

CS 6001 Homework 2

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1 Find an x that satisfies the following linear congruences

$$x \equiv 2 \pmod{5}$$

$$x \equiv 3 \pmod{7}$$

$$x \equiv 8 \pmod{11}$$

To do so, use Chinese Remainder Theorem

$$x \equiv a_i \text{ mod } n_i$$

$$x = N_1 * b_1 * a_1 + N_2 * b_2 * a_2 + N_3 * b_3 * a_3$$

$$n = 5 * 7 * 11 = 385$$

$$a_1 = 2$$

$$n_1 = 5$$

$$N_1 = n/n_1 = 77$$

$$a_2 = 3$$

$$n_2 = 7$$

$$N_2 = n/n_2 = 55$$

$$a_3 = 8$$

$$n_3 = 11$$

$$N_3 = n/n_3 = 35$$

$$N_i * b_i \equiv 1 \text{ mod } n_i$$

$$N_1 * b_1 \equiv 1 \text{ mod } n_1$$

$$77 * b_1 \equiv 1 \text{ mod } 5$$

$$b_1 \equiv 3$$

$$55 * b_2 \equiv 1 \text{ mod } 7$$

$$b_2 = 6$$

$$35 * b_3 \equiv 1 \text{ mod } 11$$

$$b_3 = 6$$

$$x = 77 * 3 * 2 + 55 * 6 * 3 + 35 * 6 * 8$$

$$x = 3132$$

2 Discuss some useful applications of the Chinese Remainder Theorem

CRT is useful for secret sharing. If you have some secret code x , and you want it to be shared among n people, however you also wish that any subset of the n people cannot decipher the code x without all n people. To do this, you give each member some function $f_i()$, which is one of congruence equations in Problem 1. To find the x that satisfies all the equations $f_i()$ you must have all the $f_i()$. Having only a subset does not spoil the secret, as they cannot calculate x .

3 Under the RSA encryption scheme, suppose $p = 89$ and $q = 113$.

Let $e = 17$, show how to derive the private key d .

$$\begin{aligned}\varphi(n) &= (p-1)(q-1) \\ &= (89-1)(113-1) = 9856\end{aligned}$$

$$GCD(e, \varphi(n)) = GCD(17, 9856) = 1$$

$$\begin{aligned}d * e \bmod \varphi(n) &= 1 \\ d * 17 \bmod 9856 &= 1\end{aligned}$$

To find d from here, we can use Euclid's algorithm

$$17 * d = 1 \pmod{9856}$$

$$\begin{aligned}9856 &= 17 * 579 + 13 \rightarrow 13 = 9856 - 17 * 579 \\ 17 &= 13 * 1 + 4 \rightarrow 4 = 17 - 13 * 1 \\ 13 &= 4 * 3 + 1 \rightarrow 1 = 13 - 4 * 3 \\ 4 &= 1 * 4\end{aligned}$$

$$GCD(9856, 17) = 1$$

$$\begin{aligned}
1 &= 13 - 4 * 3 \\
&= (9856 - 17 * 579) - (17 - 13) * 3 \\
&= 9856 - 17 * 579 - (17 - (9856 - 17 * 579)) * 3 \\
&= 9856 - 17 * 579 - (17 * 3 - (9856 * 3 - 17 * 579 * 3)) \\
&= 9856 - 17 * 579 - (17 * 3 - 9856 * 3 + 17 * 579 * 3) \\
&= 9856 - 17 * 579 - 17 * 3 + 9856 * 3 - 17 * 579 * 3 \\
&= 9856 * 4 - 17 * 2319 \\
d &= 9856 - 2319 \\
d &= 7537
\end{aligned}$$

Given $m = 65$, compute the encryption of m and verify the encryption is correct by decrypting the encrypted value.

$$E(m) = m^e \text{ mod } n$$

$$n = 89 * 113 = 10057$$

$$E(65) = 65^{17} \text{ mod } 10057$$

$$E(65) = 6619$$

$$e = 6619$$

$$D(e) = e^d \text{ mod } n$$

$$D(e) = 6619^{7537} \text{ mod } 10057$$

$$D(e) = 65$$

4 Show $f_a(x) = ax \text{ mod } n$ is a permutation of Z_n^*

Because $f_a(x) : Z_n^* \mapsto Z_n^*$, $x \in Z_n^*$ must be true. By applying the closure property of multiplicative operators it is known that,

$$x, a \in Z_n^* \implies xa \in Z_n^*$$

For $f_a(x)$ to be a permutation of Z_n^* , it must be a one-to-one function such that

$$a_i x \text{ mod } n \neq a_j x \text{ mod } n : a_i \neq a_j \text{ and } a_i, a_j \in Z_n^*$$

To prove this assume $a_i x \text{ mod } n = a_j x \text{ mod } n$

$$a_i x \text{ mod } n = a_j x \text{ mod } n \implies a_i x + kn = a_j x + k'n : k, k' \text{ are some integer}$$

$$\implies a_i x - a_j x = n(k' - k)$$

$$\implies a_i x \equiv a_j x \text{ mod } n$$

However,

$$\begin{aligned}x \in Z_n^* &\implies \gcd(x, n) = 1 \\ \gcd(x, n) = 1 &\implies xi \not\equiv xj \pmod n : 0 \leq i < j < n \\ xi \not\equiv xj \pmod n &\implies a_i x \not\equiv a_j x \pmod n\end{aligned}$$

Thus $f_{a_i}(x) \neq f_{a_j}(x)$ is one-to-one and since $\forall a, x \in Z_n^* \implies ax \in Z_n^*$ it can be said that $f_a(x)$ is a permutation of Z_n^* .