

Internet of Things Protocols - 6LoWPAN and CoAP

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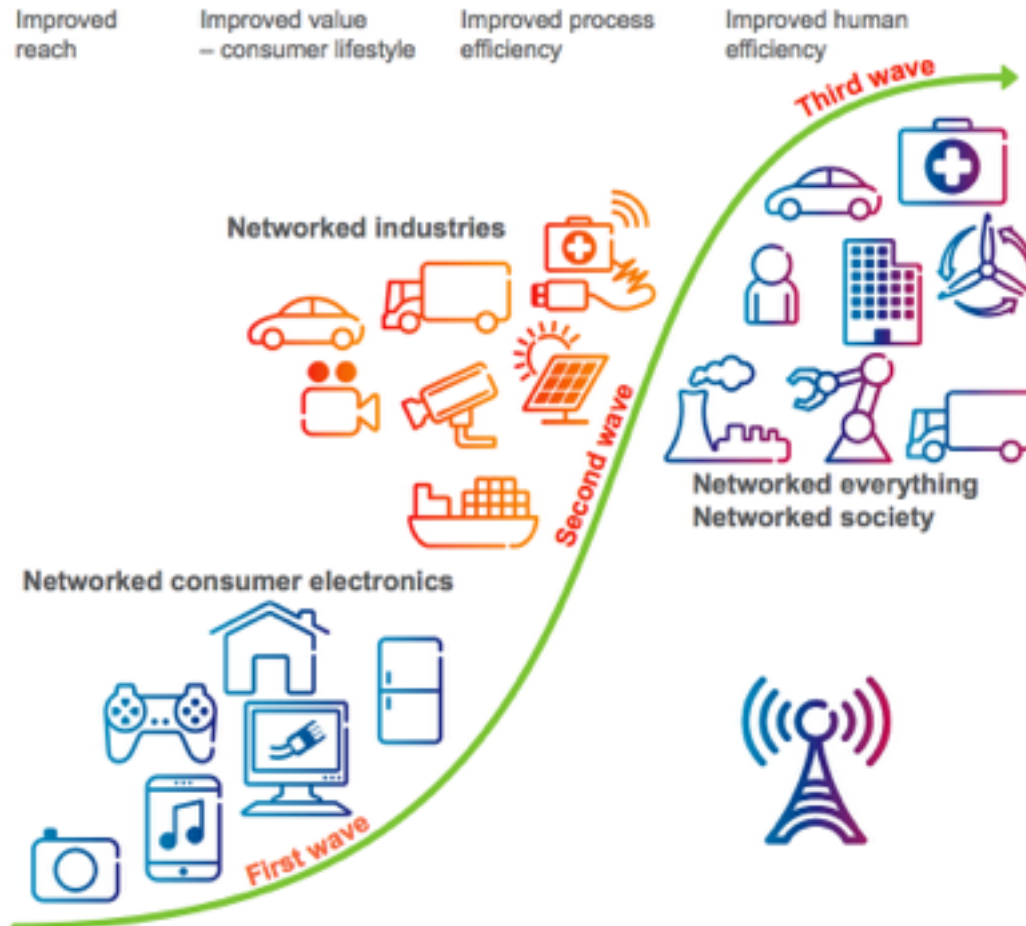
About me

- Graduated with NordSecMob M.Sc in 2010.
- Doctoral Student in DPS Group since 2011.
- Research related to energy efficiency and energy modeling.
- Started with smartphones, now focusing on Internet of Things devices.
- Investigate how to do energy management on a network level in IoT.

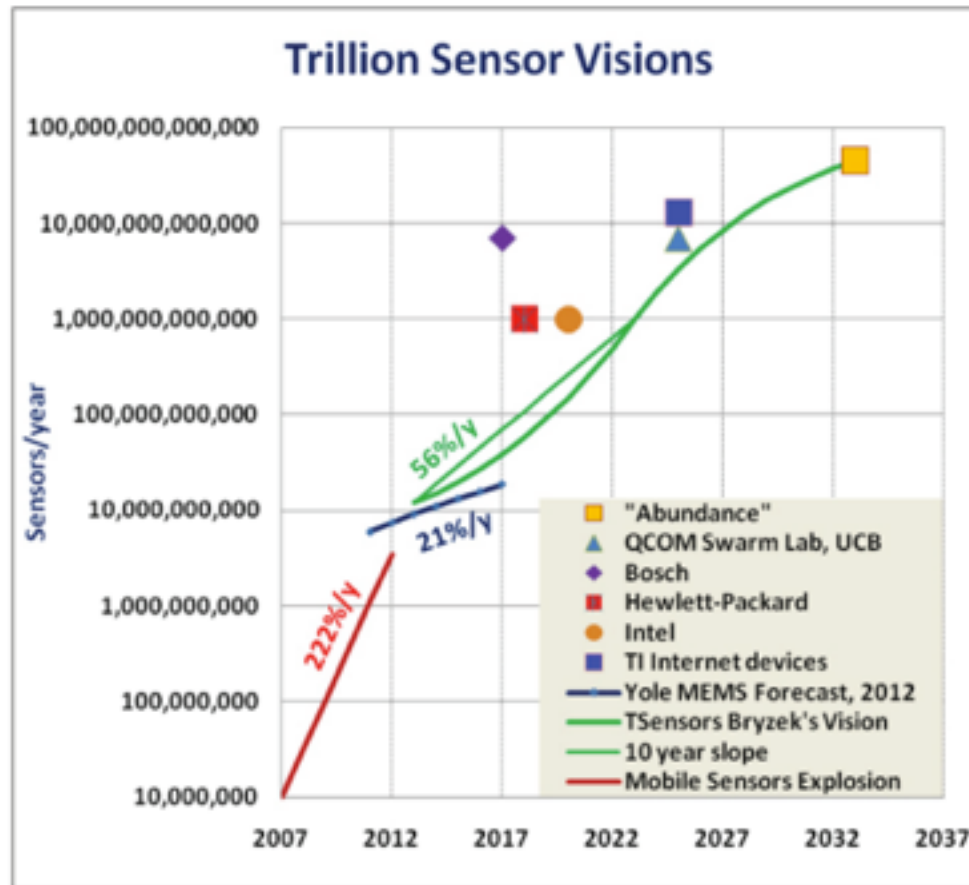
Outline

- Internet of Things
- Hardware
- Event-driven networking
- 6LoWPAN at a glance
- CoAP in depth
- Messages
- Observe
- Conclusions
- References

