## 091M4041H-Assignment 4

#### Algorithm Design and Analysis

#### Notice:

- 1. Please submit your answers in hard copy **AND** submit a digital version to UCAS website https://sep.ucas.ac.cn.
- 2. Hard copy should be submitted before 9 am. Dec. 15. and digital version should be submitted before 12pm. Dec. 15.
- 3. Please choose Problem 7 and at least two problems from Problem 1-6.
- 4. Integer Linear Programming is different from the classic Linear Programming that some extra constraints such as

$$x_i$$
 is an integer, for all  $i = 1, 2, \dots, n$ 

or

$$x_i \in \{0,1\}, \text{ for all } i = 1,2,\cdots,n$$

are added.

5. When you give the formulation of an LP or ILP, you should explain all mathematical symbols you are using if not appearing in the problem, and interpret the constraints if necessary.

## 1 Linear-inequality feasibility

Given a set of m linear inequalities on n variables  $x_1, x_2, \dots, x_n$ , the linear-inequality feasibility problem asks if there is a setting of the variables that simultaneously satisfies each of the inequalities.

Show that if we have an algorithm for linear programming, we can use it to solve the linear-inequality feasibility problem. The number of variables and constraints that you use in the linear-programming problem should be polynomial in n and m.

# 2 Airplane Landing Problem

With human lives at stake, an air traffic controller has to schedule the airplanes that are landing at an airport in order to avoid airplane collision. Each airplane i has a time window  $[s_i, t_i]$  during which it can safely land. You must compute the exact time of landing for each airplane that respects these time windows. Furthermore, the airplane landings should be stretched out as much as possible so that the minimum time gap between successive landings is as large as possible.

For example, if the time window of landing three airplanes are [10:00-11:00], [11:20-11:40], [12:00-12:20], and they land at 10:00, 11:20, 12:20 respectively, then the smallest gap is 60 minutes, which occurs between the last two airplanes.

Given n time windows, denoted as  $[s_1, t_1], [s_2, t_2], \dots, [s_n, t_n]$  satisfying  $s_1 < t_1 < s_2 < t_2 < \dots < s_n < t_n$ , you are required to give the exact landing time of each airplane, in which the smallest gap between successive landings is maximized.

Please formulate this problem as an LP, construct an instance and use GLPK or Gurobi or other similar tools to solve it.

### 3 Interval Scheduling Problem

A teaching building has m classrooms in total, and n courses are trying to use them. Each course i  $(i = 1, 2, \dots, n)$  only uses one classroom during time interval  $[S_i, F_i]$   $(F_i > S_i > 0)$ . Considering any two courses can not be carried on in a same classroom at any time, you have to select as many courses as possible and arrange them without any time collision. For simplicity, suppose 2n elements in the set  $\{S_1, F_1, \dots, S_n, F_n\}$  are all different.

- 1. Please use ILP to solve this problem, then construct an instance and use GLPK or Gurobi or other similar tools to solve it.
- 2. If you relax the integral constraints and change ILP to an LP (e.g. change  $x \in \{0,1\}$  to  $0 \le x \le 1$ ), will solution of the LP contains only integers, regardless of values of all  $S_i$  and  $F_i$ ? If it's true, prove it; if it's false, give a counter example. You can use the following lemma for help.

LEMMA If matrix A has only 0, +1 or -1 entries, and each column of A has at most one +1 entry and at most one -1 entry. In addition, the vector b has only integral entries. Then the vertex of polytope  $\{x|Ax\leqslant b,x\geqslant 0\}$  contains only integral entries.

#### 4 Gas Station Placement

Let's consider a long, quiet country road with towns scattered very sparsely along it. Sinopec, largest oil refiner in China, wants to place gas stations along the road. Each gas station is assigned to a nearby town, and the distance between any two gas stations being as small as possible. Suppose there are n towns with distances from one endpoint of the road being  $d_1, d_2, \dots, d_n$ . n gas stations are to be placed along the road, one station for one town. Besides, each station is at most r far away from its correspond town.  $d_1, \dots, d_n$  and r have been given and satisfied  $d_1 < d_2 < \dots < d_n$ ,  $0 < r < d_1$  and  $d_i + r < d_{i+1} - r$  for all i. The objective is to find the optimal placement such that the maximal distance between two successive gas stations is minimized.

Please formulate this problem as an LP.

$$\min \max(x2-x1, x3-x2,...)$$
  
s.t.  $|x_i-d_i| <= r$ 

### 5 Stable Matching Problem

 $n \text{ men } (m_1, m_2, \dots, m_n)$  and  $n \text{ women } (w_1, w_2, \dots, w_n)$ , where each person has ranked all members of the opposite gender, have to make pairs. You need to give a stable matching of the men and women such that there is no unstable pair. Please choose one of the two following known conditions, formulate the problem as an ILP (hint: Problem 1.1 in this assignment), construct an instance and use GLPK or Gurobi or other similar tools to solve it.

- 1. You have known that for every two possible pairs (man  $m_i$  and woman  $w_j$ , man  $m_k$  and woman  $w_l$ ), whether they are stable or not. If they are stable, then  $S_{i,j,k,l} = 1$ ; if not,  $S_{i,j,k,l} = 0$ .  $(i,j,k,l \in \{1,2,\cdots,n\})$
- 2. You have known that for every man  $m_i$ , whether  $m_i$  likes woman  $w_j$  more than  $w_k$ . If he does, then  $p_{i,j,k} = 1$ ; if not,  $p_{i,j,k} = 0$ . Similarly, if woman  $w_i$  likes man  $m_j$  more than  $m_k$ , then  $q_{i,j,k} = 1$ , else  $q_{i,j,k} = 0$ .  $(i, j, k \in \{1, 2, \dots, n\})$

### 6 Duality

Please write the dual problem of the MULTICOMMODITYFLOW problem in *Lec8.pdf*, and give an explanation of the dual variables.

Please also construct an instance, and try to solve both primal and dual problem using GLPK or Gurobi or other similar tools.

## 7 Dual Simplex Algorithm

For the problem

minimize

$$-7x_1 + 7x_2 - 2x_3 - x_4 - 6x_5$$

subject to:

$$3x_1 - x_2 + x_3 - 2x_4 = -3$$
$$2x_1 + x_2 + x_4 + x_5 = 4$$
$$-x_1 + 3x_2 - 3x_4 + x_6 = 12$$
$$x_i \geqslant 0, (i = 1, ..., 6)$$

Implement dual simplex algorithm with your favorate language to solve this problem, and make comparison result with GLPK or Gurobi or other similar tools.