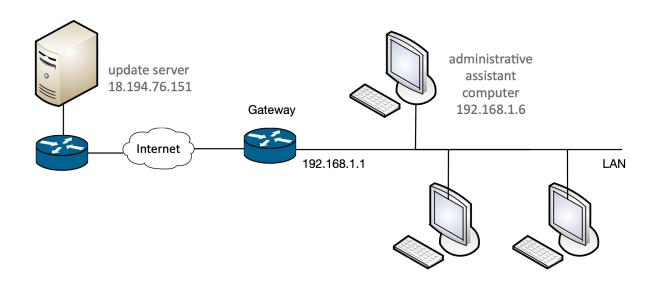
# **Exercises**Network Protocols Attacks

Computer Security

### Question

As part of your job as a security analyst, one of your clients discovers that their network is compromised. In particular, from an early analysis, they have ground to suspect that the start of the compromise was a <u>network attack against the computer of the administrative assistant</u>.

Consider the following (simplified) schema of the company network.



1. Your client managed to capture the network traffic on the administrative assistant's computer (IP address 192.168.1.6 and MAC address dc:a9:04:7a:ce:29) when the attack was taking place. During the traffic capture, the computer was automatically updating a well-known accounting software from the software vendor's web server (IP address 18.194.76.151 and MAC address dc:a6:03:01:02:fe). You also know that the IP address of the LAN interface of the company's network gateway is 192.168.1.1, and its MAC address is b6:28:97:ca:b7:48.

```
1 dc:a9:04:7a:ce:29 \rightarrow ff:ff:ff:ff:ff: ARP Who has 192.168.1.1? Tell 192.168.1.6 2 38:60:77:b9:79:98 \rightarrow dc:a9:04:7a:ce:29 ARP 192.168.1.1 is at 38:60:77:b9:79:98 3 b6:28:97:ca:b7:48 \rightarrow dc:a9:04:7a:ce:29 ARP 192.168.1.1 is at b6:28:97:ca:b7:48 4 192.168.1.6 (dc:a9:04:7a:ce:29) \rightarrow 18.194.76.151 (38:60:77:b9:79:98) TCP SYN 5 18.194.76.151 (38:60:77:b9:79:98) \rightarrow 192.168.1.6 (dc:a9:04:7a:ce:29) TCP SYN, ACK 6 38:60:77:b9:79:98 \rightarrow dc:a9:04:7a:ce:29) \rightarrow 18.194.76.151 (38:60:77:b9:79:98) TCP ACK 8 38:60:77:b9:79:98 \rightarrow dc:a9:04:7a:ce:29) \rightarrow 18.194.76.151 (38:60:77:b9:79:98) TCP ACK 8 38:60:77:b9:79:98 \rightarrow dc:a9:04:7a:ce:29) \rightarrow 18.194.76.151 (38:60:77:b9:79:98) TCP HTTP GET /downloads/software-update.exe
```

10 18.194.76.151 (38:60:77:b9:79:98)  $\rightarrow$  192.168.1.6 (dc:a9:04:7a:ce:29) **TCP** HTTP 200 OK ...

(assume that the traffic is captured directly from the network interface card of the employee's PC)

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**1.1 [2 points].** Describe the attack going on in the network, specifying the name and providing a short explanation of how the attack works *in general*.

ARP spoofing, see slide.

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Sniffing or manipulation of the traffic to\from the compromised machine. We can rule out DOS as the traffic passes (there are responses from the server).

Likely, given the scenario (malware infection), the attack is targeted at tampering with the data in transit rather than (or in addition to) sniffing.

1.3 [1 point]. Can you tell the IP address of the attacker? And the MAC address?

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IP address: no

MAC address: the attacker is using 38:60:77:b9:79:98, but it could very well be a spoofed address.

**1.4 [1 point].** Given only the above packet capture, can you tell whether the attacker is located (i.e., on the LAN, on the same network of the web server, or on an arbitrary Internet-connected network)? Why?

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The attacker is located on the same network of the target machine, i.e., on the LAN.

