### The Mathematics of Origami

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Carleton College

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# Origami? Math?

Image of traditional crane; image of fortune teller

# Origami? Math?

Image of solar panel deployment; image of chemistry microstructure stuff

### Overview

- Origami Constructions
  - Axioms, Euclid &c.
- 2 The Basics of Foldability
  - Local Flat Foldability
  - General Foldability
  - Knots
- Map Folding: An Open Problem
  - Overview
  - Linear Orderings
  - Formal Proof
- Combinatorics of Origami
  - Finding a good problem
  - Meanders
  - Our Method



#### General

- i. Brief History
- ii. Independence

### Seven Axioms of Origami

- i. Given two points  $p_1$  and  $p_2$ , we can fold a line connecting them.
- ii. Given two points  $p_1$  and  $p_2$ , we can fold  $p_1$  onto  $p_2$ .
- iii. Given two lines  $l_1$  and  $l_2$ , we can fold line  $l_1$  onto  $l_2$ .
- iv. Given a point  $p_1$  and a line  $l_1$ , we can make a fold perpendicular to  $l_1$  passing through the point  $p_1$ .

### Seven Axioms of Origami

- v. Given two points  $p_1$  and  $p_2$  and a line  $l_1$ , we can make a fold that places  $p_1$  onto  $l_1$  and passes through the point  $p_2$ .
- vi. Given two points  $p_1$  and  $p_2$  and two lines  $l_1$  and  $l_2$ , we can make a fold that places p1 onto line  $l_1$  and places  $p_2$  onto line  $l_2$ .
- vii. Given a point  $p_1$  and two lines  $l_1$  and  $l_2$ , we can make a fold perpendicular to  $l_2$  that places  $p_1$  onto line  $l_1$ .

### Seven Axioms of Origami

- i. The first six axioms allow all quadratic and cubic equations with rational coefficients to be solved.
- ii. They also allow two of the three problems of antiquity, the trisection of an angle and the doubling of the cube, to be constructed.

# Greek Problems of Antiquity

#### Problems of Antiquity

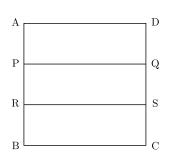
These were a trio of geometric problems whose solutions were attempted solely through the use of compass and straight-edge.

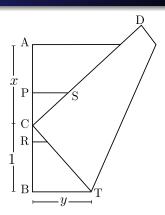
- i. Angle Trisection
- ii. Cube Duplication
- iii. Circle Squaring

### **Angle Trisection**

INCLUDE PARTIALLY FOLDED DIAGRAMS

# **Cube Duplication**





i. 
$$\frac{\alpha}{\beta} = \sqrt[3]{2}$$

ii. Thus, a cube with a side length  $\alpha$  will have twice the volume of a cube with side length  $\beta$ .

## Squaring the Circle

- i. Impossible
- ii. π

#### Constructible Numbers

- i. Given two points  $p_0$  and  $p_1$ , construct a third point  $p_1'$  a distance  $|p_0p_1|$  from point  $p_0$  such that  $\overline{p_0p_1'}$  is perpendicular to  $\overline{p_0p_1}$ .
- ii. Given two points  $p_0$  and  $p_1$ , a third point  $p_2$  can be constructed such that  $p_2$  is collinear with  $p_0$  and  $p_1$ , thus  $|p_0p_1|=|p_1p_2|$ .

#### Constructible Numbers

- iii. Given two constructible numbers  $\alpha$  and  $\beta$ , we can construct  $\frac{\alpha}{\beta}$ , their ratio.
- iv. Given two constructible numbers  $\alpha$  and  $\beta$ , we can construct their sum  $\alpha + \beta$  or their difference  $\alpha \beta$ .
- v. Given two constructible numbers  $\alpha$  and  $\beta$ , we can construct  $\alpha\beta$ , their product.

Given two constructible numbers  $\alpha$  and  $\beta$ , we can construct  $\frac{\alpha}{\beta}$ , their ratio.

