

Notes 9: Numerics

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Introduction

This document is the ninth of a set of notes, this document focusing on numerical questions, random number generation, and linear algebra. The notes are not meant to be particularly complete in terms of useful functions (Google and LLMs can now provide that quite well), but rather to introduce the language and consider key programming concepts in the context of Julia.

Given that, the document heavily relies on demos, with interpretation in some cases left to the reader.

Random number generation

The seed

As usual, it's best to set the random number seed.

```
using Random
Random.seed!(1234);      # Seed number 1234
println(rand(3))
```

```
[0.32597672886359486, 0.5490511363155669, 0.21858665481883066]
```

```
println(rand(5))
```

```
[0.8942454282009883, 0.35311164439921205, 0.39425536741585077, 0.9531246272848422, 0.7955469475347194]
```

```
Random.seed!(1234);      # Re-seed with same number - will give the same sequence of random numbers
println(rand(2))
```

```
[0.32597672886359486, 0.5490511363155669]
```

```
println(rand(6))
```

```
[0.21858665481883066, 0.8942454282009883, 0.35311164439921205, 0.39425536741585077, 0.9531246272848422, 0.7955469475347194]
```

rand has a variety of methods.

```
rand(1:10, 3)
```

3-element Vector{Int64}:

5

8

6

```
rand(['a','b','c'], 5)
```

5-element Vector{Char}:

'c': ASCII/Unicode U+0063 (category Ll: Letter, lowercase)

'a': ASCII/Unicode U+0061 (category Ll: Letter, lowercase)

'a': ASCII/Unicode U+0061 (category Ll: Letter, lowercase)

'b': ASCII/Unicode U+0062 (category Ll: Letter, lowercase)

'c': ASCII/Unicode U+0063 (category Ll: Letter, lowercase)

Generators

The manual discusses the [available generators](#).

The default RNG is Xoshiro256++, but there are details related to random number streams when working in parallel that I have not absorbed yet.

```
Random.default_rng();
Random.seed!(1234)
println(rand(3))
```

```
[0.32597672886359486, 0.5490511363155669, 0.21858665481883066]
```

```
rng = Random.Xoshiro(1234)
println(rand(rng, 3))
```

```
[0.32597672886359486, 0.5490511363155669, 0.21858665481883066]
```

The default used to be the Mersenne Twister (still the default in R and formerly the default in Python/numpy).

```
rng = Random.MersenneTwister(1234);
println(rand(rng, 3))
```

```
[0.5908446386657102, 0.7667970365022592, 0.5662374165061859]
```

Distributions

As with distributions in SciPy in Python, we first define the distribution of interest and then carry out distributional operations with it.

```
using Distributions
beta_dist = Beta(2, 5)
beta_samples = rand(beta_dist, 10)
```

```
10-element Vector{Float64}:
 0.4542940512832609
 0.27004146470644036
 0.3870672630509891
 0.21116529046475746
 0.5041249425689261
 0.2784321400052276
 0.6829933994934323
 0.1295297689806119
 0.43795698919772685
 0.39395847238977155
```

```
using Plots
grid = 0:.01:1
plot(grid, pdf.(beta_dist, grid))
pdf.(beta_dist, beta_samples[1:5])
```

```
5-element Vector{Float64}:
 1.2086291289041646
 2.300088115277408
 1.6389245754521693
 2.4529451430316644
 0.9144258514892248
```

```
logpdf.(beta_dist, beta_samples[1:5])
```

```
5-element Vector{Float64}:
```

```
0.1894867660152757
0.8329474331914342
0.4940402800647119
0.8972894018105007
-0.08945889536877027
```

Floating point issues

Integer and floating point types

64 bit integers can represent $-2^{63} \dots 2^{63}$. 32 bit integers can represent $-2^{31} \dots 2^{31}$.

Similarly for 16 and 128 bit integers.

Values outside that range *overflow*.

```
xi::Int64 = 2^62
```

```
4611686018427387904
```

```
xi::Int64 = 2^70 # Overflows
```

```
0
```

```
xi::Int64 = 2^63 # Just overflows.
```

```
-9223372036854775808
```

```
yi::Int128 = 2^63 # Hmmm.
```

```
-9223372036854775808
```

```
yi::Int128 = Int128(2)^63
```

```
9223372036854775808
```

```
Int64(yi)
```

```
LoadError: InexactError: trunc(Int64, 9223372036854775808)
```

```
InexactError: trunc(Int64, 9223372036854775808)
```

Stacktrace:

```
[1] throw_inexacterror(f::Symbol, ::Type{Int64}, val::Int128)
    @ Core ./boot.jl:634
[2] checked_trunc_sint
    @ ./boot.jl:656 [inlined]
[3] toInt64
    @ ./boot.jl:705 [inlined]
[4] Int64(x::Int128)
    @ Core ./boot.jl:784
[5] top-level scope
    @ In[19]:1
```


Computation with `BigFloats` will be slow, so you wouldn't want to do matrix operations with a matrix of them.