50 billion connected devices



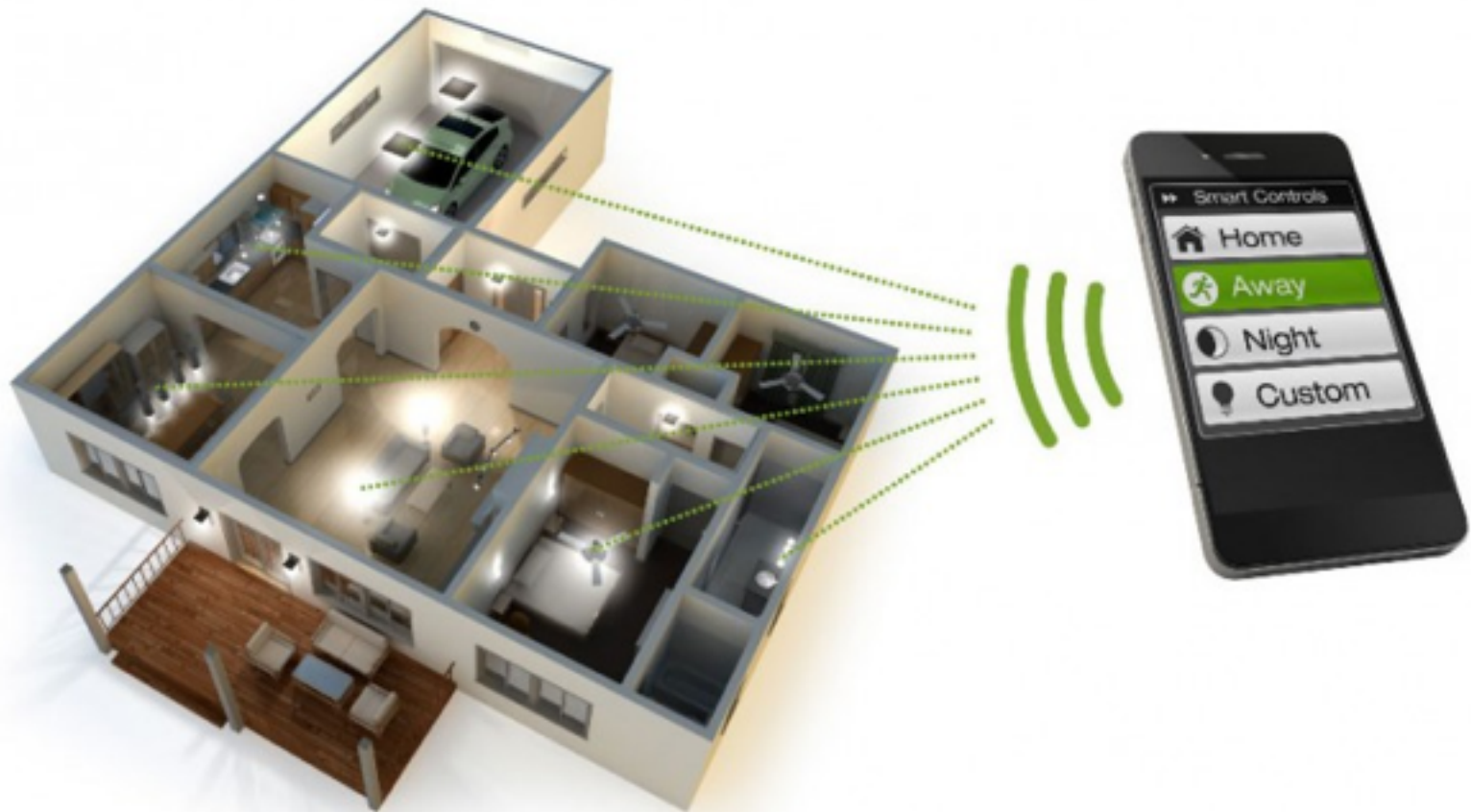
1 Trillion Sensors (yes, with a “T”)



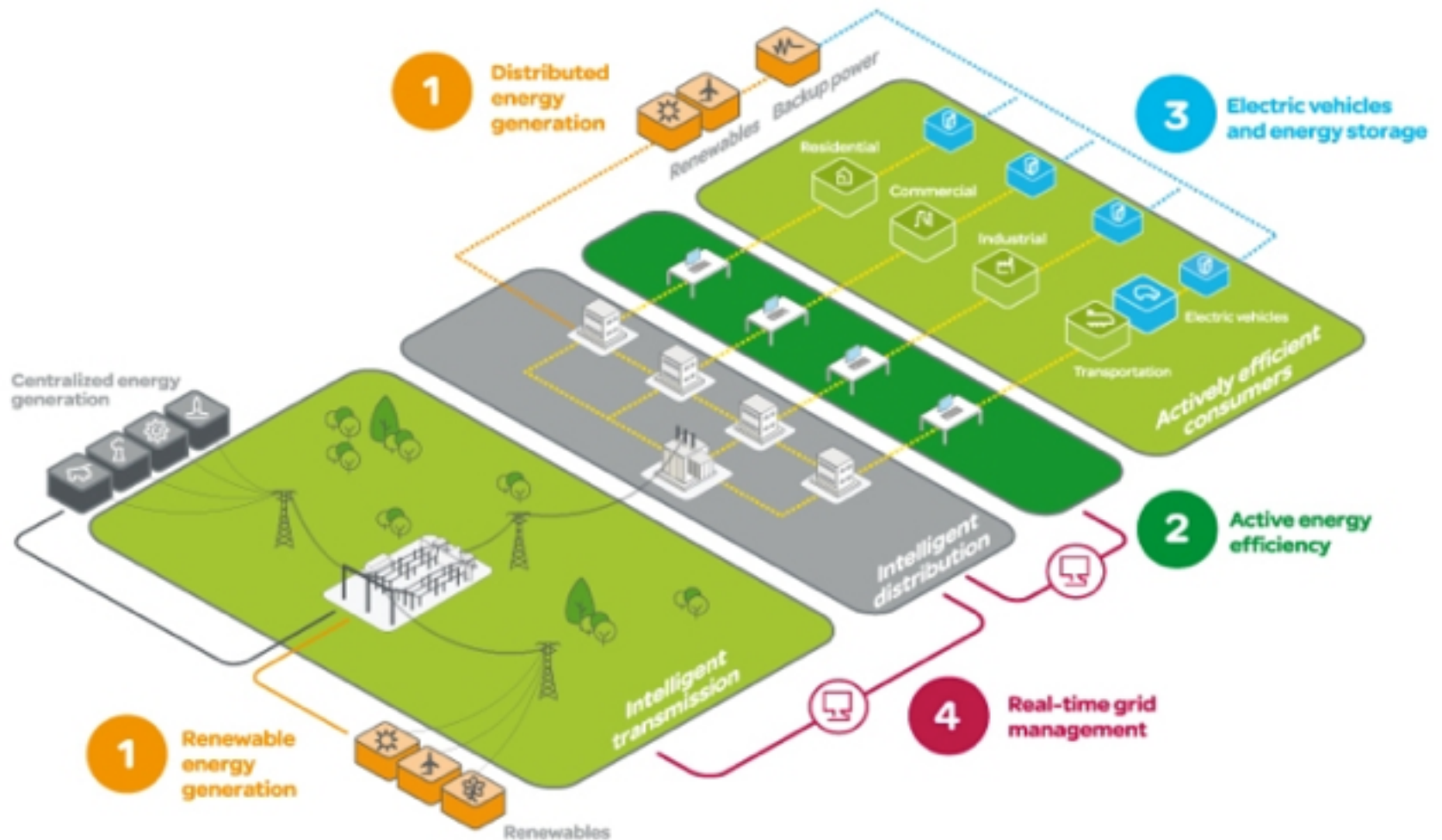
1 Trillion Sensors (yes, with a “T”)



Consumer IoT



Industry IoT - Machine to machine (M2M)



Enevo ONE

Enevo ONE is a complete waste monitoring solution that brings up to 50% savings in waste collection costs. Works with **any type of container** and **any type of waste** mixed, glass, bio, metals or fluids such as oils and waste water etc.

What's included

- ✓ Sensor Device
- ✓ Access to the ONE Collect Web Portal
- ✓ Real Time Fill-Up Measurements
- ✓ Collection Forecasts
- ✓ Daily Collection Lists
- ✓ Alert Service
- ✓ Sensor Warranty and Replacement
- ✓ All Data Transmissions
- ✓ Unlimited Number Of User Accounts
- ✓ 24/7 Email Support



Node hardware - previous gen.



