Numerical Simulation and Scientific Computing II

Lecture 1: Introduction and Distributed Parallel Computing I



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Outline

- Introduction to the Lecture
- Distributed Parallel Computing I
 - Primer
 - Overview
 - Process Model and Language Bindings
 - Messages and Point-to-Point Communication
 - Examples
- Quiz

Team Presentation

Core Team (Organization, MPI, ODE, FVM)

- Josef Weinbub
- Paul Manstetten
- Luiz Felipe Aguinsky
 Institute for Microelectronics

Extended Interdisciplinary Team

- Francesco Zonta (Fluid Dynamics)
 Institute of Fluid Mechanics and Heat Transfer
- Heinz Pettermann and Marius Schasching (Finite Elements)
 Institute of Lightweight Design and Structural Biomechanics
- Jesús Carrete Montana (Molecular Dynamics)
 Institute of Materials Chemistry
- Kevin Sturm (Partial Differential Equations)
 Institute of Analysis and Scientific Computing





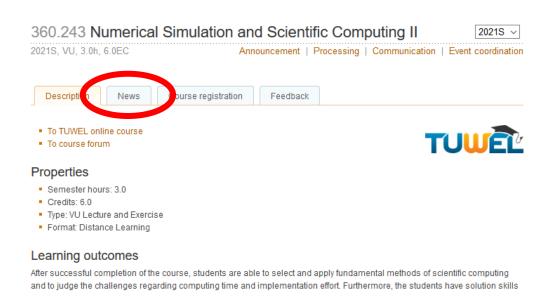






TISS

- Acceptance to course via TISS: Check today
 https://tiss.tuwien.ac.at/course/courseAnnouncement.xht
 ml?dswid=9250&dsrid=235&courseNumber=360243&courseSemester=2021S
- News will be sent via TISS' notification feature: Regularly monitor your TISS News (at least weekly)!

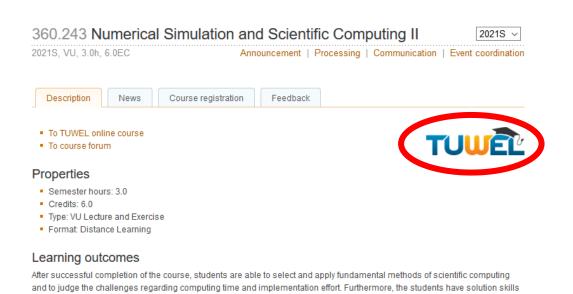


TUWEL

 Course material (slides, handouts, zoom link) provided via TUWEL course:

https://tuwel.tuwien.ac.at/course/view.php?idnumber=36 0243-2021S

 TUWEL course also "reachable" via TISS course, watch out for the TUWEL icon:



General Goals

- Introduction to advanced methods of CSE
 - Distributed Parallel Computing
 - Methods for Ordinary Differential Equations
 - Classification and Analyses of Partial Differential Equations
 - Finite Volume Method
 - Finite Element Method
 - Fluid Dynamics
 - Molecular Dynamics

Course Calendar

March 4	Lecture
March 11	Lecture with Exercise Handout
March 18	Lecture with Exercise Support
March 25	Lecture with Exercise Support
April 13	Exercise 1 Submission
April 15	Lecture
April 22	Lecture with Exercise Handout
April 29	Lecture with Exercise Support
May 4	Exercise 2 Submission
May 6	Lecture with Exercise Handout
May 12	Exercise Support
May 20	Lecture with Exercise Handout and Support
May 26	Exercise 3 Submission
May 27	Lecture with Exercise Support
June 2	Exercise Support
June 10	Exercise 4 Submission
June 30	Written (Online) Exam

- Every slot starts at 14:00 sharp!
- Will put lecture slides/exercise material on TUWEL right before slot.
- Exercise submissions until 8:00!
- Schedule also in TISS!

Necessary Background

- C++ and Python! (see TISS)
 - If not experienced: Checkout online tutorials asap!
- Not a formal requirement yet the expected level of expertise: Numerical Simulation and Scientific Computing I 360.242

Exercise Rules and Course Grade

- 4 mandatory exercises over the whole lecture
- Groups of 2-3 students must be formed; student-picked.
- Code copied? 0 points for both groups!
- Each exercise will be graded separately from 0-10 points
- Access for the final exam: sum of the points ≥ 28.0 (same procedure as with NSSC I)
- Course grade: 1/4th Exercises, 3/4th Exam
- Submission deadlines are hard deadlines
- Submissions: per email to <u>nssc@iue.tuwien.ac.at</u> as a single compressed file (More information with each exercise)

Rules - Quizzes

- In the end of each lecture, you will receive 5 questions
- 3 questions reviewing the current lecture
- 2 questions preparing for the next
- Discussion in the beginning of next class
- Participation is voluntary but encouraged
- These questions might help you learn for the exam!