INCLUDE PARTIALLY FOLDED DIAGRAMS AS WELL AS FINAL FOLDED DIAGRAM

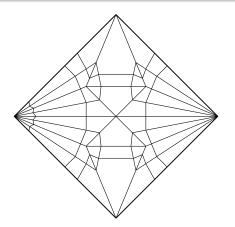
#### Constructible Numbers

- i. Since the set of constructible numbers is closed under addition, subtraction, multiplication, and division, it can be concluded that the set of constructible numbers form a field.
- ii. Ultimately, the field of origami constructible numbers are closed under taking both square roots and cube roots.
- iii. The construction of the square root of any constructible number implies the field of origami constructible numbers contains the field of compass and straightedge constructible numbers.

### Future Work

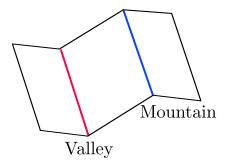
- i. Efficiency
- ii. Optimality

### An example of a crease pattern:

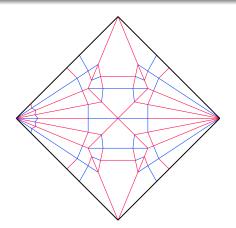


### Crease Assignments

A crease pattern doesn't contain all the information about a model, however.



A crease pattern with mountain-valley assignment:



#### Two questions

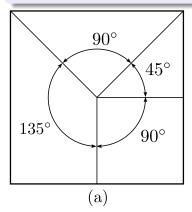
For a given crease pattern,

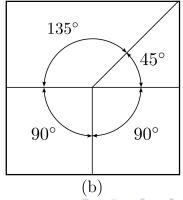
- i. is there a way to fold it flat?
- ii. is there a way to fold it at all?

### Flat-Foldability Conditions

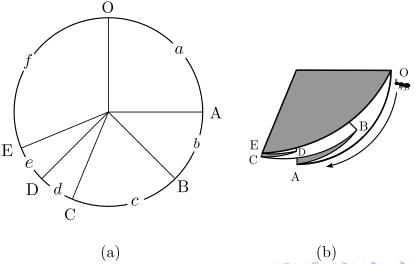
#### Kawasaki's Theorem (1989)

The alternate angles around a vertex must sum to  $180^{\circ}$ .





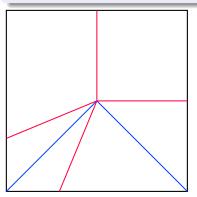
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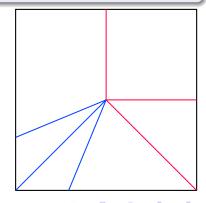


### Flat-Foldability Conditions

#### Maekawa's Theorem (1986)

The sum of mountain + valley is  $\pm 2$ .



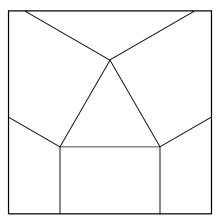


## More Complicated Crease Patterns

How can we determine flat foldability for a more complicated pattern?

# More Complicated Crease Patterns

Here's a crease pattern that can't fold flat!



# **Foldings**

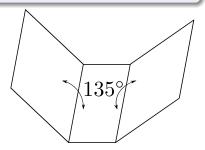
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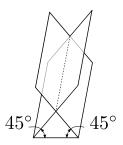
### **Foldings**

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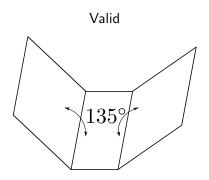
# An Invalid Folding

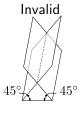
This folding has self-intersection.



## Valid Foldings

A valid folding is one that doesn't cause any self-intersecton.





### Drawing pictures on Crease Patterns

#### Question

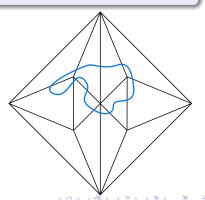
If we draw something on the flat piece of paper before folding it up, what can happen when we fold it?

