

**Exercise 8.5** Let  $(x, y) \notin \text{gph } M$  so that  $y \neq Qx$ . Let  $e$  be a unit eigen vector of  $Q$  with corresponding eigen value  $\lambda \geq 0$  (we know it's non-negative because  $Q$  is positive). Fix  $0 < \epsilon < \lambda^{-1}|\langle e, y - Qx \rangle|$  (possible because  $y \neq Qx$ ). Also put

$$q = -\epsilon \text{sgn}(\langle e, y - Qx \rangle) e$$

Then we have

$$\langle x - (x - q), y - Q(x - q) \rangle = \langle q, y - Mx \rangle + \langle q, Qq \rangle = -\epsilon |\langle e, y - Qx \rangle| + \epsilon^2 \lambda = \epsilon(\epsilon \lambda - |\langle e, y - Qx \rangle|) < 0$$

so that including  $(x, y)$  in the graph breaks monotonicity, hence  $M$  is maximal.

**Exercise 8.11** By hypothesis  $\exists \bar{x} \in \text{argmin } f$ . Let  $\phi$  be the forward complete solution to (8.10). We have

$$\frac{d}{dt} \|\phi(t) - \bar{x}\|^2 = 2 \langle \phi(t) - \bar{x}, \dot{\phi}(t) \rangle \leq 2(f(\bar{x}) - f(\phi(t))) \leq 0$$

where the last inequality comes from the fact that  $\bar{x}$  is a minimizer and the first from the fact that  $\phi$  is a solution, i.e.

$$\dot{\phi}(t) \in -\partial f(\phi(t)) \implies f(\bar{x}) \geq f(\phi(t)) - \langle \dot{\phi}(t), \bar{x} - \phi(t) \rangle$$

I claim  $\exists t_k \rightarrow \infty$  so that  $f(\phi(t_k)) \rightarrow f(\bar{x})$ . If there isn't, then for large  $t$   $f(\phi(t)) \geq f(\bar{x}) + \delta$  for some  $\delta > 0$ . This would mean for large  $t$   $f(\bar{x}) - f(\phi(t)) \leq -\delta$ , and so from above, for large  $t$ , we'd have  $\frac{d}{dt} \|\phi(t) - \bar{x}\|^2 \leq -2\delta$ . This draws a contradiction because that means  $\|\phi(t) - \bar{x}\| \rightarrow -\infty$ , but Fact 8.7 tells us  $\partial f$  is maximal montone, and so Prop 8.9 (a) tells us  $\|\phi(t) - \bar{x}\|$  is bounded (because  $\bar{x}$  as a constant function is a solution to (8.10)).

Quick question: Do I really need to invoke Prop 8.9 here? In Prop 1.12 you state that going to  $-\infty$  is absurd, but I can't see how to show it's "absurd" in this case without effectively showing  $\phi$  is bounded, which I think I need Prop 8.9 for? Not sure, seems pretty heavy weight. Anyway, back to the proof..

Since  $\|\phi(t) - \bar{x}\|$  is bounded, we know  $\phi(t_k)$  is bounded and hence  $\phi(t_{k_i}) \rightarrow \bar{y}$  for some subsequence  $t_{k_i}$ . Because  $f$  is lsc we know

$$f(\bar{y}) = f\left(\lim_{i \rightarrow \infty} \phi(t_{k_i})\right) \leq \liminf_{i \rightarrow \infty} f(\phi(t_{k_i})) = f(\bar{x})$$

and so  $\bar{y} \in \text{argmin } f$ . All that's left is to show  $\phi(t) \rightarrow \bar{y}$  (i.e. without a sequence). Again we can use Fact 8.7 combined with Prop 8.9 (a), this time to see  $\|\phi(t) - \bar{y}\|$  is nonincreasing, since  $\bar{y}$  as a constaint function is a solution to (8.10). Combined with  $\phi(t_{k_i}) \rightarrow \bar{y}$  gives us  $\phi(t) \rightarrow \bar{y}$ .

**Exercise 8.12** In 1.12 we have  $f(x, y) = 2x + y$  and

$$C = \{(x, y) \mid x \geq 0, y \geq 0, 2x + y \geq 2\}$$

This gives us

$$\nabla f(x, y) = (2, 1), N_C(x, y) = \begin{cases} \{(0, 0)\} & (x, y) \in \text{int } C, \\ \{0\} \times (-\infty, 0] & x > 1, y = 0, \\ (-\infty, 0] \times \{0\} & x = 0, y > 2, \\ (-2, -1)[0, \infty) & 2x + y = 2, x \neq 0, y \neq 0, \\ \text{con } N_C(1.5, 0) \cup N_C(0.5, 1) & x = 1, y = 0, \\ \text{con } N_C(0, 2.5) \cup N_C(0, 5, 1) & x = 0, y = 2, \\ \emptyset & (x, y) \notin C \end{cases}$$