Computer Resources

- It is expected that you have access to a modern, x86based, personal computer (laptop, home desktop, etc.)
- Linux (virtual machine, dual boot or full OS) is required
 - More information on the development environment for a particular exercise will be made available with the exercise handout

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Acknowledgments

Thanks to Rolf Rabenseifner's (HLRS)
2014 Parallel Programming Workshop Lecture Material on:
Introduction to the Message Passing Interface (MPI)

Sources

- High Performance Computing Center Stuttgart (HLRS)
 Online Courses
 https://www.hlrs.de/about-us/media-publications/teaching-training-material/
- YouTube -- search for "Introduction to MPI", e.g., <u>https://youtu.be/RoQJNx5npF4</u> -- Part I of III

Additional Courses at TU Wien

ECTS Courses (count as free electives)

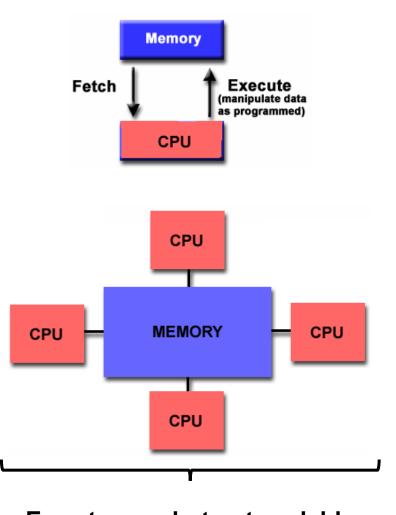
- 057.020 (Winter term)
 VSC-School I Courses in High Performance Computing
- 057.021 (Summer term)
 VSC-School II Courses in High Performance Computing

Non-ECTS Trainings (don't count as free electives)

- Node-Level Performance Engineering
- OpenMP
- MPI
- Deep-Learning und GPU programming (OpenACC)
- Hybrid-programming MPI+X

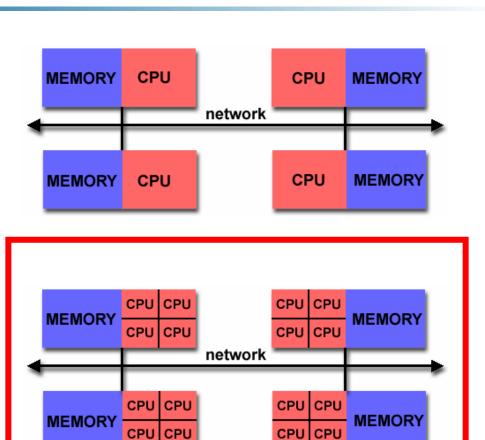
http://typo3.vsc.ac.at/research/vsc-research-center/vsc-school-seminar/

Parallel Computer Architectures



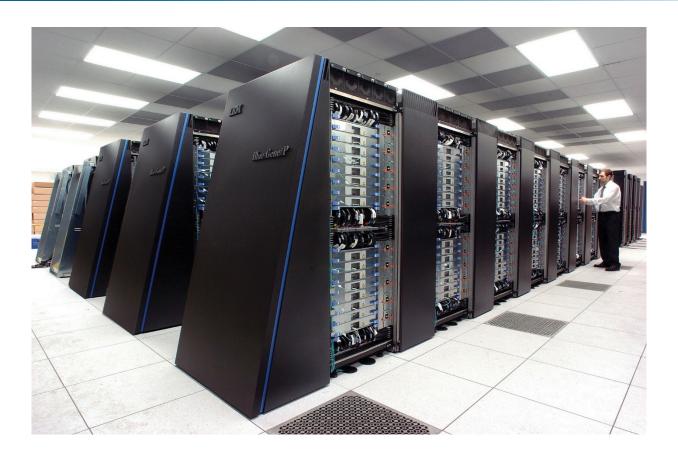
Easy to use, but not scalable

Source: Lyle N. Long, PSU



Difficult to use, but scalable

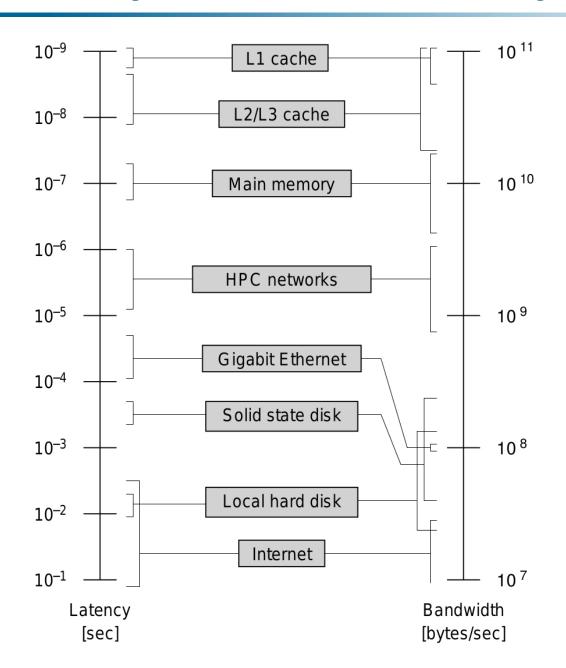
Distributed Parallel Computing Clusters



IBM Blue Gene/P: Intrepid Argonne National Laboratory

Source: Wikipedia

Latency / Bandwidth Hierarchy



Processes and Threads

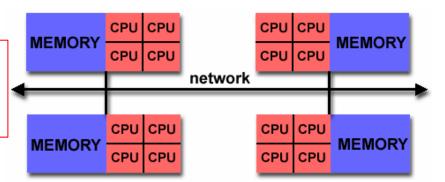
Process

A process is a program in execution

Represented by a process control block: all necessary

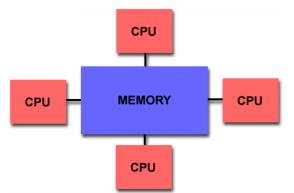
information to run a program

Data exchange happens explicitly via sending messages between processes!



