the total space of a  $\mathbb{T}$ -cover of an affine. Moreover,  $\sum_t V_t \to \sum_t \mathrm{fib}_{\parallel}(t) \to U/R$  will be a  $\mathbb{T}$ -cover, as  $V_t \in \mathbb{T}$ . This is what we wanted to show.

#### 6 n-stacks

**Definition 6.1.** Let  $\mathbb{T}$  be a subcanonical topology finer than the Zariski topology. Let  $n \geq -2$ . A type X

- is a (smooth) -2-stack if it is contractible
- is A (n+1)-stack, if
  - -X is a  $\mathbb{T}$ -sheaf
  - For any  $x, y : X \ x =_X y$  is a *n*-stack
  - There exists an n-atlas, i.e. a T-surjective map  $\operatorname{Spec} A \to X$  fibered in
    - \*  $\mathbb{T}$ , if n < 0
    - \* smooth n-stacks, if n > 0.
- X is a smooth n+1-stack if
  - -X is a (n+1)-stack
  - There exists a *n*-atlas Spec  $A \to X$  with Spec  $A \in \mathbb{T}$

**Lemma 6.2.** One could only alternatively talk about (smooth) n-stacks for  $n \ge 1$ , define them by induction as above. Then later define:

- A (smooth) -1-stack is a (smooth) 1- stack is a proposition.
- A (smooth) 0-stack is a (smooth) 1- that is a 0-type.

Proof.

**Lemma 6.3.** A (smooth) n-stack is a (smooth) n + 1-stack.

*Proof.* Induction. Be aware of the induction start, where maybe no atlas is assumed! We need, that  $\mathbb{T}$  is subcanonical to conclude that affines are  $\mathbb{T}$ -sheaves.

**Remark 5.** If one changes the definition of atlas to be a map out of a scheme, then smooth -1 atlas will be scheme in T. Otherwise propositional -1-stack are not 0-stacks.

# 7 Stability results

**Theorem 7.1.** Let  $n \geq -2$ . Smooth / n-stacks are stable by dependent sums.

*Proof.* Induction. For n=-2 its okay. Let  $B:X\to \mathcal{U}$  be a family of n+1-stacks indexed over a n+1-stack X, then surely the total space  $\sum_{x:X}Bx$  is a  $\mathbb{T}$ -sheaf as  $\mathbb{T}$ -sheaves are stable under dependent sum. The identity types in a  $\sum$  type are  $\sum$  of identity types. Admitting an n-atlas is stable under dependent sum: We apply 4.4 to the class of (smooth) n-atlasses, which is stable under depent sum by induction.

Corollary. n-atlasses are stable under composition.

**Lemma 7.2.** n+1-stacks are closed under taking closed (open) subtypes.

*Proof.* First we show:if X has an n-atlas and Y is a closed (open) subtype of X, then Y has an n-atlas. Choose an n-atlas Spec  $A \to X$ . The pullback to Y has have the same fibers. If Y is closed, and the total space is a closed subtype of Spec A, hence it will be affine. if Y is an open subtype of X, then the pullback is an open subtype of Spec A, hence by zariski local choice merely of the form  $\bigcup_{i=1}^n D(a_i) \subset A$ . As n-atlasses are stable under composition T, it suffices to show, that the map  $f: \bigsqcup_i D(a_i) \to \bigcup_{i=1}^n D(a_i)$  is a Zariski-atlas, because then it will be an n-atlas as well. Let  $x: \bigcup_{i=1}^n D(a_i)$ , i.e. there merely exists an i, such that  $a_i(x)$  is invertible. The fiber is exactly  $D(a_1(x)) + \ldots + D(a_n(x))$ . thus we are done. (MAYBE OUTSOURCE THIS and say open subschemes of affines have zariski atlas)

Corollary. Let X be a quasi-projective scheme that is a sheaf. Then X is a 0-stack.

*Proof.* It suffices to see that X has a zariski atlas. Use .

