

# AM 120 Notes

Bannus Van der Kloot

October 18, 2012

## Contents

<b>1</b>	<b>September 6 Lecture</b>	<b>3</b>
1.1	Geometric interpretation . . . . .	3
1.2	Algorithmic approach . . . . .	4
<b>2</b>	<b>September 11 Lecture</b>	<b>6</b>
2.1	General Example . . . . .	8
<b>3</b>	<b>September 13 Lecture</b>	<b>9</b>
<b>4</b>	<b>September 18 Lecture</b>	<b>12</b>
<b>5</b>	<b>September 20 Lecture</b>	<b>14</b>
5.1	Finite precision . . . . .	15
<b>6</b>	<b>September 25 Lecture</b>	<b>16</b>
6.1	Floating point nonsense . . . . .	16
6.2	Catastrophic cancellation . . . . .	17
<b>7</b>	<b>September 27 Lecture</b>	<b>18</b>
7.1	Ill-conditioned Matrices . . . . .	18
<b>8</b>	<b>October 2 Lecture</b>	<b>20</b>
8.1	Over/underconstrained Systems . . . . .	20
8.2	Vector Spaces . . . . .	21
<b>9</b>	<b>October 4 Lecture</b>	<b>23</b>
<b>10</b>	<b>Oct 9 Lecture</b>	<b>26</b>
<b>11</b>	<b>Oct 11</b>	<b>30</b>
11.1	Fundamental Theorem of Linear Algebra . . . . .	30

<b>12 Oct 16</b>	<b>33</b>
12.1 Fourier series . . . . .	34

# 1 September 6 Lecture

There are two main problems that we will learn how to handle in this class.

1. Find  $x \in R^n$  such that  $Ax = b$ .  $A$  is  $m$  by  $n$  matrix,  $b \in R^n$  vector
2. Find  $x$  and  $\lambda$  such that  $Ax = \lambda x$

**Example** 
$$\begin{aligned} x + 2y &= 3 \\ 4x + 5y &= 6 \end{aligned}$$

$$A = \begin{bmatrix} 1 & 2 \\ 4 & 5 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} 3 \\ 6 \end{bmatrix}$$

There are 3 ways to solve:

1.  $\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 2 & 3 \\ 0 & -3 & -6 \end{bmatrix} \Rightarrow y = 2, x = -1$
2.  $A^{-1} = \frac{1}{\det(A)} \begin{bmatrix} 5 & -2 \\ -4 & -1 \end{bmatrix} \det A = -3$   
 $x = A^{-1}b \Rightarrow \frac{1}{-3} \begin{bmatrix} +3 \\ -6 \end{bmatrix}$
3. Kramer's rule

**Summary** Topics covered in next 3 classes:

1. Geometric interpretation of solving linear systems
2. Matrix notation (LU factorization)
3. Singular cases (no solution, multiple soln's)
4. Efficient way to solve  $Ax = b$  using computers

## 1.1 Geometric interpretation

**Example** Graphical method:

Row interpretation (plot lines on coordinate system):

$$\begin{aligned} 2x - y &= 1 \\ x + y &= 5 \end{aligned} \quad \text{Solution: } x = 2, y = 3$$

Column interpretation:

$$\begin{bmatrix} 2 \\ 1 \end{bmatrix} x + \begin{bmatrix} -1 \\ 1 \end{bmatrix} y = \begin{bmatrix} 1 \\ 5 \end{bmatrix}$$

**Example** 3 by 3 system:

Each row represents a plane:

$$2u + v + w = 5$$

$$4u - 6v + 0 = -2$$

$$-2u + 7v + 2w = 9$$

Remember: inner product of vector with another vector equals 0  $\Rightarrow$  orthogonal.

Column interpretation:

$$\begin{bmatrix} 2 \\ 4 \\ 2 \end{bmatrix} u + \begin{bmatrix} 1 \\ -6 \\ 7 \end{bmatrix} v + \begin{bmatrix} 1 \\ 0 \\ 2 \end{bmatrix} w = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix}$$

**Example** Overdetermined system:

$$\begin{bmatrix} 1 & 1 \\ 2 & 3 \\ 3 & 4 \end{bmatrix} \begin{bmatrix} c \\ d \end{bmatrix} = \begin{bmatrix} 2 \\ 5 \\ 7 \end{bmatrix}$$

Solution:  $c = 1, d = 1$

In 4 dimensions, the rows represent 3-spaces, which are ‘flat’ relative to 4 dimensional space. If we intersect  $(x, y, z, t = 0)$  with  $(x, y, z = 0, t)$ , two three spaces, we get  $(x, y)$  plane.

$$a_1u + a_2v + a_3w + a_4z = b$$

$$A = (a_1|a_2|a_3|a_4)$$

## 1.2 Algorithmic approach

Generalizing to  $n$  by  $n$ . How to solve  $Ax = b$  in a way that scales well? Gaussian elimination (row reduction).

$$2u + v + w = 5$$

$$4u - 6v + 0 = -2$$

$$-2u + 7v + 2w = 9$$

$$\hline 2u + v + w = 5$$

$$\Rightarrow -8v - 2w = -12$$

$$8v + 3w = 14$$

$$\hline 2u + v + w = 5$$

$$\Rightarrow -8v - 2w = -12$$

$$w = 2$$

$$\hline \Rightarrow v = 1, u = 1$$

We need a process that takes:

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}$$

$$x = \begin{bmatrix} u \\ v \\ w \end{bmatrix}$$

$$b = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix}$$

...this  $Ax = b$  problem and transforms it to a  $Ux = \hat{b}$  problem. We can get an upper triangular matrix, and obtain solution by back substitution.

**Problems** One issue that could arise is if the bottom row is all 0s: infinitely many solutions.

## 2 September 11 Lecture

Last class:

- Introduced first central problem of linear algebra: solving linear equations
- Studied column and row interpretation of linear systems
- Introduced Gaussian elimination

**Example** (from previous class) Row/Column interpretation.

$$2u + v + w = 5$$

$$4u - 6v + 0 = -2$$

$$-2u + 7v + 2w = 9$$

Row: Three planes intersecting. Column: linear combination of three vectors

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}$$

We were trying to figure out how to transform matrix  $A$  into an upper triangular matrix.

$$\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix} = \begin{bmatrix} 5 \\ -12 \\ 9 \end{bmatrix}$$
$$\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ -2 & 7 & 2 \end{bmatrix}$$

**Matrix operations** Addition is associative:  $A+B+C = (A+B)+C = A+(B+C)$

Multiplication: dimension  $m \times n$  multiplied by  $n \times p$  results in  $m \times p$  matrix.  $AB \neq BA$ .