### Drawing pictures on Crease Patterns

#### Question

If we draw something on the flat piece of paper before folding it up, what can happen when we fold it?

We'll see what happens when we draw Jordan curves on the crease patterns, like the picture on the right.

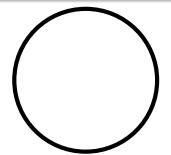


## Knot Theory 101

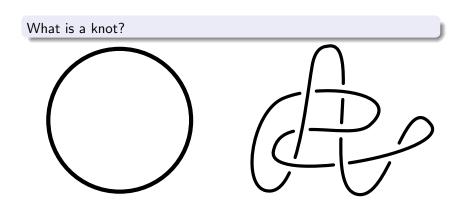
What is a knot?

## Knot Theory 101

#### What is a knot?

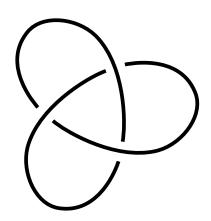


## Knot Theory 101



## Knot Theory 101

This is the *trefoil knot*. We can't "untangle" it, no matter how hard we try.

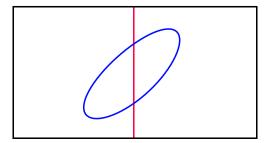


#### **Theorem**

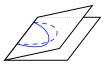
A folding of a crease pattern is valid if and only if every Jordan curve embedded in the paper before folding is mapped to the unknot after folding.



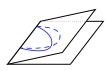






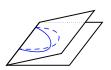










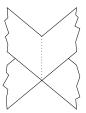






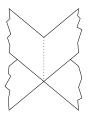


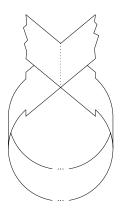
Invalid folding means intersection.





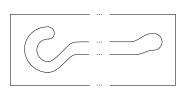
We can connect up the paper like this.

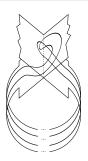






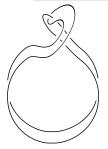
We can draw a curve like this.

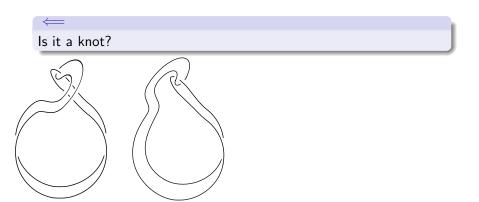


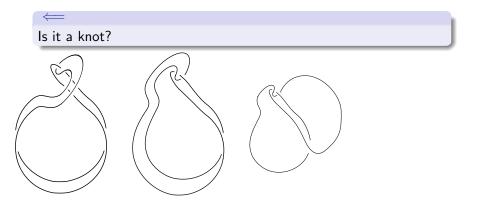


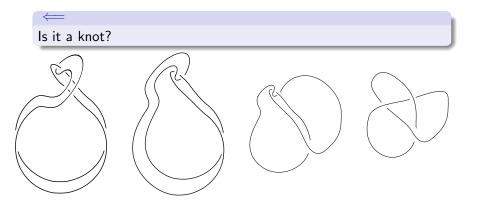


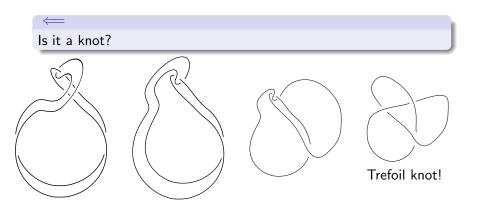
Is it a knot?











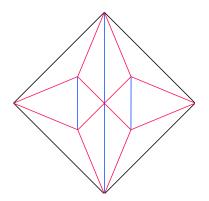
#### Now what?

- Maybe a characteristic for crease patterns
- Finding ways to describe the complexity of a crease pattern or a folding (e.g. the number of self-intersections)

# Determining flat-foldability

### Question:

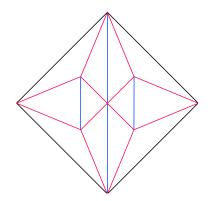
Is this crease pattern flat-foldable?

