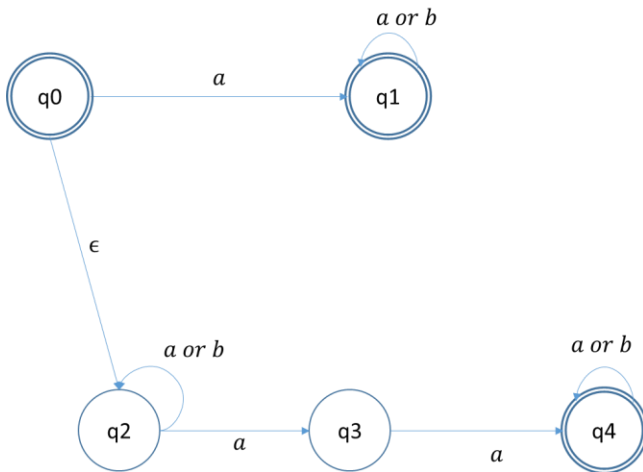


## Question 1

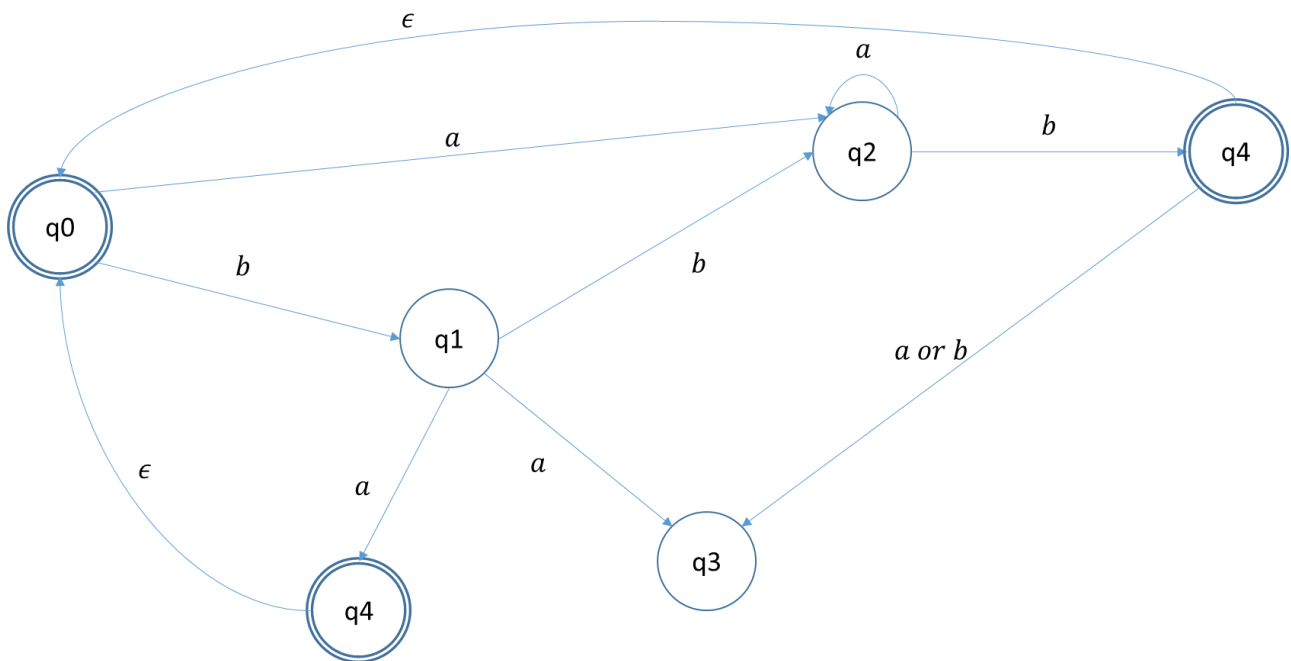
1.  $a(a+b)^*a$
2.  $(b(a+b))^*b^*$
3.  $a^*baaa^* + a^*abaa^* + a^*aaba^* + aaa^*$
4.  $(a+b)(a+b)^*$

