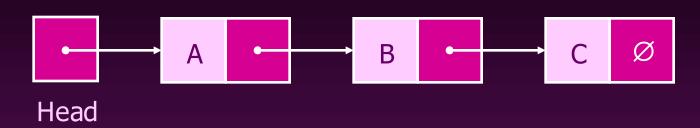
Linked Lists

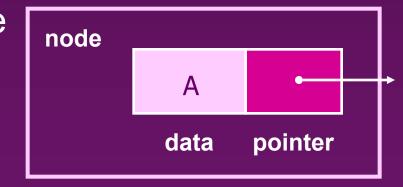
List Overview

- Linked lists
 - Abstract data type (ADT)
- Basic operations of linked lists
 - Insert, find, delete, print, etc.
- ► Variations of linked lists
 - Circular linked lists
 - Doubly linked lists

Linked Lists



- A linked list is a series of connected nodes
- Each node contains at least
 - A piece of data (any type)
 - Pointer to the next node in the list
- **►** Head: pointer to the first node
- **► The last node points to NULL**



A Simple Linked List Class

- ► We use two classes: Node and List
- ► Declare Node class for the nodes
 - data: double-type data in this example
 - next: a pointer to the next node in the list

A Simple Linked List Class

- Declare List, which contains
 - head: a pointer to the first node in the list.

 Since the list is empty initially, head is set to NULL
 - Operations on List

```
class List {
public:
      List(void) { head = NULL; } // constructor
      ~List(void);
                                        // destructor
      bool IsEmpty() { return head == NULL; }
      Node* InsertNode(int index, double x);
      int FindNode(double x);
      int DeleteNode(double x);
      void DisplayList(void);
private:
      Node* head;
};
```

A Simple Linked List Class

- Operations of List
 - IsEmpty: determine whether or not the list is empty
 - InsertNode: insert a new node at a particular position
 - FindNode: find a node with a given value
 - DeleteNode: delete a node with a given value
 - DisplayList: print all the nodes in the list

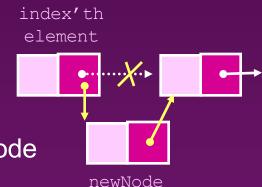
- ► Node* InsertNode(int index, double x)
 - Insert a node with data equal to x after the index'th elements. (i.e., when index = 0, insert the node as the first element; when index = 1, insert the node after the first element, and so on)
 - If the insertion is successful, return the inserted node.

 Otherwise, return NULL.

 (If index is < 0 or > length of the list, the insertion will fail.)

Steps

- 1. Locate index'th element
- Allocate memory for the new node
- 3. Point the new node to its successor
- 4. Point the new node's predecessor to the new node



- ► Possible cases of InsertNode
 - 1. Insert into an empty list
 - 2. Insert in front
 - 3. Insert at back
 - 4. Insert in middle
- But, in fact, only need to handle two cases
 - Insert as the first node (Case 1 and Case 2)
 - Insert in the middle or at the end of the list (Case 3 and Case 4)

```
Try to locate
Node* List::InsertNode(int index, double x)
                                                 index'th node. If it
       if (index < 0) return NULL;</pre>
                                                 doesn't exist,
       int currIndex =
                             1;
                                                 return NULL.
       Node* currNode =
                            head;
       while (currNode && index > currIndex) {
              currNode = currNode->next;
              currIndex++;
       if (index > 0 && currNode == NULL) return NULL;
```

```
Node* List::InsertNode(int index, double x) {
       if (index < 0) return NULL;</pre>
       int currIndex = 1;
       Node* currNode = head;
       while (currNode && index > currIndex) {
              currNode = currNode->next;
              currIndex++;
       if (index > 0 && currNode == NULL) return NULL;
       Node* newNode =
                                   Node;
                            new
       newNode->data =
                            Х;
                                                Create a new node
```

```
Node* List::InsertNode(int index, double x) {
       if (index < 0) return NULL;</pre>
       int currIndex = 1;
       Node* currNode = head;
       while (currNode && index > currIndex) {
              currNode = currNode->next;
              currIndex++;
       if (index > 0 && currNode == NULL) return NULL;
       Node* newNode =
                                    Node;
                             new
                                              Insert as first element
       newNode->data =
                             х;
                                                       head
       if (index == 0)  {
              newNode->next =
                                    head;
              head
                                    newNode;
                                                         newNode
```

```
Node* List::InsertNode(int index, double x) {
       if (index < 0) return NULL;</pre>
       int currIndex = 1;
       Node* currNode = head;
       while (currNode && index > currIndex) {
              currNode = currNode->next;
              currIndex++;
       if (index > 0 && currNode == NULL) return NULL;
       Node* newNode = new
                                   Node;
       newNode->data =
                           х;
       if (index == 0) {
              newNode->next =
                                   head;
                                           Insert after currNode
                                   newNode 🔏
              head
                                                   currNode
       else {
              newNode->next =
                                   currNode->next;
              currNode->next =
                                   newNode;
       return newNode;
                                                       newNode
```

Finding a node

- ➡ int FindNode (double x)
 - Search for a node with the value equal to x in the list.
 - If such a node is found, return its position. Otherwise, return 0.

```
int List::FindNode(double x) {
    Node* currNode = head;
    int currIndex = 1;
    while (currNode && currNode->data != x) {
        currNode = currNode->next;
        currIndex++;
    }
    if (currNode) return currIndex;
    return 0;
}
```

- int DeleteNode(double x)
 - Delete a node with the value equal to x from the list.
 - If such a node is found, return its position. Otherwise, return 0.

