

CSCI 2270

Data Structures and Algorithms

Lecture 17

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Office hours: ECCS 128

Wed 1-2pm

Fri 2-3pm

Administrivia

Thursday 2-3 office hours have moved to Friday 2-3

HW2 will post on Tuesday (part 1)

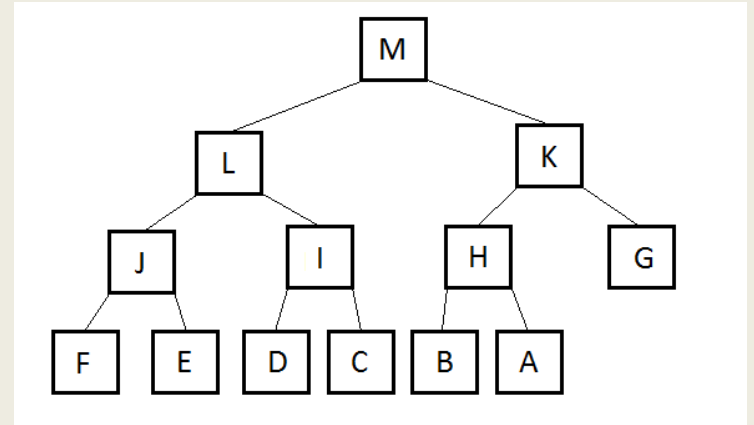
linked list implementation of a biiig integer in any base

Lab this week: review for test

One bug in doubly linked list code: posted correction

Tree traversals

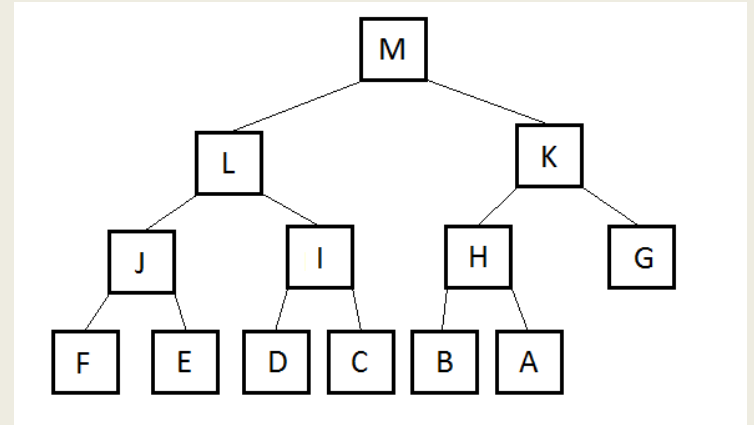
Recursive process



```
preorder_print (const binary_tree_node* bintree)
{
    if (bintree is not empty)
    {
        print out data at root
        preorder_print(left subtree of bintree's root)
        preorder_print(right subtree of bintree's root)
    }
}
```

