As this one has multiple end points, we might need to Write for recurrence relations which is not advisable . so instead of that we will pick up whatever is fixed. That is the starting point and we will write the record recurrence relation based on the starting point so that we can we need to write only one recurrence relation.

Base case - in Should be last row 
$$n \rightarrow (n-1, n+1)$$

## LC120- Triongle

```
if (i == triangle. &ite()-1) Recursion

return triangle. get(i).get(j)

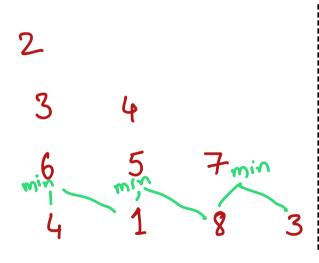
int down = current Val + recursion (i+1,j)

int diag = current Val + recursion (i+1,j+1)

return min (down # diag)
```

Any recursion problem with overlapping sub-problems, can be "memoized" using "dp" array

## Tabulation:



Tabulation is reverse of memoization. In memoization we traversed from top-bottom, here well do revent.

Create a list called "below Row" and fill in the last row details.

- → Traverse from "n-2" row
- → Store the subproblem computations in a new LinkedList called "currentRow"
- → Find the cost at each "col" in the current row
- → When we traverse through the top row, only one element remains. i.e the min cost of the path to reach from first row to last Row.

```
private int tabulationSpaceOptimisedHelper(List<List<Integer>>
triangle) {
    int n = triangle.size();
    List<Integer> belowRow = triangle.get(n-1);

for(int row = n-2;row >= 0;row--) {
    LinkedList<Integer> currentRow = new LinkedList<>();
    for(int col = row;col >= 0;col--) {
        int currentVal = triangle.get(row).get(col);
        int down = currentVal + belowRow.get(col);
        int diag = currentVal + belowRow.get(col+1);

        currentRow.addFirst(Math.min(down, diag));
    }
    belowRow = new ArrayList<>(currentRow);
}
```