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6	Matching 6.1 KM	<b>10</b> 10	2.1 RDTICC	
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7	Math	12	Struct Node (	
	7.1 Ax+by=gcd         7.2 FFT         7.3 GaussElimination         7.4 Karatsuba         7.5 LinearPrime         7.6 Miller-Rabin         7.7 Mobius         7.8 PollardRho         7.9 Simplex	12 13 13 13 13 14 14 14	<pre>int id,f; Node *L, *R; }tree[MXN]; int n; Node *root;  long long dis2(int x1, int y1, int x2, int y2) {   long long dx = x1-x2;   long long dy = y1-y2;   return dx*dx+dy*dy; }</pre>	
	7.10 Sprague-Grundy	15 15	static bool clipx (Node a a, Node a b) { return a.x\b.x;	

```
static Splay nil, mem[MEM], *pmem;
   n = ip.size();
    for (int i=0; i<n; i++) {</pre>
                                                            Splay *ch[2], *f;
     tree[i].id = i;
                                                            int val, rev, size;
                                                            Splay () : val(-1), rev(0), size(0) {
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
                                                             f = ch[0] = ch[1] = &nil;
   root = build tree(0, n-1, 0);
                                                            Splay (int val) : val( val), rev(0), size(1) {
                                                             f = ch[0] = ch[1] = &nil;
  Node* build tree(int L, int R, int dep) {
   if (L>R) return nullptr;
                                                            bool isr() {
   int M = (L+R)/2;
                                                              return f->ch[0] != this && f->ch[1] != this;
   tree[M].f = dep%2;
   nth element(tree+L, tree+M, tree+R+1, tree[M].f ?
                                                            int dir() {
        cmpy : cmpx);
                                                              return f->ch[0] == this ? 0 : 1;
    tree[M].x1 = tree[M].x2 = tree[M].x;
   tree[M].y1 = tree[M].y2 = tree[M].y;
                                                            void setCh(Splay *c, int d) {
                                                              ch[d] = c;
                                                              if (c != &nil) c->f = this;
   tree[M].L = build tree(L, M-1, dep+1);
   if (tree[M].L) {
                                                              pull();
     tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
                                                            void push() {
     tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
                                                              if (rev) {
                                                                swap(ch[0], ch[1]);
      tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
                                                                if (ch[0] != &nil) ch[0]->rev ^= 1;
                                                                if (ch[1] != &nil) ch[1]->rev ^= 1;
   tree[M].R = build tree(M+1, R, dep+1);
   if (tree[M].R) {
                                                              }
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
                                                            void pull() {
     tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
                                                              size = ch[0] -> size + ch[1] -> size + 1;
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
                                                              if (ch[0] != &nil) ch[0]->f = this;
                                                              if (ch[1] != &nil) ch[1]->f = this;
   return tree+M;
                                                          } Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::
                                                              mem;
  int touch(Node* r, int x, int y, long long d2){
                                                          Splay *nil = &Splay::nil;
   long long dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis || y<r->y1-dis || y>
                                                          void rotate(Splay *x) {
       r->y2+dis)
                                                            Splay *p = x->f;
                                                            int d = x->dir();
      return 0;
   return 1;
                                                            if (!p->isr()) p->f->setCh(x, p->dir());
                                                            else x->f = p->f;
  void nearest(Node* r, int x, int y, int &mID, long
                                                            p->setCh(x->ch[!d], d);
     long &md2) {
                                                            x->setCh(p, !d);
    if (!r || !touch(r, x, y, md2)) return;
                                                           p->pull(); x->pull();
   long long d2 = dis2(r->x, r->y, x, y);
   if (d2 < md2 \mid | (d2 == md2 && mID < r->id)) {
     mID = r->id;
                                                          vector<Splay*> splayVec;
     md2 = d2;
                                                          void splay(Splay *x) {
                                                            splayVec.clear();
    // search order depends on split dim
                                                            for (Splay *q=x;; q=q->f) {
   if ((r->f == 0 \&\& x < r->x) | |
                                                              splayVec.push back(q);
       (r->f == 1 && y < r->y)) {
                                                              if (q->isr()) break;
      nearest(r->L, x, y, mID, md2);
                                                            reverse(begin(splayVec), end(splayVec));
     nearest(r\rightarrow R, x, y, mID, md2);
    } else {
                                                            for (auto it : splayVec) it->push();
      nearest(r->R, x, y, mID, md2);
                                                            while (!x->isr()) {
      nearest(r->L, x, y, mID, md2);
                                                              if (x->f->isr()) rotate(x);
                                                              else if (x->dir()==x->f->dir()) rotate(x->f), rotate
 int query(int x, int y) {
                                                              else rotate(x), rotate(x);
   int id = 1029384756;
   long long d2 = 102938475612345678LL;
   nearest(root, x, y, id, d2);
   return id;
                                                          Splay* access(Splay *x) {
 }
                                                           Splay *q = nil;
}tree;
                                                            for (;x!=nil;x=x->f) {
                                                             splav(x);
                                                             x->setCh(q, 1);
2.2 Link Cut Tree
                                                              q = x;
                                                            return q;
```

struct Splay {

```
const int MXN = 100005;
const int MEM = 100005;
```

void init(vector<pair<int,int>> ip) {

void evert(Splay \*x) {

access(x);

splav(x);

```
x->rev ^= 1;
  x \rightarrow push(); x \rightarrow pull();
void link(Splay *x, Splay *y) {
// evert(x);
 access(x);
 splav(x);
 evert(y);
 x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
// evert(x);
 access(y);
 splay(y);
  v->push();
  y - ch[0] = y - ch[0] - f = nil;
int N, Q;
Splay *vt[MXN];
int ask(Splay *x, Splay *y) {
 access(x);
 access(y);
  splav(x);
 int res = x->f->val;
  if (res == -1) res=x->val;
  return res;
int main(int argc, char** argv) {
  scanf("%d%d", &N, &Q);
  for (int i=1; i<=N; i++)</pre>
   vt[i] = new (Splay::pmem++) Splay(i);
  while (Q--) {
    char cmd[105];
    int u, v;
    scanf("%s", cmd);
    if (cmd[1] == 'i') {
      scanf("%d%d", &u, &v);
      link(vt[v], vt[u]);
    } else if (cmd[0] == 'c') {
      scanf("%d", &v);
      cut(vt[1], vt[v]);
    } else {
      scanf("%d%d", &u, &v);
      int res=ask(vt[u], vt[v]);
      printf("%d \setminus n", res);
  return 0;
```

#### 2.3 Pb Ds Heap

```
#include <bits/extc++.h>
typedef __gnu_pbds::priority_queue<int> heap_t;
heap_t a,b;
int main() {
    a.clear();b.clear();
    a.push(1);a.push(3);b.push(2);b.push(4);
    // merge two heap
    a.join(b);
    assert(a.top() == 4);
    assert(b.empty());
    return 0;
}
```

#### 2.4 Pb Ds Rbtree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
```

```
#define ordered_set tree<int, null_type,less<int>,
    rb_tree_tag,tree_order_statistics_node_update>
int main() {
    ordered_set o_set;
    o_set.insert(5);
    o_set.insert(1);
    o_set.insert(2);
    cout << *(o_set.find_by_order(1)) << endl; // 2
    cout << o_set.order_of_key(4) << endl; // 2
    cout << o_set.order_of_key(5) << endl; // 2
    if (o_set.find(2) != o_set.end())
        o_set.erase(o_set.find(2));
    cout << *(o_set.find_by_order(1)) << endl; // 5
    cout << *(o_set.order_of_key(4) << endl; // 5
    cout << *(o_set.order_of_key(4) << endl; // 5</pre>
```

## 2.5 SparseTable

```
const int MAXN = 200005;
const int lqN = 20;
struct SP{ //sparse table
  int Sp[MAXN][lqN];
  function<int(int,int)> opt;
  void build(int n, int *a){ // 0 base
    for (int i=0 ;i<n; i++) Sp[i][0]=a[i];</pre>
    for (int h=1; h<1gN; h++) {</pre>
      int len = 1<<(h-1), i=0;</pre>
      for (; i+len<n; i++)</pre>
        Sp[i][h] = opt(Sp[i][h-1], Sp[i+len][h-1]);
      for (; i<n; i++)</pre>
        Sp[i][h] = Sp[i][h-1];
  int query(int 1, int r){
    int h = __lg(r-l+1);
    int len = 1<<h;
    return opt( Sp[l][h] , Sp[r-len+1][h] );
  }
};
```

#### 2.6 Treap

```
#include <bits/stdc++.h>
using namespace std;
template<class T,unsigned seed>class treap{
 public:
   struct node{
     T data;
      int size:
      node *1, *r;
     node(T d){
       size=1;
       data=d;
        l=r=NULL;
      inline void up(){
       size=1:
        if(l)size+=l->size;
       if(r)size+=r->size;
      inline void down() {
    } *root;
    inline int size(node *p) {return p?p->size:0;}
    inline bool ran(node *a, node *b) {
      static unsigned x=seed;
      x=0xdefaced*x+1;
      unsigned all=size(a)+size(b);
      return (x%all+all)%all<size(a);
   void clear(node *&p) {
```

```
~treap(){clear(root);}
    void split(node *o, node *&a, node *&b, int k) {
      if(!k) a=NULL, b=o;
      else if(size(o) == k) a=o, b=NULL;
      else{
        o->down();
        if(k<=size(o->l)){
          split(o->l,a,b->l,k);
          b->up();
        }else{
          split(o->r,a->r,b,k-size(o->l)-1);
          a - \sup();
      }
    void merge(node *&o, node *a, node *b) {
      if(!a||!b)o=a?a:b;
        if(ran(a,b)){
         a->down();
          o=a;
          merge(o->r,a->r,b);
        }else{
         b->down();
          o=b;
          merge(o->1,a,b->1);
        o->up();
      }
    void build(node *&p,int l,int r,T *s) {
     if(l>r)return;
     int mid=(l+r)>>1;
     p=new node(s[mid]);
      build(p->1,1,mid-1,s);
     build(p->r,mid+1,r,s);
     p->up();
    inline int rank(T data){
     node *p=root;
      int cnt=0:
      while(p) {
        if (data<=p->data)p=p->l;
        else cnt+=size(p->1)+1,p=p->r;
      return cnt;
    inline void insert(node *&p,T data,int k) {
     node *a, *b, *now;
      split(p,a,b,k);
     now=new node(data);
     merge(a,a,now);
     merge(p,a,b);
    inline void remove(node *&p, int k) {
     node *a, *b, *res, *die;
      split(p, a, res, k);
      if (res == NULL) return;
      split(res, die, b, 1);
     merge(a, a, b);
     if (size(a) > size(b)) p = a;
      else p = b;
     clear(die);
};
treap<T ,20141223>bst;
int main(){
 bst.remove(bst.root, bst.rank(E));
 bst.insert(bst.root, E, bst.rank(E));
```

if(p)clear(p->1),clear(p->r),delete p,p=NULL;

## 3 Flow

#### 3.1 Dinic