**Definition 7.3.** A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

**Lemma 7.4.** Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

**Lemma 7.5.** Given a local property P of morphisms of modal n-types, a morphism  $f: X \to Y$  has P if there exists an n-atlas of f having P.

The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n = 0, 1.

#### 8 Descent

**Theorem 8.1.** Let T be a modal n-type. The Proposition, that P is a (smooth) n-stack, is modal.

# 9 Fundamental Theorem of algebraic spaces

### 9.1 For groupoids

**Lemma 9.1.** If  $R \to X \to X$  is a  $\mathbb{T}$ -htpy-coequalizer diagram of two  $\mathbb{T}$ -covers between affines, then X is a 1-stack.

#### 9.2 For sets

**Lemma 9.2.** Denote  $\mathbb{T}Set$  for the sets that are  $\mathbb{T}$ -sheaves. Assume given a  $\mathbb{T}set\ X$  then the following maps are mutually inverse

$$\begin{split} \sum_{R:X\to X\to\mathbb{T}\,\text{Prop}} R \text{ equivalence relation} &\simeq \sum_{Y:\mathbb{T}\mathsf{Set}} \sum_{p:X\to Y} p \text{ } \mathbb{T} \text{surjective} \\ R &\mapsto (X/R,[\ \ ]) \\ \lambda x,y.(p(x)=p(y)) &\leftarrow (Y,p) \end{split}$$

where X/R is defined by applying  $L_T\|_{-}\|_{0}$  at the higher inductive type X//R.

- *Proof.* Well-definedness: The map  $[\_]: X \to ||X//R||_0 \to L_T ||X//R||_0$  is the composition of a surjective with a  $\mathbb{T}$ -surjective map [ref?], hence its  $\mathbb{T}$ -surjective. Conversely given (Y,p) as Y is a sheaf, we have for all x,y:X that  $p(x)=_Y p(y)$  is a sheaf.
  - If x, y : X then we have a chain of equivalences

$$R(x,y) \simeq (\bar{x} =_{\|X//R\|_0} \bar{y}) \to ([x] =_{L_T\|X//R\|_0} [y])$$

where the first map is plain HoTT and the second map is ap, i.e. the unit of the modality [ref?], but as the  $\bar{x} = \|X//R\|_0$   $\bar{y}$  is already a sheaf, it is an isomorphism as well.

Let (Y, p) be in the RHS. Let R(x, y) = (p(x) = p(y)) : T Prop. By plain HoTT, There is a map η : X//R → Y (defined by the universal property of the set truncation and by induction on the higher inductive type X//R on canonical terms through the map p: X → Y). I claim η exhibits Y as the localization for T Set-modality of X//R. Let T be another T Set equipped with a map X//R → T. By precomposition we obtain a map X → T. Claim: it factors uniquely through p: X → Y.

$$X \longrightarrow X//R \longrightarrow_{\exists !} T$$

Proof:

Existence: We want to define a map  $Y \to T$ . Let y: Y. As p is  $\mathbb{T}$ -surjective and T is a sheaf, we may assume we merely have some element in the fiber of p over y. Now push this element through

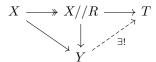
$$\|\operatorname{fib}_p y\| \to \|X//R\|_0 \to T$$

where the first map is by Plain HoTT and the second one is induced from  $X//R \to T$  by assumption and the fact that T is a set.. One can easily check this makes the diagram commute. Uniqueness follows from  $X \to Y$  beeing  $\mathbb{T}$ -surjective and the following Fact: Two parellel maps  $Y \rightrightarrows T$  into a  $\mathbb{T}$  Set T are already equal if the become equal after

precomposition with a T-surjection  $X \to Y$ .

Proof of the fact: Let y:Y. The goal is an identity type of a  $\mathbb{T}$  Set, hence a  $\mathbb{T}$  Prop. Hence As the fiber over y in X is  $\mathbb{T}$ -merely inhabited, we may assume an actual term in the fiber. As  $X \to Y$  equalizes the arrows, this term allows us to conclude.  $\Box$ (fact)  $\Box$ (Claim)

We apply the fact to the (T-)surjectivity of  $X \to X//R$  to get a unique factorization



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making the right triangle commute. This is what we wanted to show.

