

**Assumptions:**

- Fairness (weakly fair behaviour process)
- Cowns cannot become overloaded while muted.
- Mute map entries will eventually be removed and unmuted.
  - Modeled by having overloaded cowns eventually become not overloaded.

**Note:**

- Each while iteration is an atomic step.

EXTENDS *TLC*, *Integers*, *FiniteSets*

*Cowns*  $\triangleq 1 \dots 4$

*Range*(*f*)  $\triangleq \{f[x] : x \in \text{DOMAIN } f\}$

*Min*(*s*)  $\triangleq \text{CHOOSE } x \in s : \forall y \in s \setminus \{x\} : y > x$

*Subsets*(*s*, *min*, *max*)  $\triangleq$   
 $\{cs \in \text{SUBSET } s : \text{Cardinality}(cs) \geq \text{min} \wedge \text{Cardinality}(cs) \leq \text{max}\}$

--algorithm *backpressure*

**variables**

*available* = *Cowns*,  
*overloaded* = {},  
*muted* = {},  
*unmutable* = {},  
*mute\_map* = [*c* ∈ *Cowns* ↦ {}],  
*refcount* = [*c* ∈ *Cowns* ↦ 0],  
*rc\_barrier* = 0;

**define**

*BehaviourCount*  $\triangleq 3$   
*MutedInv*  $\triangleq \text{available} \cap \text{muted} = \{\}$   
*TODO*:  $\Box \Diamond (\text{unmutable} \cap \text{muted} = \{\})$   
*UnmutableInv*  $\triangleq \text{overloaded} \cap \text{muted} = \{\}$   
*RefcountInv*  $\triangleq \forall c \in \text{Cowns} : \text{refcount}[c] \geq 0$   
*MuteMapInv*  $\triangleq \forall m \in \text{muted} : m \in \text{UNION } \text{Range}(\text{mute\_map})$   
*TypeInvariant*  $\triangleq \text{MutedInv} \wedge \text{UnmutableInv} \wedge \text{RefcountInv} \wedge \text{MuteMapInv}$   
*RefcountDrop*  $\triangleq \Diamond \Box (\forall c \in \text{Cowns} : \text{refcount}[c] = 0)$   
*WillUnmute*  $\triangleq$   
 $\Box \Diamond (\forall k \in \text{DOMAIN } \text{mute\_map} : \text{mute\_map}[k] = \{\} \vee k \in \text{overloaded})$   
*TemporalProp*  $\triangleq \text{RefcountDrop} \wedge \text{WillUnmute}$   
*ZeroRC*(*c*)  $\triangleq \text{rc\_barrier} = \text{BehaviourCount} \wedge \text{refcount}[c] = 0$   
*TriggersUnmute*(*mutor*)  $\triangleq \text{mutor} \in \text{overloaded} \vee \text{ZeroRC}(\text{mutor})$

**end define ;**

**fair process** *behaviour* ∈ 1 .. *BehaviourCount*

**variables**

*required* ∈ *Subsets*(*Cowns*, 0, 3),  
*next*, *acquired* = {}, *mutor*, *muting* = {}, *unmute\_set*

**begin**

*Send*:

*refcount* :=  
 $[c \in \text{Cowns} \mapsto \text{IF } c \in \text{required} \text{ THEN } \text{refcount}[c] + 1 \text{ ELSE } \text{refcount}[c]]$ ;  
*rc\_barrier* := *rc\_barrier* + 1;

Empty required set used to represent fewer behaviours in the system.

**if** *required* = {} **then goto** *Done* **; end if ;**

*Unmute*:

**while** ( $\exists r \in \text{required} : r \notin \text{unmutable}$ )  
 $\wedge (\text{overloaded} \cap \text{required} \neq \{\})$

```

do
  next := Min( $\{r \in \text{required} : r \notin \text{unmutable}\}$ );
  unmutable := unmutable  $\cup$  {next};
  if next  $\in$  muted then
    muted := muted  $\setminus$  {next};
    available := available  $\cup$  {next};
  end if ;
end while ;

```

Acquire:

```

while required  $\neq$  {} do
  next := Min(required);
  await next  $\in$  available ;
  acquired := acquired  $\cup$  {next};
  required := required  $\setminus$  {next};
  available := available  $\setminus$  {next};
end while ;

```

Action:

```

assert required = {};
assert acquired  $\cap$  muted = {};
if (overloaded  $\neq$  {})  $\wedge$  (acquired  $\cap$  overloaded = {}) then
  either
    with mutor_  $\in$  overloaded do mutor := mutor_end with ;
    muting := acquired  $\setminus$  unmutable ;
  or
    skip ;
  end either ;
end if ;

```

Complete:

Arbitrarily toggle overloaded state of some acquired cowns.

```

with overloading  $\in$  Subsets(acquired  $\setminus$  muting, 0, 3) do
  with unoverloading  $\in$  Subsets(acquired  $\cap$  overloaded, 0, 3) do
    overloaded := (overloaded  $\cup$  overloading)  $\setminus$  unoverloading ;
  end with ;
end with ;

if mutor  $\neq$  defaultInitValue then
  muted := muted  $\cup$  muting ;
  mute_map[mutor] := mute_map[mutor]  $\cup$  muting ;
end if ;

```

```

available := available  $\cup$  (acquired  $\setminus$  muting);
muting := {};
refcount :=
  [c  $\in$  Cowns  $\mapsto$  IF c  $\in$  acquired THEN refcount[c] - 1 ELSE refcount[c]];
acquired := {};
assert acquired  $\cup$  required = {};

```

MuteMapScan:

```

unmute_set :=
  UNION Range([c  $\in$  {k  $\in$  Cowns : TriggersUnmute(k)}  $\mapsto$  mute_map[c]]);
mute_map := [c  $\in$  Cowns  $\mapsto$  IF TriggersUnmute(c) THEN {} ELSE mute_map[c]];
muted := muted  $\setminus$  unmute_set ;
available := available  $\cup$  unmute_set ;

```

end process ;

end algorithm ;

BEGIN TRANSLATION — the hash of the PCal code: PCal-bc8e2f6d1cef2a2ee2ced52555aa4e69  
 CONSTANT defaultInitValue  
 VARIABLES available, overloaded, muted, unmutable, mute\_map, refcount,

$rc\_barrier, pc$

define statement

$BehaviourCount \triangleq 3$

$MutedInv \triangleq available \cap muted = \{\}$

$UnmutableInv \triangleq overloaded \cap muted = \{\}$

$RefCountInv \triangleq \forall c \in Cowns : refcount[c] \geq 0$

$MuteMapInv \triangleq \forall m \in muted : m \in \text{UNION } Range(mute\_map)$

$TypeInvariant \triangleq MutedInv \wedge UnmutableInv \wedge RefcountInv \wedge MuteMapInv$

$RefCountDrop \triangleq \Diamond \Box (\forall c \in Cowns : refcount[c] = 0)$

$WillUnmute \triangleq$

$\Box \Diamond (\forall k \in \text{DOMAIN } mute\_map : mute\_map[k] = \{\} \vee k \in overloaded)$

$TemporalProp \triangleq RefcountDrop \wedge WillUnmute$

$ZeroRC(c) \triangleq rc\_barrier = BehaviourCount \wedge refcount[c] = 0$

$TriggersUnmute(mutor) \triangleq mutor \in overloaded \vee ZeroRC(mutor)$

VARIABLES  $required, next, acquired, mutor, muting, unmute\_set$

$vars \triangleq \langle available, overloaded, muted, unmutable, mute\_map, refcount, \\ rc\_barrier, pc, required, next, acquired, mutor, muting, \\ unmute\_set \rangle$

$ProcSet \triangleq (1 \dots BehaviourCount)$

$Init \triangleq$  Global variables  
 $\wedge available = Cowns$   
 $\wedge overloaded = \{\}$   
 $\wedge muted = \{\}$   
 $\wedge unmutable = \{\}$   
 $\wedge mute\_map = [c \in Cowns \mapsto \{\}]$   
 $\wedge refcount = [c \in Cowns \mapsto 0]$   
 $\wedge rc\_barrier = 0$   
 Process behaviour  
 $\wedge required \in [1 \dots BehaviourCount \rightarrow Subsets(Cowns, 0, 3)]$   
 $\wedge next = [self \in 1 \dots BehaviourCount \mapsto defaultInitValue]$   
 $\wedge acquired = [self \in 1 \dots BehaviourCount \mapsto \{\}]$   
 $\wedge mutor = [self \in 1 \dots BehaviourCount \mapsto defaultInitValue]$   
 $\wedge muting = [self \in 1 \dots BehaviourCount \mapsto \{\}]$   
 $\wedge unmute\_set = [self \in 1 \dots BehaviourCount \mapsto defaultInitValue]$   
 $\wedge pc = [self \in ProcSet \mapsto \text{"Send"}]$

$Send(self) \triangleq$   $\wedge pc[self] = \text{"Send"}$   
 $\wedge refcount' = [c \in Cowns \mapsto \text{IF } c \in required[self] \text{ THEN } refcount[c] + 1 \text{ ELSE } refcount[c]]$   
 $\wedge rc\_barrier' = rc\_barrier + 1$   
 $\wedge \text{IF } required[self] = \{\}$   
     THEN  $\wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Done"}]$   
     ELSE  $\wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Unmute"}]$   
 $\wedge \text{UNCHANGED } \langle available, overloaded, muted, unmutable, \\ mute\_map, required, next, acquired, mutor, \\ muting, unmute\_set \rangle$