```
//Dinic
#define V 1000
struct edge{
    edge(){}
    edge(int a,int b,int c):to(a),cap(b),rev(c){}
    int to,cap,rev;
};
vector<edge> g[V];
int level[V];
int iter[V];
void add edge(int from,int to,int cap) {
    g[from].push_back(edge(to,cap,g[to].size()));
    g[to].push back(edge(from, 0, g[from].size()-1));
void bfs(int s) {
    memset(level, -1, sizeof(level));
    queue<int>que;
    level[s]=0;
    que.push(s);
    while(!que.empty()){
        int v=que.front();
        que.pop();
        for (int q=0; q<g[v].size(); q++) {</pre>
            edge &e=g[v][q];
             if(e.cap>0&&level[e.to]<0){
                level[e.to]=level[v]+1;
                 que.push(e.to);
        }
int dfs(int v,int t,int f){
    if(v==t)return f;
    for(int &q=iter[v];q<g[v].size();++q){</pre>
        edge &e=g[v][q];
        if(e.cap>0&&level[v]<level[e.to]){</pre>
             int d=dfs(e.to,t,min(f,e.cap));
             if(d>0){
                 e.cap-=d;
                 g[e.to][e.rev].cap+=d;
                 return d;
        }
    return 0;
int max flow(int s,int t) {
    int flow=0;
    for(;;){
        bfs(s);
        if(level[t]<0)return flow;</pre>
        memset(iter, 0, sizeof(iter));
        int f:
        while((f=dfs(s,t,1e9))>0)
               flow+=f:
    }
int main(){
  for (int i=0;i<V;++i)</pre>
    g[i].clear();
```

## 3.2 Gomory Hu

```
Construct of Gomory Hu Tree

1. make sure the whole graph is clear
2. set node 0 as root, also be the parent of other nodes.
```

```
3. for every node i > 0, we run maxflow from i to
    parent[i]
4. hense we know the weight between i and parent[i]
5. for each node j > i, if j is at the same side with i
  make the parent of j as i
int e[MAXN][MAXN];
int p[MAXN];
Dinic D; // original graph
void gomory_hu() {
    fill(p, p+n, 0);
    fill(e[0], e[n], INF);
    for ( int s = 1 ; s < n ; s++ ) {</pre>
        int t = p[s];
        Dinic F = D;
        int tmp = F.max_flow(s, t);
        for ( int i = 1 ; i < s ; i++ )</pre>
            e[s][i] = e[i][s] = min(tmp, e[t][i]);
        for ( int i = s+1 ; i <= n ; i++ )</pre>
            if ( p[i] == t && F.side[i] ) p[i] = s;
```

#### 3.3 Min Cost Flow

```
#include<bits/stdc++.h>
using namespace std;
#define int long long
typedef pair<int,int> P;
struct edge{
    edge(){}
    edge(int a,int b,int c,int d):to(a),cap(b),cost(c),
       rev(d){}
    int to,cap,cost,rev;
};
#define V 210
#define inf 10000000000000000
vector<edge> g[V];
int h[V],dist[V],prev v[V],prev e[V];
void add edge(int from, int to, int cap, int cost) {
    g[from].push back(edge(to,cap,cost,g[to].size()));
    g[to].push back(edge(from, 0, -cost, g[from].size()-1)
       );
int min costflow(int s,int t,int f) {
    int res=0;
    memset(h, 0, sizeof(h));
    while (f>0) {
        priority queue<P, vector<P>, greater<P> >que;
        fill(dist, dist+V, inf);
        dist[s]=0;
        que.push(P(dist[s],s));
        while(!que.empty()){
           P p=que.top();
            que.pop();
            int v=p.second;
            if(dist[v]<p.first)continue;</pre>
            for (int i=0;i<g[v].size();++i){</pre>
                edae &e=a[v][i];
                if(e.cap>0&&dist[e.to]>dist[v]+e.cost+h
                     [v]-h[e.to]){
                    dist[e.to]=dist[v]+e.cost+h[v]-h[e.
                         to];
                    prev v[e.to]=v;
                    prev e[e.to]=i;
                    que.push(P(dist[e.to],e.to));
            }
```

```
if(dist[t]==inf) return -1;
        for (int v=0; v<V; ++v) h[v] +=dist[v];</pre>
         int d=f;
         for (int v=t; v!=s; v=prev v[v]) d=min(d,g[prev v[
             v]][prev e[v]].cap);
         f-=d;
        res+=d*h[t];
         for(int v=t;v!=s;v=prev v[v]){
             edge &e=g[prev v[v]][prev e[v]];
             e.cap-=d;
             g[v][e.rev].cap+=d;
    return res:
#undef int
int main()
#define int long long
    int T,n,m,cost,l,s,t,ans;
    cin>>T:
    while(T--) {
      cin>>n>>m;
        for(int q=0;q<V;++q)g[q].clear();</pre>
         s=m+n;
        t=m+n+1;
        for (int i=0;i<n;++i)</pre>
          for (int j=0; j<m; ++j) {</pre>
             cin>>cost;
             if (cost>0)
               add edge(n+j,i,1,cost);
         for (int i=0; i<m; ++i) {</pre>
           cin>>l;
           add edge(s,n+i,l,0);
         for (int i=0; i<n; ++i)</pre>
          add_edge(i,t,1,0);
         ans=min costflow(s,t,n);
        cout<<ans<<endl:
    return 0;
```

#### 3.4 SW-mincut

```
// all pair min cut
// global min cut
struct SW{ // O(V^3)
  static const int MXN = 514;
  int n, vst[MXN], del[MXN];
  int edge[MXN][MXN], wei[MXN];
  void init(int _n){
    n = _n; FZ(edge); FZ(del);
  void addEdge(int u, int v, int w) {
    edge[u][v] += w; edge[v][u] += w;
  void search(int &s, int &t){
    FZ(vst); FZ(wei);
    s = t = -1;
    while (true) {
      int mx=-1, cur=0;
       for (int i=0; i<n; i++)</pre>
         if (!del[i] && !vst[i] && mx<wei[i])</pre>
          cur = i, mx = wei[i];
       if (mx == -1) break;
       vst[cur] = 1;
       s = t; t = cur;
       for (int i=0; i<n; i++)</pre>
         if (!vst[i] && !del[i]) wei[i] += edge[cur][i];
  int solve(){
    int res = 2147483647;
     for (int i=0,x,y; i<n-1; i++) {</pre>
```

```
search(x,y);
  res = min(res,wei[y]);
  del[y] = 1;
  for (int j=0; j<n; j++)
      edge[x][j] = (edge[j][x] += edge[y][j]);
  }
  return res;
}
}graph;</pre>
```

## 4 Geometry

## 4.1 2Dpoint

```
typedef double Double;
struct Point {
  Double x, y;
  bool operator < (const Point &b)const{</pre>
    //return tie(x,y) < tie(b.x,b.y);
    //return atan2(y,x) < atan2(b.y,b.x);
    assert(0 && "choose compare");
  Point operator + (const Point &b) const{
    return (Point) {x+b.x,y+b.y};
  Point operator - (const Point &b) const{
    return (Point) {x-b.x,y-b.y};
  Point operator * (const Double &d)const{
   return Point(d*x,d*y);
  Double operator * (const Point &b) const{
   return x*b.x + y*b.y;
  Double operator % (const Point &b) const{
    return x*b.y - y*b.x;
  friend Double abs2(const Point &p) {
   return p.x*p.x + p.y*p.y;
  friend Double abs(const Point &p) {
    return sqrt( abs2(p) );
typedef Point Vector;
struct Line{
  Point P; Vector v;
 bool operator < (const Line &b)const{</pre>
    return atan2(v.y,v.x) < atan2(b.v.y,b.v.x);</pre>
};
```

#### 4.2 Circumcentre

```
#include "2Dpoint.cpp"

Point circumcentre(Point &p0, Point &p1, Point &p2) {
   Point a = p1-p0;
   Point b = p2-p0;
   Double c1 = abs2(a)*0.5;
   Double c2 = abs2(b)*0.5;
   Double d = a % b;
   Double x = p0.x + (c1*b.y - c2*a.y) / d;
   Double y = p0.y + (c2*a.x - c1*b.x) / d;
   return {x,y};
}
```

#### 4.3 ConvexHull

#### 4.4 Half Plane Intersection

```
| bool OnLeft(const Line& L, const Point& p) {
  return Cross(L.v,p-L.P)>0;
Point GetIntersection(Line a, Line b) {
  Vector u = a.P-b.P;
  Double t = Cross(b.v,u)/Cross(a.v,b.v);
  return a.P + a.v*t;
int HalfplaneIntersection(Line* L,int n,Point* poly) {
  sort(L,L+n);
  int first, last;
  Point *p = new Point[n];
  Line *q = new Line[n];
  q[first=last=0] = L[0];
  for (int i=1;i<n;i++) {</pre>
    while(first < last && !OnLeft(L[i],p[last-1])) last</pre>
    while(first < last && !OnLeft(L[i],p[first])) first</pre>
        ++:
    q[++last]=L[i];
    if (fabs (Cross (q[last].v, q[last-1].v)) < EPS) {</pre>
       if (OnLeft(q[last],L[i].P)) q[last]=L[i];
    if(first < last) p[last-1]=GetIntersection(q[last</pre>
         -1],q[last]);
  while(first<last && !OnLeft(q[first],p[last-1])) last</pre>
  if(last-first<=1) return 0;</pre>
  p[last]=GetIntersection(q[last],q[first]);
  for(int i=first;i<=last;i++) poly[m++]=p[i];</pre>
  return m;
```

#### 4.5 Intersection Of Two Circle

#### 4.6 Intersection Of Two Lines

