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Pay One, Get Hundreds for Free: **Reducing Cloud Costs through Shared Query Execution**

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Motivation: Query-as-a-service

- Pros
 - No need to move data outside the cloud
 - No infrastructure deployment
 - No database maintenance
- Pricing model
 - Pay-per-byte-processed
- Popular Systems

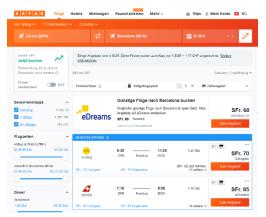
- Disadvantages
 - Expensive with frequent usage
 - No intuitive way to optimize
 - Cost
 - Throughput



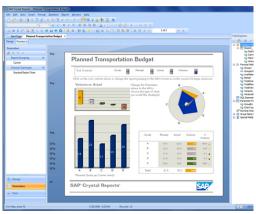




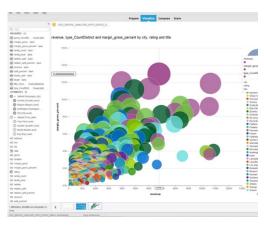
Motivation: Use cases



Parameter exploration



BI reporting



Ad-hoc analytics

- Observations:
 - High overlap among concurrent queries
 - Query burst from single user
 - Similar workload from multiple users

Work sharing: carry out redundant work only once.





Our Goal

Reduce query-as-a-service costs using work sharing techniques through query rewriting







Outline

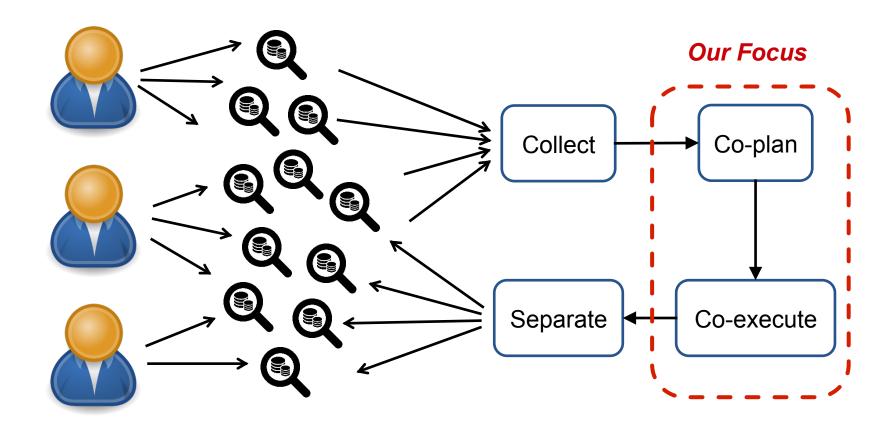
Related work







Multi Query Execution: General sharing workflow







Shared execution techniques

1988 - 2005

2005 - 2010

2010 - 2017

MQO

Scan Sharing Work Sharing

1st generation: Crescando

2nd generation: SharedDB,

MQJoin, BatchDB

- Annotate tuples with query_id attribute
- Supports all relational operators

Name	Age	Sex	query_id
Renato	30	m	Q1,Q3
Darko	30	m	Q1,Q2,Q3
Ingo	32	m	Q1



Shared execution techniques

1988 - 2005

2005 - 2010

2010 - 2017

MQO

Scan Sharing Work Sharing

- Previous approaches need dedicated engine
 - High implementation effort
 - Needs vendor support

This work: rewrite shared query plans as SQL.





Outline

- Related work
- Multi-query execution







Multi Query Execution: Data-query model

SELECT * FROM Emp WHERE ...

Q1:	•••	age	< 4	10				
Q2:	•••	name	<u> </u>	[ΚΕ	' %k%	,		
Q3:	•••	age	<=	30	AND	age	>	20

Name	Age	Sex
Renato	30	m
Darko	30	m
Ingo	32	m
Gustavo	50	m

Name	Age	Sex	query_id
Renato	30	m	{1, 3}
Darko	30	m	{ 1, 2, 3 }
Ingo	32	m	1
Gustavo	50	m	{}





Multi Query Execution: The query_id attribute

Name:	Age:	Sex:	query_id:
Renato	30	m	{ 1, 3 }
Darko	30	m	{ 1, 2, 3 }
Ingo	32	m	{1}

Result of Qx is given by

SELECT * FROM R WHERE x = ANY(query id)

This is standard SQL!



