

# Weak Hiding for C++ Concepts and a Generic Way to Understand Name Binding

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# Objectives

- ❖ A simpler way to describe name binding.
- ❖ A more powerful way to describe name binding.
- ❖ Applying this description to C++ concepts.



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# Our Tools

- ❖ scope combinators,
- ❖ an abstraction of programming languages, and
- ❖ a distinction between **name lookup** and **name resolution**.



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# Outline

## ① Current name binding mechanisms:

1. Across languages.
2. Specific to C++:
  - a. Argument-dependent lookup (ADL).
  - b. Uses of operators.

## ② A new scoping rule: **Weak hiding**

- Implementation: two-stage name binding (**Bind<sup>x2</sup>**)
- Application: Preventing code breakage when switching to concepts-enabled libraries.



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# Name Binding

- ❖ Deciding what a particular use of a name refers to.
- ❖ Use of a name = **a *reference***; What it refers to = **a *declaration***.

```
void foo();  
void foo(int, double);
```

```
void test() {  
    foo();  
}
```



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# Name Binding

- ❖ Binds a *reference* to a *declaration*.

```
void foo();  
void foo(int, double);
```

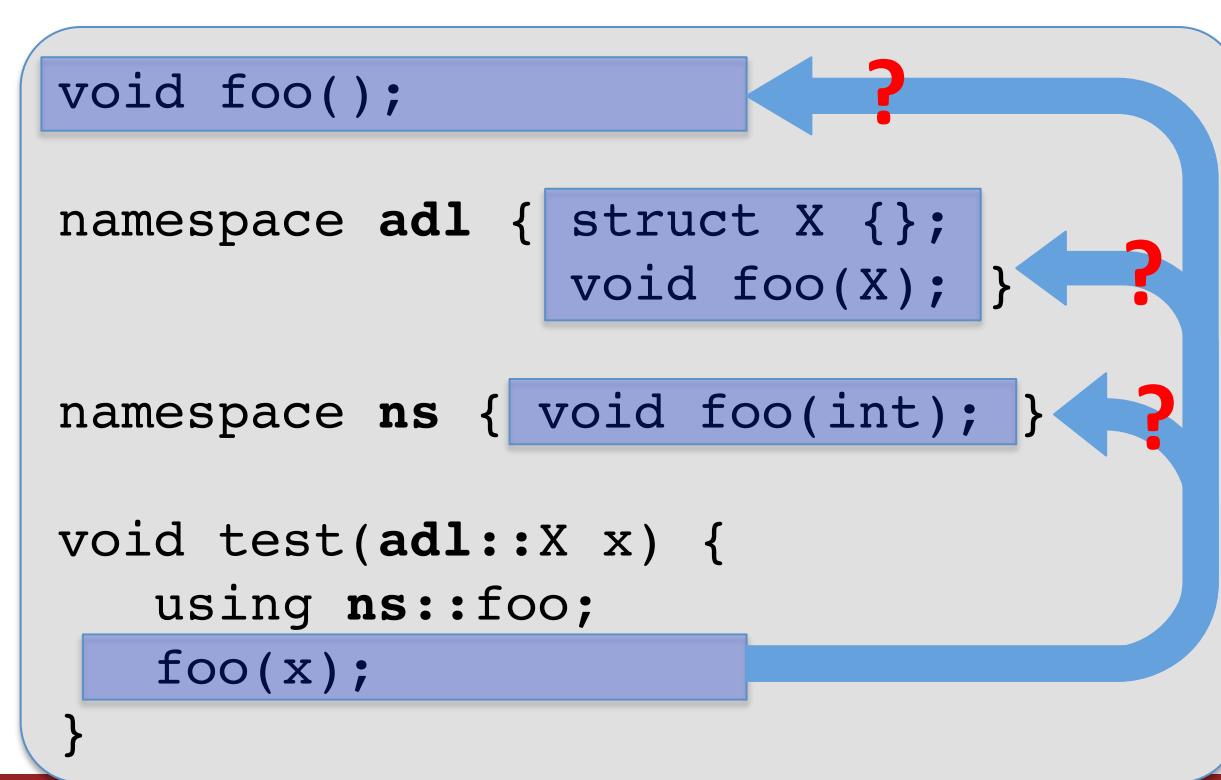
```
namespace ns { void foo(int); }
```

```
void test() {  
    using ns::foo;  
    foo();  
}
```



# Name Binding

- ❖ Binds a *reference* to a *declaration*.



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# Name Binding

- ❖ Binds a *reference* to a *declaration*.

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); } ?  
  
void foo(int) { } ?  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

The diagram illustrates the process of name binding. It shows three code snippets: a concept definition, a function declaration, and a template definition. Blue arrows point from each declaration to its corresponding definition, with red question marks indicating the binding relationship.



# Name Binding

- ❖ Depends on a **scope**, nested or combined with other scopes.

```
void foo(); ::  
  
namespace adl { struct X {}; adl  
    void foo(X); }  
  
namespace ns { void foo(int); ns  
  
void test(adl::X x) {  
    using ns::foo; test  
    foo(x); }  
}
```



# Name Binding

- ❖ Differs between **languages, designs, or kinds of references**:

- argument-dependent lookup (ADL),
- simple function calls,
- uses of operators,
- uses of types,
- C++ Multimethods,
- C++,
- Haskell,
- etc...

```
void foo();  
  
namespace adl { struct X {};  
    void foo(X); }  
  
namespace ns { void foo(int); }  
  
void test(adl::X x, adl::X y) {  
    using ns::foo;  
    foo(x + y); }  
}
```



# Standards for Name Binding

- ❖ Can be challenging to understand, analyze, and compare.
- ❖ Are intertwined with other sections of the standard.

## ▪ C++ Name Binding Specification, at a glance:

▪ 3.4 – Name lookup	13 pages, 45+ paragraphs
▪ 14.6 – Name resolution	13 pages, 50+ paragraphs
▪ 13 – Overloading	30 pages, 140+ paragraphs
▪ 11 – Member access	10 pages, 35+ paragraphs



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# Standards for Name Binding

- ❖ Typically depend on details of languages and the varied features they support.

## ConceptC++:

```
template<Foo2 T>
void test(T) {
    foo();
}
```

- ⇒ C++ templates + Concepts
- ⇒ (Non)dependent names
- ⇒ Overload resolution
- ⇒ Namespaces + using declarations
- ⇒ C++ compilers, e.g., GCC, Clang, ...
- ⇒ Etc...

## Haskell:

```
test::Foo2 t => t -> ()
test _ = foo()
```

- ⇒ Hindley-Milner type system + Type classes
- ⇒ Type inference + Dependency analysis
- ⇒ Dictionary-passing style
- ⇒ Modules and import declarations
- ⇒ Haskell compilers, e.g., GHC, Hugs, NHC,
- ...
- ⇒ Etc...



# Standards for Name Binding

- ❖ Are understood differently by different compilers.

```
namespace adl { struct X {};  
                 void foo(X); }  
  
namespace ns { void foo(int); }  
  
void bar(adl::X x) {  
    using ns::foo;  
    foo(x);  
}  
  
void baz(adl::X x) {  
    void foo(int);  
    foo(x);  
}
```

In bar():

Success. ADL enabled.

Binds **foo(x)** to **adl::foo()**.

In baz():

GCC, Clang, Comeau:

Failure. ADL Disabled.

**foo(x)** does not match **baz :: foo(int)**.

Intel:

Success.



# Standards for Name Binding

- ❖ Are understood differently by different compilers.

```
struct X {}; struct Y {};  
  
void operator+(X, X) { }  
void operator+(X, Y) { }  
  
void test(X x, Y y) {  
    void operator+(X, X);  
    x + x;  
    x + y;  
    operator+(x, x);  
    operator+(x, y);  
}
```

GCC, Clang, MS Visual Studio editor :

**x + y succeeds.**

**operator+(x, y) fails.**

Intel, MS Visual Studio:

**x + y succeeds.**

**operator+(x, y) succeeds.**

Comeau:

**x + y fails.**

**operator+(x, y) fails.**



# Standards for Name Binding

- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.

```
concept Foo2<typename T> =  
    requires (T a, T b) { foo(a, b); } ?  
  
void foo(int) { } ?  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```



# Name Binding Specifications

- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.

