Topological Sort

Ju-Won Seo 2019.08.21

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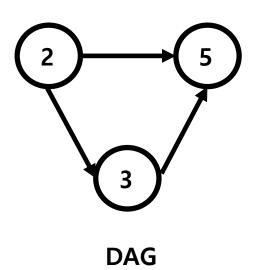
1. Introduction

What is Topological sort?

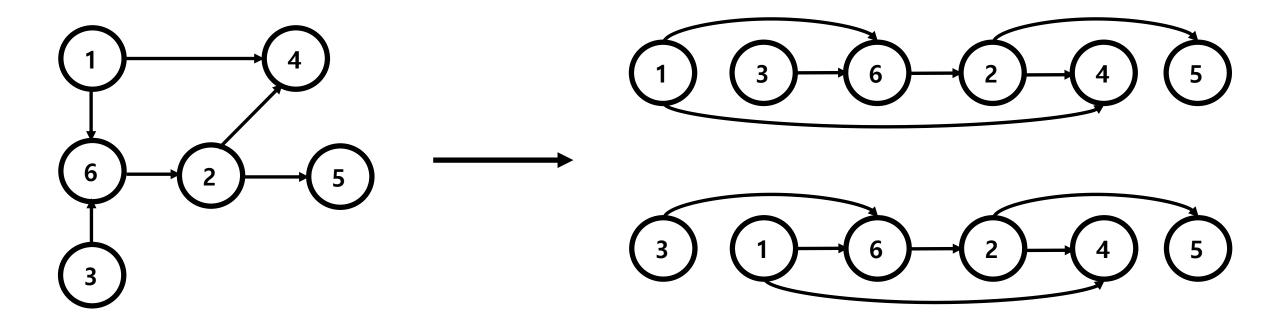
- Listing all vertices without violating the linear ordering of each vertex.

- The graph is a directed acyclic graph called 'DAG'.

- After sorting, we can view a horizontal line so that all directed edges go from left to right.



What is Topological sort?



Applications

- To indicate precedence among events.

1. A project that must be completed before each task can be completed.

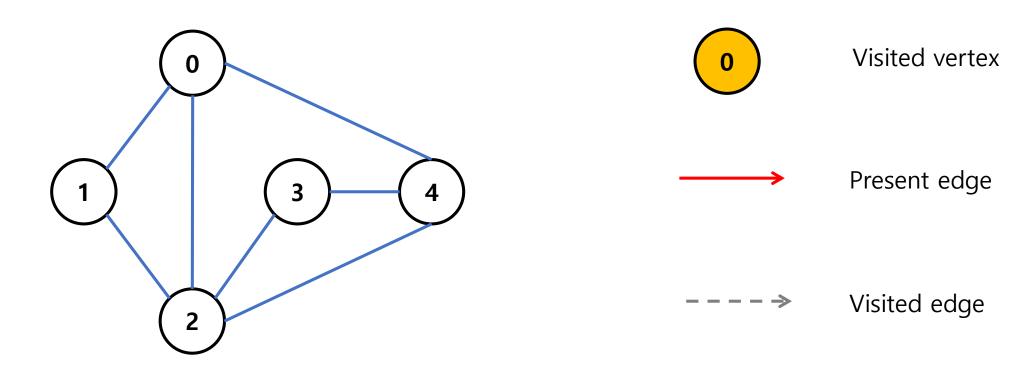
2. Prerequisite subject.

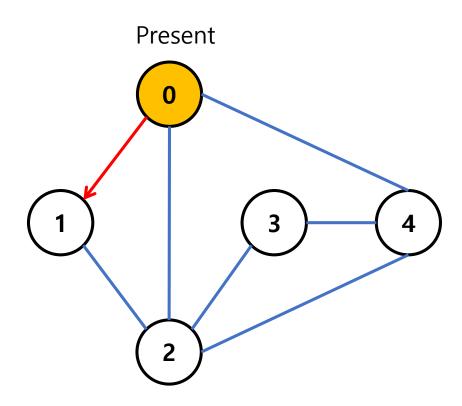
2. DFS, BFS

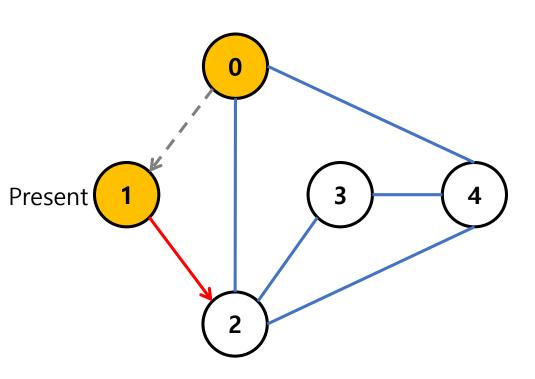
- DFS (Depth First Search)

- DFS has a form of recursive algorithm.
- If you want to visit all nodes, you can select this method.
- Must be checked whether any nodes have been visited. If you not, there is a risk into an infinite loop.

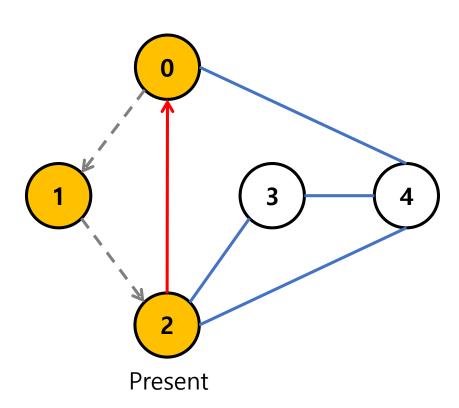
```
1 void DFS(Node v){
       if(v == null) return;
       // 1. visit node
       v.visited = true; //check visited node
       // 2. visit all adjacent vertices.
 9
       for each (Node n in v.adjacent){
           if(n.visited==false){
10
               DFS(n); // 3. recursive DFS
11
12
13
       }
```



