

EXTENDS *FiniteSets, Integers, Sequences, TLC*

$Null \triangleq 0$   
 $Cowns \triangleq 1..3$   
 $BehaviourLimit \triangleq 3$   
 $OverloadThreshold \triangleq 2$   
 $PriorityLevels \triangleq \{2, 1, 0\}$   
 $Min(s) \triangleq \text{CHOOSE } x \in s : \forall y \in s \setminus \{x\} : y > x$   
 $Max(s) \triangleq \text{CHOOSE } x \in s : \forall y \in s \setminus \{x\} : y < x$   
 $Range(f) \triangleq \{f[x] : x \in \text{DOMAIN } f\}$

VARIABLES *fuel, queue, scheduled, running, priority, blocker, mutor*  
 vars  $\triangleq \langle fuel, queue, scheduled, running, priority, blocker, mutor \rangle$

$Messages \triangleq \text{UNION } \{Range(queue[c]) : c \in Cowns\}$   
 $EmptyQueue(c) \triangleq Len(queue[c]) = 0$

$Init \triangleq$   
 $\wedge fuel = BehaviourLimit$   
 $\wedge queue = [c \in Cowns \mapsto \langle \{c\} \rangle]$   
 $\wedge scheduled = [c \in Cowns \mapsto \text{TRUE}]$   
 $\wedge running = [c \in Cowns \mapsto \text{FALSE}]$   
 $\wedge priority = [c \in Cowns \mapsto 0]$   
 $\wedge blocker = [c \in Cowns \mapsto Null]$   
 $\wedge mutor = [c \in Cowns \mapsto Null]$

$Terminating \triangleq$   
 $\wedge \forall c \in Cowns : EmptyQueue(c)$   
 $\wedge \text{UNCHANGED } vars$

$Acquire(cown) \triangleq$   
 $\# \text{ Preconditions}$   
 $\wedge scheduled[cown]$   
 $\wedge \neg running[cown]$   
 $\wedge \neg EmptyQueue(cown)$   
 $\wedge cown < Max(Head(queue[cown]))$   
 $\# \text{ Forward the message to the next } cown.$   
 $\wedge \text{LET}$   
 $msg \triangleq Head(queue[cown])$   
 $next \triangleq Min(\{c \in msg : c > cown\})$   
 $\text{IN}$   
 $queue' =$   
 $(next \rightarrow Append(queue[next], msg)) @@$   
 $(cown \rightarrow Tail(queue[cown])) @@$

*queue*

$\wedge$  UNCHANGED  $\langle \text{fuel}, \text{scheduled}, \text{running}, \text{priority}, \text{blocker}, \text{mutor} \rangle$

$\text{PreRun}(c) \triangleq$

# Preconditions

$\wedge \text{scheduled}[c]$

$\wedge \neg \text{running}[c]$

$\wedge \neg \text{EmptyQueue}(c)$

$\wedge c = \text{Max}(\text{Head}(\text{queue}[c]))$

# Set max *cown* in current message to running

$\wedge \text{running}' = (c :> \text{TRUE}) @ @ \text{running}$

$\wedge$  UNCHANGED  $\langle \text{fuel}, \text{queue}, \text{scheduled}, \text{priority}, \text{blocker}, \text{mutor} \rangle$

$\text{Send}(c) \triangleq$

# Preconditions

$\wedge \text{running}[c]$

$\wedge \text{fuel} > 0$

# Select set of receivers

$\wedge \exists \text{receivers} \in \{cs \in \text{SUBSET } \text{Cowns} : \text{Cardinality}(cs) > 1\} :$

# place message for receivers in the first receiver's queue

LET  $\text{next} \triangleq \text{Min}(\text{receivers})$  IN

$\text{queue}' = (\text{next} :> \text{Append}(\text{queue}[\text{next}], \text{receivers})) @ @ \text{queue}$

$\wedge \text{fuel}' = \text{fuel} - 1$

$\wedge$  UNCHANGED  $\langle \text{scheduled}, \text{running}, \text{priority}, \text{blocker}, \text{mutor} \rangle$

$\text{PostRun}(c) \triangleq$

# Preconditions

$\wedge \text{running}[c]$

# Transition

$\wedge \text{running}' = (c :> \text{FALSE}) @ @ \text{running}$

# Remove message from queue

$\wedge \text{queue}' = (c :> \text{Tail}(\text{queue}[c])) @ @ \text{queue}$

$\wedge$  UNCHANGED  $\langle \text{fuel}, \text{scheduled}, \text{priority}, \text{blocker}, \text{mutor} \rangle$

$\text{RunStep}(c) \triangleq$

$\vee \text{Acquire}(c)$

$\vee \text{PreRun}(c)$

$\vee \text{Send}(c)$

$\vee \text{PostRun}(c)$

$\text{Next} \triangleq \exists c \in \text{Cowns} : \text{RunStep}(c)$

$\text{Spec} \triangleq$

$\wedge \text{Init}$

$\wedge \Box [\text{Next} \vee \text{Terminating}]_{\text{vars}}$

$$\wedge \forall c \in \text{Cowns} : \text{WF}_{\text{vars}}(\text{RunStep}(c))$$

# Properties

# Ensure that the termination condition is reached by the model.

$$\text{Termination} \triangleq \Diamond \Box (\forall c \in \text{Cowns} : \text{EmptyQueue}(c))$$

# Invariants

# Ensure that the model produces finite messages.

$$\text{MessageLimit} \triangleq \text{Cardinality}(\text{Messages}) \leq (\text{Cardinality}(\text{Cowns}) + \text{BehaviourLimit})$$

# A message must contain at least one *cow*.

$$\text{MessagesAreNonEmpty} \triangleq \forall m \in \text{Messages} : m \neq \{\}$$

# A running *cow* must be scheduled and be the max *cow* in the message at the head of its queue.

$$\text{RunningImplication} \triangleq \forall c \in \text{Cowns} : \text{running}[c] \Rightarrow \text{scheduled}[c] \wedge (c = \text{Max}(\text{Head}(\text{queue}[c])))$$


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