

Step-by-step Logic

1. Initialize pointers:

- `left = 0`
- `right = n - 1`

2. Loop while `left < right`:

- Find `mid = Math.floor((left + right) / 2)`

3. Compare `arr[mid]` with `arr[right]`:

- **Case 1:** `arr[mid] > arr[right]`
 - This means the smallest element is **to the right of mid**.
 - So, set `left = mid + 1`.
- **Case 2:** `arr[mid] <= arr[right]`
 - This means the smallest element is **at mid or to the left of mid**.
 - So, set `right = mid`.

4. Loop ends when `left == right`

- Both will point to the smallest element.

5. Return `left` (or `right`) as the pivot index.

Example Walkthrough

For `[4, 5, 6, 7, 0, 1, 2]`:

- Initial: `left=0, right=6`
- 1st Iteration: `mid=3` (`arr[3]=7, arr[6]=2`)
 - `7 > 2` \Rightarrow move `left` to `mid+1 = 4`
- 2nd Iteration: `left=4, right=6, mid=5`
 - `arr[5]=1, arr[6]=2` $\Rightarrow 1 < 2 \Rightarrow$ move `right = mid = 5`
- 3rd Iteration: `left=4, right=5, mid=4`
 - `arr[4]=0, arr[5]=1` $\Rightarrow 0 < 1 \Rightarrow$ move `right = mid = 4`
- Now `left=right=4`
 - Pivot is at index 4.