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 2 & 3 \\ 7 & 8 \end{bmatrix} = \begin{bmatrix} 7 & 8 \\ 2 & 3 \end{bmatrix}$$
$$\begin{bmatrix} 2 & 3 \\ 7 & 8 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 3 & 2 \\ 8 & 7 \end{bmatrix}$$

Matrix multiplication:

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ a_{21} & a_{22} & a_{23} & \dots & a_{2n} \\ \dots & & & & \\ a_{m1} & a_{m2} & a_{m3} & \dots & a_{mn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \dots \\ x_n \end{bmatrix} = \begin{bmatrix} \sum_{i=1}^n a_{1i}x_i \\ \sum_{i=1}^n a_{2i}x_i \\ \dots \\ \sum_{i=1}^n a_{ni}x_i \end{bmatrix}$$

$$\begin{bmatrix} | & | & & | \\ a_1 & a_2 & \dots & a_n \\ | & | & & | \end{bmatrix} \begin{bmatrix} x_1 \\ \dots \\ x_n \end{bmatrix} = a_1x_1 + a_2x_2 + \dots + a_nx_n$$

$$\begin{bmatrix} | & | & & | \\ a_1 & a_2 & \dots & a_n \\ | & | & & | \end{bmatrix} \begin{bmatrix} | & | \\ b_1 & b_2 \\ | & | \end{bmatrix} = \begin{bmatrix} | & | \\ Ab_1 & Ab_2 \\ | & | \end{bmatrix}$$

**Row reduction** In matrix form

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}$$

1. Subtract 2 times row 1 to row 2

$$\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}_{E_{21}} \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix}_A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ -2 & 7 & 2 \end{bmatrix}$$

2. Subtract -1 times row 1 to row 3

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}_{E_{31}} E_{21}A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 8 & 3 \end{bmatrix}$$

3. Subtract -1 times row 2 to row 3

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix}_{E_{32}} E_{31}E_{21}A = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix}$$

Originally we wanted to solve  $Ax = b$ . Now we have:

$$E_{32}E_{31}E_{21}A = U$$

where  $U$  is an upper triangular matrix.

$$E_{32}E_{31}E_{21}Ax = Ux$$

Let's let  $E_{32}E_{31}E_{21} = L$ . Then, we have

$$\begin{aligned} L^{-1}A &= U \\ A &= LU \\ Ux &= C = E_{32}E_{31}E_{21}b \end{aligned}$$

Now we can solve by back substitution.

$$L^{-1} = E_{32}E_{31}E_{21} = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ -1 & 1 & 1 \end{bmatrix}$$

Matrix inverse properties:

$$(AB)^{-1} = B^{-1}A^{-1}$$

$$(A_1A_2\dots A_n)^{-1} = A_n^{-1}\dots A_2^{-1}A_1^{-1}$$

So we have:

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix} = E_{21}^{-1}E_{31}^{-1}E_{32}^{-1}$$

**Row reduction matrices** A matrix that subtracts  $l$  times row  $j$  from row  $i$  is such that it includes  $-l$  in row  $i$ , column  $j$ .

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix}_L \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix}_U$$

$L$  is lower triangular and  $U$  is upper triangular.

1. Compute LU factorization
2. Solve for  $c$  in  $Lc = b$  (forward substitution)
3. Solve for  $x$  in  $Ux = c$  (back substitution)

We want to solve  $Ax = b$ . We factor to get  $LUx = b$ . First we find  $c$  such that  $Lc = b$

## 2.1 General Example

$$\begin{bmatrix} l_{11} & 0 & 0 \\ l_{21} & l_{22} & 0 \\ l_{31} & l_{32} & l_{33} \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ c_3 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix} \Rightarrow \begin{aligned} c_1 &= b_1/l_{11} \\ c_2 &= b_2 - b_1l_{21}/l_{11} \\ c_3 &= b_3 - l_{31}b_1 - l_{32}(b_2 - b_1l_{21}) \end{aligned}$$

$$\begin{bmatrix} u_{11} & u_{12} & u_{13} \\ 0 & u_{22} & u_{23} \\ 0 & 0 & u_{33} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} c_1 \\ c_2 \\ c_3 \end{bmatrix} \Rightarrow \begin{aligned} x_3 &= c_3/u_{33} \\ x_2 &= \frac{1}{u_{22}}(c_2 - u_{23}c_3/u_{33}) \\ x_1 &= \dots \end{aligned}$$



### 3 September 13 Lecture

#### Announcements

- Matlab tutorials (sections)
- Final projects
  - Adjustment based on class size
  - Pairs
- Assignment 1 due Fri @ 7pm in Pierce 303
- Collaboration policy

From last time:

- Linear equations  $\rightarrow$  Matrix notation
- Column  $j$  of  $AB = Ab_j$

$$A \begin{bmatrix} | & | & \dots & | \\ b_1 & b_2 & \dots & b_n \\ | & | & \dots & | \end{bmatrix} = [Ab_1 \quad Ab_2 \quad \dots \quad Ab_n]$$

- Introduced the  $LU$  factorization of square matrix  $A$  (see general example at end of last lecture)

$$Ax = b \Rightarrow LUx = b$$

1. Find  $LU$
2. Solve for  $c$  in  $Lc = b$
3. Solve for  $x$  in  $Ux = c$

**Example**  $LU$  factorization

$$A = \begin{bmatrix} 1 & 0 & 1 \\ 2 & 2 & 2 \\ 3 & 4 & 5 \end{bmatrix}$$

1. Subtract 2 times row 1 to row 2

$$\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}_{E_{21}} A = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 2 & 0 \\ 3 & 4 & 5 \end{bmatrix}$$

$$E_{21}^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

2. Subtract 3 times row 1 to row 2

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -3 & 0 & 1 \end{bmatrix}_{E_{31}} E_{21} A = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 2 & 0 \\ 0 & 4 & 2 \end{bmatrix}$$

3. Subtract 2 times row 2 to row 3

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -2 & 1 \end{bmatrix}_E 32 E_{31} E_{21} A = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 2 & 0 \\ 0 & 0 & 2 \end{bmatrix}_U$$

$$L^{-1} = E_{32} E_{31} E_{21} = \begin{bmatrix} 1 & 0 & 0 & -2 & 1 & 0 & -3 & -2 & 1 \end{bmatrix}$$

$$\begin{aligned} L^{-1} A &= U \\ L &= E_{21}^{-1} E_{31}^{-1} E_{32}^{-1} \\ L &= \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 3 & 2 & 0 \end{bmatrix} \end{aligned}$$

