Saturn: a library of verified concurrent data structures for OCaml 5

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Plan

Intro

Formal verification

Testing a concurrent algorithm is hard due to the number of potential interleavings.

Formal verification is here to help us!





From OCaml to Coq

let rec push t v =

```
let old = Atomic.get t in
  let new_ = v :: old in
  if not (Atomic.compare and set t old new ) then (
   Domain.cpu relax () ;
   push t v
                                              OCAML
                                                Coo
Definition stack_push : val :=
  rec: "stack push" "t" "v" =>
    let: "old" := !"t" in
    let: "new" := 'Cons( "v", "old" ) in
    ifnot: CAS "t" "old" "new" then (
      Yield ::
      "stack_push" "t" "v"
```

Proving linearizability

Theorems for free from separation logic specifications
Birkedal, Dinsdale-Young, Guéneau, Jaber, Svendsen & Tzevelekos

Relaxed memory model (future work)

Cosmo: a concurrent separation logic for multicore OCAML Mével, Jourdan & Pottier

Writing concurrent protocols in Iris

```
Definition stack inv t \iota : iProp \Sigma :=
  \exists 1 \gamma,
  \lceil t = \#1 \rceil * meta l nroot \gamma *
  inv \iota (
     \exists vs, 1 \mapsto 1st to val vs * stack model<sub>2</sub> \gamma vs
Lemma stack push spec t \iota v :
  <<< stack inv t \iota
   \forall \forall vs, stack\_model t vs >>>
             stack_push t v @ ↑ι
   <<< stack_model t (v :: vs)</pre>
   RET (); True
                                        >>>.
Proof.
Oed.
```

Thank you for your attention!