

# Mathematical Foundation for Machine Learning

51-22 AIMLCZC 416

Assignment 1

by

Group 18

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Q1 We have studied that during Gaussian Elimination we can write  $U = E_m E_{m-1} \cdots E_1 A$  where the matrices  $E_i$  are elementary transformations. Is this true for any arbitrary invertible  $n \times n$  matrix? If it is true for an arbitrary invertible  $n \times n$  matrix, provide a justification. If it is not true for an arbitrary  $n \times n$  matrix, explain why and show what modifications you will make to the equation  $U = E_m E_{m-1} \cdots E_1 A$  to make it work for any arbitrary invertible  $n \times n$  matrix.

We know that we can solve a linear system of equations  $Ax=b$  using Gaussian Elimination

We can do this because what Gaussian Elimination does is, it reduces the linear system of equations  $Ax=b$  to a much simpler form  $A'x=b'$  (i.e. the row-echelon form) using a series of elementary transformations while keeping the solution set the same. Such that after applying Gaussian Elimination, we can easily use back substitution to find the value of  $x$ .

The elementary operations allowed are -

- ① Exchange of any two rows.
- ② Multiplying any row by a constant
- ③ Adding a multiple of one row to another

Note -

④ All the elementary operations can be represented as an elementary matrix. where the  $n \times n$  elementary matrix  $E$  is the result of applying one elementary operation to a  $n \times n$  identity matrix

⑤ The elementary matrices are invertible  $E^{-1}E = I = EE^{-1}$

④ An elementary matrix where only addition of rows is performed in a way such that  $R_i \leftarrow R_i - kR_j$  where  $i > j$  is a lower triangular matrix

$$\begin{bmatrix} 1 & 0 & \dots & 0 & 0 \\ * & 1 & \dots & 0 & 0 \\ : & : & \ddots & : & : \\ * & * & \dots & 1 & 0 \\ * & * & \dots & * & 1 \end{bmatrix}$$

lower

⑤ If we multiply two triangular matrices the result is also a lower triangular matrix.

Now let's assume that we are able to reduce a matrix  $A$  to row-echelon form (i.e. some upper triangular matrix)  $U$  using Gaussian Elimination without performing any row exchanges.

In that case, we can write Gaussian Elimination as a series of elementary matrix multiplication on  $A$  such that:

$$E_m E_{m-1} \dots E_1 A = U$$

$$\Rightarrow L^{-1} A = U \quad [\text{where } E_m E_{m-1} \dots E_1 = L^{-1}]$$

$$\Rightarrow LL^{-1} A = LU \quad [\text{multiplying both sides with } L]$$

$$\Rightarrow A = LU \quad [L L^{-1} = I \text{ as elementary matrices are invertible}]$$

However, if at least one row exchange is needed  $A$  cannot be factorized to  $LU$

This is because the elementary matrix required to perform a row swap is achieved by swapping the columns of a Identity matrix

Example → The elementary matrix

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \quad \text{will swap row 2 and 3}$$

And these elementary matrices are not lower triangular in nature. As a result when they are multiplied with a lower triangular matrix, the result is no longer lower triangular.

An elementary matrix that represents a row exchange is called a Permutation Matrix  $P$ .

Therefore, if in order to reduce a matrix  $A$  to row-echelon form  $U$  requires some row exchanges, we can write the Gaussian Elimination as:  $A = P L U$

We also know that for a  $n \times n$  invertible matrix  $A$  (i.e.  $\det(A) \neq 0$ ) a pure LU decomposition exists if and only if all its leading principal minors are non-zero. However, the PLU decomposition always exists.

We can easily prove this with an example: —

Let,  $A = \begin{bmatrix} 1 & 1 & 2 \\ 1 & 1 & 1 \\ 2 & 3 & 4 \end{bmatrix}$  be a  $3 \times 3$  invertible matrix

By using the Gauss Jordan Elimination we can easily find  $A^{-1}$  as

$$[A|I] \xrightarrow{\dots} [I|A^{-1}]$$

$$\left[ \begin{array}{ccc|ccc} 1 & 1 & 2 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 2 & 3 & 4 & 0 & 0 & 1 \end{array} \right] \xrightarrow{\dots} \left[ \begin{array}{ccc|ccc} 1 & 0 & 0 & 1 & 2 & -1 \\ 0 & 1 & 0 & -2 & 0 & 1 \\ 0 & 0 & 1 & 1 & -1 & 0 \end{array} \right] \quad \text{check } A^{-1} A = I$$

$$\left[ \begin{array}{ccc} 1 & 2 & -1 \\ -2 & 0 & 1 \\ 1 & -1 & 0 \end{array} \right] \left[ \begin{array}{ccc} 1 & 1 & 2 \\ 1 & 1 & 1 \\ 2 & 3 & 4 \end{array} \right] = \left[ \begin{array}{ccc} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{array} \right]$$

We can see that  $A$  can not be factorized to LU all the leading principle minors are not non-zero.

$$|a_{11}| = 1 \quad |a_{33}| = (4-3) - (4-2) + 2(3-2) = 1 \quad \text{but} \quad |a_{22}| = 0$$

Meaning  $A$  cannot be reduced to row-echelon form  $U$  without row exchange (in this case  $R_2 \leftrightarrow R_3$ )

Therefore, we can factorize  $A$  as PLU

$$\left[ \begin{array}{ccc} 1 & 1 & 2 \\ 1 & 1 & 1 \\ 2 & 3 & 4 \end{array} \right] = \left[ \begin{array}{ccc} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{array} \right] \left[ \begin{array}{ccc} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 1 & 0 & 1 \end{array} \right] \left[ \begin{array}{ccc} 1 & 1 & 2 \\ 0 & 1 & 0 \\ 0 & 0 & -1 \end{array} \right]$$

Therefore, we can conclude that  $A = (E_m E_{m-1} \dots E_1)U$  is not true for any arbitrary invertible  $n \times n$  matrix as sometimes elementary row exchanges are required to reduce to row-echelon form. In such cases, we must re-write Gaussian Elimination as  $\star$

$$A = (P_m P_{m-1} \dots P_1) (E_m E_{m-1} \dots E_1) U$$

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Q-2 To solve systems like  $Ax=b$  where  $A$  is a invertible  $n \times n$  matrix we write a program  $\text{Solve}(A, b)$  that takes a matrix  $A$  and right-hand  $b$  as input & computes the solution to  $Ax=b$ . Suppose that algo. used by  $\text{Solve}(A, b)$  is the augmented matrix method. Let us say we need to solve  $k$  system of type  $Ax=b$ , where the right hand side changes, but the left hand side stays the same. We can do this by making  $k$  invocations to the procedure  $\text{Solve}(A, b)$ . Can you come up with a better way of solving such systems & characterize the improvement in operation count compared with making  $k$  calls to  $\text{Solve}(A, b)$ ?

Answer:

Observation

1)  $A$  is an invertible  $n \times n$  (square) matrix.

Inference

This means the below property holds true for  $A_{n \times n}$  -

$$A^{-1} A = I_n$$

And in this case, this property can also be applied to simplify a linear system of equations as follows

$$Ax = b$$

$$A^{-1} A x = A^{-1} b$$

$$I_n x = A^{-1} b$$

$$x = A^{-1} b$$

2) The function  $\text{Solve}(A, b)$  uses the augmented matrix method to solve  $Ax = b$ .

This means the function  $\text{Solve}(A, b)$  uses Gaussian Elimination with back-substitution (row-echelon form) to solve  $Ax = b$ .