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Shared operators: Shared scan

```
SELECT * FROM Emp WHERE ...
Q1: ... age < 40
Q2: ... name LIKE '%k%'
Q3: ... age <= 30 AND age > 20
```

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Name	Age	Sex
Renato	30	m
Darko	30	m
Ingo	32	m
Gustavo	50	m

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```
SELECT *,

ARRAY_REMOVE(

ARRAY[

CASE WHEN age < 40 THEN 1 ELSE 0 END

CASE WHEN name LIKE '%k%' THEN 2 ELSE 0 END

CASE WHEN age <= 30 AND age > 20 THEN 3 ELSE 0 END

]

, 0)

FROM Emp

WHERE (age < 40) OR (name LIKE '%k%') OR (age <= 30 AND age > 20)
```



Shared operators: Shared join

Employees

Name	Age	Sex	D_ID	query_id
Renato	30	m	1	{ 1, 3 }
Darko	30	m	1	{ 1, 2, 3 }
Ingo	32	m	1	{1}

Departments

D_ID	D_Name	query_id
1	Systems	{ 2, 3 }
2	Algorithms	{2}
3	ML	{ 1, 2, 3 }

Name	Age	Sex	D_ID	D_Name	query_id
Renato	30	m	1	Systems	{3}
Darko	30	m	1	Systems	{ 2, 3 }





Shared operators: Shared join

Employees

Name	Age	Sex	D_ID	query_id
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Departments

D_ID	D_Name	query_id
1	Systems	{ 2, 3 }
2	Algorithms	{2}
3	ML	{ 1, 2, 3 }