```
Point interPnt(Point p1, Point p2, Point q1, Point q2,
    bool &res) {
    Double f1 = cross(p2, q1, p1);
    Double f2 = -cross(p2, q2, p1);
    Double f = (f1 + f2);

    if(fabs(f) < EPS) {
       res = false;
       return {};
    }

    res = true;
    return (f2 / f) * q1 + (f1 / f) * q2;
}</pre>
```

#### 4.7 Smallest Circle

```
#include "circumcentre.cpp"
pair<Point,Double> SmallestCircle(int n, Point p[]) {
  Point *p = new Point[n];
  memcpy(p,_p,sizeof(Point)*n);
  random shuffle(p,p+n);
  Double r2=0;
  Point cen;
  for (int i=0; i<n; i++) {</pre>
    if ( abs2(cen-p[i]) <= r2)continue;</pre>
    cen = p[i], r2=0;
    for (int j=0; j<i; j++) {</pre>
      if ( abs2(cen-p[j]) <= r2)continue;</pre>
      cen = (p[i]+p[j])*0.5;
      r2 = abs2(cen-p[i]);
      for (int k=0; k<j; k++) {</pre>
        if ( abs2(cen-p[k]) <= r2)continue;</pre>
        cen = circumcentre(p[i],p[j],p[k]);
        r2 = abs2(cen-p[k]);
    }
  }
  delete[] p;
  return {cen,r2};
// auto res = SmallestCircle(,);
```

# 5 Graph

#### 5.1 BCC Edge

static const int MXN = 100005;

struct Edge { int v,eid; };

```
邊雙連通
```

```
任意兩點問至少有兩條不重疊的路徑連接,找法:
1. 標記出所有的橋
2. 對全圖進行 DFS,不走橋,每一次 DFS 就是一個新的邊雙連通

// from BCW

struct BccEdge {
```

```
int n,m,step,par[MXN],dfn[MXN],low[MXN];
  vector<Edge> E[MXN];
  DisjointSet djs;
  void init(int _n)
    n = _n; m = 0;
    for (int i=0; i<n; i++) E[i].clear();</pre>
    djs.init(n);
  void add edge(int u, int v) {
    E[u].PB({v, m});
    E[v].PB({u, m});
    m++;
  void DFS(int u, int f, int f_eid) {
    par[u] = f;
    dfn[u] = low[u] = step++;
    for (auto it:E[u]) {
      if (it.eid == f eid) continue;
      int v = it.v;
      if (dfn[v] == -1) {
       DFS(v, u, it.eid);
        low[u] = min(low[u], low[v]);
      } else {
        low[u] = min(low[u], dfn[v]);
    }
  void solve() {
    step = 0;
    memset(dfn, -1, sizeof(int)*n);
    for (int i=0; i<n; i++) {</pre>
      if (dfn[i] == -1) DFS(i, i, -1);
    djs.init(n);
    for (int i=0; i<n; i++) {</pre>
      if (low[i] < dfn[i]) djs.uni(i, par[i]);</pre>
  }
}graph;
```

## 5.2 Dijkstra

```
from heapq import *
INF = 2*10**10000
t = input()
for pp in range(t):
  n, m = map(int, raw_input().split())
  g, d, q = [[] for _ in range(n+1)], [0] + [INF] * n,
      [(0, 0)]
  #for i in range(1, m):
  # a[i], b[i], c[i], l[i], o[i] = map(int, input().
      split())
  for _ in range(m):
    u, v, c, l, o = map(int, raw_input().split())
    g[u] += [(o, v, c, 1)]
  while q:
    u = heappop(q)[1]
    for e in g[u]:
      k = d[u] / e[2]
      if k < 0:
       k = 0
      else:
        k = k * e[3]
      t, v = d[u] + e[0] + k, e[1]
      if t < d[v]:
        d[v] = t
        heappush(q, (d[v], v))
 print(d[n])
```

#### 5.3 Directed MST

```
template<typename T>
struct zhu_liu{
```

```
static const int MAXN=110,MAXM=10005;
  struct node{
    int u, v;
    T w, tag;
    node *1, *r;
    node(int u=0,int v=0,T w=0):u(u),v(v),w(w),tag(0),1
         (0), r(0) \{ \}
    void down() {
      w+=tag;
      if(1)1->tag+=tag;
      if(r)r->tag+=tag;
      tag=0;
    }
  }mem[MAXM];//靜態記憶體
  node *pq[MAXN*2],*E[MAXN*2];
  int st[MAXN*2],id[MAXN*2],m;
  void init(int n){
    for (int i=1;i<=n;++i) {</pre>
      pq[i]=E[i]=0;
      st[i]=id[i]=i;
    m=0;
  node *merge(node *a, node *b) { //skew heap
    if(!a||!b)return a?a:b;
    a->down(),b->down();
    if(b->w<a->w)return merge(b,a);
    swap(a->1,a->r);
    a->l=merge(b,a->l);
    return a;
  void add edge(int u,int v,T w) {
    if(u!=v)pq[v]=merge(pq[v], &(mem[m++]=node(u,v,w)));
  int find(int x,int *st) {
    return st[x] == x?x:st[x] = find(st[x], st);
  T build(int root, int n) {
    T ans=0; int N=n, all=n;
    for (int i=1;i<=N;++i) {</pre>
      if (i==root | |!pq[i]) continue;
      while (pq[i]) {
        pq[i]->down(), E[i]=pq[i];
         pq[i] = merge(pq[i] -> 1, pq[i] -> r);
         if (find (E[i] ->u, id) !=find(i, id))break;
      if (find(E[i]->u,id) == find(i,id)) continue;
      ans+=E[i] \rightarrow w;
      if(find(E[i]->u,st)==find(i,st)){
        if (pq[i])pq[i]->tag-=E[i]->w;
         pq[++N] = pq[i], id[N] = N;
         for (int u=find(E[i]->u,id);u!=i;u=find(E[u]->u,
           if(pq[u])pq[u]->tag-=E[u]->w;
           id[find(u,id)]=N;
          pq[N]=merge(pq[N],pq[u]);
         st[N] = find(i, st);
        id[find(i,id)]=N;
      }else st[find(i,st)]=find(E[i]->u,st),--all;
    return all==1?ans:-INT MAX;//圖不連通就無解
};
```

#### 5.4 LCA

```
| //lv紀錄深度
| //father[多少冪次][誰]
| //已經建好每個人的父親是誰 (father[0][i]已經建好)
| //已經建好深度 (lv[i]已經建好)
| void makePP() {
| for(int i = 1; i < 20; i++) {
| for(int j = 2; j <= n; j++) {
| father[i][j]=father[i-1][ father[i-1][j] ];
| }
```

```
}
int find(int a, int b) {
    if(lv[a] < lv[b]) swap(a,b);
    int need = lv[a] - lv[b];
    for(int i = 0; need!=0; i++) {
        if(need&1) a=father[i][a];
        need >>= 1;
    }
    for(int i = 19 ;i >= 0 ;i--) {
        if(father[i][a] != father[i][b]) {
            a=father[i][a];
            b=father[i][b];
        }
    }
    return a!=b?father[0][a] : a;
}
```

## 5.5 MaximumClique

```
const int MAXN = 105;
int best;
int m ,n;
int num[MAXN];
// int x[MAXN];
int path[MAXN];
int g[MAXN][MAXN];
bool dfs( int *adj, int total, int cnt ){
    int i, j, k;
    int t[MAXN];
    if( total == 0 ){
        if( best < cnt ) {</pre>
             // for(i = 0; i < cnt; i++) path[i] = x[i]
            best = cnt; return true;
        return false;
    for( i = 0; i < total; i++) {</pre>
        if( cnt+(total-i) <= best ) return false;</pre>
        if( cnt+num[adj[i]] <= best ) return false;</pre>
        // x[cnt] = adj[i];
        for( k = 0, j = i+1; j < total; j++ )</pre>
            if( g[ adj[i] ][ adj[j] ] )
                 t[k++] = adj[j];
        if( dfs( t, k, cnt+1 ) ) return true;
    } return false;
int MaximumClique() {
    int i, j, k;
    int adi[MAXN];
    if( n <= 0 ) return 0;</pre>
    best = 0;
    for( i = n-1; i >= 0; i-- ){
        // x[0] = i;
        for (k = 0, j = i+1; j < n; j++)
            if( g[i][j] ) adj[k++] = j;
        dfs( adj, k, 1 );
        num[i] = best;
    return best;
```

#### 5.6 MinimumSteinerTree

```
// Minimum Steiner Tree
// O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
   int n , dst[V][V] , dp[1 << T][V] , tdst[V];</pre>
```

```
void init( int n ) {
    for( int i = 0 ; i < n ; i ++ ) {</pre>
      for( int j = 0 ; j < n ; j ++ )</pre>
        dst[ i ][ j ] = INF;
      dst[i][i] = 0;
  void add edge( int ui , int vi , int wi ){
    dst[ui][vi] = min(dst[ui][vi], wi);
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  void shortest_path(){
    for( int k = 0 ; k < n ; k ++ )
      for( int i = 0 ; i < n ; i ++ )</pre>
        for( int j = 0 ; j < n ; j ++ )</pre>
          dst[ i ][ j ] = min( dst[ i ][ j ],
                dst[ i ][ k ] + dst[ k ][ j ] );
  int solve( const vector<int>& ter ){
    int t = (int)ter.size();
    for( int i = 0 ; i < ( 1 << t ) ; i ++ )</pre>
      for( int j = 0 ; j < n ; j ++ )</pre>
        dp[i][j] = INF;
    for( int i = 0 ; i < n ; i ++ )</pre>
      dp[0][i] = 0;
    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
      if (msk == (msk & (-msk))) 
        int who = __lg( msk );
        for( int i = 0 ; i < n ; i ++ )</pre>
         dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
        continue;
      for( int i = 0 ; i < n ; i ++ )</pre>
        for( int submsk = ( msk - 1 ) & msk ; submsk ;
                 submsk = (submsk - 1) \& msk)
            dp[ msk ][ i ] = min( dp[ msk ][ i ],
                             dp[submsk][i] +
                             dp[ msk ^ submsk ][ i ] );
      for( int i = 0 ; i < n ; i ++ ){</pre>
        tdst[i] = INF;
        for( int j = 0 ; j < n ; j ++ )</pre>
          tdst[i] = min(tdst[i],
                     dp[ msk ][ j ] + dst[ j ][ i ] );
      for( int i = 0 ; i < n ; i ++ )</pre>
        dp[ msk ][ i ] = tdst[ i ];
    int ans = INF;
    for( int i = 0 ; i < n ; i ++ )</pre>
      ans = min(ans, dp[(1 << t) - 1][i]);
    return ans;
} solver;
```

#### 5.7 Min Mean Cycle