**Generalizing LU factorization** To  $n \times n$  matrix:

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\ a_{21} & a_{22} & a_{23} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots & \dots \\ a_{n1} & a_{n2} & a_{n3} & \dots & a_{nn} \end{bmatrix}$$

1. Introduce zeros below  $a_{11}$  by subtracting multiples of row 1
2. Use multipliers  $l = \frac{a_{i1}}{a_{11}}$
3. Repeat 1 and 2 for  $a_{22}^*, a_{33}^*, \dots$

Step 1:

$$\begin{array}{c|cccc}
 a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\
 \hline
 0 & a_{22}^* & a_{23}^* & \dots & a_{2n}^* \\
 \dots & \dots & \dots & \dots & \dots \\
 0 & a_{n2}^* & a_{n3}^* & \dots & a_{nn}^*
 \end{array}$$

Step 2:

$$\begin{array}{c|c|ccc}
 a_{11} & a_{12} & a_{13} & \dots & a_{1n} \\
 \hline
 0 & a_{22}^* & a_{23}^* & \dots & a_{2n}^* \\
 \dots & 0 & a_{33}^* & \dots & a_{3n}^* \\
 \dots & \dots & \dots & \dots & \dots \\
 0 & 0 & a_{n3}^* & \dots & a_{nn}^*
 \end{array}$$

How many operations does this algorithm use?

$$\sum_{k=1}^n k^2 - \sum_{k=1}^n k = \frac{n(n+1)(2n+1)}{6} - \frac{n(n+1)}{2}$$

## 4 September 18 Lecture

To review: solving  $Ax = b$ :

1. Find  $LU = A$
2. Solve for  $c$  in  $Lc = b$  (forward substitution)
3. Solve for  $x$  in  $Ux = c$  (back-subst)

Multipliers to find  $U$  are entries of  $L$ .

What is the # of operations needed to get  $LU$  factorization?

$$\approx \frac{n^3 - n}{3}$$

**Forward substitution** Number of operations:

$$(n-1) + (n-2) + \dots + 1 \approx O(n^2)$$

Back substitution is similar process (also  $O(n^2)$ ). Most time consuming place is step 1.

**Algorithm Failure** This  $Ax = b$ :

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

has solution  $\begin{bmatrix} 1 \\ 1 \end{bmatrix}$ . However, our algorithm won't find the answer because it can't switch rows. If the algorithm fails we have two options:

1. We need to rearrange rows
2. No solution
3. Infinitely many solutions

Example of (2):

$$\begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

Example of (3):

$$\begin{bmatrix} 0 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

**Fact**  $\det(A) = \det(LU) = \det(L)\det(U)$

$$\det(U) = \prod_{i=1}^n u_{ii}$$

**Example** Consider this:

$$\begin{aligned} \begin{bmatrix} 0.0001 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} &= \begin{bmatrix} 1 \\ 2 \end{bmatrix} \\ \begin{bmatrix} 0.0001 & 1 \\ 0 & -9999 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} &= \begin{bmatrix} 1 \\ -9998 \end{bmatrix} \\ \Rightarrow x_2 &= \frac{9998}{9999} \\ 0.0001x_1 + \frac{9998}{9999} &= 1 \\ \Rightarrow x_1 &= \frac{10000}{9999} \end{aligned}$$

If we do all of this with limited precision (say 3 digits), we do the following:

$$\begin{aligned} \begin{bmatrix} 0.0001 & 1 \\ 0 & -10^4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} &= \begin{bmatrix} 1 \\ -10^4 \end{bmatrix} \\ \Rightarrow x_2 &= 1 \end{aligned}$$

Then if we use the first equation, we get

$$\Rightarrow x_1 = 0$$

This is called **catastrophic cancellation**.

## 5 September 20 Lecture

First part of AM120 is to solve  $Ax = b$  for arbitrary  $\|A\| = n$ .

$$u_{11} = a_{11}, u_{22} = a_{22}$$

Pseudocode did not have 0s in  $L$  and  $U$ . Second part of code is given  $L$  and  $b$ , should output  $c$ . Third part takes  $U$  and  $c$  and outputs  $x$ .

Assignment 2 Due on Monday morning (9am).

This Doolittle algorithm can fail:

1. If there is a pivot = 0
  - (a) System is singular  $\Rightarrow \det(A) = 0$ . This means there is no solution or infinitely many solutions.
  - (b) We can exchange rows and 'cure' system.

$$\det(A) = \det(L) \det(U) = 1 \prod_{k=1}^n u_{kk}$$

**Example** From last class:

$$\begin{bmatrix} 0.0001 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$$

This had true solution:

$$x_1 = \frac{10000}{9999}, x_2 = \frac{9998}{9999}$$

But with limited precision (three digit arithmetic), we got'

$$x_1 = 0, x_2 = 1$$

What if we switch the rows?

$$\begin{bmatrix} 1 & 1 \\ 0.0001 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \\ 10^{-4} & 1 \end{bmatrix}_L \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 - 2 \cdot 10^{-4}t \end{bmatrix}$$

## 5.1 Finite precision

The computer represents a floating point number with a sign, exponent, and digits for the value itself. When we are talking about  $n$ -digit arithmetic, we are referring to the number of digits storing the value.

$$\begin{aligned} U &= \text{max exponent} \\ L &= \text{lowest exponent} \\ P &= \text{mantissa number of digits} \\ \beta &= \text{base} \end{aligned}$$

For example, for  $L = -1, U = 1, p = 2$  and  $\beta = 10$

$$(\text{sign})d_0.d_1d_2\dots d_p \times 10^{\text{exponent}}$$

$$\begin{array}{ll} \text{largest} & 9.9 \times 10^1 = 99 \\ \text{smallest (non-zero)} & 1.0 \times 10^{-1} = 0.1 \end{array}$$

Examples of real values in floating point systems:

	$\beta$	$P$	$L$	$U$	
IEEE	2	24	-126	123	single
	2	53	-1022	1023	double
HP	10	12	-499	499	

What is the total number of floating point numbers?

$$2(\beta - 1)\beta^{p-1}(u - l + 1) + 1$$

Largest representable number:

$$(\beta - 1).(\beta - 1)\dots(\beta - 1) \cdot \beta^U$$

Smallest number (absolute value):

$$\beta^L(\text{underflow})$$

**Machine precision** Note that the difference between the real number and the floating point number chosen depends on exponent.

$$\forall x \in R, \exists \text{fl}(x) = \hat{x} \text{ such that } |x - \hat{x}| \leq \sum |x|$$

Note that this is not really true for all  $x \in R$  – only within a certain range.  $\epsilon_{\text{mach}}$  is the largest number s.t.  $\text{fl}(1 + \epsilon_{\text{mach}}) > 1$ .