**Definition 9.3.** An equivalence relation R on a type X is called:

- redundant if for all x, y : X the proposition R(x, y) is a -1-stack.
- smooth if its and for any y: X its fibers:

$$R_y :\equiv \sum_{x:X} R(x,y)$$

are affine in  $\mathbb{T}$ .

**Lemma 9.4.** Assume that  $\mathbb{T}$  satisfies descent for propositions and for sets 8.1, i.e. that a modal proposition being a (-1)-stack is a sheaf. Assume that a modal set beeing affine in  $\mathbb{T}$  is a sheaf. Assume given a  $\mathbb{T}$ set X, then the following types are equivalent:

- ullet The type of redundant smooth equivalence relations over X.
- The type of  $\mathbb{T}$  sets Y with identity types beeing stacks and an -1-atlas X to Y (in V2 a  $\mathbb{T}$ -cover).

*Proof.* By the equivalence in 9.2, it is enough to check that:

• The identity types in X/R are (-1)-stacks if and only if the relation R is redundant . For any x,y:X we know that:

$$R(x,y) \simeq [x] =_{X/R} [y]$$

so the direct direction is immediate. For the converse we use the assumption that a modal proposition being a (-1)-stack is a sheaf and that the map [ $\_$ ]:  $X \to X/R$  is  $\mathbb{T}$ -surjective.

• The fibers of:

$$[\_]: X \to X/R$$

are affine in  $\mathbb{T}$  if and only if the relation R is smooth. For any y:X we have that:

$$\sum_{x \in X} R(x, y) \simeq \mathrm{fib}_{[-]}([y])$$

so the direct direction is immediate. Here as well the converse follows from  $\mathbb{T}$ -surjectivity of  $[\_]$  and that the topology has descent.

**Corollary.** Assume  $\mathbb{T}$  satisfies descent for propositions and for sets. A type is a 0-stack iff its merely the  $\mathbb{T}$ -quotient of an affine scheme by a smooth equivalence relation.

**Theorem 9.5.** Assume  $\mathbb{T}$  satisfies descent for propositions. The quotient of a 0-stack  $X \in \mathbb{T}$  Set by an 0-smooth equivalence relation R is a 0-stack. TODO

*Proof.* The identity types in X/R are propositional 0-stacks, hence (-1)-Truncations of -1-stacks by 11.2 as desired.

How to find an atlas: todo. How to proceed, if we could choose all atlasses we want at the same time?

**Remark 6.** This is equivalent to saying that 1-stacks that are 0-types are geomeric 0-stacks: One direction we prove later. If R is a 0-smooth equivalence relation on a 0-stack X, then X/R is a 1-stack by observing that any -1-atlas  $X' \to X$  gives a 0-atlas  $X' \to X \to X/R$ . Moreover, X/R is a 0-type, hence by assumption a 0-stack.

**Example 9.6.** There are open affine subschemes U of affine schemes  $\operatorname{Spec} A$ , which are not (disjoints unions of) principal open

*Proof.* Consider  $A = R[x, y, u, v]/(xy + ux^2 + vy^2), X = \operatorname{Spec} A$  and consider the open U = D(x, y).

We cant expect U to be a disjoint union of principal opens (todo). However, D(x,y) is affine: We have maps  $U \to R$  given by  $f = -v/x = (y + ux)/y^2, g = -u/y = (x + vy)/x^2$ . Then  $D(f) \cup D(g) = \operatorname{Spec} R^X$ , as yf + xg = 1 in  $R^U$ . Taking preimages under the affinization map,  $U_f \cup U_g = X$  and one checks this defines an open affine cover (for example :  $U_f \simeq \operatorname{Spec} R[x, u, f^{\pm 1}, g]/(xy + ux^2 + uy^2)$  with y := (1 - gx)/f.) But on both of this open subsets the affinization map is an isomorphism hence the affinization of X is an isomorphism. compare (Hartshorne II.2.17)

**Lemma 9.7.** Let  $f: X \to Y$  be surjective. There exists a Zariski Cover  $X' \to X$  such that  $X' \to Y$  is a Zariski cover iff there exists a Zariski Cover  $X' \to X$ , some  $n: \mathbb{N}$  and an open affine embedding  $X' \hookrightarrow Y^n$  over Y.

