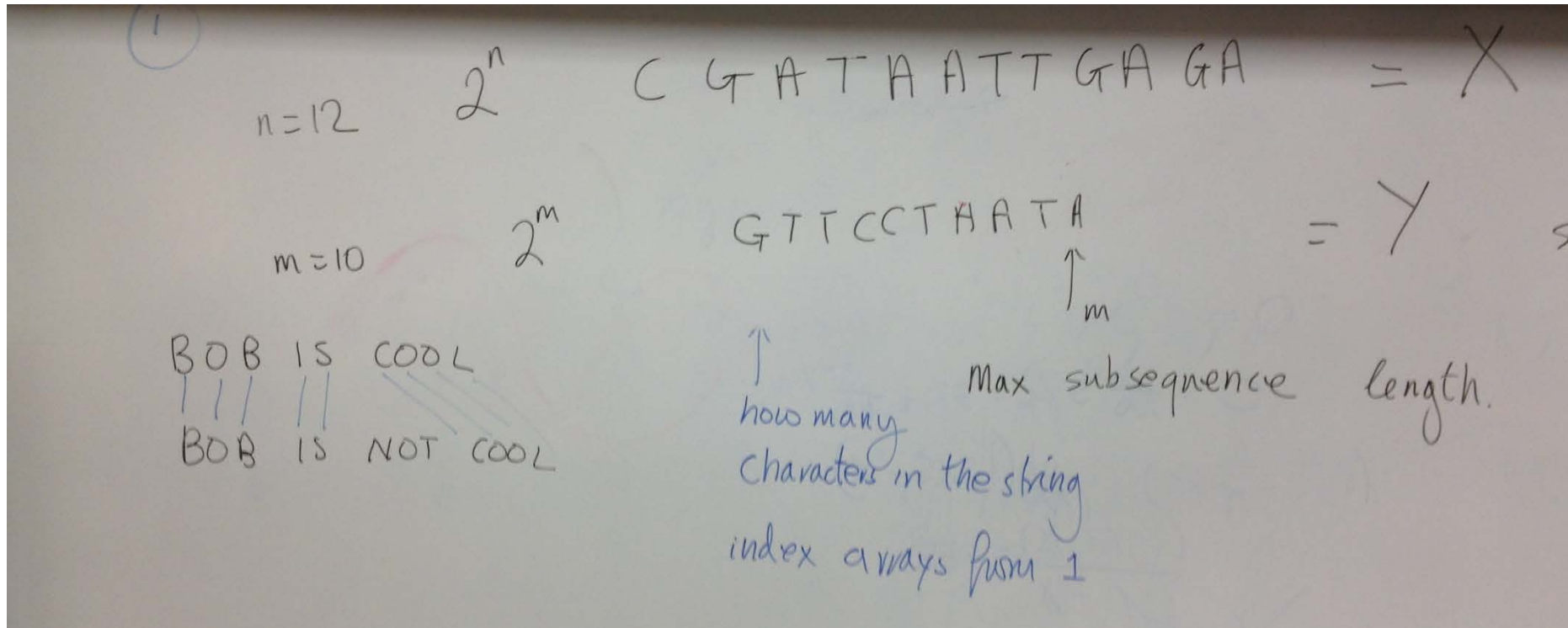


Cs5050 notes 01 21 14



Longest common subsequence problem.

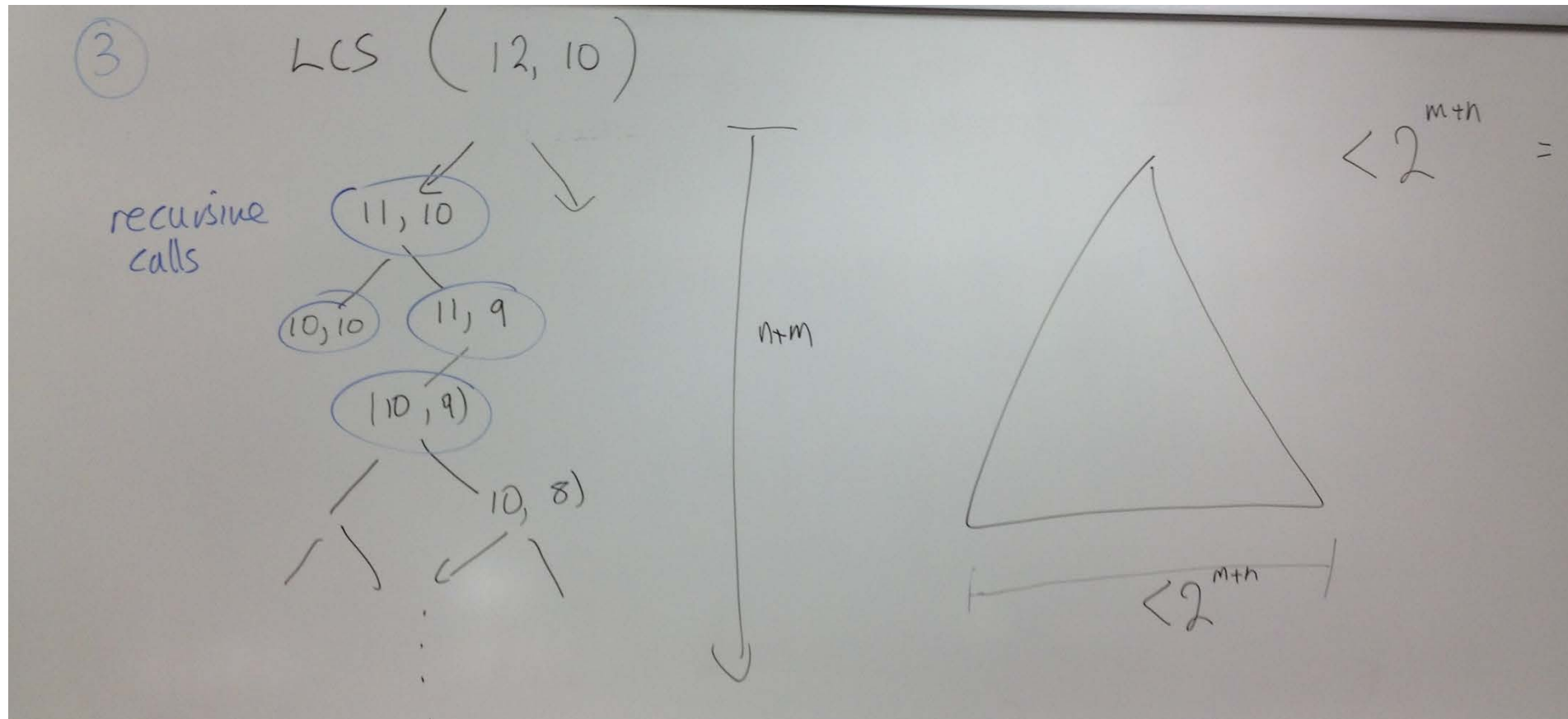
A subsequence of a string A is a string contains a subset of characters from A preserving the order