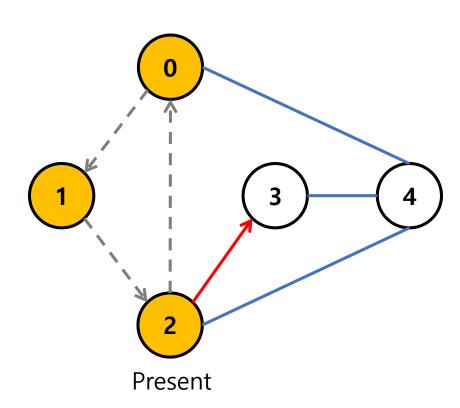


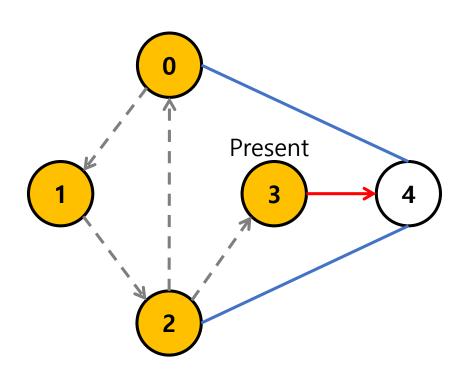


0	1		

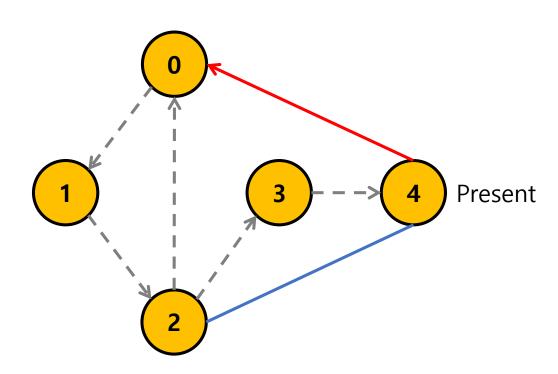


0	1	2	
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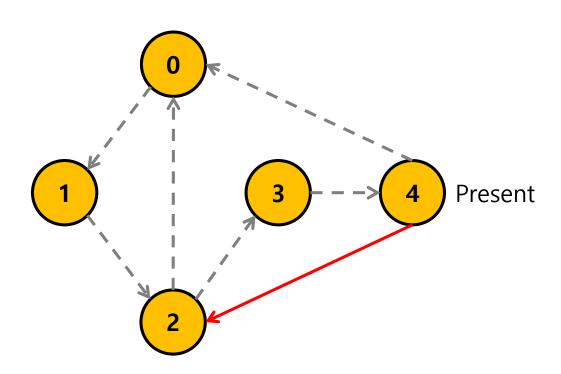




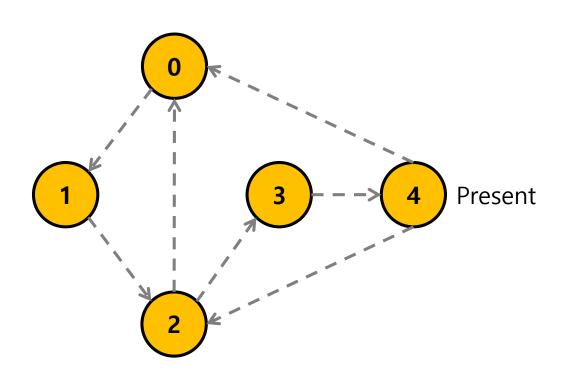
0 1 2 3

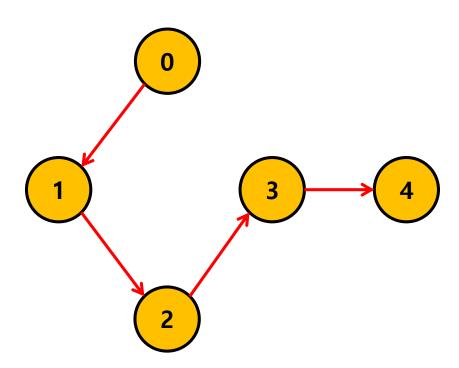


0	1	2	3	4



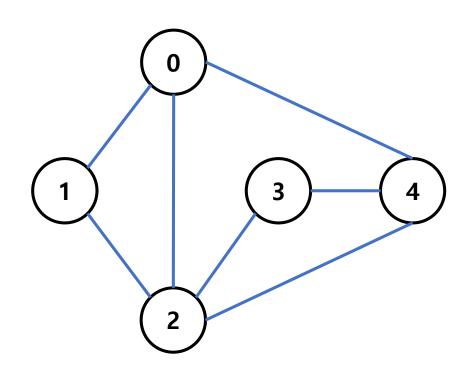
0 1 2 3 4	
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0 1 2 3 4	
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- Analysis

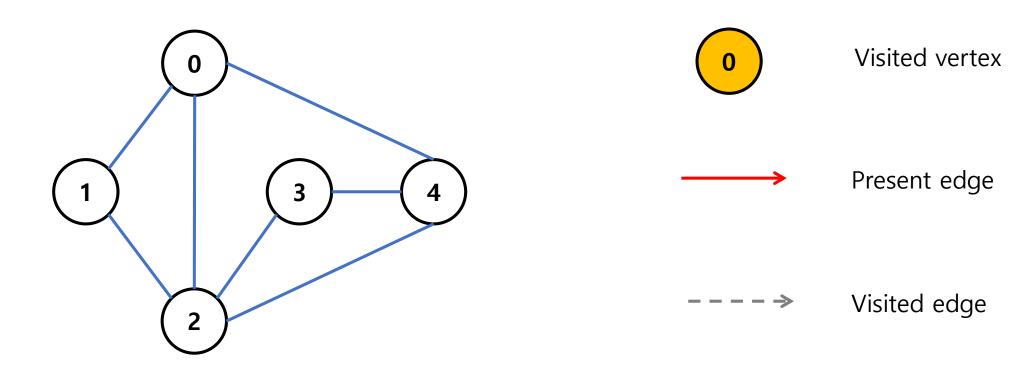


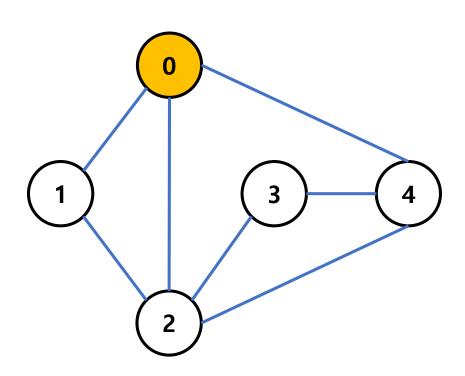
Adjacent list: O(V + E)(V is vetex and E is edge)