And so we have

$$-\nabla f(x, y) - N_C(x, y) = \begin{cases} \{(-2, -1)\} & (x, y) \in \text{int } C, \\ \{-2\} \times [-1, \infty) & x > 1, y = 0, \\ [-2, \infty) \times \{-1\}, & x = 0, y > 2, \\ (-2, -1) + (2, 1)[0, \infty) & 2x + y = 2, x \neq 0, y \neq 0, \\ (-2, -1) - \text{con } N_C(1.5, 0) \cup N_C(0.5, 1) & x = 1, y = 0, \\ (-2, -1) - \text{con } N_C(0, 2.5) \cup N_C(0, 5, 1) & x = 0, y = 2, \\ \emptyset & (x, y) \notin C \end{cases}$$

In each case above we have

$$\begin{aligned} m(\{(-2, -1)\}) &= (-2, -1) \\ m(\{-2\} \times [-1, \infty)) &= (-2, 0) \\ m([-2, \infty) \times \{-1\}) &= (0, -1) \\ m((-2, -1) + (2, 1)[0, \infty)) &= (0, 0) \end{aligned}$$

And the last two cases are the same as  $m((-2, -1) + (2, 1)[0, \infty))$  (because of the union). Plugging along this gives us

$$m(-\nabla f(x) - N_C(x)) = \begin{cases} (-2, -1) & (x, y) \in \text{int } C \\ (-2, 0) & x > 1, y = 0, \\ (0, -1) & x = 0, y > 2, \\ (0, 0) & 2x + y = 2, \\ \infty & x \notin C, \end{cases}$$

which matches exactly the diff eq in 1.13, so indeed the dynamics match (8.12).

Now to show (8.13) matches as well we need to compute the tangent cone, which I've done below with the help of Proposition 7.11 (tangent and normal cones are polar to one another)

$$T_C(x, y) = \begin{cases} \mathbb{R}^n & (x, y) \in \text{int } C, \\ \mathbb{R} \times [0, \infty) & x > 1, y = 0, \\ [0, \infty) \times \mathbb{R} & x = 0, y > 2, \\ \{(x, y) \mid y \geq -2x\} & 2x + y = 2, x \neq 0, y \neq 0, \\ T_C(1.5, 0) \cap T_C(0.5, 1) & x = 1, y = 0, \\ T_C(0, 2.5) \cap T_C(0, 5, 1) & x = 0, y = 2, \\ \emptyset & (x, y) \notin C, \end{cases}$$

Thankfully  $-\nabla f(x) = (-2, -1)$  is constant, so it's easy to see (after drawing some pictures, which I've elided because tikz is hard)

$$P_{T_C(x, y)}(-\nabla f(x)) = P_{T_C(x, y)}(-2, -1) = \begin{cases} (-2, -1) & (x, y) \in \text{int } C, \\ (-2, 0) & x > 1, y = 0, \\ (0, -1) & x = 0, y > 2, \\ (0, 0) & 2x + y = 2, \end{cases}$$

This matches exactly the dynamics in 1.13, so we've shown the dynamics match (8.13).

**Exercise 8.16** If we consider (8.5) a system of two equations, fixing  $u_1 = z, u_2 = \dot{z}$  then solving (8.5) is the same as solving

$$\begin{pmatrix} \dot{u}_1 \\ \dot{u}_2 \end{pmatrix} \in \left( -\frac{k}{m}u_1 - \frac{d}{m}u_2 - \frac{c_{max}}{m}\partial|u_2| \right) := F(u_1, u_2)$$

It's easy to check the opposite of the righthand side isn't monotone, hence not maximal monotone.

To show a solution exists we need to show  $F$  satisfies the basic assumptions:

(osc)  $F$  is a matrix plus a subdifferential of a finite lsc convex function composed with a projection: Fact 8.7 combined with Prop 8.6 gives us the subdifferential is osc, the rest of the operations are osc and locally bounded as they're single valued (classically) continuous operations, so the composition is osc, so Prop 2.26 gives us the whole operation is osc.

(loc bdd) This is trivial to see as, again,  $F$  is a matrix plus a subdifferential, the latter being bounded by  $\frac{c_{max}}{m}\mathbb{B}$ .

(convex) The righthand side gets its mass (i.e. is non-single-valued) from the subdifferential, which is convex, hence  $F$  takes on convex values.

With this, we can invoke Theorem 5.11 to get a solution  $\phi : [0, T] \rightarrow \mathbb{R}^2$  from any  $u_0 \in \mathbb{R}^2$ , for some  $T > 0$ .