We can also solve it with Gauss Jordan elimination.

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### Observation

3. We are attempting to solve K system of the type  $Ax=b$ .  
With constraint right-hand side changes while the left hand side does not.

### Inference

This means that we are attempting to solve a linear system of equations K times.

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$
$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2$$
$$\dots$$

$$a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n$$

where over k iterations:

- (a) the coefficients  $a_{ij}$  ( $i \geq 1, j \leq n$ ) (i.e. matrix A) are not changing
- (b) the variables  $x_i$  ( $i \geq 1, j \leq n$ ) are changing.
- (c) and as a result, the values  $b_j$  ( $1 \geq i \leq n$ ) are also changing.

4. We have to suggest an optimal implementation for  $\text{Solve}(Ax, b)$  in terms of operations performed. When  $\text{Solve}(A, b)$  is called k times under the above mentioned constraint

This means that we have to show that when  $\text{Solve}(A, b)$  is called k times, the Gauss Jordan elimination method does a lot fewer operations.

Case I: Calculating the number of operations performed when  $\text{Solve}(A, b)$  is called k times using augmented matrix method.

↳ To solve a linear equation  $Ax=b$ , using augmented matrix or gaussian elimination with back-substitution method our goal is to reduce the augmented matrix  $[A|b]$  using elementary operations to the row-echelon form.

↳ A matrix is said to be in row-echelon form when:

- \* All rows containing at least one non-zero element is on top of rows containing zero.

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↳ looking at the non-zero rows only, the first non-zero element is on to the left (i.e. pivot or leading coefficient) is always strictly on the right of the pivot above it.

↳ The steps of Gaussian elimination are:

1) Write  $Ax=b$  as a augmented matrix  $[A \mid b]$

2) get 1 in the  $i^{th}$  row of the  $i^{th}$  column

3) Use row i to get 0's in the  $i^{th}$  column of rows  $i+1$  to  $n$

4) Repeat steps 2 & 3 from  $i=1$  to  $n$ .

5) Change the augmented matrix back into a linear system of equation

6) Use back substitution to solve for the variables.

Augmented Matrix

$$\left[ \begin{array}{cccc|c} * & * & \dots & * & * \\ * & * & \dots & * & * \\ * & * & \dots & * & * \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ * & * & \dots & * & * \end{array} \right] \rightarrow \left[ \begin{array}{cccc|c} 1 & * & \dots & * & * \\ 0 & * & \dots & * & * \\ 0 & * & \dots & * & * \\ 0 & * & \dots & * & * \\ 0 & * & \dots & - & * \\ 0 & * & \dots & * & * \end{array} \right]$$

Get 1 in the 1<sup>st</sup> row of 1<sup>st</sup> column  
Use row 1 to get 0 in the 1<sup>st</sup> column  
of rows 2 to n

Total operation performed: n.

$$x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$

$$x_2 + \dots + a_{2n}x_n = b_2$$

$$x_n = b_n$$

Use  $x_n = b_n$  to back  
substitute the values of  $x$ , to  
 $x_{n-1}$

Total operation performed: n.

Therefore, we can conclude that to solve  $Ax=b$  just once the Gaussian elimination we have to perform:

$(n + (n-1) + \dots + 1)$  operations to reduce to the row-echelon form +  $(n$  operations to back-substitute the values of  $x$ )

$$\Rightarrow \frac{n(n+1)}{2} + n \text{ operations}$$

Now, since we intend to solve  $Ax=b$  k times whether the right hand side changes, but the left hand side stays the same, we must perform the complete Gaussian elimination process from start to finish k numbers of time.

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We need to do this as the augmented matrix  $[A|b_i]$  ( $1 \leq i \leq k$ ) used in Gaussian elimination process from start to finish  $k$  number of times.

We need to do this as the augmented matrix  $[A|b_i]$  ( $1 \leq i \leq k$ ) used in Gaussian elimination is depended on  $b_i$  for being reduced to the row-echelon form.  $\therefore$  Our total number of operation =  $K \left( \frac{n(n+1)}{2} + n \right)$

Case 2: Calculating the number of operations performed when solve  $(A, b)$  is called  $k$  times using the Gauss-Jordan elimination method to solve  $Ax=b$  under the given constraint.

We know that, since the matrix  $A$  is invertible, we can solve the linear system of equations  $Ax=b$  as  $x = A^{-1}b$

We also know that we can find  $A^{-1}$  using the Gauss-Jordan elimination method. To do the same, we have to start with the augmented matrix,  $[A|I_n]$  and by using elementary operations reduce it to the reduced-row-echelon form to get  $[I_n | A^{-1}]$ .

A matrix is said to be in reduced-row echelon form when:

- \* if it is already in row-echelon.
- \* Every pivot is 1.
- \* The pivot is the only non-zero entry in the column.

The steps for Gaussian-Jordan elimination are:

1. Write the augmented matrix as  $[A|I_n]$ .
2. Get a 1 in the  $i^{\text{th}}$  row of the  $i^{\text{th}}$  column.
3. Use row  $i$  to get 0's in the  $i^{\text{th}}$  column of rows ( $1 \text{ to } i-1$ ) and ( $i+1 \text{ to } n$ ).
  - a) Repeat steps 2 & 3 from  $(i+1) \text{ to } n$
- 5) When augmented matrix becomes  $[I_n | A^{-1}]$  get  $A^{-1}$ .

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For example, Augmented Matrix.

$$\left[ \begin{array}{cccc|ccc} * & * & \dots & * & 1 & 0 & \dots & 0 \\ * & * & \dots & * & 0 & 1 & \dots & 0 \\ \vdots & & & & 0 & 0 & \dots & 0 \\ * & * & \dots & * & 0 & 0 & \dots & 1 \end{array} \right]$$

$$\left[ \begin{array}{cccc|ccc} 1 & * & \dots & * & 1 & 0 & \dots & 0 \\ 0 & * & \dots & * & * & * & \dots & * \\ \vdots & & & & * & * & \dots & * \\ 0 & * & \dots & * & * & * & \dots & * \end{array} \right]$$

get 1 in the 1<sup>st</sup> row of the 1<sup>st</sup> column  
use row 1 to get 0's in the 1<sup>st</sup> column  
of rows 2 to n  
Total operation performed : n.

Finally doing all this information once we have  $A^{-1}$  we can compute  $A^{-1}$  by performing n operation to get each row of the matrix.

Therefore, we can conclude that to solve  $Ax = b$  using Gauss Jordan elimination we have to perform:

$$\Rightarrow (n + n + \dots + n) \text{ operations to find } A^{-1} + (n \text{ operations to compute } A^{-1}b)$$
$$\Rightarrow n^2 + n \text{ operations.}$$

Now, since we are intending to solve  $Ax = b$  k times where the right hand side changes, but the left hand side stays the same, i.e., matrix A remains constant & matrix b changes, we can:

- \* Calculate  $A^{-1}$  just once using the Gauss Jordan elimination method & store the value of  $A^{-1}$  for subsequent (k-1) calls.
- \* Calculate  $A^{-1}b_i$  ( $1 \leq i \leq k$ ) k times to get values of  $x_i$ .

Therefore, our Total number of operations would be :

$$n^2 + kn \text{ operations.}$$

Thus, in conclusion, we can see that CASE 2 requires a significantly lesser number of operations to solve  $Ax = b$ .

