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CSCI 532 Problem 1-1

Collaborators: *Dana Parker*

1. Write this homework in LaTeX. (You can use this document as a starting point!)
2. Update your photo on D2L to be a recognizable headshot of you.

Done!

CSCI 532 Problem 1-2Collaborators: *Dana Parker*

Consider the third function defined in EPI, Section 13.1 (Compute the Intersection of Two Sorted Arrays).

1. When we design an algorithm, we design the algorithm to solve a problem or answer a question. What is the problem that this algorithm solves?

This while loop solves the problem of finding the intersection of two arrays A, B . That is to say, that a new array "intersection_A_B" is created and contains the elements present in both array A and array B

2. Prove that the while loop terminates using a decrementing function.

Proof. Let there be a decrementing function $D(x)$ such that x equals a possible state of the while loop. Let $x = (\text{len}(A) + \text{len}(B)) - (i + j)$

Note: the function $D(x)$ is a valid decrementing function for this while loop because it includes the conditions for termination present within the while loop, $i < \text{len}(A)$ and $j < \text{len}(B)$. therefore, if i, j increment, the function $D(x)$ decrements. If x reaches 0 it must mean that either i or j or both have increased to a point where they are greater than $\text{len}(A)$ or $\text{len}(B)$.

Now, we must show that either i or j , or both, increment during each iteration of the loop.

If we look at the while loop. We notice that there are 3 cases that can occur. An "if", an "elif", and an "else".

Case 1: If the "if" statement is triggered, $i = i + 1$ (i increments) and $j = j + 1$ (j increments).

Case 2: If the "elif" statement is triggered, $i = i + 1$ (i increments)

Case 3: If the "else" statement is triggered, $j = j + 1$ (j increments)

Therefore, since all possible cases within the while loop increment either i or j or both, the decrementing function must reach 0. Thus, the while loop must terminate. \square

CSCI 532 Problem 1-3

Collaborators: *Dana Parker*

Prove the following statement: Every tree with one or more nodes/vertices has exactly $n - 1$ edges.

Proof. Let n be the number of vertices in a tree (T).

Base Case: If $n = 1$, then the number of edges is 0, because there are no other nodes to connect to.

Assume: Assume that $n=k$ holds: A tree with k nodes has $k - 1$ edges.

Show: We must show that $n = k+1$ holds: A tree $k+1$ nodes has k edges

Now, let V_1 and V_2 be vertices of a tree (M) that has k and let E be the edge connecting them. Since M is a tree, then there must be only one path, or edge, connecting vertices V_1, V_2 . Thus, if edge E is deleted, there are now two independent trees containing the vertices V_1, V_2 respectively. Let these two trees be G_1, G_2 respectively. It can be assumed that G_1, G_2 have no cycle, since M had no cycles.

Let k_1 and k_2 be the number of vertices in G_1, G_2 respectively such that $k_1 + k_2 = n$. therefore, because of the properties of addition $k_1 < k$ and $k_2 < k$ Thus, based on the inductive hypothesis, G_1, G_2 must have $k_1 - 1$ and $k_2 - 1$ nodes respectively.

therefore, the number of edges in $M = k$

$$= (k_1 - 1) + (k_2 - 1) + 1$$

$$= k_1 + k_2 - 1$$

$$= k - 1$$

$$k_1 + k_2 = k \text{ by substitution}$$

□

CSCI 532 Problem 1-4Collaborators: *Dana Parker*

Consider the following statement: If a and b are both odd numbers, then ab is an odd number.

1. What is the definition of an odd number?

An odd number can be defined as $k = 2n + 1$ where $n \in \mathbb{Z}$

2. What is the definition of an even number?

An even number can be defined as $k = 2n$ where $n \in \mathbb{Z}$

3. What is the contrapositive of this statement?

The contrapositive of this statement is, if ab is an even number, then a and b are not both odd.

4. Prove this statement, using the contrapositive.

Proof. Notice that there are 3 possible cases for this statement

Case 1: a, b are both even, where $a, b \in \mathbb{Z}$

Let $a = 2n$ by definition of an even integer
 $b = 2m$ by definition of an even integer

$$\begin{aligned} ab &= (2n)(2m) \\ &= 2(mn) && \text{Let } k = mn, \text{ where } k \text{ is an integer by the multiplicative properties of integers} \\ &= 2k \end{aligned}$$

ab can be written in the form $2k$ where k is an integer. Therefore, by the definition of an even integer, ab is even. Thus, the statement, if ab is an even number, then a and b are not both odd, is true for this case.

Case 2: a or b is even, but not both. where $a, b \in \mathbb{Z}$

Let $a = 2n$ by definition of an even integer
 $b = 2m + 1$ by definition of an odd integer

$$\begin{aligned} ab &= (2n)(2m + 1) \\ &= 4(mn) + 2n \\ &= 2(2mn + n) && \text{Let } k = 2mn + n, \text{ where } k \text{ is an integer by the multiplicative and additive properties of integers} \\ &= 2k \end{aligned}$$

ab can be written in the form $2k$ where k is an integer. a, b are arbitrary integers, so this proof is valid for the case where a is odd and b is even, and the case where a is even and b is odd. Therefore, by the definition of an even integer, ab is even. Thus, the statement, if ab is an even number, then a and b are not both odd, is true for this case.

Case 3: a, b are both odd

Let $a = 2n + 1$ by definition of an odd integer

$b = 2m + 1$ by definition of an odd integer

$$ab = (2n + 1)(2m + 1)$$

$$= 4mn + 2n + 2m + 1$$

$$= 2(2mn + n + m) + 1 \quad \text{Let } k = 2mn + n + m, \text{ where } k \text{ is an integer by the multiplicative and additive properties of integers.}$$

$$= 2k + 1$$

ab can be written in the form $2k + 1$ where k is an integer. Therefore, by the definition of an odd integer, ab is odd. Thus, the statement, if ab is an even number, then a and b are not both odd, is false for this case.

Having shown that for every case, except where a and b are both odd, that ab is even, we have shown that the contrapositive is true. Therefore, by proof by contraposition, the original statement, if a and b are both odd numbers, then ab is an odd number, is true.

□