

Dynamic Programming Problem Solving

- Priyansh Agarwal

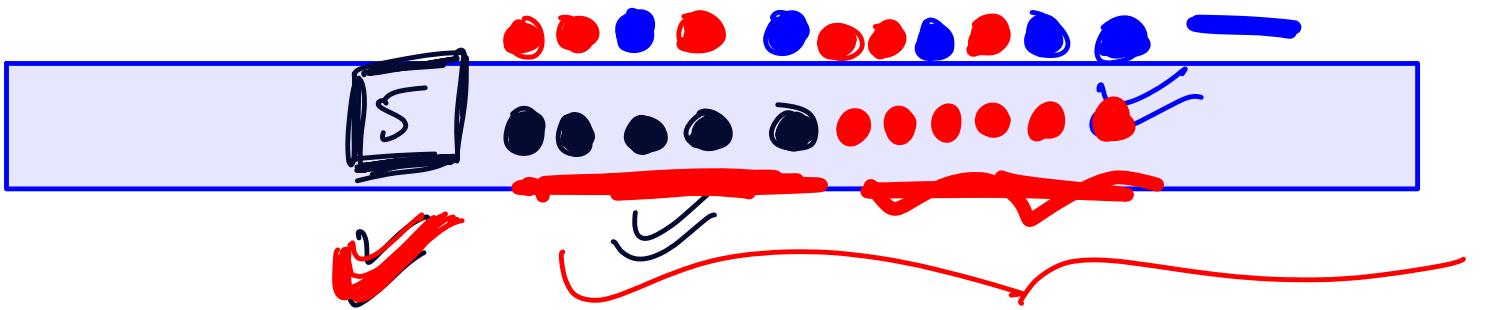
Problem 1: Block Sequence

[CP-31 1500 P3]

Observations

- ① We want to remove minimum no. of elements from the array to make it Beautiful

Removing min elements \Leftrightarrow Retaining max # elements

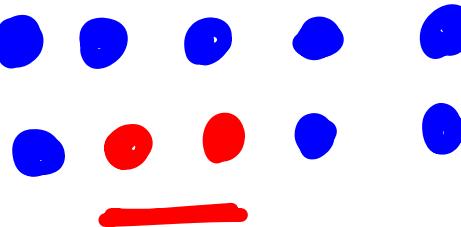


1 5 - - - -

4 - - - -

6 - - - -

1 S



3

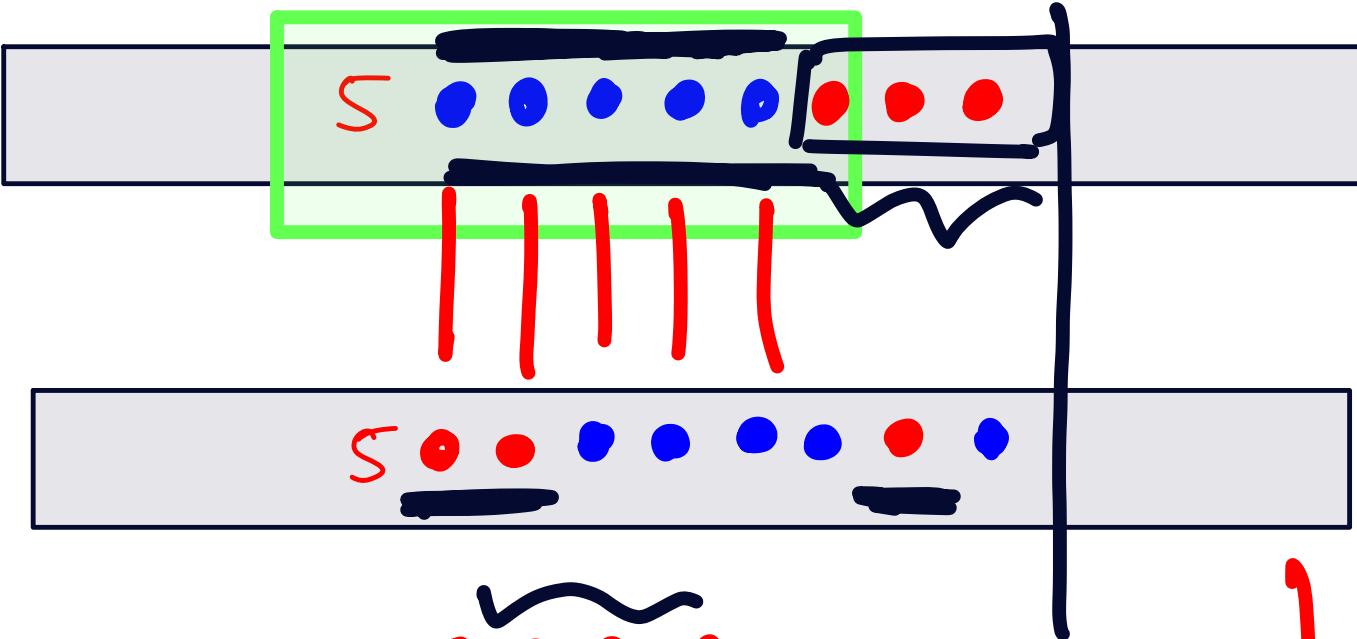
5 4 3 2 2 2 3 4
=====
~~✓~~ ~~✓~~ ~~✓~~ ~~✓~~ ~~✓~~