Answer

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Q3 Consider  $n \times n$  elementary matrices where  $E_{ij}$  represents the elementary matrix where there is a non-zero value at the  $i^{\text{th}}$  row &  $j^{\text{th}}$  column in addition to is diagonal. Given a particular elementary matrix  $E_{ij}$ , for which other elementary matrices  $E_{pq}$  is it the case that

$$E_{pq} E_{ij} = E_{ij} E_{pq} ?$$

Answer →

We know that an elementary matrix is a matrix that has:

\* Is on the diagonal and

\* a non-zero element at the  $i^{\text{th}}$  row and  $j^{\text{th}}$  column.

For example:

$$\begin{matrix} c_1 & c_2 & c_3 & c_4 & c_5 \\ \begin{matrix} R_1 \\ R_2 \\ R_3 \\ R_4 \\ R_5 \end{matrix} & \left[ \begin{matrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & \alpha & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{matrix} \right] & 5 \times 5 \\ E_{42} & & \end{matrix}$$

$$\left[ \begin{matrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \\ a_{41} & a_{42} & a_{43} \\ a_{51} & a_{52} & a_{53} \end{matrix} \right] \quad A \quad 5 \times 3$$

$$E_{42} \times A$$

$$\left[ \begin{matrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \\ a_{41} + a_{42}, & a_{42} + a_{41}, & a_{43} + a_{42}, \\ a_{51} & a_{52} & a_{53} \end{matrix} \right]$$

Matrix multiplication is possible when no. of columns in the first matrix is the same as the number of rows in the second matrix.

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Let us consider  $E_{ij}$  as an Elementary matrix, such that:

- \* it has non-zero element at position  $\alpha_{ij}$  and
- \* the 1 in the  $i^{\text{th}}$  row is in the  $j^{\text{th}}$  column

Now when  $E_{ij}$  is multiplied with any  $A_{n \times m}$  matrix, the following happens:

- \* the  $j^{\text{th}}$  row of  $A$  is scaled by non-zero value  $\alpha$ ,
- \* the scaled row is added to the  $i^{\text{th}}$  row of  $A$ .
- \* all other elements in  $A$  remain unchanged.

Or we can see in other words, it is the same as taking the  $i^{\text{th}}$  row of  $A$  and adding it to the  $\alpha$  times  $j^{\text{th}}$  row of  $A$ :

$$A : R_i = R_i + \alpha R_j$$

By this logic,

if elementary matrix  $E_{pq}$  having a non-zero value  $\alpha_{pq}$  is multiplied with another elementary matrix  $E_{ij}$  having a non-zero value  $\beta_{ij}$  then

if  $p = i$  and  $q = j$

$$E_{pq} = E_{24}$$

$$\Rightarrow E_{ij} = E_{24}$$

$$\Rightarrow E_{pq} \times E_{ij} = E_{24} \times E_{24} \quad (\text{same as rewriting } 2^{\text{nd}} \text{ row of } E_{ij} \text{ as } R_{p=2} + \alpha R_{q=4})$$

$$\Rightarrow E_{ij} E_{pq} = E_{24} \times E_{24}$$

(same as re-writing 2<sup>nd</sup> row of  $E_{pq}$  as  $R_{i=2} + \beta R_{j=4}$ )

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Example

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & \alpha \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & \beta \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & \beta + \alpha \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & \alpha + \beta \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

\* if  $p = i$  &  $q \neq j$

$$E_{pq} = E_{24}$$

$$E_{ij} = E_{23}$$

$E_{pq} E_{ij} = E_{24} E_{(23)}$  (Same as rewriting 2<sup>nd</sup> row of  $E_{ij}$  as  $R_p = R_2 + \alpha R_q$ )

$E_{ij} E_{pq} = E_{23} E_{24}$  (Same as rewriting 2<sup>nd</sup> row of  $E_{pq}$  as  $R_i = R_2 + \beta R_j$ )

Example

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & \alpha \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & \alpha \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & \alpha \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

if  $p \neq i$  &  $q = j$

$$E_{pq} = E_{43}$$

$$E_{ij} = E_{23}$$

$E_{pq} E_{ij} = E_{43} E_{23}$  (Same as rewriting 4<sup>th</sup> row of  $E_{ij}$  as  $R_p = R_4 + \alpha R_q$ )

$E_{ij} E_{pq} = E_{23} E_{43}$  (Same as rewriting 2<sup>nd</sup> row of  $E_{pq}$  as  $R_i = R_2 + \beta R_j$ )

Example

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & \alpha & 0 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

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if  $p \neq i$  and  $q \neq j$

$$E_{pq} = E_{12}$$

$$E_{ij} = E_{23}$$

$$E_{pq} E_{ij} = E_{12} E_{23} \quad (\text{same as re-writing 1st row of } E_{ij} \text{ as } R_{p=1} + R_{q=2})$$

$$E_{ij} E_{pq} = E_{23} E_{12} \quad (\text{same as re-writing 2nd row of } E_{pq} \text{ as } R_{i=2} + R_{j=3})$$

Example:

$$\begin{bmatrix} 1 & \alpha & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & \alpha & \alpha\beta & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & \alpha & 0 & 0 \\ 0 & 1 & \beta & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

Therefore, we can conclude that

$$E_{pq} E_{ij} = E_{ij} E_{pq} \text{ when}$$

$p = i$  and  $q = j$ , or

$p = i$  and  $q \neq j$ , or

$p \neq i$  and  $q = j$

\*\*\* — End — \*\*\*

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Q-4 As a part of computer application, a sub-routine needs to be written whose input parameter  $p$  has to be used in the computation  $A^p$  where  $A$  is  $100 \times 100$  symmetric, positive definite matrix  $A$ . Note that  $A$  is a fixed matrix and it is only  $p$  which is the input parameter. What is the most efficient way you can come up with to perform the required computation, if the sub-routine is called millions of times for arbitrary value of  $p$ ? Your solution needs to be efficient in terms of both time & space taken by the algorithm?

Answers:

There are multiple ways to perform this tasks such as:

- \* Eigen decomposition
- \* Spectral theorem
- \* Conjugate Gradient Method

We will perform the required computation efficiently with conjugate gradient Method.

The conjugate gradient method is an iterative algorithm that can be used to solve systems of linear equations, such as

$Ax = b$  where  $A$  is a symmetric positive definite matrix.

The conjugate gradient method has a time complexity of  $O(n^2)$  where  $A$  is a symmetric matrix and  $n$  is the size of the matrix and it requires only  $O(n)$  space, making it very efficient solution for this problem.

To use the conjugate gradient matrix method, we will first need to initialize a vector  $x$  with an initial guess for the solution

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and set a residual vector  $r = b - Ax$ , Then, you can iterate the following steps until the residual vector  $r$  is sufficiently small:

steps

- 1 choose a search direction  $p$  based on the residual vector  $r$ .
- 2 perform a line search along the direction  $p$  to find the step size alpha that minimizes the residual vectors.
- 3 update the solution vector  $x$  by  $x = x + \alpha p$
- 4 Update the residual vector  $r$  by  $r = r - \alpha Ap$ .

The conjugate gradient method can be improved & implemented in a subroutine that takes the input vector  $p$  as an argument and returns the solution vector  $x$ . This subroutine can be called multiple times with different values of  $p$  to perform the required computation efficiently.

Now using the Diagonal matrix  $D$  we can easily find

$$Ap = P D_p P^T$$

In order to find the matrix  $D$ , we have to find the  $\lambda_1, \dots, \lambda_n$  eigenvalues of matrix  $A$  by computing the determinant of Matrix  $(A - \lambda I)$ .

