Analysis Report

cuda_kernel(int*, int*, unsigned long, unsigned long)

Duration	61.308 ms (61,307,621 ns)
Grid Size	[1,1251,1]
Block Size	[8,8,1]
Registers/Thread	29
Shared Memory/Block	0 B
Shared Memory Requested	96 KiB
Shared Memory Executed	96 KiB
Shared Memory Bank Size	4 B

[0] GeForce GTX 970

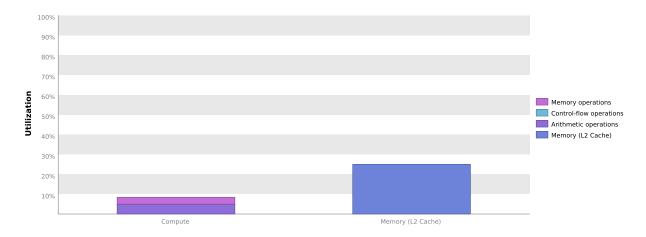
[U] Gerorce G1X 9/U						
GPU UUID	GPU-92461dda-c508-5f62-ee5f-85fc2de903e6					
Compute Capability	5.2					
Max. Threads per Block	1024					
Max. Threads per Multiprocessor	2048					
Max. Shared Memory per Block	48 KiB					
Max. Shared Memory per Multiprocessor	96 KiB					
Max. Registers per Block	65536					
Max. Registers per Multiprocessor	65536					
Max. Grid Dimensions	[2147483647, 65535, 65535]					
Max. Block Dimensions	[1024, 1024, 64]					
Max. Warps per Multiprocessor	64					
Max. Blocks per Multiprocessor	32					
Single Precision FLOP/s	4.17 TeraFLOP/s					
Double Precision FLOP/s	130.312 GigaFLOP/s					
Number of Multiprocessors	13					
Multiprocessor Clock Rate	1.253 GHz					
Concurrent Kernel	true					
Max IPC	6					
Threads per Warp	32					
Global Memory Bandwidth	224.32 GB/s					
Global Memory Size	3.942 GiB					
Constant Memory Size	64 KiB					
L2 Cache Size	1.75 MiB					
Memcpy Engines	2					
PCIe Generation	2					
PCIe Link Rate	5 Gbit/s					
PCIe Link Width	16					

1. Compute, Bandwidth, or Latency Bound

The first step in analyzing an individual kernel is to determine if the performance of the kernel is bounded by computation, memory bandwidth, or instruction/memory latency. The results below indicate that the performance of kernel "cuda_kernel" is most likely limited by instruction and memory latency. You should first examine the information in the "Instruction And Memory Latency" section to determine how it is limiting performance.

1.1. Kernel Performance Is Bound By Instruction And Memory Latency

This kernel exhibits low compute throughput and memory bandwidth utilization relative to the peak performance of "GeForce GTX 970". These utilization levels indicate that the performance of the kernel is most likely limited by the latency of arithmetic or memory operations. Achieved compute throughput and/or memory bandwidth below 60% of peak typically indicates latency issues.



2. Instruction and Memory Latency

Instruction and memory latency limit the performance of a kernel when the GPU does not have enough work to keep busy. The results below indicate that the GPU does not have enough work because instruction execution is stalling excessively.

2.1. Kernel Profile - PC Sampling

The Kernel Profile - PC Sampling gives the number of samples for each source and assembly line with various stall reasons. Using this information you can pinpoint portions of your kernel that are introducing latencies and the reason for the latency. Samples are taken in round robin order for all active warps at a fixed number of cycles regardless of whether the warp is issuing an instruction or not.

Instruction Issued - Warp was issued

Instruction Fetch - The next assembly instruction has not yet been fetched.

Execution Dependency - An input required by the instruction is not yet available. Execution dependency stalls can potentially be reduced by increasing instruction-level parallelism.

Memory Dependency - A load/store cannot be made because the required resources are not available or are fully utilized, or too many requests of a given type are outstanding. Data request stalls can potentially be reduced by optimizing memory alignment and access patterns.

Texture - The texture sub-system is fully utilized or has too many outstanding requests.

Synchronization - The warp is blocked at a __syncthreads() call.

Constant - A constant load is blocked due to a miss in the constants cache.

Pipe Busy - The compute resource(s) required by the instruction is not yet available.

