# halcheck

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## halcheck — Overview

Motivation

Design

 ${\sf Summary}$ 

# **Motivation**

### halcheck — Motivation

Why halcheck? (and why not RapidCheck)?

- 1. Simpler API
- 2.
- 3.

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### A central generator data type:

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#### A set of basic combinators:

```
choose :: (Int, Int) \rightarrow Gen Int suchThat :: (a \rightarrow Bool) \rightarrow Gen a \rightarrow Gen a frequency :: [(Int, Gen a)] \rightarrow Gen a ...
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```

- Users must be comfortable reasoning about higher-order functions.
- Users must know the distinction between generator code and non-generator code.

**Example:** Write a generator combinator that produces std::vectors up to (but not including) a given length.

```
Gen<std::vector<int>> example(int N) {
   return gen::container<std::vector<int>>(
    *gen::inRange(0, N),
      gen::arbitrary<int>);
}
```

**Example:** Write a generator combinator that produces std::vectors up to (but not including) a given length.

```
Gen<std::vector<int>> example(int N) {
  return gen::container<std::vector<int>>(
    *gen::inRange(0, N), // WRONG
    gen::arbitrary<int>);
}
```

**Problem:** \*gen::inRange(0, N) is evaluated *before* example returns, resulting in a generator that always produces std::vectors of the same length.

**Example:** Write a generator combinator that produces std::vectors up to (but not including) a given length.

```
Gen<std::vector<int>> example(int N) {
  return gen::exec([=] {
    return gen::container<std::vector<int>>(
         *gen::inRange(0, N),
         gen::arbitrary<int>);
  });
}
```

**Solution:** Need to delay computation of \*gen::inRange(0, N) using gen::exec.

- Generators should only be invoked (operator \*) in the correct context.
- Haskell's type system completely rules out this kind of error.
- C++'s type system is not powerful enough!
- Problem goes away if we simply get rid of the generator type and just write first-order functions!

```
std::vector<int> example(int N) {
  // container(int, () → int) → std::vector<int>
  return gen::container<std::vector<int>>(
    gen::range(0, N),
    gen::arbitrary<int>);
}
```

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Why halcheck? (and why not RapidCheck)?

- 1. Simpler API
- 2. Support for test-case generation strategies
- 3.

There are various desirable strategies for generating data:

- Exhaustive (SmallCheck/LeanCheck)
- Learning-based (RLCheck)
- Coverage-guided (FuzzTest)
- Random (almost everything)

Most PBT frameworks (including RapidCheck) use a fixed strategy.

halcheck provides combinators for specifying strategies:

```
// random(int) → strategy
// v Executes random test cases forever or until a bug is found.
test::random(seed)([] { run_test(); });
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// limited(strategy, int) → strategy
// v Executes at most 100 test cases.
test::limited(test::random(), 100)([] { run_test(); });
```

halcheck provides combinators for specifying strategies:

```
// random(int) → strategy
// ▼ Executes random test cases forever or until a bug is found.
test::random(seed)([] { run test(); });
// limited(strategy, int) → strategy
// ▼ Executes at most 100 test cases.
test::limited(test::random(), 100)([] { run test(); });
// shrinking(strategy) → strategy
// ¬ Performs test-case shrinking if a bug is found.
test::shrinking(test::random())([] { run test(); });
```

### halcheck — Motivation

Why halcheck? (and why not RapidCheck)?

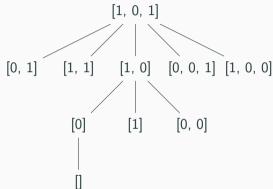
- 1. Simpler API
- 2. Support for test-case generation strategies
- 3. Better space complexity

### How does shrinking work?

Internally, every generator returns a "shrink tree" of values.

Shrink trees can be **very large** so they must be computed lazily.

### Shrink tree for a list:

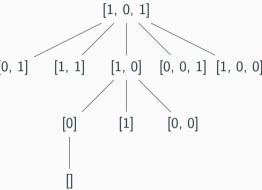


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Internally, every generator returns a "shrink tree" of values.

Shrink trees can be **very large** so they must be computed lazily.

### Shrink tree for a list:



This implementation strategy does not work for C++!

### **Example:**

```
auto xs = *gen::arbitrary<std::vector<int>>();
auto x = *gen::elementOf(xs);
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- gen::elementOf must store a copy of xs in order to avoid creating a dangling reference.