Priyanshu

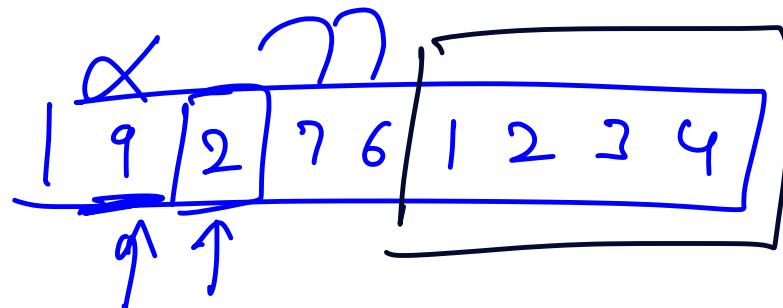
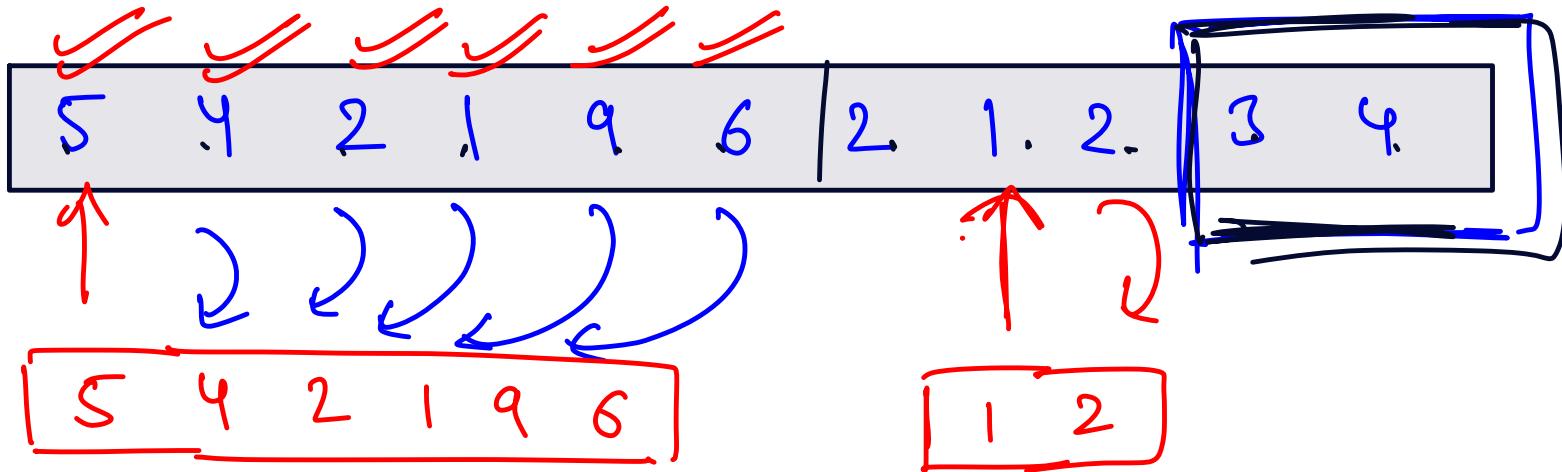
5 4 3 2 2 2 3 4

5 ✓ ✓ ✓ ✓
=====
4 3 2 2 [1 3] 4

② if $\underline{a[i]} = x$ and we have decided
to pick up $\underline{a[i]}$ as the first element in
a block then we need x more elements
to be picked \rightarrow we should pick up
 $a[i+1], a[i+2], \dots, a[i+x-1]$ as our
next x elements.