- BFS (Breadth First Search)

- It visits vertices near the starting vertex first.
- When you want to find the shortest path, you can choose this method.
- It used data structure called 'Queue'
- Must be checked whether any nodes have been visited. If you not, there is a risk into an infinite loop.

```
1 void BFS(Node* start){
       q.push(satart); // push the start vertex in queue
       //repeat until queue is empty
       while(!q.empty()){
           node* tmp = q.front(); // pop first vertex in queue
           q.pop();
           tmp.visited = true; // checked present vertex
10
11
12
           //visit all adjacent vertices of tmp vertex
           for each (node v in tmp.adjacent){
13
               if(v.visited == false){
14
                   v.visited = true; // checked v vertex
15
16
                   q.push(v); // push the v vertex in queue
17
18
19
20 }
```

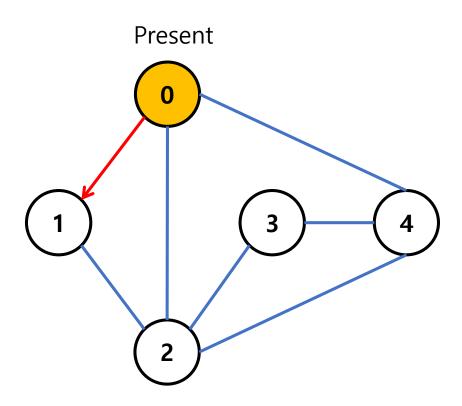


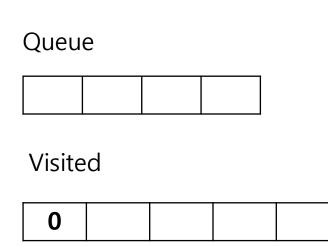


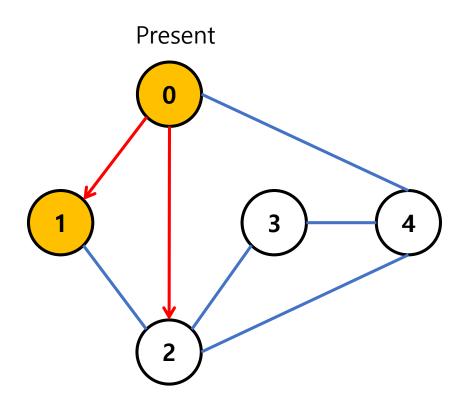