Then we can write as

$$\begin{bmatrix} \lambda_1 & 0 & 0 & 0 & 0 \\ 0 & \ddots & & & \\ 0 & 0 & \ddots & & 0 \\ 0 & 0 & 0 & \ddots & 0 \\ 0 & 0 & 0 & 0 & \lambda_n \end{bmatrix}$$

Once the eigen values are found, we can find matrix  $P$  by finding the corresponding  $p_1, \dots, p_n$  eigenvectors of matrix  $A$ . This can be done by solving the linear system of equations

$$(A - \lambda_i I) p_i = 0 \text{ where } n \leq i \leq 1$$

Q5 We need to send a  $1000 \times 1000$  matrix of numbers across a channel, and would like to minimize the total amount of data sent on the channel for reasons having to do with both the possibility of data corruption and the time taken to send the data. Can you think of a way of minimizing the amount of data to be sent across the channel, so that we can represent the most important information in the matrix?

We know that any matrix  $A \in \mathbb{R}^{m \times n}$  of rank  $r \in [0, \min(m, n)]$  can always be factored into a Singular Value Decomposition (SVD) as follows:

$$\begin{matrix} n \\ A \\ \left\{ \right. \end{matrix} = \begin{matrix} m \\ U \\ \left\{ \right. \end{matrix} + \begin{matrix} n \\ \Sigma \\ \left\{ \right. \end{matrix} + \begin{matrix} m \\ V^T \\ \left\{ \right. \end{matrix}$$

$$U = \begin{bmatrix} \vdots & \vdots & \vdots & \vdots \\ u_1 & u_2 & \cdots & u_m \end{bmatrix}$$

$$\Sigma = \begin{bmatrix} \sigma_1 & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & \sigma_m \\ 0 & \cdots & 0 \\ \vdots & & \vdots \\ 0 & \cdots & 0 \end{bmatrix}$$

$$V^T = \begin{bmatrix} \cdots & v_1^T & \cdots \\ \cdots & v_2^T & \cdots \\ \vdots & \vdots & \vdots \\ \cdots & v_n^T & \cdots \end{bmatrix}$$

- Where,
- ①  $U$  is a  $m \times m$  orthogonal matrix of column vectors  $u_i (1 \leq i \leq m)$ . Where  $U$  essentially contains information about the column space of  $A$  and is also called the left-singular vectors.
  - ②  $V$  is a  $n \times n$  orthogonal matrix of column vectors  $v_i (1 \leq i \leq n)$ . Where  $V$  essentially contains information about the row space of  $A$  and is also called the right-singular vectors.
  - ③  $\Sigma$  is a  $m \times n$  matrix with  $\Sigma_{ii} = \sigma_i \geq 0$  and  $\Sigma_{ij} = 0, i \neq j$ . Where  $\Sigma$  essentially contains information about how important the columns of  $U$  &  $V$  are and is also called the singular values.

The diagonal entries of  $\Sigma$  are ordered as  $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_m \geq 0$ .

$\Sigma$  has the same size as  $A$ . Meaning :

- ④ If  $m > n$  then  $\Sigma$  has diagonal structure upto row  $n$  & then consists of  $0_s$  from  $n+1$  to  $m$ .
- ⑤ If  $m < n$  then  $\Sigma$  has diagonal structure upto column  $m$  & then consists of  $0_s$  from  $m+1$  to  $n$ .

- \* The columns of  $U$  are hierarchically arranged such that column  $u_1$  is more important than  $u_2$  and so on.

The rows of  $V$  are hierarchically arranged such that row  $v_1$  is more important than  $v_2$  and so on.

And their importance is encoded in the singular values  $\sigma$ .

We know the computing the full SVD of a large  $m \times n$  matrix can be quite taking. So, instead we will now demonstrate how SVD allows us to represent matrix  $A$  as a sum of simpler (low-rank) matrices  $A_i$ , which leads itself to a matrix approximation scheme that is cheaper than the full SVD.

We will now try to demonstrate the full SVD as a sum of rank 1 matrices  $A_i$ :

$$A_{m \times n} = U \Sigma V^T = \begin{bmatrix} \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ u_1 & u_2 & \dots & u_n & \dots & u_m \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{bmatrix} \begin{bmatrix} \sigma_1 & & & & & \\ & \ddots & & & & \\ & & \sigma_n & & & \\ & & & 0 & \dots & 0 \\ & & & & \vdots & \\ & & & & 0 & \dots & 0 \end{bmatrix} \begin{bmatrix} \dots & v_1^T & \dots \\ \dots & v_2^T & \dots \\ \vdots & \vdots & \vdots \\ \dots & v_n^T & \dots \end{bmatrix}$$

$$= \sigma_1 u_1 v_1^T + \sigma_2 u_2 v_2^T + \dots + \sigma_n u_n v_n^T + 0$$

$$= \hat{U} \hat{\Sigma} \hat{V}^T$$

- \* Since  $\Sigma$  is a diagonal matrix when we multiply  $U \Sigma$  column  $u_1$  essentially gets scaled by  $\sigma_1$ , column  $u_2$  by  $\sigma_2$  and so on; and similarly, when we multiply  $U \Sigma V^T$  the first column  $\sigma_1 u_1$  only multiplies the  $v_1^T$  column,  $\sigma_2 u_2$  column only multiplies the  $v_2^T$  column and so on.
- \* Even though the  $V$  matrix has  $n$  columns, there are only  $m$  non-zero singular values in the  $\Sigma$  matrix. So everything after the first  $m$  columns in  $V$  &  $V^T$  becomes 0.

Essentially what this means is that we can select just the  $n$  columns of  $U$  i.e.  $\hat{U}$ , the first  $m \times n$  block in  $\Sigma$  i.e.  $\hat{\Sigma}$  and the  $m \times n$  matrix  $V^T$  and write that as  $\hat{U} \hat{\Sigma} \hat{V}^T$  and that is exactly the same as  $A_{m \times n}$ .

Now that we have represented matrix  $A$  as a sum of rank 1 matrices  $A_i$ . We can intuitively see that —

- (\*) The best rank 1 approximation of  $A$  is  $\sigma_1 u_1 v_1^T$
- (\*) The best rank 2 approximation of  $A$  is  $\sigma_1 u_1 v_1^T + \sigma_2 u_2 v_2^T$   
and so on (and this is what SVD essentially means)

What we will now do is truncate out approximation of matrix  $A$  at rank  $K$ .

What this means is if we have a lot of small singular values  $\sigma_i$  ( $i+1 \geq i \geq n$ ) are negligibly small and most of the information about matrix  $A$  is captured in the  $K$  singular values and singular vectors. We can keep  $\sigma_1 u_1 v_1^T + \sigma_2 u_2 v_2^T + \dots + \sigma_K u_K v_K^T$  i.e. the first  $K$  columns of  $U$  and  $V^T$ , the first  $K \times K$  submatrix of  $\Sigma$  and ignore the rest.

$$A_{m \times n} = U \Sigma V^T = \hat{U} \hat{\Sigma} \hat{V}^T \quad \text{where } \hat{U} \text{ has } K \text{ columns}$$

A formal definition of this type of approximation of  $A$  can be found in the Eckart-Young Theorem.