## Question

Suppose that you are the network security administrator of the network 131.168.0.0/24, with gateway 131.168.0.1 and DNS servers 131.168.0.100 and 131.168.0.101. While examining the network activity, you notice a DHCP offer packet coming from IP 131.168.0.10 with gateway set to 131.168.0.5. Answer the following questions and provide a reason. Answers with no reason will not give any point.

•	What kind of attack do you suspect, and how does it work?	

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DHCP poisoning. Someone is trying to trick a client connected to 131.168.0.0/24 into believing that 131.168.0.5 is the gateway, by sending a crafted DHCP offer before, which comes before the real offer sent out by the real DHCP server.

• Why such an attack works?

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Because the DHCP protocol does not support authentication, so the client must blindly believe any DHCP offer that it sees, and because an arbitrary client can race (and win) against the real DHCP server. • Can you tell the IP address of the host where those packets come from?

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Not really. 131.168.0.10 is the sender of the DHCP offer, but the address may be spoofed. We could look at the MAC address of the scender, but it could be spoofed as well.

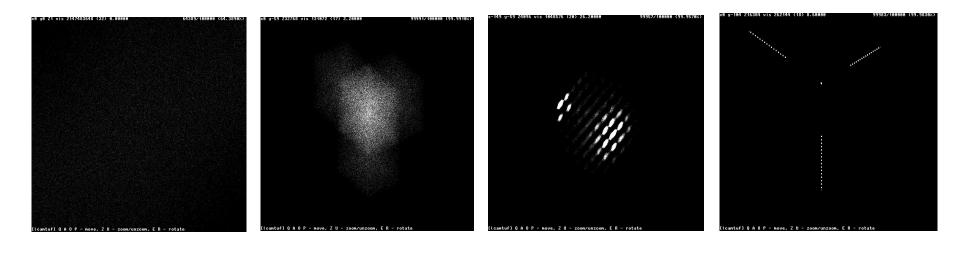
•	Can you tell the I	P address or n	etwork addres	ss of the victim?

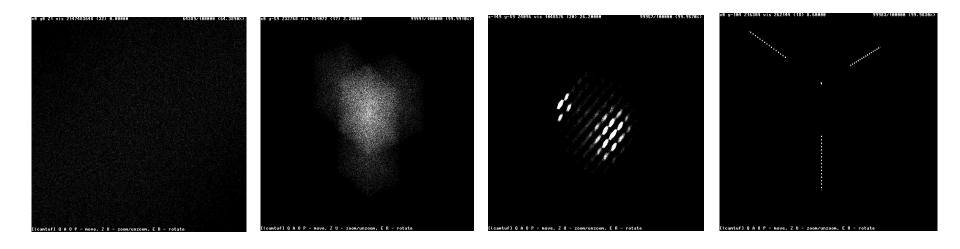
• Can you tell the IP address or network address of the victim?

From the above information there is little evidence to say that, although the potential victims are those that will receive and accept the spoofed DHCP offers. So, likely, 131.168.0.0/24.

# Question

- Explain why the Initial Sequence Number (ISN) in a TCP connection should be random.
- Which of the following TCP stacks is more vulnerable to the above attack, and why (in brief)?