Queue







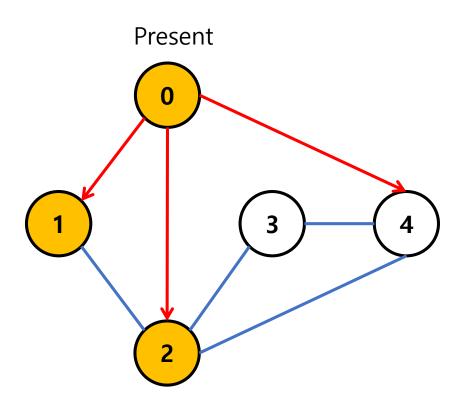




Queue

1



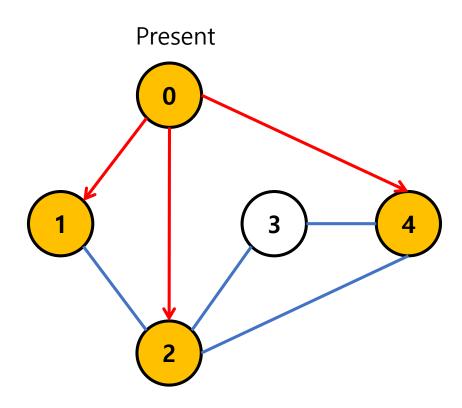


Queue

1 2

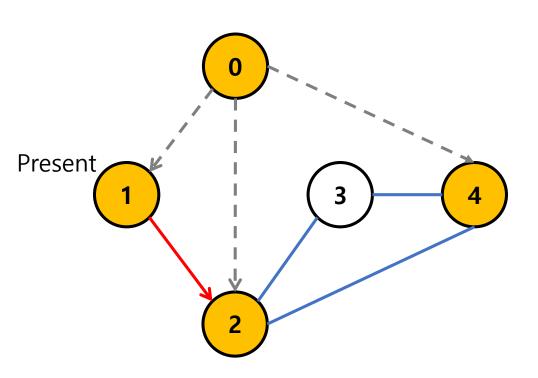
Visited

0 1 2



Queue

1	2	4	

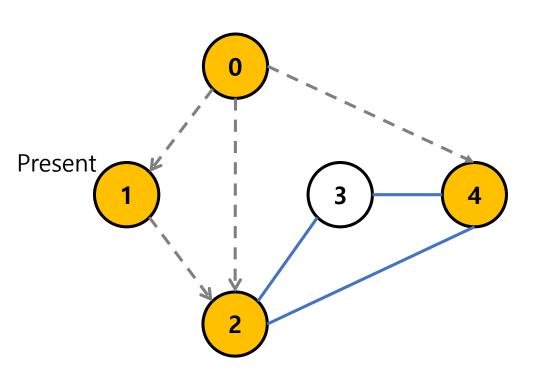


Queue

2 4

Visited

0 1 2 4

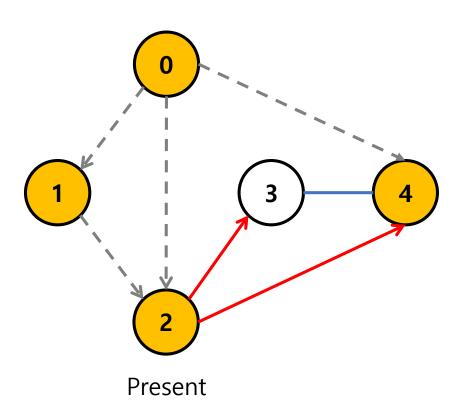


Queue

2 4

Visited

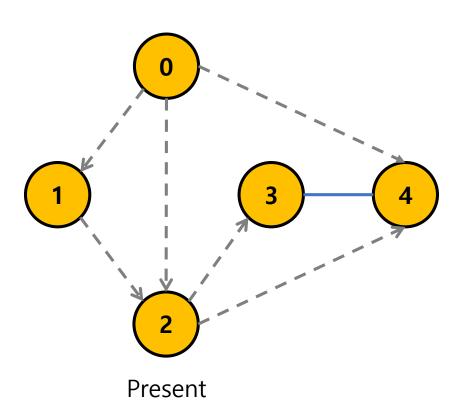
0 1 2 4



Queue

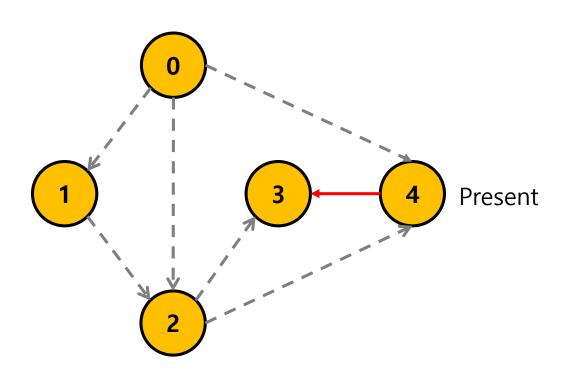


0	1	2	3	4



Queue

1 9 1 - 1 - 1 - 1 -

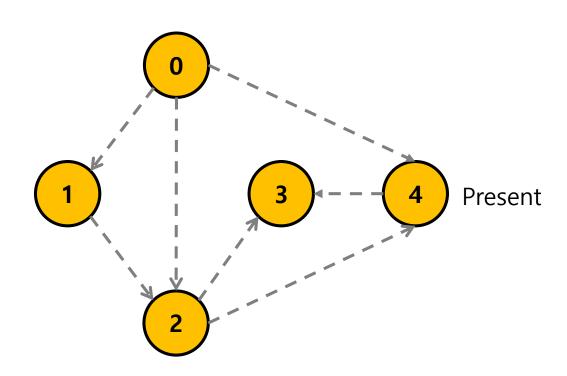


Queue

3

Visited

0 1 2 3 4

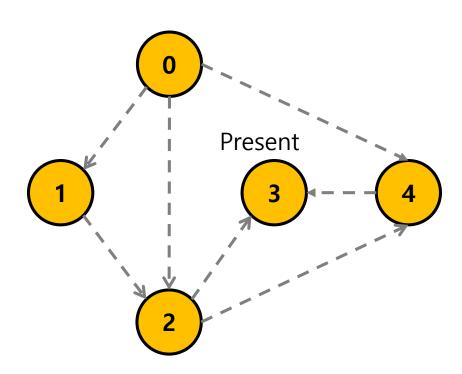


Queue

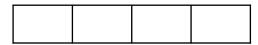
3

Visited

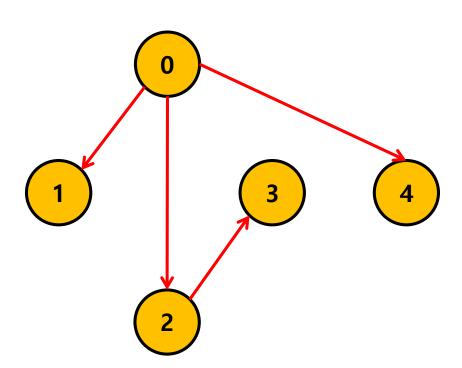
0 1 2 3 4



Queue



0 1 2 3 4	C)	1	2	3	4
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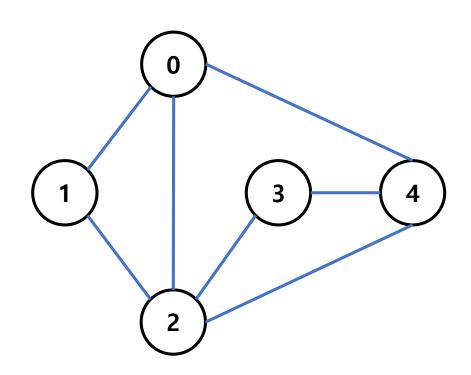


Queue



0 1 2 3 4	C)	1	2	3	4
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- Analysis



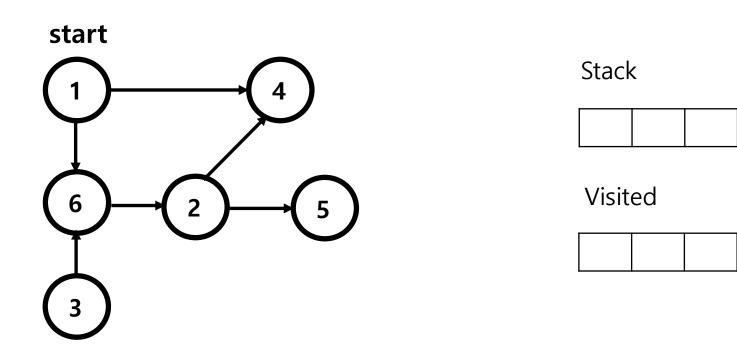
Adjacent list: O(V + E)(V is vetex and E is edge)