Where a matrix  $A \in \mathbb{R}^{m \times n}$  of rank  $r$  and a matrix  $B \in \mathbb{R}^{m \times n}$  of rank  $K$  for any  $K \leq r$  with  $\hat{A}(K) = \sum_{i=1}^K \sigma_i u_i v_i^T$  it holds that

$$\hat{A}(K) = \underset{\text{rank}(B)=K}{\text{argmin}} \|A - B\|_F, \|A - \hat{A}(K)\|_F = \|\hat{U} \hat{\Sigma} \hat{V}^T\|_F$$

The Eckart-Young theorem implies that we can use SVD to reduce a rank  $r$  matrix  $A$  to a rank  $K$  matrix  $\hat{A}$  in a principled, optimal (in the spectral norm sense) manner.

Therefore, in conclusion, by using the Eckart-Young Theorem we can approximate the most important information of matrix  $A$  by a rank  $K$  matrix (ie the first  $K$  columns of  $U$  and  $V$ , the first  $K \times K$  submatrix of  $\Sigma$ ) as a form of lossy compression.

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- 6) Consider a linear system  $Ax = b$ . Assume that column vectors  $a_1, \dots, a_n \in \mathbb{R}^n$  are columns of matrix  $A$  i.e.  $A = [a_1 \ a_2 \ a_3 \ \dots \ a_n]$ . Let  $C = [A|b]$  be the augmented matrix associated with this linear system. Let us consider 2 different scenarios
- Suppose I interchange column  $a_i$  and  $a_k$  of the augmented matrix  $C$  giving a new matrix  $C_1$  ( $i, k \leq n$ ). Now I solve the problem assuming  $C_1$  is my augmented matrix. How are the solutions of augmented matrix  $C$  and  $C_1$  related.
  - Suppose I scale the  $i^{th}$  column of the augmented matrix  $C$  by  $a$  giving a new matrix  $C_2$  ( $i, k \leq n$ ). Now I solve the problem assuming  $C_2$  is my augmented matrix. How are the solutions of augmented matrix  $C$  and  $C_2$  related.

Ans) Since,  $Ax = b$  is a linear system of equations, where  $[a_1, a_2, a_3, \dots, a_n]$  are the column vectors. And  $C = [A|b]$  is its augmented matrix.

The general form  $Ax = b$  and the augmented matrix  $C = [A|b]$  would look like:

General Form:

$$\Delta_{11}x_1 + \Delta_{12}x_2 + \dots + \Delta_{1i}x_i + \dots + \Delta_{1k}x_k + \dots + \Delta_{1n}x_n = b_1$$

$$\Delta_{21}x_1 + \Delta_{22}x_2 + \dots + \Delta_{2i}x_i + \dots + \Delta_{2k}x_k + \dots + \Delta_{2n}x_n = b_2$$

...

...

$$\Delta_{n1}x_1 + \Delta_{n2}x_2 + \dots + \Delta_{ni}x_i + \dots + \Delta_{nk}x_k + \dots + \Delta_{nn}x_n = b_n$$

Augmented Matrix :

$$\left[ \begin{array}{cccc|c} \Delta_{11} & \Delta_{12} & \dots & \Delta_{1i} & \dots & \Delta_{1k} & \dots & \Delta_{1n} & | & b_1 \\ \Delta_{21} & \Delta_{22} & \dots & \Delta_{2i} & \dots & \Delta_{2k} & \dots & \Delta_{2n} & | & b_2 \\ \ddots & | & \vdots \\ \ddots & | & \vdots \\ \Delta_{ni} & \Delta_{n2} & \dots & \Delta_{ni} & \dots & \Delta_{nk} & \dots & \Delta_{nn} & | & b_n \end{array} \right]$$

Note: Where the column vector  $a_i = [\Delta_{1i}, \Delta_{2i}, \dots, \Delta_{ni}]^T$

### Scenario A

Now, if we interchange the columns  $a_i$  and  $a_{ik}$  of the augmented matrix  $C$  giving a new matrix  $C_i$ .

What we have essentially done is re-written our equations as follows :

### General Form :

$$\Delta_{11}x_1 + \Delta_{12}x_2 + \dots + \Delta_{1k}x_i + \dots + \Delta_{1i}x_k + \dots + \Delta_{1n}x_n = b_1$$

$$\Delta_{21}x_1 + \Delta_{22}x_2 + \dots + \Delta_{2k}x_i + \dots + \Delta_{2i}x_k + \dots + \Delta_{2n}x_n = b_2$$

$\ddots$

$$\Delta_{ni}x_1 + \Delta_{n2}x_2 + \dots + \Delta_{nk}x_i + \dots + \Delta_{ni}x_k + \dots + \Delta_{nn}x_n = b_n$$

Augmented Matrix:

$$\left[ \begin{array}{cccc|c} \Delta_{11} & \Delta_{12} & \dots & \Delta_{1k} & \dots & \Delta_{1i} & \dots & \Delta_{1n} & b_1 \\ \Delta_{21} & \Delta_{22} & \dots & \Delta_{2k} & \dots & \Delta_{2i} & \dots & \Delta_{2n} & b_2 \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots & \vdots \\ \Delta_{n1} & \Delta_{n2} & \dots & \Delta_{nk} & \dots & \Delta_{ni} & \dots & \Delta_{nn} & b_n \end{array} \right]$$

Therefore, we can conclude that :

- Only the  $i^{\text{th}}$  and  $k^{\text{th}}$  entries of the solution  $C$  will be swapped in the solution  $C_1$ .
- All other entries will be the same.

### Scenario B

Now, if we scale the  $i^{\text{th}}$  column of the augmented matrix  $C$  by  $\alpha$  giving a new matrix  $C_2$

What we have essentially done is re-written our equations as follows :

General Form:

$$\Delta_{11}x_1 + \Delta_{12}x_2 + \dots + \alpha\Delta_{1i}x_i + \dots + \Delta_{1n}x_n = b_1$$

$$\Delta_{21}x_1 + \Delta_{22}x_2 + \dots + \alpha\Delta_{2i}x_i + \dots + \Delta_{2n}x_n = b_2$$

....

$$\Delta_{n1}x_1 + \Delta_{n2}x_2 + \dots + \alpha\Delta_{ni}x_i + \dots + \Delta_{nn}x_n = b_n$$

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Augmented Matrix:

$$\left[ \begin{array}{cccc|c} \Delta_{11} & \Delta_{12} & \dots & \alpha \Delta_{1i} & \dots & \Delta_{1n} & b_1 \\ \Delta_{21} & \Delta_{22} & \dots & \alpha \Delta_{2i} & \dots & \Delta_{2n} & b_2 \\ \vdots & \vdots & & \vdots & & \vdots & \vdots \\ \Delta_{n1} & \Delta_{n2} & & \alpha \Delta_{ni} & & \Delta_{nn} & b_n \end{array} \right]$$

Therefore, we can conclude that :

- Only the  $i^{th}$  entry of the solution  $C$  will be scaled by  $\frac{1}{\alpha}$  times of the solution of  $C_2$ .
- All the other entries will be same.

7) In a Class, a professor informed students that  $M$  is a real 3 by 3 real matrix such that  $M^3 = I$ . Using the given information, students were asked whether the matrix  $M$  is invertible and to find the eigenvalues of  $M$ . Is  $M$  invertible and find the eigenvalues of  $M$ .