**Conclusion:** all combinators must make copies of their arguments!

**Problem:** by default, copies in C++ are deep  $(\mathcal{O}(n))$  instead of  $\mathcal{O}(1)$ .

#### Generators cannot return references:

```
// Generates a random reference
// to an element of xs.
rc::Gen<int &> referenceOf(??? xs):
// Example: assign a
// random element to 0.
*referenceOf(xs) = 0:
```

#### **Generators cannot return references:**

```
// Generates a random reference
// to an element of xs.
rc::Gen<int &> referenceOf(??? xs);

// Example: assign a
// random element to 0.
*referenceOf(xs) = 0;

What type should referenceOf have?
```

#### exec arguments must capture by value:

```
auto bound = *gen::arbitrary<int>();
auto ranges = *gen::container<...>(
  gen::exec([&] {
    // bound does not exist
    // during shrinking!
    auto min = *gen::inRange(0, bound);
    auto max = *gen::inRange(min, bound);
    return std::make_pair(min, max);
}));
```

- halcheck never makes copies of generated data.
- halcheck only records decision made during the generation process.
- All the previous examples work in halcheck without any unnecessary copying!

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# Design

#### Goals:

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- First-order API (as much as possible)
- C++ 11 support
- Thread safety

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- Support user-defined generation strategies
- First-order API (as much as possible)

To support these goals, every generator needs to be built from a set of **overridable core functions**:

gen::next

gen::discard

gen::shrink

■ gen::group

## halcheck — Design — gen∷next

gen::next(int 
$$x = 1$$
, int  $y = 1$ )  $\rightarrow$  bool

**Intuition:** gen::next(x, y) performs a biased coin flip with  $p = \frac{y}{x+y}$ .

- This is our *only* source of randomness.
- Is this enough to define all generators?

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

**Example 1:** generate a number in the range [a, b).

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```

**Example 1:** generate a number in the range [a, b).

 Solution: use binary search, replace comparison with biased coin flip.

```
int range(int a, int b) {
  if (a + 1 >= b)
    return a;

int c = std::midpoint(a, b);
  if (gen::next(c - a, b - c))
    return range(mid, b);
  else
    return range(a, mid);
}
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## Example 2: generate a enum dim.

```
enum class dim {
   x = 0,
   y = 3,
   z = 8
};
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

#### **Example 2:** generate a **enum dim**.

 Solution 1: consult gen::range to determine which value to return.

```
enum class dim {
  x = 0,
  v = 3,
  z = 8
};
dim arbitrarv(gen::tag<dim>) {
  switch (gen::range(0, 3)) {
  case 0: return dim::x:
  case 1: return dim::y;
  case 2: return dim::z:
  } }
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 2:** generate a **enum dim**.

- Solution 1: consult gen::range to determine which value to return.
- Solution 2: define gen::element (implemented via gen::range).

```
enum class dim {
    x = 0,
    y = 3,
    z = 8
};
dim arbitrary(gen::tag<dim>) {
    return gen::element({x, y, z});
}
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

**Example 3:** generate a **struct vec**.

```
struct vec {
   int x;
   int y;
};
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 3:** generate a **struct vec**.

Trivial.

```
struct vec {
   int x;
   int y;
};

vec arbitrary(gen::tag<vec>) {
   vec output;
   output.x = gen::arbitrary<int>();
   output.y = gen::arbitrary<int>();
   return output;
}
```

gen::next(int 
$$x = 1$$
, int  $y = 1$ )  $\rightarrow$  bool

**Example 4:** generate a std::string.

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 4:** generate a std::string.

 Attempt: generate random size n, then generate n elements.

```
std::string gen_string() {
  std::string xs;
  auto size = gen::arbitrary<size_t>();
  while (size-- > 0)
     xs.push_back(gen::arbitrary<char>());
  return xs;
}
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 4:** generate a std::string.

- Attempt: generate random size n, then generate n elements.
- Problem: gen::arbitrary<size\_t> is uniformly distributed — high chance of generating billions of GBs

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std::string gen_string() {
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   while (size-- > 0)
       xs.push_back(gen::arbitrary<char>());
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}
```

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gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 4:** generate a std::string.

 Attempt: add elements as long as gen::next returns true.

```
std::string gen_string() {
  std::string xs;
  while (gen::next())
     xs.push_back(gen::arbitrary<char>());
  return xs;
}
```

```
gen::next(int x = 1, int y = 1) \rightarrow bool
```

## **Example 4:** generate a std::string.

- Attempt: add elements as long as gen::next returns true.
- Problem: distribution biased towards shorter lists.

```
std::string gen_string() {
  std::string xs;
  while (gen::next())
     xs.push_back(gen::arbitrary<char>());
  return xs;
}
```

**Problem:** how do we define gen\_string with a sensible distribution?

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**Standard Solution:** define a incrementally increasing size parameter and choose lengths uniformly up to the current size.