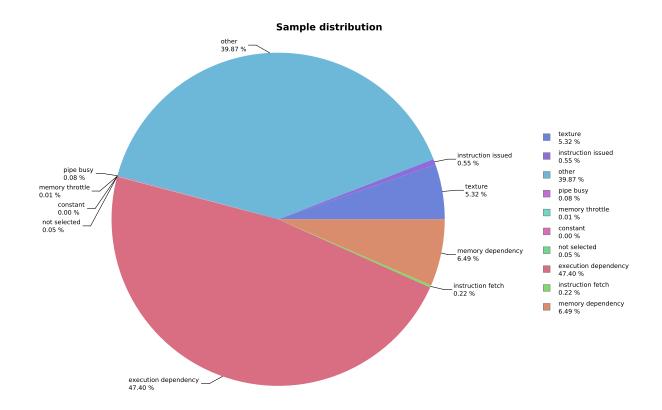
Memory Throttle - Large number of pending memory operations prevent further forward progress. These can be reduced by combining several memory transactions into one.

Not Selected - Warp was ready to issue, but some other warp issued instead. You may be able to sacrifice occupancy without impacting latency hiding and doing so may help improve cache hit rates.

Other - The warp is blocked for a uncommon reason.

Examine portions of the kernel that have high number of samples to know where the maximum time was spent and observe the latency reasons for those samples to identify optimization opportunities.

Cuda Functions	Sample Count	% of Kernel Samples
cuda_kernel(int*, int*, unsigned long,	4678761	100.0
unsigned long)		



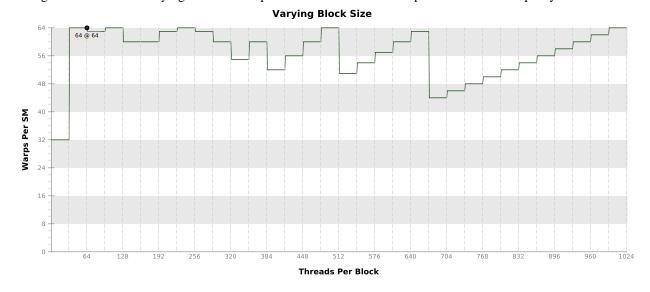
2.2. Occupancy Is Not Limiting Kernel Performance

The kernel's block size, register usage, and shared memory usage allow it to fully utilize all warps on the GPU.

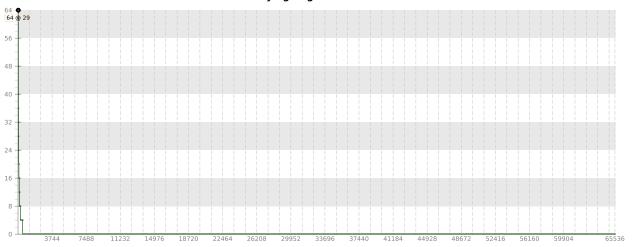
Variable	Achieved	Theoretical	Device Limit	Grid Size:	[1,1251	,1](1	L251 bl	ocks) Block	Size:	[8,8,1
Occupancy Per SM											
Active Blocks		32	32	0 4	8	12	16	20	24	28	32
Active Warps	63.06	64	64	0 9	9 18	2	7 36	5	45	54	664
Active Threads		2048	2048	0	512		1024		1536		2048
Occupancy	98.5%	100%	100%	0%	25%		50%		75%)	100%
Warps											
Threads/Block		64	1024	0	256		512		768		1024
Warps/Block		2	32	0 4	8	12	16	20	24	28	32
Block Limit		32	32	0 4	8	12	16	20	24	28	32
Registers											
Registers/Thread		29	65536	0	16384		32768		49152	2	65536
Registers/Block		2048	65536	0	16k		32k		48k		64k
Block Limit		32	32	0 4	8	12	16	20	24	28	32
Shared Memory											
Shared Memory/Block		0	98304	0	3	• 32k		6	4k		96k
Block Limit			32								

2.3. Occupancy Charts

The following charts show how varying different components of the kernel will impact theoretical occupancy.

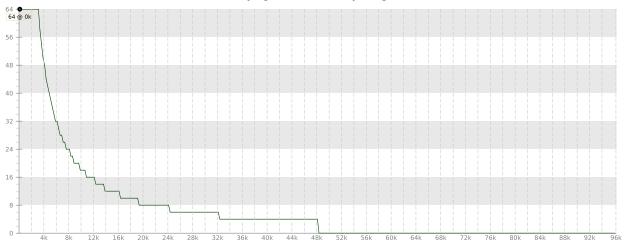


Varying Register Count



Registers Per Thread

Varying Shared Memory Usage



Shared Memory Per Block (bytes)

3. Compute Resources

GPU compute resources limit the performance of a kernel when those resources are insufficient or poorly utilized. Compute resources are used most efficiently when all threads in a warp have the same branching and predication behavior. The results below indicate that a significant fraction of the available compute performance is being wasted because branch and predication behavior is differing for threads within a warp.

