Saturn: a library of verified concurrent data structures for OCaml 5

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Why Saturn?

Github: ocaml-multicore/saturn

- A collection of concurrent-safe data structures for OCAML 5:
 - ⋄ well-tested
 - ⋄ benchmarked
 - ⋄ optimized
 - ⋄ verified

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- Writing concurrent code is hard

Buggy queue with size

```
module Oueue = Saturn.Oueue
type 'a t = {size : int Atomic.t; queue : 'a Queue.t}
let create () =
    {size = Atomic.make 0; queue = Queue.create ()}
let push t msg =
 Atomic.incr t.size;
 Queue.push t.queue msq
let pop_opt t =
 match Queue.pop_opt t.queue with
  | Some elt -> Atomic.decr t.size; Some elt
  | None -> Atomic.set t.size 0; None
let size t = Atomic.get t.size
```

Buggy queue with size

```
let test () =
  let queue = create () in
  let d1 = Domain.spawn (fun () -> push queue 1) in
  let d2 = Domain.spawn (fun () -> pop_opt queue |>
      ignore) in
  Domain.join d1;
  Domain.join d2;
  pop_opt queue |> ignore;
  size queue
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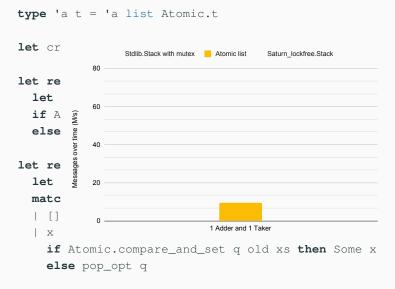
In 10 to 20% of the tries, the test returns a size of -1.

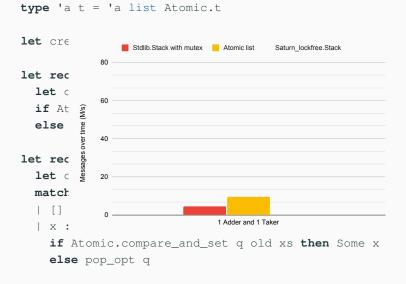
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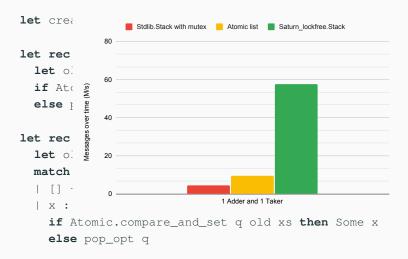
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- Writing concurrent code is hard
 - bugs that can be hard to reproduce and understand

```
type 'a t = 'a list Atomic.t
let create () = Atomic.make []
let rec push q a =
  let old = Atomic.get q in
  if Atomic.compare_and_set q old (a :: old) then ()
  else push q a
let rec pop_opt q =
  let old = Atomic.get q in
  match old with
  | [] -> None
  | x :: xs ->
    if Atomic.compare_and_set q old xs then Some x
    else pop_opt q
```





type 'a t = 'a list Atomic.t



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 - ♦ Queues

- A collection of concurrent-safe data structures for OCAML 5
 - ♦ Queues
 - multi-producer, multi-consumer
 (based on Michael-Scott queue algorithm)
 - o single-producer, single-consumer
 - o single-producer, multi-consumer
 - bounded, blocking (opened PRs)

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 - ⋄ Work-stealing deque
 - ♦ Stacks
 - ♦ Hashtable (almost in a PR)
 - \diamond Skiplist: a sorted linked list with $o(\log(n))$ operations
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- Need something else?
 - → Open an issue on ocaml-multicore/saturn

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- Most available data structures are lock-free
 - → Two libraries: SATURN and SATURN LOCKFREE

Testing in Saturn

What should (and can) be tested?

- Correctness
- Linearizability
- Progress (i.e. lock freedom)

How?

- Lin / STM (ocaml-multicore/multicoretests):
 - ⋄ Correctness
 - Linearizability
- Dscheck (ocaml-multicore/dscheck):
 - ♦ Correctness
 - ♦ Lock-freedom

About optimizations

Numerous micro-optimizations:

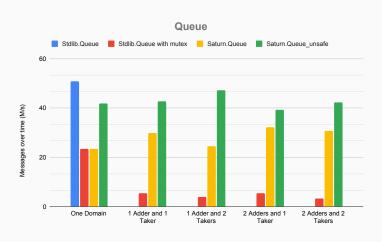
- Improving algorithms,
- Preventing false sharing,
- ⋄ Removing indirections etc...
- Experimental: some optimizations use **Obj.magic** (e.g. Saturn.Queue_unsafe)

About optimizations

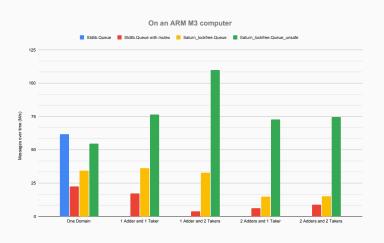
Numerous micro-optimizations:

- Improving algorithms,
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- Experimental: some optimizations use **Obj.magic** (e.g. Saturn.Queue_unsafe)
 - ightarrow a source of inspiration for langage features for efficient concurrent code (e.g. opened PR on atomic record fields)

Benchmarks



Benchmarks



Formal verification

Testing a concurrent algorithm is hard due to the number of potential interleavings.

Formal verification is here to help us!





From OCaml to Coq

let rec push t v =

```
let old = Atomic.get t in
  let new_ = v :: old in
  if not (Atomic.compare and set t old new ) then (
   Domain.cpu relax () ;
   push t v
                                              OCAML
                                                Coo
Definition stack_push : val :=
  rec: "stack push" "t" "v" =>
    let: "old" := !"t" in
    let: "new" := 'Cons( "v", "old" ) in
    ifnot: CAS "t" "old" "new" then (
      Yield ::
      "stack_push" "t" "v"
```

Proving linearizability

Theorems for free from separation logic specifications
Birkedal, Dinsdale-Young, Guéneau, Jaber, Svendsen & Tzevelekos

Writing concurrent protocols in Iris

```
Definition stack inv t \iota : iProp \Sigma :=
  \exists 1 \gamma,
  \lceil t = \#1 \rceil * meta l nroot \gamma *
  inv \iota (
     \exists vs, 1 \mapsto 1st to val vs * stack model<sub>2</sub> \gamma vs
Lemma stack push spec t \iota v :
  <<< stack inv t \iota
   \forall \forall vs, stack\_model t vs >>>
             stack_push t v @ ↑ι
   <<< stack_model t (v :: vs)</pre>
   RET (); True
                                        >>>.
Proof.
Oed.
```





Relaxed memory model (future work)

Cosmo: a concurrent separation logic for multicore OCAML, Mével, Jourdan & Pottier

Formal verification of a concurrent bounded queue in a weak memory model, Mével & Jourdan

Stack benchmarks

