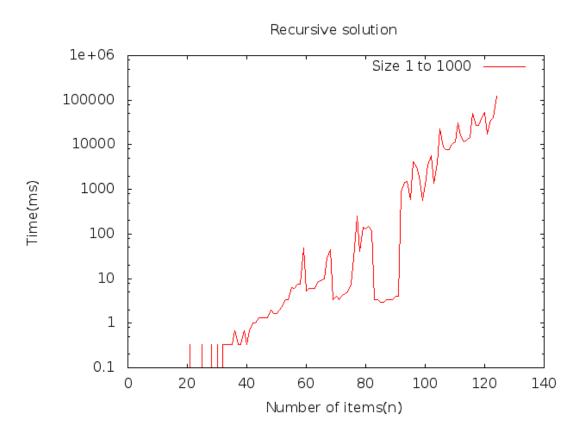
Assignment1 - Knapsack

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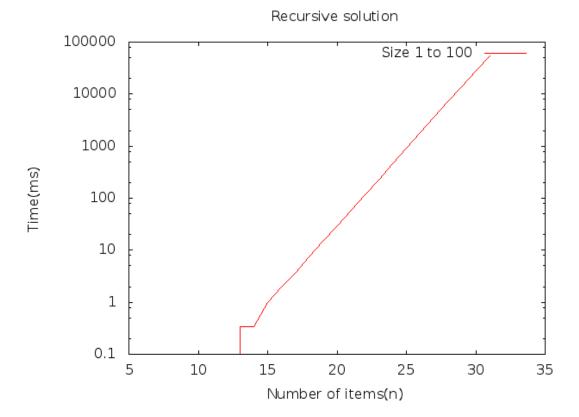
Recursive Solution



This graph shows the recursive algorithm with the Time on the y axis in log scale and the number of items it tried to put in the bag on the x axis in regular scale. Since the line exponentially increases the solution is not very efficient. The slope of the line appears to be one log over 20 items. $10^{(1/20)} \sim 1.1220$

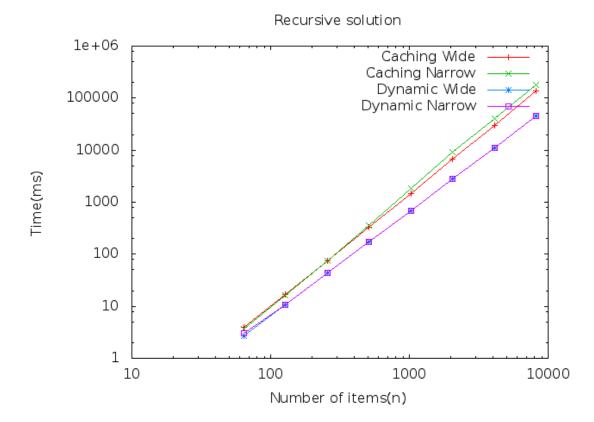
We were expecting that the problem would result in a runtime of about 2^n. That means that it ran a little faster than our worst case approximation and ran at about 1.12^n. So don't count on a worst case as actual run time but it is a good approximation and upper bound.

The first graph was using items sizes from 1-1000 and I tested item sizes 1-100 to see the difference and I ended up with a similar result:



This curve ends up more like the worst case. Which it is a worst case. The smaller items always require more calls to fill the bag. With item sizes maxing at 100 it takes 10 max items to fill the 1000 size bag. The curve of this graph is about $10^{(1/3)} \sim 2$. Meaning we have the 2ⁿ case.

Dynamic Programming and Caching Solutions



I plotted the dynamic and caching together to see which is faster and the results show that it is very similar. This also shows the wide range of sizes verses the narrow range of sizes. All of the solutions are so similar on this graph. It seems in the caching the narrow graph better as the number of items increase. I would say this is because we have less items to cache. As we try the narrow items it is more probable that we would try items of the same size or that combinations of items would hit the same part of the cache. The graph seems to show that this could have happened.

The dynamic solution must fill the entire array every time so the wide vs. narrow range or sizes makes no apparent difference. The runtime depends completely on the size of the cache.

Raw Data

Recursive

Bag Size: 1000

Runs: 30

n	time(ms) 1000	time(ms)	n	time(ms) 1000	time(ms)	n	time(ms) 1000	time(ms) 100
6	0	0	46	1		86	3	
7	0	0	47	1		87	3	
8	0	0	48	2		88	3	

9	0	0	49	2	89	3
10	0	0	50	2	90	4
11	0	0	51	2	91	4
12	0	0	52	2	92	865
13	0	0	53	3	93	1,440
14	0	0	54	3	94	1,503
15	0	1	55	6	95	602
16	0	2	56	6	96	4,233
17	0	4	57	8	97	3,233
18	0	7	58	8	98	1,930
19	0	15	59	50	99	558
20	0	30	60	5	100	1,232
21	0	59	61	6	101	3,434
22	0	115	62	6	102	5,718
23	0	225	63	6	103	1,403
24	0	459	64	9	104	3,820
25	0	906	65	9	105	23,224
26	0	1,808	66	10	106	8,949
27	0	3,596	67	29	107	7,942
28	0	7,218	68	44	108	7,909
29	0	13,858	69	3	109	10,296
30	0	27,992	70	4	110	11,740
31	0	53,455	71	3	111	31,287
32	0		72	4	112	16,008
33	0		73	5	113	11,606
34	0		74	5	114	12,709
35	0		75	8	115	14,664

36	1	7	76	29	116	50,501	
37	0	7	77	255	117	27,346	
38	0	7	78	41	118	27,611	
39	1	7	79	143	119	42,892	
40	0	8	30	130	120	53,576	
41	1	8	31	150	121	17,157	
42	1	8	32	118	122	36,204	
43	1	8	33	3	123	41,299	
44	1	8	34	3	124	125,932	
45	1	8	35	3			

The columns with time(ms) 1000 were tested with items size 1-1000 The columns with time(ms) 100 were tested with items size 1-100

Caching and Dynamic

n (#items)	Caching Wide	Caching Narrow	Dynamic Wide	Dynamic Narrow
64	4	3.66667	2.66667	3
128	16.6667	16.3333	10.6667	10.6667
256	74	73.6667	43	43.3333
512	326.667	357.667	174	173.333
1024	1483.67	1813.67	694.333	690
2048	6714.33	9165.33	2779.33	2765.33
4096	30419.7	40860.7	10939	11065.3
8192	136268	177094	44735	44815.3