**Exercise 8.18** To show  $\bar{x}$  is constant I'll show the derivative vanishes:

$$\dot{\bar{x}} = \frac{1}{l} \sum_{i=1}^l \dot{x}_i = \frac{1}{l} \sum_{i=1}^l \left( \frac{1}{l} \sum_{j=1}^l x_j - x_i \right) = \frac{1}{l} \left( \sum_{j=1}^l x_j - \sum_{i=1}^l x_i \right) = 0$$

For the second part of the problem fix  $\psi_i(t) = \|\phi_i(t) - \bar{x}\|^2$  then when  $\phi_i(t) \neq \bar{x}$

$$\frac{d}{dt}\psi_i(t) = 2\langle \phi_i(t) - \bar{x}, \dot{\phi}_i(t) \rangle = 2\langle \phi_i(t) - \bar{x}, \bar{x} - \phi_i(t) \rangle = -2\psi_i(t)$$

and solving this we get  $\psi_i(t) = \|\phi_i(0) - \bar{x}\|^2 e^{-2t}$  so

$$\|\phi_i(t) - \bar{x}\|^2 = \psi_i(t) = \psi_i(0)e^{-2t} = \|\phi_i(0) - \bar{x}\|^2 e^{-2t} \implies \|\phi_i(t) - \bar{x}\| = \|\phi_i(0) - \bar{x}\| e^{-t}$$

**Exercise 8.19** I'll show  $\bar{x}$  is constant by showing its derivative is 0:

$$l\dot{\bar{x}} = \sum_{i=1}^l \sum_{j=1}^l a_{ij}(x_j - x_i) = \sum_{i=1}^l \sum_{j=1}^l a_{ij}x_j - \sum_{i=1}^l x_i \sum_{j=1}^l a_{ij} = \sum_{j=1}^l x_j \sum_{i=1}^l a_{ij} - \sum_{i=1}^l x_i \sum_{j=1}^l a_{ij} = \sum_{j=1}^l x_j \sum_{i=1}^l a_{ij} - \sum_{i=1}^l x_i \sum_{j=1}^l a_{ji} = 0$$

where the second to last equality comes from  $a_{ij} = a_{ji}$  and the last equality can be seen by swapping i/j in the right most double sum.

**Exercise 8.20** Before jumping in I want to introduce some notation:  $e_q$  is the  $q$ th standard basis vector,  $x_{k,q}$  denotes the  $q$ th coordinate of the  $k$ th vector in  $x$  and  $\delta_{i,j}$  is the kronecker delta, i.e. it vanishes when  $i \neq j$  and is 1 otherwise. Using these together for some  $x \in \mathbb{R}^n$   $\frac{\partial}{\partial x_{k,q}} \|x_i\|^2 = 2\delta_{i,k}x_{i,q}$ . Now I'm ready to calculate the following:

$$\begin{aligned}
\frac{\partial f(x)}{\partial x_{k,q}} &= \frac{1}{4} \sum_{i,j=1}^l a_{ij} \frac{\partial}{\partial x_{k,q}} (\|x_i\|^2 + \|x_j\|^2 - 2\langle x_i, x_j \rangle) \\
&= \frac{1}{4} \sum_{i=1}^l \sum_{j=1}^l a_{ij} (2\delta_{i,k}x_{i,q} + 2\delta_{j,k}x_{j,q} - 2\delta_{i,k}x_{j,q} - 2\delta_{j,k}x_{i,q}) \\
&= \frac{1}{2} \sum_{i=1}^l \sum_{j=1}^l a_{ij} (\delta_{i,k}(x_{i,q} - x_{j,q}) + \delta_{j,k}(x_{j,q} - x_{i,q})) \\
&= \frac{1}{2} \left( \sum_{j=1}^l a_{kj}(x_{k,q} - x_{j,q}) + \sum_{i=1}^l a_{ik}(x_{k,q} - x_{i,q}) \right) \\
&= \frac{1}{2} \left( \sum_{i=1}^l a_{ik}(x_{k,q} - x_{i,q}) + \sum_{i=1}^l a_{ik}(x_{k,q} - x_{i,q}) \right) = \left( \sum_{i=1}^l a_{ik}(x_k - x_i) \right) \cdot e_q
\end{aligned}$$

Now if we solve  $\dot{x} = -\nabla f(x)$  then we're solving for every  $k = 1, 2, \dots, l$  and  $q = 1, 2, \dots, m$

$$\dot{x}_{k,q} = -\frac{\partial f(x)}{\partial x_{k,q}} = \left( \sum_{j=1}^l a_{jk}(x_j - x_k) \right) \cdot e_q$$

which equivalently can be put as solving for every  $i = 1, 2, \dots, l$

$$\dot{x}_i = \sum_{j=1}^l a_{ji}(x_j - x_i) = \sum_{j=1}^l a_{ij}(x_j - x_i)$$

which is exactly (8.26).