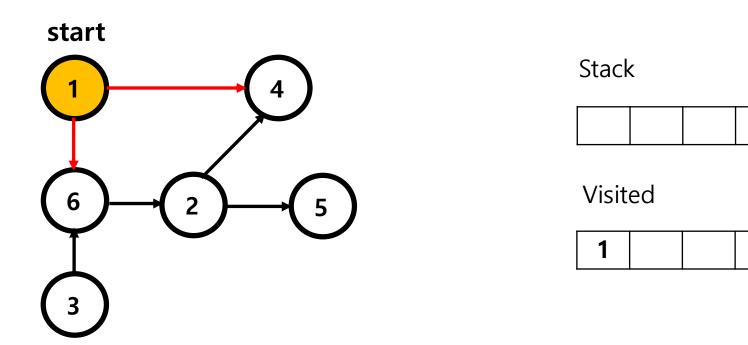
3. Topological Sort

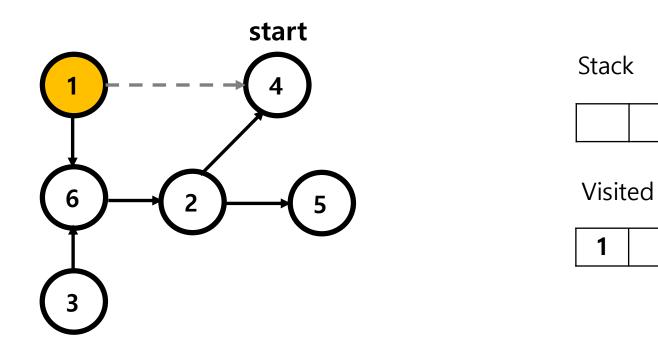
- Topological sort (DFS)

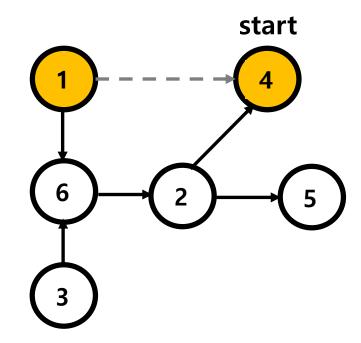
```
1 void Topological sort(){
       //About all vertices, if not visited execute DFS
       for(int i=1; i<=n; i++){</pre>
           if(visited[i]==false){
               dfs(i);
6
8
       //print stack elements
10
      while(!stack.empty()){
11
           cout << stack.top();</pre>
           stack.pop();
13
14
```

```
1 void dfs(vertex start){
       visited[start] = true; // checked start vertex
       // About adjacent vertices, if not visited execute dfs
       for(each vertex v in adjacent start){
           if(visited[v] == false){
               dfs(v);
10
11
       stack.push(start); // push the start vertex in stack
12
13 }
```



