

# ① Fast Fourier Transform

Problem: Evaluate a polynomial  
$$\text{Poly}(x) = \sum_{i=0}^{n-1} p_i x^i$$
 at  $n$  distinct points

$n$  is a power of 2

Given  $p[0 \dots n-1]$  array of doubles  
 $x[0 \dots n-1]$  array of values  
 $p[i] = p_i$   
 $x[i] = x_i$

Find  $\text{Poly}[\text{poly}(x[0]), \text{poly}(x[1]), \dots, \text{poly}(x[n-1])]$

Write the  $n^2$  method for computing  
 $\text{Poly}[i] = \text{poly}(x[i])$  for all  $0 \leq i \leq n-1$

## ① FFT

Goal: Find an  $n \log n$  algorithm

D+C: Identify a way of splitting the polynomial into 2  $\frac{1}{2}$  sized problems

$$\text{Poly}(x) = \sum_{i=0}^{n-1} p_i x^i$$

split into odd and even polynomials

$$\text{Poly E}(x) = p_0 + p_2 x^2 + p_4 x^4 + \dots + p_{n-2}$$

even term poly  $\sum_{i=0}^{n/2-1} p_{2i} x^{2i}$

odd term poly  $\sum_{i=0}^{n/2-1} p_{2i+1} x^{2i+1}$

$$\text{Poly}(x) = \sum_{i=0}^{n/2-1} p_{2i} (x^2)^i + x \sum_{i=0}^{n/2-1} p_{2i+1} (x^2)^i$$

$$= \text{Poly E}(x^2) + x \text{Poly O}(x^2)$$

Two smaller problems of  $\frac{1}{2}$  size  $n/2$

one with just the even coefficients Poly E

— " — odd — " — Poly O

② start writing the pseudo code  
 Problem decomposition code for odd/even  
 split

Given  $p[0 \dots n-1]$   
 and  $start = 0$  and  $end = n-1$

write code to "shuffle" the coefficients in  
 $p[]$  such that

$PE[0 \dots n/2-1]$  contains the even coefficients

$PO[0 \dots n/2-1]$  contains the odd coefficients  
 Do this with newly allocated arrays.

③ Deal with the evaluate values  $X[0 \dots n-1]$

A problem instance is  $n$  coefficients  $n$  values.  
 think of a subproblem  $1/2$  size  
 currently we have  $1/2$  coefficients  
 BUT NOT  $1/2$  values!

Need to compute subproblems with  $1/2$  points!

How?

Note: subproblems evaluate at  $x^2$   
 and  $x$  and  $-x$  are distinct

so if  $x_i$   $0 \leq i \leq n/2 - 1$  and

$$x_{i+n/2} = -x_i$$

$$(x_i)^2 = (x_{i+n/2})^2 \quad \text{so only need half}$$

$$\text{Poly}(x_i) = \text{Poly} E(x_i^2) + x_i \text{Poly} O(x_i^2)$$

$$\text{Poly}(x_{i+n/2}) = \text{Poly} E(x_i^2) - x_i \text{Poly} O(x_i^2)$$

compute  $\text{Poly} E(x_i^2)$  once - use twice

compute  $\text{Poly} O(x_i^2)$  once - use twice

④ Write the FFT code

$\text{Poly}[] \leftarrow \text{FFT}(\text{P}[], \text{int } n, \text{X}[])$

// return an array of end-start + 1 values.

// base case

// shuffle odd / even into  $\text{PE}[]$  and  $\text{PO}[]$

// compute  $z_i^2$  from  $z_i$       $\text{X2}[]$  from  $\text{X}[]$

// Call recursion on  $2 \times \frac{1}{2}$

// solution construction  $\text{Poly}[i] =$

⑤ Use correct values of  $x$

initial  $n=8$

$$x_0 \quad x_1 \quad x_2 \quad x_3 \quad -x_0 \quad -x_1 \quad -x_2 \quad -x_3$$

next subproblem

$$x_0^2 \quad x_1^2 \quad x_2^2 \quad x_3^2$$

$$\text{so } x_0^2 = -x_2^2, \quad x_1^2 = -x_3^2$$

next  $x_0^4 \quad x_1^4$

$$\text{so } x_0^4 = -x_1^4$$

solution  $x_0 = 1, \quad x_1 = \omega$

$$\omega^8 = 1 \quad \omega \neq 1$$

In general

$$\omega = e^{i\sqrt{-1} \frac{2\pi}{n}} = \cos \frac{2\pi}{n} + i \sin \frac{2\pi}{n}$$

$$\omega^i = e^{i\sqrt{-1} \frac{2\pi}{n} i} = \cos \frac{2i\pi}{n} + i \sin \frac{2i\pi}{n}$$

Write code to precompute  $x[]$