```
// from BCW
/* minimum mean cycle */
const int MAXE = 1805;
const int MAXN = 35;
const double inf = 1029384756;
const double eps = 1e-6;
struct Edge {
 int v,u;
  double c;
};
int n,m,prv[MAXN][MAXN], prve[MAXN][MAXN], vst[MAXN];
Edge e[MAXE];
vector<int> edgeID, cycle, rho;
double d[MAXN][MAXN];
inline void bellman_ford() {
  for (int i=0; i<n; i++) d[0][i]=0;</pre>
  for(int i=0; i<n; i++) {</pre>
    fill(d[i+1], d[i+1]+n, inf);
```

```
for(int j=0; j<m; j++) {</pre>
      int v = e[j].v, u = e[j].u;
      if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
        d[i+1][u] = d[i][v]+e[j].c;
        prv[i+1][u] = v;
        prve[i+1][u] = j;
   }
 }
double karp mmc() {
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
  int st = -1;
  bellman ford();
  for (int i=0; i<n; i++) {</pre>
    double avg=-inf;
    for(int k=0; k<n; k++) {</pre>
      if(d[n][i] \le inf-eps) avg=max(avg,(d[n][i]-d[k][i])
      else avg=max(avg,inf);
    if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  for(int i=0; i<n; i++) vst[i] = 0;</pre>
  edgeID.clear(); cycle.clear(); rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
    vst[st]++;
    edgeID.PB(prve[i][st]);
    rho.PB(st);
  while (vst[st] != 2) {
    int v = rho.back(); rho.pop back();
    cycle.PB(v);
    vst[v]++;
  reverse (ALL (edgeID));
  edgeID.resize(SZ(cycle));
  return mmc;
```

### 5.8 Tarjan

```
割點
點 u 為割點 if and only if 滿足 1. or 2.
1. u 爲樹根,且 u 有多於一個子樹。
2. u 不爲樹根,且滿足存在 (u, v) 爲樹枝邊 (或稱父子邊,
    即 u 爲 v 在搜索樹中的父親),使得 DFN(u) <= Low(v)
一條無向邊 (u,v) 是橋 if and only if (u,v) 爲樹枝邊,且
    滿足 DFN(u) < Low(v)。
// 0 base
struct TarjanSCC{
 static const int MAXN = 1000006;
  int n, dfn[MAXN], low[MAXN], scc[MAXN], scn, count;
 vector<int> G[MAXN];
 stack<int> stk;
 bool ins[MAXN];
 void tarjan(int u) {
   dfn[u] = low[u] = ++count;
   stk.push(u);
   ins[u] = true;
   for(auto v:G[u]){
     if(!dfn[v]){
       tarjan(v);
       low[u] = min(low[u], low[v]);
     }else if(ins[v]){
       low[u] = min(low[u], dfn[v]);
```

```
if(dfn[u] == low[u]){
      int v:
      do {
      v = stk.top();
      stk.pop();
      scc[v] = scn;
      ins[v] = false;
      } while (v != u);
      scn++;
    }
  }
  void getSCC() {
   memset(dfn,0,sizeof(dfn));
    memset(low, 0, sizeof(low));
    memset(ins,0,sizeof(ins));
   memset(scc, 0, sizeof(scc));
    count = scn = 0;
    for(int i = 0 ; i < n ; i++ ){</pre>
      if(!dfn[i]) tarjan(i);
  }
}SCC:
```

#### 5.9 TwoSAT

```
const int MAXN = 2020;
struct TwoSAT{
   static const int MAXv = 2*MAXN;
    vector<int> GO[MAXv], BK[MAXv], stk;
   bool vis[MAXv];
    int SC[MAXv];
    void imply(int u,int v){ // u imply v
        GO[u].push_back(v);
        BK[v].push back(u);
    int dfs(int u,vector<int>*G,int sc){
        vis[u]=1, SC[u]=sc;
        for (int v:G[u])if (!vis[v])
            dfs(v,G,sc);
        if (G==GO) stk.push back(u);
    int scc(int n=MAXv) {
        memset(vis, 0, sizeof(vis));
        for (int i=0; i<n; i++)if (!vis[i])</pre>
            dfs(i,GO,-1);
        memset(vis, 0, sizeof(vis));
        int sc=0;
        while (!stk.empty()){
            if (!vis[stk.back()])
                dfs(stk.back(),BK,sc++);
            stk.pop_back();
        }
}SAT;
int main(){
    SAT.scc(2*n);
    bool ok=1;
    for (int i=0; i<n; i++) {</pre>
        if (SAT.SC[2*i] == SAT.SC[2*i+1]) ok=0;
    if (ok) {
        for (int i=0; i<n; i++) {</pre>
            if (SAT.SC[2*i]>SAT.SC[2*i+1]) {
                cout << i << endl;</pre>
        }
    else puts("NO");
```

## 6 Matching

#### 6.1 KM

```
#define MAXN 100
  #define INF INT MAX
  int g[MAXN][MAXN], lx[MAXN], ly[MAXN], slack y[MAXN];
  int px[MAXN],py[MAXN],match y[MAXN],par[MAXN];
  int n;
  void adjust (int y) { //把增廣路上所有邊反轉
    match y[y]=py[y];
    if(px[match y[y]]!=-2)
      adjust(px[match y[y]]);
  bool dfs(int x){//DFS找增廣路
    for (int y=0; y<n; ++y) {</pre>
      if (py[y]!=-1) continue;
      int t=lx[x]+ly[y]-g[x][y];
      if(t==0){
        py[y]=x;
        if (match y[y] ==-1) {
          adjust(y);
          return 1;
        if (px[match_y[y]]!=-1) continue;
        px[match y[y]]=y;
        if (dfs (match_y[y])) return 1;
      }else if(slack_y[y]>t){
        slack_y[y]=t;
        par[y]=x;
    return 0;
  inline int km(){
    memset(ly,0,sizeof(int)*n);
    memset (match_y, -1, sizeof (int) *n);
    for (int x=0; x<n; ++x) {</pre>
      lx[x] = -INF;
      for (int y=0; y<n; ++y) {</pre>
        lx[x]=max(lx[x],g[x][y]);
    for (int x=0; x<n; ++x) {</pre>
      for (int y=0; y<n; ++y) slack y[y]=INF;</pre>
      memset(px,-1,sizeof(int)*n);
      memset(py,-1,sizeof(int)*n);
      px[x] = -2;
      if(dfs(x))continue;
      bool flag=1;
      while(flag){
        int cut=INF;
        for (int y=0; y<n; ++y)</pre>
           if(py[y] ==-1&&cut>slack_y[y])cut=slack_y[y];
        for (int j=0; j<n; ++j) {</pre>
          if (px[j]!=-1)lx[j]-=cut;
           if (py[j]!=-1)ly[j]+=cut;
           else slack_y[j]-=cut;
        for (int y=0; y<n; ++y) {</pre>
          if(py[y] ==-1&&slack_y[y] ==0) {
             py[y]=par[y];
             if (match_y[y] ==-1) {
               adjust(y);
               flag=0;
               break;
             px[match_y[y]]=y;
             if (dfs(match y[y])) {
               flag=0;
               break;
            }
          }
        }
     }
10
```

```
int ans=0;
for(int y=0;y<n;++y)if(g[match_y[y]][y]!=-INF)ans+=g[
    match_y[y]][y];
return ans;
}</pre>
```

## 6.2 Maximum General Matching

```
// Maximum Cardinality Matching
struct Graph {
  vector<int> G[MAXN];
  int pa[MAXN], match[MAXN], st[MAXN], S[MAXN], vis[
  int t, n;
  void init(int n) {
    n = n;
    for ( int i = 1 ; i <= n ; i++ ) G[i].clear();</pre>
  void add edge(int u, int v) {
    G[u].push back(v);
    G[v].push back(u);
  int lca(int u, int v) {
    for ( ++t ; ; swap(u, v) ) {
      if ( u == 0 ) continue;
      if ( vis[u] == t ) return u;
      vis[u] = t;
      u = st[ pa[ match[u] ] ];
  void flower(int u, int v, int l, queue<int> &q) {
    while ( st[u] != l ) {
      pa[u] = v;
      if (S[v = match[u]] == 1) {
        q.push(v);
        S[v] = 0;
      st[u] = st[v] = 1;
      u = pa[v];
  bool bfs(int u) {
    for ( int i = 1 ; i <= n ; i++ ) st[i] = i;</pre>
    memset(S, -1, sizeof(S));
    queue<int>q;
    q.push(u);
    S[u] = 0;
    while ( !q.empty() ) {
      u = q.front(); q.pop();
      for ( int i = 0 ; i < (int)G[u].size(); i++) {</pre>
        int v = G[u][i];
        if (S[v] == -1) {
          pa[v] = u;
          S[v] = 1;
          if ( !match[v] ) {
            for ( int lst ; u ; v = lst, u = pa[v] ) {
              lst = match[u];
              match[u] = v;
              match[v] = u;
            return 1;
          q.push(match[v]);
          S[match[v]] = 0;
        } else if ( !S[v] && st[v] != st[u] ) {
          int l = lca(st[v], st[u]);
          flower(v, u, l, q);
          flower(u, v, l, q);
      }
    }
    return 0;
  int solve(){
```

```
memset(pa, 0, sizeof(pa));
memset(match, 0, sizeof(match));
int ans = 0;
for ( int i = 1 ; i <= n ; i++ )
    if ( !match[i] && bfs(i) ) ans++;
    return ans;
}
graph;</pre>
```

## 6.3 Minimum General Weighted Matching