Floating point details

With `Float64`, any number is stored as a base 2 number of the form:

$$(-1)^S \times 1.d \times 2^{e-1023} = (-1)^S \times 1.d_1 d_2 \dots d_{52} \times 2^{e-1023}$$

where the computer uses base 2, $b = 2$, (so $d_i \in \{0, 1\}$) because base-2 arithmetic is faster than base-10 arithmetic. The leading 1 normalizes the number; i.e., ensures there is a unique representation for a given computer number. This avoids representing any number in multiple ways, e.g., either $1 = 1.0 \times 2^0 = 0.1 \times 2^1 = 0.01 \times 2^2$. For a double, we have 8 bytes=64 bits. Consider our representation as (S, d, e) where S is the sign. The leading 1 is the *hidden bit* and doesn't need to be stored because it is always present. In general e is represented using 11 bits ($2^{11} = 2048$), and the subtraction takes the place of having a sign bit for the exponent. (Note that in our discussion we'll just think of e in terms of its base 10 representation, although it is of course represented in base 2.) This leaves $p = 52 = 64 - 1 - 11$ bits for d .

The representations for floating point numbers with more or fewer bits than 64 is similar in structure but with a different split of bits used for the magnitude and the precision.

53 bits of precision in base 10 is about 16 digits. And 11 bits for the magnitude corresponds with when over/underflow occur. For 32-bit floats, we have both less precision and we more easily over/underflow.

Overflow

Integer numbers can be represented exactly by `Float64` up to 2^{53} .

```
function pri(x)
    @sprintf("%.20i", x)
end
```

```
pri(2.0^52)
```

```
"4503599627370496"
```

```
pri(2.0^52 + 1)
```

```
"4503599627370497"
```

```
pri(2.0^53)
```

```
"9007199254740992"
```

```
pri(2.0^53 + 1)
```

```
"9007199254740992"
```

```
pri(2.0^53 + 2)
```

```
"9007199254740994"
```

```
pri(2.0^63)
```

```
"9223372036854775808"
```

```
pri(2.0^70)      # No overflow here unlike Int64.
```

```
"1180591620717411303424"
```

```
pri(12345678123456781234.0)  # Not exact.
```

```
"12345678123456780288"
```

Float64 overflow is not until $\sim 2^{1023} \approx 10^{308}$.

```
function prf(x)
    @sprintf("%.20f", x)
end
```

```
prf(2.0^1022)
```

```
"4494232837155789769323262976972561834044942447355766431835752028943316895137524078317711933060188400"
```

```
prf(2.0^1023)
```

```
"8988465674311579538646525953945123668089884894711532863671504057886633790275048156635423866120376801"
```

```
prf(10.0^308)
```

```
"100000000000000000010979063629440455417404923096773118463368106829031575854049114915371633289784946888"
```

```
prf(2.0^1024)
```

```
"Inf"
```

But Int64s that big do overflow.

```
pri(10^307)
```

```
"00000000000000000000"
```

```
pri(10^310)
```

```
"00000000000000000000"
```

Implications for comparisons and calculations

```
prf(1-2/3)
```

```
"0.33333333333333337034"
```

```
prf(1/3)
```

```
"0.33333333333333331483"
```

```
prf(2/3-1/3)
```

```
"0.33333333333333331483"
```

```
prf(0.3 - 0.2)
```

```
"0.09999999999999997780"
```

```
prf(0.1)
```

```
"0.10000000000000000555"
```

```
0.3 - 0.2 == 0.1
```

```
false
```

Here's an example of catastrophic cancellation:

```
prf(123456781234.56)
```

```
"123456781234.55999755859375000000"
```

```
prf(123456781234.00)
```

```
"123456781234.00000000000000000000"
```

```
prf(123456781234.56 - 123456781234.00)
```

```
"0.55999755859375000000"
```

And here the precision is that of the larger magnitude number:

```
prf(1.0 + 1e-8)
```

```
"1.00000000999999993923"
```

```
prf(1.0 + 1e-17)
```

```
"1.000000000000000000000000"
```

```
1.0 + 1e-17 == 1.0
```

```
true
```

Avoid multiplying/dividing many numbers

How large does n need to be to see underflow here? In other cases it could overflow.

```
using Distributions
normd = Normal(0, 1)
samples = rand(normd, n);
prod(pdf.(normd, samples)) # Log-likelihood/log-density
```

```
sum(logpdf.(normd, samples))
```


Linear algebra

We'll see that linear algebra operations heavily exploit Julia's multiple dispatch system, calling the most appropriate method for different kinds of input matrices.