```
std::string gen_string() {
   std::string xs;
   auto size = gen::range(0, gen::size());
   while (size-- > 0)
      xs.push_back(gen::arbitrary<char>());
   return xs;
}
```

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- gen::size cannot be defined in terms of gen::next.
- Do we really need an extra combinator?

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   while (size-- > 0)
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   return xs;
}
```

- gen::size cannot be defined in terms of gen::next.
- Do we really need an extra combinator? NO!

```
gen::next(gen::weight x = 1, gen::weight y = 1) \rightarrow bool
```

- Actual definition of gen::next takes gen::weight parameters.
- A gen::weight is a symbolic value parameterized by an unknown size.
- A gen::weight is semantically a function in  $\mathbb{N} \to \mathbb{N}$ .
- Bonus: generators automatically behave uniformly with respect to size!

```
std::string gen_string() {
   std::string xs;
   auto size = gen::weight::current; // current is the symbolic variable
   while (gen::next(1, size--))
        xs.push_back(gen::arbitrary<char>());
   return xs;
}
```

Proof of correctness left as an exercise to the reader ©

Summary: gen::next alone allows us to write generators for

- weighted distributions,
- enums and unions (i.e., sum types),
- structs (i.e., product types), and
- recursive data structures via sized generation.

# halcheck — Design — gen∷discard

```
gen::discard()
```

**Intuition:** gen::discard() terminates the current test case (no failure).

- Used to implement precondition checking (RC\_PRE).
- Also useful for implementing rejection sampling (gen::suchThat).

# halcheck — Design — gen∷shrink

```
gen::shrink(int size = 1) → std::optional<int>
```

**Intuition:** gen::shrink(n) returns an integer if shrinking should occur at the call site, or std::nullopt otherwise.

```
gen::shrink(int size = 1) → std::optional<int>
```

## **Example:**

#### **Behaviour:**

```
std::string gen string shrink() {
 auto xs = gen string();
 for (size t i = 0; i < xs.size();) {</pre>
   if (auto c = gen::shrink(2))
     if (*c == 0) xs.erase(i);
     else xs[i] = ' 0':
   else
     i++:
  return xs;
```

```
gen::shrink(int size = 1) → std::optional<int>
```

## **Example:**

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std::string gen string shrink() {
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#### **Behaviour:**

gen::shrink(n) returns
 std::nullopt until a test
 case failure occurs.

```
gen::shrink(int size = 1) → std::optional<int>
```

## **Example:**

```
std::string gen string shrink() {
 auto xs = gen string();
 for (size t i = 0; i < xs.size();) {
    if (auto c = gen::shrink(2))
      if (*c == 0) xs.erase(i);
     else
                xs[i] = ' \setminus 0':
    else
      i++:
  return xs;
```

#### Behaviour:

- 1. gen::shrink(n) returns
   std::nullopt until a test
   case failure occurs.
- On subsequent test cases, gen::shrink(n) returns an integer < n indicating which value to shrink to.

# halcheck — Design — gen::group

```
gen::group()
```

## **Example:**

```
int x = 0;
if (gen::next() && // a
    !gen::shrink())
    x += gen::next() ? 1 : 0; // b
x += gen::next() ? 1 : 0; // c
```

- Suppose a test case failure occurs with
   a == b == true and c == false.
- Then !gen::shrink() subsequently evaluates to false and b is not evaluated.
- We expect that c should still evaluate to false.

## halcheck — Design — gen∷group

```
gen::group()
```

#### **Example:**

```
int x = 0;
if (gen::next() && // a
    !gen::shrink())
    x += gen::next() ? 1 : 0; // b
x += gen::next() ? 1 : 0; // c
```

- Problem: halcheck is just a library. It cannot inspect the internal structure of a program!
- Internally, halcheck only saves and replays the results of gen::next.
- Sequence 1: {a: true, b: true, c: false}
- Sequence 2: {a: true, c: true, [unused]: false}

# halcheck — Design — gen::group

```
gen::group()
```

# Example:

```
int x = 0;
{
   auto _ = gen::group();
   if (gen::next() && // a
        !gen::shrink())
        x += gen::next() ? 1 : 0; // b
}
x += gen::next() ? 1 : 0; // c
```

Intuition: gen::group informs halcheck
that all subsequent calls to gen::next (and
gen::shrink) should be treated as a single
"step".

- Sequence 1:
  {{a: true, b: true}, c: false}
- Sequence 2:
  {{a: true, [unused]: true}, c: false}

All generators are implemented using the core functions.

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- How are the core functions themselves implemented?

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- How are the core functions themselves implemented? It depends!
- Core functions can be **overriden**.

```
Example: Overriding gen::next
CHECK THROWS(gen::next()); // Throws by default
 // ¬ gen::next calls the lambda as long as this is in scope
 auto = gen::next.handle([](int x, int y) {
    CHECK THROWS(gen::next()); // Calling gen::next within the handler
                               // invokes the previous behaviour.
    return x == 0 \&\& y > 0;
 });
 CHECK EQ(gen::next(0, 1), true);
CHECK THROWS(gen::next()); // Original behaviour restored
```

halcheck's core functions are implemented using a primitive system of **effect** handlers.

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  - This property is what makes strategies composable.
- Effects are thread local.

**Summary** 

# halcheck — Design — Summary

- Every generator in halcheck is built from four first-order core functions, enabling a "direct-style" API.
- User-defined generation strategies are supported: every core function can be overriden.
- Shrinking is performed by changing generator inputs (gen::shrink) instead of generator outputs, resulting in less memory usage.