## Question 2

1.  $\epsilon + a(a+b)^* + (a+b)^*aa(a+b)^*a(a+b)^*$



2.  $[ba + (a + bb)a^*b]^*$



### Question 3

In fact the only condition for it to work is to satisfies this expression which correspond to having at least an  $a$  and  $b$

$$\Sigma^*$$

Let take a string that satisfies the expression

Let split this string in two pieces  $x$  and  $y$  with  $n$   $a$  in  $x$  and  $m$   $a$  in  $y$  and  $n, m \geq 0$

If  $n = m$  then we are done the string is in the set

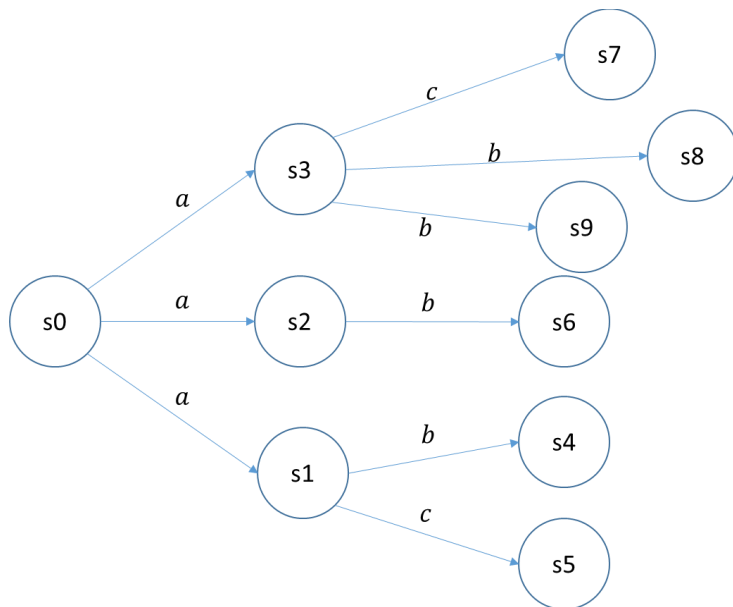
Let suppose  $n < m$

Then we can move the split of the string one letter further. If the letter is  $a$  then we have  $n + 1$   $a$ 's in  $x$  and still  $m$   $b$ 's in  $y$ , if it's a  $b$  then we have still  $n$   $a$ 's in  $x$  and  $m - 1$   $b$ 's in  $y$  Then we check if the amount are the same now, if not then we can repeat this step until they match

Similarly with  $n > m$  we move the split left

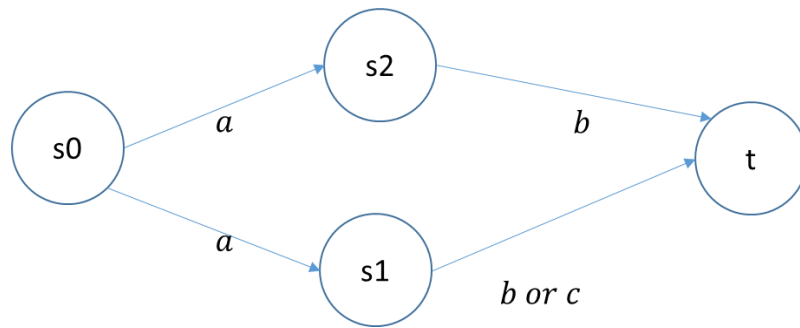
Then this expression define this language and thus it's a regular language.

### Question 4



We see that  $s_1$  and  $s_3$  are bisimilar as we have the same outgoing transition (b and c)

We also have  $s_4$  to  $s_9$  bisimilar as they don't have any outgoing transitions



## Question 5

	$\epsilon$	$a$	$b$	$aa$	$bb$
$\epsilon$	Eq				
$a$		Eq			
$b$			Eq		
$aa$				Eq	
$bb$			Eqq		Eq

Nothing starts with  $b$  that is in  $L$ . All equivalent classes on  $a$  are equivalent to themselves.  $ax$  is equivalent to  $ax$  for any  $x$ .

The only equivalent classes in  $L$  are on  $a^*$

Because  $a^n, n \in \mathbb{N}$  there are infinite equivalence classes as there is an equivalent classes for each  $a^n$  where  $n \in \mathbb{N}$ ,  $\mathbb{N}$  is distinct from itself