Ans) We know that if a matrix  $A \in R^{n \times n}$  is invertible, then  $A^{-1}A = I_n = AA^{-1}$

Now, since

$$\begin{aligned} M^3 &= I_3 \\ \Rightarrow M^2M &= I_3 \quad [\text{Therefore, the matrix } M \text{ is invertible and}] \\ \Rightarrow M^{-1}M &= I_3 \quad [M^2 = M^{-1}] \quad [\text{Where } M^2 = M^{-1}] \end{aligned}$$

We will next try to find the eigen values of matrix  $M$

We know that a square matrix is invertible iff it does not have a zero eigen value.

$\therefore$  the matrix  $M$  will not have 0 as an eigen value.

Next, let  $\lambda$  and  $x$  be the eigenvalue and eigen vectors of matrix  $M$ , such that :

$$\begin{aligned} Mx &= \lambda x \quad \dots \text{①} \\ \Rightarrow M^2Mx &= M^2\lambda x \quad (\text{Multiplying } M^2 \text{ on both sides}) \\ \Rightarrow M^3x &= \lambda MMx \\ \Rightarrow M^3x &= \lambda^2 Mx \quad (\text{from ①}) \\ \Rightarrow M^3x &= \lambda^3 x \quad (\text{from ①}) \\ \Rightarrow M^3x &= \lambda^3 x \end{aligned}$$

$$\Rightarrow I_3 x = \lambda^3 x \quad [\text{since } M^3 = I_3]$$

$$\Rightarrow x = \lambda^3 x$$

$$\Rightarrow \lambda^3 x - x = 0$$

$$\Rightarrow (\lambda^3 - 1)x = 0$$

$$\Rightarrow (\lambda - 1)(\lambda^2 + \lambda + 1)x = 0$$

∴ the eigenvalues  $\lambda$  of matrix  $M$  are

$$1, \frac{-1 + i\sqrt{3}}{2} \quad \text{and} \quad \frac{-1 - i\sqrt{3}}{2}$$

8) A student from Linear Algebra class received a matrix of the form given below.

$$A = \begin{bmatrix} 0 & \Delta & 1 & 0 \\ \Delta & 0 & \Delta & 0 \\ 1 & \Delta & 0 & \Delta \\ 0 & 0 & \Delta & 1 \end{bmatrix}, \Delta > 1$$

Help him evaluate the column rank, trace and the determinant and then using trace, rank and determinant, what conclusion can be drawn about the signs (+ or -) of eigenvalues of A.

Ans) From the above problem statement, we can see that A is  $4 \times 4$  square matrix.

We know that the trace of a matrix  $A \in \mathbb{R}^{n \times n}$  is the sum of the diagonals of A.

Therefore

$$\text{tr}(A) = \sum_{i=1}^4 a_{ii} = (0+0+0+1) = 1$$

We will next find the determinant of matrix A:

$$\begin{aligned} \det(A) &= 0 - \Delta \begin{vmatrix} \Delta & \Delta & 0 \\ 1 & 0 & \Delta \\ 0 & \Delta & 1 \end{vmatrix} + 1 \begin{vmatrix} \Delta & 0 & 0 \\ 1 & \Delta & \Delta \\ 0 & 0 & 1 \end{vmatrix} - 0 \\ &= -\Delta (\Delta(0 \cdot 1 - \Delta \cdot \Delta) - \Delta(1 \cdot 1 - 0 \cdot \Delta)) + (\Delta(\Delta \cdot 1 - 0 \cdot \Delta)) \end{aligned}$$

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$$\begin{aligned} &= -\Delta(-\Delta^3 - \Delta) + \Delta^2 \\ &= \Delta^4 + \Delta^2 + \Delta^2 \\ &= \Delta^4 + 2\Delta^2 \end{aligned}$$

Therefore

$$\det(A) = \Delta^4 + 2\Delta^2$$

We will next find the rank of a matrix A:

Since,  $\Delta > 1$  then  $\det(A) \neq 0$

Now we know that a  $A \in R^{n \times n}$  matrix has rank  $n$  if and only if its determinant is not equal to zero.

Therefore

$$rk(A) = 4$$

Now using the trace, rank and determinant we will try to draw conclusion about the signs of the eigenvalues of A:

We know that,

$\lambda$  is an eigenvalue of a matrix  $A \in R^{n \times n}$  if and only if  $\lambda$  is a root of the characteristic polynomial  $p_A(\lambda)$  of degree  $n$ .

Since, A is  $4 \times 4$  matrix then A must have 4 eigenvalues  $\lambda_1, \lambda_2, \lambda_3$  and  $\lambda_4$ .

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We also know that, the determinant of a matrix  $A \in \mathbb{R}^{n \times n}$  is the product of eigenvalues.

Therefore,

$$\det(A) = \prod_{i=1}^4 \lambda_i = \lambda_1 \cdot \lambda_2 \cdot \lambda_3 \cdot \lambda_4 = \Delta^4 + 2\Delta^2 \neq 0$$

Therefore, we can conclude that either all 4 eigenvalues will have the same sign or at least two of them will have different signs.

Q.9) Harry is the team lead in a company but is new to Linear Algebra. While working on his project, he arrived at a problem. He got three vectors  $v_1, v_2, v_3 \in \mathbb{R}^5$

Let  $S = \text{span}\{v_1, v_2, v_3\}$  and  $W = \{u_1, u_2, \dots, u_r\}$

The problem is to find all those  $u_i$ 's that belong to  $S$  and if  $u_i \in S$ , find linear combination of  $u_i$  in terms of  $v_1, v_2, v_3$ . Explain the method to solve Harry's problem with proper justification. Is the method efficient?

Using the above method solve it

$$v_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \\ 0 \end{bmatrix}, v_2 = \begin{bmatrix} 1 \\ 2 \\ 0 \\ 0 \\ 1 \end{bmatrix}, v_3 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \text{ and } W = \left\{ \begin{bmatrix} 3 \\ 3 \\ 3 \\ 0 \\ 2 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 2 \\ 0 \\ -1 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} \right\}$$

→

Answer :-

In the above problem statement, we see that  $S$  is the span of vector  $v_1, v_2, v_3$  in space  $\mathbb{R}^5$

This means that  $S$  is the set of all vectors that can be derived from all possible linear combinations of vectors  $v_1, v_2, v_3$  for all choices of scalars  $\lambda_1, \lambda_2, \lambda_3$  such that

Any element in set  $S = \lambda_1 v_1 + \lambda_2 v_2 + \lambda_3 v_3$

We also set that  $W$  is a set of vectors  $\{u_1, u_2, \dots, u_r\}$ . Now, to determine if a vector  $u_i \in S$ , we can write a linear system of equations as follows:

$$u_i = \lambda_1 v_1 + \lambda_2 v_2 + \lambda_3 v_3$$

Since matrixes can be used to compactly represent a linear system of equations  $Ax=b$ , we can re-write the above equation as follows :

$$\begin{bmatrix} v_1 & v_2 & v_3 \end{bmatrix} \begin{bmatrix} \lambda_1 \\ \lambda_2 \\ \lambda_3 \end{bmatrix} = u_i$$

$$Ax=b$$

$$\begin{bmatrix} v_1 & v_2 & v_3 & | & u_i \end{bmatrix}$$

Augmented Matrix  $[A|b]$

From here, we can use Gaussian elimination to reduce the augmented matrix to the row-echelon form using elementary transformations. And from there, we can perform back-substitution to get the values of  $\lambda_1, \lambda_2, \lambda_3$ .