### 10 Saturated Topologies revisited

**Lemma 10.1** (1.1). We want that every n-1-atlas of a smooth n-atlas has the additional requirement in the definition of smooth n-atlas. It turns out, that for this topology needs to be saturated: The following are equivalent

- 1. Beeing in the topology descents along  $\mathbb{T}$ -covers between affines, i.e.  $\mathbb{T}$  is saturated.
- 2. A smooth n -stack X that is an affine scheme lies in the Topology  $\mathbb{T}$ .
- 3. Let  $n \geq 0$ . If T is a smooth n-stack, then any n-1-atlas  $U \to T$  satisfies  $U \in \mathbb{T}$ .
- 4. If  $U \xrightarrow{f} V \xrightarrow{g} W$  are maps between affines and f and gf are  $\mathbb{T}$  covers, then g is a  $\mathbb{T}$  Cover

Proof.  $1 \Rightarrow 2$ 

Induction. This holds for n=-1. Assume it holds for n-1. Choose a n-1-atlas with T source, i.e.  $T\ni\operatorname{Spec} A\to X$  fibered in smooth n-1-stacks. As it is affine, all the fibers of the atlas are affine smooth n-1-stacks, hence by induction they lie in  $\mathbb{T}$ , thus the atlas is a  $\mathbb{T}$ -cover between affines, hence  $X\in\mathbb{T}$ .

 $2 \Rightarrow 3$ 

As  $U \to T$  is fibered in smooth n-1 stacks, all the fibers are in particular smooth n-stacks by 6.3. By stability under dependent sum  $U = \sum_{t:T} U_t$  is a smooth n-stack that is affine, hence by assumption (2) it lies in the topology.

 $3 \Rightarrow 1$ 

Let  $X \to Y$  be a  $\mathbb{T}$ -cover with X affine in  $\mathbb{T}$  and Y affine. Then Y is a smooth 0-stack, But  $Y \to Y$  is a -1-atlas, hence by assumption  $Y \in T$ .

 $4 \Rightarrow 1$ 

Obvious

 $1 \Rightarrow 4$ 

Check fiberwise  $\Box$ 

If  $n \geq$ , replacing  $\mathbb{T}$  by its saturation  $\mathbb{T}'$  does change the notion of (smooth) n-stack, but we have the following statement, that tells us, that if we start with 0- $\mathbb{T}$ -stacks then the notion of smoothness does not see the difference between  $\mathbb{T}$  and its saturation.

**Proposition 10.2.** Let X be a 0-stack that is a weak smooth 0-stack, i.e. there exists a  $\mathbb{T}'$ -atlas  $\mathbb{T}' \ni X' \to X$  (i.e. fibered in  $\mathbb{T}'$ ). Then X is a smooth 0-stack.

*Proof.* Wlog  $X' \in \mathbb{T}$ . Choose a -1-atlas Spec  $A \to X$  (i.e. fibered in  $\mathbb{T}$ ). As the fibers of  $X' \to X$  merely have smooth atlasses  $\tilde{X}'_x \to X'_x$ , we can use Local choice to obtain a commutative diagram  $Y = \sum_{x':X'} \tilde{X}'_x$ 

$$\tilde{X} \xrightarrow{\mathbb{T}} \operatorname{Spec} A$$

$$\mathbb{T} \downarrow \qquad \qquad \downarrow$$

$$X' \xrightarrow{\mathbb{T}'} X$$

As  $Y \to X'$  is a  $\mathbb{T}$ -cover and  $X' \in \mathbb{T}$  we conclude  $Y \in \mathbb{T}$ . Hence we found a smooth  $\mathbb{T}$ -atlas of X.

