

Sandbox

Ben Johnson Weitang Liu Agnieszka Łupińska
Muhammad Osama John D. Owens Yuechao Pan
Leyuan Wang Xiaoyun Wang Carl Yang
UC Davis

Contents

| | | |
|---|---------------------------------------|---|
| 1 | Sandbox for playing with pandoc/slate | 2 |
|---|---------------------------------------|---|

Chapter 1

Sandbox for playing with pandoc/slate

Normal table

| Operation | Self | Peer | Host | All |
|----------------|--------|-------|---------|-------|
| Regular read | 448.59 | 14.01 | 444.74 | 12.17 |
| Regular write | 442.98 | 16.21 | 16.18 | 12.17 |
| Regular update | 248.80 | 11.71 | 0.0028 | 6.00 |
| Random read | 6.78 | 1.43 | 2.39 | 4.04 |
| Random write | 6.63 | 1.14 | 3.47E-5 | 3.82 |
| Random update | 3.44 | 0.83 | 1.92E-5 | 2.08 |

Pretty wide table

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|-------------------------|-----------------|------------|----------------------|-------------|----------------------------|
| Modularity optimization | $10(E + V) / p$ | 20V bytes | $E/p : 2V$ | Okay | $88E/p + 12V$ bytes |
| Graph contraction | $5E / p + E'$ | 8E' bytes | $5E/p + E' : 8E'$ | Hard | 16E' bytes |
| Louvain | $10(E + V) / p$ | 20V bytes | $E/p : 2V$ | Okay | $88E/p + 12V + 16E'$ bytes |

Pretty wide table, with math

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|-------------------------|---------------|-------------|----------------------|-------------|----------------------------|
| Modularity optimization | $10(E + V)/p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V$ bytes |
| Graph contraction | $5E/p + E'$ | $8E'$ bytes | $5E/p + E' : 8E'$ | Hard | $16E'$ bytes |
| Louvain | $10(E + V)/p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V + 16E'$ bytes |

JDO hacked version of above

| Parts | Comp cost | Comm cost | Comp/comm ratio | Scalability | Memory usage (B) |
|-------------------|-----------------|-------------|-------------------|-------------|----------------------|
| Modularity optim. | $10(E + V) / p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V$ |
| Graph contraction | $5E / p + E'$ | $8E'$ bytes | $5E/p + E' : 8E'$ | Hard | $16E'$ |
| Louvain | $10(E + V) / p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V + 16E'$ |

and with math!

| Parts | Comp cost | Comm cost | Comp/comm ratio | Scalability | Memory usage (B) |
|-------------------|---------------|-------------|-------------------|-------------|----------------------|
| Modularity optim. | $10(E + V)/p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V$ |
| Graph contraction | $5E/p + E'$ | $8E'$ bytes | $5E/p + E' : 8E'$ | Hard | $16E'$ |
| Louvain | $10(E + V)/p$ | $20V$ bytes | $E/p : 2V$ | Okay | $88E/p + 12V + 16E'$ |

Even wider table

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|------------------|------------|------------|----------------------|-------------|--------------|
| Wedge generation | dE/p | | | | |

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|--------------------------------|---|------------------------|-------------------------------------|-------------|--------------|
| Wedge communication | 0 | $aE/p \times 12$ bytes | | | |
| Wedge checking | $aE/p \times \log(d)$ | | | | |
| AllReduce | $2V$ | $2V \times 4$ bytes | | | |
| Triangle Counting | $(d + a \times \log(d))E/p + 8V + 2V$ | $aE/p \times 12$ bytes | $\sim (d + a \times \log(d)) : 12a$ | Okay | |
| Scan Statistics (wedge checks) | $(d + a \times \log(d))E/p + 8V + 2V + V/p$ | $12aE/p + 8V$ bytes | $\sim (d + a \times \log(d)) : 12a$ | Okay | |
| Scan Statistics (intersection) | $Vdd + V/p$ | $8V$ bytes | $dd : 8$ | Perfect | |

Even wider table, with math!

