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EXTENDS FiniteSets, Integers, Sequences, TLC
Null \triangleq 0
Cowns \triangleq 1 \dots 3 \# TODO: 4
MaxMessageCount \triangleq 4 \# TODO: 4
MaxMessageSize \stackrel{\triangle}{=} 3
OverloadThreshold \stackrel{\triangle}{=} 2
 PriorityLevels \stackrel{\Delta}{=} \{-1, 0, 1, 2\}
Pick(s) \stackrel{\Delta}{=} \text{CHOOSE } x \in s : \text{TRUE}
Min(s) \stackrel{\triangle}{=} CHOOSE \ x \in s : \forall y \in s \setminus \{x\} : y > x
Max(s) \triangleq \text{CHOOSE } x \in s : \forall y \in s \setminus \{x\} : y < x
Range(f) \triangleq \{f[x] : x \in DOMAIN f\}
Subsets(s, min, max) \triangleq
  \{x \in \text{SUBSET } s : (Cardinality(x) \ge min) \land (Cardinality(x) \le max)\}
RECURSIVE Concat(_)
Concat(s) \stackrel{\triangle}{=} \text{ if } s = \{\} \text{ THEN } \langle \rangle \text{ ELSE } \text{ LET } x \stackrel{\triangle}{=} Pick(s) \text{IN } x \circ Concat(s \setminus \{x\}) \}
Variables fuel, queue, scheduled, running, mutor, muted, blocker
vars \triangleq \langle fuel, queue, scheduled, running, mutor, muted, blocker \rangle
Messages \stackrel{\Delta}{=} UNION \{Range(queue[c]) : c \in Cowns\}
EmptyQueue(c) \triangleq Len(queue[c]) = 0
Overloaded(c) \triangleq Len(queue[c]) \geq OverloadThreshold
Enqueue(c, m) \stackrel{\triangle}{=} c :> Append(queue[c], m)
Dequeue(c) \stackrel{\triangle}{=} c :> Tail(queue[c])
RECURSIVE Blockers(_)
Blockers(c) \stackrel{\Delta}{=}
  IF blocker[c] = Null THEN \{\}
   ELSE \{blocker[c]\} \cup Blockers(blocker[c])
Running(c) \stackrel{\triangle}{=} \exists k \in Cowns : running[k] \land c \in Head(queue[k])
AcquiredBy(a, b) \triangleq
   \land a < b
   \land \exists c \in Cowns : a \in UNION \ Range(queue[b])
   \land b = Min(\{c \in Cowns : a \in UNION \ Range(queue[b])\})
Acquired(c) \stackrel{\Delta}{=} \exists k \in Cowns : AcquiredBy(c, k)
MutedBy(a, b) \triangleq
   \land muted[a]
   \land \exists m \in Range(queue[b]) : (b \notin m) \land (a \in m)
Init \triangleq
   \land fuel = MaxMessageCount
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\land queue = [c \in Cowns \mapsto \langle \{c\} \rangle]
  \land scheduled = [c \in Cowns \mapsto TRUE]
  \land running = [c \in Cowns \mapsto FALSE]
  \land mutor = [c \in Cowns \mapsto Null]
  \land muted = [c \in Cowns \mapsto FALSE]
  \land blocker = [c \in Cowns \mapsto Null]
Terminating \triangleq
  \land \forall c \in Cowns : EmptyQueue(c)
  ∧ UNCHANGED vars
ExternalReceive(cown) \stackrel{\Delta}{=}
  \land fuel > 0
  ∧ UNCHANGED ⟨scheduled, running, mutor, muted, blocker⟩
  \wedge fuel' = fuel - 1
   \# Receive a message from an external source
  \land \exists others \in Subsets(\{c \in Cowns : c > cown\}, 0, MaxMessageSize - 1):
    queue' = Enqueue(cown, \{cown\} \cup others) @@ queue
Acquire(cown) \triangleq
  \land scheduled[cown]
  \land \neg running[cown]
  \land \neg EmptyQueue(cown)
  \land cown \in Head(queue[cown])
  \land cown < Max(Head(queue[cown]))