Thread

- Unit of CPU utilization
- A single process can contain multiple threads
- Threads belonging to the same process can share resources



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Message Passing Interface (MPI)

#include "mpi.h"

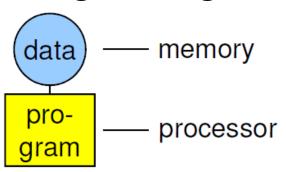
I am 0 out of 4

I am 2 out of 4

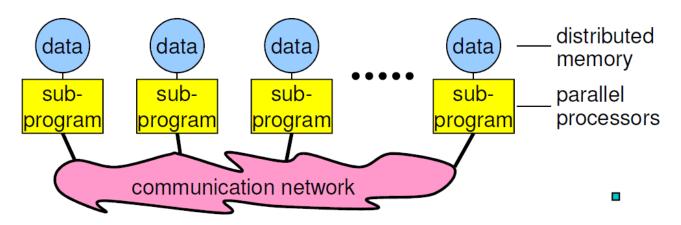
```
#include <stdio.h>
int main( int argc, char *argv[] )
 int myid, numprocs;
 MPI Init(&argc,&argv);
 MPI Comm_size(MPI_COMM_WORLD,&numprocs);
 MPI_Comm_rank(MPI_COMM_WORLD,&myid);
 printf("I am %d out of %d\n", myid, numprocs);
 MPI Finalize();
 return 0;
> mpiexec -n 4 ./hello
I am 3 out of 4
I am 1 out of 4
```

Message Passing Programming Paradigm

Sequential Programming Paradigm



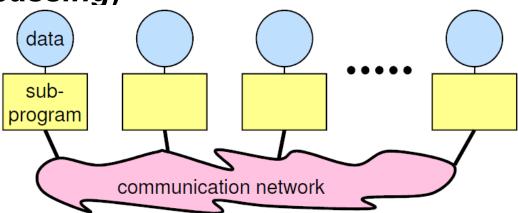
Message Passing Programming Paradigm



Message Passing Programming Paradigm

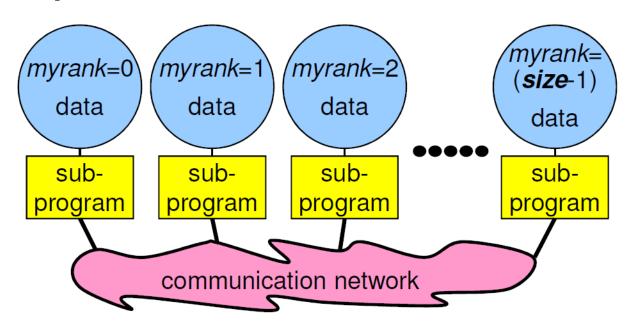
- Each processor in a message passing program runs a sub-program:
 - written in a conventional sequential language, e.g., C or Fortran,
 - typically the same on each processor,
 - the variables of each sub-program have
 - the same name
 - but different locations (distributed memory) and different data!
 - i.e., all variables are private

 communicate via special send & receive routines (message passing)



Data and Work Distribution

- the value of myrank is returned by special library routine
- the system of size processes is started by special MPI initialization
- program (mpirun or mpiexec)
- all distribution decisions are based on myrank
- i.e., which process works on which data



What is SPMD?

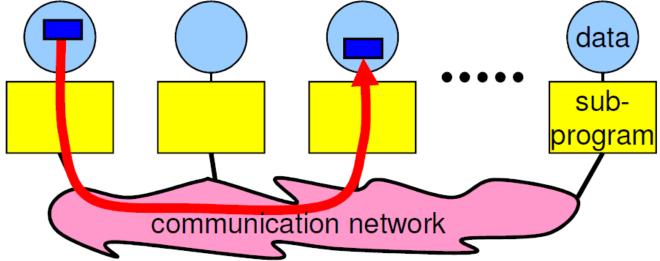
- Single Program, Multiple Data
- Same (sub-)program runs on each processor
- MPI allows also <u>MP</u>MD, i.e., <u>Multiple Program</u>, ...
- but some vendors may be restricted to SPMD
- MPMD can be emulated with SPMD

Emulation of Multiple Program (MPMD)

```
main(int argc, char **argv)
{
  if (myrank < .... /* process should run the ocean model */)
  {
    ocean( /* arguments */ );
  }
  else{
    weather( /* arguments */ );
  }
}</pre>
```

Messages

- Messages are packets of data moving between subprograms
- Necessary information for the message passing system:
 - sending process receiving process
 source location destination location
 source data type destination data type
 source data size destination buffer size



Access

- A sub-program needs to be connected to a message passing system
- A message passing system is similar to:
 - mail box
 - phone line
 - fax machine
 - · etc.