#### 10.1 Zariski Topology is not saturated

**Example 10.3** (Weird Zariski Atlasses). Assume those equivalent conditions on the Zariski topology. There exist Zariski atlasses of affines Spec A = X which are not of the form  $D(a_1) + \ldots + D(a_n) \to \operatorname{Spec} A$  for  $(a_1, \ldots, a_n) \in Um(A)$ 

*Proof.* Indeed, using the first example, choose  $U \subset \operatorname{Spec} A$  affine not principal open, then choosing a Zariski atlas  $V \to U$  gives  $V + X \to U + X \to X$  where  $V + X \to X$  is a Zariski cover and  $V + X \to U + X$  is a Zariski cover. From (4), we deduce that  $U + X \to X$  is a Zariski cover, but U is not a disjoint union of principal opens in  $\operatorname{Spec} A$ .

**Example 10.4.** Assume those equivalent conditions on the Zariski topology. Every affine open proposition U is principal open!

*Proof.* Let  $V \to U$  be a Zariski atlas. Then  $V+1 \to U+1$  is a Zariski atlas with  $V+1 \in \mathbb{T}$  and U+1 affine, hence by (1)  $U+1 \in \mathbb{T}$ , hence U is a disjoint union of principal opens hence, as it is a proposition, its principal open.

### 11 beeing a stack is indepent of the truncation level

**Lemma 11.1.** Let  $n \ge 0$ . A n-stack is an modal n-type.

*Proof.* The n-Truncationis an n-type. Now conclude by induction.

We want to show that the notion of stack makes sense, i.e. beeing a stack should not depend on the truncation level.

**Lemma 11.2.** Assume  $\mathbb{T}$  is saturated and satisfies descent for propositions. Let P be a modal proposition. Then TFAE

- 1. For some m > 0, P is a m stack
- 2. There exists some fp algebra A such that  $\operatorname{Spec} A \to P$  and P is logically equivalent to  $(\operatorname{Spec} A \in \mathbb{T})$ .
- 3. P is equivalent to  $\|\operatorname{Spec} A\|_{\mathbb{T}}$  for some fp A, i.e. P is a -1-stack.

Proof.

- 1. ⇒ 2. Let Spec  $A \to P$  be a m-1 atlas. Assume Spec  $A \in \mathbb{T}$ . Then  $\|\operatorname{Spec} A\| \to P$  so as P is a sheaf, we have P. Conversely, if x : P, then the fiber over x is Spec A and a smooth m-1 stack, hence belongs to the topology by 10.1.
- 2.  $\Rightarrow$  3. We have to show: There exists some flat algebra such that P is logically equivalent to  $\|\operatorname{Spec} A\|_{\mathbb{T}}$ . By assumption we have  $\operatorname{Spec} A \to P \to (\operatorname{Spec} A \in \mathbb{T})$ , so we deduce  $\|\operatorname{Spec} A\|_{\mathbb{T}} \to P \to (\operatorname{Spec} A \in \mathbb{T})$ , as P is a modal proposition. In particular A is flat. Conversely  $P \to (\operatorname{Spec} A \in \mathbb{T}) \to \|\operatorname{Spec} A\|_{\mathbb{T}}$ , where the first arrow is by assumption.

 $3. \Rightarrow 1. 6.3$ 

**Lemma 11.3.** A smooth -1-stack P is contractible.

*Proof.* Choose a  $\mathbb{T}$ -cover  $\mathbb{T} \ni \operatorname{Spec} A \to P$ . As P is a proposition we have  $\|\operatorname{Spec} A\| \to P$ . As P is a sheaf we have P.

**Example 11.4.** A 0-stack is a  $\mathbb{T}$ -sheaf whose identity types are (-1)- $\mathbb{T}$ runcations of ((affine ?)) schemes and there exists a  $\mathbb{T}$ -atlas Spec  $A \to X$ .

