# COMS 633:

## Advanced Topics in Computational Randomness Lecture Notes - Fall 2017

Professor: Dr. Jack Lutz Notes by: Alex Scheel

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### 1 Introduction

These notes were prepared by Alex Scheel.

### 2 Finite-State Gambling

Let  $\Sigma$  be an alphabet with  $2 \leq |\Sigma| < \infty$ . That is,  $\Sigma$  is finite. A probability measure on  $\Sigma$  is a function:

$$\pi: \Sigma \to [0,\infty)$$

satisfying:

$$\sum_{a \in \Sigma} \pi(a) = 1$$

A rational probability measure on  $\Sigma$  is a function:

$$\pi: \Sigma \to \mathbb{Q} \cap [0, \infty)$$

that is a probability measure on  $\Sigma$ . We write:

 $\Delta(\Sigma) = \{\text{probability measures on } \Sigma\}$   $\Delta_{\mathbb{Q}} = \{\text{rational probability measures on } \Sigma\}$   $\Delta^{+}(\Sigma) = \{\pi \in \Delta(\Sigma) | (\forall a \in \Sigma)\pi(a) > 0\}$   $\Delta_{\mathbb{Q}}^{+}(\Sigma) = \{\pi \in \Delta_{\mathbb{Q}}(\Sigma) | (\forall a \in \Sigma)\pi(a) > 0\}$ 

Note that  $\Delta(\Sigma)$  is a  $(|\Sigma|-1)$ -dimensional simplex in  $\mathbb{R}^{|\Sigma|}$ .

**Definition 2.1.** A finite-state automaton (FSA) on  $\Sigma$  is a triple:

$$A = (Q, \delta, s)$$

where:

Q is a finite set of states,  $\delta: Q \times \Sigma \to Q \text{ is a transition function,}$   $s \in Q \text{ is a start state.}$ 

Example 2.2. DIAGRAM

Given an FSA,  $A = (Q, \delta, s)$  on  $\Sigma$ , we define the extended-transition function:

$$\hat{\delta}: Q \times \Sigma^* \to Q$$

by the recursion:

$$\hat{\delta}(q, \lambda) = q$$

$$\hat{\delta}(q, wa) = \delta(\hat{\delta}(q, w), a)$$

 $\forall q \in Q, w \in \Sigma^*, \text{ and } a \in \Sigma.$ 

#### **Notational Conventions:**

- 1. We write  $\delta(q, w)$  for  $\hat{\delta}(q, w)$ .
- 2. We write  $\delta(w)$  for  $\delta(s, w)$ .

Our next objective is to endow FSAs with the ability to gamble.

**Definition 2.3.** A bet on  $\Sigma$  is a rational probability measure  $\beta \in \Delta_{\mathbb{Q}}(\Sigma)$ .

**Intuition:** Assume that:

- You have  $d \in \mathbb{Q} \cap [0, \infty)$  dollars
- You are confronting an experiment whose outcome is some element of  $\Sigma$
- You place the bet  $\beta \in \Delta_{\mathbb{O}}(\Sigma)$

This means that  $\forall a \in \Sigma$ , you are betting  $d \cdot \beta(a)$  dollars that the outcome is a. Since  $\sum_{a \in \Sigma} d\beta(a) = d$ , you have to bet all your money in this scenario. After the bet, if the martingale's outcome was a, you will have d(a), an amount that we now specify.

Suppose that the outcomes of this experiment occur according to a probability measure  $\pi \in \Delta(\Sigma)$ . What is your expected amount of money after the bet, given that you had d dollars?

Answer:

$$\sum_{a \in \Sigma} \pi(a)d(a) = \mathbb{E}_{\pi}[d(a)|d]$$

Therefore "the payoffs are fair" if:

$$d = \sum_{a \in \Sigma} \pi(a) d(a)$$

**Recall:** A bet on  $\Sigma$  is a probability measure  $\beta \in \Delta_{\mathbb{Q}}(\Sigma)$ . For now, a payoff rule on  $\Sigma$  is a probability measure  $\rho \in \Delta(\Sigma)$ .