• MPI:

- sub-program must be linked with an MPI library
- sub-program must use include file of this MPI library
- the total program (i.e., all sub-programs of the program) must be started with the MPI startup tool

Addressing

- Messages need to have addresses to be sent to.
- Addresses are similar to:
 - mail addresses
 - phone number
 - fax number
 - etc.
- MPI: addresses are ranks of the MPI processes (sub-programs)

Reception

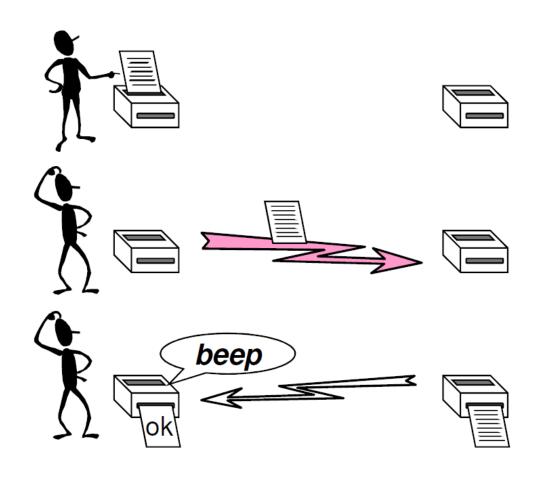
• All messages must be received.

Point-to-Point Communication

- Simplest form of message passing.
- One process sends a message to another.
- Different types of point-to-point communication:
 - synchronous send
 - buffered = asynchronous send

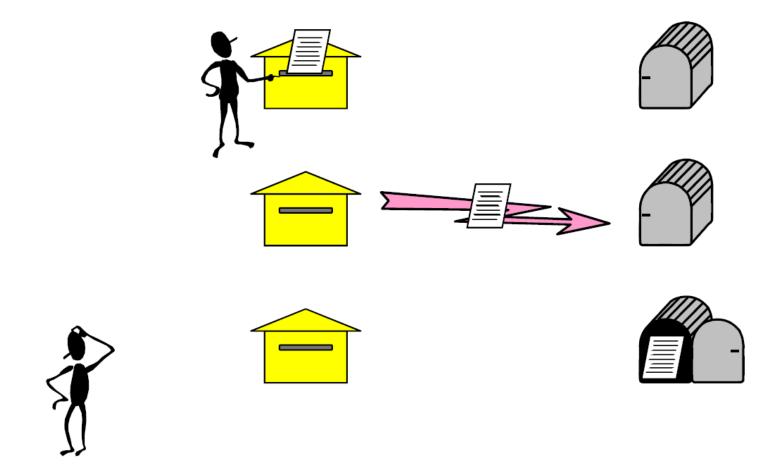
Synchronous Sends

- The sender gets an information that the message is received.
- Analogue to the beep or okay-sheet of a fax.



Buffered = Asynchronous Sends

Only know when the message has left.

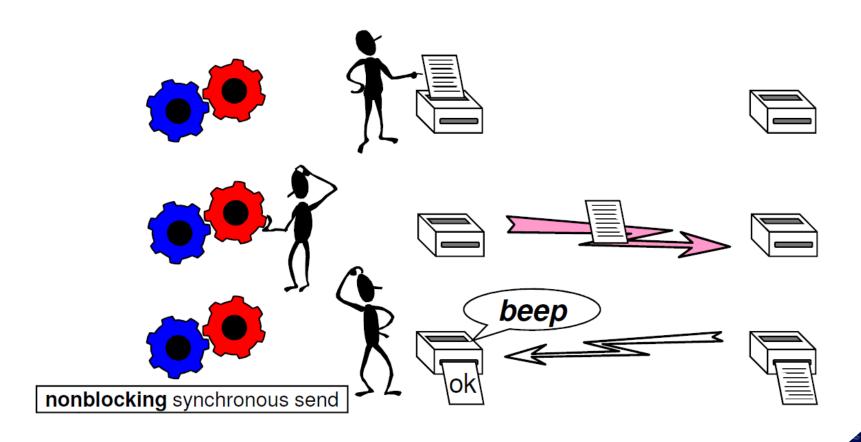


Blocking Operations

- Operations are local activities, e.g.,
 - sending (a message)
 - receiving (a message)
- Some operations may block until another process acts:
 - synchronous send operation blocks until receive is posted;
 - receive operation blocks until message was sent.
- Relates to the completion of an operation.
- Blocking subroutine returns only when the operation has completed.

Non-Blocking Operations

- Non-blocking operation: returns immediately and allow the sub-program to perform other work.
- At some later time the sub-program must test or wait for the completion of the non-blocking operation.



Non-Blocking Operations (cont'd)

- All non-blocking operations must have matching wait (or test) operations.
 (Some system or application resources can be freed only when the non-blocking operation is completed.)
- A non-blocking operation immediately followed by a matching wait is equivalent to a blocking operation.
- Non-blocking operations are not the same as sequential subroutine calls:
 - the operation may continue while the application executes the next statements!