S



5 4 2 1 2 3 4 5 6 7 8 9 2 1

↓ X X X ↓ ↓ X

1 5 4 2 1 2 3

11 2

State : $d[i]$ = maximum no. of elements we can retain from $[i \text{ to } n-1]$ in the array such that all new blocks will be formed from this range

transition : $d[i] \Rightarrow$

- start a block $\rightarrow d[i+1 + \underline{d[i]}]$
- skip $\rightarrow d[i+1 + \underline{d[i]}]$

Base Case :

$$d[n-1] = 0, d[n] = 0$$

final subproblem $\rightarrow \underline{n - d[0]}$

Problem 2: Color with Occurrences

$$t = \underline{abcabc}$$

$$\underline{\underline{s_1}} = abc$$

$$\underline{\underline{s_2}} = \underline{abcabc}$$

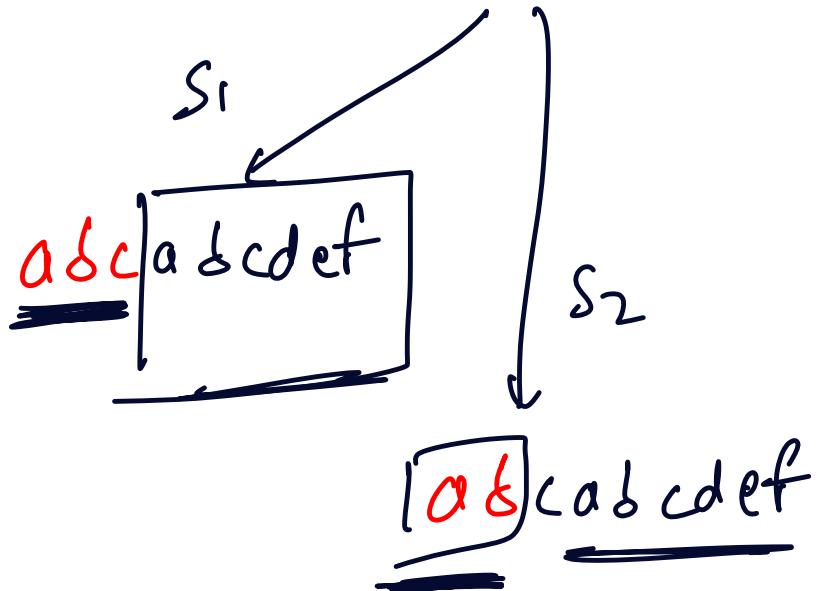
$$t = \underline{\underline{adcd}} \underline{\underline{adcd}}$$

$$\underline{\underline{s_1}} = abc$$

$$\underline{\underline{s_2}} = d$$

$$\underline{\underline{s_3}} = adcadbc$$

$t = abc\alpha bcd\beta ef$



$$s_1 = abc$$

$$s_2 = \beta ef$$

$$s_3 = \alpha bcd$$

$$s_4 = de$$

$t = \underline{abcdef}$ s_1 $\underline{abc} \ def$ $ad\cancel{c}\cancel{d}\cancel{e}\cancel{f}$ $s_1 = abc$ $s_2 = cdef \times$ $s_2 = \cancel{d}c\cancel{d}ef \times$ $s_1 = ad\cancel{c}def$ $t =$