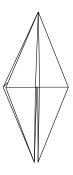


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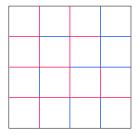
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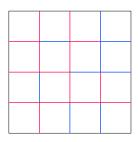


# Introduction to Maps

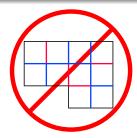


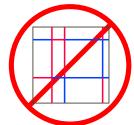


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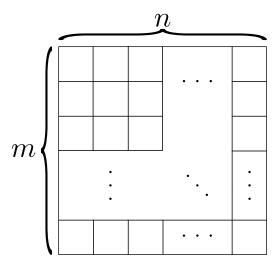








## Introduction to Maps



### Map Folding: An Open Problem

### Open Problem:

How hard is it to determine whether or not a map is flat-foldable?

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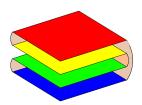
How hard is it to determine whether or not a map is flat-foldable?

#### Easier Problem:

How hard is it to determine whether or not a map is *flat-folded*?

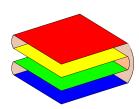
### Question:

How can we represent a folded form?



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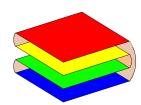
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 $\mathsf{red} \to \mathsf{yellow} \to \mathsf{green} \to \mathsf{blue}$ 

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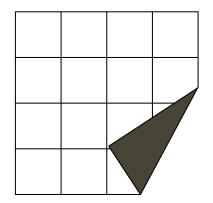
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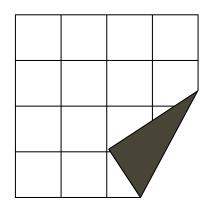
### Note:

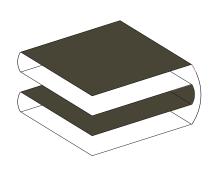
Each linear ordering corresponds with exactly one folded form.

### Question:

How can we tell if a linear ordering is realizable by a given crease pattern?





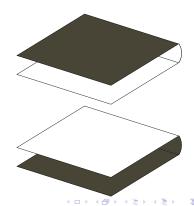


### Goal:

	?	?	?
	?	?	?
?	?	?	?
?	?	?	?

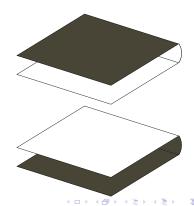
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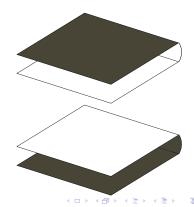
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?	?	?	?
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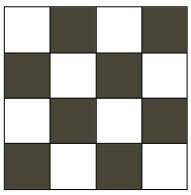


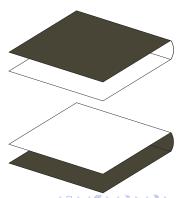
### Goal:

			?
		?	?
	?	?	?
?	?	?	?



### Goal:





### Partial Ordering

### Definition:

A partial ordering P on a set S is a set of ordered pairs of elements of S that orders some of the elements of S.

### Example:

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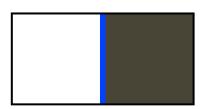
$$c \rightarrow d \rightarrow a \rightarrow b$$

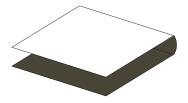


#### Goal:

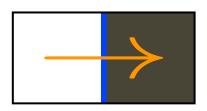


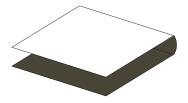
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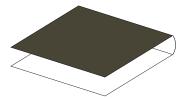


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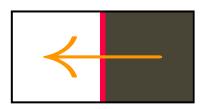


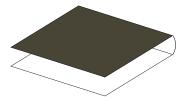
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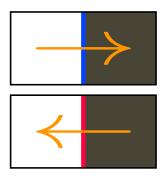