And once we have the values of  $\lambda_1, \lambda_2, \lambda_3$ , we can check their correctness using :

$$u_i = \lambda_1 v_1 + \lambda_2 v_2 + \lambda_3 v_3$$

Thus in conclusion, this is the method Harry should follow to check if a vector  $u_i \in S$

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Now, using this method we will try to see if vectors in set WES, where:

$$v_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \\ 0 \end{bmatrix}, v_2 = \begin{bmatrix} 1 \\ 2 \\ 0 \\ 0 \\ 1 \end{bmatrix}, v_3 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \text{ and } W = \left\{ \begin{bmatrix} 3 \\ 3 \\ 3 \\ 0 \\ 2 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 2 \\ 0 \\ -1 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} \right\}$$

### CASE 1:

Given that,

$$v_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \\ 0 \end{bmatrix}, v_2 = \begin{bmatrix} 1 \\ 2 \\ 0 \\ 0 \\ 1 \end{bmatrix}, v_3 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \text{ and } u_1 = \begin{bmatrix} 3 \\ 3 \\ 3 \\ 0 \\ 2 \end{bmatrix}$$

Therefore,

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 1 & 2 & 0 & 3 \\ 2 & 0 & 0 & 3 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 2 \end{array} \right] \quad \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 1 & 2 & 0 & 3 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & 1 & 0 \\ 2 & 0 & 0 & 3 \end{array} \right]$$

swap R<sub>3</sub>, R<sub>5</sub>

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 0 & 1 & -1 & 0 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & 1 & 0 \\ 0 & -2 & -2 & -3 \end{array} \right] \quad \boxed{\text{R}_2 - \text{R}_1} \quad \boxed{\text{R}_5 - 2\text{R}_1}$$

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 0 & 1 & -1 & 0 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right] \quad \boxed{\text{R}_5 + 2\text{R}_3}$$

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If we now convert the row-echelon form of the augmented matrix to a linear system of equations we get :

$$\lambda_1 + \lambda_2 + \lambda_3 = 3 \dots (1)$$

$$\lambda_2 - \lambda_3 = 0 \dots (2)$$

$$\lambda_2 + \lambda_3 = 2 \dots (3)$$

$$\lambda_3 = 0 \dots (4)$$

$$0\lambda_1 + 0\lambda_2 + 0\lambda_3 = 1 \dots (5)$$

We do not even need to perform back-substitution here to get the values of  $\lambda_1, \lambda_2, \lambda_3$  as the last equation is a false statement. Hence,  $u_1 \notin S$  or in other words vector  $u_1$  cannot be represented as linear combination of vectors  $v_1, v_2, v_3$ .

### CASE 2 :-

Given that,

$$v_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \\ 0 \end{bmatrix}, v_2 = \begin{bmatrix} 1 \\ 2 \\ 0 \\ 0 \\ 1 \end{bmatrix}, v_3 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \text{ and } u_1 = \begin{bmatrix} 0 \\ 1 \\ 2 \\ 0 \\ -1 \end{bmatrix}$$

Therefore,

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 1 & 2 & 0 & 1 \\ 2 & 0 & 0 & 2 \\ 0 & 0 & 1 & -1 \\ 0 & 1 & 1 & -1 \end{array} \right] \xrightarrow{\text{swap R}_3, \text{R}_5} \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 1 & 2 & 0 & 1 \\ 0 & 0 & 1 & -1 \\ 0 & 0 & 1 & 0 \\ 2 & 0 & 0 & 2 \end{array} \right]$$

swap R<sub>3</sub>, R<sub>5</sub>

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 1 & 1 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & -2 & -2 & 2 \end{array} \right] \xrightarrow{\text{R}_2 - \text{R}_1, \text{R}_5 - 2\text{R}_1} \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 2 & 1 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right]$$

R<sub>2</sub> - R<sub>1</sub>  
R<sub>5</sub> - 2R<sub>1</sub>

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 2 & 1 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right] \xrightarrow{\text{R}_5 + 2\text{R}_3} \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 2 & 1 & -1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right]$$

R<sub>5</sub> + 2R<sub>3</sub>

If we now convert the row-echelon of the augmented matrix to a linear system of equations we get.

$$\lambda_1 + \lambda_2 + \lambda_3 = 0 \dots (1)$$

$$\lambda_2 + \lambda_3 = 1 \dots (2)$$

$$\lambda_2 + \lambda_3 = -1 \dots (3)$$

$$\lambda_3 = 0 \dots (4)$$

$$0\lambda_1 + 0\lambda_2 + 0\lambda_3 = 0 \dots (5)$$

From equation (4), we get  $\lambda_3 = 0$

From equation (3), we get  $\lambda_2 = -1$ , but this does not satisfy equation (2)

Hence,  $u_2 \notin S$  or vector  $u_2$  cannot be represented as linear combination of vectors  $v_1, v_2, v_3$ .

### CASE 3 :

Given that,

$$v_1 = \begin{bmatrix} 1 \\ 1 \\ 2 \\ 0 \\ 0 \end{bmatrix}, v_2 = \begin{bmatrix} 1 \\ 2 \\ 0 \\ 0 \\ 1 \end{bmatrix}, v_3 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 1 \\ 1 \end{bmatrix} \text{ and } u_3 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

Therefore,

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 1 & 2 & 0 & 1 \\ 2 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \end{array} \right], \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 1 & 2 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 2 & 0 & 0 & 1 \end{array} \right] \quad \boxed{\text{Swap R}_3, \text{R}_5}$$

$$\left[ \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & 1 & -1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & -2 & -2 & -1 \end{array} \right] \quad \boxed{\text{R}_2 - \text{R}_1} \quad , \quad \left[ \begin{array}{ccc|c} 1 & 1 & 1 & 1 \\ 0 & 1 & -1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{array} \right] \quad \boxed{\text{R}_5 + 2\text{R}_3}$$

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If we now convert the row-echelon form of the augmented matrix to a linear system of equations we get :

$$\lambda_1 + \lambda_2 + \lambda_3 = 1 \quad \dots (1)$$

$$\lambda_2 - \lambda_3 = 0 \quad \dots (2)$$

$$\lambda_2 + \lambda_3 = 1 \quad \dots (3)$$

$$\lambda_3 = 1 \quad \dots (4)$$

$$0\lambda_1 + 0\lambda_2 + 0\lambda_3 = 1 \quad \dots (5)$$

We do not even need to perform back-substitution here to get the values of  $\lambda_1, \lambda_2, \lambda_3$  as the last equation is a false statement.

Hence,  $u_3$  or vector  $u_3$  cannot be represented as linear combination of vectors  $v_1, v_2, v_3$ .

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Q.10) Consider system of equations in the matrix form as  $Ax=b$ , where A is a matrix of order  $m \times n$

Case 1.  
If  $m=5, n=6$ , then will the system have solutions for every choice of  $b$ ? Discuss with explanation.

Case 2.  
If  $m=6, n=8$  and rank of A is 6, then is it possible to make the system have no solution by changing  $b$ ? Discuss with explanation.

Case 3.  
If  $m=10, n=12, b_i=0$  for all  $i$  from 1 to 12 then is it possible that all solutions are multiples of one fixed non-zero solution? Discuss with explanation.

Answer →

We know that the rank of a matrix A can be calculated as follows:

If the columns of matrix  $A \in \mathbb{R}^{m \times n}$  span a vector subspace  $V \subseteq \mathbb{R}^m$  then  $\text{rk}(A) = \dim(V)$ .

where,

- $\text{rk}(A)$  is the rank of matrix A
- $\dim(V)$  is the number of pivot columns in matrix A found by reducing it to the row-echelon form using Gaussian elimination.