  \land UNCHANGED \langle fuel, running, mutor \rangle
   # Unschedule and forward the message to the next cown.
  \wedge LET
      msg \stackrel{\triangle}{=} Head(queue[cown])
      next \stackrel{\triangle}{=} Min(\{c \in msg : c > cown\})
    \land queue' = Enqueue(next, msg) @@ Dequeue(cown) @@ queue
    \land blocker' = (cown:> next)@@blocker
    \land IF (Overloaded(cown) \lor Overloaded(next)) \land muted[next] THEN
       \land scheduled' = (next:> TRUE) @@ (cown:> FALSE) @@ scheduled
       \land muted' = (next :> FALSE) @@ muted
       ELSE
       \land UNCHANGED \langle muted \rangle
       \land scheduled' = (cown:> FALSE) @@ scheduled
Unmute(cown) \triangleq
  \land \ scheduled[cown]
  \land \neg running[cown]
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\wedge \neg EmptyQueue(cown)
   \land cown \notin Head(queue[cown])
   \land UNCHANGED \langle fuel, running, mutor, blocker <math>\rangle
    # Remove message from queue.
   \land queue' = Dequeue(cown) @@ queue
   # Reschedule muted cowns.
   \land muted' = [c \in Head(queue[cown]) \mapsto FALSE]@@muted
   \land scheduled' = [c \in Head(queue[cown]) \mapsto TRUE]@@ scheduled
PreRun(cown) \triangleq
   \land scheduled[cown]
   \land \neg running[cown]
  \wedge \neg EmptyQueue(cown)
   \wedge cown = Max(Head(queue[cown]))
   \land UNCHANGED \langle fuel, queue, scheduled, mutor, muted <math>\rangle
    # Set max cown in current message to running
   \land \mathit{running'} = (\mathit{cown} :> \mathtt{TRUE}) @@ \mathit{running}
   \land blocker' = [c \in Head(queue[cown]) \mapsto Null] @@ blocker
Send(cown) \triangleq
   \land running[cown]
   \wedge \mathit{fuel} > 0
   \land UNCHANGED \langle running, blocker \rangle
   \wedge fuel' = fuel - 1
   # Select set of receivers
   \land \exists receivers \in Subsets(Cowns, 1, MaxMessageSize):
    LET
       next \stackrel{\triangle}{=} Min(receivers)
       senders \stackrel{\triangle}{=} Head(queue[cown])
       mutors \stackrel{\triangle}{=} \{c \in receivers : Overloaded(c)\}
    IN
      # Place message for receivers in the first receiver's queue.
     \land queue' = Enqueue(next, receivers) @@ queue
     \land IF (Overloaded(cown) \lor Overloaded(next)) \land muted[next] THEN
        \land UNCHANGED \langle mutor \rangle
        # Unmute next.
        \land scheduled' = (next:> TRUE) @@ scheduled
        \land muted' = (next :> FALSE) @@ muted
        ELSE
        \land UNCHANGED \langle scheduled, muted \rangle
        \# Set mutor if any receiver is overloaded and there are no receivers in the set of senders.
        \wedge IF
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\land mutors \neq \{\}
          \land mutor[cown] = Null
          \land (senders \cap receivers) = \{\}
          THEN mutor' = (cown:> Min(mutors)) @@ mutor
          ELSE UNCHANGED \langle mutor \rangle
PostRun(cown) \triangleq
  \land running[cown]
  \land UNCHANGED \langle fuel, blocker \rangle
  \land running' = (cown :> FALSE) @@ running
  \land mutor' = (cown :> Null) @@ mutor
  \wedge \text{ LET } msg \stackrel{\triangle}{=} Head(queue[cown])IN
      # Mute if mutor is set and no running cowns are overloaded.
     \land IF (mutor[cown] \neq Null) \land (\forall c \in msg : \neg Overloaded(c)) THEN
       \land muted' = [c \in msg \mapsto TRUE] @@ muted
       \land scheduled' = [c \in msg \mapsto FALSE] @@ scheduled
        \# Send unmute message to mutor
       \land queue' = Enqueue(mutor[cown], msg) @@ Dequeue(cown) @@ queue
       ELSE
       \land UNCHANGED \langle muted \rangle
       \land scheduled' = [c \in msg \mapsto TRUE] @@ scheduled
       \land queue' = Dequeue(cown) @@ queue
RunStep(cown) \triangleq
    ∨ ExternalReceive(cown) \*# Very expensive check
   \vee Acquire(cown)
   \vee Unmute(cown)
  \vee PreRun(cown)
  \vee Send(cown)
  \vee PostRun(cown)
Next \stackrel{\Delta}{=} \exists c \in Cowns : RunStep(c)
Spec \triangleq
  \land Init
  \wedge \Box [Next \vee Terminating]_{vars}
  \land \forall c \in Cowns : WF_{vars}(RunStep(c))
 # Properties
 # Ensure that the termination condition is reached by the model.
Termination \triangleq \Diamond \Box (\forall c \in Cowns : EmptyQueue(c))
 # Invariants
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# Ensure that the model produces finite messages.