In order to show  $\text{argmin } f = A$  we only need to show  $\text{argmin } f \subset A$ . Let  $x \in \text{argmin } f$ , then  $f(x) = 0$ .

I'm unsure why, but this fact escaped me for far too long. But, I'm glad to admit: this wasn't a reading comprehension issue, this time I was in fact struggling to do math.

For  $j \neq 1$  there is a finite path  $k_1, k_2, \dots, k_q$  from  $k_1 = 1$  to  $k_q = j$  where  $a_{k_p k_{p+1}} > 0$  (by connectedness). Because  $f(x) = 0$  we must have  $\|x_{k_p} - x_{k_{p+1}}\| = 0$  (since the sum is over non-negative values, each summed element must vanish, so the positivity of the scalar causes the norm to vanish). This means that  $x_1 = x_{k_2} = \dots = x_j$ . Since  $j$  was arbitrary we know  $x_1 = x_2 = \dots = x_l$  so that  $x \in A$ .

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Leaving old proof here for historical record:

In order to show  $\text{argmin } f = A$  we only need to show  $\text{argmin } f \subset A$ , which we'll do inductively. Specifically we'll induct on the following hypothesis for  $2 < n \leq l$

**Claim** If  $n$  agents are connected, then  $x \in \text{argmin } f \implies x_i = x_j$  for all  $n$  connected agents indexed by  $i, j$

For the base case  $a_{12} > 0$  and  $a_{ij} = 0$  otherwise. By above equivalence

$$x \in \operatorname{argmin} f \implies \nabla f(x) = 0 \implies 0 = \dot{x}_1 = \sum_{j=3}^l a_{1j}(x_j - x_1) + a_{12}(x_2 - x_1) = a_{12}(x_2 - x_1) \implies x_1 = x_2$$

For the inductive step we assume the claim holds for  $n$  connected agents and must show the claim holds for  $n+1$ . W.l.o.g. suppose the first  $n+1$  agents are connected so that  $a_{jk} = 0 \ \forall j > n+1, \forall k$  and there's some  $k < n+1$  so that  $a_{(n+1)k} > 0$ . Now, again using the above equivalence if  $x \in \operatorname{argmin} f$

$$\begin{aligned} 0 = \dot{x}_{n+1} &= \sum_{j=1}^l a_{(n+1)j}(x_j - x_{n+1}) \\ &= \sum_{j=1}^n a_{(n+1)j}(x_j - x_{n+1}) \\ &= \sum_{j=1}^n a_{(n+1)j}(x_j - x_k) + a_{(n+1)j}(x_k - x_{n+1}) \\ &= (x_k - x_{n+1}) \sum_{j=1}^n a_{(n+1)j} \end{aligned}$$

where the last equality is found by the inductive hypothesis (i.e. the first  $n$  agents are connected, so  $x_j = x_k \ \forall j \leq n$ ). Since  $a_{(n+1)k} > 0$  and  $a_{(n+1)j} \geq 0$  otherwise  $\sum_{j=1}^n a_{(n+1)j} > 0$ , so  $x_k = x_{n+1}$ .

Since the above holds for all  $n \leq l$ , it holds in particular for  $n = l$ , and so if all  $l$  agents are connected,

$$x \in \operatorname{argmin} f \implies x_1 = x_2 = \dots = x_l \implies x \in A$$

**Exercise 8.26** The solution is given by

$$\phi(t) = \begin{cases} (2 + \frac{t}{5}, 1 - \frac{2t}{5}) & 0 \leq t \leq \frac{5}{2} \\ (t, 0) & t > \frac{5}{2} \end{cases}$$

To show this, first note that

$$-N_{C(t)}(x, y) = \{w \mid w \cdot (x' + t - x, y' - y) \geq 0 \ \forall (x', y') \in C\}$$

Then for  $0 \leq t \leq \frac{5}{2}$   $\dot{\phi}(t) = \frac{1}{5}(1, -2)$  and we have

$$\frac{1}{5}(1, -2) \cdot \left(x' + t - 2 - \frac{t}{5}, y' - 1 + \frac{2t}{5}\right) = \frac{1}{5}\left(x' + t - 2 - \frac{t}{5} - 2y' + 2 - \frac{4t}{5}\right) = \frac{1}{5}(x' - 2y') \geq 0$$

whenever  $x', y' \in C$ , and so  $\dot{\phi}(t) \in -N_{C(t)}(\phi(t))$ . When  $t > \frac{5}{2}$   $\dot{\phi}(t) = (1, 0)$  and

$$(1, 0) \cdot (x' + t - t, y') = x' \geq 0$$

for any  $x' \in C$ , so  $\dot{\phi}(t) \in -N_{C(t)}(\phi(t))$  again, thus  $\phi$  is indeed the full solution.