```
concept Foo2<typename P> =
    requires (P a, P b) { foo(a, b); } Foo2<P>

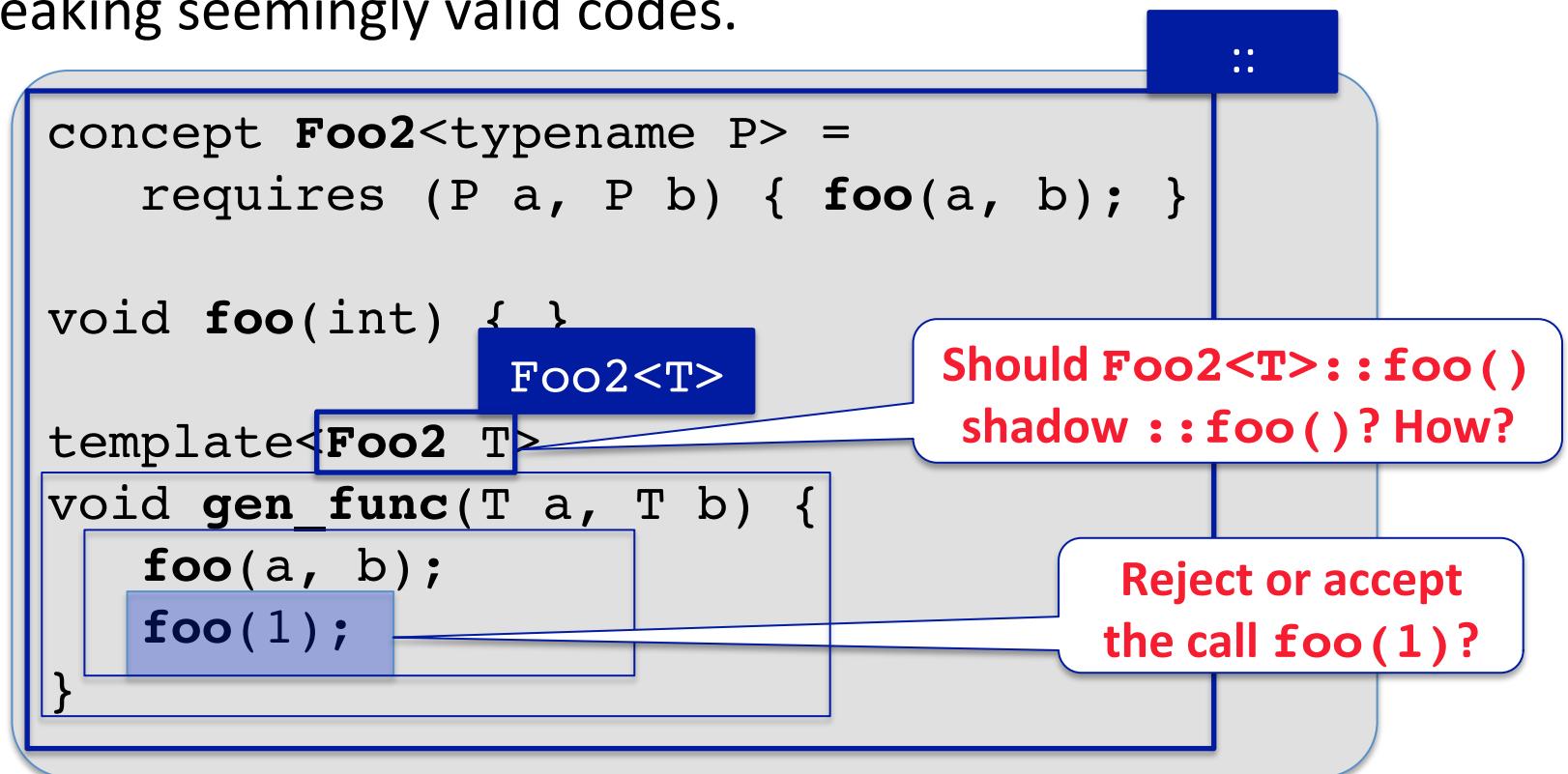
void foo(int) {}

template<Foo2 T>
void gen_func(T a, T b) { template...
    foo(a, b);
    foo(1);
}
```



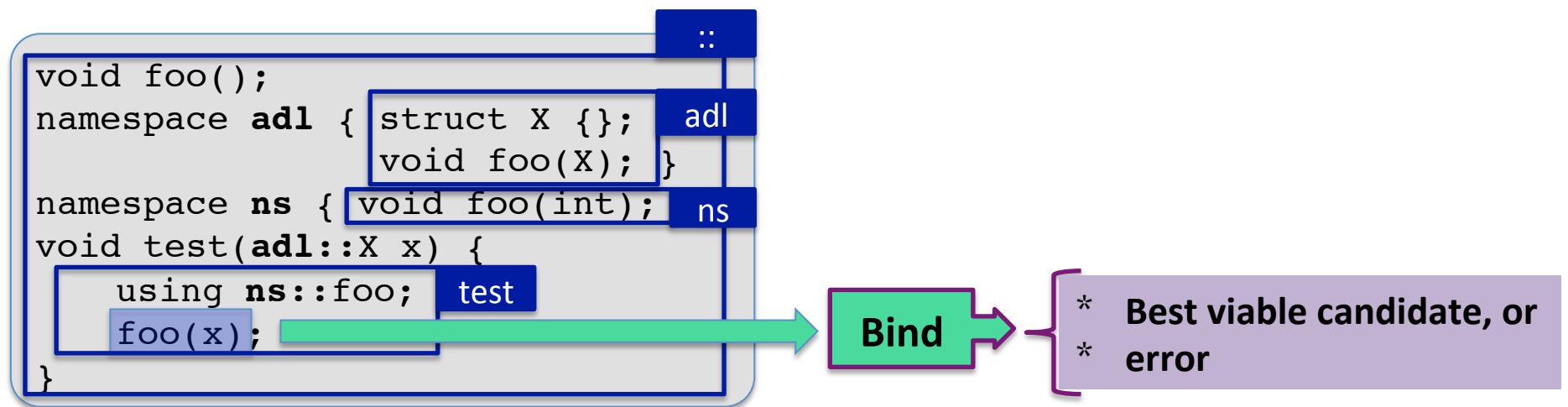
# Name Binding Specifications

- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.



# Our Name Binding Framework

- ❖ Specifies name binding.

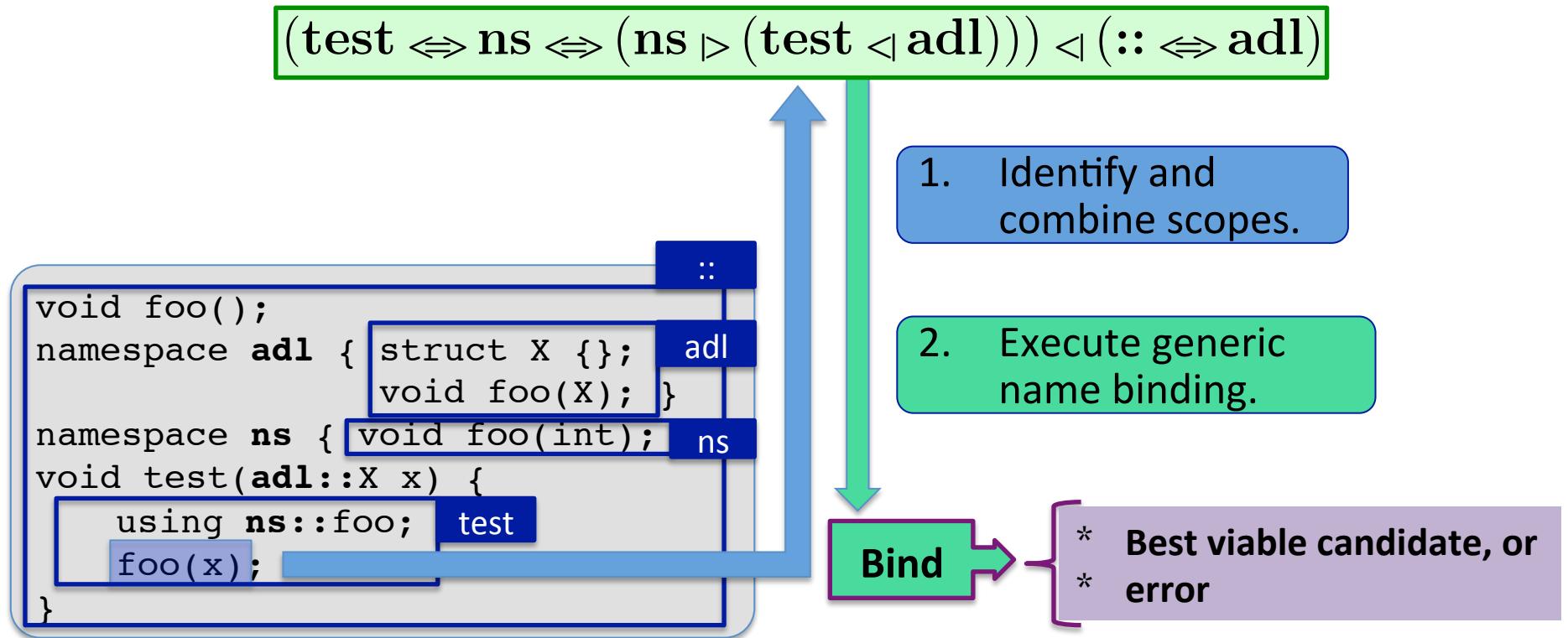


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# Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.

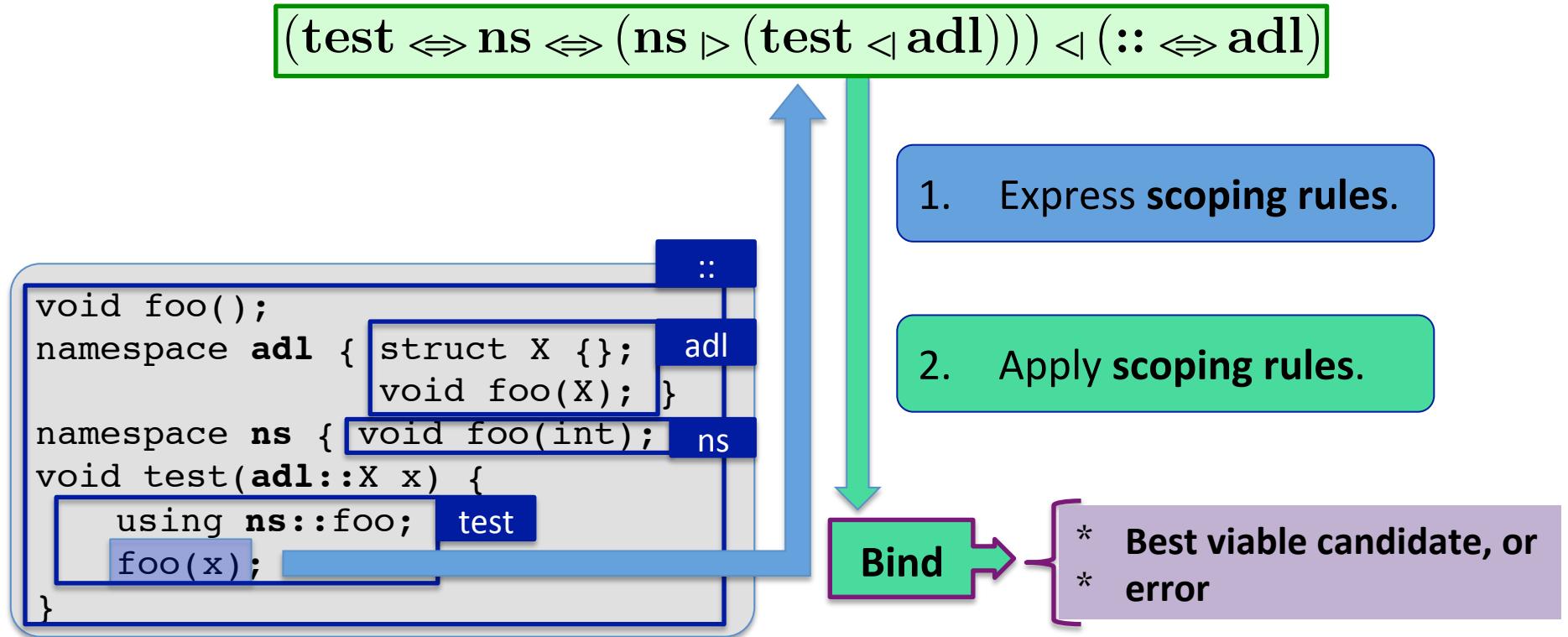


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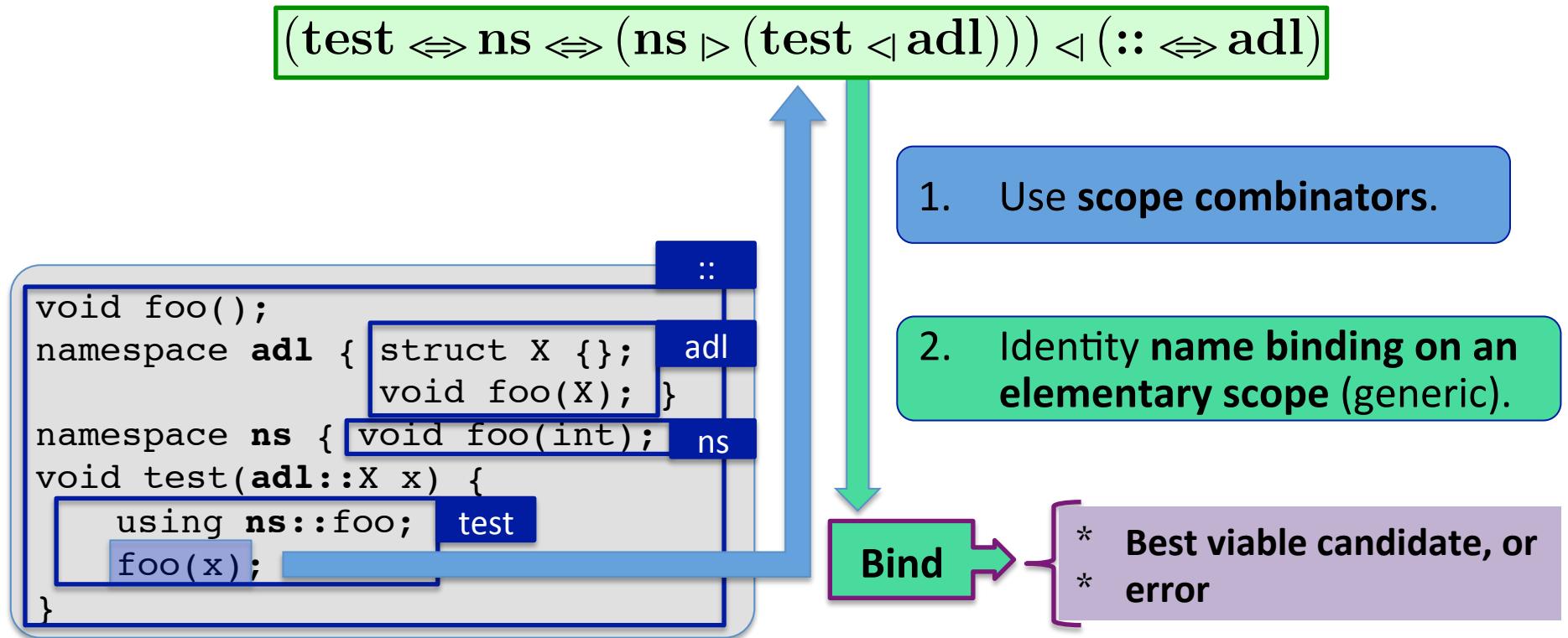


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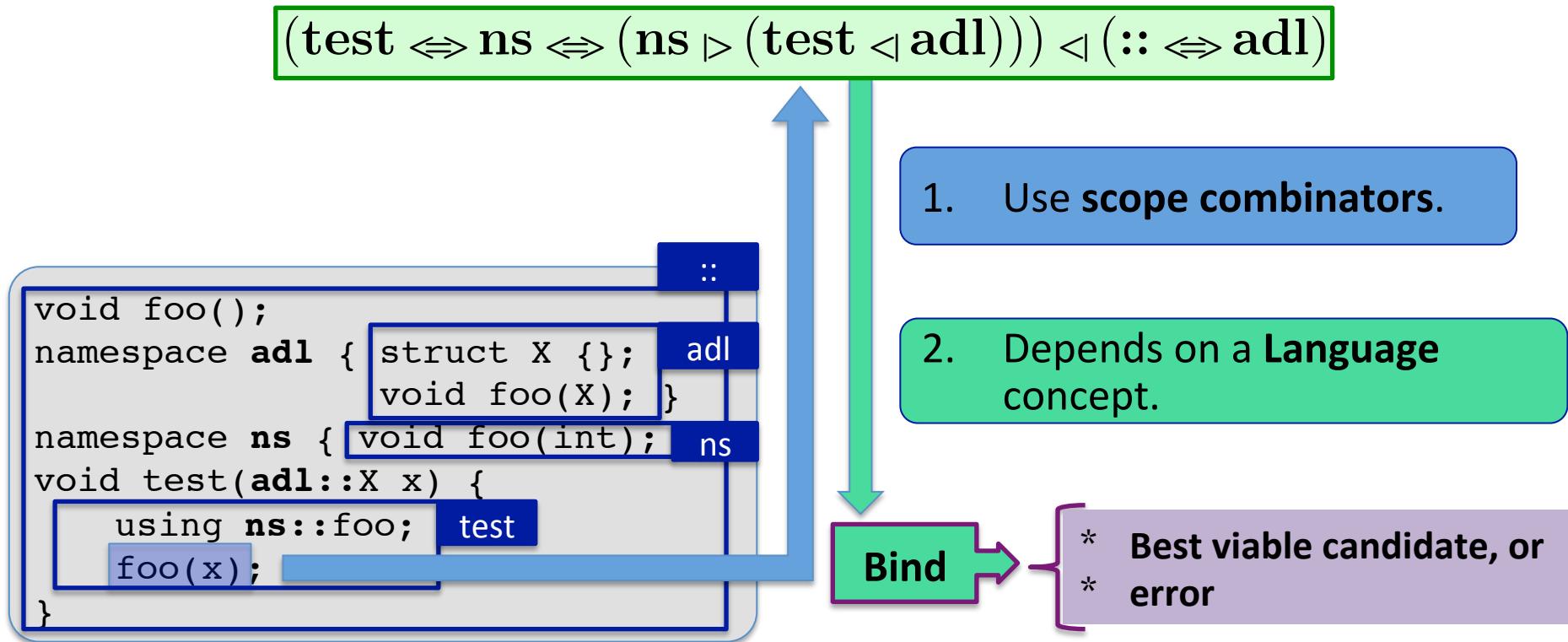


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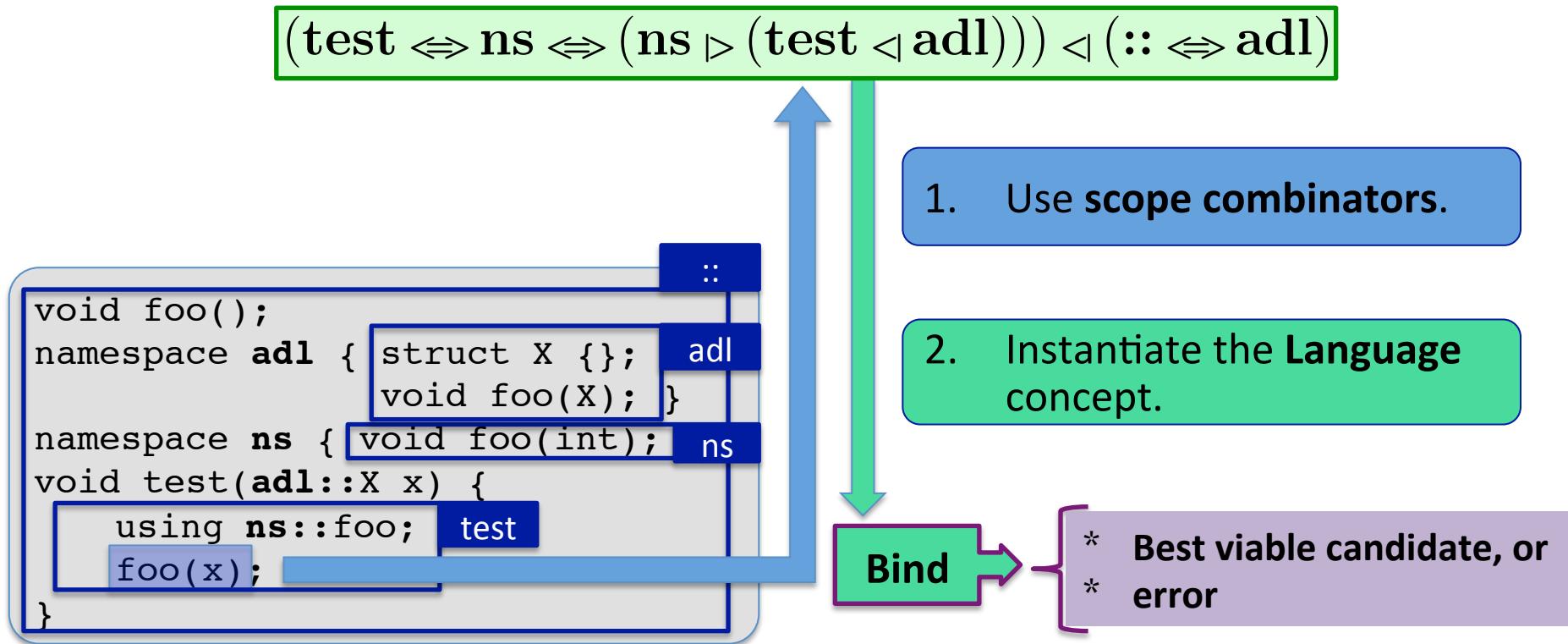


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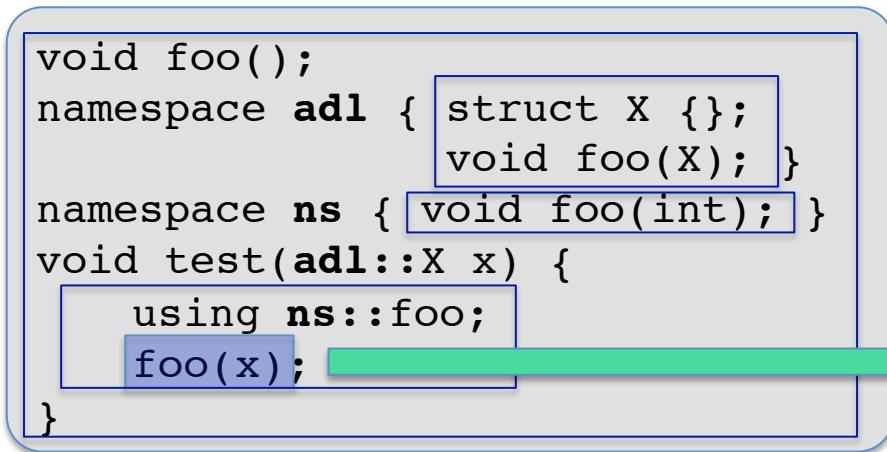


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# Our Name Binding Framework

- ❖ Abstracts from *declarations*, *references*, and *scopes*.

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$


Scopes as maps of **references** to  
**sets of matching declarations**.

Bind

- \* Best viable candidate, or
- \* error



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# Our Name Binding Framework

- ❖ Views name binding as composed of **name lookup** and **resolution**.

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$
$$bind (ref, scope) = ((\text{resolve } ref) \circ (\text{lookup } ref)) \text{ scope}$$