**Intuition:** Assume that a gambler has  $d \in [0, \infty)$  dollars and places a bet,  $\beta$  on an experiment that will have an outcome that is an element of  $\Sigma$ . Placing this bet means that,  $(\forall a \in \Sigma)$ , the gambler is betting  $d\beta(a)$  dollars that the outcome is a. Note that  $d = \sum_{a \in \Sigma} d\beta(a)$ , so the gambler is required to bet all its money.

If the payoff rule is  $\rho \in \Delta(\Sigma)$ , and the actual outcome is  $a \in \Sigma$ , then the gambler will have:

$$d(a) := \frac{\beta(a)}{\rho(a)} \tag{1}$$

dollars after the bet.

Now assum that the outcome of the experiment occurs according to a probability measure,  $\pi \in \Delta(\Sigma)$ , which we call the actual probability measure of the experiment. The expected value of the gambler's amount of money after the bet (given that it has d dollars before the bet) is:

$$\mathbb{E}_{a \sim \pi}[d(a)] := \sum_{a \in \Sigma} d(a)\pi(a), \tag{2}$$

where  $a \sim \pi$  means "a is drawn according to  $\pi$ ". By (2.1) and (2.2),

$$\mathbb{E}_{a \sim \pi}[d(a)] = d \sum_{a \in \Sigma} \frac{\beta(a)\pi(a)}{\rho(a)}$$
(3)

**Observation:** If  $\rho = \pi$ , i.e., the payoff rule is the actual probability measure on  $\Sigma$ , then the payoffs are fair in the sense that  $\mathbf{E}_{a \sim \pi}[d(a) = d]$ .

**Definition 2.4.** A finite state gambler (FSA) on  $\Sigma$  is a 5-tuple,  $G = (Q, \delta, s, \beta, c)$ , where:

- $(Q, \delta, s)$  is a FSA,
- $\beta: Q \to \Delta_{\mathbb{Q}}(\Sigma)$  is the betting function,
- $c \in \mathbb{Q} \cap [0, \infty)$  is the initial capital.

### Example 2.5. DIAGRAM

#### Intuition:

Before defining the semantics of FSGs formally, let us use example 1.3 to gain some insight. Assume that G is given an input string  $w \in \{0,1\}^*$ , whose bits are chosen by independent tosses of a fair coin, and assume that the payoff rule coincides with this. Let  $d_G(w)$  denote the amount of money that G has after betting on w. Then:

$$d_G(\lambda) = c = 1 \qquad d_G(1) = \frac{4}{3} = 2\frac{2}{3}d_G(\lambda)d_G(11) = 2\frac{1}{3}d_G(1) = \frac{8}{9} \qquad d_G(110) = 2\frac{2}{3}d_G(11) = frac3227...$$

**Definition 2.6.** If  $G = (Q, \delta, s, \beta, c)$  is an FSG and  $\pi \in \Delta^+(\Sigma)$ , then the  $\pi$ -martingale of G is the function:

$$d_G^{\pi}: \Sigma^* \to [0, \infty)$$

defined by the recursion:

$$d_G^{\pi}(\lambda) = cd_G^{\pi}(wa) = d_G^{\pi}(w) \frac{\beta(\delta(w))(a)}{\pi(a)}$$

**Definition 2.7.** (Ville - 1939). Let  $\pi \in \Delta(\Sigma)$ . A  $\pi$ -martingale on  $\Sigma$  is a function:

$$d: \Sigma^* \to [0, \infty)$$

that satisfies,  $(\forall w \in \Sigma^*)$ :

$$d(w) = \sum_{a \in \Sigma} d(wa)\pi(a) \tag{4}$$

**Obervation:** For every FSG G on  $\Sigma$  and every probability measure  $\pi \in \Delta^+(\Sigma)$ ,  $d_G^{\pi}$  is a  $\pi$ -martingale.