Project Report Group Badgers

Java and C# in depth, Spring 2013

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1 Introduction

This document describes the design and implementation of the *Personal Virtual File System* of group *Badgers*. The project is part of the course *Java and C# in depth* at ETH Zurich. The following sections describe each project phase, listing the requirements that were implemented and the design decisions taken. The last section describes a use case of using the *Personal Virtual File System*.

2 VFS Core

VFS Core is a library that provides an implementation of a virtual file system. The API that a client of this library can use consists of three interfaces that are described in 2.2. The VFS Core provides functionality to create/open/dispose new virtual disks and allows the management of files and directories within such a disk. Furthermore it provides a simple way to import/export files from/to the host file system.

The library internally works with a virtual disk that is divided into a header, index and data section, having the index represented as B-tree. Such a design allows performant access to the data inside the disk.

2.1 Requirements

Below one will find a list of the requirements implemented in this project.

2.1.1 disk management

- The virtual disk must be stored in a single file in the working directory in the host file system.
- VFS must support the creation of a new disk with the specified maximum size at the specified location in the host file system.
- VFS must support several virtual disks in the host file system.
- VFS must support disposing of the virtual disk.
- VFS must support querying of free/occupied space in the virtual disk.

These requirements are met with the classes

```
ch.eth.jcd.badgers.vfs.core.VFSDiskManagerImpl
ch.eth.jcd.badgers.vfs.core.config.DiskConfiguration
```

allow creation/deletion and opening of the disk on the host file system. Clients of the library have to pass a DiskConfiguration to the VFSDiskManagerImpl when calling create or open

2.1.2 file management

- VFS must support creating/deleting/renaming directories and files.
- VFS must support navigation: listing of files and folders, and going to a location expressed by a concrete path.
- VFS must support moving/copying directories and files, including hierarchy.
- VFS must support importing files and directories from the host file system.
- VFS must support exporting files and directories to the host file system.

These requirements are met with the classes

```
ch.eth.jcd.badgers.vfs.core.VFSEntryImpl
ch.eth.jcd.badgers.vfs.core.VFSPathImpl
ch.eth.jcd.badgers.vfs.core.VFSFileInputStream
ch.eth.jcd.badgers.vfs.core.VFSFileOutputStream
```

The *VFSEntryImpl* allows copy, move (and rename), delete, listing of children and going to the parent (navigation). Together with the streams it also supports importing/exporting to any location clients of the vfs library wish to. The classes

```
ch.eth.jcd.badgers.vfs.ui.VFSConsole
ch.eth.jcd.badgers.vfs.ui.VFSUIController
```

demonstrate this by importing and exporting to the host file system.

2.1.3 bonus features

• Elastic disk: Virtual disk can dynamically grow or shrink, depending on its occupied space.

The implementation only allows growing if more files are imported. Shrinking is not supported.

- Compression, if implemented with 3d party library.
- Compression, if implemented by hand (you can take a look at the arithmetic compression)

The classes

```
ch.eth.jcd.badgers.vfs.compression.BadgersLZ77CompressionInputStream ch.eth.jcd.badgers.vfs.compression.BadgersLZ77CompressionOutputStream ch.eth.jcd.badgers.vfs.compression.BadgersRLECompressionInputStream ch.eth.jcd.badgers.vfs.compression.BadgersRLECompressionOutputStream
```

implement streams that can be wrapped around VFSFileInputStream and VFSFileOutputStream. The DiskConfiguration allows to switch compression on and to declare which algorithm shall be chosen. This allows easy configuration of any 3d party compression streams (which was not chosen to implement, because of the implementation of our own compression algorithm).

- Encryption, if implemented with 3d party library.
- Encryption, if implemented by hand.

The classes

ch.eth.jcd.badgers.vfs.encryption.CaesarInputStream
ch.eth.jcd.badgers.vfs.encryption.CaesarOutputStream

show how encryption can be implemented in the library. It was chosen to implement encryption similar to compression with streams, which allows easy configuration via *DiskConfiguration*. These streams are mainly for demonstration how encryption should work and shall not be used in high security environments:-)

2.2 Core Classes

Figure 1 gives an overview of the main interfaces and classes that were implemented. The interfaces VFSDiskManager, VFSEntry and VFSPath can be used by clients using the library implemented here.

- VFSDiskManager The implementations of this interface provide mainly a way to open, create and dispose new virtual disks. Additionally one can get the root entry of the file system and get additional information about an opened disk.
- VFSEntry Represents a directory or file on the file system. A VFSEntry provides all the required methods to manipulate files and directories and importing/exporting files into the virtual file system. The general meaning of VFSEntry is, that such objects usually exist on the filesystem.
- VFSPath Represents a path on the file system to a given VFSEntry. It has a slightly looser coupling to the file system as a path does not imperatively need to exist.