Arithmetic

```
A = rand(3, 3);
B = rand(3, 3);
A + B;    # Element-wise addition
A .+ B;   # Element-wise addition
A * B;    # Matrix multiplication
A .* B;   # Element-wise multiplication
```

Solving systems of equations / inversion

```
using LinearAlgebra

A = rand(3, 3);

AtA = A' * A;    # A' * A
b = rand(3);
AtA \ b          # Solve system of equations (AtA^{-1} b)
```

```
3-element Vector{Float64}:
-1.3254878978759739
 5.5034206079487324
 0.30748299868104345
```

```
inv(AtA) * b    # Not as efficient
```

```
3-element Vector{Float64}:
-1.3254878978759748
 5.503420607948734
 0.30748299868104345
```

```
eigvals(AtA)
```

```
3-element Vector{Float64}:
 0.03992762277861558
 0.22181922851410624
 2.0138442890779333
```

```
det(AtA)        # Be careful of over/underflow!
```

```
0.017836043877867285
```

```
logdet(AtA)
```

```
-4.026533932146638
```

```
tr(A)
```

```
1.5834783240781707
```

```
chol = cholesky(AtA)
```

```
Cholesky{Float64, Matrix{Float64}}
```

```
U factor:
```

```
3×3 UpperTriangular{Float64, Matrix{Float64}}:
```

```
 0.957996  0.309353  0.873741
           0.213918  0.168095
           0.651687
```

We have $(U^{\text{top}}U)^{-1}b = U^{-1}(U^{\text{top}})^{-1}b$, which in Julia is implemented internally.

```
chol \ b          # Automatically exploits triangularity.
```

```
3-element Vector{Float64}:
```

```
-1.3254878978759743
 5.503420607948734
 0.30748299868104323
```

```
chol.U \ (chol.L \ b)  # Manual equivalent solution.
```

```
3-element Vector{Float64}:
```

```
-1.325487897875974
 5.503420607948734
 0.3074829986810431
```

```
typeof(chol)          # Special type of object
```

```
Cholesky{Float64, Matrix{Float64}}
```

```
typeof(chol.U)         # Special kind of matrix.
```

```
UpperTriangular{Float64, Matrix{Float64}}
```

This works in Julia but is having problems via Quarto, so I'll just paste in some timings.

```
using BenchmarkTools
```

```
n = 5000;
```

```
A = rand(n, n);
```

```
AtA = A'A;
```

```
b = rand(n);
```

```
@btime AtA \ b;          # Solve system of equations (AtA^{-1} b) via Gaussian elimination (LU)
# 633.046 ms (6 allocations: 190.81 MiB)
```

```
@btime inv(AtA) * b;     # Not as efficient
```

```
# 2.538 s (8 allocations: 193.25 MiB)
```

```
@btime cholesky(AtA) \ b; # Best if matrix is positive definite.
```

```
# 546.957 ms (5 allocations: 190.77 MiB)
@btime chol = cholesky(AtA); chol.U \ (chol.L \ b); # As good but verbose.
# 520.021 ms (3 allocations: 190.73 MiB)
```

In principle, the Cholesky approach should involve $n^3/6$ calculations and the Gaussian elimination $n^3/3$, but we don't see a two-fold difference here in practice.

Spectral (eigen) decomposition

```
Γ = eigvecs(AtA);
    = eigvals(AtA);
result = Γ * diagm() * Γ';

result[1:3,1:3]
```

```
3×3 Matrix{Float64}:
 0.917756  0.296359  0.83704
 0.296359  0.14146  0.306253
 0.83704   0.306253  1.21638
```

```
AtA[1:3,1:3]
```

```
3×3 Matrix{Float64}:
 0.917756  0.296359  0.83704
 0.296359  0.14146  0.306253
 0.83704   0.306253  1.21638
```

```
result == AtA
```

```
false
```

```
isapprox(result, AtA)
```

```
true
```

```
result AtA # \approx TAB
```

```
true
```

Don't forget floating point issues

```
n=500;
A = rand(n, n);
det(A'A)
```

```
Inf
```

```
logdet(A'A)
```

```
1364.3425695000474
```