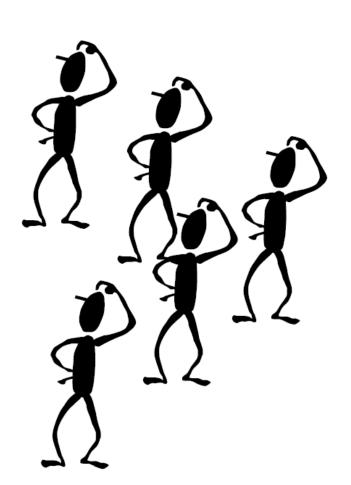
Collective Communications

- Collective communication routines are higher level routines.
- Several processes are involved at a time.
- May allow optimized internal implementations, e.g., tree based algorithms.
- Can be built out of point-to-point communications.

Broadcast

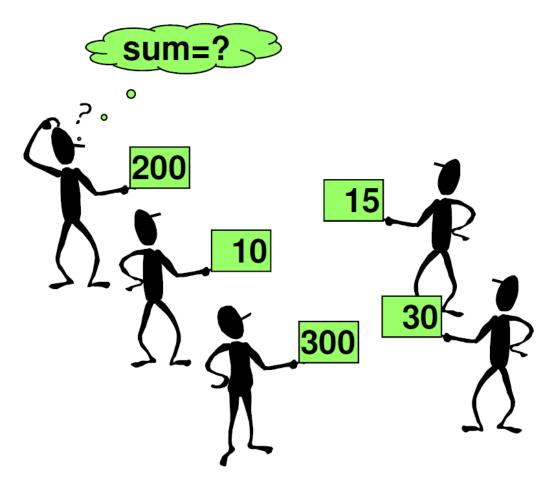
A one-to-many communication.





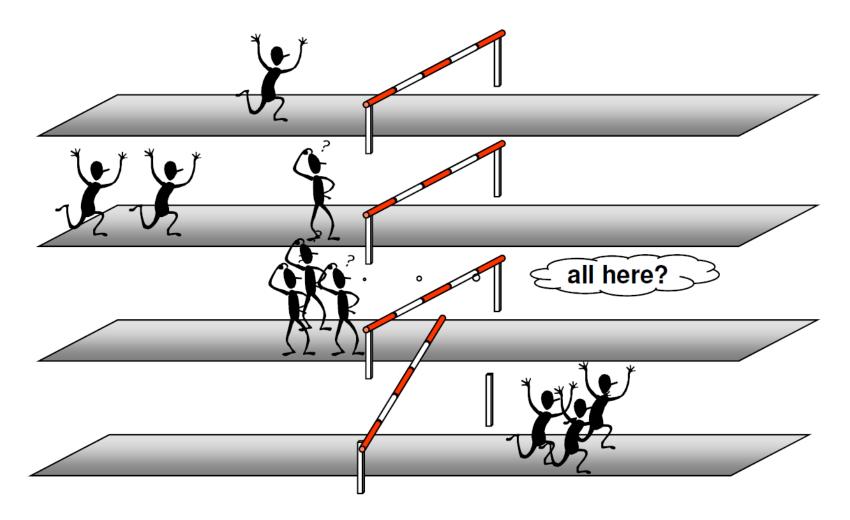
Reduction Operations

 Combine data from several processes to produce a single result.



Barriers

• Synchronize processes.



MPI Forum

- MPI is a standard
 - Each MPI routine is defined
- MPI Forum: the standardization forum www.mpi-forum.org/
- MPI-1.0 June 1994
- MPI-3.1 June 2015
- MPI-4.0 under development
- Many MPI libraries ((mostly) adhering to the standard) are available
 - Open source: OpenMPI, MPICH
 - Native support for C/C++, Fortran (this course focuses only on C/C++)

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Header Files and Function Format

Header Files

```
#include <mpi.h>
```

- MPI Namespace
 - MPI_... namespace is reserved for MPI constants and routines
 - Application routines and variable names must not begin with MPI_

Function Format

```
error = MPI_Xxxxxx( parameter, ... );
MPI_Xxxxxx( parameter, ... );
```

Initializing MPI

First routine to be called

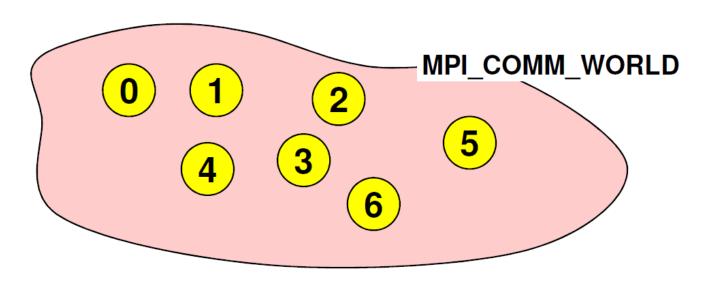
```
int MPI_Init( int *argc, char ***argv)
#include <mpi.h>
int main(int argc, char **argv)
{
    MPI_Init(&argc, &argv);
    ...
}
```

Starting the MPI Program

- Start mechanism is implementation dependent
- mpirun –np number_of_processes ./executable (most implementations)
- mpiexec –n number_of_processes ./executable (with MPI-2 and later)