```
void foo();  
namespace adl { struct X {};  
                 void foo(X); }  
namespace ns { void foo(int); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x); }
```

Name lookup returns the **set of matching declarations**, for a given reference.

Bind

- \* Best viable candidate, or
- \* error



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# 1. Expressing Scoping Rules

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \triangleleft (\text{::} \Leftrightarrow \text{adl})$$

```
void foo();  
namespace adl { struct X {}; adl  
    void foo(X); }  
namespace ns { void foo(int); ns  
void test(adl::X x) {  
    using ns::foo; test  
    foo(x); }  
}
```

1. Use **scope combinators**.

2. Instantiate the **Language concept**.

Bind →

\* Best viable candidate, or  
\* error

$$\text{Scope}_{\text{Ref},\text{Decl}} \times \text{Scope}_{\text{Ref},\text{Decl}} \rightarrow \text{Scope}_{\text{Ref},\text{Decl}}$$


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# The Combinators

## ❖ *Hiding:*

<|

- Commonly known as “shadowing”.

## ❖ *Merging:*

↔

- Usually the alternative option to “shadowing”.

## ❖ *Opening:*

|>

- [New name]      Necessary to describe ADL.
- A dual of *hiding*.

## ❖ *Weak Hiding:*

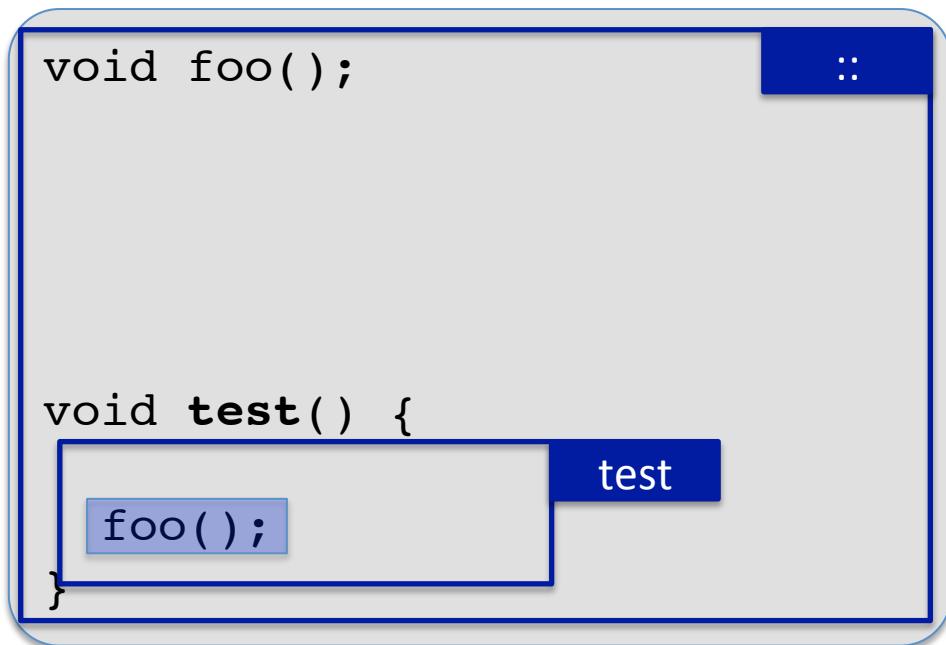
↔

- [New rule]      Necessary for (C++) concepts.
- A sweet middle between *hiding* and merging.



# The *Hiding* Combinator ( $\text{<|}$ )

$$s_1 \text{ <| } s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_1 : \text{lookup}_{ref} s_2$$



test  $\text{<| } ::$

Result:

Binds **foo()** to **::foo()**.

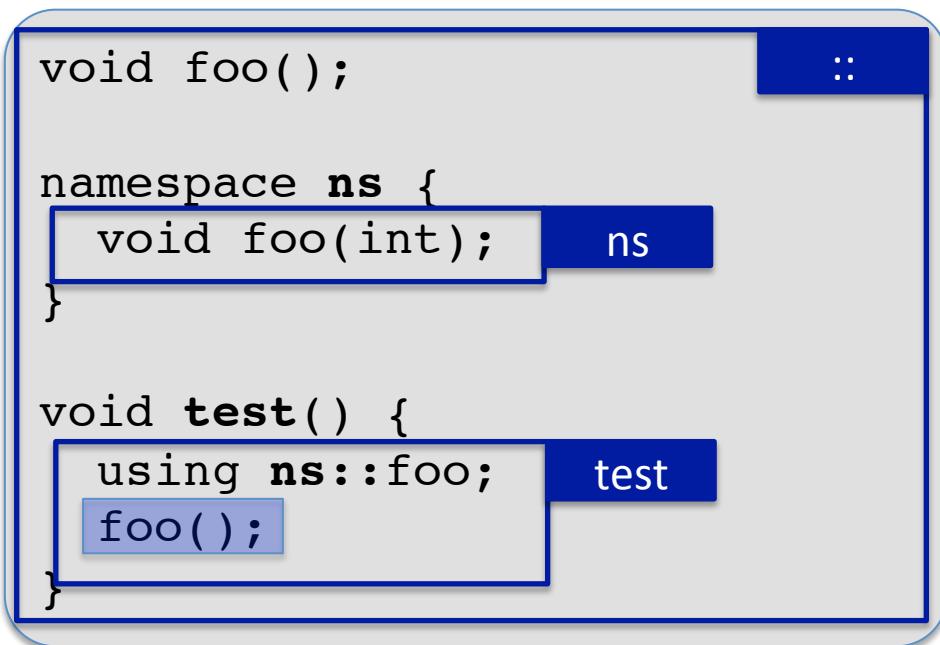


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# The *Merging Combinator* ( $\Leftrightarrow$ )

$$s_1 \Leftrightarrow s_2 = \text{lookup}_{\text{ref}} s_1 \cup \text{lookup}_{\text{ref}} s_2$$



(`test`  $\Leftrightarrow$  `ns`)  $\triangleleft ::$

Result:

Finds `ns::foo()`  
Fails to bind `foo()`.

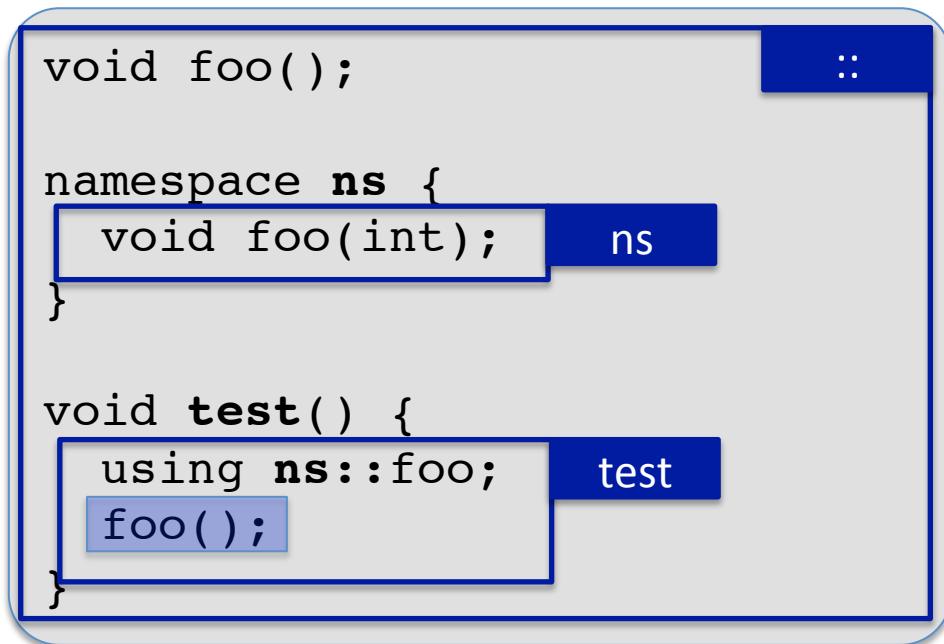


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# The *Weak Hiding* Combinator ( $\Leftarrow\Rightarrow$ )

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ? \\ lookup_{ref} s_1 : lookup_{ref} s_2$$



Result:

Binds `foo()` to `::foo()`.

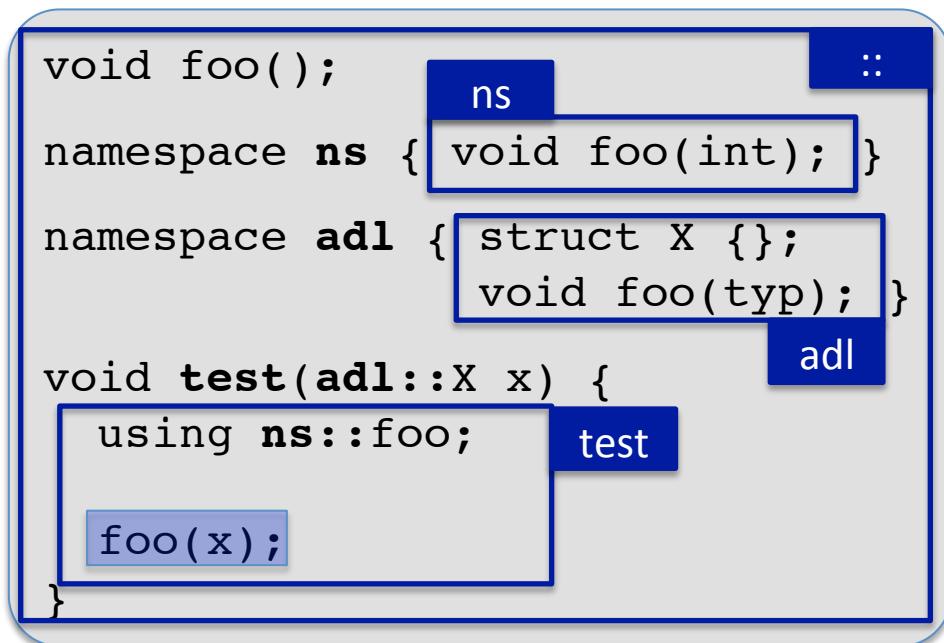


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# The *Opening* Combinator ( $\triangleright$ )

$$s_1 \triangleright s_2 = \text{lookup}_{\text{ref}} s_1 ? \text{lookup}_{\text{ref}} s_2 : \text{empty}$$



$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl})))$   
 $\triangleleft (\text{::} \Leftrightarrow \text{adl})$

Result:

Finds `ns::foo()`;

Enables ADL;

Binds `foo(x)` to `adl::foo()`.

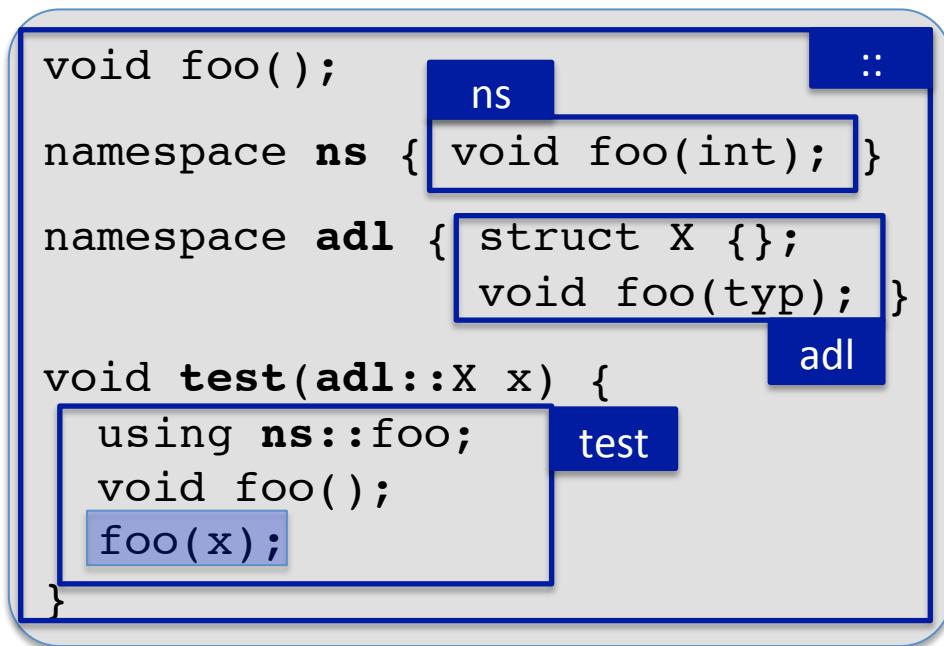


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# The *Opening* Combinator ( $\triangleright$ )

$$s_1 \triangleright s_2 = \text{lookup}_{\text{ref}} s_1 ? \text{lookup}_{\text{ref}} s_2 : \text{empty}$$



$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl})))$   
 $\triangleleft (\text{::} \Leftrightarrow \text{adl})$

Result:

Finds `test::foo()`;  
Disables ADL;  
Fails to bind `foo(x)`.



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# The Combinators – Recap

$$\langle\! , \Leftrightarrow\! , \triangleright\! , \Leftarrow\! \rangle : \text{Scope}_{\text{Ref},\text{Decl}} \times \text{Scope}_{\text{Ref},\text{Decl}} \rightarrow \text{Scope}_{\text{Ref},\text{Decl}}$$
$$s_1 \Leftrightarrow s_2 = \text{lookup}_{ref} s_1 \cup \text{lookup}_{ref} s_2$$
$$s_1 \triangleleft s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_1 : \text{lookup}_{ref} s_2$$
$$s_1 \triangleright s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_2 : \text{empty}$$
$$s_1 \Leftarrow s_2 = (\text{resolve}_{ref} \circ \text{lookup}_{ref}) s_1 ?$$
$$\text{lookup}_{ref} s_1 : \text{lookup}_{ref} s_2$$

Special cases of a “*conditional*” combinator.



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# Applications

- ❖ Understanding current name binding mechanisms:
  - Argument-dependent lookup (ADL).
  - C++ operators.
  - A cross-language analysis.
- ❖ Exploring concepts designs
  - Understanding the current limitations.
  - Exploring new solutions:
    - weak hiding,
    - 2-Stage name binding, and
    - parameterized weak hiding.
- ❖ Simplifying compiler designs.



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# ADL Example