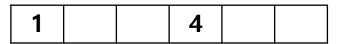


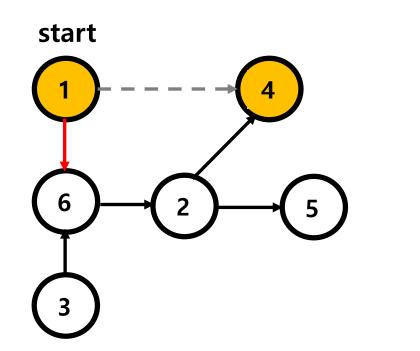




Stack

4

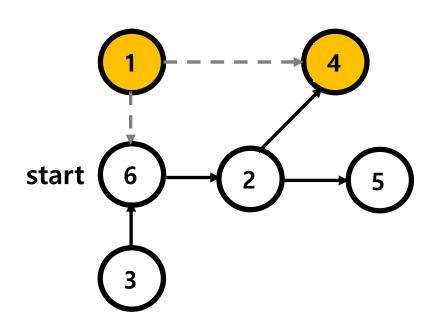




Stack

4



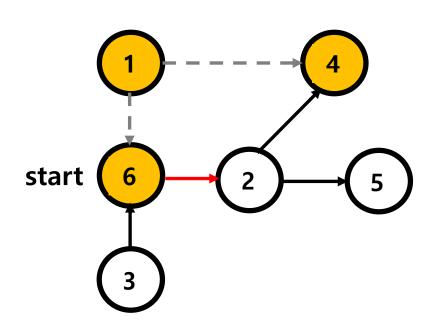


Stack

4

Visited

1 4

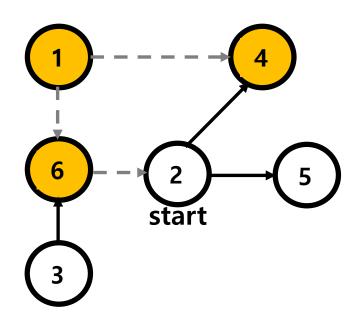


Stack

4

Visited

1 4 6

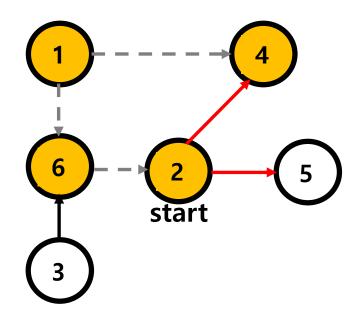


Stack

4

Visited

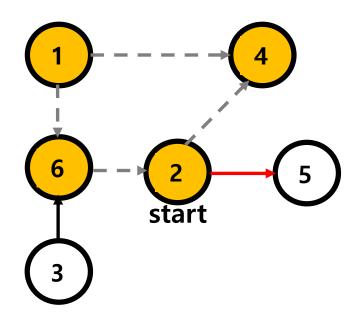
1 4 6



Stack



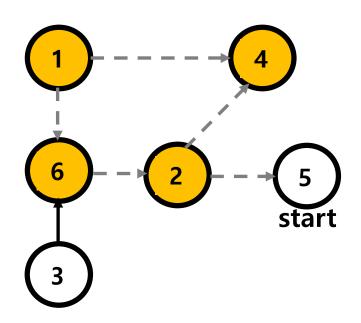
1	2	4	6



Stack



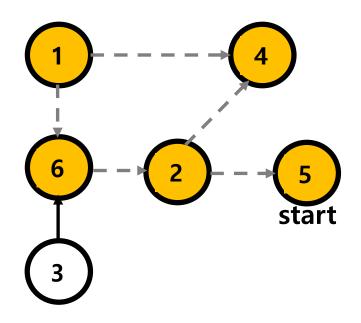
1	2	4	6



Stack



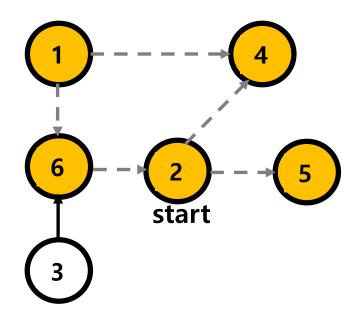
1	2	4	6



Stack



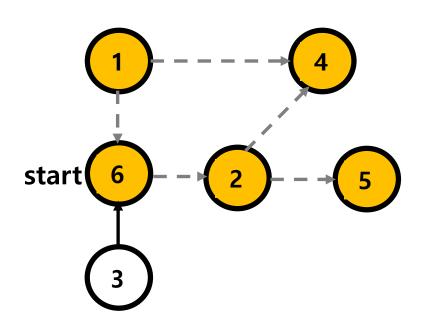
1	2	4	5	6



Stack

4	5		

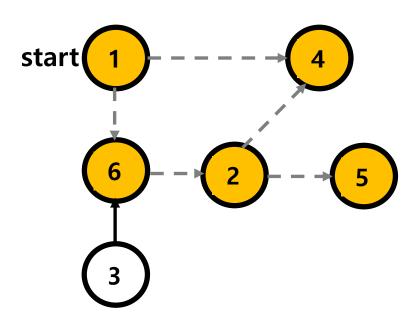
1	2	4	5	6
				i



Stack

7 3 2	4 5
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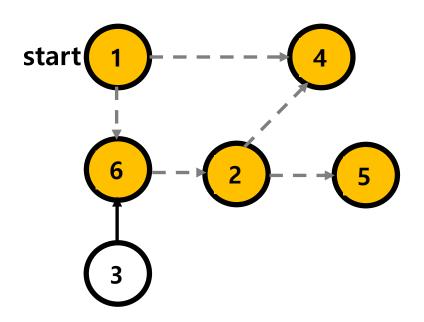
1	2	4	5	6



Stack

4 5 2	6		
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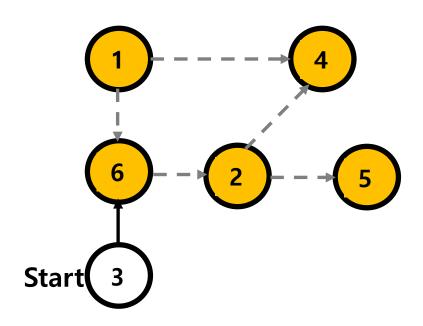
1	2	4	5	6



Stack

4	5	2	6	1	

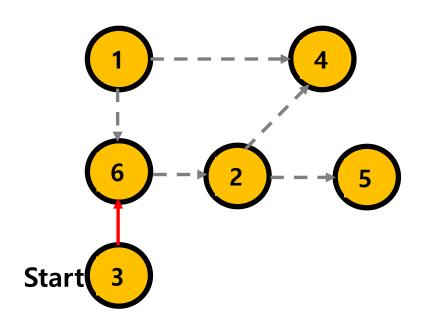
1	2	4	5	6



Stack

	4	5	2	6	1	
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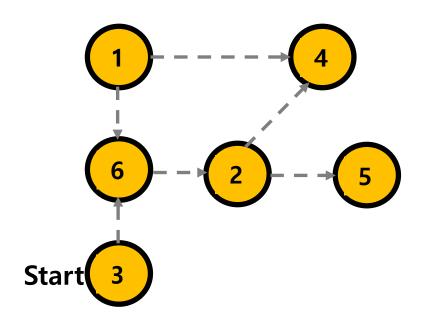
1	2	4	5	6



Stack

4	5	2	6	1	

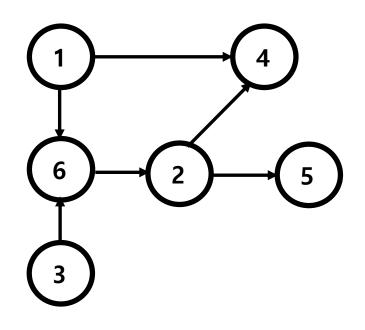
1	2	3	4	5	6

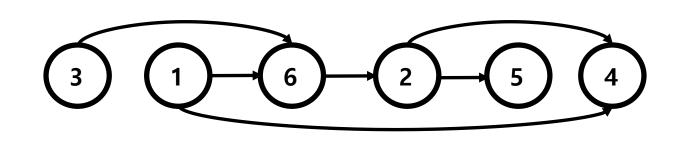


Stack

4	5	2	6	1	3
	1		1		

1	2	3	4	5	6





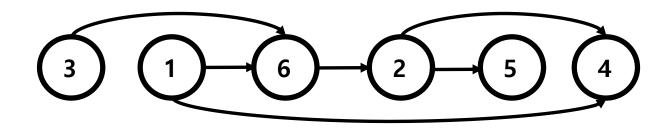
Stack

4 5 2 6 1 3

sorting

3 1 6 2 5 4

- Analysis

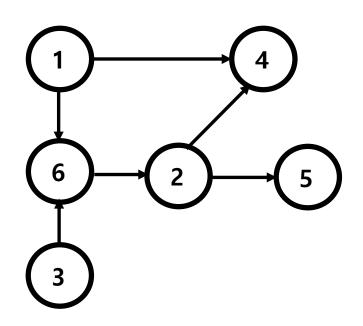


Adjacent list: O(V + E)(V is vetex and E is edge)

- Topological sort (BFS)

```
1 void Topological sort(){
2
3    for(int i=1; i<=n; i++){
4        if(indegree[i]==0){
5            q.push(i);
6        }
7     }
8     bfs();
9 }</pre>
```