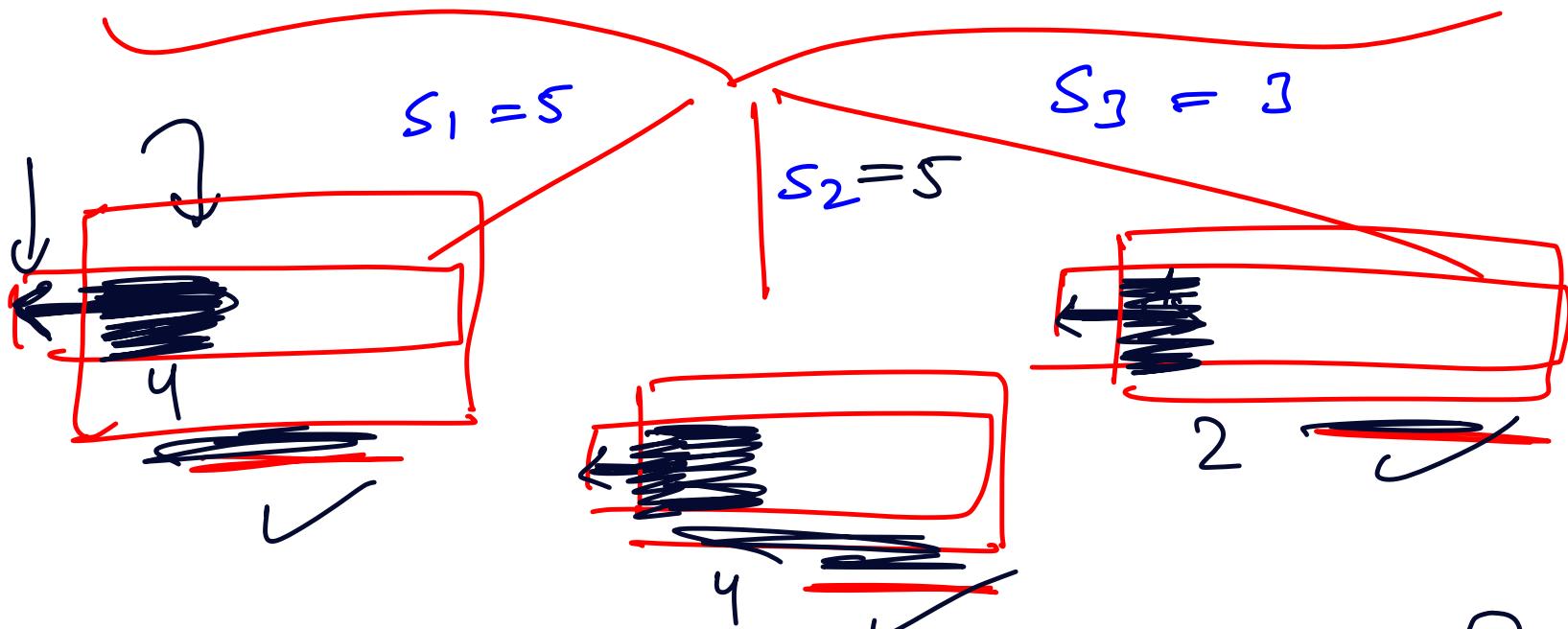
A horizontal line with a double-headed arrow above it, and a bracket underneath, indicating a range or set.

$t = abcdef$

$t = \underline{abcdef}$

$\underline{\quad}$

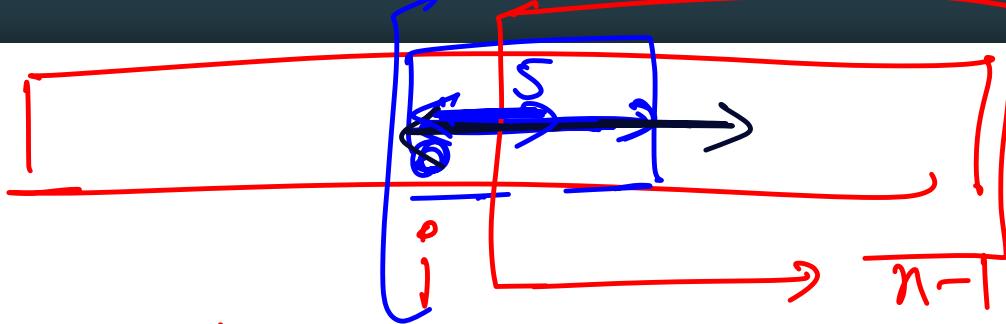
 $t = abcdef$ $s_1 = adc$ $s_2 = adcd$ s_1 $\underline{adc def}$ s_2 $\underline{adcd \underline{ef}}$



→ no. of letters left in total
 → no. of letters in section colored already

State :

$dp[i][k] = \min \text{ no of operations to color all characters from } i \text{ to } n-1 \text{ in red such that } k \text{ characters from } (i) \text{ to } (i+k-1)^{\text{th}} \text{ index are already colored.}$