Floating point numbers are not associative:

$$A + (B + C) \neq (A + B) + C$$

## 6 September 25 Lecture

Need to solve  $Ax = b$ . In theory, we can find  $A^{-1}$ . In practice, this is not necessary – it takes too many operations!

$$A \begin{bmatrix} | & | & \cdots & | \\ x_1 & x_2 & \cdots & x_n \\ | & | & \cdots & | \end{bmatrix} = \begin{bmatrix} | & | & \cdots & | \\ e_1 & e_2 & \cdots & e_n \\ | & | & \cdots & | \end{bmatrix} =$$

$$AA^{-1} = I$$

$$x = A^{-1}b$$

In practice, we find  $LU$ . Then we solve  $Lc = b$

$$Lc = b$$

$$c = L^{-1}b$$

Then, we solve  $Ux = c$ :

$$Ux = cL^{-1}Ax = L^{-1}b$$

This algorithm fails if any of the pivots are 0.

**Singular matrix** A matrix is singular if  $\det(A) \neq 0 \iff Ax = b$  has a unique solution.

### 6.1 Floating point nonsense

$$fl(1 + \epsilon) > 1 \Rightarrow \epsilon = \frac{1}{2}\beta^{1-\rho}$$

Machine precision ( $\epsilon_{\text{mach}}$ ): single  $10^{-8}$ ; double  $10^{-16}$ .

$$a = 1.23456 \times 10^2 a = -1.23455 \times 10^{-2}a = -1.11123 \times 10^{-3}a + b + c = (1 \times 10^{-3}) - 1.11123 \times 10^{-3}$$

This is the result if we calculate  $(a+b)+c$ . Notice that if we use 6-digit arithmetic and calculate  $a + (b + c)$ , we get 0! Floating point arithmetic is not associative.

$$\sum_{n=1}^{\infty} \frac{1}{n} \rightarrow$$



1.  $\frac{1}{n}$  i UFL
2.  $\sum_{n=1}^{k_2} \frac{1}{n}$  j OFL
3.  $\text{fl}\left(\left(\sum_{n=1}^{k_3} \frac{1}{n}\right) + \frac{1}{k_3+1}\right) = \sum_{n=1}^{k_3} \frac{1}{n}$

## 6.2 Catastrophic cancellation

If we add a very small number to a big number, sometimes we get the same big number! Like  $1.0 \times 10^8 + 1.0 \times 10^{-9} = 1.0 \times 10^8$ .

$$A + b = A$$

$$A - b = A$$

Partial pivoting minimized catastrophic cancellation. \*Bunch of matrices about the steps of partial pivoting, that I don't even think were accurate\*

**Theorem** For a non-singular and square matrix  $A$ ,  $\exists P$  (permutation matrix) that reorders rows of  $A$  to avoid zeros in the pivot positions.  $Ax = b$  has a unique solution and with the rows ordered "in advance."

$$PA = LU \text{ and } L \text{ and } U \text{ are unique}$$

## 7 September 27 Lecture

**Last Class** Theory: For a non-singular and square matrix  $A$ ,  $\exists P$  (permutation matrix) that reorders rows of  $A$  to avoid zeroes in the pivot positions

$Ax = b$  has a unique solution, and with the rows ordered “in advance.”

$$PA = LU \text{ where } L \text{ and } U \text{ are unique.}$$

Note: If  $A$  is singular, no  $P$  can produce a full set of pivots and elimination fails. Gaussian elimination with partial pivoting: if a pivot is zero, then  $A$  is singular

$$Ax = \begin{bmatrix} 0 & 4 & 1 \\ 1 & 1 & 3 \\ 2 & -2 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 9 \\ 6 \\ -1 \end{bmatrix} = b$$

$$PA = \begin{bmatrix} 2 & -2 & 1 \\ 0 & 4 & 1 \\ 1 & 1 & 3 \end{bmatrix}, P = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

$$\tilde{A} = PA = LU = \begin{bmatrix} 1 & 0 & 0 & 1 & 0.5 & 0.5 & 1 \\ 2 & -2 & 10 & 4 & 10 & 0 & 6 \end{bmatrix}_{\tilde{U}}$$

### 7.1 Ill-conditioned Matrices

The presence of round-off error makes it difficult to identify singular matrices.

$$A = \begin{bmatrix} 1000 & 999 \\ 999 & 998 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 0 \\ .999 & 1 \end{bmatrix}_L \begin{bmatrix} 1000 & 999 \\ 0 & -.001 \end{bmatrix}_U$$

$$Ax = b = \begin{bmatrix} 1999 \\ 1997 \end{bmatrix} \rightarrow x = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

With limited precision (5-digit arithmetic),  $A$  appears singular, because  $0.999(999) = 998.00$ .

$$Ax = \hat{b} = \begin{bmatrix} 1998.99 \\ 1997.01 \end{bmatrix} = b + \delta b = b + 10^{-2} \begin{bmatrix} -1 \\ 1 \end{bmatrix} \Rightarrow x = \begin{bmatrix} 20.97 \\ -18.99 \end{bmatrix}$$

Small change in  $b$  and same  $A$ ,  $x$  changes a lot. We call  $A$  ‘ill-conditioned,’ meaning that its ‘condition number,’  $k(A)$  is big. This is the definition of condition number:

$$\frac{\|\delta x\|}{\|x\|} \leq k(A) \frac{\|\delta b\|}{\|b\|}$$

Hilbert matrices are of this kind: changing  $b$  a little bit,  $x$  changes a lot (they are ill conditioned). Condition number of a singular matrix  $A$  is  $\infty$ .

If  $\frac{\|\delta x\|}{\|x\|} > 1$  we don’t expect to find a solution close to the one we were looking for.

In numerical calculations, singular matrices are indistinguishable from ill-conditioned matrix.

$$\begin{aligned} -\frac{d^2u}{dx^2} &= f(x), 0 \leq x \leq 1 \\ u(0) &= c_1 u(1) = c_2 \end{aligned}$$

$$u(x+h) = u(x) + hu'(x) + h^2 \frac{u''(x)}{2} + \dots + h^k \frac{d^k u}{dx^k}$$

Discretize interval into points  $x_i$ , solve for value of  $u$  at each point. At each point we can solve the problem as a linear algebra problem. We ignore higher terms, and say:

$$u'(x_i) \approx \frac{u(x_i+h) - u(x_i)}{h}$$

$$-u''(x_i) \approx -\frac{u(x_{i+1}) - 2u(x_i) + u(x_{i-1}))}{h^2} = f(x_i)$$

This results in a large matrix:

$$\begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ 0 & -1 & 2 & -1 & \\ & & & \ddots & \\ & & & & 2 & -1 \\ & & & & -1 & 2 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \\ \vdots \\ u_m \end{bmatrix} = \begin{bmatrix} f(x_1) \\ \vdots \\ f(x_m) \end{bmatrix}$$

## 8 October 2 Lecture

**Last class** Condition numbers. Blah blah blah blah.

We have implemented solving  $Ax = b$  the same way that MATLAB's "\ " function works. Now we move on to other things.