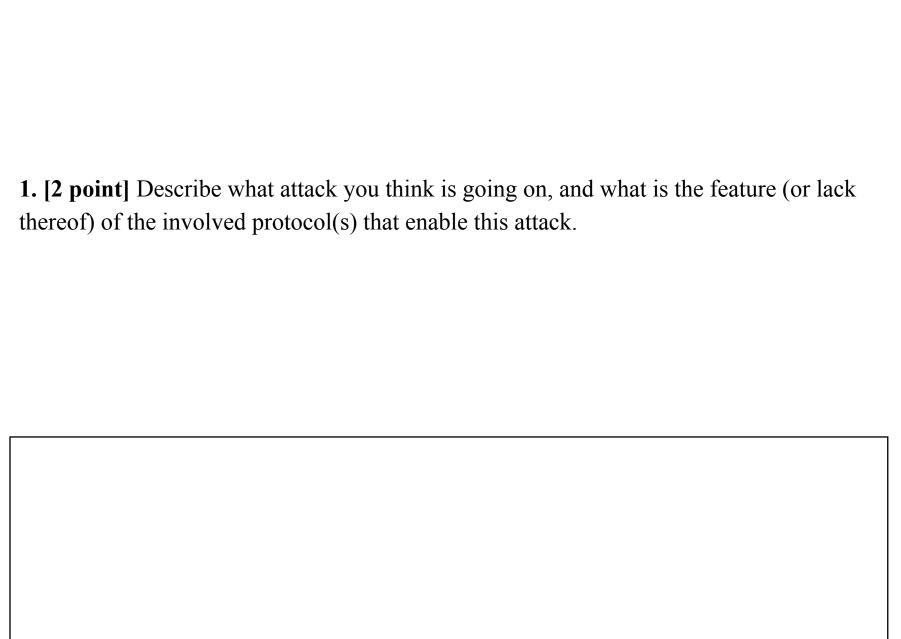
- 1. If an attacker can guess the ISN value of a TCP connection, they can forge a response TCP packet that will be valid (because the acknowledgment of the ISN is correct). This means that the attacker can spoof a TCP connection, even if they are blind, i.e. not on the route of packets (MITM)
- 2. Bottom right is easiest, as only a few values are used and there's no uniform spreading in the space

Ref: https://www.ietf.org/rfc/rfc1948.txt

#### Question

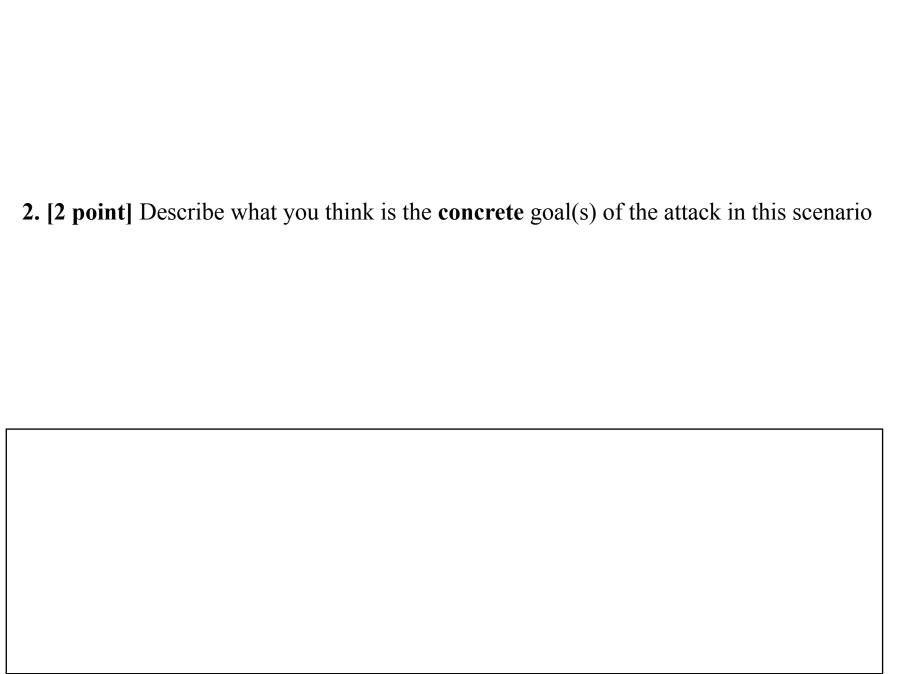
A network analyst is analyzing some traffic captured from a network belonging to Politecnico di Milano. The network is: 131.175.0.0/16. In particular, we suspect that a database server, whose IP address is 131.175.14.12, is victim of an attack. Indeed, observing the network traffic, we notice the following pattern:

```
131.175.14.12 \rightarrow 131.175.255.255
                                           [ICMP] Echo (ping) request
              (broadcast)
131.175.0.2 \rightarrow 131.175.14.12
                                       [ICMP] Echo (ping) reply
7a:ce:29 \rightarrow ff:ff:ff (broadcast) [ARP] Who has 131.175.14.12?
Tell 131,175,0,2
4b:74:28 \rightarrow 7a:ce:29
                                        [ARP] 131.175.14.12 is at
4b:74:28
131.175.0.3 \rightarrow 131.175.14.12
                                        [ICMP] Echo (ping) reply
131.175.0.4 \rightarrow 131.175.14.12
                                        [ICMP] Echo (ping) reply
131.175.255.251 \rightarrow 131.175.14.12
                                        [ICMP] Echo (ping) reply
131.175.255.252 \rightarrow 131.175.14.12
                                        [ICMP] Echo (ping) reply
131.175.255.253 \rightarrow 131.175.14.12
                                        [ICMP] Echo (ping) reply
```