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# Cowns are acquired by one running message at a time.
UniqueAcquisition \triangleq
  LET msgs \triangleq Concat(\{\langle Head(queue[c]) \rangle : c \in \{k \in Cowns : running[k]\}\})
       Cardinality(Range(msgs)) = Len(msgs)
 # Each queue has at most one token message.
LoneToken \stackrel{\triangle}{=} \forall c \in Cowns : Len(SelectSeq(queue[c], LAMBDA m : m = \{\})) \leq 1
 \# A running cown must be scheduled and be the max cown in the message at the head of its queue.
RunningImplication \stackrel{\triangle}{=} \forall c \in Cowns : running[c] \Rightarrow
   \land scheduled[c]
   \wedge c = Max(Head(queue[c]))
   \land \, \forall \, k \in \mathit{Head}(\mathit{queue}[\mathit{c}]) : (k < \mathit{c}) \Rightarrow \mathit{AcquiredBy}(k, \, \mathit{c})
 \# An acquired cown is not scheduled.
AcquiredImplication \stackrel{\triangle}{=} \forall c \in Cowns : Acquired(c) \Rightarrow
   \land \neg scheduled[c]
 # A muted cown is not scheduled or running.
MutedImplication \stackrel{\triangle}{=} \forall c \in Cowns : muted[c] \equiv
   \land \exists k \in Cowns : MutedBy(c, k)
   \land \neg scheduled[c]
   \wedge \neg Running(c)
 \# A muted cown exists in an unmute message in the queue of at least one mutor.
MutedOnce \triangleq
  \forall m \in \{c \in Cowns : muted[c]\}:
     Cardinality(\{c \in Cowns : MutedBy(m, c)\}) > 0
 # A cown may only be acquired by one message.
AcquiredOnce \triangleq
  \forall a \in \{c \in Cowns : Acquired(c)\}:
     Cardinality(\{c \in Cowns : AcquiredBy(a, c)\}) = 1
 # An acquired cown is acquired by a cown in its blocker set.
AcquiredByBlocker \stackrel{\Delta}{=} \forall \langle a, b \rangle \in Cowns \times Cowns :
  AcquiredBy(a, b) \Rightarrow b \in Blockers(a)
 # An overloaded cown doesn't exist in a muted cown's queue.
OverloadedNotInMutedQueue \stackrel{\triangle}{=} \forall \langle o, m \rangle \in Cowns \times Cowns :
  Overloaded(o) \land muted[m] \Rightarrow o \notin UNION \ Range(queue[m])
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 $MessageLimit \stackrel{\triangle}{=} Cardinality(Messages) < (Cardinality(Cowns) + MaxMessageCount)$ 

 $^{+}$  https://github.com/tlaplus/Examples/blob/master/specifications/TransitiveClosure/TransitiveClosure.tla#L114

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TC(R) \stackrel{\triangle}{=} LET S \stackrel{\triangle}{=} \{r[1]: r \in R\} \cup \{r[2]: r \in R\} RECURSIVE TCR(-) TCR(T) \stackrel{\triangle}{=} IF T = \{\} THEN R ELSE LET r \stackrel{\triangle}{=} \text{CHOOSE } s \in T \text{: TRUE } RR \stackrel{\triangle}{=} TCR(T \setminus \{r\}) IN RR \cup \{\langle s, t \rangle \in S \times S : \langle s, r \rangle \in RR \wedge \langle r, t \rangle \in RR\} IN TCR(S) CylcicTransitiveClosure(R(-, -)) \stackrel{\triangle}{=} \text{LET } s \stackrel{\triangle}{=} \{\langle a, b \rangle \in Cowns \times Cowns : R(a, b)\} IN \exists c \in Cowns : \langle c, c \rangle \in TC(s)
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