```
namespace ns1 { struct X {};  
    ns1 void foo(X); }  
  
namespace ns2 { void foo(int); }  
  
void bar(ns1::X x) {  
    using ns2::foo; ns2  
    foo(x); bar  
}  
  
void baz(ns1::X x) {  
    void foo(int); baz  
    foo(x);  
    {  
        using ns2::foo; // H  
        foo(x); H  
    }  
}
```

In bar():

In baz():

In baz()'s inner scope, H:



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# ADL Example

```
namespace ns1 { struct X {};  
    ns1 void foo(X); }  
  
namespace ns2 { void foo(int); }  
  
void bar(ns1::X x) {  
    using ns2::foo; bar  
    foo(x); }  
  
void baz(ns1::X x) {  
    void foo(int);  
    foo(x);  
    {  
        using ns2::foo;  
        foo(x);  
    }  
}
```

In bar():

$$(\text{bar} \Leftrightarrow \text{ns2} \Leftrightarrow (\text{ns2} \triangleright (\text{bar} \triangleleft \text{ns1}))) \\ \triangleleft (:\Leftrightarrow \text{ns1})$$

In baz():

In baz()'s inner scope, H:

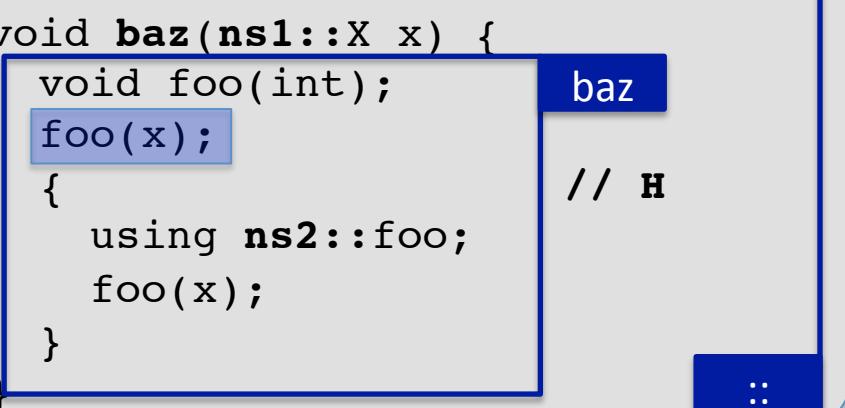


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# ADL Example

```
namespace ns1 { struct X {};  
    ns1 void foo(X); }  
  
namespace ns2 { void foo(int); }  
  
void bar(ns1::X x) {  
    using ns2::foo;  
    foo(x);  
}  
  
void baz(ns1::X x) {  
    void foo(int);  
    foo(x);  
    {  
        using ns2::foo;  
        foo(x);  
    }  
}
```



In bar():

(bar  $\Leftrightarrow$  ns2  $\Leftrightarrow$  (ns2  $\triangleright$  (bar  $\triangleleft$  ns1)))  
 $\triangleleft$  ( ::  $\Leftrightarrow$  ns1 )

In baz():

baz  $\triangleleft$  ( ::  $\Leftrightarrow$  ns1 )

In baz()'s inner scope, H:



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# ADL Example

```
namespace ns1 { struct X {};  
    ns1 void foo(X); }  
  
namespace ns2 { void foo(int); }  
  
void bar(ns1::X x) { ns2  
    using ns2::foo;  
    foo(x);  
}  
  
void baz(ns1::X x) {  
    void foo(int);  
    foo(x);  
    {  
        using ns2::foo;  
        foo(x);  
    }  
}
```

In bar():

$$(\text{bar} \Leftrightarrow \text{ns2} \Leftrightarrow (\text{ns2} \triangleright (\text{bar} \triangleleft \text{ns1}))) \\ \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$

In baz():

$$\text{baz} \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$

In baz()'s inner scope, H:

$$(\text{H} \Leftrightarrow \text{ns2} \Leftrightarrow (\text{ns2} \triangleright (\text{H} \triangleleft \text{ns1}))) \\ \triangleleft \text{baz} \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$


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# ADL Example

```

namespace ns1 { struct X {};
    void foo(X); }
namespace ns2 { void foo(int); }

void bar(ns1::X x) {
    using ns2::foo;
    foo(x);
}

void baz(ns1::X x) {
    void foo(int);
    foo(x);
    {
        using ns2::foo;
        foo(x);
    }
}

```

The diagram illustrates the scope resolution of identifiers in the provided C++ code. The code is contained within a blue rounded rectangle. Identifiers are highlighted in blue boxes. Labels are placed near the identifiers to indicate their scope or context:

- ns1**: A blue box containing the identifier `ns1`.
- ns2**: A blue box containing the identifier `ns2`.
- bar**: A blue box containing the identifier `bar`.
- baz**: A blue box containing the identifier `baz`.
- H**: A blue box containing the identifier `H`, which is a label for the innermost scope of the `baz` function.
- ::**: A blue box containing the double colon operator `::`.

In `bar()`:

$$f(\text{bar}) \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$

In `baz()`:

$$f(\text{baz}) \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$

In `baz()`'s inner scope, H:

$$f(H) \triangleleft f(\text{baz}) \triangleleft (\text{::} \Leftrightarrow \text{ns1})$$

----- with -----

$$f(X) = X \Leftrightarrow U_X \Leftrightarrow (U_X \triangleright (X \triangleleft \text{ns1}))$$

$$U_{\text{bar}} = U_H = \text{ns2}$$

$$U_{\text{baz}} = \text{empty}$$



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# ADL Scoping Rules

$$f_{ADL}(H) \triangleleft \langle\!\langle \underset{i=1}{\overset{s}{\cup}} f_{ADL}(S_i) \rangle\!\rangle \triangleleft \langle\!\langle \underset{i=1}{\overset{n-1}{\cup}} fn_{ADL}(N_i) \rangle\!\rangle \triangleleft \left( N_n \Leftrightarrow U_{N_n} \Leftrightarrow \left( \left( \tilde{N}_n \Leftrightarrow \tilde{U}_{N_n} \right) \triangleleft ADL \right) \right)$$

$$f_{ADL}(X) = X \Leftrightarrow U_X \Leftrightarrow \left( U_X \triangleright \left( \left( X \Leftrightarrow \tilde{U}_X \right) \triangleleft ADL \right) \right),$$

$S_1 \cdots S_s$  = surrounding non-namespace scopes,

$$fn_{ADL}(N) = N \Leftrightarrow U_N \Leftrightarrow \left( \left( N \Leftrightarrow U_N \right) \triangleright \left( \left( \tilde{N} \Leftrightarrow \tilde{U}_N \right) \triangleleft ADL \right) \right),$$

$N_1 \cdots N_n$  = surrounding namespaces,

$H$  = scope where name binding is triggered from,

$U_X$  = using declarations in scope  $X$ ,

$\tilde{X}$  = non-function (template) declarations in scope  $X$ , and

$ADL$  = associated namespaces of reference's arguments.



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# ADL Scoping Rules

when scope **H** is an inner namespace scope

$$\text{fn}_{\text{ADL}}(\mathbf{H}) \triangleleft \langle \!\! \langle \bigcup_{i=1}^{n-1} \text{fn}_{\text{ADL}}(\mathbf{N}_i) \triangleleft \left( \mathbf{N}_n \Leftrightarrow \mathbf{U}_{\mathbf{N}_n} \Leftrightarrow \left( \left( \tilde{\mathbf{N}}_n \Leftrightarrow \tilde{\mathbf{U}}_{\mathbf{N}_n} \right) \triangleleft \text{ADL} \right) \right) \rangle \!\! \rangle$$

$$\text{fn}_{\text{ADL}}(\mathbf{N}) = \mathbf{N} \Leftrightarrow \mathbf{U}_{\mathbf{N}} \Leftrightarrow \left( \left( \mathbf{N} \Leftrightarrow \mathbf{U}_{\mathbf{N}} \right) \triangleright \left( \left( \tilde{\mathbf{N}} \Leftrightarrow \tilde{\mathbf{U}}_{\mathbf{N}} \right) \triangleleft \text{ADL} \right) \right),$$

$\mathbf{N}_1 \cdots \mathbf{N}_n$  = surrounding namespaces,

$\mathbf{H}$  = scope where name binding is triggered from,

$\mathbf{U}_X$  = using declarations in scope  $X$ ,

$\tilde{\mathbf{X}}$  = non-function (template) declarations in scope  $X$ , and

**ADL** = associated namespaces of reference's arguments.



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# ADL Scoping Rules

when scope **H** is the outermost namespace scope

$$(N_n \Leftrightarrow U_{N_n} \Leftrightarrow ((\tilde{N}_n \Leftrightarrow \tilde{U}_{N_n}) \triangleleft ADL))$$

**H = N<sub>n</sub>**

**H** = scope where name binding is triggered from,  
**U<sub>X</sub>** = using declarations in scope **X**,  
**˜X** = non-function (template) declarations in scope **X**, and  
**ADL** = associated namespaces of reference's arguments.



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# C++ Operators Example