```
WITH shared_emp AS (...), -- shared scan on emp
     shared dep AS (...), -- shared scan on dep
shared join helper AS (
   SFI FCT
      R.name, R.age, R.sex, S.d id, S.d name
      ARRAY INTERSECT(R.query id, S.query id) AS query id
   FROM shared emp e JOIN shared dep d ON e.d id = d.d id
SELECT * FROM shared join helper
WHERE CARDINALITY(query id) > 0
```



Shared operators: Other shared operators

- GROUP BY
 - Use UNNEST to duplicate tuples for each query
 - Then group by query_id
- LIMIT / TOP K
 - PARTITION BY query_id
 - Then filter by RANK()

Enough to express all TPC-H queries!





Outline

- Related work
- Multi-query execution
- **Evaluation**
- Summary





Evaluation: Query-as-a-Service systems

- Systems-under-test
 - Amazon Athena
 - Google BigQuery
- Run SQL queries against files in cloud storage
 - Apache Parquet
 - Google internal columnar storage
- Micro-benchmarks
- End-to-end query execution





Evaluation: Shared scan

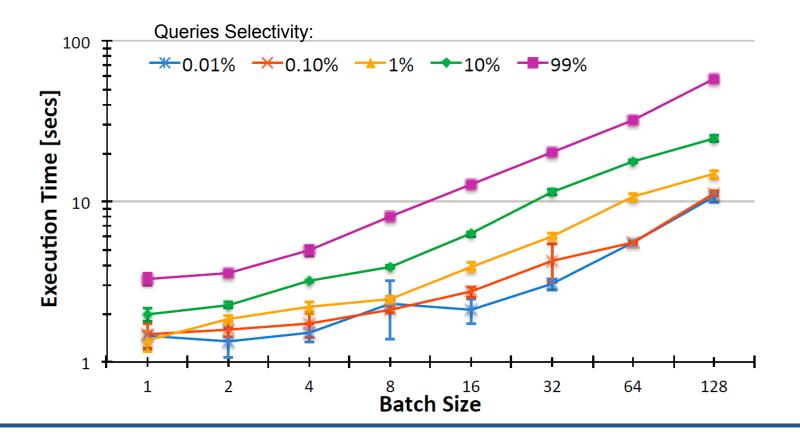
- Data
 - Lineitem from TPC-H Scale Factor 100
- Workload
 - Parameterized queries
 - Queries in a group/batch have equal selectivity







Evaluation: Shared scan

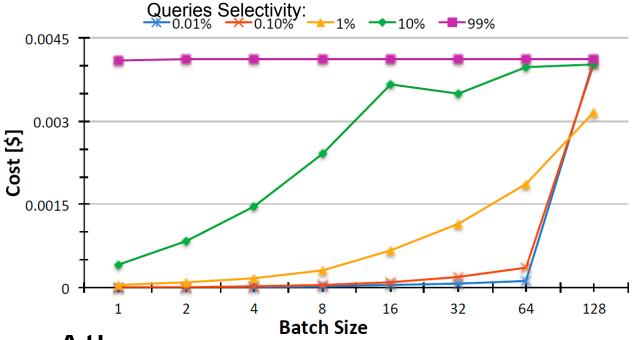


Execution time increases sublinearly with query count.





Evaluation: Shared scan



Amazon Athena

27GB, Parquet format

Pay per processed byte

Monetary cost (almost) independent of batch size.





Evaluation: TPC-H

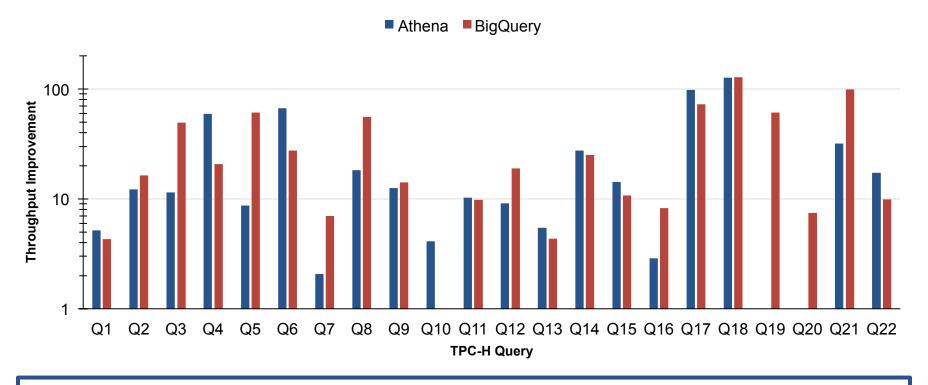
- Data
 - TPC-H Scale Factor 100
- Simulate multiple apps interacting with QaaS
 - Each app emitting different TPC-H queries
- Workload
 - 128 instantiations of each TPC-H query type





Evaluation: TPC-H

Throughput improvement over Query-at-a-Time (batch size = 128)



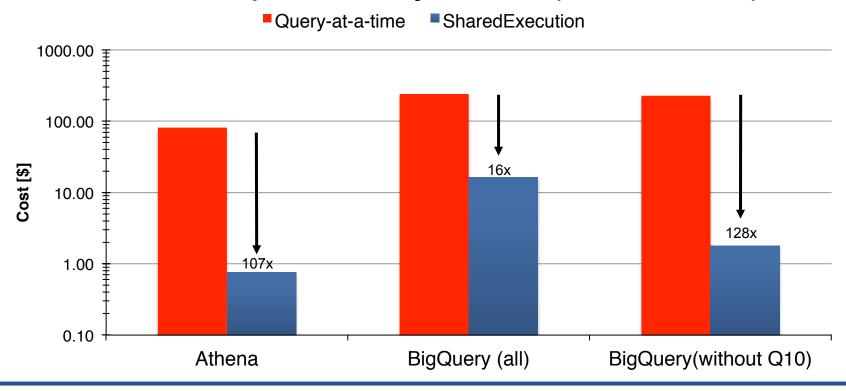
Considerable speedup: up to ~128x.





Evaluation: TPC-H

Cost compared to Query-at-a-Time (batch size = 128)



Again: cost (almost) independent of batch size.





Outline

- Related work
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- Conclusions



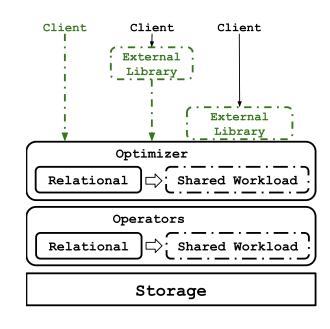


Summary

- Shared query execution through query rewriting
 - Rewriting can be done with standard SQL
 - Evaluate our approach by executing queries end-to-end
- Show improvements in execution cost and throughput

Conclusion

- Sharing can be implemented in:
 - Client
 - Library
 - Middleware







Thanks!