### 8.1 Over/underconstrained Systems

$$-\frac{d^2u}{dx^2} = f(x), u(0) = \alpha, u(1) = \beta$$

Approximation of second derivative at discrete point  $x_i$ :

$$-\frac{d^2u(x_i)}{dx^2} \approx \frac{u(x_{i+1}) - 2u(x_i) + u(x_{i-1}))}{h^2}$$

$$h^2 f(x_i) \approx u(x_{i+1}) - 2u(x_i) + u(x_{i-1}))$$

At the boundary:

$$h^2 f(x_1) \approx u(x_2) - 2u(x_1) + \alpha$$

We got this from this, ignoring smaller terms:

$$u(x+h) = u(x) + \frac{du}{dx}(x)h + \frac{d^2u}{dx^2} \frac{h^2}{2} + \dots$$

This yields:

$$\begin{bmatrix} 2 & -1 & & \\ -1 & 2 & -1 & \\ 0 & -1 & 2 & -1 \\ & & \ddots & \ddots \end{bmatrix} \begin{bmatrix} u(x_1) \\ u(x_2) \\ \vdots \\ u(x_m) \end{bmatrix} = \begin{bmatrix} -\alpha + f(x_1)h^2 \\ f(x_2)h^2 \\ \vdots \\ -\beta + f(x_m)h^2 \end{bmatrix}$$

Need to solve  $Ax = b$ . Now we are studying methods that have to do with matrices that are not square! First, we'll say  $m < n$ :

$$\begin{bmatrix} & A & \end{bmatrix}_{m \times n} \begin{bmatrix} x \\ \end{bmatrix}_{n \times 1} = \begin{bmatrix} b \\ \end{bmatrix}_{m \times 1}$$

The left side of  $A$  has an  $m \times m$  square section. This is underdetermined, so there are many solutions. What is  $m > n$ :

$$\begin{bmatrix} A \end{bmatrix}_{m \times n} \begin{bmatrix} x \end{bmatrix}_{n \times 1} = \begin{bmatrix} b \end{bmatrix}_{m \times 1}$$

This is an overconstrained system, which may not have a solution. A real example of this is fitting a line to points in a least-squares sense. If each point is at  $(t_i, y_i)$ , and we want to find a line  $y = mt + b$ , then we solve:

$$\begin{bmatrix} t_1 & 1 \\ t_2 & 1 \\ \vdots & \vdots \\ t_m & 1 \end{bmatrix}_A \begin{bmatrix} m \\ b \end{bmatrix}_x = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}_b$$

## 8.2 Vector Spaces

**Vector spaces** have two operations:

[Example:  $\mathbb{R}^n$ ]

1. if  $x, y \in \mathbb{R}^n$ ,  $x + y = z \in \mathbb{R}^n$
2. if  $x \in \mathbb{R}^n, \alpha \in \mathbb{R}$ ,  $\alpha x \in \mathbb{R}^n$

These properties must hold:

**Addition:** for  $x, y, z \in \mathbb{R}^n$

1. Commutativity:  $x + y = y + x$
2. Associativity:  $x + (y + z) = (x + y) + z$
3.  $\exists!$  zero vector s.t.  $x + 0 = x \forall x \in \mathbb{R}^n$
4.  $\exists! -x \in \mathbb{R}^n$  s.t.  $x + (-x) = 0 \forall x \in \mathbb{R}^n$

**Scalar multiplication**

5.  $1x = x$  : 1 is scalar
6.  $(c_1 c_2)x = c_1(c_2 x)$

---

<sup>1</sup>! indicates there exists some unique zero vector

$$7. \ c(x + y) = cx + cy$$

$$8. \ (c_1 + c_2)x = c_1x + c_2x$$

A **subspace** is a non-empty subset of a vector space that satisfies all these properties and all linear combinations stay in the subspace.

**Example**

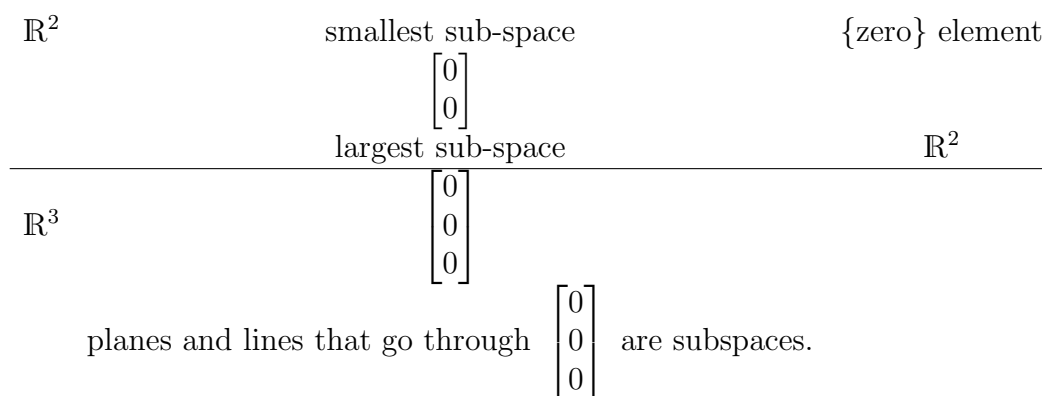
$$\begin{bmatrix} 1 & 0 \\ 5 & 4 \\ 4 & 4 \end{bmatrix} \begin{bmatrix} u \\ v \end{bmatrix} = \begin{bmatrix} 1 \\ 5 \\ 4 \end{bmatrix} u + \begin{bmatrix} 0 \\ 4 \\ 4 \end{bmatrix} v$$

## 9 October 4 Lecture

Use ‘\’ operator in MATLAB from now on unless we specify otherwise. Assignment 4 is now due Tuesday in class.

**Last class** We discussed vector spaces, with the intention of being able to solve any linear algebra problem, whether it be overdetermined or underdetermined. Vector spaces have two operations, addition and scalar multiplication, with 8 axioms. Note: no notion of proximity or distance (no topology).