```
struct X {}; struct Y {};  
  
void operator+(X, X) { }  
void operator+(X, Y) { }  
  
void test(X x, Y y) {  
    void operator+(X, X); test  
    x + x;  
    x + y;  
    operator+(x, x);  
    operator+(x, y);  
}
```

(test  $\Leftrightarrow$  builtin-ops  $\Leftrightarrow$  X)  $\triangleleft$  ::



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# C++ Operators Scoping Rules

## without ADL

$$\left( H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \right)$$
$$\triangleleft \langle \rangle_{i=1}^s S_i \triangleleft \langle \rangle_{i=1}^n N_i$$

$S_1 \dots S_s$  = surrounding non-namespace scopes,

$N_1 \dots N_n$  = surrounding namespaces,

$H$  = scope where name binding is triggered from,

$O_B$  = built-in operators,

$M$  = member scope of operator's first argument,

$U_X$  = using declarations in scope  $X$ , and



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# C++ Operators Scoping Rules with ADL

$$\begin{aligned} & \left( H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow \left( U_H \triangleright \left( \left( H \Leftrightarrow \tilde{U}_H \Leftrightarrow O_B \Leftrightarrow M \right) \triangleleft ADL \right) \right) \right) \\ & \triangleleft \langle \langle_{i=1}^s f_{ADL}(S_i) \triangleleft \langle \langle_{i=1}^{n-1} fn_{ADL}(N_i) \triangleleft \left( N_n \Leftrightarrow U_{N_n} \Leftrightarrow \left( \left( \tilde{N}_n \Leftrightarrow \tilde{U}_{N_n} \right) \triangleleft ADL \right) \right) \end{aligned}$$

$$f_{ADL}(X) = X \Leftrightarrow U_X \Leftrightarrow \left( U_X \triangleright \left( \left( X \Leftrightarrow \tilde{U}_X \right) \triangleleft ADL \right) \right),$$

$S_1 \dots S_s$  = surrounding non-namespace scopes,

$$fn_{ADL}(N) = N \Leftrightarrow U_N \Leftrightarrow \left( \left( N \Leftrightarrow U_N \right) \triangleright \left( \left( \tilde{N} \Leftrightarrow \tilde{U}_N \right) \triangleleft ADL \right) \right),$$

$N_1 \dots N_n$  = surrounding namespaces,

$H$  = scope where name binding is triggered from,

$O_B$  = built-in operators,

$M$  = member scope of operator's first argument,

$U_X$  = using declarations in scope  $X$ ,

$\tilde{X}$  = non-function (template) declarations in scope  $X$ , and

$ADL$  = associated namespaces of reference's arguments.

Empty, since  
**operator** is a  
reserved keyword.



# C++ Operators Scoping Rules with ADL

$$\left( \mathbf{H} \Leftrightarrow \mathbf{U}_\mathbf{H} \Leftrightarrow \mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \Leftrightarrow \left( \mathbf{U}_\mathbf{H} \triangleright \left( \left( \mathbf{H} \Leftrightarrow \mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \right) \triangleleft \mathbf{ADL} \right) \right) \right) \\ \triangleleft \langle \!\! \begin{array}{c} s \\ i=1 \end{array} \!\! \rangle \mathbf{f}_{\mathbf{ADL}} (\mathbf{S}_i) \triangleleft \langle \!\! \begin{array}{c} n-1 \\ i=1 \end{array} \!\! \rangle \mathbf{fn}_{\mathbf{ADL}} (\mathbf{N}_i) \triangleleft (\mathbf{N}_n \Leftrightarrow \mathbf{U}_{\mathbf{N}_n} \Leftrightarrow \mathbf{ADL})$$

$\mathbf{f}_{\mathbf{ADL}} (\mathbf{X}) = \mathbf{X} \Leftrightarrow \mathbf{U}_\mathbf{X} \Leftrightarrow (\mathbf{U}_\mathbf{X} \triangleright (\mathbf{X} \triangleleft \mathbf{ADL}))$ ,

$\mathbf{S}_1 \cdots \mathbf{S}_s$  = surrounding non-namespace scopes,

$\mathbf{fn}_{\mathbf{ADL}} (\mathbf{N}) = \mathbf{N} \Leftrightarrow \mathbf{U}_\mathbf{N} \Leftrightarrow \left( (\mathbf{N} \Leftrightarrow \mathbf{U}_\mathbf{N}) \triangleright \mathbf{ADL} \right)$ ,

$\mathbf{N}_1 \cdots \mathbf{N}_n$  = surrounding namespaces,

$\mathbf{H}$  = scope where name binding is triggered from,

$\mathbf{O}_\mathbf{B}$  = built-in operators,

$\mathbf{M}$  = member scope of operator's first argument,

$\mathbf{U}_\mathbf{X}$  = using declarations in scope  $\mathbf{X}$ , and

$\mathbf{ADL}$  = associated namespaces of reference's arguments.



# C++ Operators Scoping Rules with ADL

$$\begin{aligned} & \left( H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow \left( U_H \triangleright \left( \left( H \Leftrightarrow O_B \Leftrightarrow M \right) \triangleleft ADL \right) \right) \right) \\ & \triangleleft \langle\!\langle_{i=1}^s f_{ADL}(S_i) \triangleleft \left( \langle\!\langle_{i=1}^n (N_i \Leftrightarrow U_{N_i}) \Leftrightarrow ADL \right) \right. \end{aligned}$$

$f_{ADL}(X) = X \Leftrightarrow U_X \Leftrightarrow (U_X \triangleright (X \triangleleft ADL))$ ,

$S_1 \cdots S_s$  = surrounding non-namespace scopes,

$N_1 \cdots N_n$  = surrounding namespaces,

$H$  = scope where name binding is triggered from,

$O_B$  = built-in operators,

$M$  = member scope of operator's first argument,

$U_X$  = using declarations in scope  $X$ , and

$ADL$  = associated namespaces of reference's arguments.



# C++ Operators Scoping Rules

when scope **H** is an inner namespace scope

$$\begin{aligned} & \left( H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow \left( \left( H \Leftrightarrow U_H \right) \triangleright \left( \left( O_B \Leftrightarrow M \right) \triangleleft ADL \right) \right) \right) \\ & \quad \triangleleft \left( \bigwedge_{i=1}^n (N_i \Leftrightarrow U_{N_i}) \Leftrightarrow ADL \right) \end{aligned}$$

**N<sub>1</sub>** … **N<sub>n</sub>** = surrounding namespaces,

**H** = scope where name binding is triggered from,

**O<sub>B</sub>** = built-in operators,

**M** = member scope of operator's first argument,

**U<sub>X</sub>** = using declarations in scope **X**, and

**ADL** = associated namespaces of reference's arguments.



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# C++ Operators Scoping Rules

when scope **H** is the outermost namespace scope

$$(H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow ( ((O_B \Leftrightarrow M) \triangleleft ADL) ))$$

**H** = scope where name binding is triggered from,

**O<sub>B</sub>** = built-in operators,

**M** = member scope of operator's first argument,

**U<sub>X</sub>** = using declarations in scope **X**, and

**ADL** = associated namespaces of reference's arguments.



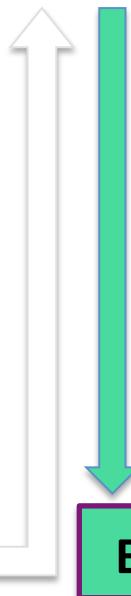
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## 2. Executing Scoping Rules

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \triangleleft (\text{::} \Leftrightarrow \text{adl})$$

```
void foo();  
namespace adl { struct X {};  
                 void foo(X); }  
namespace ns { void foo(int); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x); }
```



1. Use scope combinators.

2. Instantiate the **Language** concept.

Bind

- \* Best viable candidate, or
- \* error

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$


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# Elementary Name Binding

Scopes as sets  
of declarations.

```
lookup :: (...) => ref -> Set decl -> OverloadSet decl
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (...) => ref -> OverloadSet decl -> Maybe decl
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (...) => BestViableSet decl -> Maybe decl
assess decls = case (Set.elems decls) of
    []      -> Nothing
    [decl]  -> (Just decl)
    _        -> ambiguity decls
```



# Elementary Name Binding

## ❖ What is the necessary minimal abstraction?

```
lookup :: (...) => ref -> Set decl -> OverloadSet decl  
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (...) => ref -> OverloadSet decl -> Maybe decl  
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (...) => BestViableSet decl -> Maybe decl  
assess decls = case (Set.elems decls) of  
    [ ]      -> Nothing  
    [decl]   -> (Just decl)  
    _         -> ambiguity decls
```



# An Abstraction of PLs

## ❖ The **Language** concept:

- **match**: How references match declarations.
- **select\_best\_viable**: How best viable candidates are determined.
- **ambiguity**: How to handle ambiguous results.



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# An Abstraction of PLs

## ❖ The **Ambiguity** concept:

- ambiguity: How to handle ambiguous results.

## ❖ The **Language** concept:

- match: How references match declarations.
- select\_best\_viable: How best viable candidates are determined.
- refines **Ambiguity**.



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# An Abstraction of PLs, Ext'd

## ❖ The **Ambiguity** concept:

- ambiguity: How to handle ambiguous results.

## ❖ The **Language** concept:

- match: How references match declarations.
- `select_best_viable`: How best viable candidates are determined.
- refines **Ambiguity**.

## ❖ The **Parameterized** concept:

- expresses changes in the bind environment during name binding.
- e.g., **changes in ambiguity** for a variant of *weak hiding*.



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# An Abstraction of PLs

```
class Ambiguity decl where
    ambiguity :: ViableSet decl -> Maybe decl
    ambiguity _ = Nothing                                -- default.
```

```
class Ambiguity decl => Language ref decl where
    match :: ref -> decl -> bool
    select_best_viable :: ref -> OverloadSet decl -> BestViableSet decl
```



# An Abstraction of PLs, Ext'd

```
class Ambiguity decl where
  ambiguity :: ViableSet decl -> Maybe decl
  ambiguity _ = Nothing                                -- default.
```

```
class Ambiguity decl => Basic.Language ref decl where
  match :: ref -> decl -> bool
  select_best_viable :: ref -> OverloadSet decl -> BestViableSet decl
```

```
class (Parameterized decl, Basic.Language ref decl) =>
  Language ref decl
```



# Elementary Name Binding

```
lookup :: (Language ref decl) => ref -> Set decl -> OverloadSet decl
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (Language ref decl) => ref -> OverloadSet decl -> Maybe decl
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (Ambiguity decl) => BestViableSet decl -> Maybe decl
assess decls = case (Set.elems decls) of
    []      -> Nothing
    [decl]  -> (Just decl)
    _       -> ambiguity decls
```

- ❖ **assess** requires the **Ambiguity** concept only,  
gets the meaning of ambiguity from **the language**.



# Elementary Name Binding, Ext'd

```
resolve :: (Basic.Language ref decl) =>
    BindEnv decl -> ref -> OverloadSet decl -> Maybe decl
resolve env ref decls = assess env $ select_best_viable ref decls
```

```
assess :: BindEnv decl -> BestViableSet decl -> Maybe decl
assess env decls = case (Set.elems decls) of
    []      -> Nothing
    [decl] -> (Just decl)
    _       -> ambiguity env decls
```

- ❖ **assess** requires **no** concept,  
gets the meaning of ambiguity from **the bind environment**.



# Applications

- ❖ Understanding current name binding mechanisms:
  - Argument-dependent lookup.
  - C++ operators.
  - A cross-language analysis.
- ❖ Exploring concepts designs
  - Understanding the current limitations.
  - Exploring new solutions:
    - weak hiding,
    - 2-Stage name binding, and
    - parameterized weak hiding.
- ❖ Simplifying compiler designs.



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# A cross language analysis

- ❖ Instantiate the **Language** concept with different languages, designs and kinds of references:
  - C++: function calls
    - **match** = have\_same\_name + filter out types.
    - **select\_best\_viable** = overload resolution (incl. viability check).
    - **ambiguity** = ambiguity\_is\_error.
  - C++: uses of types
    - **match** = have\_same\_name.
    - **select\_best\_viable** = viability check if single result, identity otherwise.
    - **ambiguity** = ambiguity\_is\_error.
  - C++: Multimethods proposal (cf. N2216)
    - **match** = (same as function calls).
    - **select\_best\_viable** = (same as function calls).
    - **ambiguity** = ambiguity\_is\_not\_error.



# C++ Multimethods (Std. Doc. N2216)

titled “Open Multi-Methods for C++”, Pirkelbauer et al., 2007

- ❖ Relaxes the rules for ambiguity,
  - allowing ambiguous calls when all best viable candidates “**have a unique-base method** through which the call can be dispatched”.