Zolertia Z1 (Contiki OS)

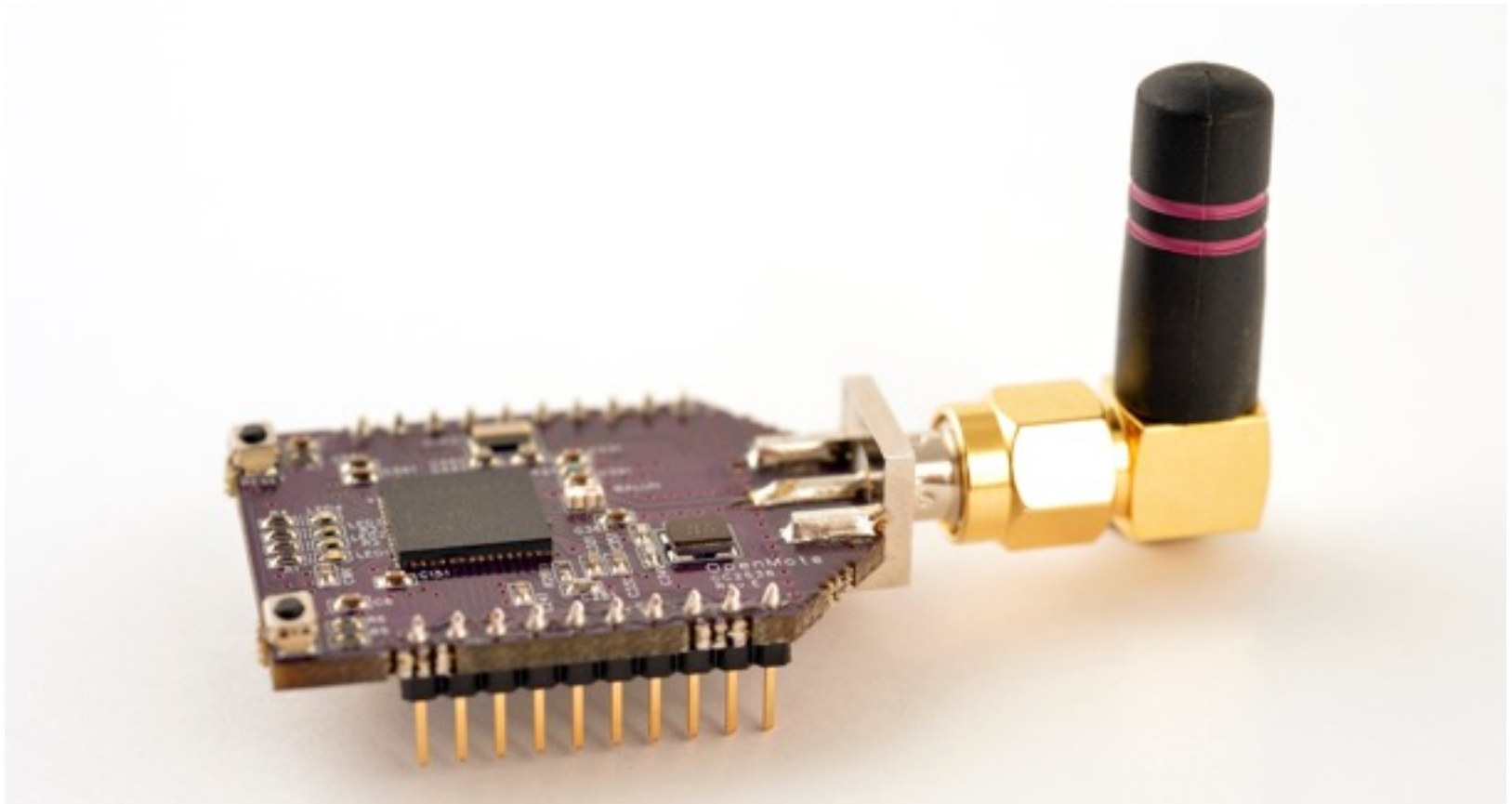


Arduino

Node hardware - previous gen.

- Limited resources:
 - 8KB RAM, 92KB ROM, 16-bit CPU
- Wireless communication: 6lowPAN over 802.15.4, GPRS, 4G
- Long-lasting battery

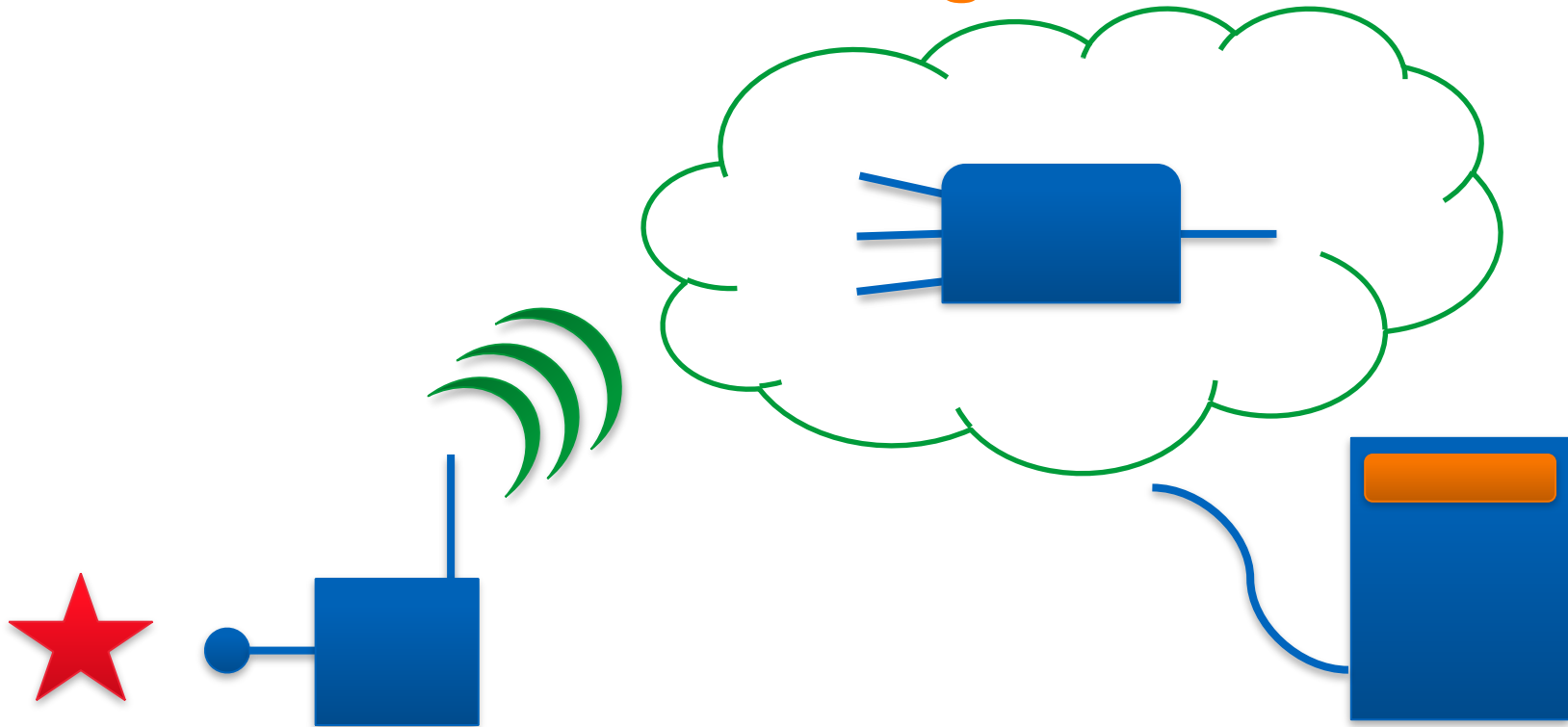
Node hardware - current gen.