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|--------------------------------|---|-----------------------------|-------------------------------|-------------|--------------|
| Wedge generation | dE/p | | | | |
| Wedge communication | 0 | $aE/p \times 12$ bytes | | | |
| Wedge checking | $aE/p \times \log(d)$ | | | | |
| AllReduce | $2V$ | $2V \times 4$ bytes | | | |
| Triangle Counting | $(d + a \times \log(d))E/p + 8V + 2V$ | $aE/p \times 12 + 8V$ bytes | $(d + a \cdot \log(d)) : 12a$ | Okay | |
| Scan Statistics (wedge checks) | $(d + a \times \log(d))E/p + 8V + 2V + V/p$ | $12aE/p + 8V$ bytes | $(d + a \cdot \log(d)) : 12a$ | Okay | |
| Scan Statistics (intersection) | $Vdd + V/p$ | $8V$ bytes | $dd : 8$ | Perfect | |

JDO hacked version of above

| Parts | Comp. cost | Comm. cost (B) | Comp/comm ratio | Scalability |
|-------------------------------------|--|-----------------------|------------------------------------|-------------|
| Wedge generation | dE/p | | | |
| Wedge communication | 0 | $aE/p \times 12$ | | |
| Wedge checking | $aE/p \times \log(d)$ | | | |
| AllReduce | $2V$ | $2V \times 4$ | | |
| Triangle Counting | $(d + a \times \log(d))E/p + 2V$ | $aE/p \times 12 + 8V$ | $\sim(d + a \times \log(d)) : 12a$ | Okay |
| Scan Statistics | $(d + a \times \log(d))E/p + 2V + V/p$ | $12aE/p + 8V$ | $\sim(d + a \times \log(d)) : 12a$ | Okay |
| (with wedge checks) | $Vdd + V/p$ | $8V$ | $dd : 8$ | Perfect |
| Scan Statistics (with intersection) | | | | |

Table with line breaks

| Application | Computation to communication ratio | Scalability | Implementation difficulty |
|---------------------------------|---|----------------------|---------------------------|
| Louvain | $E/p : 2V$ | Okay | Hard |
| Graph SAGE | $\sim CF : \min(C, 2p) \times 4$ | Good | Easy |
| Random walk | Duplicated graph: infinity Distributed graph: $1 : 24$ | Perfect Very poor | Trivial Easy |
| Graph search: Uniform | $1 : 24$ | very poor | Easy |
| Graph search: Greedy | Straight forward: $d : 24$ Pre-visit: $1:24$ | Poor very poor | Easy Easy |
| Graph search: Stochastic greedy | Straight forward: $d : 24$ Pre-visit: $\log(d) : 24$ | Poor very poor | Easy Easy |
| Geo location | Explicit movement: $25E/p : 4V$ UVM or peer access: $25 : 1$ | Okay Good | Easy Easy |

| Application | Computation to communication ratio | Scalability | Implementation difficulty |
|-------------------------------|--|----------------------|------------------------------|
| Vertex nomination | $E : 8V \times \min(d, p)$ | Okay | Easy |
| Scan statistics | Duplicated graph: infinity | Perfect | Trivial |
| | Distributed graph: $\sim (d + a \cdot \log(d)) : 12$ | Okay | Easy |
| Sparse fused lasso | $\sim a:8$ | Less than okay | Hard |
| Graph projection | Duplicated graph : infinity | Perfect | Easy |
| | Distributed graph : $dE/p + E' : 6E'$ | Okay | Easy |
| Local graph clustering | $(6 + d)/p : 4$ | Good | Easy |
| Seeded graph matching | | | |
| Application classification | | | |

Table with line breaks, laid out as a grid table

| Application | Computation to communication ratio | Scalability | Implementation difficulty |
|------------------------------------|---|-------------|------------------------------|
| Louvain | $E/p : 2V$ | Okay | Hard |
| Graph SAGE | $\sim CF : \min(C, 2p) \cdot 4$ | Good | Easy |
| Random walk | Duplicated graph: infinity | Perfect | Trivial |
| | Distributed graph: $1 : 24$ | Very poor | Easy |
| Graph search: Uniform | $1 : 24$ | Very poor | Easy |
| Graph search: Greedy | Straightforward: $d : 24$ | Poor | Easy |
| | Pre-visit: $1:24$ | Very poor | Easy |
| Graph search: Stochastic greedy | Straightforward: $d : 24$ | Poor | Easy |
| | Pre-visit: $\log(d) : 24$ | Very poor | Easy |
| Geolocation | Explicit movement: $25E/p : 4V$ | Okay | Easy |
| | UVM or peer access: $25 : 1$ | Good | Easy |
| Vertex nomination | $E : 8V \cdot \min(d, p)$ | Okay | Easy |
| Scan statistics | Duplicated graph: infinity | Perfect | Trivial |
| | Distributed graph: $\sim (d + a \cdot \log(d)) : 12$ | Okay | Easy |

| Application | Computation to communication ratio | Scalability | Implementation difficulty |
|----------------------------|---------------------------------------|----------------|---------------------------|
| Sparse fused lasso | $\sim a : 8$ | Less than okay | Hard |
| Graph projection | Duplicated graph : infinity | Perfect | Easy |
| | Distributed graph : $dE/p + E' : 6E'$ | Okay | Easy |
| Local graph clustering | $(6 + d)/p : 4$ | Good | Easy |
| Seeded graph matching | | | |
| Application classification | | | |