→ Steps

- Find the desirable node (similar to FindNode)
- Release the memory occupied by the found node
- Set the pointer of the predecessor of the found node to the successor of the found node
- ► Like InsertNode, there are two special cases
 - Delete first node
 - Delete the node in middle or at the end of the list

```
int List::DeleteNode(double x) {
                                           Try to find the node with
       Node* prevNode =
                              NULL;
                                           its value equal to x
       Node* currNode =
                              head;
       int currIndex =
                              1;
       while (currNode && currNode->data != x) {
               prevNode
                                      currNode;
               currNode
                                      currNode->next;
               currIndex++;
```

```
int List::DeleteNode(double x) {
       Node* prevNode =
                         NULL;
       Node* currNode = head;
       int currIndex =
                            1;
       while (currNode && currNode->data != x) {
              prevNode
                                    currNode;
                                    currNode->next;
              currNode
              currIndex++;
                                        prevNode currNode
          (currNode) {
              if (prevNode) {
                     prevNode->next =
                                            currNode->next;
                      delete currNode;
       return 0;
```

```
int List::DeleteNode(double x) {
      Node* prevNode = NULL;
      Node* currNode = head;
       int currIndex = 1;
      while (currNode && currNode->data != x) {
             prevNode
                           = currNode;
             currNode
                                  currNode->next;
             currIndex++;
         (currNode) {
             if (prevNode) {
                    prevNode->next = currNode->next;
                     delete currNode;
              else {
                    head
                                          currNode->next;
                    delete currNode;
              return currIndex;
                                         head
                                              currNode
      return 0;
```

Printing all the elements

- ► void DisplayList (void)
 - Print the data of all the elements
 - Print the number of the nodes in the list

Destroying the list

- ~ ~List(void)
 - Use the destructor to release all the memory used by the list.
 - Step through the list and delete each node one by one.

```
List::~List(void) {
  Node* currNode = head, *nextNode = NULL;
  while (currNode != NULL)
  {
    nextNode = currNode->next;
    // destroy the current node
    delete currNode;
    currNode = nextNode;
}
```

Using List

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Number of nodes in the list: 3

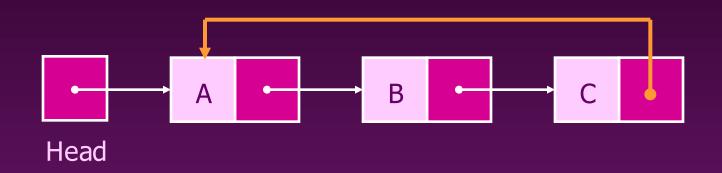
result

```
5.0 found
                                                4.5 not found
int main(void)
                                                Number of nodes in the list: 2
       List list;
       list.InsertNode(0, 7.0); // successful
       list.InsertNode(1, 5.0); // successful
       list.InsertNode(-1, 5.0); // unsuccessful
       list.InsertNode(0, 6.0); // successful
       list.InsertNode(8, 4.0); // unsuccessful
       // print all the elements
       list.DisplayList();
       if(list.FindNode(5.0) > 0) cout << "5.0 found" << endl;
       else
                                   cout << "5.0 not found" << endl;</pre>
       if(list.FindNode(4.5) > 0) cout << "4.5 found" << endl;
                                    cout << "4.5 not found" << endl;</pre>
       else
       list.DeleteNode(7.0);
       list.DisplayList();
       return 0;
```

Variations of Linked Lists

Circular linked lists

■ The last node points to the first node of the list

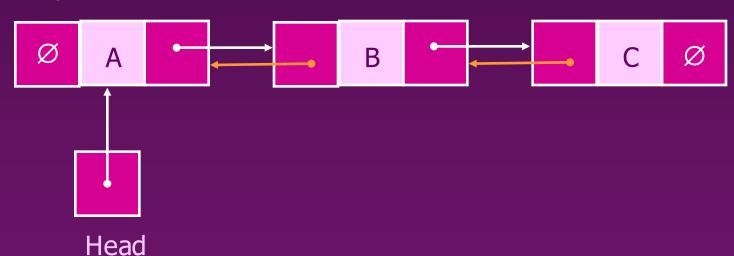


■ How do we know when we have finished traversing the list? (Tip: check if the pointer of the current node is equal to the head.)

Variations of Linked Lists

Doubly linked lists

- Each node points to not only successor but the predecessor
- There are two NULL: at the first and last nodes in the list
- Advantage: given a node, it is easy to visit its predecessor. Convenient to traverse lists backwards



Array versus Linked Lists

- Linked lists are more complex to code and manage than arrays, but they have some distinct advantages.
 - **Dynamic**: a linked list can easily grow and shrink in size.
 - The We don't need to know how many nodes will be in the list. They are created in memory as needed.
 - □ In contrast, the size of a C++ array is fixed at compilation time.
 - Easy and fast insertions and deletions