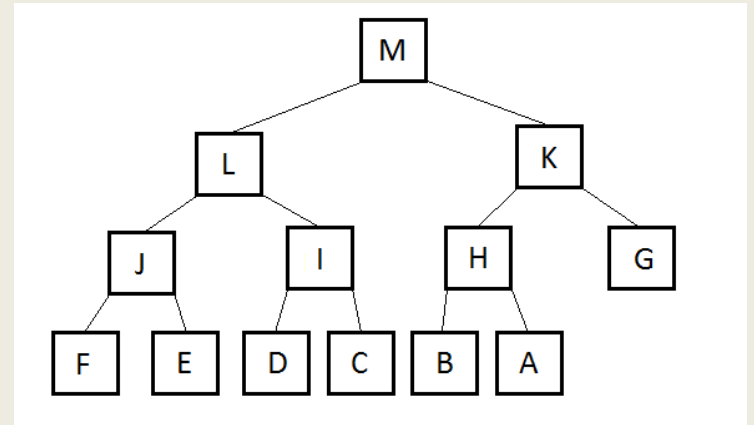
Tree traversals

Recursive process



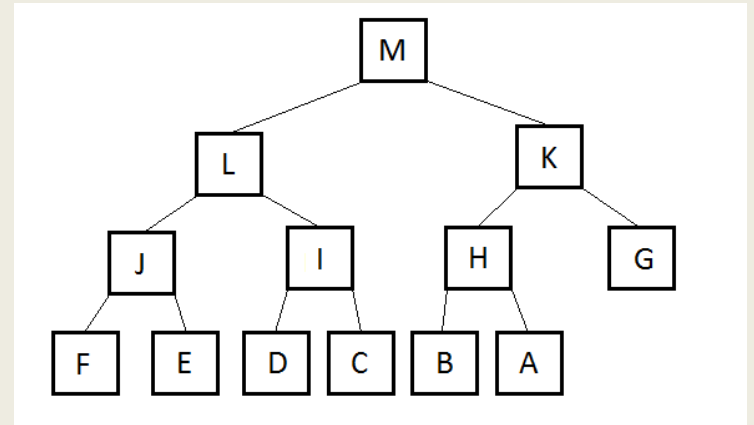
```
preorder_print (const binary_tree_node* bintree)
{
    if (bintree is not empty)
    {
        print out data at root
        preorder_print(left subtree of bintree's root)
        preorder_print(right subtree of bintree's root)
    }
}
```

Tree traversals



```
inorder_print (const binary_tree_node* bintree) :  
    if (bintree is not empty)  
    {  
        inorder_print(left subtree of bintree's root)  
        print out data at root  
        inorder_print(right subtree of bintree's root)  
    }
```

Tree traversals



```
postorder_print (const binary_tree_node* bintree) :  
    if (bintree is not empty)  
    {  
        postorder_print(left subtree of bintree's root)  
        postorder_print(right subtree of bintree's root)  
        print out data at root  
    }
```

Destroying a binary tree

```
void tree_clear(binary_tree_node*& root_ptr) {  
    binary_tree_node* child;  
    if (root_ptr != nullptr)  
    {  
        child = root_ptr->left;  
        tree_clear( child );  
        child = root_ptr->right;  
        tree_clear( child );  
        delete root_ptr;  
        root_ptr = nullptr;  
    }  
}
```

Destroying a binary tree

```
void tree_clear(binary_tree_node*& root_ptr) {  
    binary_tree_node* child;  
    if (root_ptr != nullptr)    // what happens if it is nullptr?  
    {        // destroy child nodes first (else we lose them)  
        child = root_ptr->left;  
        tree_clear( child );  
        child = root_ptr->right;  
        tree_clear( child );  
        delete root_ptr;  
        root_ptr = nullptr;  
    }  
}
```


Copying a binary tree

```
binary_tree_node* tree_copy(const binary_tree_node*  
    root_ptr) {  
    binary_tree_node* l_ptr; binary_tree_node* r_ptr;  
    if (root_ptr == nullptr) return nullptr;  
    else {  
        l_ptr = tree_copy( root_ptr->left );  
        r_ptr = tree_copy( root_ptr->right );  
        binary_tree_node* new_tree = new binary_tree_node;  
        new_tree->data = root_ptr->data;  
        new_tree->left = l_ptr; new_tree->right = r_ptr;  
        return new_tree;  
    }  
}
```

Traversal order matters (*post order*)

Tree_clear: child = root_ptr->left; tree_clear(child);
 child = root_ptr->right; tree_clear(child);
 delete root_ptr; root_ptr = nullptr;

Tree_copy: l_ptr = tree_copy(root_ptr->left);
 r_ptr = tree_copy(root_ptr->right);
 binary_tree_node* new_tree = new
 binary_tree_node;
 new_tree->data = root_ptr->data;
 new_tree->left = l_ptr; new_tree->right = r_ptr;
 return new_tree;

Trees offer the chance for $\log(n)$ performance if you're lucky

For a binary search tree of n nodes,

Search: $O(\log n)$ average, $O(n)$ worst case

Insert: $O(\log n)$ average, $O(n)$ worst case

Remove: $O(\log n)$ average, $O(n)$ worst case

Traverse: $O(n)$

Another tree ADT, the B-tree, ensures that the tree is perfectly full—guarantees the $O(\log n)$ performance, worst case