JDO hacked version of above

| Application | Computation to communication ratio | Scalability | Implementation diff. |
|-------------------------|------------------------------------|-------------|----------------------|
| Louvain | $E/p : 2V$ | Okay | Hard |
| Graph SAGE | $\sim CF : \min(C, 2p) \times 4$ | Good | Easy |
| Random walk | Duplicated graph: infinity | Perfect | Trivial |
| Random walk | Distrib. graph: $1 : 24$ | Very poor | Easy |
| Graph search: Uniform | $1 : 24$ | Very poor | Easy |
| Graph search: Greedy | Straightforward: $d : 24$ | Poor | Easy |
| Graph search: Greedy | Pre-visit: $1:24$ | Very poor | Easy |
| G.S.: Stochastic greedy | Straightforward: $d : 24$ | Poor | Easy |
| G.S.: Stochastic greedy | Pre-visit: $\log(d) : 24$ | Very poor | Easy |
| Geolocation | Explicit movement: $25E/p : 4V$ | Okay | Easy |
| Geolocation | UVM or peer access: $25 : 1$ | Good | Easy |
| Vertex nomination | $E : 8V \times \min(d, p)$ | Okay | Easy |

| Application | Computation to communication ratio | Scalability | Implementation diff. |
|---------------------------|--|----------------------|-------------------------|
| Scan statistics | Duplicated graph: infinity | Perfect | Trivial |
| Scan statistics | Distrib. graph: \sim (d + a * log(d)) : 12 | Okay | Easy |
| Sparse fused lasso | \sim a:8 | Less than okay | Hard |
| Graph projection | Duplicated graph : infinity | Perfect | Easy |
| Graph projection | Distrib. graph : dE/p + E' : 6E' | Okay | Easy |
| Local graph clustering | (6 + d)/p : 4 | Good | Easy |

Really wide table

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|----------------------------|---|---|-------------------------------|-------------|-----------------|
| <i>Feature duplication</i> | | | | | |
| Children selection | BC | 8BC bytes | 1 : 8 | Poor | |
| Child-centric comp. | BCF x (2 + L + Wf1.y + Wa1.y) | 4B x (F + Wf1.y + Wa1.y) x min(C, 2p) bytes | \sim CF : min(C, 2p) x 4 | Good | |
| Source-centric comp. | B x (CF + (Wf1.y + Wa1.y) x (C + F + Wf2.y + Wa2.y) | 0 bytes | N.A. | N.A. | |

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|------------------------------|---|---|--|-------------|--------------|
| Graph SAGE | $B \times (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \times \min(C, 2p) + Wa1.y) \times (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \times (F + Wf1.y + Wa1.y) \times \min(C, 2p)$ bytes | at least $\sim CF : \min(C, 2p) \times 4$ | Good | |
| <i>Direct feature access</i> | | | | | |
| Child-centric comp. | $BCF \times (2 + L + Wf1.y + Wa1.y)$ | $4B \times ((F + Wf1.y + Wa1.y) \times \min(C, 2p) + CLF)$ bytes | $\sim (2 + L + Wf1.y + Wa1.y) : 4L$ | poor | |
| Graph SAGE | $B \times (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \times \min(C, 2p) + Wa1.y) \times (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \times (F + Wf1.y + Wa1.y) \times \min(C, 2p) + 4BCFL$ bytes | $\sim (2 + L + Wf1.y + Wa1.y) : 4L$ | poor | |
| <i>Feature in UVM</i> | | | | | |
| Child-centric comp. | $BCF \times (2 + L + Wf1.y + Wa1.y)$ | $4B \times (F + Wf1.y + Wa1.y) \times \min(C, 2p)$ bytes over GPU-GPU + $4BCLF$ bytes over GPU-CPU | $\sim (2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor | |

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|------------|---|--|--|-------------|-----------------|
| Graph SAGE | $B \times (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \times (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \times (F + Wf1.y + Wa1.y) \times \min(C, 2p)$ bytes over GPU-GPU + $4BCFL$ bytes over GPU-CPU | $\sim (2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor | |

Really wide table, with math

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|------------------------------|---|--|-------------------------------------|-------------|-----------------|
| <i>Feature duplication</i> | | | | | |
| Children selection | BC | $8BC$ bytes | $1 : 8$ | Poor | |
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes | $CF : \min(C, 2p) \cdot 4$ | Good | |
| Source-centric comp. | $B \cdot (CF + (Wf1.y + Wa1.y) \cdot (C + F + Wf2.y + Wa2.y))$ | 0 bytes | N.A. | N.A. | |
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes | at least $CF : \min(C, 2p) \cdot 4$ | Good | |
| <i>Direct feature access</i> | | | | | |

| Parts | Comp. cost | Comm. cost | Comp. to comm. ratio | Scalability | Memory usage |
|------------------------|---|---|---|--------------|-----------------|
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot ((F + Wf1.y + Wa1.y) \cdot \min(C, 2p) + CLF)$ bytes | $(2 + L + Wf1.y + Wa1.y) : 4L$ | poor | |
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p) + 4BCFL$ bytes | $(2 + L + Wf1.y + Wa1.y) : 4L$ | poor | |
| <i>Feature in UVM</i> | | | | | |
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes over GPU-GPU + $4BCFL$ bytes over GPU-CPU | $(2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor | |
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes over GPU-GPU + $4BCFL$ bytes over GPU-CPU | $(2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor | |

Really wide table, with math, nicely laid out

| Parts | Computation cost | Communication cost (Bytes) | Comp. to comm. ratio | Scalability |
|------------------------------|---|--|---|-------------|
| <i>Feature duplication</i> | | | | |
| Children selection | BC | $8BC$ | 1 : 8 | Poor |
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ | $CF : \min(C, 2p) \cdot 4$ | Good |
| Source-centric comp. | $B \cdot (CF + (Wf1.y + Wa1.y) \cdot (C + F + Wf2.y + Wa2.y))$ | 0 | N.A. | N.A. |
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ | at least $CF : \min(C, 2p) \cdot 4$ | Good |
| <i>Direct feature access</i> | | | | |
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot ((F + Wf1.y + Wa1.y) \cdot \min(C, 2p) + CLF)$ | $(2 + L + Wf1.y + Wa1.y) : 4L$ | poor |
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p) + 4BCFL$ | $(2 + L + Wf1.y + Wa1.y) : 4L$ | poor |
| <i>Feature in UVM</i> | | | | |
| Child-centric comp. | $BCF \cdot (2 + L + Wf1.y + Wa1.y)$ | $4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes over GPU-GPU + $4BCFL$ bytes over GPU-CPU | $(2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor |

| Parts | Computation cost | Communication cost (Bytes) | Comp. to comm. ratio | Scalability |
|------------|---|---|---|-------------|
| Graph SAGE | $B \cdot (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \cdot (CF + C + F + Wf2.y + Wa2.y))$ | $8BC + 4B \cdot (F + Wf1.y + Wa1.y) \cdot \min(C, 2p)$ bytes over GPU-GPU + $4BCFL$ bytes over GPU-CPU | $(2 + L + Wf1.y + Wa1.y) : 4L$ over GPU-CPU | very poor |

JDO hacked version of above

| Parts | Comp. cost |
|------------------------------|---|
| <i>Feature duplication</i> | |
| Children selection | BC |
| Child-centric comp. | $BCF \times (2 + L + Wf1.y + Wa1.y)$ |
| Source-centric comp. | $B \times (CF + (Wf1.y + Wa1.y) \times (C + F + Wf2.y + Wa2.y))$ |
| Graph SAGE | $B \times (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \times (CF + C + F + Wf2.y + Wa2.y))$ |
| <i>Direct feature access</i> | |
| Child-centric comp. | $BCF \times (2 + L + Wf1.y + Wa1.y)$ |
| Graph SAGE | $B \times (C + 3CF + 3LCF + (Wf1.y + Wa1.y) \times (CF + C + F + Wf2.y + Wa2.y))$ |