```
// Minimum Weight Perfect Matching (Perfect Match)
struct Graph {
    static const int MAXN = 105;
    int n, e[MAXN][MAXN];
    int match[MAXN], d[MAXN], onstk[MAXN];
    vector<int> stk;
    void init(int n) {
        n = _n;
        for( int i = 0 ; i < n ; i ++ )</pre>
            for( int j = 0 ; j < n ; j ++ )</pre>
                e[i][j] = 0;
    void add_edge(int u, int v, int w) {
        e[u][v] = e[v][u] = w;
    bool SPFA(int u) {
        if (onstk[u]) return true;
        stk.push_back(u);
        onstk[u] = 1;
        for ( int v = 0 ; v < n ; v++ ) {</pre>
            if (u != v && match[u] != v && !onstk[v] )
                 int m = match[v];
                 if (d[m] > d[u] - e[v][m] + e[u][v])
                     d[m] = d[u] - e[v][m] + e[u][v];
                     onstk[v] = 1;
                     stk.push back(v);
                     if (SPFA(m)) return true;
                     stk.pop back();
                     onstk[v] = 0;
                 }
            }
        onstk[u] = 0;
        stk.pop back();
        return false;
    int solve() {
        for ( int i = 0 ; i < n ; i += 2 ) {</pre>
            match[i] = i+1;
            match[i+1] = i;
        while (true) {
            int found = 0;
            for ( int i = 0 ; i < n ; i++ )</pre>
                onstk[ i ] = d[ i ] = 0;
             for ( int i = 0 ; i < n ; i++ ) {</pre>
                stk.clear();
                 if ( !onstk[i] && SPFA(i) ) {
                     found = 1;
                     while ( stk.size() >= 2 ) {
                         int u = stk.back(); stk.
                             pop back();
                         int v = stk.back(); stk.
                             pop_back();
                         match[u] = v;
                         match[v] = u;
                     }
             if (!found) break;
        int ret = 0;
```

11

## 6.4 Stable Marriage

```
#define F(n) Fi(i, n)
#define Fi(i, n) Fl(i, 0, n)
#define Fl(i, l, n) for(int i = l ; i < n ; ++i)
#include <bits/stdc++.h>
using namespace std;
int D, quota[205], weight[205][5];
int S, scoretodep[12005][205], score[5];
int P, prefer[12005][85], iter[12005];
int ans[12005];
typedef pair<int, int> PII;
map<int, int> samescore[205];
typedef priority queue<PII, vector<PII>, greater<PII>>
    000;
000 pri[205];
void check(int d) {
 PII t = pri[d].top();
 if (pri[d].size() - samescore[d][t.first] + 1 <=</pre>
      quota[d]) return;
  while (pri[d].top().first == t.first) {
   v = pri[d].top().second;
    ans[v] = -1;
    --samescore[d][t.first];
    pri[d].pop();
void push(int s, int d) {
 if (pri[d].size() < quota[d]) {</pre>
   pri[d].push(PII(scoretodep[s][d], s));
    ans[s] = d;
    ++samescore[s][scoretodep[s][d]];
  } else if (scoretodep[s][d] >= pri[d].top().first) {
   pri[d].push(PII(scoretodep[s][d], s));
    ans[s] = d;
    ++samescore[s][scoretodep[s][d]];
    check(d);
 }
void f() {
 int over;
  while (true) {
    over = 1;
      if (ans[q] != -1 || iter[q] >= P) continue;
     push(q, prefer[q][iter[q]++]);
      over = 0;
    if (over) break;
 }
main() {
 ios::sync_with_stdio(false);
 cin.tie(NULL);
 int sadmit, stof, dexceed, dfew;
  while (cin >> D, D) { // Beware of the input format
      or judge may troll us.
    sadmit = stof = dexceed = dfew = 0;
    memset(iter, 0, sizeof(iter));
    memset(ans, 0, sizeof(ans));
    Fi (q, 205) {
     pri[q] = QQQ();
      samescore[q].clear();
    cin >> S >> P;
    Fi (q, D) {
      cin >> quota[q];
      Fi (w, 5) cin >> weight[q][w];
```

## 7 Math

## 7.1 Ax+by=gcd

```
pair<int,int> extgcd(int a, int b) {
   if (b==0) return {1,0};
   int k = a/b;
   pair<int,int> p = extgcd(b,a-k*b);
   return { p.second, p.first - k*p.second };
}
```

#### 7.2 FFT

```
// use llround() to avoid EPS
typedef double Double;
const Double PI = acos(-1);
// STL complex may TLE
typedef complex<Double> Complex;
#define x real()
#define y imag()
template<typename Iter> // Complex*
void BitReverse(Iter a, int n) {
    for (int i=1, j=0; i<n; i++) {</pre>
        for (int k = n>>1; k>(j^=k); k>>=1);
        if (i<j) swap(a[i],a[j]);</pre>
    }
template<typename Iter> // Complex*
void FFT(Iter a, int n, int rev=1) { // rev = 1 or -1
    assert( (n&(-n)) == n); // n is power of 2
    BitReverse(a,n);
    Iter A = a:
    for (int s=1; (1<<s) <=n; s++) {</pre>
        int m = (1 << s);
        Complex wm(cos(2*PI*rev/m), sin(2*PI*rev/m));
        for (int k=0; k<n; k+=m) {</pre>
            Complex w(1,0);
            for (int j=0; j<(m>>1); j++) {
                Complex t = w * A[k+j+(m>>1)];
                Complex u = A[k+j];
                A[k+j] = u+t;
                A[k+j+(m>>1)] = u-t;
                w = w*wm;
            }
```

```
}

if (rev==-1) {
    for (int i=0; i<n; i++) {
        A[i] /= n;
    }
}
</pre>
```

#### 7.3 GaussElimination

```
// by bcw codebook
const int MAXN = 300;
const double EPS = 1e-8;
int n:
double A[MAXN][MAXN];
void Gauss() {
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
    for(int j = i; j < n; j++) {</pre>
      if(fabs(A[j][i]) > EPS) {
        swap(A[j], A[i]);
        ok = 1;
        break;
      }
    if(!ok) continue;
    double fs = A[i][i];
    for(int j = i+1; j < n; j++) {</pre>
      double r = A[j][i] / fs;
      for(int k = i; k < n; k++) {</pre>
        A[j][k] -= A[i][k] * r;
    }
  }
```

### 7.4 Karatsuba

```
// N is power of 2
template<typename Iter>
void DC(int N, Iter tmp, Iter A, Iter B, Iter res) {
    fill(res,res+2*N,0);
    if (N<=32) {
        for (int i=0; i<N; i++) {</pre>
            for (int j=0; j<N; j++) {</pre>
                 res[i+j] += A[i]*B[j];
        }
        return;
    int n = N/2;
    auto a = A+n, b = A;
    auto c = B+n, d = B;
    DC(n, tmp+N, a, c, res+2*N);
    for (int i=0; i<N; i++) {</pre>
        res[i+N] += res[2*N+i];
        res[i+n] -= res[2*N+i];
    DC(n,tmp+N,b,d,res+2*N);
    for (int i=0; i<N; i++) {</pre>
        res[i] += res[2*N+i];
        res[i+n] -= res[2*N+i];
    auto x = tmp;
    auto y = tmp+n;
    for (int i=0; i<n; i++) x[i] = a[i]+b[i];</pre>
```

```
for (int i=0; i<n; i++) y[i] = c[i]+d[i];
  DC(n,tmp+N,x,y,res+2*N);
  for (int i=0; i<N; i++) {
     res[i+n] += res[2*N+i];
  }
}
// DC(1<<16,tmp.begin(),A.begin(),B.begin(),res.begin()
);</pre>
```

#### 7.5 LinearPrime

```
const int MAXP = 100; //max prime
vector<int> P; // primes
void build_prime() {
    static bitset<MAXP> ok;
    int np=0;
    for (int i=2; i<MAXP; i++) {
        if (ok[i]==0) P. push_back(i), np++;
        for (int j=0; j<np && i*P[j]<MAXP; j++) {
            ok[ i*P[j] ] = 1;
            if (i%P[j]==0) break;
        }
    }
}</pre>
```

#### 7.6 Miller-Rabin

```
typedef long long LL;
inline LL bin mul(LL a, LL n,const LL& MOD) {
 LL re=0;
 while (n>0) {
   if (n&1) re += a;
    a += a; if (a>=MOD) a-=MOD;
   n >> = 1;
 return re%MOD;
inline LL bin_pow(LL a, LL n,const LL& MOD) {
 LL re=1;
 while (n>0) {
   if (n&1) re = bin mul(re,a,MOD);
    a = bin mul(a,a,MOD);
   n >> = 1:
 return re;
bool is prime(LL n) {
  //static LL sprp[3] = { 2LL, 7LL, 61LL};
  static LL sprp[7] = { 2LL, 325LL, 9375LL,
    28178LL, 450775LL, 9780504LL,
    1795265022LL };
 if (n==1 || (n&1) ==0 ) return n==2;
  int u=n-1, t=0;
 while ( (u&1) ==0 ) u>>=1, t++;
  for (int i=0; i<3; i++) {</pre>
   LL x = bin_pow(sprp[i]%n, u, n);
    if (x==0 | x==1 | x==n-1) continue;
    for (int j=1; j<t; j++) {</pre>
     x=x*x%n;
      if (x==1 || x==n-1)break;
    if (x==n-1)continue;
   return 0;
  return 1;
```

#### 7.7 Mobius

#### 7.8 PollardRho

```
// does not work when n is prime
inline LL f(LL x , LL mod) {
  return (x * x % mod + 1) % mod;
}
inline LL pollard_rho(LL n) {
  if(!(n&1)) return 2;
  while(true) {
    LL y = 2 , x = rand() % (n - 1) + 1 , res = 1;
    for(int sz = 2; res == 1; sz *= 2) {
      for(int i = 0; i < sz && res <= 1; i++) {
         x = f(x , n);
         res = __gcd(abs(x - y) , n);
      }
      y = x;
    }
  if (res != 0 && res != n) return res;
}</pre>
```

### 7.9 Simplex

```
// Two-phase simplex algorithm for solving linear
   programs of the form
      maximize
                  C^T X
11
      subject\ to\ Ax <= b
                   x >= 0
11
//
// INPUT: A -- an m x n matrix
        b -- an m-dimensional vector
11
         c -- an n-dimensional vector
         x -- a vector where the optimal solution will
     be stored
// OUTPUT: value of the optimal solution (infinity if
   unbounded
          above, nan if infeasible)
11
// To use this code, create an LPSolver object with A,
   b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
```