Why are schemes 0-stacks? This holds in special case, for example if the scheme is quasi projective.

**Theorem 11.5.** Let  $\mathbb{T}$  be saturated. Assume the topology satisfies descent Let  $m, n \geq -2$ . Given an n-type T that is a (smooth) m-stack then T is a (smooth) n-stack.

*Proof.* By 6.3 we may assume  $m \ge n \ge -2$ .

If  $m \le 1$  this is clear. Now assume  $m \ge 2$ . Induction. Inductionstart m = 2. Let us prove the case of m = 2, n = 1, the cases  $-2 \le n < 1$  are immediate from this.

Choose a 1-atlas  $X' \to T$ , i.e. its fibered in smooth 1-stacks. As T is a groupoid and X' is a set, the fibers are actually sets, i.e. smooth 0-stacks.

Now consider  $R := X' \times_T X'$ . As X' is in particular a 0-stack and 0-stacks are stable under dependent sums, R will be a 0-stack. Choose a a  $\mathbb{T}$ -cover  $R' \to R$  with R' affine. Now  $R' \to R \to X'$  is a map between affine schemes i.e it is fibered in smooth 0-stacks that are affine. As  $\mathbb{T}$  is saturated, the fibers of  $R' \to X'$  are in  $\mathbb{T}$ . As X'//R' is a 1-stack by ??, it suffices to show that  $X'//R' \to X'//R$  is a  $\mathbb{T}$ -cover. Pick a term in X'//R. As the fiber beeing in  $\mathbb{T}$  is sheaf If additionally T is assumed to be a smooth 2-stack, then we can assume X' to be in the topology. This will force R to be a smooth 0-stack, so we may choose R' Assume m > 2 and the statement is proven for all (n', m') < (n, m) in lexicographical ordering. As the identity types of T are n-1-types and m-1 stacks by induction they are n-1 stacks. Let  $X \to T$  be an m-1-atlas, i.e. fibered in smooth m-1-stacks with X affine. The fibers are in particular n-1-types, so by induction they are smooth n-1-stacks. Hence  $X \to T$  is an n-1-atlas. If, additionally T is assumed to be a smooth m-stack, we can choose  $X \in \mathbb{T}$ , hence  $X \to T$  witnesses that T is a smooth n-stack.

## 12 Stability under Quotients

**Definition 12.1.** A morphism between n-stacks is smooth if it is fibered in

- $\mathbb{T}$  if n < 0
- smooth n-stacks if n > 0.

**Lemma 12.2.** Let C be a class of types stable under finite limits, i.e. containing 1, stable under dependent sums and identity types. The class  $\mathsf{HasAtlas}_C$  of types Y which admit a map  $\mathsf{Spec}\,A \to Y$  fibered in C is stable under finite limits

*Proof.* Obviously 1 has an atlas, and the class of types admitting an atlas is stable by  $\sum$  by 4.4. It remains to show, that identity types in Y have an atlas provided that Y has an atlas

By assumption we can choose a map  $p:V\to Y$  out of an affine fibered in C. Let y,y':Y. Then we have the map

$$(\operatorname{fib}_p y) \times_V (\operatorname{fib}_p y') \to y = y'$$
 
$$(v, q: y = pv), (v', q': y' = pv'), (h: v = v') \mapsto q \cdot h \cdot q'^{-1}$$

The fiber over j: y = y' looks like

$$\sum_{v} (\underbrace{\sum_{v'} (h:v=v')}) \times (q:y=pv) \times (q':y'=pv') \times (q\cdot h\cdot q'^{-1}=j) \simeq \sum_{v} (v=py) \simeq \operatorname{fib}_{p} y$$

Hence the map is fibered in C. It suffices to show, that  $(\operatorname{fib}_p y) \times_V (\operatorname{fib}_p y')$  has an atlas, because then we can compose such an atlas with the above map to obtain an atlas of y = y'. By assumption the fibers of p have an atlas, so we can choose  $q:W\to\operatorname{fib}_p y,q':W'\to\operatorname{fib}_p y'$  atlasses. Then  $W\times_V W'\to (\operatorname{fib}_p y)\times_V (\operatorname{fib}_p y')$  is an atlas: The domain is a fiber product of affines, hence affine. The fiber over (x,x') is equivalent to the product of fibers  $(\operatorname{fib}_q x)\times (\operatorname{fib}_{q'} x')$  which is in C by stability under dependent sums (so in particular under finite products).