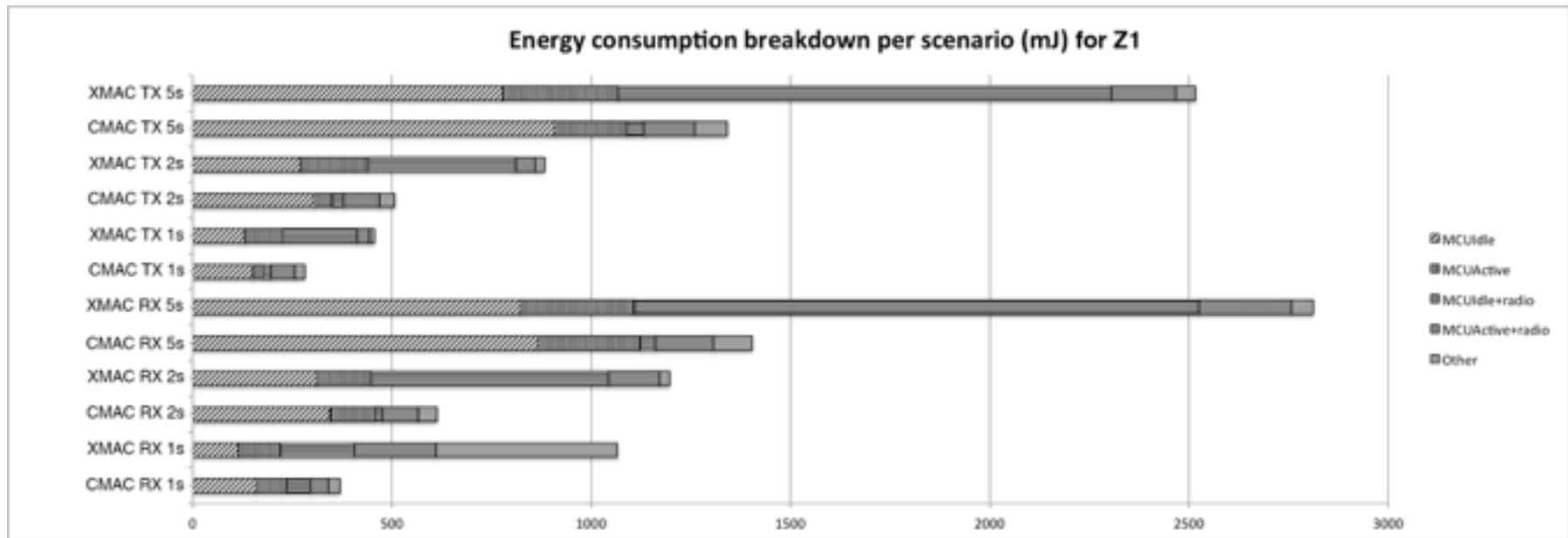
Node hardware - current gen.

- Still a bit limited:
 - 32 MHz 32-bit Cortex-M3 ARM CPU
 - 32 KB RAM, 512 KB Flash
- Connectivity:
 - 6LoWPAN over 802.15.4, 4G, Low-Power WiFi

Event-driven networking



Radio dictates energy consumption



IoT Network Stack

| <i>TCP/IP Protocol stack</i> | <i>IoT protocol stack</i> |
|-------------------------------------|---|
| <i>HTTP</i> | <i>CoAP</i> |
| <i>TCP, UDP</i> | <i>UDP</i> |
| <i>IPv4, IPv6</i> | <i>6LoWPAN</i> |
| <i>Ethernet MAC</i> | <i>ContikiMAC, X-MAC, TSCH</i> |
| <i>Ethernet PHY</i> | <i>IEEE 802.15.4 PHY, Low-power Wifi, 4G</i> |

IPv6 over Low-Power Wireless Personal Area Networks (6LoWPAN)

Features

- “The Internet protocol should be applied to the smallest low-power devices with limited processing capabilities” - *Shelby & Bormann*
- Stateless header compression
- Network autoconfiguration using neighbour discovery
- Support for 802.15.4:
 - addressing schemes (16-bit short address within the PAN)
 - fragmentation

Network autoconfiguration

- Link-layer connectivity between nodes(commissioning):
 - Channel, modulation, data rate
 - Addressing mode (64- or 16-bits)
 - MAC mode (ContikiMAC, TSCH)
- Network layer address configuration, discovery of neighbors, registrations (bootstrapping)
- Routing algorithm sets up paths (route initialization)
- Continuous maintenance of

802.15.4 support: fragmentation

- MTU size: 802.15.4 leaves only ~100 bytes per frame for payload, while max MTU size in IPV6 can be 1280 bytes
- Fragmentation is done by assigning tags to IP packets at the source of fragmentation
- The payloads are assembled at the destination once all the fragments are received (order of arrival is not important)

Constrained Application Protocol (CoAP)

Constrained Application Protocol (CoAP)

- Application level protocol over UDP
- Designed to be used with constrained nodes and lossy networks
- Designed for M2M applications, such as home and infrastructure monitoring
- Built-in resource discovery and observation (“push notification”)
- Block-wise transfer

Constrained Application Protocol (CoAP)

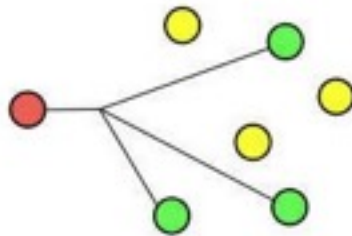
- RESTful for easy interfacing with HTTP
- Multicast support
- Low overhead and simple
- Should fit into a single UDP packet or IEEE 802.15.4 frame

(Helper slide) RESTful

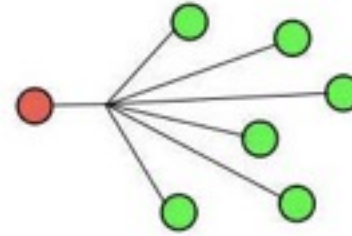
- REpresentational State Transfer is a software architectural style for Web client-server
- Resources are represented as URL:
 - “example.com/profile/johnny”
 - “example.com/domain/sensor3/temp1”
- Resources can be retrieved and manipulated using VERBS:
 - GET, POST, PUT, DELETE
- Example protocol: HTTP