Feature in UVM

| Parts | Comp. cost |
|---------------------|--|
| Child-centric comp. | $\text{BCF} \times (2 + L + \text{Wf1.y} + \text{Wa1.y})$ |
| Graph SAGE | $\text{B} \times (\text{C} + 3\text{CF} + 3\text{LCF} + (\text{Wf1.y} + \text{Wa1.y}) \times (\text{CF} + \text{C} + \text{F} + \text{Wf2.y} + \text{Wa2.y}))$ |

| Parts | Comm. cost |
|------------------------------|--|
| <i>Feature duplication</i> | |
| Children selection | 8BC bytes |
| Child-centric comp. | $4\text{B} \times (\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p})$ bytes |
| Source-centric comp. | 0 bytes |
| Graph SAGE | $8\text{BC} + 4\text{B} \times (\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p})$ bytes |
| <i>Direct feature access</i> | |
| Child-centric comp. | $4\text{B} \times ((\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p}) + \text{CLF})$ bytes |
| Graph SAGE | $8\text{BC} + 4\text{B} \times (\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p}) + 4\text{BCFL}$ bytes |
| <i>Feature in UVM</i> | |
| Child-centric comp. | $4\text{B} \times (\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p})$ bytes over GPU-GPU + 4BCLF bytes over GPU-CPU |
| Graph SAGE | $8\text{BC} + 4\text{B} \times (\text{F} + \text{Wf1.y} + \text{Wa1.y}) \times \min(\text{C}, 2\text{p})$ bytes over GPU-GPU |

| Parts | Comm. cost |
|-------|-------------------------------|
| | + 4BCFL bytes over GPU-CPU |

| Parts | Comp/comm ratio | Scalability |
|------------------------------|--|-------------|
| <i>Feature duplication</i> | | |
| Children selection | 1 : 8 | Poor |
| Child-centric comp. | $\sim CF : \min(C, 2p) \times 4$ | Good |
| Source-centric comp. | N.A. | N.A. |
| Graph SAGE | at least $\sim CF : \min(C, 2p) \times 4$ | Good |
| <i>Direct feature access</i> | | |
| Child-centric comp. | $\sim (2 + L + W_{fl,y} + W_{a1,y}) : 4L$ | poor |
| Graph SAGE | $\sim (2 + L + W_{fl,y} + W_{a1,y}) : 4L$ | poor |
| <i>Feature in UVM</i> | | |
| Child-centric comp. | $\sim (2 + L + W_{fl,y} + W_{a1,y}) : 4L$ over GPU-CPU | very poor |
| Graph SAGE | $\sim (2 + L + W_{fl,y} + W_{a1,y}) : 4L$ over GPU-CPU | very poor |

Let's try a grid table with backslashes.

| Fruit | Price | Advantages |
|---------|-------------------------|-------------------------|
| Bananas | first line next line | first line next line |
| Bananas | first line next line | first line next line |

Multi-row (col) grid table without using backslashes.

| Fruit | Price | Advantages |
|---------|-------------------------|----------------------|
| Bananas | first line next line | first line next line |
| Bananas | first line next line | first line next line |

this won't work:

```
+-----+-----+-----+-----+
| Header row, column 1 | Header 2 | Header 3 | Header 4 |
| (header rows optional) |          |          |          |
+=====+=====+=====+=====+
| body row 1, column 1 | column 2 | column 3 | column 4 |
+-----+-----+-----+-----+
| body row 2          | Cells may span columns. |
+-----+-----+-----+-----+
| body row 3          | Cells may | - Table cells |
+-----+-----+-----+-----+
| body row 4          |          | - contain     |
|                     |          | - body elements. |
+-----+-----+-----+-----+
```

this may:

| Header row, column 1 (header rows optional) | Header 2 | Header 3 | Header 4 |
|--|-------------|--|----------|
| body row 1, column 1 | column 2 | column 3 | column 4 |
| body row 2 | Cells may s | an columns | . |
| body row 3 | Cells may | <ul style="list-style-type: none"> • Table ce | lls |
| body row 4 | | <ul style="list-style-type: none"> • body ele | ments. |

Some care must be taken with grid tables to avoid undesired interactions with cell text in rare cases. For example, the following table contains a cell in row 2 spanning from column 2 to column 4:

| row 1, col 1 | column 2 | column 3 | column 4 |
|--------------|----------|----------|----------|
| row 2 | | | |
| row 3 | | | |

If a vertical bar is used in the text of that cell, it could have unintended effects

if accidentally aligned with column boundaries:

| | | | |
|--------------|------------|-----------|----------|
| row 1, col 1 | column 2 | column 3 | column 4 |
| row 2 | Use the co | mmand “ls | more“. |
| row 3 | | | |

Several solutions are possible. All that is needed is to break the continuity of the cell outline rectangle. One possibility is to shift the text by adding an extra space before:

| | | | |
|--------------|-----------|------------|----------|
| row 1, col 1 | column 2 | column 3 | column 4 |
| row 2 | Use the c | ommand “ls | more“. |
| row 3 | | | |