```
1. randomly choose two prime numbers p and q
       p=11; q=13
2. compute n=pq
      n=143
3. randomly choose an odd number e in the range 1<e<phi(n)
       phi(n) is the totient function (number of positive integers <=n that are relatively prime to n)
       relatively prime means they do not contain any common factors
       1 is relatively prime to all numbers
       phi(24)=8 since there are 8 totatives of 24: 1,5,7,11,13,17,19,23
       phi(n)=phi(p)*phi(q)
       phi(11)=10 (1,2,3,4,5,6,7,8,9,10)
       phi(13)=12 (1,2,3,4,5,6,7,8,9,10,11,12)
       phi(n)=10*12=120
       1<e<120; we pick 7
4. compute d=e^-1 (mod phi(n)) by Euclid's algorithm
       thus de=1 \pmod{phi(n)}
       d=7^-1 \pmod{phi(143)}
       d=7^-1 \pmod{120}
       d is the inverse of 7 (mod 120)
       so 7d = 1 \pmod{120}
       so we can do euclidean's algorithm on 120 and 7 (we find that t=-17)
              -17+120=103; so t=103
5. publish (n,e) as the public key; keep d as the secret ket
       (143,7) is the public key
       (143,103) is the secret key
encryption
       E = encryption algorithm
       A = user
       m = message (0 \le m \le n A)
       c = ciphertext (0 \le c \le n A)
       c = E \ a(m) = m^{(e)} \ a) \% n \ A
       encrypt m=3
              c=m^{(e A)} % n A
              c=3^7 % 143
              c=2187 % 143
              c = 42
```

decryption

```
 \begin{aligned} m &= D_A(c) = c^{\prime}(d_A) \% \text{ n\_A} \\ \text{decrypt } c &= 42 \\ \text{m=}c^{\prime}(d_A) \% \text{ n\_A} \\ \text{m=}42^{\prime}103 \% 143 \\ \text{too large so use:} & a = bc \text{ mod n} = (b \text{ mod n})(c \text{ mod n}) \text{ mod n} \\ & 42^{\prime}103 = 42^{\prime}(2 + 2 + ... + 2 + 1) \text{ (power of 2, 51 times)} \\ \text{m=}([(42^{\prime}2 \% 143)^{\prime}51]^*42) \% 143 \\ \text{m=}3 \end{aligned}
```

the euclidean algorithm gives us the gcd of 2 integers we repeatedly divide the divisor by the remainder until the remainder is 0 gcd is the last non-zero remainder

```
gcd(81,57): 81 = 1(57) + 24

57 = 2(24) + 9

24 = 2(9) + 6

9 = 1(6) + 3

6 = 2(3) + 0
```

so gcd(81,57)=3

if gcd(a,b)=r then there exists integers s and t such that s(a)+t(b)=r we can get s and t by reversing the steps of euclidean's algorithm start with the non-0 remainder line and rewrite: 3 = 9 - 1(6) substitute for 6 by using the line above that: 3 = 9 - 1(24 - 2(9)) = 3(9) - 1(24) substitute for 9 by using the line above that: 3 = 3(57 - 2(24)) - 1(24) = 3(57) - 7(24)

3 = 3(57) - 7(81 - 1(57))

= 10(57) - 7(81)

substitute for 24 by using the line above that: so s=-7 and t=10

find gcd(120,23):
$$120 = 5(23) + 5$$
$$23 = 4(5) + 3$$
$$5 = 1(3) + 2$$
$$3 = 1(2) + 1$$
$$2 = 2(1) + 0$$

so gcd(120,23)=1 (coprime) now find s and t: 1=3

$$1 = 3 - 1(2)$$

$$1 = 3 - 1(5 - 1(3)) = 2(3) - 1(5)$$

$$1 = 2(23 - 4(5)) - 1(5) = 2(23) - 9(5)$$

$$1 = 2(23) - 9(120 - 5(23)) = 47(5) - 9(120)$$

so s=-9 and t=47

now we can say that: s is inverse of a (mod b)

 $as = 1 \pmod{b}$

and: t is the inverse of b (mod a)

 $bt = 1 \pmod{a}$