Communicator MPI_COMM_WORLD

- All processes (= sub-programs) of one MPI program are combined in the communicator MPI_COMM_WORLD.
- MPI_COMM_WORLD is a predefined <u>handle</u> in mpi.h.
- Each process has its own rank in a communicator:
 - starting with 0
 - ending with (size-1)



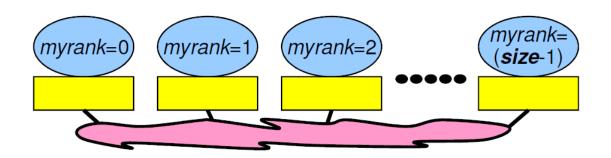
Handles

- Handles identify MPI objects.
- For the programmer, handles are
 - predefined constants in mpi.h
 - Example: MPI_COMM_WORLD
 - Can be used in initialization expressions or assignments.
 - The object accessed by the predefined constant handle exists and does not change only between MPI_Init and MPI_Finalize.
 - values returned by some MPI routines, to be stored in variables, that are defined as

Rank

- The rank identifies different processes.
- The rank is the basis for any work and data distribution.

int MPI_Comm_rank(MPI_Comm comm, int *rank)



CALL MPI_COMM_RANK(MPI_COMM_WORLD, myrank, ierror)

Size

 How many processes are contained within a communicator?

int MPI_Comm_size(MPI_Comm comm, int *size)

Exiting MPI

- Must be called last by all processes.
- User must ensure the completion of all pending communications (locally) before calling finalize
- After MPI_Finalize:
 - Further MPI-calls are forbidden
 - Especially re-initialization with MPI_Init is forbidden
 - May abort all processes except "rank==0" in MPI_COMM_WORLD

```
int MPI_Finalize()
```

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Messages

- A message contains a number of elements of some particular datatype.
- MPI datatypes:
 - Basic datatype.
 - Derived datatypes.
- Derived datatypes can be built up from basic or derived datatypes.
- Datatype handles are used to describe the type of the data in the memory.

Example: message with 5 integers

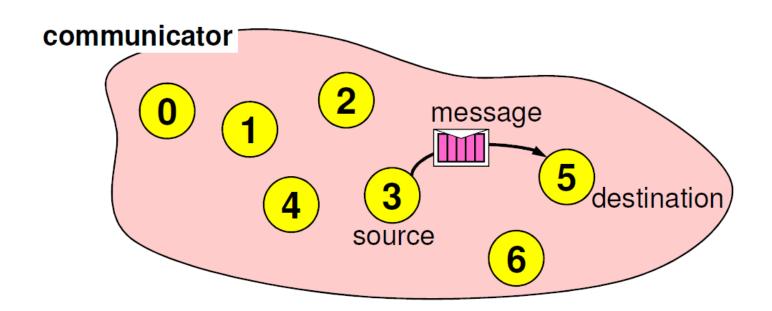
2345	654	96574	-12	7676

MPI Basic Datatypes

MPI Datatype	C datatype	Remark	s	
MPI_CHAR	char	Treated	Treated as printable character	
MPI_SHORT	signed short int			
MPI_INT	signed int			
MPI_LONG	signed long int			
MPI_LONG_LONG	signed long long			
MPI_SIGNED_CHAR	signed char	Treated as integral value		
MPI_UNSIGNED_CHAR	unsigned char	Treated as integral value		
MPI_UNSIGNED_SHORT	unsigned short int			
MPI_UNSIGNED	unsigned int		Further datatypes,	
MPI_UNSIGNED_LONG	unsigned long int		see, e.g., MPI-3.0,	
MPI_UNSIGNED_LONG_LONG	unsigned long long		Annex A.1	
MPI_FLOAT	float			
MPI_DOUBLE	double	(Includes also special C++ types,	
MPI_LONG_DOUBLE	long double			
MPI_BYTE			e.g., bool,	
MPI_PACKED			see page 666	

Point-to-Point Communication

- Communication between two processes.
- Source process sends message to destination process.
- Communication takes place within a communicator, e.g., MPI_COMM_WORLD.
- Processes are identified by their ranks in the communicator.



Sending a Message

```
int MPI_Send( void* buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)
```

- <u>buf</u> is the starting point of the message with <u>count</u> elements, each described with <u>datatype</u>.
- dest is the rank of the destination process within the communicator comm.
- <u>tag</u> is an additional nonnegative integer piggyback information, additionally transferred with the message.
- The tag can be used by the program to distinguish different types of messages.

Receiving a Message

```
int MPI_Recv( void* buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Status* status)
```

- <u>buf/count/datatype</u> describe the receive buffer.
- Receiving the message sent by process with rank source in comm.
- Envelope information is returned in <u>status</u>.
- On can pass MPI_STATUS_IGNORE instead of a status argument.
- Output arguments are printed blue-cursive.
- Only messages with matching tag are received.

Requirements for Point-to-Point Communications

For a communication to succeed:

- Sender must specify a valid destination rank.
- Receiver must specify a valid source rank.
- The communicator must be the same.
- Tags must match.
- Buffer's type must match with the datatype handle (in the send and receive call)
- Message datatypes must match.
- Receiver's buffer must be large enough.