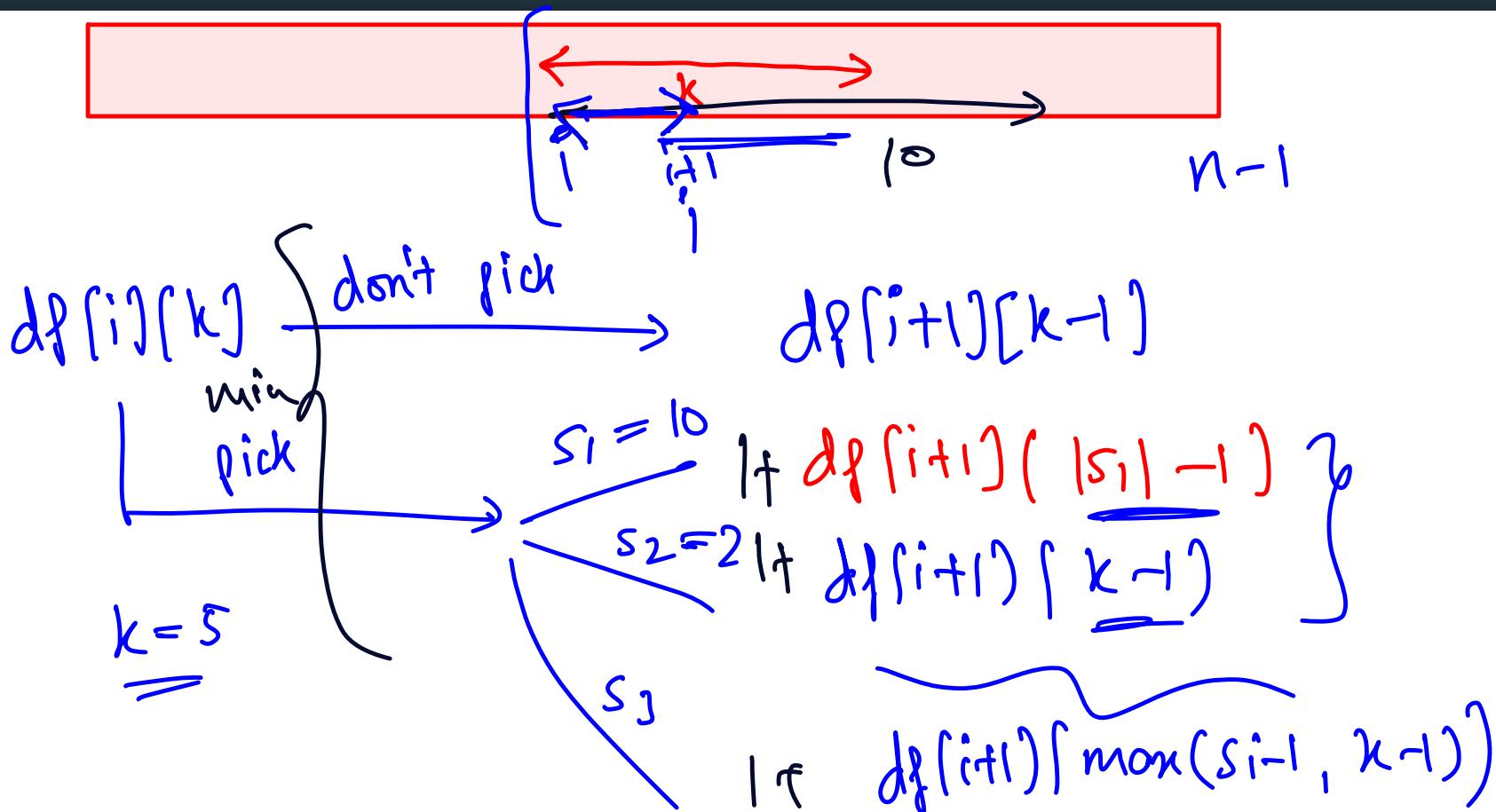
dp[i][k]

don't put anything $\Rightarrow \text{dp}[i+1][k-1]$

put

$1 + \min \{ \text{dp}[i+1][\max(k-1, s_j-1)] \}$

+ all j



Base Case

$dp[i][k]$

$$\left\{ \begin{array}{l} \text{if } k = n-i \\ dp[i][k] = 0 \end{array} \right.$$