(Helper slide) Uni-, multi-, broadcast

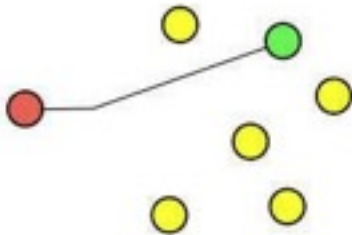
Multicast



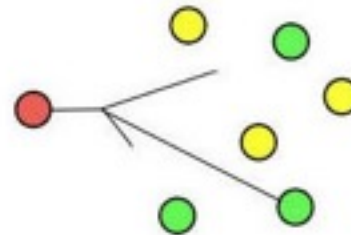
Broadcast



Unicast

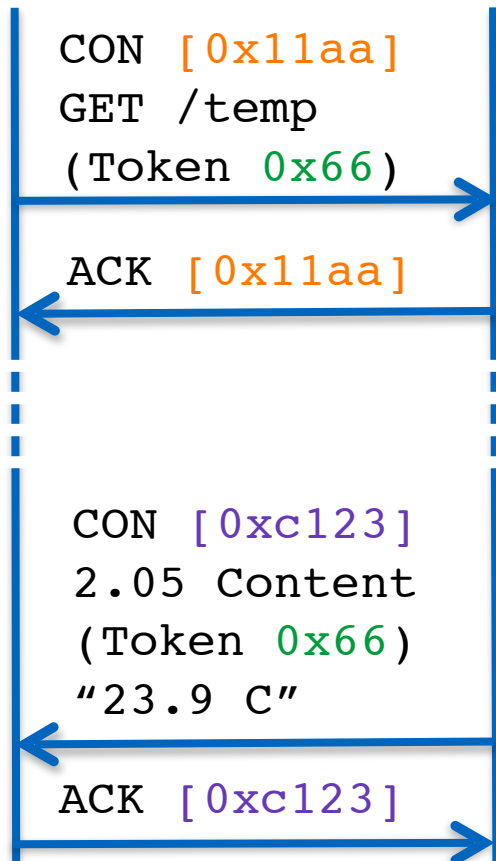


Anycast

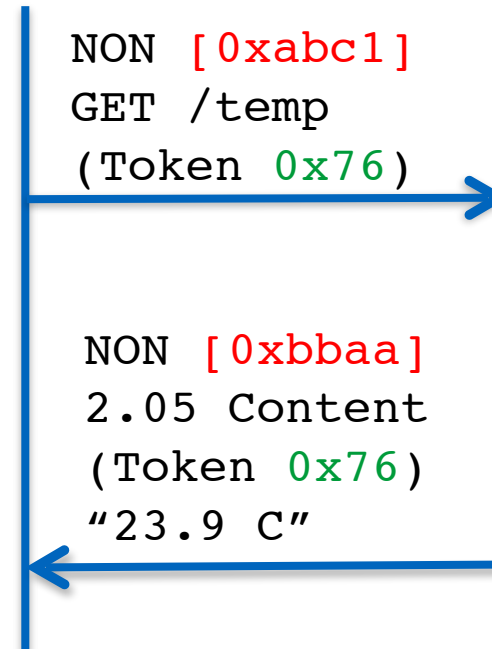


Request/Response model

Client Server



Client Server



CoAP message

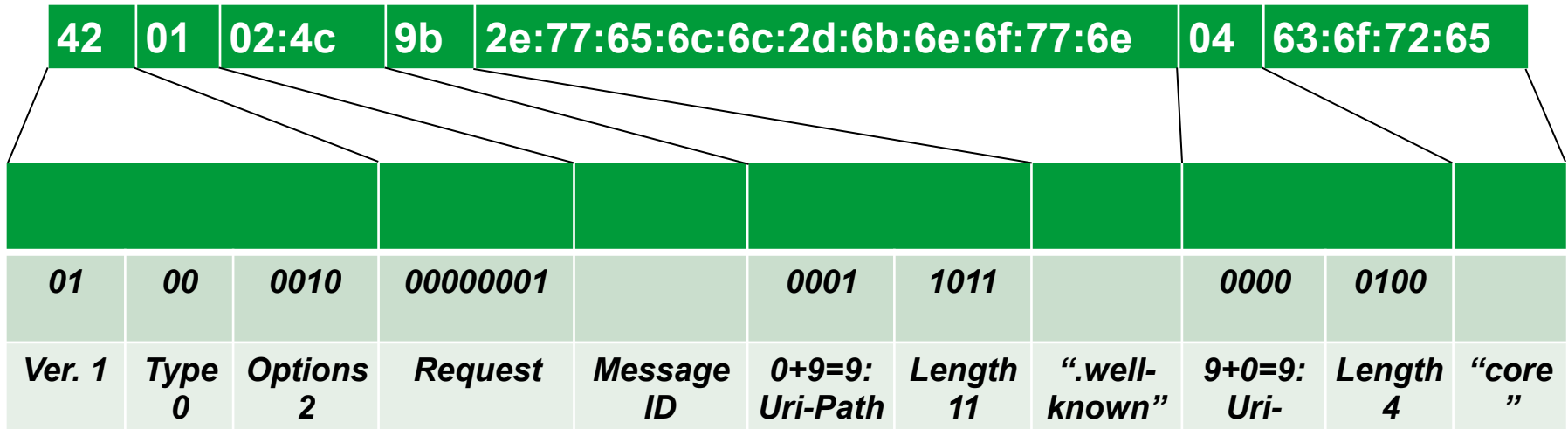
UDP packet

| SRC port | DST port | Length | CRC | CoaP MSG |
|----------|----------|--------|-----|----------|
|----------|----------|--------|-----|----------|

CoAP message

| Ver | T | OC | Code | Message ID |
|---------------------|---|----|---------------|--------------|
| <i>Option Delta</i> | | | <i>Length</i> | <i>Value</i> |
| ... | | | | |
| <i>Payload</i> | | | | |

CoAP message



Message types

- Confirmable (0)
- Non-Confirmable (1)
- Acknowledgement (2)
- Reset (3)

Message Code

- Request (1-31)
- Response (64-191)
- Empty (0)