Wildcarding

- Receiver can wildcard.
- To receive from any <u>source</u>: source = MPI_ANY_SOURCE
- To receive from any tag: tag = MPI_ANY_TAG
- Actual source and tag are returned in the receiver's status parameter.

Communication Envelope

Envelope information is returned from MPI_RECV in status.

```
MPI_Status status;
status.MPI_SOURCE
status.MPI_TAG
status.MPI_ERROR
```

Receive Message Count

```
int MPI_Get_count( MPI_Status* status, MPI_Datatype datatype, int* count)
```

Communication Modes

Send communication modes:

synchronous send MPI_SSEND

buffered [asynchronous] send → MPI_BSEND

standard send -> MPI_SEND

Ready send -> MPI_RSEND

Receiving all modes -> MPI_RECV

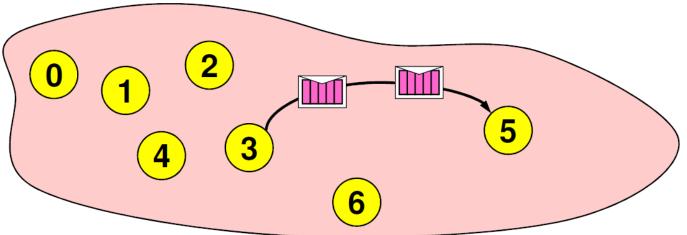
Sender mode	Definition	Notes
Synchronous send MPI_SSEND	Only completes when the receive has started	
Buffered send MPI_BSEND	Always completes (unless an error occurs), irrespective of receiver	needs application-defined buffer to be declared with MPI_BUFFER_ATTACH
Standard send MPI_SEND	Either synchronous or buffered	uses an internal buffer
Ready send MPI_RSEND	May be started only if the matching receive is already posted!	highly dangerous!
Receive MPI_RECV	Completes when a message has arrived	same routine for all communication modes

Communication Modes Rules

- Standard send (MPI_SEND)
 - minimal transfer time
 - may block due to synchronous mode
- Synchronous send (MPI_SSEND)
 - risk of deadlock
 - risk of waiting → idle time
 - high latency / best bandwidth
- Buffered send (MPI_BSEND)
 - low latency / bad bandwidth
- Ready send (MPI_RSEND)
 - use never, except you have a 200% guarantee that Recv is already called in the current version and all future versions of your code
 - may be the fastest

Message Order Preservation

- Rule for messages on the same connection,
- i.e., same communicator, source, and destination rank:
- Messages do not overtake each other.
- This is true even for non-synchronous sends.
- If both receives match both messages, then the order is preserved.



Outlook Collective Communications

- Two kinds
 - Data movement (broadcast, scatter, gather, etc.)
 - Collective computation (min, max, sum, logical OR etc.)
- Advantages
 - More convenient
- Can be used as short cut for an ensemble of P2P operations
 - More efficient
- Encapsulate sophisticated algorithms
- Implementation can take advantage of the structure of a machine to optimize and increase parallelism in these operations

Broadcast

- Broadcasts a message from the process with rank root to all other processes of the group.
 - buf = starting address of buffer
 - count = number of entries in buffer
 - datatype = data type of buffer
 - root = rank of broadcast root
 - comm = communicator

```
int MPI_Bcast( void* buf, int count, MPI_Datatype datatype, int root, MPI_Comm comm)
```

Reduce

- Combines the elements in the input buffer of each process using the operation op and returns the combined value in the output buffer of the process with rank root
 - sendbuf = address of send buffer
 - recvbuf = address of receive buffer
 - count = number of elements in send buffer
 - datatype = data type of elements of send buffer
 - op = reduce operation
 - root = rank of root process
 - comm = communicator

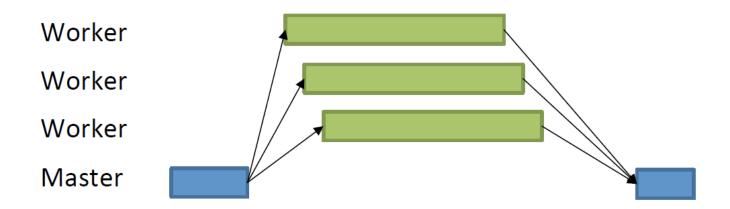
```
int MPI_Reduce( void* sendbuf, void* recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm)
```

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Master Worker

- Idea: self-scheduling algorithm
 - Master coordinates processing of tasks by providing input data to workers and collecting results
- Suitable if
 - Workers need not communicate with one another
 - Amount of work each worker must perform is difficult to predict
- Example: matrix-vector multiplication



Matrix-Vector Multiplication

$$A \cdot \vec{b} = \vec{c}$$

Unit of work:

dot product of one row of matrix A with vector b

Master

- Broadcasts b to each worker
- Sends one row to each worker
- Loop
 - Receives dot product from whichever worker sends one
 - Sends next task to that worker
 - Termination if all tasks are handed out