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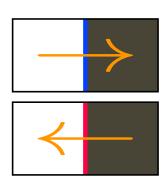


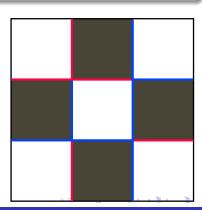


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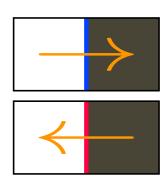


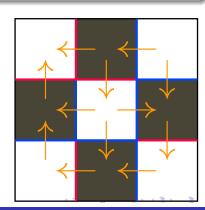
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#### Goal:





## Linear Orderings

Is satisfying this partial ordering enough to ensure foldability?

# Linear Orderings

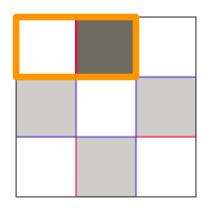
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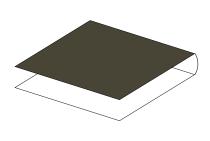
# Linear Orderings

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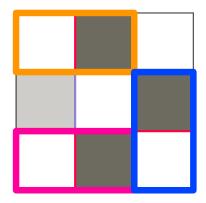


#### **Butterflies**

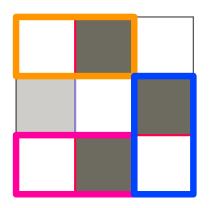


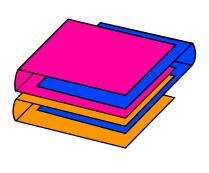


#### Butterflies



#### Butterflies





### **Butterfly Condition**

#### Goal:

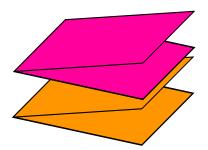
Enumerate the realizable configurations of twin butterfly pairs.



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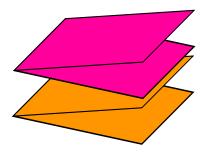
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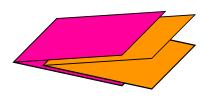


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Enumerate the realizable configurations of twin butterfly pairs.





#### Theorem: (Nishat and Whitesides 2013)

A linear ordering  $\mathcal{L}$  of faces is flat-foldable if and only if (i)  $\mathcal{L}$  satisfies the partial ordering given by the map and (ii) every pair of twin butterflies stacks or nests in  $\mathcal{L}$ .

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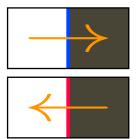
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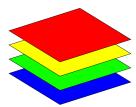
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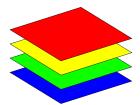
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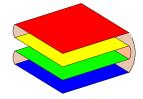


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Algorithm for determining whether or not a linear ordering is valid:

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**Summary:** If someone says "This map can be flat-folded, here is the folding," we can quickly test whether or not they were correct.

# **Counting Problems**

# Counting Problems in Origami

- How many ways can we fold an  $m \times n$  map?
- How many  $m \times n$  map crease patterns have a valid folding?

# Counting Problems in Origami

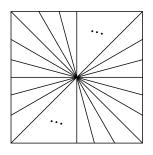
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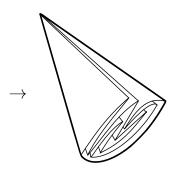
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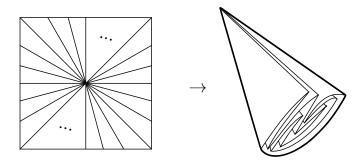
#### Star Patterns

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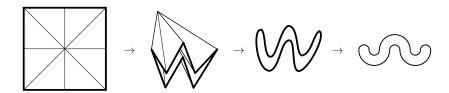
#### Star Patterns



**Question:** How many foldings are generated by star patterns with 2n creases?

# Representing foldings of Star Patterns

## Representing foldings of Star Patterns



#### Meanders

#### Definition (Closed Meanders)

An *closed meander* of order n has two collections of n arches such that when they are placed opposite each other on a line, they form a Jordan curve.