```
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
 VVD D;
 LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
       VD(n + 2)) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n;</pre>
       j++) D[i][j] = A[i][j];
    for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]</pre>
        = -1; D[i][n + 1] = b[i]; 
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -</pre>
        c[j]; }
   N[n] = -1; D[m + 1][n] = 1;
 void Pivot(int r, int s) {
   double inv = 1.0 / D[r][s];
    for (int i = 0; i < m + 2; i++) if (i != r)</pre>
      for (int j = 0; j < n + 2; j++) if (j != s)</pre>
       D[i][j] = D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]</pre>
        *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]</pre>
         *=-inv;
   D[r][s] = inv;
   swap(B[r], N[s]);
 bool Simplex(int phase) {
   int x = phase == 1 ? m + 1 : m;
   while (true) {
      int s = -1;
      for (int j = 0; j <= n; j++) {</pre>
       if (phase == 2 && N[j] == -1) continue;
        if (s == -1 || D[x][j] < D[x][s] || D[x][j] ==
            D[x][s] && N[j] < N[s]) s = j;
      if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n +
             1] / D[r][s] ||
          (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r]
              [s] && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  DOUBLE Solve(VD &x) {
   int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][</pre>
        n + 1) r = i;
   if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return</pre>
           -numeric limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {</pre>
       int s = -1;
        for (int j = 0; j <= n; j++)</pre>
         if (s == -1 || D[i][j] < D[i][s] || D[i][j]</pre>
              == D[i][s] \&\& N[j] < N[s]) s = j;
        Pivot(i, s);
     }
   if (!Simplex(2)) return numeric limits<DOUBLE>::
       infinity();
   x = VD(n);
```

```
for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = | state = []</pre>
        D[i][n + 1];
    return D[m][n + 1];
 }
};
int main() {
 const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] = {
   \{ 6, -1, 0 \},
   \{-1, -5, 0\},
   { 1, 5, 1 },
   \{-1, -5, -1\}
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE c[n] = \{ 1, -1, 0 \};
 VVD A(m);
 VD b(_b, _b + m);
  VD c(c, c+n);
  for (int i = 0; i < m; i++) A[i] = VD( A[i], A[i] +</pre>
 LPSolver solver(A, b, c);
 DOUBLE value = solver.Solve(x);
 cerr << "VALUE: " << value << endl; // VALUE: 1.29032
 cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size t i = 0; i < x.size(); i++) cerr << " " <<</pre>
    x[i];
 cerr << endl;
 return 0;
```

#### 7.10 Sprague-Grundy

```
Anti Nim (取走最後一個石子者敗)
先手必勝 if and only if
1. 「所有」堆的石子數都為 1 且遊戲的 SG 值為 0。
2. 「有些」堆的石子數大於 1 且遊戲的 SG 值不為 0。
Anti-SG (決策集合為空的遊戲者贏)
定義 SG 值為 0 時,遊戲結束,
則先手必勝 if and only if
1. 遊戲中沒有單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
2. 遊戲中某個單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
  不為 0。
______
Sprague-Grundy
1. 雙人、回合制
2. 資訊完全公開
3. 無隨機因素
4. 可在有限步內結束
5. 沒有和局
6. 雙方可採取的行動相同
SG(S) 的值為 0:後手(P)必勝
不為 0:先手(N)必勝
int mex(set S) {
 // find the min number >= 0 that not in the S
 // e.g. S = \{0, 1, 3, 4\} \max(S) = 2
```

```
int SG(A) {
 if (A not in state) {
   S = sub states(A)
   if( len(S) > 1 ) state[A] = reduce(operator.xor, [
       SG(B) for B in S])
   else state[A] = mex(set(SG(B) for B in next states(
 return state[A]
```

#### 7.11 Theorem

Lagrange multiplier

f(x,y) 求極值。必須滿足 g(x,y) = 0。

```
Lucas's Theorem
 For non-negative integer n,m and prime P,
 C(m,n) \mod P = C(m/M,n/M) * C(m%M,n%M) \mod P
 = mult i ( C(m i, n i) )
where m i is the i-th digit of m in base P.
Kirchhoff's theorem
 A_{\{ii\}} = deg(i), A_{\{ij\}} = (i,j) \setminus in E ? -1 : 0
 Deleting any one row, one column, and cal the det(A)
Nth Catalan recursive function:
C \ 0 = 1, \ C \ \{n+1\} = C \ n \ * \ 2(2n + 1)/(n+2)
______
Mobius Formula
u(n) = 1 , if n = 1
      (-1)^m , 若 n 無平方數因數,且 n = p1*p2*p3
          *...*pk
              ,若 n 有大於 1 的平方數因數
- Property
1. (積性函數) u(a)u(b) = u(ab)
2. \sum_{\{d|n\}} u(d) = [n == 1]
_____
Mobius Inversion Formula
if f(n) = \sum_{i=1}^{n} \{d \mid n\} g(d)
      g(n) = \sum_{i=1}^{n} \{d \mid n\} \ u(n/d) f(d)= \sum_{i=1}^{n} \{d \mid n\} \ u(d) f(n/d)
- Application
the number/power of gcd(i, j) = k
- Trick
分塊, O(sqrt(n))
Chinese Remainder Theorem (m i 兩兩互質)
 x = a 1 \pmod{m} 1
 x = a_2 \pmod{m} 2
 x = a i \pmod{m} i
construct a solution:
  Let M = m_1 * m_2 * m_3 * ... * m_n
 Let M_i = M / m^i
  t i = 1 / M i
 t i * M i = 1 \pmod{m i}
 solution x = a_1 * t_1 * M_1 + a_2 * t_2 * M_2 + ...
   + a n * t n * M n + k * M
  = k*M + \sum a i * t i * M i, k is positive integer.
 under mod M, there is one solution x = \sum a i * t i *
  M_{\_}i
Burnside's lemma
|G| * |X/G| = sum(|X^q|) where q in G
總方法數:每一種旋轉下不動點的個數總和 除以 旋轉的方法
```

```
凑得 f(x,y) = f(x,y) + \lambda g(x,y)
定義 s(x,y,\lambda) = f(x,y) + \lambda g(x,y)
f(x,y) 的極值,等同 s(x,y,\lambda) = f(x,y) + \lambda g(x,y) 的極值。
欲求極值:
對 x 偏微分,讓斜率是 0 \circ
對 y 偏微分,讓斜率是 0 \circ
不管 \lambda 如何變化,\lambda g(x,y) 都是零,s(x,y,\lambda) 永遠不變。
欲求永遠不變的地方:
對 \lambda 偏微分,讓斜率是 0 \circ
三道偏微分方程式聯立之後,其解涵蓋了(不全是)所有符合約束條件的極值。
{ \partial/\partial x \ s(x,y,\lambda) = 0}
{ \partial/\partial y \ s(x,y,\lambda) = 0}
{ \partial/\partial \lambda \ s(x,y,\lambda) = 0}
```

## 8 Other

int sgn(double x)

## 8.1 Count Spanning Tree

```
新的方法介绍
下面我们介绍一种新的方法——Matrix-Tree定理(Kirchhoff矩
   阵-树定理)。
Matrix-Tree定理是解决生成树计数问题最有力的武器之一。它
   首先于1847年被Kirchhoff证明。在介绍定理之前,我们首
   先明确几个概念:
1、G的度数矩阵D[G]是一个n*n的矩阵,并且满足:当i≠j时,
   dij=0;当i=j时,dij等于vi的度数。
2、G的邻接矩阵A[G]也是一个n*n的矩阵, 并且满足:如果vi
   、vi之间有边直接相连,则aij=1,否则为0。
我们定义G的Kirchhoff矩阵(也称为拉普拉斯算子)C[G]为C[G]=
   D[G]-A[G] ,
则Matrix-Tree定理可以描述为:G的所有不同的生成树的个数
   等于其Kirchhoff矩阵C[G]任何一个n-1阶主子式的行列式
   的绝对值。
所谓n-1阶主子式,就是对于r(1≤r≤n),将C[G]的第r行、第r列
   同时去掉后得到的新矩阵,用Cr[G]表示。
生成树计数
算法步骤:
1、 构建拉普拉斯矩阵
   Matrix[i][j] =
degree(i) , i==j
       -1,i-j有边
       ○,其他情况
2、 去掉第r行,第r列(r任意)
3、 计算矩阵的行列式
/* **************
MYID : Chen Fan
T<sub>i</sub>ANG
     : G++
     : Count Spaning Tree From Kuangbin
#include <stdio.h>
#include <string.h>
#include <algorithm>
#include <iostream>
#include <math.h>
using namespace std;
const double eps = 1e-8;
const int MAXN = 110;
```

```
if(fabs(x) < eps)return 0;</pre>
    if(x < 0) return -1;
    else return 1;
double b[MAXN][MAXN];
double det(double a[][MAXN],int n)
    int i, j, k, sign = 0;
    double ret = 1;
    for(i = 0;i < n;i++)</pre>
    for(j = 0; j < n; j++) b[i][j] = a[i][j];</pre>
    for(i = 0; i < n; i++)
        if(sgn(b[i][i]) == 0)
             for(j = i + 1; j < n; j++)</pre>
            if(sgn(b[j][i]) != 0) break;
            if(j == n)return 0;
             for(k = i; k < n; k++) swap(b[i][k], b[j][k]);
             sign++;
        ret *= b[i][i];
        for(k = i + 1; k < n; k++) b[i][k]/=b[i][i];
        for(j = i+1; j < n; j++)</pre>
        for(k = i+1;k < n;k++) b[j][k] -= b[j][i]*b[i][</pre>
            k1;
    if(sign & 1)ret = -ret;
    return ret;
double a[MAXN][MAXN];
int g[MAXN][MAXN];
int main()
    int T;
    int n,m;
    int u, v;
    scanf("%d",&T);
    while (T--)
        scanf("%d%d",&n,&m);
        memset(g,0,sizeof(g));
        while (m--)
             scanf("%d%d",&u,&v);
            u--;v--;
            g[u][v] = g[v][u] = 1;
        memset(a,0,sizeof(a));
        for(int i = 0;i < n;i++)</pre>
        for(int j = 0; j < n; j++)</pre>
        if(i != j && g[i][j])
             a[i][i]++;
            a[i][j] = -1;
        double ans = det(a,n-1);
        printf("%.01f\n",ans);
    return 0;
```

#### 8.2 CYK

```
// 2016 NCPC from sunmoon

// 轉換

#define MAXN 55
struct CNF{
   int s,x,y;//s->xy | s->x, if y==-1
   int cost;
   CNF() {}
   CNF(int s,int x,int y,int c):s(s),x(x),y(y),cost(c) {}
};