```
struct X, Y, Z;
void foo(virtual X&, virtual Y&); // (1)
void foo(virtual Y&, virtual Y&); // (2) - ** unique-base method ***
void foo(virtual Y&, virtual Z&); // (3)
struct XY : X, Y {}
struct YZ : Y, Z {}
void foo(virtual XY&, virtual Y&); // (4) - overrider for (1) and (2)
void foo(virtual Y&, virtual YZ&); // (5) - overrider for (2) and (3)

XY xy; YZ yz;
foo(xy,yz); // both (4) and (5) are equally viable matches,
// w/ unique base (2)
```



# A cross language analysis

- ❖ Instantiate the **Language** concept with different languages, designs and kinds of references:
  - Haskell: name uses, excluding type inference
    - **match** = have\_same\_name.
    - **select\_best\_viable** = identity.
    - **ambiguity** = ambiguity\_is\_error.
  - Haskell: name uses, including type inference
    - **match** = have\_same\_name.
    - **select\_best\_viable** = viability check if single result, identity otherwise.
    - **ambiguity** = ambiguity\_is\_error.
  - Alternatives have different implications with weak hiding or **Bind**<sup>x2</sup>.



# Applications

- ❖ Understanding current name binding mechanisms:
  - Argument-dependent lookup.
  - C++ Operators.
  - A cross-language analysis.
- ❖ Exploring concepts designs
  - Understanding the current limitations.
  - Exploring new solutions:
    - weak hiding,
    - 2-Stage name binding, and
    - parameterized weak hiding.
- ❖ Simplifying compiler designs

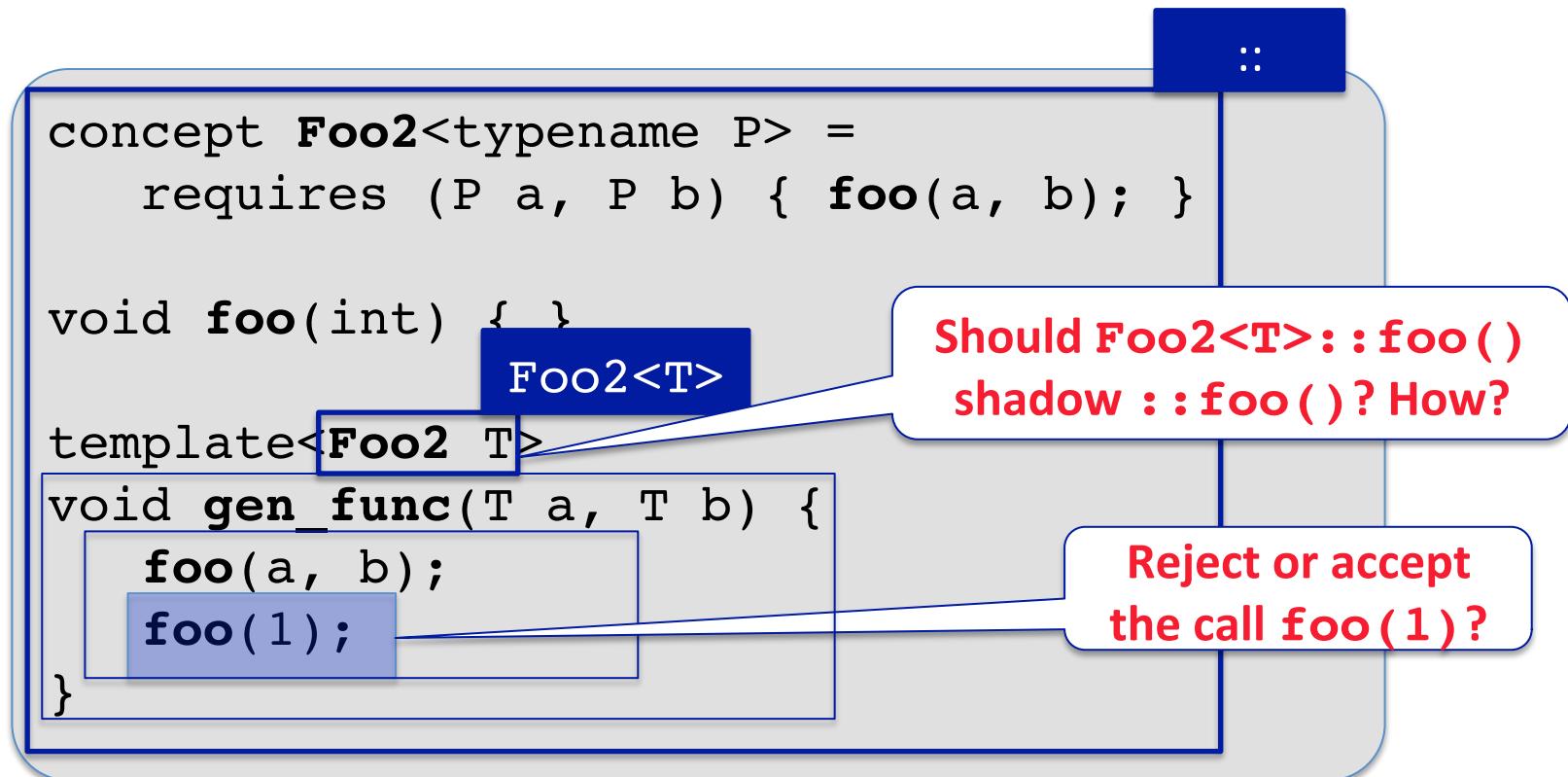


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# Problem: Current Limitations

- ❖ Current scoping rules break seemingly valid codes.



# Practical Examples

- ❖ STL: `rotate()` and `move()`
  - Two function declarations, different number of parameters.
  - Name lookup only finds one, type-check fails.
- ❖ Plenoptic photography: Image rendering
  - Two function declarations, different parameter types.
  - Name lookup only finds one, type-check fails.
- ❖ STL: `common_type`
  - Two type (function) declarations, different type parameters.
  - Name lookup only finds one, type-check fails.



# STL: rotate( ) and move( )

```
// Specialization adapted from latest release of libstdc++
template<RandomAccessIterator I>
    requires Permutable<I>           // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = std::move(*p);                      //
        std::move(p+1, p+last-first, p);                      //
        *(p + last - first -1) = std::move(t);   //
    }                                         // Ok w/ constrained std::move(), But...
    ...                                         // Does not use the constraints.
}
```



# STL: rotate( ) and move( )

```
// A novice's approach

template<RandomAccessIterator I>
    requires Permutable<I>           // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {

    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = move(*p);                      // Ok
        move(p+1, p+last-first, p);                      // ...
        *(p + last - first -1) = move(t);               // Ok
    }          // Compile error... or... Possible ADL issues...
    ...
}
```



# STL: rotate( ) and move( )

```
// The right way...

template<RandomAccessIterator I>
    requires Permutable<I> // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {

    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = move(*p); // Ok
        std::move(p+1, p+last-first, p); // Ok
        *(p + last - first -1) = move(t); // Ok
    } // Shadow or No Shadow... But...
    ... // Requires change in existing implementation.
}
```



# STL: rotate( ) and move( )

```
// Update MoveWritable concept ...
template<RandomAccessIterator I>
    requires Permutable<I> // has std::move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>)) && (middle - first == 1)) {
        ValueType<I> t = std::move(*p); // Ok
        std::move(p+1, p+last-first, p); // 
        *(p + last - first -1) = std::move(t); // Ok
    } // Compile error... or... No ADL issues...
    ... // Maybe OK to not shadow?
} // std::move() is not a customization point.
```



# Plenoptic photo.: Image rendering

```
template<typename PixelType>
struct Radiance {
    typedef typename
        boost::multi_array<PixelType, 4> RadianceType;
    RadianceType pixels;      // 4D array of pixels
    ...
    void Read(const string &ImageFile, ...);
    Radiance<PixelType> Render_Basic(...);
    Radiance<PixelType> Render_Blended(...);
    ... // Several rendering variants ...
    void Print(const string& OutputFile, ...);
    ...
}
```



# Plenoptic photo.: Image rendering

```
template<typename PixelType>
Radiance<PixelType> Radiance<PixelType>::Render_Blended(...) {
    ...
    RadianceType Rendered(boost::extents[Iy][Ix][1][1]); ...
    for (...) { ... // for each image pixel
        PixelType pixel_avg; ...
        for (...) { ... // for each direction
            pixel_avg += pixels[ri][rj][rl][rm]; ... // integrate pixel
        } ...
        Rendered[i][j][0][0] = move(pixel_avg); ...
    }
    return move(new Radiance<PixelType>(move(Rendered), Ix, Iy));
}
```



# Plenoptic photo.: Image rendering

```
template<typename PixelType>
using MultiArrayIterator =
    boost::multi_array<PixelType, 4>::iterator

// Add safety
template<Regular PixelType>
    requires IndirectlyCopyable<MultiArrayIterator<PixelType>,
                           MultiArrayIterator<PixelType>>

struct Radiance {
    typedef typename
        boost::multi_array<PixelType, 4> RadianceType;
    ...
}
```



# Plenoptic photo.: Image rendering

```
template<Regular PixelType> ...
Radiance<PixelType>&& Radiance<PixelType>::Render_Blended(...) {
    ...
    RadianceType Rendered(boost::extents[Iy][Ix][1][1]); ...
    for (...) { ... // for each image pixel
        PixelType pixel_avg; ...
        for (...) { ... // for each direction
            pixel_avg += pixels[ri][rj][rl][rm]; ... // integrate pixel
        } ...
        Rendered[i][j][0][0] = move(pixel_avg); ... // ok.
    } // Compile error.
    return move(Radiance<PixelType>(move(Rendered), Ix, Iy));
}
```



# STL: common\_type

```
concept Common<typename T, typename U> =
    requires { common_type<T, U>::type;      axiom (...) { ... } };

template<typename T, typename U>
requires Common<T, U> && ...
void gen_func(...) {
    common_type<T, U>::type ...           // Ok.
    common_type<int>::type ...           // Error.
    common_type<int, double>::type ... // Error.
    common_type<int, double, char>::type ... // Error.
}
```



# Solution: Weak Hiding

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); }  
  
void foo(int) {}  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

**Foo2<T>::foo()**  
should weakly hide  
::foo()!

Accept the call  
foo(1)!



# The Weak Hiding Scoping Rule

$$s_1 \Leftrightarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$
$$lookup_{ref} s_1 : lookup_{ref} s_2$$

```
concept Foo2<typename P> =
    requires (P a, P b) { ... }

void foo(int) { }

template<Foo2 T>
void gen_func(T a, T b) {
    foo(a, b);
    foo(1);
}
```

gen\_func  $\triangleleft$  P  
 $\triangleleft$  (Concept<T>  $\Leftrightarrow$  T)  $\Leftrightarrow ::$

Result:

Binds **foo()** to  $::::\text{foo}()$ .



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# Implementing Weak Hiding

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$
$$lookup_{ref} s_1 : lookup_{ref} s_2$$
$$bind (ref, scope) = ((resolve\ ref) \circ (lookup\ ref))\ scope$$

❖ Implementation = **Two-Stage Name Binding (Bind<sup>x2</sup>)**

1. Bind with inner scope:      **s1.**
2. Bind with outer scope:      **s2.**

❖ Bind<sup>x2</sup> repeats name binding under different contexts.



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# Bind<sup>x2</sup> for C++ Concepts

## 1. Within restricted scope:

- up to the outermost restricted scope.
- Disables ADL and some qualified name lookups.

## 2. In surrounding scope:

- normal lookup – including ADL.