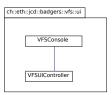
The intention of those interfaces is to hide the real implementation of the virtual file system from a client. With that in mind it should be simple to add a network layer upon the real implementation without changing client code.

2.3 Mocking

For discussing the semantics of the virtual file system and the development of the interfaces explained in 2.2 it was decided to implement a mock that works against the host file system. The mock classes implement all the interfaces and were very helpful for acquiring a common understanding of how the interface shall be used by clients. In a further step it was very useful to have the mock classes while developing the console application which could be developed independently from the real implementation.

2.4 Test

During the development a bunch of test cases came to life. The tests solely depend on the interfaces and thus they can run against the mock classes and the real implementation. This was a huge help in finding bugs in the slightly more complicated real implementation.



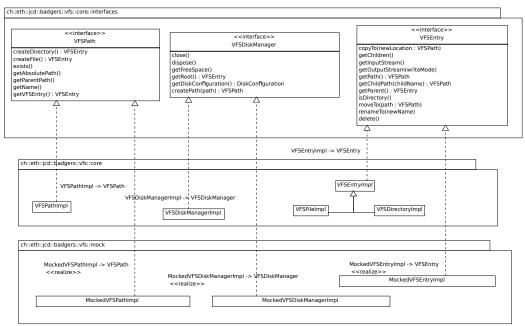


Figure 1: core classes

2.5 The *real* implemenation

Eventually some code that handles a virtual disk was developed and is described in this section.

2.5.1 The File Format

This section describes the binary file format used by the file system inside a virtual disk. The file is separated into three major parts. The header, index and the data section. Each of them is described below.

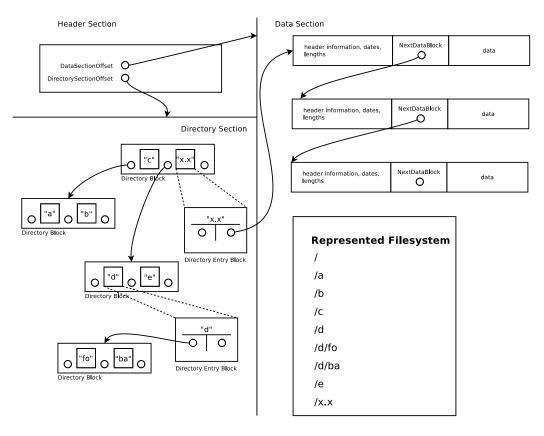


Figure 2: Overview of a disk

Header Section The header section contains some general information about the currently opened virtual disk. The details can be found in the following table:

Name	Length	Description
Info	50 byte UTF-8 String	Contains something like
		Badger VFS 2013 V1.0
Version	10 byte UTF-8 String	Contains something like
		"1.0"
Compression used	20 byte UTF-8 String	null or indicates compres-
		sion used for this file
Encryption used	20 byte UTF-8 String	null or indicates encryption
		used for this file
DirectorySectionOffset	long (8 byte)	File offset where our direc-
		tory section starts
DataSectionOffset	long (8 byte)	File offset where our data
		section starts
SaltString	8 bytes	Salt used to hash username
		and password randomly
		string generated while
		creating this file
Password	xxx bytes	CryptoHash (SHA-
		whatever) of Pass-
		word+SaltString

Directory Section The directory section describes which files and folders belong to which parent directory. This section has a fixed size and contains so called DirectoryBlocks which also have a fixed size. This makes management an manipulation easy. To each directory belongs a B-Tree structure which lists all contained entries.

Directory Block One Directory Block represents a Node in our B-Tree of order 2.

Name	Length	Description
DirectoryHeader	1 byte	Header information. This header makes it easy to determine whether a DirecotryBlock is in used or not (memory management)
DirectoryEntryBlock1	128 byte	The smaller key inserted into our B-Tree
DirectoryEntryBlock2	128 byte	The bigger key inserted into our B-Tree
DirectoryBlockLink1	8 byte	Points to another DirectoryBlock which contains keys smaller than DirectoryEntryBlock1
DirectoryBlockLink2	8 byte	Points to another DirectoryBlock which contains keys bigger than DirectoryEntry-Block1 but smaller than DirectoryEntryBlock2
DirectoryBlockLink3	8 byte	Points to another DirectoryBlock which contains keys bigger than DirectoryEntry-Block2

Directory Entry Block Represents a single directory or file.