16
```

```
int state; //規則數量
map<char,int> rule;//每個字元對應到的規則,小寫字母為終
    端字符
vector<CNF> cnf;
inline void init(){
  state=0;
  rule.clear();
  cnf.clear();
inline void add to cnf(char s,const string &p,int cost)
  if (rule.find(s) == rule.end()) rule[s] = state++;
  for (auto c:p) if (rule.find(c) == rule.end()) rule[c] =
      state++;
  if(p.size()==1){
    cnf.push back(CNF(rule[s],rule[p[0]],-1,cost));
  }else{
    int left=rule[s];
    int sz=p.size();
    for (int i=0;i<sz-2;++i) {</pre>
      cnf.push back(CNF(left,rule[p[i]],state,0));
    cnf.push back(CNF(left, rule[p[sz-2]], rule[p[sz-1]],
        cost));
// 計算
vector<long long> dp[MAXN][MAXN];
vector<bool> neg_INF[MAXN][MAXN];//如果花費是負的可能會
    有無限小的情形
inline void relax(int 1,int r,const CNF &c,long long
    cost,bool neg c=0) {
  \textbf{if} (!neg\_INF[l][r][c.s] \& (neg\_INF[l][r][c.x]||cost < dp[
      l][r][c.s])){
    if (neg_c||neg_INF[l][r][c.x]){
      dp[1][r][c.s]=0;
      neg INF[1][r][c.s]=true;
    }else dp[l][r][c.s]=cost;
inline void bellman(int l,int r,int n) {
  for (int k=1; k<=state; ++k)</pre>
    for (auto c:cnf)
      if(c.y==-1)relax(l,r,c,dp[l][r][c.x]+c.cost,k==n)
inline void cyk(const vector<int> &tok){
  for (int i=0; i < (int) tok.size(); ++i) {</pre>
    for (int j=0; j < (int) tok.size(); ++j) {</pre>
      dp[i][j]=vector<long long>(state+1,INT MAX);
      neg INF[i][j]=vector<bool>(state+1, false);
    dp[i][i][tok[i]]=0;
    bellman(i,i,tok.size());
  for (int r=1; r < (int) tok.size(); ++r) {</pre>
    for (int l=r-1; l>=0; --1) {
      for (int k=1; k<r; ++k)</pre>
        for (auto c:cnf)
          if(~c.y)relax(l,r,c,dp[l][k][c.x]+dp[k+1][r][
               c.y]+c.cost);
      bellman(l,r,tok.size());
  }
```

### 8.3 DigitCounting

```
int dfs(int pos, int state1, int state2 ....., bool
    limit, bool zero) {
    if ( pos == -1 ) return 是否符合條件;
    int &ret = dp[pos][state1][state2][....];
```

```
if ( ret != -1 && !limit ) return ret;
int ans = 0;
int upper = limit ? digit[pos] : 9;
for ( int i = 0 ; i <= upper ; i++ ) {
    ans += dfs(pos - 1, new_state1, new_state2,
        limit & ( i == upper), ( i == 0) && zero);
}
if ( !limit ) ret = ans;
return ans;
}
int solve(int n) {
    int it = 0;
    for (; n ; n /= 10 ) digit[it++] = n % 10;
    return dfs(it - 1, 0, 0, 1, 1);
}</pre>
```

## 8.4 DP-optimization

```
Monotonicity & 1D/1D DP & 2D/1D DP
Definition xD/yD
1D/1D DP[j] = min(0 \le i < j) { DP[i] + w(i, j) }; DP[0] = k
2D/1D DP[i][j] = min(i < k \le j) \{ DP[i][k - 1] + DP[k][j] \}
    + w(i, j); DP[i][i] = 0
Monotonicity
               d
     C
a | w(a, c) w(a, d)
b | w(b, c) w(b, d)
Monge Condition
Concave(凹四邊形不等式): w(a, c) + w(b, d) >= w(a, d) +
Convex (凸四邊形不等式): w(a, c) + w(b, d) <= w(a, d) +
    w(b, c)
Totally Monotone
Concave(凹單調): w(a, c) <= w(b, d) ----> w(a, d) <= w
Convex (凸單調): w(a, c) >= w(b, d) -----> w(a, d) >= w
1D/1D DP O(n^2) -> O(nlgn)
**CONSIDER THE TRANSITION POINT**
Solve 1D/1D Concave by Stack
Solve 1D/1D Convex by Deque
2D/1D Convex DP (Totally Monotone) O(n^3) \rightarrow O(n^2)
h(i, j - 1) \le h(i, j) \le h(i + 1, j)
```

## 8.5 Dp1D1D

```
#include <bits/stdc++.h>
  int t, n, L;
  int p;
  char s[MAXN][35];
  11 \text{ sum}[MAXN] = \{0\};
  long double dp[MAXN] = {0};
  int prevd[MAXN] = {0};
  long double pw(long double a, int n) {
      if ( n == 1 ) return a;
      long double b = pw(a, n/2);
      if ( n & 1 ) return b*b*a;
      else return b*b;
 long double f(int i, int j) {
       cout \ll (sum[i] - sum[j]+i-j-1-L) \ll endl;
      return pw(abs(sum[i] - sum[j]+i-j-1-L), p) + dp[j];
17
```

```
struct INV {
   int L, R, pos;
INV stk[MAXN*10];
int top = 1, bot = 1;
void update(int i) {
    while (top > bot && i < stk[top].L && f(stk[top].L
        , i) < f(stk[top].L, stk[top].pos) ) {
        stk[top - 1].R = stk[top].R;
        top--;
    int lo = stk[top].L, hi = stk[top].R, mid, pos =
        stk[top].pos;
    //if ( i >= 10 ) 10 = i + 1;
    while ( lo != hi ) {
        mid = lo + (hi - lo) / 2;
        if (f(mid, i) < f(mid, pos) ) hi = mid;</pre>
        else lo = mid + 1;
    if ( hi < stk[top].R ) {
        stk[top + 1] = (INV) { hi, stk[top].R, i };
        stk[top++].R = hi;
    }
}
int main() {
   cin >> t;
    while ( t-- ) {
        cin >> n >> L >> p;
        dp[0] = sum[0] = 0;
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
            cin >> s[i];
            sum[i] = sum[i-1] + strlen(s[i]);
            dp[i] = numeric limits<long double>::max();
        stk[top] = (INV) \{1, n + 1, 0\};
        for ( int i = 1 ; i <= n ; i++ ) {
            if ( i >= stk[bot].R ) bot++;
            dp[i] = f(i, stk[bot].pos);
            update(i);
//
              cout << (ll) f(i, stk[bot].pos) << endl;</pre>
        if ( dp[n] > 1e18 ) {
            cout << "Too hard to arrange" << endl;</pre>
        } else {
            vector<PI> as;
            cout << (11)dp[n] << endl;</pre>
    return 0:
```

### 8.6 ManhattanMST

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 100005;
const int OFFSET = 2000; // y-x may < 0, offset it, if</pre>
   y-x too large, please write a unique function
const int INF = 0xFFFFFFF;
int n:
int x[MAXN], y[MAXN], p[MAXN];
typedef pair<int, int> pii;
pii bit[MAXN]; // [ val, pos ]
struct P {
   int x, y, id;
    bool operator<(const P&b ) const {</pre>
        if (x == b.x) return y > b.y;
        else return x > b.x;
    }
};
vector<P> op;
```

```
struct E {
    int x, y, cost;
    bool operator<(const E&b ) const {</pre>
        return cost < b.cost;</pre>
};
vector<E> edges;
int find(int x) {
    return p[x] == x ? x : p[x] = find(p[x]);
void update(int i, int v, int p) {
    while ( i ) {
        if ( bit[i].first > v ) bit[i] = {v, p};
        i -= i & (-i);
pii query(int i) {
    pii res = {INF, INF};
    while ( i < MAXN ) {</pre>
        if ( bit[i].first < res.first ) res = {bit[i].</pre>
            first, bit[i].second);
        i += i & (-i);
    return res;
}
void input() {
    cin >> n;
    for ( int i = 0 ; i < n ; i++ ) cin >> x[i] >> y[i
        ], op.push back((P) \{x[i], y[i], i\});
void mst() {
    for ( int i = 0 ; i < MAXN ; i++ ) p[i] = i;</pre>
    int res = 0;
    sort(edges.begin(), edges.end());
    for ( auto e : edges ) {
        int x = find(e.x), y = find(e.y);
        if ( x != y ) {
            p[x] = y;
            res += e.cost;
    }
    cout << res << endl;
void construct() {
    sort(op.begin(), op.end());
    for ( int i = 0 ; i < n ; i++ ) {</pre>
        pii q = query(op[i].y - op[i].x + OFFSET);
        update(op[i].y - op[i].x + OFFSET, op[i].x + op
            [i].y, op[i].id);
        if ( q.first == INF ) continue;
        edges.push_back((E) {op[i].id, q.second, abs(x[
             op[i].id]-x[q.second]) + abs(y[op[i].id]-y[
             q.second]) });
void solve() {
    // [45 ~ 90 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,</pre>
        TNF}:
    construct();
    // [0 ~ 45 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,</pre>
        TNF}:
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i</pre>
        ].y);
    construct();
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i</pre>
```

```
// [-90 ~ -45 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,</pre>
        TNF };
    for ( int i = 0 ; i < n ; i++ ) op[i].y *= -1;</pre>
    construct();
    // [-45 ~ 0 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,</pre>
        INF };
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i</pre>
       1.v);
    construct();
    // mst
   mst();
int main () {
   input();
    solve();
    return 0;
```

## 9 String

#### 9.1 AC

```
// remember make fail() !!!
// notice MLE
const int sigma = 62;
const int MAXC = 200005;
inline int idx(char c){
   if ('A'<= c && c <= 'Z')return c-'A';</pre>
    if ('a'<= c && c <= 'z') return c-'a' + 26;
    if ('0'<= c && c <= '9')return c-'0' + 52;</pre>
struct ACautomaton{
    struct Node{
        Node *next[sigma], *fail;
        int cnt; // dp
        Node(){
           memset(next, 0, sizeof(next));
            fail=0;
            cnt=0:
    } buf[MAXC], *bufp, *ori, *root;
    void init() {
       bufp = buf;
        ori = new (bufp++) Node();
        root = new (bufp++) Node();
    void insert(int n, char *s) {
        Node *ptr = root;
        for (int i=0; s[i]; i++){
            int c = idx(s[i]);
            if (ptr->next[c]==NULL)
               ptr->next[c] = new (bufp++) Node();
            ptr = ptr->next[c];
        ptr->cnt=1;
    Node* trans(Node *o, int c) {
        while (o->next[c]==NULL) o = o->fail;
        return o->next[c];
    }
```

```
void make_fail() {
    static queue<Node*> que;

    for (int i=0; i<sigma; i++)
        ori->next[i] = root;
    root->fail = ori;

    que.push(root);
    while ( que.size() ) {
        Node *u = que.front(); que.pop();
        for (int i=0; i<sigma; i++) {
            if (u->next[i]==NULL) continue;
            u->next[i]->fail = trans(u->fail,i);
            que.push(u->next[i]);
        }
        u->cnt += u->fail->cnt;
    }
}
ac;
```