Options

1 If-Match

3 Uri-Host

4 ETag

5 If-None-Match

7 Uri-Port

8 Location-Path

11 Uri-Path

12 Content-Format

14 Max-Age

15 Uri-Query

17 Accept

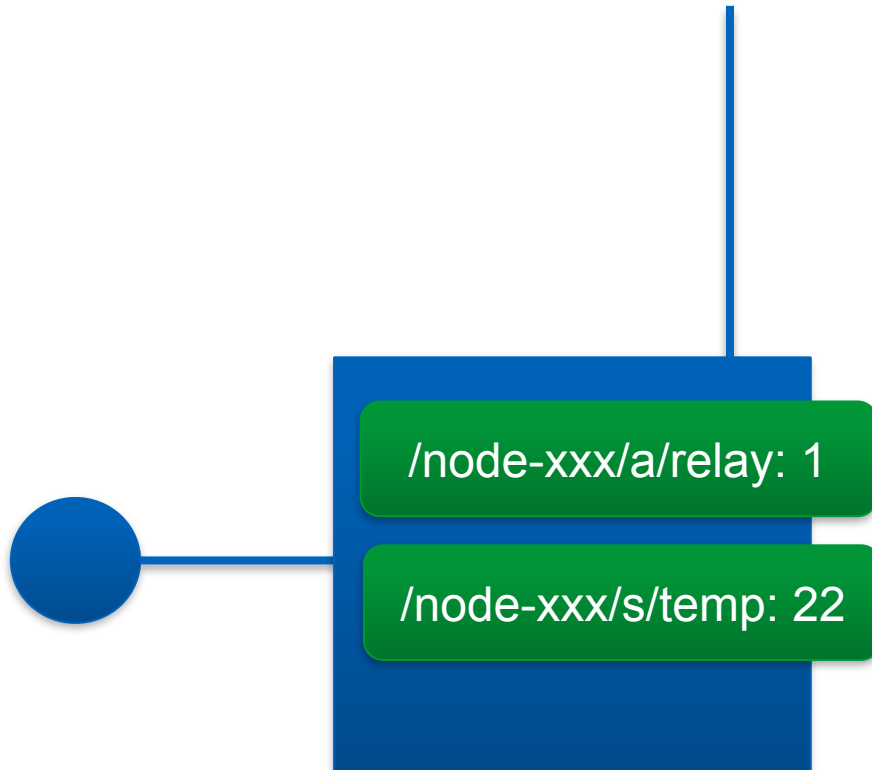
20 Location-Query

35 Proxy-Uri

39 Proxy-Scheme

60 Size1

RESTful protocol



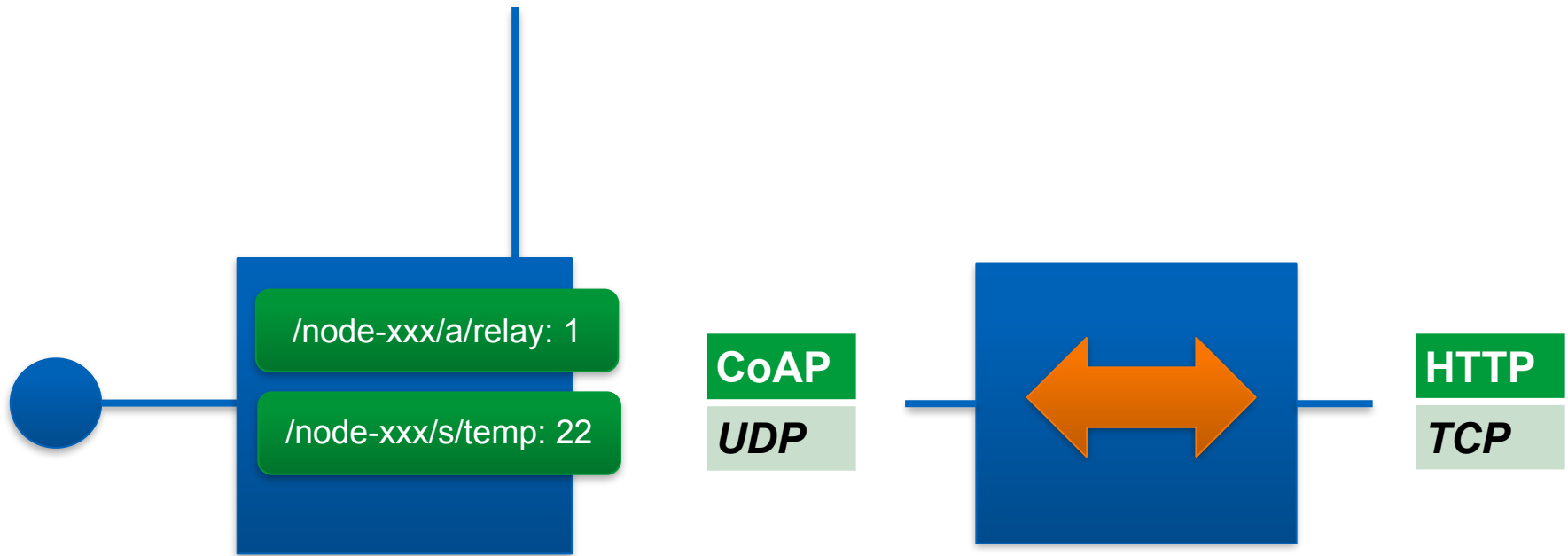
GET /node-xxx/a/relay
-> 1

GET /node-xxx/s/temp
-> 22

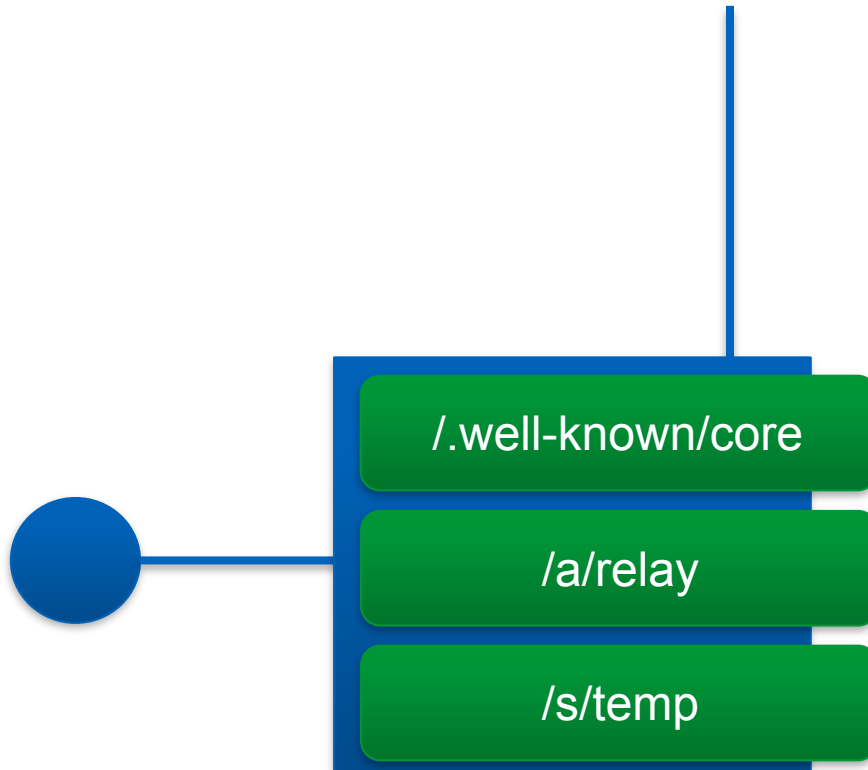
POST /node-xxx/a/relay/1
-> OK

DELETE /node-xxx/a/relay
-> OK

CoAP to HTTP proxy



Resource discovery

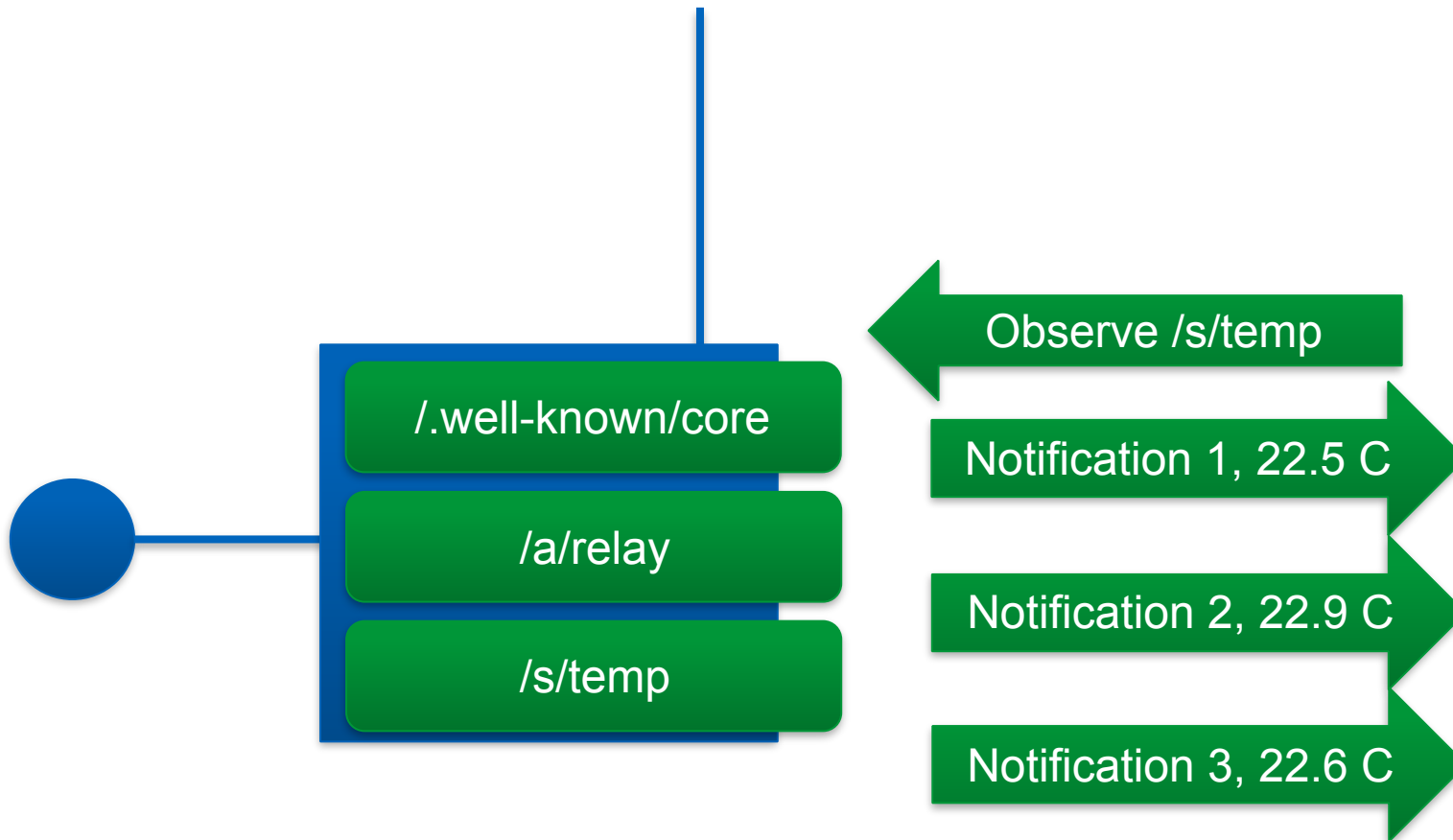


GET `/.well-known/core`

->

`</a/relay>;rt="Relay";ct=0,`
`</s/temp>;rt="Temperature";ct=0`

Observe functionality



Multicast

- One-to-many has obvious application in IoT networks:
 - Data and code dissemination
 - Delivering event data to a set of interested servers
 - Measurement data to more than one subscriber
- Current multicast proposals have drawbacks:
 - SMRF: one-way directional flooding; only useful for data flowing from border router to nodes
 - PIM-WSN: too many forwarding states inherited from PIM-SSM alleviated by Bloom filters