```
concept Foo2<typename P> =
    requires (P a, P b) { ... }

void foo(int) { }

template<Foo2 T>
void gen_func(T a, T b) {
    foo(a, b);
    foo(1);
}
```



# Applications

- ❖ Understanding current name binding mechanisms:
  - Argument-dependent lookup.
  - C++ Operators.
  - A cross-language analysis.
- ❖ Exploring concepts designs
  - Understanding the current limitations.
  - Exploring new solutions:
    - weak hiding,
    - 2-Stage name binding, and
    - parameterized weak hiding.
- ❖ Simplifying compiler designs



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# Inner Scope Ambiguity

when ambiguity IS an error

Current C++ Concepts:  
**Reject!**

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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# Inner Scope Ambiguity

when ambiguity IS an error

Proposed Extension:  
**Accept!**

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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# Inner Scope Ambiguity

when ambiguity IS an error

False

Repeat bind

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$
$$lookup_{ref} s_1 : lookup_{ref} s_2$$
$$bind (ref, scope) = ((resolve\ ref) \circ (lookup\ ref))\ scope$$


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# Inner Scope Ambiguity

when ambiguity IS an error

False

$s_1 \leqslant s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$

Repeat bind

$lookup_{ref} s_1 : lookup_{ref} s_2$

Not Current C++ Concepts:  
Reject or bind to unintended!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}
```

```
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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# Inner Scope Ambiguity

when ambiguity IS an error

True

temporary change in ambiguity

Do not repeat bind

$$s_1 \Leftrightarrow s_2 = (resolve_{ref} \circ (update \circ lookup_{ref})) \quad ?$$

$lookup_{ref} s_1 : lookup_{ref} s_2$

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}
```

```
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```

Current C++ Concepts:  
**Reject!**



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# Inner Scope Ambiguity

when ambiguity IS NOT an error

True

Do not repeat bind

$$S_1 \Leftarrow S_2 = (resolve_{ref} \circ lookup_{ref}) S_1 ?$$

$lookup_{ref} S_1 : lookup_{ref} S_2$

Not Current C++ Concepts:  
Accept!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}
```

```
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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# Inner Scope Ambiguity

when ambiguity IS NOT an error

False

$s_1 \leqslant s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$

Repeat bind

$lookup_{ref} s_1 : lookup_{ref} s_2$

```
concept Foo2<typename P> {
    void foo(P, int);
    void foo(int, P);
}
```

```
template<Foo2 T>
void gen_func(T a, int b) {
    foo(b);
}
```

Not Current C++ Concepts:  
Accept unintended ambiguity!



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# Ambiguity for C++ Concepts:

What is the most desirable?

1. Ambiguity IS an error, always?

Rejects desirable, or binds to undesirable

2. Ambiguity IS NOT an error, always?

Accepts undesirable

3. A middle ground option?

Proposed Extension:  
Accept only desirable! (?)

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}
```

```
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b); foo(b);  
}
```



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# Ambiguity for C++ Concepts

Our proposed extension

- ❖ Ambiguity IS NOT an error, when in restricted scope.
  - Similar to multimethods proposal N2216.
- ❖ Ambiguity IS an error, otherwise.

Proposed Extension:  
Accept only desirable!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b); foo(b);  
}
```



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# Revisiting Hinnant's Example

- ❖ Cf. C++ standard documents:
  - C++ Library Reflector message c++std-lib-20050
  - N2576: “*Type-Soundness and Optimization in the Concepts Proposal*”, Douglas Gregor, March 2008.
- ❖ **Question: How to type-check constrained templates?**

```
template <class T>
    requires std::DefaultConstructible<T> && std::Addable<T> &&
    std::Convertible<T::result_type, T> && std::CopyConstructible<T>
void test() {
    T s1;
    T s2;
    T s3 = s1 + s2 + T() + T() + T();
}

int main() { test<string>(); }
```



# Hinnant's Example: Problem

## ❖ The desired (optimized) output

- is observed with unconstrained templates, but
- is not observed with constrained templates.

```
auto concept Addable<typename T, typename U=T> {
    typename result_type;
    result_type operator+(T const&, U const&);
}

template <class T>
    requires ... std::Addable<T> ...
void test() {
    T s1;
    T s2;
    T s3 = s1 + s2 + T() + T() + T();
}

int main() { test<string>(); }
```

### Unconstrained:

lv string + lv string  
rv string += rv string  
rv string += rv string  
rv string += rv string

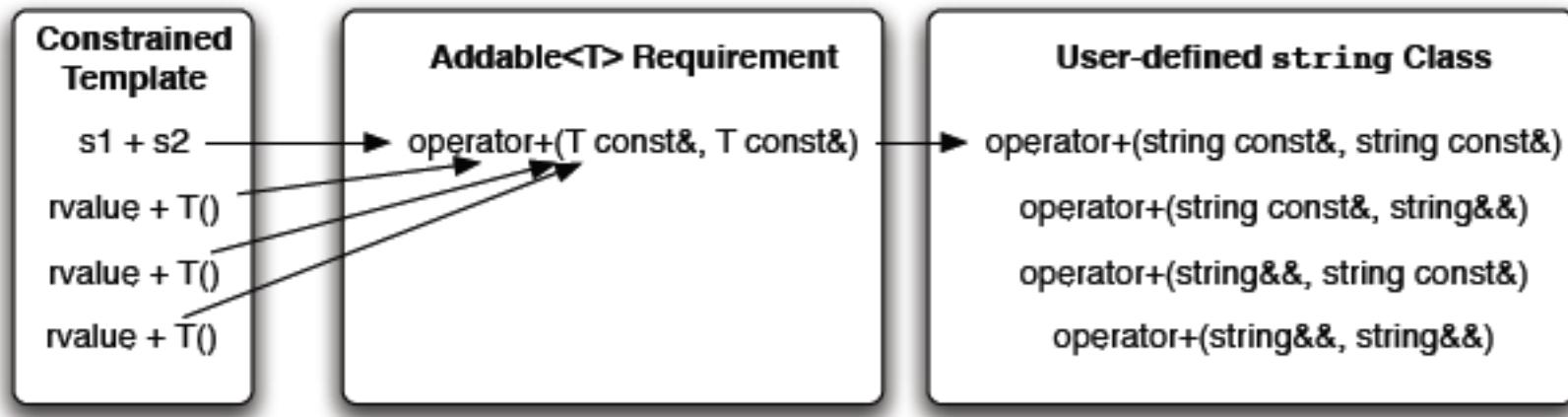
### Constrained:

lv string + lv string  
lv string + lv string  
lv string + lv string  
lv string + lv string



# Hinnant's Example in Pictures

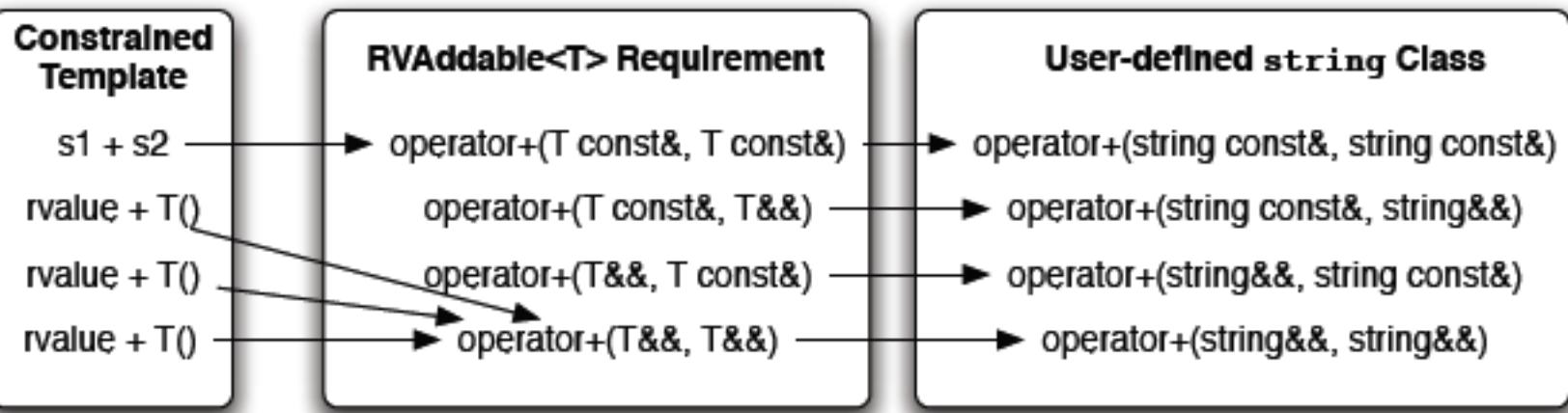
```
auto concept Addable<typename T, typename U = T> {
    typename result_type;
    result_type operator+(T const&, U const&);
}
```



# Hinnant's Example: Solution 1

## Manual Introduction of Overloads

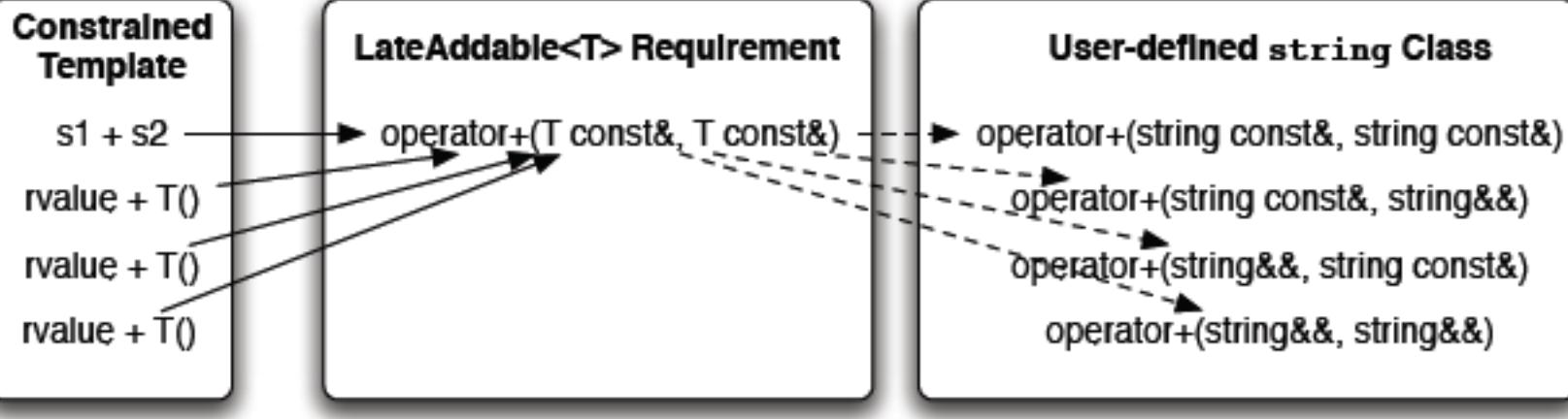
```
auto concept RVAddable<typename T, typename U = T> {
    typename result_type;
    result_type operator+(T const&, U const&);
    result_type operator+(T const&, U&&);
    result_type operator+(T&&, U const&);
    result_type operator+(T&&, U&&);
}
```



# Hinnant's Example: Solution 2

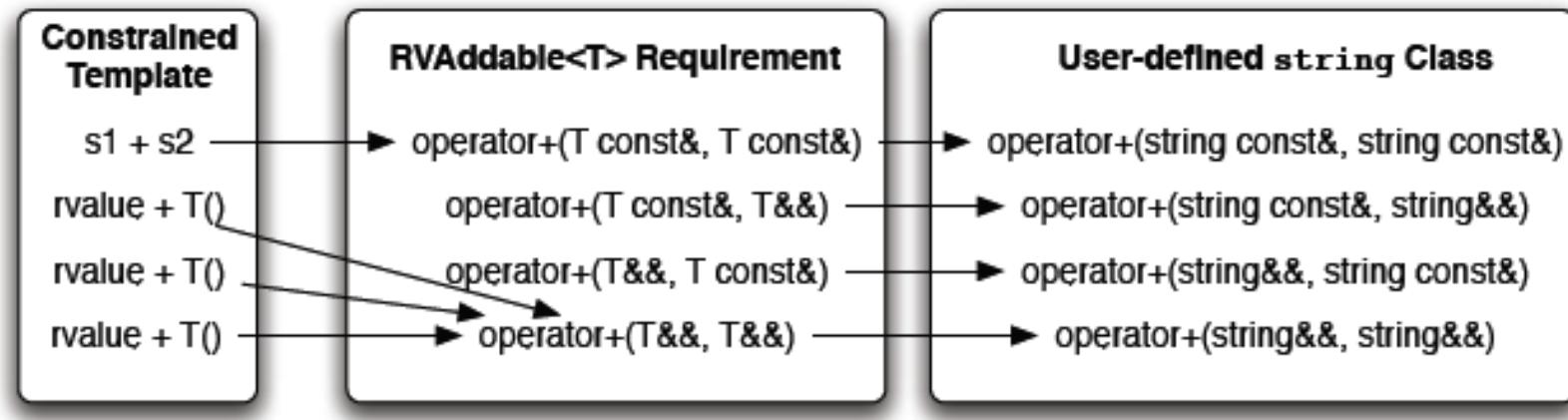
## “Eliminating” Forwarding Functions

```
auto concept LateAddable<typename T, typename U = T> {
    typename result_type;
    result_type operator+(T const&, U const&);
}
concept_map LateAddable<string> {
    typedef string result_type;
    string operator+(T const& x, T const& y) { return x + y; }
}
```

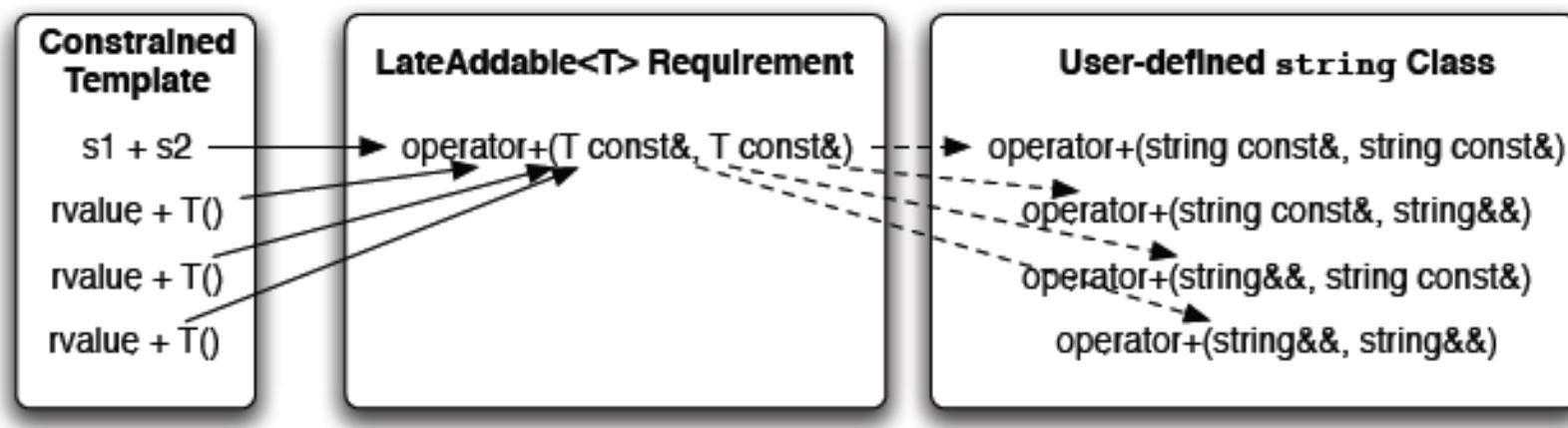


# Revisiting the Proposed Solutions

1. Manual introduction of overloads:



2. “Eliminating” forwarding functions:



# A “Sweet Middle” Alternative?