**Subspace** A non-empty subset of a vector space. Closed under addition and scalar multiplication. ( $x + y \in \text{subspace}$ ,  $ax \in \text{subspace}$ ).



$$A = \begin{bmatrix} 1 & 0 \\ 5 & 4 \\ 2 & 4 \end{bmatrix}, x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}, Ax = \begin{bmatrix} 1 \\ 5 \\ 2 \end{bmatrix} x_1 + \begin{bmatrix} 0 \\ 4 \\ 4 \end{bmatrix} x_2$$

Column spaces of  $A$  (denoted  $C(A)$ ) is the spaces that contains all linear combinations of the columns of  $A$ .

If we have a matrix  $[A]_{M \times N}$  (with  $m$  rows,  $n$  columns), then  $C(A) \in \mathbb{R}^m$ .

$b$  and  $\tilde{b} \in C(A)$ ,  $\exists x$  and  $\tilde{x}$

$$\begin{aligned} Ax &= b & A(x + \tilde{x}) &= b + \tilde{b} \\ A\tilde{x} &= \tilde{b} \\ cb & & A(cx) &= cAx = cb \end{aligned}$$

Null space:

$$A = \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 0 & 0 \end{bmatrix}$$

Null space of  $A$  consists of all vectors  $x$  such that  $Ax = 0$ , denoted  $N(A) \in \mathbb{R}^n$ .

$$Ax = \begin{bmatrix} 1 \\ 5 \\ 2 \end{bmatrix} x_1 + \begin{bmatrix} 0 \\ 4 \\ 4 \end{bmatrix} x_2 = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} x_1 \\ x_2 \end{bmatrix} \in N(A), \begin{bmatrix} 0 \\ 0 \end{bmatrix} \in N(A)$$

**Theorem** If zero is the only element of  $N(A) \Rightarrow$  columns of  $A$  are linearly independent.

If  $N(A) = \{0\}$  and  $A$  is a square matrix  $\Rightarrow \exists! x$  such that  $Ax = b$  for any  $b$ .

Basis for a **vector space**  $V\{v_k\}$ .

1.  $v_k$ 's are linearly independent
2. they span  $V$  (any  $v \in V$  is a linear combination of the basis vectors  $\{v_k\}$ )

$\Rightarrow \exists!$  way to represent any element of  $V$

$\text{dim}(V) = \#$  of basis vectors

The **complete solution** of a linear system of equations

$Ax = b$  is given by  $x = x_p + x_n$  (if it exists)

where

$$Ax_p = b \text{ and } Ax_n = 0$$

$$A(x_p + x_n) = b + 0 = b$$

$$A = \begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix} \quad C(A) = k \begin{bmatrix} 1 \\ 2 \end{bmatrix} \text{ line}$$

$$N(A) = d \begin{bmatrix} 2 \\ -1 \end{bmatrix} \text{ line}$$

$$\begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix} \begin{bmatrix} x_1 & x_2 \end{bmatrix} = Ax = b = \begin{bmatrix} 3 \\ 6 \end{bmatrix}$$

$$x = \begin{bmatrix} 3 \\ 0 \end{bmatrix}_{x_p} + d \begin{bmatrix} 2 \\ -1 \end{bmatrix}_{x_n}$$



**Theorem** : For any  $m \times n$  matrix  $A \exists P$  (permutation) and  $L$  lower triangular matrix and an  $m \times n$  Echelon matrix  $U$  such that  $PA = LU$

$$A = \begin{bmatrix} 1 & 3 & 3 & 2 \\ 2 & 6 & 9 & 7 \\ -1 & -3 & 3 & 4 \end{bmatrix}$$

Let's find  $LU$ :

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 2 & 1 \end{bmatrix}, U = \begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Note that  $L$  is  $m \times m$  square, and  $U$  is also  $m \times n$ . Also, we see that  $U$  has 2 LI columns.

We already knew that the column space of  $A$  lives in  $\mathbb{R}^3$ , so one column had to be dependent. Now we know that only two vectors in the 4 columns are LI, so the column space is a plane. Null space is also a plane, but it lives in  $\mathbb{R}^4$ .

Another example:

$$\begin{aligned} (t_1, y_1)' &= (0, 0) \\ (t_2, y_2) &= (2, 1) \\ \begin{bmatrix} 1 & 0 \\ 1 & 2 \end{bmatrix}_A \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} &= \begin{bmatrix} 0 \\ 1 \end{bmatrix} \end{aligned}$$

$$\begin{aligned} y &= P(t) = x_1 + x_2 t \\ y_1 &= P(t_1) = x_1 + x_2 t_1 \\ y_2 &= P(t_2) = x_1 + x_2 t_2 \end{aligned}$$

We get  $P = \frac{1}{2}t$

## 10 Oct 9 Lecture

The transpose of a matrix  $A$  (denoted  $A^T$ ) is a matrix with columns directly from rows of  $A$  (the  $i$ th row becomes the  $i$ th column of  $A^T$ ).  $(AB)^T = B^T A^T$ .

**Column space of  $A$**  Denoted  $C(A)$ . Contains all linear combinations of the columns of  $A$ .  $C(A)$  is a subspace of  $\mathbb{R}^m$ .

**Complete solution** to the problem  $Ax = b$  can be expressed as  $x = x_p + x_n$ .  $x_n \in N(A)$ ,  $x_p$  a particular solution.

$$A = \begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix}$$

$$Ax_p = \begin{bmatrix} 3 \\ 6 \end{bmatrix}$$

Should know

1. What is a vector space?
2. When a set of vectors are linearly independent?
3. Dimension of a vector space
4. Basis for a vector space

**Theorem**  $N(A)$  contains only the vector iff zero the columns of  $A$  are linearly independent. Therefore, if  $A$  is square,  $\exists! x$  for any  $b$ .

$$A = \begin{bmatrix} 1 & 3 & 3 & 2 \\ 2 & 6 & 9 & 7 \\ -1 & -3 & 3 & 4 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 2 & 1 \end{bmatrix}_L \begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}_U$$

The *row space* of  $A$  is the column space of  $A^T$  ( $C(A^T)$ ). It is a subspace of  $\mathbb{R}^n$ .

The left null space of  $A$  is the null space of  $A^T$ .  $y \in N(A^T)$  if  $A^T y = 0$  iff  $y^T A = 0$ .