So for the top string above, CGAA is a subsequence, but AGT is not. For a string of length  $n$ , there are  $2^n$  substrings

Given two strings find the longest subsequence that exists in both strings

Brut force algorithm tries every combination takes  $2^{(n+m)}$  steps





Visualizing the calling tree we get a tree with a branching factor of 2 and a depth between  $\min(m, n)$  and  $m+n$ . Hence, the algorithm will take around  $2^{(n+m)}$  steps.

Cache? Input arguments:

$LCS(n, m)$

$\bigcirc \dots n$

$\bigcirc \dots m$

$(n+1) \cdot (m+1)$  unique calls

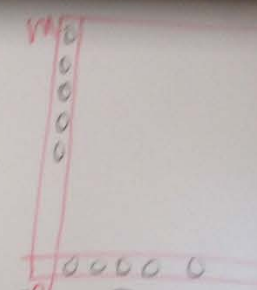
INT LCScache[n+1, m+1]

// initialize with 0

```

for l=1 to n
  for j=1 to m
    if X[l]==Y[j]
      LCScache[l,j]=1+LCScache[l-1,j-1]
    else LCScache[l,j]=max(LCScache[l,j-1],
                           LCScache[l-1,j])
  
```

return LCScache[n,m]



Recall that the arguments to the function represent the possible problem instances.

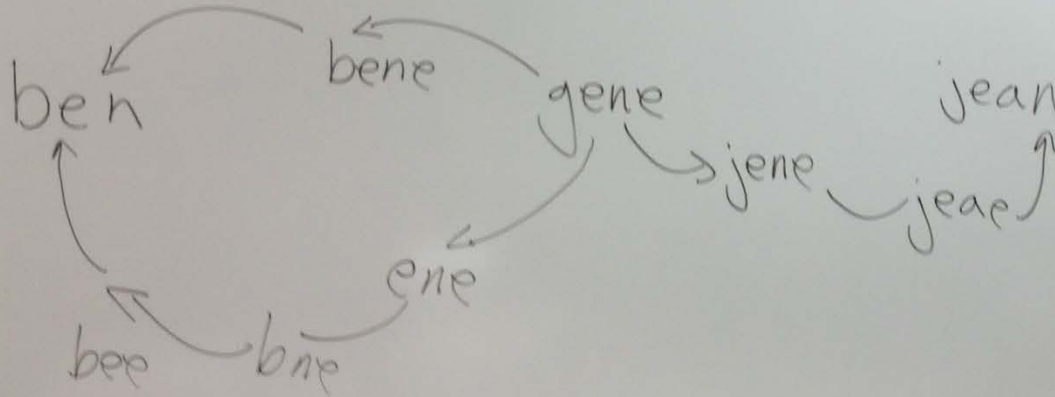
So since the first argument is the index into X and the second the index into Y, then there are only  $(n+1) \cdot (m+1)$  Unique function calls.

So we can create a simple DP algorithm that scans through an integer 2D array

⑤ Min Edit Distance

Given two strings, X Y  
size size m

Find min # of insertions, deletions or substitutions  
 $X \rightarrow Y$

$$X \rightarrow Y$$


INTMED (i, j)

//simple prob, return simple sd.

```
if (i == 0) return j
```

if ( $j=0$ ) return  $i$

Next problem is the minimum edit distance with applications in spell checking and Bioinformatics