Binomial Heaps

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Comp 401: Senior Seminar

3/05/2014

- 1. Binomial Trees
- 2. Binomial Heaps
- 3. Implementation Structure
- 4. Standard Functions
- 5. Uses of Binomial Heaps

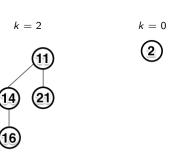
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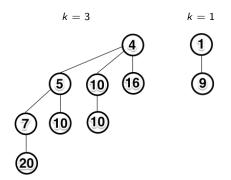
Binomial Trees

- A Binomial Tree is a specific type of tree that includes the following specifications:
 - 1. The order or rank of the binomial tree is the number of children of the root node.
 - 2. A Binomial Tree of order 0 is a single node.
 - 3. A Binomial Tree of order k has k child nodes, all of which are the roots of binomial trees of orders $k-1,\ k-2,\ ...,\ 2,\ 1,\ 0$ from left to right.

Binomial Trees: Examples

• If a binomial tree has order k, the orders of the k child nodes decrease from left to right from k-1 to 0.





- If a Binomial Tree has an order k:
 - 1. The tree has 2^k nodes.
 - 2. The height of the tree is k.
 - 3. There are $\binom{k}{d}$ nodes at depth d.
- $\binom{k}{d} = \frac{k!}{d!(k-d)!}$ is known as the Binomial Coefficient.

Example:
$$k = 3$$
, $d = 2$

5 10 16

7 10 10

$$\binom{k}{d} = \binom{3}{2} = \frac{3!}{2!(3-2)!} = \frac{6}{2*1} = \frac{6}{2} = 3$$

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- A Binomial Heap is a collection of binomial trees that satisfy the following two binomial heap properties:
 - 1. The key of any node is greater than or equal to the key of its parent (minimum-heap property).
 - 2. There cannot be two binomial trees of the same order.

Binomial Heap: Property #1 (minimum-heap)

- The first property (minimum-heap) ensures that the root is the smallest key in each binomial tree.
- The smallest key of the entire heap is one of the roots.

Min-Heap Property ✓

1 4 5 2 9 5 10 16 6 5 6 9 7 10 10 3 10 7

Min-Heap Property ×

Binomial Heap: Property #2

• The order of each binomial tree must be unique.

Property #2 × Property #2 √ Property #2 ×

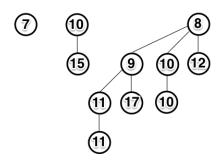
6 3 4 1 4 2 1 3

8 9 5 10 5 6 9 6 2 3 9

11 7 7 5

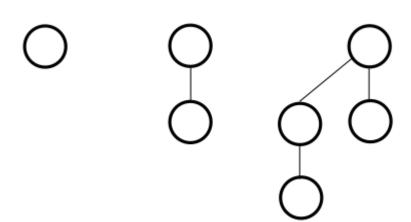
Binomial Heap: Property #2, Cont.

- The second property ensures that if a binomial heap has n nodes, then it will have at most ⌊log n⌋ + 1 binomial trees.
- The total number of nodes can also be thought of as a binary string, where each binomial tree represents a bit.



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Implementation Structure



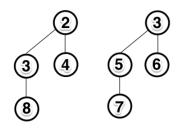
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Standard Functions: Merge

• Two k ordered binomial trees are merged into one k+1 ordered binomial tree.

Merge

```
Input: Node* root_A, Node* root_B
  if root_A.key ≤ root_B.key then
     root_B.right_sibling \leftarrow root_A.child
     root_A.child \leftarrow root_B
     root_B.parent \leftarrow root_A
     return root A
  else
     root_A.right_sibling \leftarrow root_B.child
     root B.child \leftarrow root A
     root\_A.parent \leftarrow root\_B
     return root B
  end if
```

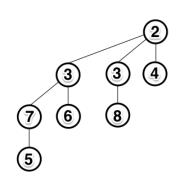


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  end if
```



Standard Functions: Join

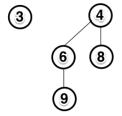
Joins one Binomial Heap into another Binomial Heap.

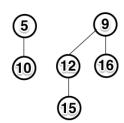
BinomialHeap: Join

Input: Binomialheap merge_heap

for each binomialtree in merge_heap do
 this_heap.insert(binomialtree.root)

end for





Standard Functions: Insert

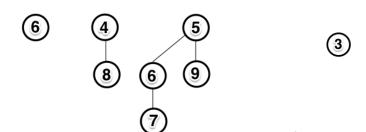
 Inserts either an integer or an existing node into a Binomial Heap.

BinomialHeap: Insert

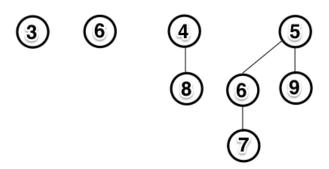
Input: int key

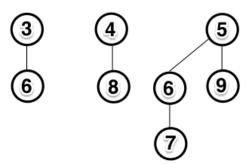
 $BinomialHeap\ temp_heap\ \leftarrow\ new\ BinomialHeap(key)$

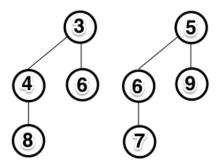
this_heap.join(temp_heap)



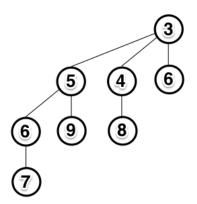
Standard Functions: Insert







Standard Functions: Insert



Standard Functions: DeleteMinimum

 Deletes the node with the smallest key from the Binomial Heap.