Let's calculate the null space of  $A$ . We've broken  $A$  into  $LU$ . Since  $U$  was obtained by adding and subtracting rows of  $A$ , it follows that  $N(A) = N(U)$ .

$$\begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}_U \begin{bmatrix} u \\ v \\ w \\ y \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$$\begin{aligned}
3w + 3y &= 0 \Rightarrow w = -y \\
u + 3v - 3y + 2y &= 0 \\
u &= y - 3v
\end{aligned}$$

$$x \in N(A) = \begin{bmatrix} y - 3v \\ v \\ -y \\ y \end{bmatrix} = y \begin{bmatrix} 1 \\ 0 \\ -1 \\ 1 \end{bmatrix} + v \begin{bmatrix} -3 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

The linear combination of these two vectors generates a plane in  $\mathbb{R}^4$ , the null space of  $A$ .

$$Ax = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}_U \begin{bmatrix} u \\ v \\ w \\ y \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 - 2b_1 \\ (b_3 + b_1) - 2(b_2 - 2b_1) \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 - 2b_1 \\ b_3 - 2b_2 + 5b_1 \end{bmatrix}$$

$$\Longleftrightarrow b_3 - 2b_2 + 5b_1 = 0 \text{ solvability condition}$$

Now, we choose:

$$\begin{aligned}
b &= \begin{bmatrix} 1 \\ 5 \\ 5 \end{bmatrix}, \tilde{b} = \begin{bmatrix} 1 \\ 3 \\ 0 \end{bmatrix} \\
\begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}_U \begin{bmatrix} u \\ v \\ w \\ y \end{bmatrix} &= \begin{bmatrix} 1 \\ 3 \\ 0 \end{bmatrix}
\end{aligned}$$

$$\begin{aligned}
3w + 3y &= 3 \Rightarrow w = 1 - y \\
u + 3v + 3(1 - y) + 2y &= 1 \Rightarrow u = -2 - 3v + y
\end{aligned}$$

Now we can write

$$\begin{bmatrix} 1 & 3 & 3 & 2 \\ 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 0 \end{bmatrix}_U \begin{bmatrix} u \\ v \\ w \\ y \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \\ 0 \end{bmatrix} \Rightarrow x = \underbrace{\begin{bmatrix} -2 \\ 0 \\ 1 \\ 0 \end{bmatrix}}_{x_p} + v \underbrace{\begin{bmatrix} -3 \\ 1 \\ 0 \\ 0 \end{bmatrix}}_{x_n} + y \begin{bmatrix} 1 \\ 0 \\ -1 \\ 1 \end{bmatrix}$$

Remarks:

- The null space of  $A$ ,  $N(A)$ , and the row space of  $A$ ,  $C(A^T)$ , are subspaces of  $\mathbb{R}^n$ .
- The left null space,  $N(A^T)$ , and column space of  $A$ ,  $C(A)$  are subspace of  $\mathbb{R}^m$ .

Transform  $A \xrightarrow[\text{using Gauss elim}]{U}$  we can immediately identify a basis for  $C(A^T)$ .

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}, A^T = \begin{bmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$$

$$C(A) = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \text{ a line in } \mathbb{R}^2$$

$$\text{Row space of } A = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \text{ c lives in } \mathbb{R}^3$$

$N(A)$  lives in  $\mathbb{R}^3 : Ax = 0$ :

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \vec{0}$$

We see that  $x_1$  must be 0, but the other values are free:

$$x_1 = 0$$

$$x_2 = a$$

$$x_3 = b$$

$$x = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} a + \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} b \text{ a plane in } \mathbb{R}^3$$

Left null space  $A^T y = 0$ :

$$\begin{bmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \vec{0}$$

Dimension of row space corresponds with # of linearly independent rows.  $C(A^T)$  dimension is  $r$  ( $r$  linearly independent rows). If we have  $A_{m \times n}$ , then  $r \leq m$  and  $r \leq n$ .  $C(A)$  dimension is  $r$  (even though they live in different spaces).

The dimension of  $N(A)$  is  $n - r$ . Null space lives in  $\mathbb{R}^n$

The dimension of  $N(A^T)$  is  $m - r$ .

## 11 Oct 11

Final projects: think about what brought you to study Applied Math. What problems do you like to solve? We'll find some linear algebra component to it.

Last class:

**Row space of  $A$**   $A \rightarrow U$ , the ' $r$ ' non zero rows are a basis for the row space  $C(A^T)$ . It has dimension  $r$  and it is a subspace of  $\mathbb{R}^n$ . The row space of  $U$  is the same as the row space of  $A$ , since they only differ by linear combinations of rows.

The row space of  $A$  and  $U$  have the same basis.

**Null space of  $A$**   $N(A) = N(U)$ . If  $r$  rows are linearly independent  $\Rightarrow$  there are  $(n - r)$  free variables the dimension of  $N(A) = n - r$ .

Null space definition:

$$x \in \mathbb{R}^n \text{ s.t. } Ax = 0$$

**Column space of  $A$**   $C(A)$ . When we transform  $A$  to  $U$ , the first non-zero elements' index in each row determines which variable of  $x$  will be a pivot variable, suggesting that those indices determine the column space of  $A$ . The column space of  $A$  is *not* the same column space of  $U$ . Dimension of  $C(A) = r$ .  $C(A)$  is a subspace of  $\mathbb{R}^n$ .

**Left null space**

$$\begin{bmatrix} y_1 & y_2 & \cdots & y_m \end{bmatrix} \begin{bmatrix} A^T \\ \end{bmatrix}_{m \times n} = \underbrace{\begin{bmatrix} 0 & \cdots & 0 \end{bmatrix}}_n \text{ it is a subspace of } \mathbb{R}^m$$

$N(A^T)$  is  $A^T y = 0$ . If  $y$  is in the null space of  $A^T$ , then  $y^T$  is in the left null space of  $A$  ( $y^T A = 0$ ).

### 11.1 Fundamental Theorem of Linear Algebra

- $\dim(C(A)) = r$
- $\dim(C(A^T)) = r$
- $\dim(N(A)) = n - r$
- $\dim(N(A^T)) = m - r$

$Ax = b$ . Case:  $(m \leq n)$

$$\begin{bmatrix} & & \\ & A & \\ & & \end{bmatrix}_{m \times n} \begin{bmatrix} \\ x \\ \end{bmatrix}_{n \times 1} = \begin{bmatrix} \\ b \\ \end{bmatrix}_{m \times 1}$$

**Existence:** If  $A$  has the maximum number of linearly independent rows( $= m$ )  $A$  is said to have full “row” rank. There exists at least one solution for any  $b$

In this case,  $A$  has a right inverse. An example:

$$A = \begin{bmatrix} 4 & 0 & 0 \\ 0 & 5 & 0 \end{bmatrix} C = \begin{bmatrix} 1/4 & 0 \\ 0 & 1/5 \\ \alpha & \beta \end{bmatrix}$$

$$AC = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

In this example, the right inverse is not unique! Any values of  $\alpha$  and  $\beta$  will work.

$$ACb = b$$

$$Ax = bx \qquad \qquad \qquad = Cb$$

**Uniqueness**  $m \geq n$

$$\begin{bmatrix} & & \\ & A & \\ & & \end{bmatrix}_{m \times n} \begin{bmatrix} \\ x \\ \end{bmatrix}_{n \times 1} = \begin{bmatrix} \\ b \\ \end{bmatrix}_{m \times 1}$$

If all columns of  $A$  are linearly independent,  $A$  is said to be full “column” rank.