```
auto concept SweetAddable<typename T, typename U = T> {
    typename result_type;
    result_type operator+(T const&, U&&);
    result_type operator+(T&&, U const&);
}
concept_map SweetAddable<string> {
    typedef string result_type;
    result_type operator+(T const& x, U&& y) { return x + y; }
    result_type operator+(T&& x, U const& y) { return x + y; }
}
```

## Constrained Templates

$s1 + T()$   
 $rvalue + T()$   
 $rvalue + s2$

## **SweetAddable<T>** Requirement

$\operatorname{operator+}(T \text{ const\&, } T \&\&)$   
 $\operatorname{operator+}(T \&\&, T \text{ const\&})$   
 $(ANY)$

## User-defined **string** class

$\operatorname{operator+}(\text{string const\&, string const\&})$   
 $\operatorname{operator+}(\text{string const\&, string\&\&})$   
 $\operatorname{operator+}(\text{string\&\&, string const\&})$   
 $\operatorname{operator+}(\text{string\&\&, string\&\&})$



# A “Sweet Middle” Example

Minimizing and clarifying template constraints

```
template <class T>
    requires ... std::SweetAddable<T> ...
void test() {
    T s1;
    T s2;
    T s3 = s1 + T() + T() + s2;
}
int main() { test<string>(); }
```

Constrained  
Templates

s1 + T()  
rvalue + T()  
rvalue + s2

SweetAddable<T> Requirement

operator+(T const&, T&&)  
operator+(T&&, T const&)

User-defined string class

operator+(string const&, string const&)  
operator+(string const&, string&&)  
operator+(string&&, string const&)  
operator+(string&&, string&&)



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# Current C++ Concepts: Problem

- ❖ Clarity is penalized.

- Minimal requirement specification leads to **ambiguity**,
- which **is an error** in a constrained template definition,
- even though **instantiation guarantees single matches**.

- ❖ Limited control over candidate sets.

- Either **one or all** possible candidates.

- ❖ The “sweet middle” alternative is not allowed.



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# C++ Concepts: Proposed Extension

- ❖ Ambiguity is not an error, when in restricted scope.
  - ❖ Ambiguity remains as defined in the language, otherwise.
- Parameterize weak hiding over the bind environment.



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# The Combinators, Revisited

- ❖ Weak hiding selectively updates lookup results

$$s_1 \Leftrightarrow s_2 = (resolve_{ref} \circ (update \circ lookup_{ref})) s_1 ? \\ (update \circ lookup_{ref}) s_1 : lookup_{ref} s_2$$

- ❖ Define *update* based on a new concept: **Parameterized**.

```
update :: (Parameterized decl) => OverloadSet decl -> OverloadSet decl
update (os, env) = case get_env of
    Nothing          -> (os, env)
    (Just new_env)  -> (os, new_env)
```

Lookup results are coupled  
with the bind environment.



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# Parameterized example

```
class Parameterized decl where
    get_env :: Maybe (ViableSet decl -> Maybe decl)
    get_env = Nothing
                                -- default, constant bind env.
```

Bind environment

```
class (Ambiguity decl, Parameterized decl,
      Basic.Language ref decl) => Language ref decl
```

```
instance Ambiguity decl where
    ambiguity = ambiguity_is_error

instance Parameterized decl where
    get_env = Just (ambiguity_is_not_error)
```

For C++ Concepts



# Exploring Concepts Designs: Recap

- ❖ Introducing a new scoping rule: **Weak Hiding**.
- ❖ Introducing a new mechanism for name binding: **Bind<sup>x2</sup>**.
- ❖ Main Motivation:
  - Software compatibility
  - In transition from unconstrained to constrained polymorphism.
    - e.g., a transition from C++ to ConceptC++.



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# Exploring Concepts Designs: Recap

- ❖ **Bind<sup>x2</sup>** is an implementation of **weak hiding, optionally parameterized** (e.g., over the **bind environment**).
- ❖ Design variants: (Experimental)
  - Constant meaning of ambiguity: from **Ambiguity** concept.
  - Changing meaning of ambiguity: from **bind environment**.
    - e.g., Relaxing restrictions for name binding in restricted scope.



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# Applications

- ❖ Understanding current name binding mechanisms:
  - Argument-dependent lookup.
  - C++ Operators.
  - A cross-language analysis.
- ❖ Exploring concepts designs
  - Understanding the current limitations.
  - Exploring new solutions:
    - weak hiding,
    - 2-Stage name binding, and
    - parameterized weak hiding.
- ❖ Simplifying compiler designs

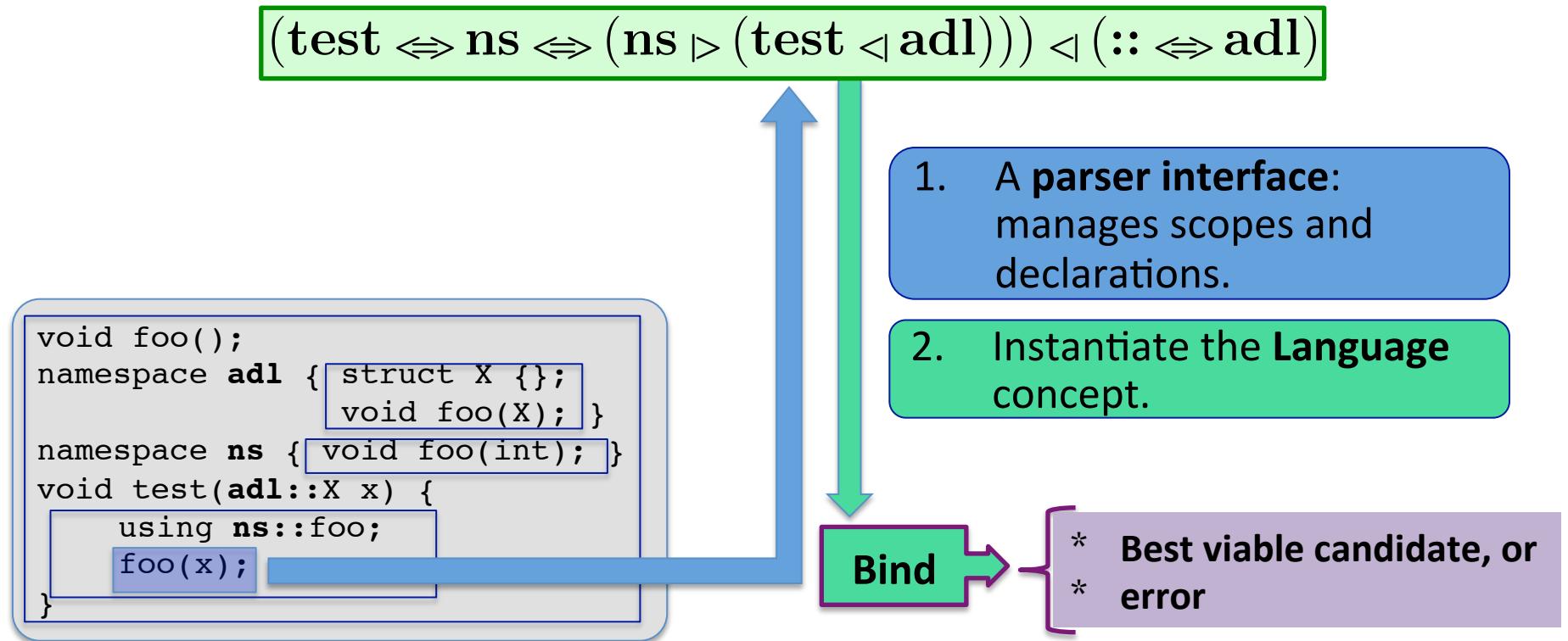


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# Using our Framework in Compilers

Option #1: For all combinator, integrate directly into parsers



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# Using our Framework in Compilers

Option #2: for weak hiding only, a conservative integration

```
concept Foo2<typename P> =  
    requires (P a, P b) { ... }  
void foo(int) { }  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

Bind<sup>x2</sup>

- \* Best viable candidate, or
- \* error



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# Prototypes Underway

[ <https://github.iu.edu/lvoufo/BindIt> ]

1. Parser interface for managing scopes:
  - Using Haskell's **Parsec** library and a **mini-C++** language.
  - Modifying existing compilers: **language-c**, **featherweight Java**, etc...
2. Executing **Bind<sup>x2</sup>**:
  - Explored in **ConceptClang**.
3. Alternative compositional view of name binding
  - as composed of some **lookup\_best** and **assess**.



# Recap

## ❖ Our name binding framework allows:

- expressing the scoping rules of a language, for a given reference, in terms of 3-4 **scope combinators**, and  $\langle , \Leftrightarrow , \triangleright , \Leftarrow$
- reasoning about the application of scoping rules generically, abstracting over the **Language** concept
  - incl. the **Ambiguity** and, optionally, **Parameterized** concepts.

## ❖ 2-stage name binding (**Bind<sup>x2</sup>**):

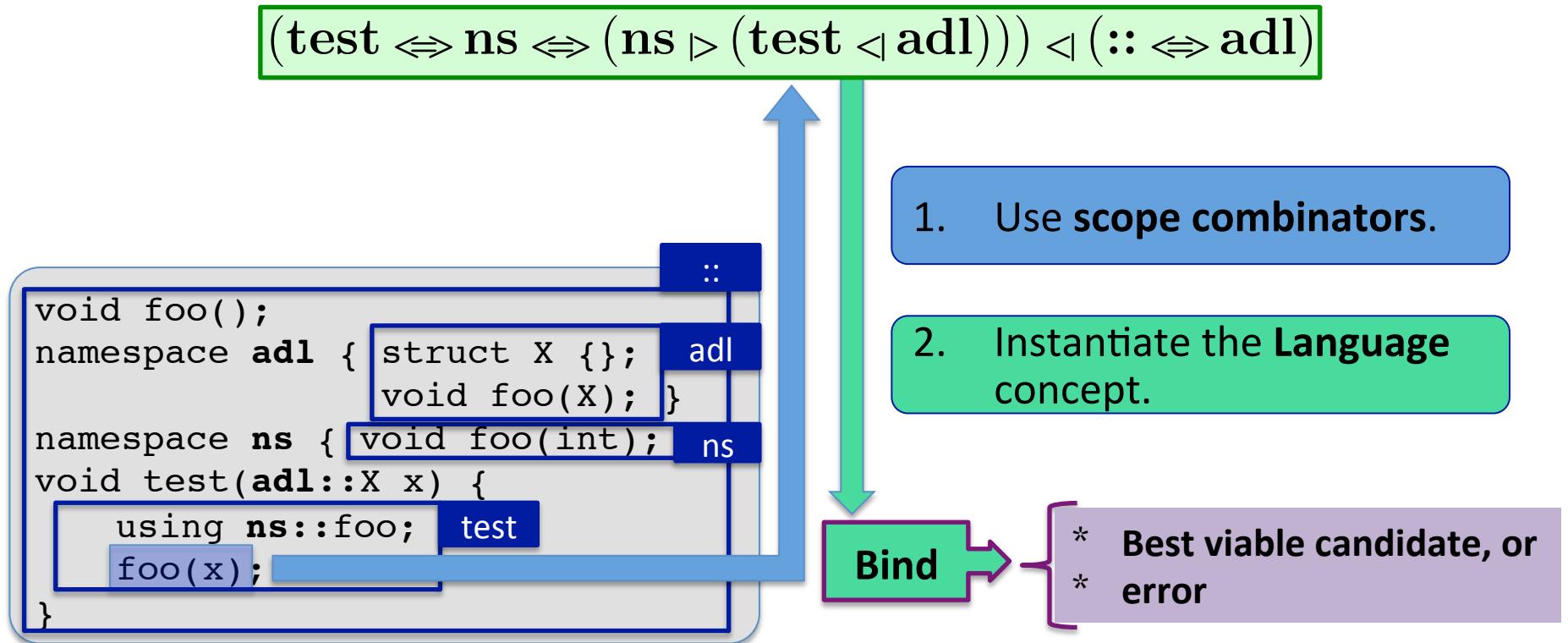
- is an implementation of **weak hiding**, optionally parameterized by the bind environment, and
- preserves valid programs in transition from C++ to ConceptC++.



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# Thank You! Questions/Comments?



[ <https://github.iu.edu/lvoufo/BindIt> ]



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