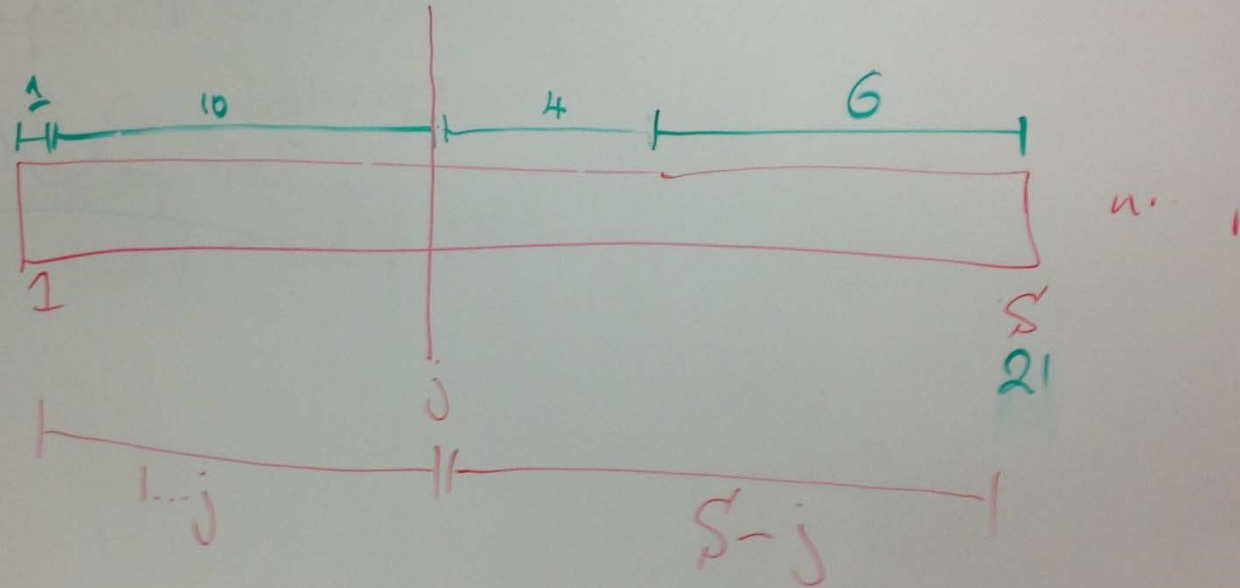
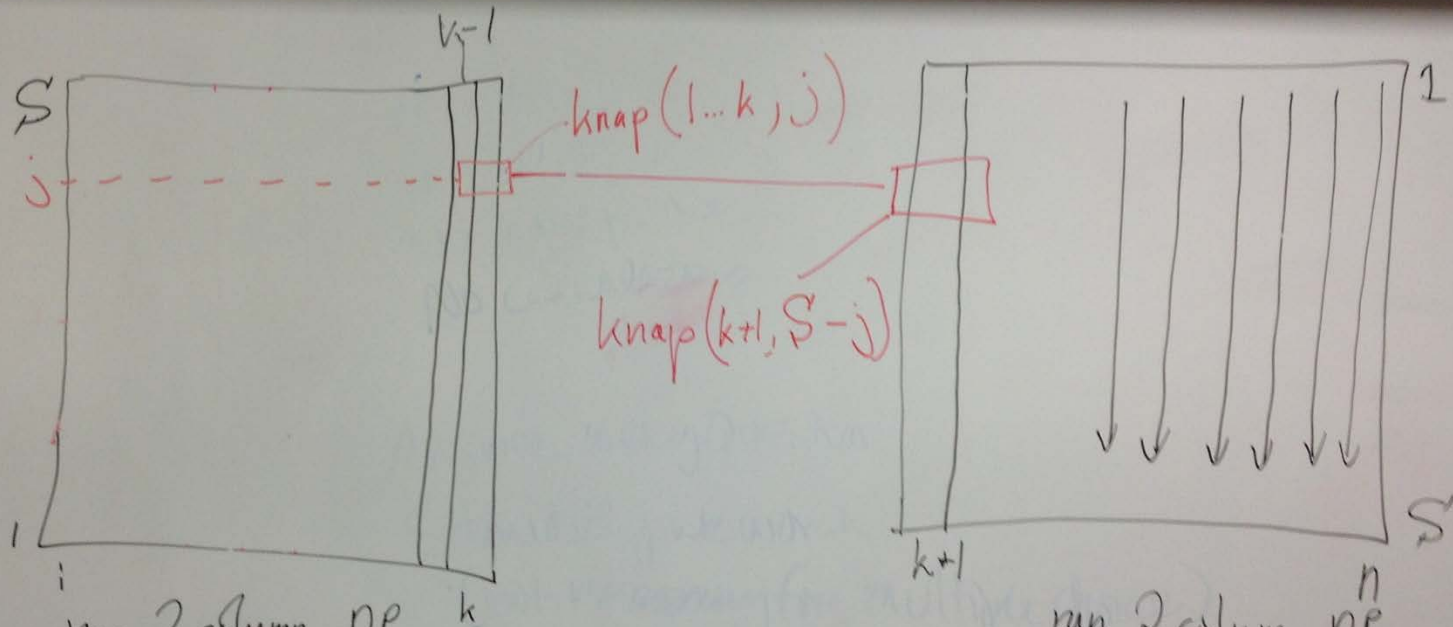