Worker

- Receives broadcast value of b
- while-Loop
 - Receives row from A
 - Forms dot product
 - Returns answer back to master

Matrix-Vector Multiplication: Common Part 1

```
#include <stdio.h>
#include <stdlib.h>
#include "mpi.h"
#define MAX ROWS 1000 // upper limits
#define MAX COLS 1000 // upper limits
#define MIN(a, b) ((a) > (b) ? (b) : (a))
#define DONE MAX ROWS+1
int main(int argc, char **argv) {
 double** A;
 double* b, * c, * buffer;
 double ans:
 int myid, master, numprocs, i, j, numsent, sender, done, anstype, row, rows, cols;
 MPI Status status;
 MPI Init(&argc, &argv);
 MPI Comm rank(MPI COMM WORLD, &myid);
 MPI Comm size(MPI COMM WORLD, &numprocs);
 A = (double **)calloc(MAX ROWS, sizeof(double *));
 for (i=0; i<MAX ROWS; i++)
  A[i] = (double *)calloc(MAX_COLS, sizeof(double));
 b = (double *)calloc(MAX COLS, sizeof(double));
 c = (double *)calloc(MAX ROWS, sizeof(double));
 buffer = (double *)calloc(MAX COLS, sizeof(double));
```

Matrix-Vector Multiplication: Common Part 2

```
master = 0;
rows = 2; // for debugging ...
cols = 2; // for debugging ...
if (myid == master) { /* master code: next slides */
                       /* worker code: next slides */
} else
// result c vector finished: considering the test values, all element values should be 3.0
for (i=0; i<MAX ROWS; i++)
 free(A[i]);
free(A);
free(b);
free(c);
free(buffer);
MPI Finalize();
return 0;
```

Matrix-Vector Multiplication: Master Part 1

Master

```
Broadcasts b to each worker
if (myid == master) {
                                           Sends one row to each worker
/* Initialize A and b (arbitrary) */
                                           Loop
A[0][0] = 1.0;

    Receives dot product from

A[0][1] = 2.0
                                                whichever worker sends one
A[1][0] = 1.0;

    Sends next task to that

A[1][1] = 2.0;
                                               worker
 b[0] = 1.0;

    Termination if all tasks are

 b[1] = 1.0;
                                                handed out
numsent = 0;
/* Send b to each worker process */
MPI_Bcast(b, cols, MPI_DOUBLE, master, MPI_COMM_WORLD);
/* Send a row to each worker process; tag with row number */
for (i = 0; i < MIN(numprocs - 1, rows); i++) {
 MPI_Send(&A[i][0], cols, MPI_DOUBLE, i+1, i, MPI_COMM_WORLD);
 numsent++;
```

Matrix-Vector Multiplication: Master Part 2

```
for (i = 0; i < rows; i++) {
 MPI_Recv(&ans, 1, MPI_DOUBLE, MPI_ANY_SOURCE, MPI_ANY_TAG,
                                           Master
           MPI COMM WORLD, &status);
                                              Broadcasts b to each worker
 sender = status.MPI SOURCE;
                                              Sends one row to each worker
 /* row is tag value */
                                              Loop
 anstype = status.MPI TAG;
                                                  Receives dot product from
 c[anstype] = ans;
                                                  whichever worker sends one
 /* send another row */
                                                 Sends next task to that
 if (numsent < rows) {</pre>
                                                  worker
  MPI_Send(&A[numsent][0], cols,
                                                 Termination if all tasks are
            MPI DOUBLE, sender,
                                                  handed out
            numsent, MPI COMM WORLD);
  numsent++;
 } else {
  /* Tell sender that there is no more work */
  MPI Send(MPI BOTTOM, 0, MPI DOUBLE, sender, DONE, MPI COMM WORLD);
} /* end of master specific part*/
```

/* start of worker code: next slide */

Matrix-Vector Multiplication: Worker Part

```
MPI_Bcast(b, cols, MPI_DOUBLE, master, MPI_COMM_WORLD);
/* Skip if more processes than work */
done = myid > rows;
while (!done) {
 MPI_Recv(buffer, cols, MPI_DOUBLE, master, MPI_ANY_TAG,
           MPI COMM WORLD, &status);
 done = status.MPI TAG == DONE;
 if (!done) {
  row = status.MPI TAG;
  ans = 0.0;
  for (i = 0; i < cols; i++) {
   ans += buffer[i] * b[i];
  MPI Send(&ans, 1, MPI DOUBLE, master, row, MPI COMM WORLD);
                               Worker
```

/* end of worker code */

- Receives broadcast value of b
- while-Loop
 - Receives row from A
 - Forms dot product
 - · Returns answer back to master

Outline

- Introduction to the Lecture
- Distributed Parallel Computing I
 - Primer
 - Overview
 - Process Model and Language Bindings
 - Messages and Point-to-Point Communication
 - Examples
- Quiz

Quiz

- Q1: How is it ensured that a specific message is received by a specific process?
- Q2: What is the first and last routine to be called in a MPI program?
- Q3: Is "MPI_Init" executed by one, several or all MPI processes?
- Q4: Name typical reduction operations?
- Q5: How can a point-to-point communication be made non-blocking? What potential advantage is there?