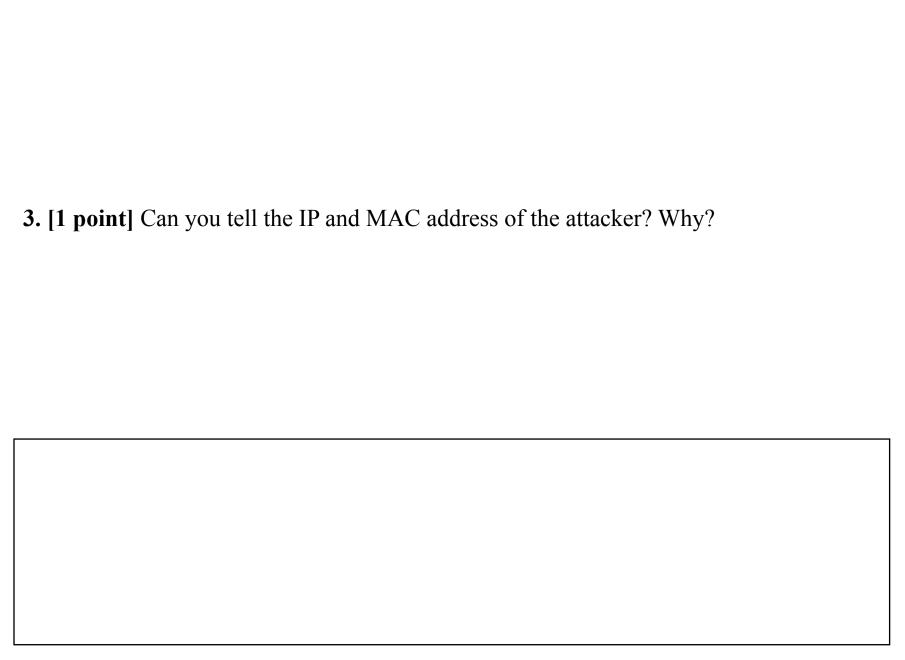
**1. [2 point]** Describe what attack you think is going on, and what is the feature (or lack thereof) of the involved protocol(s) that enable this attack.

PING Smurf - see slides ....
Authentication



2. [2 point] Describe what you think is the concrete goal(s) of the attack in this scenario

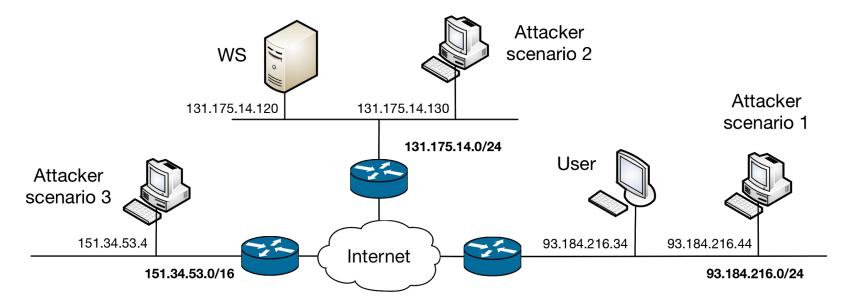
The goal is to saturate the resources of the victim using other machines on the network as an amplification mean. This results in a denial of service.