```
BinomialHeap : DeleteMinimum

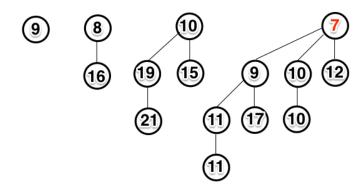
Input: None
Node* deleted_node ← this_heap.FindMinimum()
BinomialHeap temp_heap ← new BinomialHeap()
for each child in delete_node do
temp_heap.insert(child)
end for
this_heap.join(temp_heap)
```

Standard Functions: DeleteMinimum

BinomialHeap: DeleteMinimum

Input: None

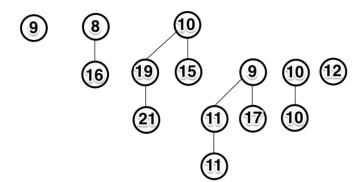
Node* deleted_node ← this_heap.FindMinimum()
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Standard Functions: DeleteMinimum

BinomialHeap: DeleteMinimum, cont.

for each child in delete_node do
 temp_heap.insert(child)
end for...



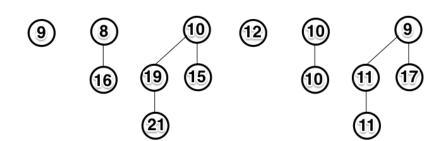
Standard Functions: DeleteMinimum, cont.

BinomialHeap: DeleteMinimum

this_heap.join(temp_heap)

this_heap

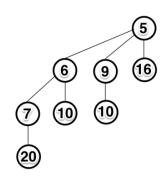
temp_heap



• Decreases the key of a given node in the Binomial Heap

```
BinomialHeap : DecreaseKey
```

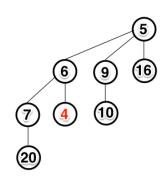
```
Input: Node* node, int new_key
  decrease_node.key ← new_key
  Node* temp ← decrease_node
  while temp.parent \neq NULL do
    if temp.key < temp.parent.key</pre>
    then
       swap(temp.key, temp.parent.key)
    else
       break
    end if
    temp \leftarrow temp.parent
  end while
```



Decreases the key of a given node in the Binomial Heap

```
BinomialHeap : DecreaseKey
```

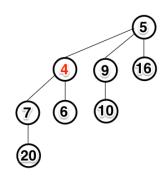
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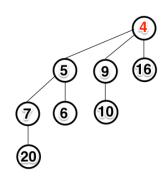
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BinomialHeap : DecreaseKey

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    else
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    end if
    temp \leftarrow temp.parent
  end while
```



Standard Functions: Delete

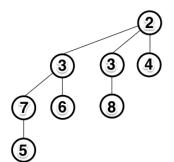
Deletes a given node from the Binomial Heap

BinomialHeap: Delete

Input: Node* delete_node

BinomialHeap.DecreaseKey(delete_node, -9999)

BinomialHeap.DeleteMinimum()



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Uses of Binomial Heaps

- Binomial Heaps are used to implement priority queues
- Priority queues can be used in applications that manage network bandwidth or manage events in a discrete even simulation.
- Priority queues can also be used in various algorithms including heapsort, Dijkstra's shortest-path algorithm, Prim's minimum spanning tree algorithm, and Huffman Encoding.

Operation	Linked List	Binary	d-ary	Binomial
FindMinimum	O(n)	O(1)	O(1)	O(log n)
Insert	O(log n)	O(log n)	$O(log_d n)$	O(log n)
Join	O(1)	O(n)	O(n)	O(log n)
DeleteMinimum	O(n)	O(log n)	$O(d \log_d n)$	O(log n)
DecreaseKey	O(1)	O(log n)	$O(log_d n)$	O(log n)
Delete	O(1)	O(log n)	$O(d \log_d n)$	O(log n)

References

- KEVIN WAYNE. Priority Queues. http://www.cs.princeton.edu/ wayne/kleinberg-tardos/pdf/BinomialHeaps.pdf, 2013.
- UNKNOWN.
 Binomial Heap.
 http://en.wikipedia.org/wiki/Binomial_heap, 2015.