#### 9.2 BWT

```
// BWT
                           // 字串長度
const int N = 8;
int s[N+N+1] = "suffixes"; // 字串,後面預留一倍空間。
                            // 後綴陣列
int sa[N];
int pivot;
int cmp(const void* i, const void* j)
    return strncmp(s+*(int*)i, s+*(int*)j, N);
// 此處便宜行事,採用 O(N² logN) 的後綴陣列演算法。
void BWT()
    strncpy(s + N, s, N);
    for (int i=0; i<N; ++i) sa[i] = i;</pre>
    qsort(sa, N, sizeof(int), cmp);
    // 當輸入字串的所有字元都相同,必須當作特例處理。
    // 或者改用stable sort。
    for (int i=0; i<N; ++i)</pre>
        cout << s[(sa[i] + N-1) % N];
    for (int i=0; i<N; ++i)</pre>
        if (sa[i] == 0)
            pivot = i;
            break;
// Inverse BWT
const int N = 8;
                           // 字串長度
char t[N+1] = "xuffessi";
                           // 字串
int pivot;
int next[N];
void IBWT()
    vector<int> index[256];
    for (int i=0; i<N; ++i)</pre>
        index[t[i]].push back(i);
    for (int i=0, n=0; i<256; ++i)</pre>
        for (int j=0; j<index[i].size(); ++j)</pre>
           next[n++] = index[i][j];
    int p = pivot;
    for (int i=0; i<N; ++i)</pre>
        cout << t[p = next[p]];
```

#### 9.3 KMP

```
template<typename T>
void build KMP(int n, T *s, int *f) { // 1 base
 f[0]=-1, f[1]=0;
  for (int i=2; i<=n; i++) {</pre>
   int w = f[i-1];
    while (w \ge 0 \&\& s[w+1]! = s[i])w = f[w];
    f[i]=w+1;
 }
template<typename T>
int KMP(int n, T *a, int m, T *b) {
 build KMP(m,b,f);
 int ans=0;
  for (int i=1, w=0; i<=n; i++) {</pre>
    while ( w>=0 && b[w+1]!=a[i] )w = f[w];
    if (w==m) {
     ans++;
      w=f[w];
  return ans;
```

### 9.4 PalindromicTree

```
// remember init()
// remember make fail() !!!
// insert s need 1 base !!!
// notice MLE
const int sigma = 62;
const int MAXC = 1000006;
inline int idx(char c){
   if ('a'<= c && c <= 'z')return c-'a';</pre>
    if ('A'<= c && c <= 'Z') return c-'A'+26;</pre>
    if ('0'<= c && c <= '9')return c-'0'+52;</pre>
struct PalindromicTree{
    struct Node{
        Node *next[sigma], *fail;
        int len, cnt; // for dp
        Node(){
            memset(next, 0, sizeof(next));
            fail=0:
            len = cnt = 0;
    } buf[MAXC], *bufp, *even, *odd;
    void init(){
        bufp = buf;
        even = new (bufp++) Node();
        odd = new (bufp++) Node();
        even->fail = odd;
        odd \rightarrow len = -1;
    void insert(char *s) {
        Node* ptr = even;
        for (int i=1; s[i]; i++) {
            ptr = extend(ptr,s+i);
    Node* extend(Node *o, char *ptr) {
        int c = idx(*ptr);
        while ( *ptr != *(ptr-1-o->len) )o=o->fail;
        Node *&np = o->next[c];
        if (!np) {
            np = new (bufp++) Node();
            np->len = o->len+2;
            Node *f = o->fail;
```

#### 9.5 SAM

```
// par : fail link
// val : a topological order ( useful for DP )
// go[x] : automata edge ( x is integer in [0,26) )
struct SAM{
  struct State{
    int par, go[26], val;
    State () : par(0), val(0) { FZ(go); }
    State (int val) : par(0), val( val) { FZ(go); }
  vector<State> vec;
  int root, tail;
  void init(int arr[], int len) {
    vec.resize(2);
    vec[0] = vec[1] = State(0);
    root = tail = 1;
    for (int i=0; i<len; i++)</pre>
      extend(arr[i]);
  void extend(int w) {
    int p = tail, np = vec.size();
    vec.PB(State(vec[p].val+1));
    for ( ; p && vec[p].go[w] == 0; p = vec[p].par)
     vec[p].go[w] = np;
    if (p == 0) {
     vec[np].par = root;
    } else {
      if (vec[vec[p].go[w]].val == vec[p].val+1) {
       vec[np].par = vec[p].go[w];
      } else {
        int q = vec[p].go[w], r = vec.size();
        vec.PB(vec[q]);
        vec[r].val = vec[p].val+1;
        vec[q].par = vec[np].par = r;
        for ( ; p && vec[p].go[w] == q; p=vec[p].par)
          vec[p].go[w] = r;
    tail = np;
  }
};
```

#### 9.6 Smallest Rotation

```
string mcp(string s) {
   int n = s.length();
   s += s;
   int i=0, j=1;
   while (i<n && j<n) {
      int k = 0;
      while (k < n && s[i+k] == s[j+k]) k++;
      if (s[i+k] <= s[j+k]) j += k+1;
      else i += k+1;
      if (i == j) j++;
   }
}</pre>
```

```
int ans = i < n ? i : j;
return s.substr(ans, n);
}</pre>
```

## 9.7 Suffix Array

```
/*he[i]保存了在後綴數組中相鄰兩個後綴的最長公共前綴長度
 *sa[i]表示的是字典序排名為i的後綴是誰(字典序越小的排
     名越靠前)
 *rk[i]表示的是後綴我所對應的排名是多少 */
const int MAX = 1020304;
int ct[MAX], he[MAX], rk[MAX];
int sa[MAX], tsa[MAX], tp[MAX][2];
void suffix array(char *ip) {
 int len = strlen(ip);
  int alp = 256;
  memset(ct, 0, sizeof(ct));
  for (int i=0;i<len;i++) ct[ip[i]+1]++;</pre>
  for(int i=1;i<alp;i++) ct[i]+=ct[i-1];</pre>
  for(int i=0;i<len;i++) rk[i]=ct[ip[i]];</pre>
  for (int i=1;i<len;i*=2) {</pre>
    for (int j=0; j<len; j++) {</pre>
      if(j+i>=len) tp[j][1]=0;
      else tp[j][1]=rk[j+i]+1;
      tp[j][0]=rk[j];
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][1]+1]++;</pre>
    for (int j=1; j<len+2; j++) ct[j]+=ct[j-1];</pre>
    for (int j=0; j<len; j++) tsa[ct[tp[j][1]]++]=j;</pre>
    memset(ct, 0, sizeof(ct));
    for (int j=0;j<len;j++) ct[tp[j][0]+1]++;</pre>
    for (int j=1; j<len+1; j++) ct[j]+=ct[j-1];</pre>
    for (int j=0; j<len; j++)</pre>
     sa[ct[tp[tsa[j]][0]]++]=tsa[j];
    rk[sa[0]]=0;
    for (int j=1; j<len; j++) {</pre>
      if( tp[sa[j]][0] == tp[sa[j-1]][0] &&
        tp[sa[j]][1] == tp[sa[j-1]][1])
        rk[sa[j]] = rk[sa[j-1]];
      else
        rk[sa[j]] = j;
  for (int i=0, h=0; i<len; i++) {</pre>
    if (rk[i] == 0) h= 0;
    else{
      int j=sa[rk[i]-1];
      h=max(0,h-1);
      for (; ip[i+h] == ip[j+h]; h++);
    he[rk[i]]=h;
```

#### 9.8 Z-value

```
| z[0] = 0;
| for ( int bst = 0, i = 1; i < len; i++ ) {
| if ( z[bst] + bst <= i ) z[i] = 0;
| else z[i] = min(z[i - bst], z[bst] + bst - i);
| while ( str[i + z[i]] == str[z[i]] ) z[i]++;
| if ( i + z[i] > bst + z[bst] ) bst = i;
| }
| // 回文版
| void Zpal(const char *s, int len, int *z) {
| // Only odd palindrome len is considered | // z[i] means that the longest odd palindrom centered at
```

```
// i is [i-z[i] .. i+z[i]]
z[0] = 0;
for (int b=0, i=1; i<len; i++) {
   if (z[b] + b >= i) z[i] = min(z[2*b-i], b+z[b]-
        i);
   else z[i] = 0;
   while (i+z[i]+1 < len and i-z[i]-1 >= 0 and
        s[i+z[i]+1] == s[i-z[i]-1]) z[i] ++;
   if (z[i] + i > z[b] + b) b = i;
}
```

# 10 無權邊的生成樹個數 Kirchhoff's Theorem

1. 定義  $n \times m$  矩陣  $E = (a_{i,j})$ ,n 為點數,m 為邊數,若 i 點在 j 邊上,i 為小點  $a_{i,j} = 1$ ,i 為大點  $a_{i,j} = -1$ ,否則  $a_{i,j} = 0$ 。 (證明省略)

4. 令  $E(E^T) = Q$ ,他是一種有負號的 kirchhoff 的矩陣,取 Q 的子矩陣即為  $F(F^T)$ 

結論:做 Q 取子矩陣算  $\det$  即為所求。(除去第一行第一列 by mz)

## 11 monge

```
i \le i' < j \le j'

m(i, j) + m(i', j') \le m(i', j) + m(i, j')

k(i, j - 1) <= k(i, j) <= k(i + 1, j)
```

## 12 四心

```
\frac{sa*A+sb*B+sc*C}{sa+sb+sc}
外心 \sin 2A : \sin 2B : \sin 2C
內心 \sin A : \sin B : \sin C
垂心 \tan A : \tan B : \tan C
```

# 13 Runge-Kutta

```
y_{n+1} = y_n + \frac{h}{6}(k_1 + 2k_2 + 2k_3 + k_4)
k_1 = f(t_n, y_n)
k_2 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_2)
k_3 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_3)
k_2 = f(t_n + h, y_n + hk_3)
```

## 14 Householder Matrix

$$I - 2 \frac{vv^T}{v^T v}$$