$Unmute(self) \triangleq$   $\wedge pc[self] = \text{"Unmute"}$   
 $\wedge \text{IF } (\exists r \in required[self] : r \notin unmutable)$   
 $\wedge (overloaded \cap required[self] \neq \{\})$   
     THEN  $\wedge next' = [next \text{ EXCEPT } ![self] = \text{Min}(\{r \in required[self] : r \notin unmutable\})]$   
      $\wedge unmutable' = (unmutable \cup \{next'[self]\})$   
      $\wedge \text{IF } next'[self] \in muted$   
         THEN  $\wedge muted' = muted \setminus \{next'[self]\}$   
          $\wedge available' = (available \cup \{next'[self]\})$   
     ELSE  $\wedge \text{TRUE}$

$$\begin{aligned}
& \wedge \text{UNCHANGED } \langle \text{available}, \text{muted} \rangle \\
& \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Unmute"}] \\
\text{ELSE } & \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Acquire"}] \\
& \wedge \text{UNCHANGED } \langle \text{available}, \text{muted}, \text{immutable}, \text{next} \rangle \\
& \wedge \text{UNCHANGED } \langle \text{overloaded}, \text{mute\_map}, \text{refcount}, \text{rc\_barrier}, \\
& \quad \text{required}, \text{acquired}, \text{mutor}, \text{muting}, \text{unmute\_set} \rangle \\
\text{Acquire}(self) \triangleq & \wedge pc[self] = \text{"Acquire"} \\
& \wedge \text{IF } \text{required}[self] \neq \{\} \\
& \quad \text{THEN } \wedge \text{next}' = [\text{next} \text{ EXCEPT } ![self] = \text{Min}(\text{required}[self])] \\
& \quad \wedge \text{next}'[self] \in \text{available} \\
& \quad \wedge \text{acquired}' = [\text{acquired} \text{ EXCEPT } ![self] = \text{acquired}[self] \cup \{\text{next}'[self]\}] \\
& \quad \wedge \text{required}' = [\text{required} \text{ EXCEPT } ![self] = \text{required}[self] \setminus \{\text{next}'[self]\}] \\
& \quad \wedge \text{available}' = \text{available} \setminus \{\text{next}'[self]\} \\
& \quad \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Acquire"}] \\
& \quad \text{ELSE } \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Action"}] \\
& \quad \wedge \text{UNCHANGED } \langle \text{available}, \text{required}, \text{next}, \\
& \quad \quad \text{acquired} \rangle \\
& \wedge \text{UNCHANGED } \langle \text{overloaded}, \text{muted}, \text{immutable}, \text{mute\_map}, \\
& \quad \text{refcount}, \text{rc\_barrier}, \text{mutor}, \text{muting}, \\
& \quad \text{unmute\_set} \rangle \\
\text{Action}(self) \triangleq & \wedge pc[self] = \text{"Action"} \\
& \wedge \text{Assert}(\text{required}[self] = \{\}, \\
& \quad \text{"Failure of assertion at line 78, column 3."}) \\
& \wedge \text{Assert}(\text{acquired}[self] \cap \text{muted} = \{\}, \\
& \quad \text{"Failure of assertion at line 79, column 3."}) \\
& \wedge \text{IF } (\text{overloaded} \neq \{\}) \wedge (\text{acquired}[self] \cap \text{overloaded} = \{\}) \\
& \quad \text{THEN } \wedge \vee \wedge \exists \text{mutor\_} \in \text{overloaded} : \\
& \quad \quad \text{mutor}' = [\text{mutor} \text{ EXCEPT } ![self] = \text{mutor\_}] \\
& \quad \quad \wedge \text{muting}' = [\text{muting} \text{ EXCEPT } ![self] = \text{acquired}[self] \setminus \text{immutable}] \\
& \quad \vee \wedge \text{TRUE} \\
& \quad \wedge \text{UNCHANGED } \langle \text{mutor}, \text{muting} \rangle \\
& \quad \text{ELSE } \wedge \text{TRUE} \\
& \quad \wedge \text{UNCHANGED } \langle \text{mutor}, \text{muting} \rangle \\
& \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Complete"}] \\
& \wedge \text{UNCHANGED } \langle \text{available}, \text{overloaded}, \text{muted}, \text{immutable}, \\
& \quad \text{mute\_map}, \text{refcount}, \text{rc\_barrier}, \text{required}, \text{next}, \\
& \quad \text{acquired}, \text{unmute\_set} \rangle \\
\text{Complete}(self) \triangleq & \wedge pc[self] = \text{"Complete"} \\
& \wedge \exists \text{overloading} \in \text{Subsets}(\text{acquired}[self] \setminus \text{muting}[self], 0, 3) : \\
& \quad \exists \text{unoverloading} \in \text{Subsets}(\text{acquired}[self] \cap \text{overloaded}, 0, 3) : \\
& \quad \quad \text{overloaded}' = (\text{overloaded} \cup \text{overloading}) \setminus \text{unoverloading} \\
& \wedge \text{IF } \text{mutor}[self] \neq \text{defaultInitValue} \\
& \quad \text{THEN } \wedge \text{muted}' = (\text{muted} \cup \text{muting}[self]) \\
& \quad \quad \wedge \text{mute\_map}' = [\text{mute\_map} \text{ EXCEPT } ![mutor[self]] = \text{mute\_map}[\text{mutor}[self]] \cup \text{muting}[self]] \\
& \quad \text{ELSE } \wedge \text{TRUE} \\
& \quad \wedge \text{UNCHANGED } \langle \text{muted}, \text{mute\_map} \rangle \\
& \wedge \text{available}' = (\text{available} \cup (\text{acquired}[self] \setminus \text{muting}[self])) \\
& \wedge \text{muting}' = [\text{muting} \text{ EXCEPT } ![self] = \{\}] \\
& \wedge \text{refcount}' = [c \in \text{Cowns} \mapsto \text{IF } c \in \text{acquired}[self] \text{ THEN } \text{refcount}[c] - 1 \text{ ELSE } \text{refcount}[c]] \\
& \wedge \text{acquired}' = [\text{acquired} \text{ EXCEPT } ![self] = \{\}] \\
& \wedge \text{Assert}(\text{acquired}'[self] \cup \text{required}[self] = \{\}, \\
& \quad \text{"Failure of assertion at line 107, column 3."}) \\
& \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"MuteMapScan"}] \\
& \wedge \text{UNCHANGED } \langle \text{immutable}, \text{rc\_barrier}, \text{required}, \text{next}, \text{mutor}, \\
& \quad \text{unmute\_set} \rangle \\
\text{MuteMapScan}(self) \triangleq & \wedge pc[self] = \text{"MuteMapScan"} \\
& \wedge \text{unmute\_set}' = [\text{unmute\_set} \text{ EXCEPT } ![self] =
\end{aligned}$$