Here's a positive definite matrix (mathematically) that has all real, positive eigenvalues. On a computer, it's not positive definite, and therefore not invertible/full rank.

```
xs = 0:99
# Compute distance matrix.
dists = abs.(xs .- xs') # Using broadcasting with ' (an "outer" operation).
# Create correlation matrix.
corr_matrix = exp.-((dists/10).^2)
# Compute eigenvalues and get last 20
eigvals(corr_matrix)[1:20]
```

20-element Vector{Float64}:

```
-5.050107193681261e-15
-3.975288045003424e-15
-3.0343039148117795e-15
-2.3309913036821875e-15
-1.5093549661154687e-15
-1.3511621460923302e-15
-1.2848549714671506e-15
-1.0472866233731733e-15
-8.644575073124488e-16
-8.364136467422076e-16
-7.989065965702987e-16
-6.829697931404996e-16
-6.091008178138255e-16
-6.030845253775599e-16
-5.700968077566563e-16
-5.37906121536584e-16
-4.615826595950015e-16
-4.4365869197860045e-16
-4.0894133652379844e-16
-3.580464402267892e-16
```

```
chol = cholesky(corr_matrix);
```

LoadError: PosDefException: matrix is not positive definite; Cholesky factorization failed.
PosDefException: matrix is not positive definite; Cholesky factorization failed.

Stacktrace:

```
[1] checkpositivedefinite
   @ /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/factorization.jl:67 [inlined]
[2] cholesky!(A::Hermitian{Float64, Matrix{Float64}}, ::NoPivot; check::Bool)
   @ LinearAlgebra /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:
[3] cholesky!
   @ /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:267 [inlined]
[4] cholesky!(A::Matrix{Float64}, ::NoPivot; check::Bool)
   @ LinearAlgebra /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:
[5] cholesky! (repeats 2 times)
```

```

@ /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:295 [inlined]
[6] cholesky
@ /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:401 [inlined]
[7] cholesky(A::Matrix{Float64})
@ LinearAlgebra /system/linux/julia-1.10.4/share/julia/stdlib/v1.10/LinearAlgebra/src/cholesky.jl:
[8] top-level scope
@ In[75]:1

```

One should be able to use an approximate pivoted Cholesky that sets diagonal elements to zero corresponding to the rank deficiency. I'm having trouble seeing how to do that.

This does seem to work to solve the system of equations. We'd have to investigate to know what Julia is doing behind the scenes, but it's probably using a pivoted LU decomposition.

```

b = rand(100);
out = corr_matrix \ b;

```

Smart factorization

```

n = 500
A = rand(n, n);

```

```

typeof(factorize(A'A))

```

```

Cholesky{Float64, Matrix{Float64}}

```

```

typeof(factorize(A))          # I wouldn't have guessed that A is invertible!

```

```

LU{Float64, Matrix{Float64}, Vector{Int64}}

```

```

typeof(factorize(A[:,1:10]))

```

```

QRPivoted{Float64, Matrix{Float64}, Vector{Float64}, Vector{Int64}}

```

The orthogonal matrices generated by certain factorizations can most efficiently be worked with by having them treated as “matrix-backed, function-based linear operators”.

```

QRresult = qr(A[:,1:10]);
typeof(QRresult.Q)

```

```

LinearAlgebra.QRCompactWYQ{Float64, Matrix{Float64}, Matrix{Float64}}

```

```

QRresult.Q * rand(10);

```

Use known structure!

```

using BenchmarkTools

```

```

n = 5000
A = rand(n, n);
AtA = A'A;

```

```

b = rand(n);
chol = cholesky(AtA);
typeof(chol.U)
@btime chol.U \ b;

6.213 ms (3 allocations: 39.12 KiB)
U_dense = Matrix(chol.U);
@btime U_dense \ b;

14.940 ms (2 allocations: 39.11 KiB)
@btime logdet(chol.U);

105.511 s (2 allocations: 32 bytes)
@btime logdet(U_dense);

644.446 ms (5 allocations: 190.77 MiB)
@btime sum(log.(diag(chol.U)));

99.066 s (7 allocations: 78.27 KiB)
using SparseArrays

n = 5000;
A = Matrix{Float64}(1.0I, n, n);
A[1,3] = 7.7;
A[5,9] = 2.3;
sA = sparse(A);

sizeof(A)

200000000

sizeof(sA) # Not helpful given the pointers involved.

40
# sA. TAB # What are the components of `sA`?
sizeof(sA.colptr) + sizeof(sA.nzval) + sizeof(sA.rowval)

120040

b = rand(n);
@btime A * b;

10.897 ms (2 allocations: 39.11 KiB)
@btime sA * b;

20.305 s (2 allocations: 39.11 KiB)

```