3.1. Kernel Profile - Instruction Execution

The Kernel Profile - Instruction Execution shows the execution count, inactive threads, and predicated threads for each source and assembly line of the kernel. Using this information you can pinpoint portions of your kernel that are making inefficient use of compute resource due to divergence and predication.

Examine portions of the kernel that have high execution counts and inactive or predicated threads to identify optimization opportunities.

Cuda Fuctions :

cuda kernel(int*, int*, unsigned long, unsigned long)

Maximum instruction execution count in assembly: 3122500 Average instruction execution count in assembly: 793468 Instructions executed for the kernel: 338017518

Thread instructions executed for the kernel: 10805670576

Non-predicated thread instructions executed for the kernel: 10805235131

Warp non-predicated execution efficiency of the kernel: 99.9%

Warp execution efficiency of the kernel: 99.9%

3.2. Divergent Branches

Compute resource are used most efficiently when all threads in a warp have the same branching behavior. When this does not occur the branch is said to be divergent. Divergent branches lower warp execution efficiency which leads to inefficient use of the GPU's compute resources.

Optimization: Each entry below points to a divergent branch within the kernel. For each branch reduce the amount of intra-warp divergence.

3.3. Function Unit Utilization

Different types of instructions are executed on different function units within each SM. Performance can be limited if a function unit is over-used by the instructions executed by the kernel. The following results show that the kernel's performance is not limited by overuse of any function unit.

Load/Store - Load and store instructions for shared and constant memory.

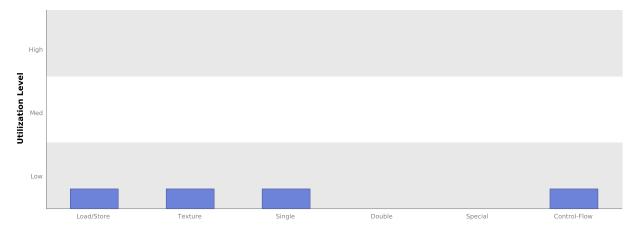
Texture - Load and store instructions for local, global, and texture memory.

Single - Single-precision integer and floating-point arithmetic instructions.

Double - Double-precision floating-point arithmetic instructions.

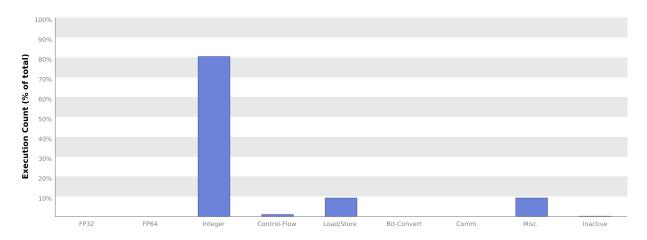
Special - Special arithmetic instructions such as sin, cos, popc, etc.

Control-Flow - Direct and indirect branches, jumps, and calls.



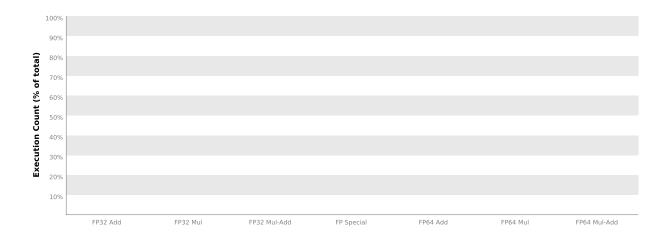
3.4. Instruction Execution Counts

The following chart shows the mix of instructions executed by the kernel. The instructions are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing instructions in that class. The "Inactive" result shows the thread executions that did not execute any instruction because the thread was predicated or inactive due to divergence.



3.5. Floating-Point Operation Counts

The following chart shows the mix of floating-point operations executed by the kernel. The operations are grouped into classes and for each class the chart shows the percentage of thread execution cycles that were devoted to executing operations in that class. The results do not sum to 100% because non-floating-point operations executed by the kernel are not shown in this chart.



4. Memory Bandwidth

Memory bandwidth limits the performance of a kernel when one or more memories in the GPU cannot provide data at the rate requested by the kernel.

4.1. Memory Bandwidth And Utilization

The following table shows the memory bandwidth used by this kernel for the various types of memory on the device. The table also shows the utilization of each memory type relative to the maximum throughput supported by the memory.