$$\begin{aligned}
& \text{UNION } \text{Range}([c \in \{k \in \text{Cowns} : \text{TriggersUnmute}(k)\} \mapsto \text{mute\_map}[c]]) \\
& \wedge \text{mute\_map}' = [c \in \text{Cowns} \mapsto \text{IF } \text{TriggersUnmute}(c) \text{ THEN } \{\} \text{ ELSE } \text{mute\_map}[c]] \\
& \wedge \text{muted}' = \text{muted} \setminus \text{unmute\_set}'[\text{self}] \\
& \wedge \text{available}' = (\text{available} \cup \text{unmute\_set}'[\text{self}]) \\
& \wedge \text{pc}' = [\text{pc} \text{ EXCEPT } ![\text{self}] = \text{"Done"}] \\
& \wedge \text{UNCHANGED } \langle \text{overloaded}, \text{immutable}, \text{refcount}, \\
& \quad \text{rc\_barrier}, \text{required}, \text{next}, \text{acquired}, \\
& \quad \text{mutor}, \text{muting} \rangle
\end{aligned}$$

$$\begin{aligned}
\text{behaviour}(\text{self}) & \triangleq \text{Send}(\text{self}) \vee \text{Unmute}(\text{self}) \vee \text{Acquire}(\text{self}) \\
& \vee \text{Action}(\text{self}) \vee \text{Complete}(\text{self}) \\
& \vee \text{MuteMapScan}(\text{self})
\end{aligned}$$

Allow infinite stuttering to prevent deadlock on termination.

$$\begin{aligned}
\text{Terminating} & \triangleq \wedge \forall \text{self} \in \text{ProcSet} : \text{pc}[\text{self}] = \text{"Done"} \\
& \wedge \text{UNCHANGED } \text{vars}
\end{aligned}$$

$$\begin{aligned}
\text{Next} & \triangleq (\exists \text{self} \in 1 \dots \text{BehaviourCount} : \text{behaviour}(\text{self})) \\
& \vee \text{Terminating}
\end{aligned}$$

$$\begin{aligned}
\text{Spec} & \triangleq \wedge \text{Init} \wedge \square[\text{Next}]_{\text{vars}} \\
& \wedge \forall \text{self} \in 1 \dots \text{BehaviourCount} : \text{WF}_{\text{vars}}(\text{behaviour}(\text{self}))
\end{aligned}$$

$$\text{Termination} \triangleq \Diamond(\forall \text{self} \in \text{ProcSet} : \text{pc}[\text{self}] = \text{"Done"})$$

END TRANSLATION – the hash of the generated TLA code (remove to silence divergence warnings): TLA-bd24460ac3e3dbc5294ac3189bb484af