**Theorem 12.3.** Let  $f: X \to Y$  be a  $\mathbb{T}$ -surjective smooth morphism between modal n-types. If X is a (smooth) stack, then Y a (smooth) stack.

(\*) This can only hold if we define -1-stacks to be modal propositions with a -2-atlas Spec  $A \to P$ , i.e. algebraic propositions 5.3

*Proof.* Induction. For n=-2 its clear. Let X be a n-stack. Lets first construct the n-1-atlas of Y. We merely find a  $V \twoheadrightarrow X$  which is an n-1-atlas. Then  $V \to X \to Y$  is an n-atlas because it is  $\mathbb{T}$ -surjective and is fibered in the correct  $\Sigma$ -stable class of types, i.e.  $\mathbb{T}$  if  $n \le 1$  and smooth n-1-stacks for n > 1. Hence Y is an n+1-stack. As Y is an n-type, Y is an n-stack 11.5.

If additionally X is assumed to be smooth, then V can be assumed to lie in  $\mathbb{T}$  which directly gives us that Y has a smooth atlas.

It remains to show that the identity types of Y are n-1-stacks. As Y has an n-1-atlas, by 12.2 we find some n-1-atlas  $p:W\to y=y'$ . The map is smooth. If  $n=0,\ y=y'$  is a -1-stack by (\*). If n>0, W is an n-1-stack and p is smooth, so by induction y=y' is an n-1-stack.

**Remark 7** (Using descent but not induction). Hugo suggested an alternative argument proving that the identity types of Y are n-1-stacks, which presumable avoids 11.5 but uses descent for n-stacks: For x:X,y:Y we have that

$$(f(x) = y) \simeq (1 \times_X \operatorname{fib}_f y)$$

is an n-stack by stability under  $\sum$ . Because it is an n-1-type, it is a n-1-stack by 11.5. Now conclude that every identity type of Y is an n-1-stack by using descent for n-1-stacks and  $\mathbb{T}$ -surjectivity of f.

# 13 Local properties

**Definition 13.1.** Let Cov be the property of morphisms of n-stacks defined by asking that the morphism is  $\mathbb{T}$ -surjective and fibered in smooth n-stacks. Its stable under basechange. A property of n-stacks is local if P(1) holds, P is stable by dependent sums and given a  $Cover\ X \to Y$  we have PX iff PY.

**Example 13.2.** beeing smooth n-stack is a local property of stacks.

*Proof.* We have to show: If  $f: X \to Y$  is a T-surjective map fibered in smooth n-stacks between n-stacks, then X is a smooth n-stack iff Y is a smooth n-stack. The only if is clear by stability under dependent sums. The other direction is 12.3.

**Definition 13.3.** A property of morphisms between n-stacks is local, if it is satisfied by identities, stable under composition and basechange/descent along Cov-maps, precomposition/right cancellability with Cov-maps.

**Lemma 13.4.** Given a local property of types P. Then beeing fibered in P is a local property of morphisms.

**Lemma 13.5** ([ref?]). Given a local property P of morphisms of n-stacks, a morphism  $f: X \to Y$  has P if there exists an n-atlas of f having P.

Example 13.6. A morphism of n-stacks is smooth iff there exists an n-atlas of f

such that  $\tilde{f}$  is a  $\mathbb{T}$ -cover.

The previous lemma tells us that we have the correct notion of smooth morphisms between n-stacks for n=0,1.