**3.** [1 point] Can you tell the IP and MAC address of the attacker? Why?

No, the IP address is spoofed for sure, and the MAC address can be spoofed as well.

Consider the following network diagram:



A user, with IP address 93.184.216.34, is attempting to download a software executable from a webserver in the Politecnico di Milano network, <a href="http://downloads.polimi.it">http://downloads.polimi.it</a>, with IP address 131.175.14.120, over the HTTP protocol (no HTTPS, no signatures, <a href="nothing">nothing</a>). Assume that the user's browser already cached the IP address of downloads.polimi.it (i.e., <a href="it does not perform any DNS request">it does not perform any DNS request</a>), and that there is <a href="nothing">nothing</a>) firewall involved. An attacker, who knows that the user is about to download this software, wants to target our user by carrying out an <a href="attack to replace the downloaded software with a piece of malware">attack to replace the downloaded software with a piece of malware</a>.

**Scenario 1** (attacker: 93.184.216.44; same network and broadcast domain of the user):

*Scenario 2* (attacker: 131.175.14.130; same network of web server, but different than user):

*Scenario 3* (attacker: 151.34.53.4; attacker, user and webserver on three different networks):

**Scenario 1** (attacker: 93.184.216.44; same network and broadcast domain of the user):

#### ARP spoofing ...

**Scenario 2** (attacker: 131.175.14.130; same network of web server, but different than user):

**Scenario** 3 (attacker: 151.34.53.4; attacker, user and webserver on three different networks):

**Scenario 1** (attacker: 93.184.216.44; same network and broadcast domain of the user):

ARP spoofing ...

**Scenario 2** (attacker: 131.175.14.130; same network of web server, but different than user):

ARP spoofing ... (same as before but this time target the server)

*Scenario 3* (attacker: 151.34.53.4; attacker, user and webserver on three different networks):

**Scenario 1** (attacker: 93.184.216.44; same network and broadcast domain of the user):

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**Scenario 2** (attacker: 131.175.14.130; same network of web server, but different than user):

ARP spoofing ... (same as before but this time target the server)

*Scenario 3* (attacker: 151.34.53.4; attacker, user and webserver on three different networks):

No attack possible, because HTTP uses TCP and, if sequence numbers are correctly implemented, not possible to perform TCP hijacking. Also, DNS poisoning not possible as DNS response already cached.

**2.** [2 points]. Describe the security issue with the proposed ISN implementation, and propose a way to solve the issue

2. [2 points]. Describe the security issue with the proposed ISN implementation, and propose a way to solve the issue *Can predict ISN -> solve by using random ISN* 

**3.** [2 points]. Describe how the attacker can perform the above attack, this time *exploiting the security issues raised by the custom ISN implementation*. Describe all the steps and assumptions that you need to perform this attack.

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TCP hijacking. We can guess the ISN of the SYN packet sent by the victim (the user), we can spoof the webserver IP and send a "correct" SYN+ACK to it and, if we can guess the content of the request (we do), subsequent packets. This way we can send a different payload. In parallel we can also use TCP hijacking to send a fake RST to the actual server spoofing the user's IP address.

*Problem\assumption:* we need to know the ephemeral port used by the client to initiate the connection.

**4.** [1 points]. Assume you are the network security administrator of the network of the attacker (in the scenario 3 and assuming the custom TCP initial sequence number implementation), and that you control the border router between the network 151.34.53.0/24 and the rest of the Internet. Propose a way to prevent the attack. Can the administrator of the other two border routers in the diagram deploy the same mitigation, and obtain the same result? Why?

**4.** [1 points]. Assume you are the network security administrator of the network of the attacker (in the scenario 3 and assuming the custom TCP initial sequence number implementation), and that you control the border router between the network 151.34.53.0/24 and the rest of the Internet. Propose a way to prevent the attack. Can the administrator of the other two border routers in the diagram deploy the same mitigation, and obtain the same result? Why?

We can filter the packets with source IPs not belonging to our network. Other routers can't do this, they can only filter out packets coming from "outside" with spoofed IPs belonging to their network, but it's of little use in this scenario.

# The End