Block-wise transfer

- Transferring large resources with minimal overhead
- Avoid IP fragmentation (datagram > MTU)
- No need to manage state at IP or adaptation layer
- Both parties agree on block size

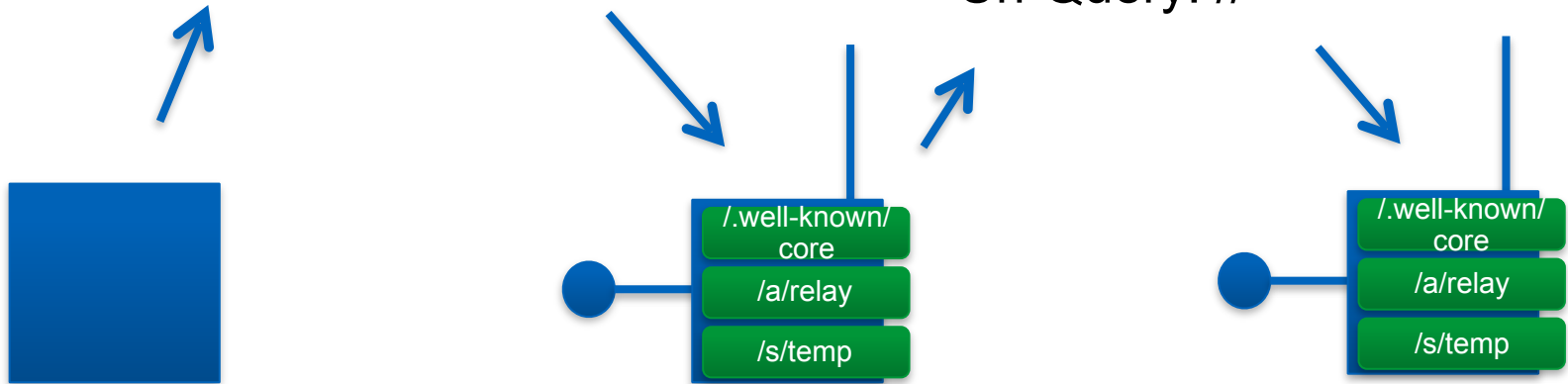
Multicast

- Our proposal (Deng Y., Doctoral student):
 - Any node can be the sender
 - Interoperability with outside domains
 - Link layer broadcast
 - One network interface per node
- RPL (Routing protocol for Low-power and Lossy networks):
 - Uses tree structure
 - Multicast messages can be sent in RPL control messages

Proxy

Proxy-URI: "coap://domain.fi:
60001/node-00x/s/temp/"

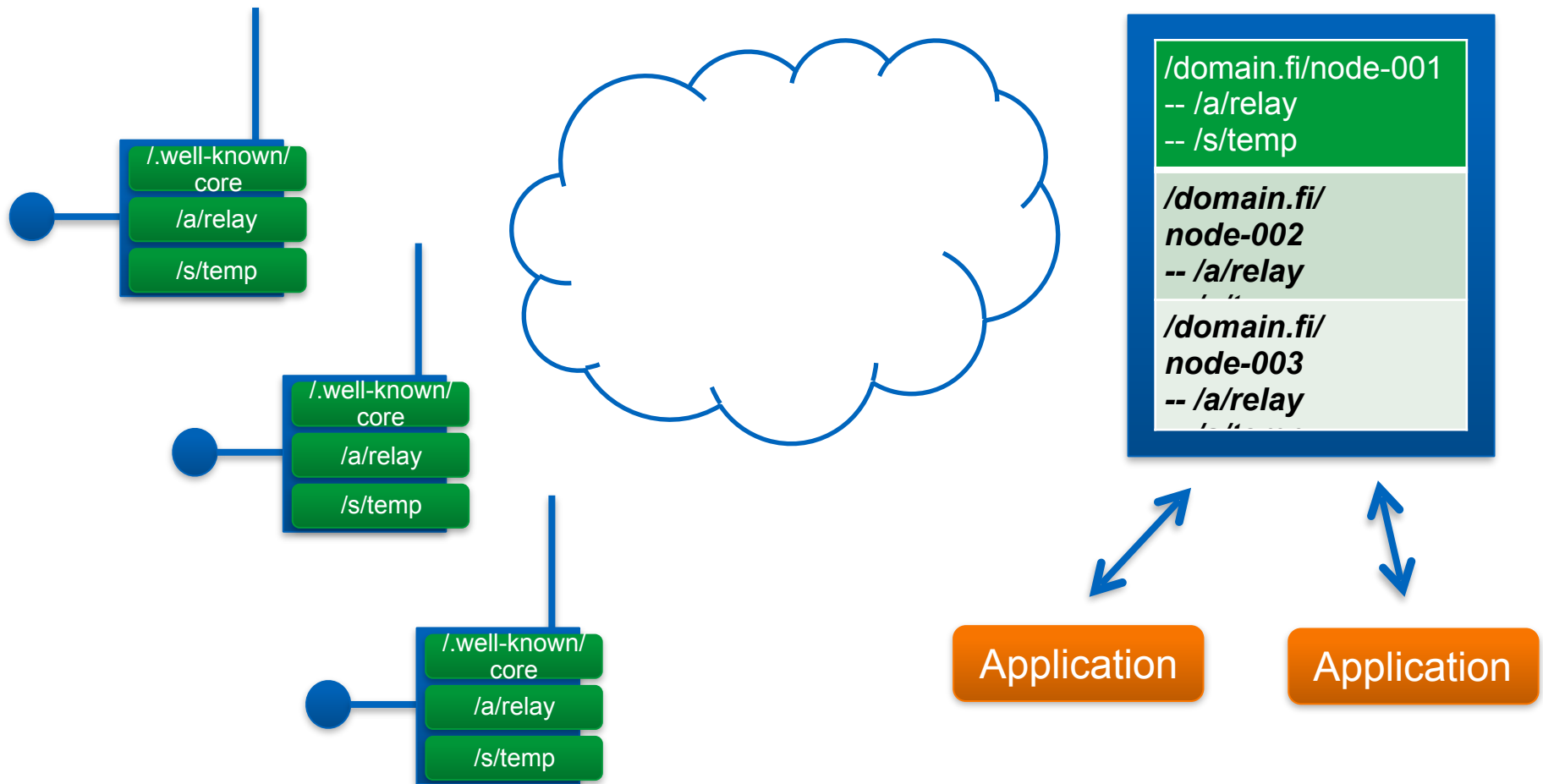
Uri-Host: domain.fi
Uri-Port: 60001
Uri-Path: s
Uri-Path: temp
Uri-Query: //