We can extend this concept to find the rank of an augmented matrix  $[A|b]$ , where  $A \in \mathbb{R}^{m \times n}$  and  $b \in \mathbb{R}^{m \times 1}$ , by considering the augmented matrix as a matrix of order  $m \times (n+1)$  and then by counting the number of pivot columns after reducing it to the row-echelon form.

We also know that a linear system of equations  $Ax=b$  is said to be consistent if  $b$  can be expressed as a linear combination of the pivot columns (i.e., linearly independent columns) of  $A$ . And since the rank of a matrix is nothing but the count of pivot columns. We can make some interesting observations:

Here,  $Ax=b$  is a linear system of equations and  $A$  is a  $3 \times 3$  matrix,  $x$  is a  $3 \times 1$  matrix and  $b$  is a  $3 \times 1$  matrix.

$$\text{rk}(A) \neq \text{rk}(A|b)$$

$$\left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ 0 & a_{22} & a_{23} & b_2 \\ 0 & 0 & 0 & b_3 \end{array} \right]$$

$$\text{rk}(A) = \text{rk}(A|b) = n$$

$$\left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ 0 & a_{22} & a_{23} & b_2 \\ 0 & 0 & a_{33} & b_3 \end{array} \right]$$

$$\text{rk}(A) < \text{rk}(A|b) < n$$

$$\left[ \begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ 0 & a_{22} & a_{23} & b_2 \\ 0 & 0 & 0 & 0 \end{array} \right]$$

- $\text{rk}(A)=2$  and  $\text{rk}(A|b)=3$

$\text{rk}(A)=3$ ,  $\text{rk}(A|b)=3$  and  $n=3$

- $\text{rk}(A)=2$ ,  $\text{rk}(A|b)=2$  and  $n=3$

- This means that  $b$  cannot be expressed as a linear combination of the pivot columns of  $A$ . And hence, no solution exists for  $Ax=b$ .

- We can also make the same argument just by looking at the equation formed by the last row of  $[A|b]$ . The equation is clearly a false statement and hence no solution exists for  $Ax=b$ .

This means that  $b$  can be expressed as a linear combination of the pivot columns of  $A$ .

And by back-substitution we can find values of  $x_i$  ( $1 \leq i \leq n$ ) that satisfy  $Ax=b$ . and if the values of  $x_i$  satisfy  $Ax=b$ .

- This means that  $b$  can be expressed as a linear combination of the pivot columns of  $A$ .

- But in this case we cannot use back-substitution we can find values of  $x_i$  ( $1 \leq i \leq n$ ) that satisfy  $Ax=b$ .

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we can say that the  
 $Ax=b$  has one  
unique solution.

- As we have  $n$  numbers of unknown in matrix  $x$  and  $>n$  number of equations.

Hence,  $Ax=b$   
will have infinite  
number of  
solutions.

Now using these observations, we will attempt to solve this problem.

### CASE 1:-

Given that,

- $Ax=b$  is a linear system of equations  $A$  is a  $5 \times 6$  matrix,  
 $x$  is a  $6 \times 1$  matrix and  $b$  is a  $5 \times 1$  matrix.

This means that rank of  $A$  can be:

$$rk(A)=5$$

$$\left[ \begin{array}{cccccc|c} a_{11} & a_{12} & a_{13} & a_{14} & a_{15} & a_{16} & b_1 \\ 0 & a_{22} & a_{23} & a_{24} & a_{25} & a_{26} & b_2 \\ 0 & 0 & a_{33} & a_{34} & a_{35} & a_{36} & b_3 \\ 0 & 0 & 0 & a_{44} & a_{45} & a_{46} & b_4 \\ 0 & 0 & 0 & 0 & a_{55} & a_{56} & b_5 \end{array} \right]$$

$$1 \leq rk(A) < 5$$

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- From the augmented beside, we can see that if  $\text{rk}(A)$  is 5 then  $\text{rk}(A|b)$  must also be 5.
- If  $\text{rk}(A)$  is 5 and  $\text{rk}(A|b)$  is also 5 and since  $n$  is 6, the linear system of equations will have infinite number of solutions.

- since  $n$  is 6, even if  $\text{rk}(A) = \text{rk}(A|b)$ , the linear system of equations will have infinite number of solutions.
- If  $\text{rk}(A) < 5$  and  $\text{rk}(A|b) = 5$ , then the linear system of equations will have 0 solutions.

Thus in conclusion, we see that  $Ax=b$ , where  $m=5$  &  $n=6$ , will not have solutions for all choices of  $b$  as there may be a case where  $\text{rk}(A) \neq \text{rk}(A|b)$ .

### CASE 2:

Given that,  
 $\bullet$   $Ax=b$  is a linear system of equations and  $A$  is a  $6 \times 8$  matrix,  $x$  is a  $8 \times 1$  matrix and  $b$  is a  $6 \times 1$  matrix.  
 $\bullet$  This means that we have  $n=8$  number of unknowns in matrix  $x$ .  
 $\bullet$   $\text{rk}(A)$  is 6

This means that:

$$\text{rk}(A)=6 \quad \left[ \begin{array}{ccccccc|c} a_{11} & & & & & & & b_1 \\ 0 & a_{22} & & & & & & b_2 \\ 0 & 0 & a_{33} & & & & & b_3 \\ 0 & 0 & 0 & a_{44} & & & & b_4 \\ 0 & 0 & 0 & 0 & a_{55} & & & b_5 \\ 0 & 0 & 0 & 0 & 0 & a_{66} & \dots & b_6 \\ & & & & & \dots & a_{68} & b_7 \end{array} \right]$$

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- From the augmented beside, we can see that if  $\text{rk}(A)$  is 6 then  $\text{rk}(A|b)$  must also be 6.
- If  $\text{rk}(A)$  is 6 and  $\text{rk}(A|b)$  is also 6 and since  $n$  is 8, the linear system of equations will have infinite number of solutions.

Thus in conclusion, we see that

- If  $b$  cannot be represented as a linear combination of the pivot columns of  $A$ , then the linear system of equations will have 0 solutions.
- Otherwise  $Ax=b$ , where  $m=6$ ,  $n=8$  and  $\text{rk}(A)=6$ , will always have infinite solutions for any given value of  $b$ .

### CASE 3:

Given that,

- $Ax=b$  is a linear system of equations and  $A$  is a  $10 \times 12$  matrix,  $x$  is a  $12 \times 1$  matrix and  $b$  is a  $10 \times 1$  matrix.
- This means that we have  $n=12$  numbers of unknowns in matrix  $x$ .
- All the elements of matrix  $b$  are zeros.  
 $b_i = 0 (1 \leq i \leq 10)$

This means that:

- The rank of matrix  $A$  can be at most 10 (i.e  $\text{rk}(A) \leq 10$ ).

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- Since all the elements of matrix  $b$  are zeros, the rank of  $A$  must always be equal to the rank of  $[A|b]$  (i.e.  $\text{rk}(A) = \text{rk}(A|b)$ ). Hence, there is no possibility of having 0 solutions for  $Ax=b$ .
- And since  $n$  is 12, the linear system of equations will have infinite number of solutions.

Thus in conclusion, we see that

- $Ax=b$ , where  $m=10$ ,  $n=12$  and  $b_i=0$  ( $1 \leq i \leq 12$ ), will have infinite solutions.
- We can hence, find a particular non-zero solution to  $Ax=b$  and represent all other solutions as multiples of it.