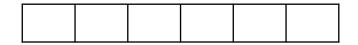
```
1 void bfs(){
       while(!q.empty()){
           int now = q.front();
           visited[now] = true; //checked now vertex
           q.pop();
           cout << now << " ";
           for(each vertex v in adjacent now){
               indegree[v]--;
               if(visited[v]==false && indegree[v]==0){
                   q.push(v);
11
12
13
14
15 }
```



Indegree

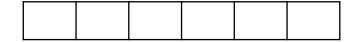
1	2	3	4	5	6
0	1	0	2	1	2

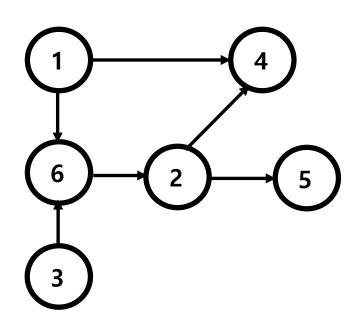
Queue



Visited







Indegree

1	2	3	4	5	6
0	1	0	2	1	2

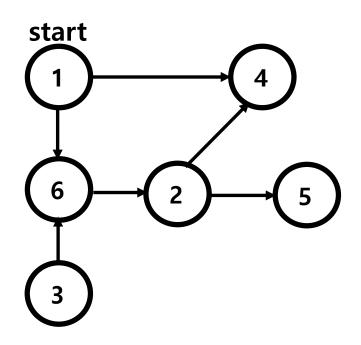
Queue

1	3		

Visited







Indegree

_ 1	2	3	4	5	6
0	1	0	2	1	2

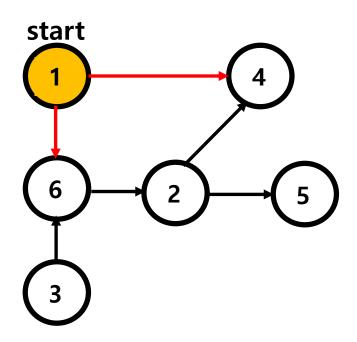
Queue



Visited







Indegree

1	2	3	4	5	6
0	1	0	2	1	2

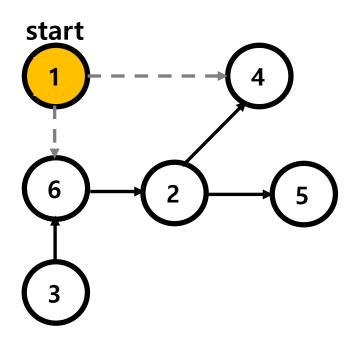
Queue

3			
	1		

Visited







Indegree

1	2	3	4	5	6
0	1	0	1	1	1

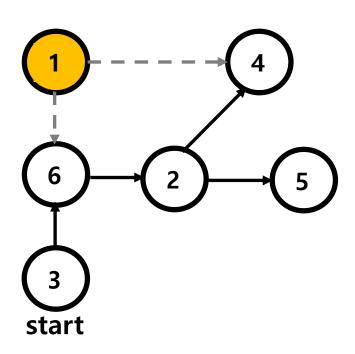
Queue



Visited



1			
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Indegree

1	2	3	4	5	6
0	1	0	1	1	1

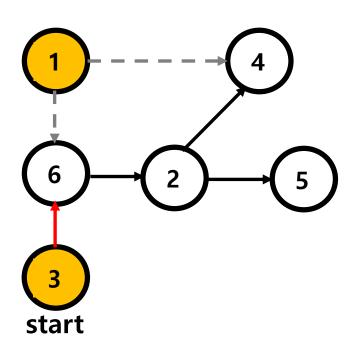
Queue



Visited



1			
•			



Indegree

1	2	3	4	5	6
0	1	0	1	1	1

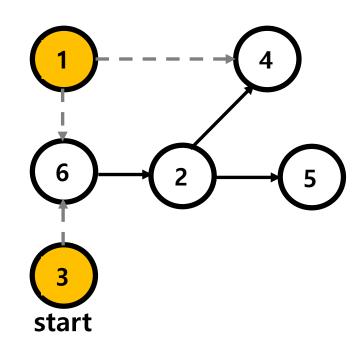
Queue



Visited

1	3		
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1



Indegree

1	2	3	4	5	6
0	1	0	1	1	0

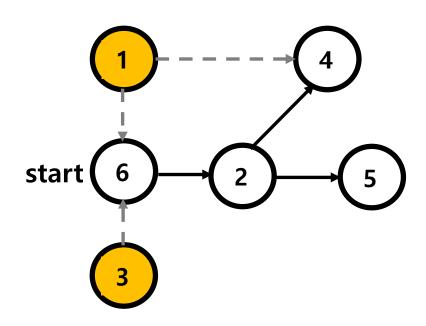
Queue



Visited



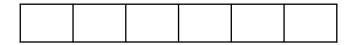
1 3				
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Indegree

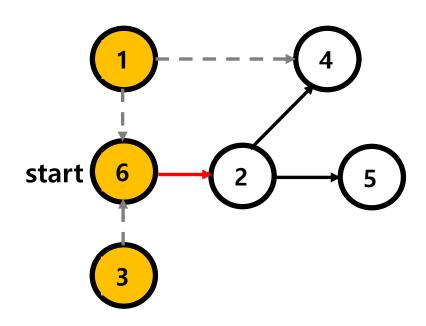
1	2	3	4	5	6
0	1	0	1	1	0

Queue



Visited

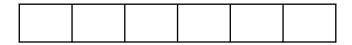
1	3		
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Indegree

1	2	3	4	5	6
0	1	0	1	1	0

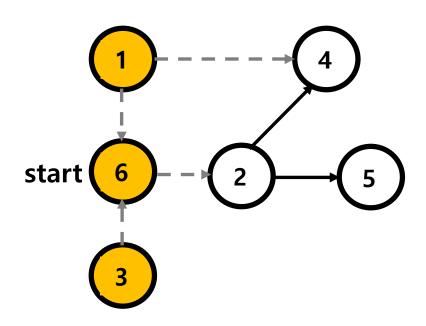
Queue



Visited

1 3 6

1 3		
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Indegree

1		3		5	6
0	0	0	1	1	0

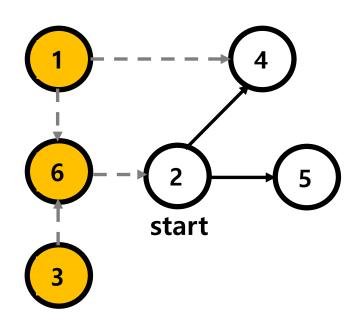
Queue

2					
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Visited

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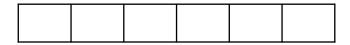
1 3 6	
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Indegree