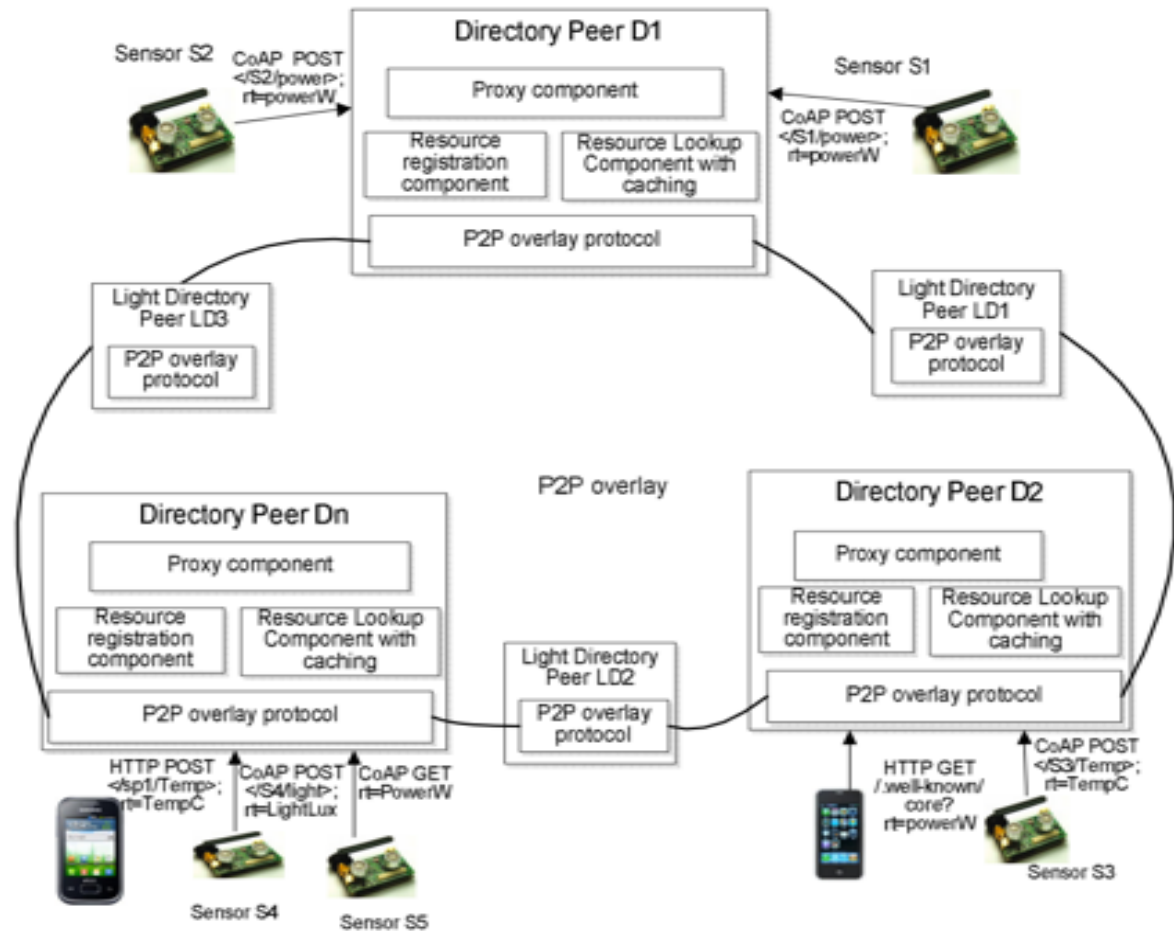
Security

- Similarly to TLS used in HTTP (-> HTTPS), CoAP is secured using Datagram TLS (DTLS)
- DTLS = TLS + features to deal with unreliability of TLS:
 - DTLS records are independent: if record N is lost N+1 can still be decrypted, while N is retransmitted
 - TLS handshake breaks if the packets are out of order; DTLS queues handshake messages until the correct one
 - Application is responsible for dealing with packet reordering, loss, data re-assembly etc.
- Minimal implementation of MUST configurations:
 - Not all cipher suits supported
 - DTLS does not work for multicast communication
- Devices should keep connection open as long as possible to avoid mutual authentication setup overhead

Resource directory



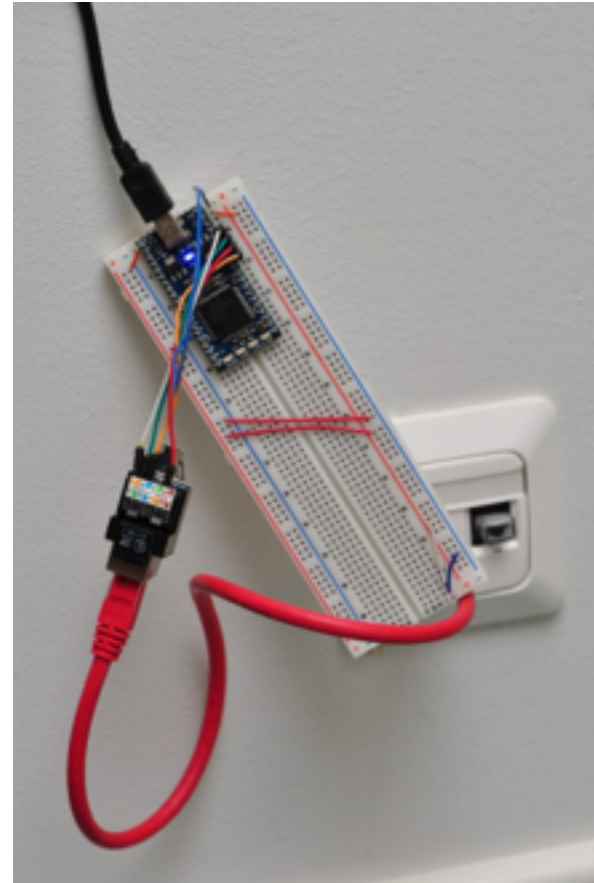
P2P Resource Directory



Hardware

Tiny CoAP sensor by Ericsson Research

- Ethernet, IPv6, UDP
- 48 lines of assembler code
- Does not support Observe



Implementation

| Software | Type | Supported Platform |
|-------------------------------|-----------------------|-----------------------------|
| <i>libcoap</i> | <i>Client, Server</i> | <i>x86, Contiki, tinyOS</i> |
| <i>Californium</i> | <i>Client, Server</i> | <i>Java</i> |
| <i>Copper</i> | <i>Client</i> | <i>Javascript</i> |
| <i>Erbium</i> | <i>Client, Server</i> | <i>Contiki</i> |
| <i>... and many many more</i> | | |

Future/baseless speculation

- CoAP will most likely peak at some point and give way to HTTP even on the most low-powered nodes
- Reasons:
 - Nodes are becoming powerful enough (see OpenMote)
 - Battery life improves
 - 802.15.4 replaced by WiFi/4G/5G
 - Industry wants to have IPv6/TCP/HTTP end-to-end
- Lesson: understand the underlying problem, don't get too attached to a solution

Energy management in large scale IoT networks (Looga et al.)

- Problem:
 - The potential for energy savings is nearing it's end when it comes to the physical layer
 - Large-scale IoT networks require different kind of energy management of nodes, that should be scalable and with minimal overhead
 - Assumptions:
 - Network traffic from M2M nodes is rather predictable
 - Network traffic determines the energy consumption of the node
 - Network topology in production systems is simple
-

Energy management in large scale IoT networks (Looga et al.)

- Our proposal: Traffic-based energy model for 802.15.4 nodes
 - requires only the traffic traces from the node and some of the nodes configuration parameters (MAC layer type, TX/RX values etc.)
 - No overhead on the node
 - Scales linearly with the number of nodes
 - Has good precision (>90%)

References

- 6LoWPAN: <http://datatracker.ietf.org/wg/6lowpan/documents/>
- IETF CoAP draft: <https://datatracker.ietf.org/doc/draft-ietf-core-coap/>
- IETF CoAP Observe draft: <http://tools.ietf.org/html/draft-ietf-core-observe>
- Jari Arikko. Tiny CoAP Sensors. http://www.arkko.com/publications/ietf81_lwip_tiny.pdf
- DTLS: <http://tools.ietf.org/html/rfc4347>
- More than 50 billion connected devices – taking connected devices to mass market and profitability (whitepaper). Ericsson
- Cisco Visual Networking Index: Global Mobile Data Traffic Forecast Update, 2011–2016. Cisco

Thank you!

Questions?