

Mixed-Integer Geometric Programming

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1 From GP to MIGP

gp standard form and convex form [1], example - aircraft wings [2]

gp with integer variables, example - gate sizing [1]

generalized GP:

$$\min \max_{S \in \mathcal{S}} \sum_{i \in S} D_i : \quad (1a)$$

$$g_i X_i^{-1} \sum_{j \in F_i} (a_j + b_j X_j) = D_i \quad \forall i \in G \setminus O \quad (1b)$$

$$g_i c_i X_i^{-1} = D_i \quad \forall i \in O \quad (1c)$$

$$X_i \geq 1 \quad \forall i \in G \quad (1d)$$

$$\sum_{i \in G} h_i X_i \leq \bar{p} \quad (1e)$$

$$\sum_{i \in G} a_i X_i \leq \bar{a} \quad (1f)$$

or GP, eliminating D_i :

$$\min D : \quad (2a)$$

$$D^{-1} \sum_{i \in S \cap O} g_i c_i X_i^{-1} + D^{-1} \sum_{i \in S \setminus O} g_i X_i^{-1} \sum_{j \in F_i} (a_j + b_j X_j) \leq 1 \quad \forall S \in \mathcal{S} \quad (2b)$$

$$X_i^{-1} \leq 1 \quad \forall i \in G \quad (2c)$$

$$\bar{p}^{-1} \sum_{i \in G} h_i X_i \leq 1 \quad (2d)$$

$$\bar{a}^{-1} \sum_{i \in G} a_i X_i \leq 1 \quad (2e)$$

now let $X_i = e^{x_i}, \forall i \in G$ and $D = e^\delta$. transform into convex form by taking log of both sides of each posynomial inequality and the objective.

$$\begin{aligned}
& \min \quad \delta : & (3a) \\
& \log \left(\sum_{i \in S \cap O} g_i c_i e^{-\chi_i} + \sum_{i \in S \setminus O} g_i e^{-\chi_i} \sum_{j \in F_i} (a_j + b_j e^{\chi_j}) \right) \leq \delta & \forall S \in \mathcal{S} \quad (3b) \\
& \chi_i \geq 0 & \forall i \in G \quad (3c) \\
& \log \sum_{i \in G} h_i e^{\chi_i} \leq \log \bar{p} & (3d) \\
& \log \sum_{i \in G} a_i e^{\chi_i} \leq \log \bar{a} & (3e)
\end{aligned}$$

where in the example problem, $O = \{6, 7\}$, $\mathcal{S} = \{\{1, 4, 6\}, \{1, 4, 7\}, \{2, 4, 6\}, \{2, 4, 7\}, \{2, 5, 7\}, \{3, 5, 6\}, \{3, 7\}\}$, $F_i = \{\dots\}$ and $h_i = f_i e_i$, $\forall i \in G$.

can impose integer constraints on the X_i : a discrete number of transistors per gate: $X_i = e^{\chi_i} \in \mathbb{N}$, which implies $\chi_i \in \{\log n | n \in \mathbb{N}, n > 0\}$. note $n \geq 1$ (valid for integer problem) implies the constraint $\chi_i \geq 0$.

2 MIDCP Formulation of MIGP

dcp [5], midcp [4]

the migp, in conic form, is... **do we actually want to write with the expo cone?**

selection of log-integer variables not MIP-representable

assuming bounds, disjunctive constraint formulations

3 Solution Algorithm

solve migp like other midcp problems with pajarito, an outer approximation alg [4]. can migp have unbounded feasible set? would cause problems for pajarito

find (and perhaps improve during iterations) bounds on the log-integer variables (tighter bounds led to significant speed up and avoids ECOS numerical issues, for gate sizing example)

4 Examples and Results

small gate-sizing problem

approximate nonconvex functions piecewise, extend Woody paper [3]

Woody aircraft example

5 Future Work

disjunctive formulations in the MILP - what encodings are best

unbounded log-integer variables
cuts to the MILP or to the GP
mixtures of cones

References

- [1] Boyd, Stephen, et al. "A tutorial on geometric programming." *Optimization and engineering* 8.1 (2007): 67-127.
- [2] Hoburg, Warren, and Pieter Abbeel. "Geometric programming for aircraft design optimization." *AIAA Journal* 52.11 (2014): 2414-2426.
- [3] Hoburg, Warren, and P. Abbeel. "Fitting Geometric Programming Models to Data." *Optimization and Engineering* (2014).
- [4] Lubin, Miles, et al. "Extended Formulations in Mixed-integer Convex Programming." *arXiv preprint arXiv:1511.06710* (2015).
- [5] Grant, Michael, Stephen Boyd, and Yinyu Ye. *Disciplined convex programming*. Springer US, 2006.