

Ans to the Q. no:

a) Huffman coding:

Huffman code is data compression algorithm.

In our computer each character stored as binary code of fixed length. The Huffman coding algorithm doesn't use fixed length code.

Huffman coding uses shorter length code and which use more number of time and a longer length code which used less number of time so we can say that Huffman

coding has greedy choice property. We know greedy choice algorithm make the choice at logically optimal and best at the moment so it is true.

4 repeat.

Huffman (6)

my id: 193002101

$$x = 193002101 \times 100$$

$$= 1$$

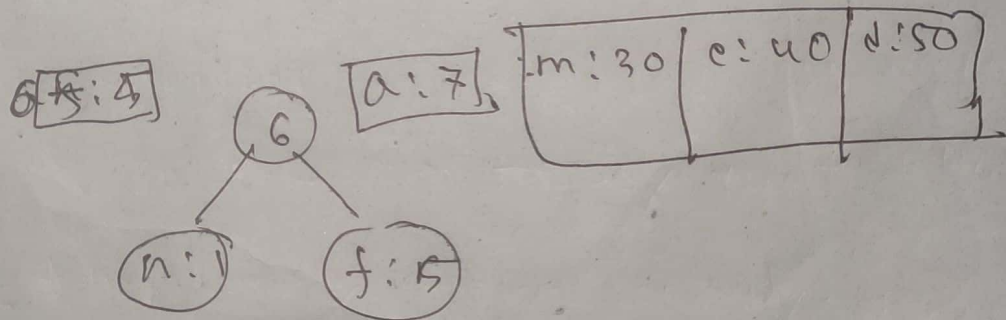
now,

a:7	e:40	d:5	f:5	m:30	n:1
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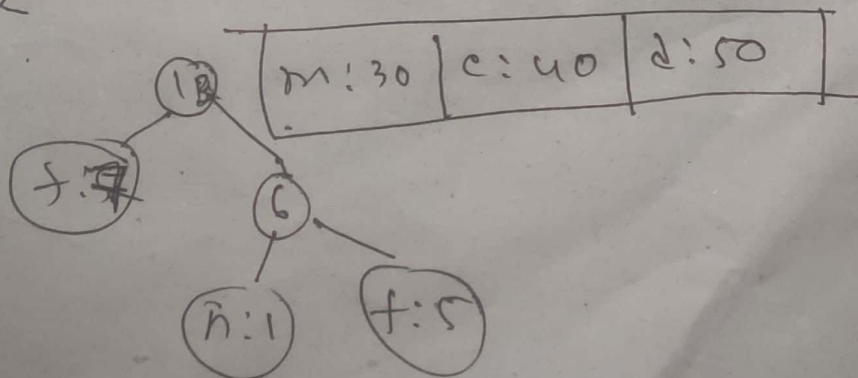
so,

n:1	f:5	a:7	m:30	e:40	d:50
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Step-1

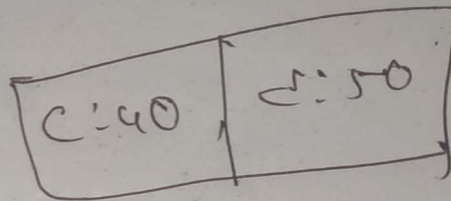
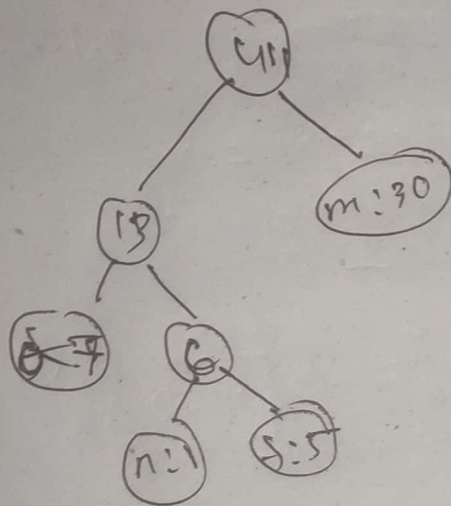


