

★ Quick Sort:

→ Quick sort is a divide and conquer sort. And it is also recursive sort.

→ In a quick sort we define three pointers.

① Pivot

② p

③ q

→ Pivot stand for center point.

→ q is used to find small elements with help of pivot.

→ ~~for~~ conditions:

→ while scanning (left to right and right to left) if "p" pointer and "q" pointer crossing each other then "Pivot" and "q" replace their position.

→ If not ~~crossing~~ p and q pointer are not crossing each other then simply "p and q" pointers change position respectively.

Ex → Perform quick sort on following data:
X = 35, 50, 15, 25, 80, 20, 90, 45

→ PASS-I:

35, 50, 15, 25, 80, 20, 90, 45
pivot p q

→ we set ^{1st} element 35 as pivot.

→ no element less than pivot and no element greater than pivot

→ Now with the help of pivot and p pointer we scan (check) Big element with comparing pivot and p.

→ As it is with the help of comparing pivot and q we find the small elements.

a) 35, 50, 15, 25, 80, 20, 90, 45
pivot p q

→ There is p and q not crossing each other we change the position.

b) 35, 20, 15, 25, 80, 50, 90, 45
pivot p q

→ we find the big element.

c) 35, 20, 15, 25, 80, 50, 90, 45
pivot p q

→ There is p and q pointers are crossing each other we change the position of pivot and q.

d) 25, 20, 15, 35, 80, 50, 90, 45

→ PASS-II :

b) $\begin{array}{c} \text{---} 25, 20, 15 \\ \text{pivot } p = 9 \end{array}$ $\begin{array}{c} 35 \quad 80, 50, 90, 45 \\ \text{pivot } p = 9 \end{array}$

→ In left sub array p and q pointers are crossing each other we change the position of pivot and q.

→ In right sub array p and q pointers are not crossing we change position of p and q.

c) $\begin{array}{ccc} & 35 & \\ \hline 15, 20, 25 & 80, 50, 45, 90 & \\ & \text{pivot} & P \end{array}$

→ In right sub array p and q pointers are crossing each other we change the position of pivot and q.

d) 15, 20, 25, 35, 45, 50, 80, 90

Now the array is sorted.

Time and Space complexity :

Space complexity :

→ Space complexity of an algorithm represents the amount of memory space required by the algorithm in its life cycle.

→ The space required by an algorithm is equal to the sum of the following two components:

(a) A fixed part that is a space required to store certain data and variables, that are independent of the size of the problem. For example, simple variables and constants used, program size, etc....

(b) A variable part is a space required by variables whose size depends on

dynamic memory allocation, recursion stack space, etc. . .

→ Space complexity of any algorithm ^{is} ~~SP(I)~~ $S(P)$ •
 $S(P) = C + SP(I)$, where C is the fixed part
and $S(I)$ is the variable part of algorithm.
 I is instance characteristic.

Ex: Algorithm : $\text{Sum}(A, B)$

Step-1 :- Start

Step-2 :- $C \leftarrow A + B + 10$

Step-3 :- Stop

→ Time complexity:

→ Time complexity of an algorithm represents the amount of time required by the algorithm to run to completion. Time requirements can be defined as a numerical function $T(n)$, where $T(n)$ can be measured as the number of steps, provided each step consumes constant time.