If  $b \in C(A) \exists$  1 unique solution. If  $b \notin C(A)$ , then no solution.

$\mathbb{R}^n$ : normed vector space (with an inner product)

$$\langle x, y \rangle = \sum_{i=1}^n x_i y_i$$

$$\|x\|^2 = \langle x, x \rangle = x^T x$$

Interesting property:

$$\langle x, y \rangle = \|x\| \|y\| \cos \theta$$

$x$  and  $y$  are said to be orthogonal if  $\langle x, y \rangle = 0$ . If we have  $k$  non-zero vectors  $(v_1, \dots, v_k)$  are mutually orthogonal, they are linearly independent. Then we can say that there is only one combination that satisfies the following:

$$c_1 v_1 + c_2 v_2 + \dots + c_k v_k = 0$$

All  $c_i$  must be 0. Proof:

$$v_1^T (c_1 v_1 + c_2 v_2 + \dots + c_k v_k) = 0$$

$$v_1^T c_1 v_1 + v_1^T c_2 v_2 + \dots + v_1^T c_k v_k = 0$$

$$c_1 \|v_1\|^2 = 0$$

Repeat for all  $v_i \Rightarrow c_i = 0 \forall i \Rightarrow \{v_k\}$  are linearly independent.

$$e_i = \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{bmatrix} \quad \text{1 at the } i\text{th element}$$

$\{e_i\}$  is an orthonormal basis of  $\mathbb{R}^n$ . Any vector in  $\mathbb{R}^n$  can be generated as a linear combination of them. If every vector in a subspace  $V$  is orthogonal to every vector in subspace  $W \Rightarrow V$  and  $W$  are said to be orthogonal subspaces.

1. The row space of  $A$ ,  $C(A^T)$  is the orthogonal complement to  $N(A)$ .
2. The column space of  $A$ ,  $C(A)$  is the orthogonal complement to the left null space  $N(A^T)$ .
1. Why is this true? By definition, the null space are vectors  $x$  such that  $Ax = 0$ . If  $A$  is  $m \times n$ . This means that the inner product of every row of  $A$  with  $x$  must be 0. In other words,  $x$  is orthogonal to every row of  $A$ . The rows define the row space  $C(A^T)$ , so (1) is true.
2. The left null space is defined as  $y$  such that  $y^T A = 0$ . If  $A$  is  $m \times n$ , then the inner product of  $y$  and each of the  $n$  columns of  $A$  must be 0.



## 12 Oct 16

### Fundamental theorem of Linear Algebra

- The null space of  $A$ ,  $N(A)$ , is the orthogonal complement of the row space of  $A$  (living in  $\mathbb{R}^n$ )
- The left null space  $N(A^T)$  is the orthogonal complement of the column space of  $A$

$$\min_{\hat{x} \in \mathbb{R}} \|a\hat{x} - b\|$$

$$\begin{aligned} b - \hat{x}a &\perp a \\ \langle a, b - \hat{x}a \rangle &= 0 \\ a^t(b - \hat{x}a) &= 0 \\ \hat{x} &= \frac{a^t b}{a^t a} \\ p &= \hat{x}a \end{aligned}$$

In  $\mathbb{R}^n$ :

$$e_i = \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{bmatrix} \quad 1 \text{ at the } i\text{th element}$$

These  $\{e_k\}$  form an orthonormal basis for  $\mathbb{R}^n$ .

$$\mathbb{R}^5 : v = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{bmatrix} = 1 \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix} + 2 \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} + 3 \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} + 4 \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{bmatrix} + 5 \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

$$v = c_1 e_1 + \dots + c_5 e_5 \text{ where } c_i = \frac{\langle v, e_i \rangle}{\|e_i\|^2}$$

## 12.1 Fourier series

Similarly, we can express an arbitrary function  $f(x)$  as a sum:

$$f(x) = a_0 + \sum_{k=1}^{\infty} a_k \cos kx + b_k \sin kx$$

The basis functions are  $\sin(kx)$  and  $\cos(kx)$ .

Define  $f(x) : [-\pi, \pi] \rightarrow \mathbb{R}$ .

$$\int_{-\pi}^{\pi} |f(x)|^2 < M$$

$M$  is finite. If  $f, g \in L^2([-\pi, \pi]) \Rightarrow \langle f, g \rangle = \int_{-\pi}^{\pi} f(x)g(x)dx$

Functions that map from finite intervals to the real numbers form a vector space.

What should  $a_k, b_k$  and  $a_o$  be?

$$a_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \cos(kx) dx$$

$$b_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \sin(kx) dx$$

Euler's formula:

$$e^{ikx} = \cos kx + i \sin kx$$

$$f(x) = \sum_{k=-\infty}^{\infty} c_k e^{ikx} \text{ with } c_k = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(x) e^{-ikx} dx$$

How do we bring this problem to finite dimension, so we can solve a problem of the form  $Ax = b$ ?

Suppose you have a discrete signal on a fixed interval.

$$f(x) = c_0 + c_1 e^{ikx} + c_2 e^{i2kx} + c_3 e^{i3kx}$$

$$f(0) = f_0 = c_0 + c_1 + c_2 + c_3$$

$$f(\pi/2) = f_1 = c_0 + c_1 i - c_2 - c_3 i$$

$$f(\pi) = f_2$$

$$f(3\pi/2) = f_3$$

This gives us matrix  $A$ :

$$\begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & i^2 & i^3 \\ 1 & i^2 & i^4 & i^6 \\ 1 & i^3 & i^6 & i^9 \end{bmatrix}$$

For use in problem  $Ac = f$ :

$$c = \begin{bmatrix} c_0 \\ c_1 \\ c_2 \\ c_3 \end{bmatrix}, f = \begin{bmatrix} f_0 \\ f_1 \\ f_2 \\ f_3 \end{bmatrix}$$

Sunspots and crazy stuff.