$$n = 1 \left\{ \bigcirc \right\}$$
 $n = 2 \left\{ \bigcirc , \bigcirc \right\}$ 
 $n = 3 \left\{ \bigcirc , \bigcirc , \bigcirc , \bigcirc , \bigcirc , \bigcirc , \bigcirc \right\}$ 

$$n=10 \left\{ \begin{array}{c} \ldots \, , \end{array} \right. \left. \begin{array}{c} \end{array} \right. \left. \left. \begin{array}{c} \end{array} \right. \left. \left. \begin{array}{c} \end{array} \right. \left. \left. \begin{array}{c} \end{array} \right. \left. \begin{array}{c} \end{array}$$

Order n	# Meanders M <sub>n</sub>
1	1
2	2
3	8
4	42
5	262
:	÷

Table: The sequence of Meandric Numbers

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Order n	# Meanders M <sub>n</sub>
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10	8,152,860
:	:

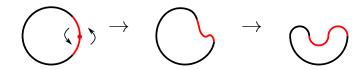
Table: The sequence of Meandric Numbers

# Our Technique

#### Game Plan

$$n=1$$
 O 
$$n=2$$

#### Game Plan



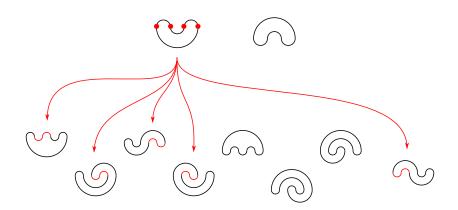
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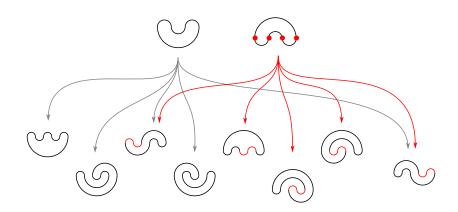


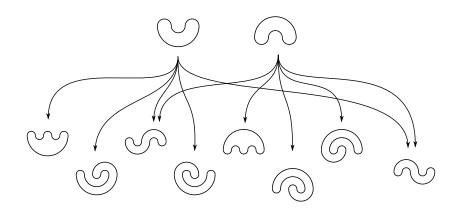
**Idea:** Produce larger meanders by adding "twists" to smaller meanders









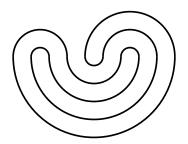


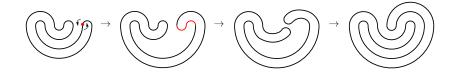
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#### Meanders

#### **Theorem**

For every meander of order n can be produced from some meander of order n-1 with a single twist and at most one shuffle.

#### Two issues to address:

- Different twists yield the same meander (double counting)
- Shuffling is necessary but greatly decreases double counting

# Simple Meanders

#### Definition (Simple Meander)

A *simple meander* is a meander of order n that can be constructed without shuffling.

## Simple Meanders

#### Theorem (Recurrence Relation for Simple Meanders)

Let 
$$\mathbb{P}(k,n) = \{(x_1,x_2,\ldots,x_n) : \sum x_i = k \text{ and } x_i \geq 0 \forall i\}$$
. Then

$$r(n) = \sum_{i=1}^{n} \sum_{P \in \mathbb{P}(i, n+1)} \prod_{k=1}^{n+1} r(P_k)$$

$$H(n) = \sum_{i=1}^{n} r(i) \sum_{P \in \mathbb{P}(i,n)} \prod_{k=1}^{n} H(P_k)$$

$$M_n^s = 2H(n)$$

# Make sure to stop by our exhibit in the Gould Library!

Opening Spring Term 2016

#### References



Rahnuma Islam Nishat and Sue Whitesides (2013)

Map Folding

Canadian Conference on Computational Geomery (2013), p49-52