$k \leq n-i$ if $k > n-i$

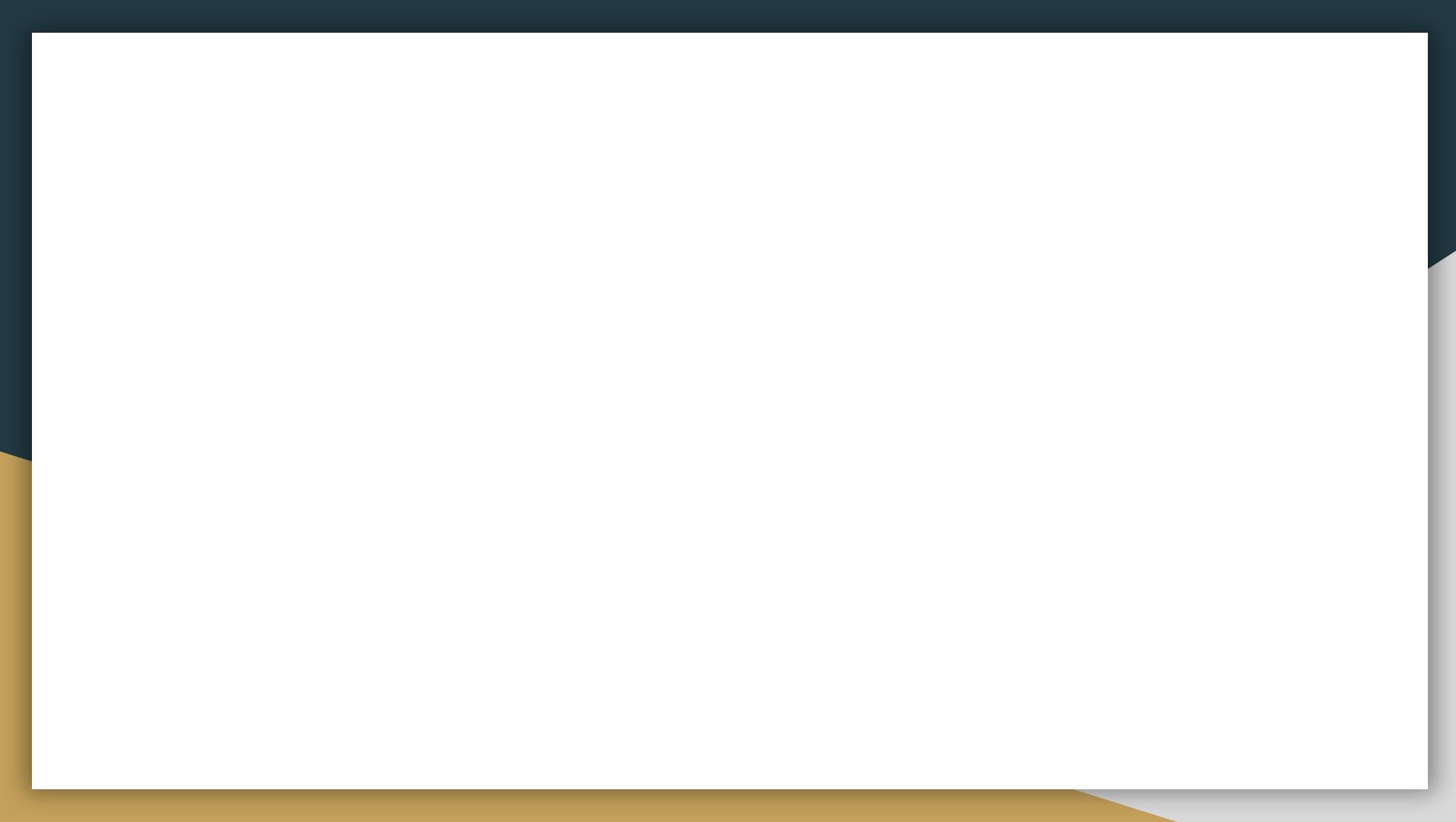
must happen
for a valid state

then
 $dp[i][k] = \infty$

final subproblem : $dp[0][0]$

Problem 3: Flowers

(next class)



Problem 4: Removal Game

(next class)

[CSES]