Name	Length	Description
Filename	112 byte	UTF-8 Filename String
DataBlockLocation	8 byte	Pointer to a DataBlock located in the Data Section. This DataBlock holds some meta information about the current directory
DirectoryEntryTreeRoot	8 bytes	Pointer to a DirectoryBlock located in the Directory Section. This referenced DirectoryBlock is the Root Block of a B-Tree containing all entries of that directory specified by the current Directory Entry Block. This field containing a 0 indicates that this entry is a file not a directory

Data Section The data section is split into blocks where each of them is X bytes long. Each block contains some amount of data and points to a subsequent block

Block layout

Name	Length	Description
BlockHeader	1 byte	
0) Header-Bit (LSB)		If set to 1 this is the first
		datablock of a file.
1) not used		
2) not used		
3) not used		
4) not used		
5) not used		
6) not used		
7) not used		
NextDataBlock	8 byte	Points to the start ad-
		dress of the next Datablock
		(linked list). 0 if this is the
		last DataBlock of a certain
		file or folder.
CreationDate	8 byte	UTC Time when this file
		was created
		This field only exists if
		Header-Bit is set to 1
DataLength	4 byte	Indicates the number of
		data saved on this Dat-
		aBlock.
Data	n byte	user data (may be en-
		crypted/compressed)

The root directory

2.5.2 compression layer

To reduce the data volume within the virtual disk, compression on each file can be enabled. Currently available compression algorithms are run length encoding [1] and LZ77 [2].

Run Length Encoding The available 8bit run length encoding(rle) algorithm is a very simple form of data compression where multiple occurrence of the same byte were stored as a single byte value and the corresponding count. It is useful for simple graphic images like line drawings and icons.

LZ77 Abraham Lempel and Jacob Ziv introduced the LZ77 lossless compression algorithm in 1977. Newer compression methods such as GZIP or DEFLATE often use LZ77-based algorithms. The compression is achieved by replacing the data with a reference to an earlier existing copy in the data input stream. For that a window of a certain size is held in memory where existing copies of the current data are searched.

3 Quick Start Guide

3.1 the eclipse project

The project requires to be compiled with JAVA 7. It also depends on the maven plugin which pulls in all the required libraries.

3.2 Command line client

The command line client allows the usage of the VFS core and is mainly intended to test the basic functionalities. The console runs either in management mode or in filesystem mode. The management mode is entered automatically when starting the command line client. It allows creating and opening virtual disks. The filesystem mode is entered as soon as a virtual disk is opened.

3.2.1 startup

The command line client can be started as follows:

```
java -jar VFSCore.jar ch.eth.jcd.badgers.vfs.ui.VFSConsole
```

or by starting ch.eth.jcd.badgers.vfs.ui.VFSConsole in eclipse. This gives a console prompt where the following commands can be used

3.2.2 commands

in.

Following commands can be used with the command line client in management mode:

- create c:\path\to\disk.bfs 1024 creates virtual disk with a maximum quota of 1024 megabytes on the host system. The file may grow up to 1024 megabytes. TODO: more parameters are needed (encryption, compression, password if there is encryption)
- open c:\path\to\disk.bfs opens filesystem mode for the given virtual disk
- exit exits the console program

follwing commands can be used in filesystem mode:

• **ls** lists the contents of the current directory

- pwd shows the path to the current directory
- cd dst changes current directory to dst which must be either a child directory of the current path or ".."
- find searchString lists absolute paths of all files containing search-String in their file name
- mkdir dirName creates a new directory dirName in the current path
- mkfile fileName creates a new empty file fileName in the current path this is rather not useful, as the "import" creates a file with content
- rm file deletes the entry denoted as *file*, it must be a child of the current path
- cp src dst copies the src file to dst as a child of the current path
- mv src dst moves the src file to dst
- import ext_src dst imports a ext_src from the host system to dst
- export src ext_src exports a src file to the host system ext_dst
- find searchString lists all filesystem entries below the current entry containing searchString
- dispose deletes the currently opened virtual disk
- **close** closes the file system mode, from now on management mode commands can be executed

4 Glossary

VFS core The main Java library, that handles all the interaction with virtual disks and importing/exporting/storing files. It is used by the command line client and the gui.

Virtual Disk A virtual disk denotes a container file that is stored on the host file system. A virtual disk can be opened with the software that is developed during this project and stores the actual files. The file extension of the virtual disk is "*.bfs".

References

- [1] Run Length Encoding.
- [2] A. Lempel and J. Ziv. A universal algorithm for sequential data compression. *IEEE Transactions on Information Theory 23, 337-343*, 1997.