1	2	3	4	5	6
0	0	0	1	1	0

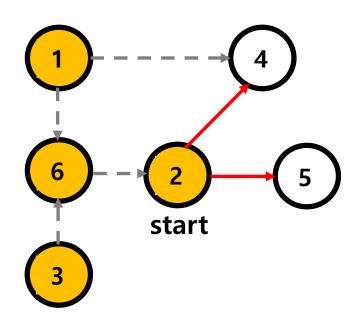
Queue



Visited

1 3	6
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1 3	6			
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Indegree

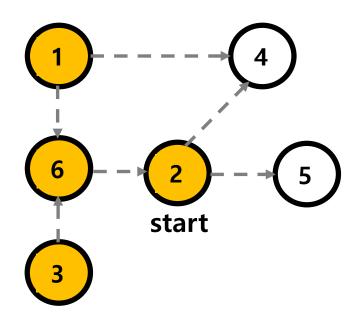
1	2	3	4	5	6
0	0	0	1	1	0

Queue



Visited

1 3



Indegree

1	2	3	4	5	6
0	0	0	0	0	0

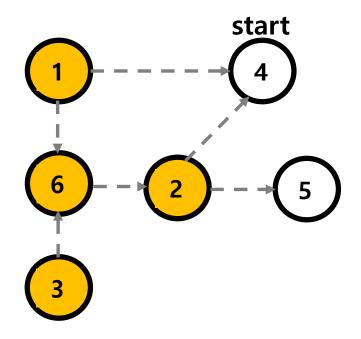
Queue

4	5				
---	---	--	--	--	--

Visited

2 3

1 3 6	2		
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Indegree

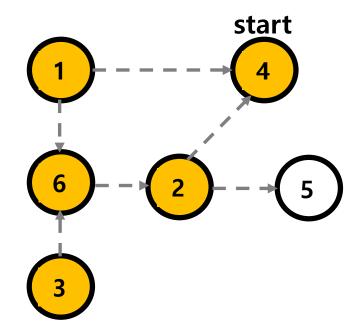
1	2	3	4	5	6
0	0	0	0	0	0

Queue



Visited

1 3 6	2		
-------	---	--	--



Indegree

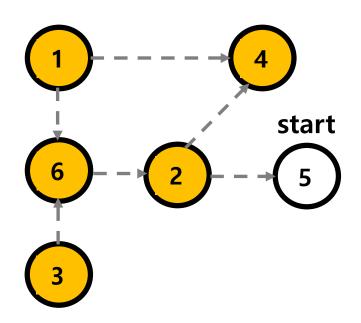
_ 1			4		
0	0	0	0	0	0

Queue



Visited

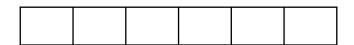
1	3	6	2	4	



Indegree

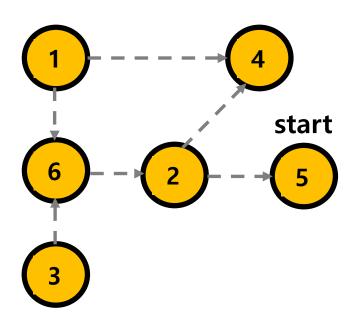
1	2	3	4	5	6
0	0	0	0	0	0

Queue



Visited

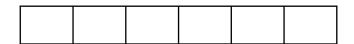
1	3	6	2	4	



Indegree

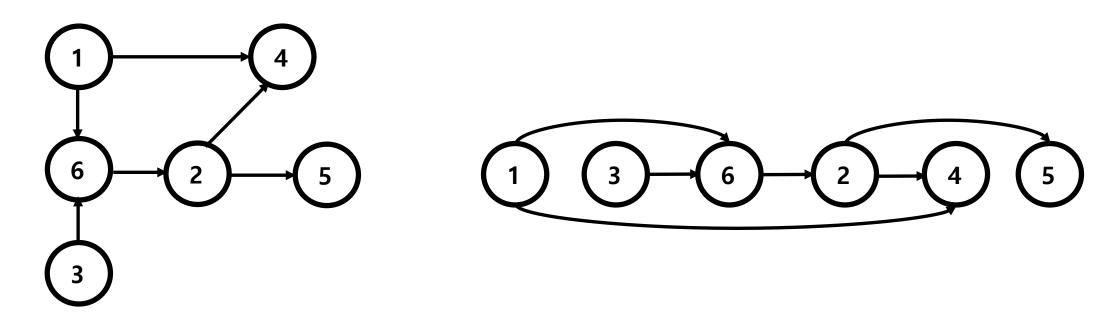
1	2	3	4	5	6
0	0	0	0	0	0

Queue



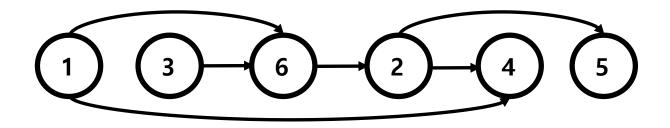
Visited

1 2 3 4 5 6



1 3	6	2	4	5
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- Analysis



Adjacent list: O(V + E)(V is vetex and E is edge)

4. Conclusion

Summary

- 1. If the graph is 'DAG', we can make topological sort by using DFS, BFS.
- 2. DFS and BFS requires O(V + E) time.
- 3. The results can appear in many ways.
- 4. When execute topological sort, you may need some data structures such as stack or queue.
- 5. No matter what method you choose, the sorting runs O(V + E) time.

Q & A

Thank you!