Step-2

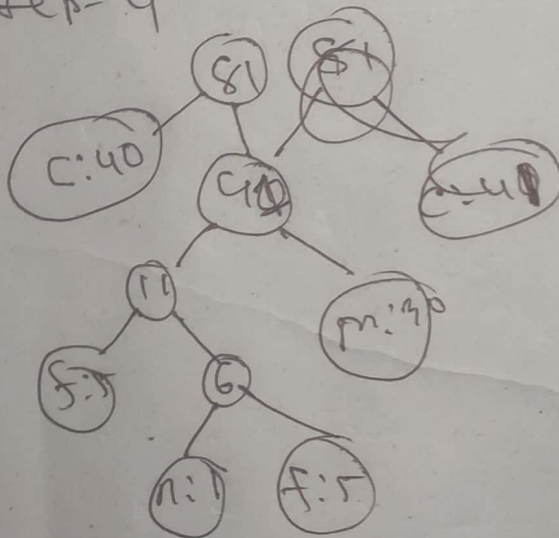


4. repeat.

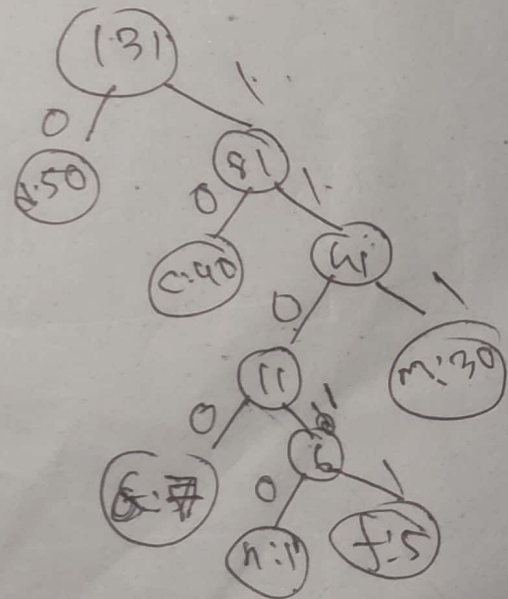
Step-3:



Step-4



Step-5



Nov 1

Character Prefix code

a ————— 1100

c ————— 10

d ————— 0

f ————— 11011

m ←———— 11

n ←———— 11010

Ans to the Q.no

Divide and conquer	Dynamic Programming
Partitions a problem into independent smaller sub problems.	1. Partitions a problem into overlapping sub problems.
2. Don't store solutions of sub-problems.	2. stores solutions of sub problems. thus avoids calculations of same quantity twice.

Ans to the Q.NO.

KMP:

KMP:

$n = \text{size of } T = 14$
 $= \text{size of } P = 5$

$T = \text{A B A A D A B A A C A B A}$
 $P = \text{A B A A C}$

Step-2 $i=2, q=0$

~~Step-2~~
A B A A D A B A A C A B A
A B A A C

Step-3
A B A A D A B A A C A B A
A B A A C

Step-4
A B A A D A A B A A C A B A
A B A A C

Step-5
A B A A D A A B A A C A B A

Step-6
A B A A D A A B A A C A B A
A B A A C

Step-7

A B A A ~~D~~ A A B A A C A B A
A B A A C

Step-8

A B A A D A A B A A C A B A
A B A A C

Step-9

A B A A D A A B A A C A B A
A B A A C

Step-10

A B A A D A A B A A C A B A
A B A A C

Step-11

A B

Difference between NP-complete and

NP-complete	NP-hard problems
NP-complete problems can be solved by Turing machine in polynomial time	NP-hard problem can be solved if and only if there is a NP-complete problem that it can be reduced into it in polynomial time
It is exclusively a decision problem	It does not have to be a NP to solve this problem

Ans to the Q.no:

Q no id: 193002101

$$= 193002101 \times 20 + 8$$

$$= 1 + 8$$

$$= 9$$

Now,

i item	weight (w_i)	value
1	2	24
2	3	27
3	4	36
4	1	12
5	2	9(n)

we know,

$$P(i, w) = \max \{ v_i + P(i-1, w - w_i), P(i-1, w) \}$$

	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	0	24	24	24	24	24	24
2	0	0	24	27	27	51	51	51
3	0	0	24	27	36	36	60	63
4	0	12	24	36	39	48	60	72
5	0	12	24	36	39	48	60	72

minimum profit : 72