Cs5050 notes 01 30 14

$$\text{knap}(1 \dots k, j) + \text{knap}(k+1 \dots n, S-j)$$





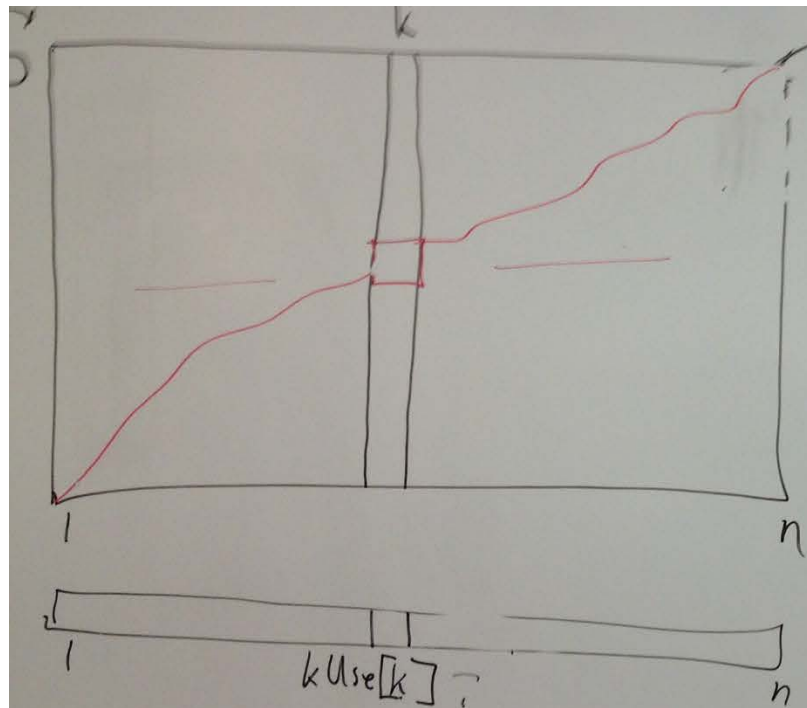
run 2 column DP
from 1 to k
compute $kCache[k, j]$
 $1 \leq j \leq S$

run 2 column DP
from n down to $k+1$
compute $kCache[k+1, j]$
 $1 \leq j \leq S$

Size j best Solution/highest total value

$$\text{bestSize} = \underset{1 \leq j \leq S}{\operatorname{argmax}} \left(k\text{Cache}[k, j] + k\text{Cache}[k+1, S'-j] \right)$$

$$\text{best Value} = \underset{1 \leq j \leq S'}{\operatorname{Max}} \left(k\text{Cache}[k, j] + k\text{Cache}[k+1, S'-j] \right)$$



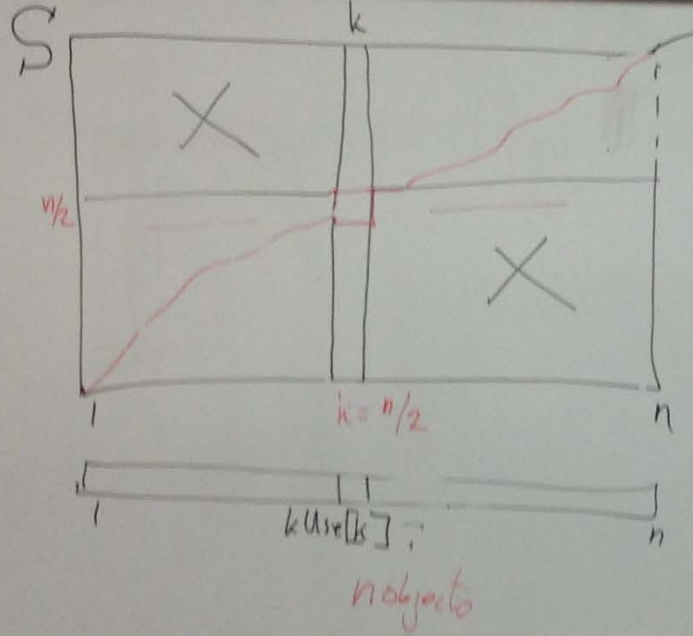
knap(1...n, S) function

kCache[n, S]

kUse[k] = (bestValue == kCache[k-1], best

Size n

Boat k Use



$f(n)$ time it takes to solve problem $n \times n$

$$f(n) = n^2 + 2f(n/2) \quad \text{Split in } 1/2$$

$$f(1) = 1$$

$$f(n) = n^2 + f(n/4) + f(3n/4) \quad k = \frac{n}{4}$$

$$f(n) = n^2 + f(1) + f(n-1)$$

Split off 1
rest