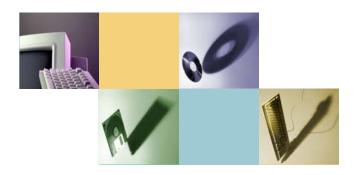
#### **PTHREADS**



**CSC 345** 

This powerpoint is based on YoLinux Tutorial: POSIX thread (pthread) libraries, http://www.yolinux.com/TUTORIALS/LinuxTutorialPosixThreads.html

# The POSIX thread (pthread) libraries

- standards based thread API for C/C++.
- spawn a new concurrent process flow.
- the process flow can be scheduled to run on another processor
- Less overhead than "forking"
- Share the same process address space no initialization for a new system virtual memory space.
- Also gain performance on uniprocessor since one thread may execute while another is waiting for I/O
- Threads are only for a single computer system. MPI processes are across multiple machines or a distributed computing environment.



#### **Basics**

- Thread operations
  - thread creation
  - termination,
  - synchronization (joins, blocking), scheduling,
  - data management
  - process interaction.







- A thread does not maintain a list of created threads, nor does it know the thread that created it.
- All threads within a process share the same address space.
  - Threads in the same process share:
  - Process instructions
  - Most data
  - open files
  - signals and signal handlers
  - current working directory
  - User and group id
- Each thread has a unique:
  - Thread ID
  - set of registers, stack pointer
  - stack for local variables, return addresses
  - signal mask
  - priority
  - Return value: errno
- pthread functions return "0" if ok.







Create an independent thread

- Normally, we wait till the created thread finishs before the main thread continues int pthread\_join(pthread\_t th, void \*\* thread\_return);
  - suspends the execution of the calling thread until the thread identified by th terminates



```
See example ptl.c in ~box/directory

• Compile:

- gcc-lpthread ptl.c

or

- g++-lpthread ptl.c

[box@oscar box]$ gcc -lpthread ptl.c

[box@oscar box]$ ./a.out

Thread 1

Thread 2

Thread 1 returns: 0

Thread 2 returns: 0
```

- Threads terminate by
  - just returning from the function or
  - explicitly calling pthread\_exit
  - by a call to the function exit which will terminate the process including any threads.







## From example

- Threads terminate by
  - just returning from the function or explicitly calling pthread\_exit

  - by a call to the function exit which will terminate the process including any threads.
- int pthread\_create(pthread\_t \* thread, const pthread\_attr\_t \* attr, void \* (\*start\_routine)(void \*), void \*arg);
- thread returns the thread id. (unsigned long int defined in bits/pthreadtypes.h)
- attr Set to NULL if default thread attributes are used. (else define members of the struct pthread\_attr\_t defined in bits/pthreadtypes.h) Attributes include:

   detached state (joinable? Default: PTHREAD\_CREATE\_JOINABLE. Other option: PTHREAD\_CREATE\_DETACHED)

   scheduling policy (real-time? PTHREAD\_INHERIT\_SCHED,PTHREAD\_EXPLICIT\_SCHED,SCHED\_OTHER)

  - scheduling parameter
  - inheritsched attribute (Default: PTHREAD\_EXPLICIT\_SCHED Inherit from parent thread: PTHREAD\_INHERIT\_SCHED)
  - scope (Kernel threads: PTHREAD\_SCOPE\_SYSTEM User threads: PTHREAD\_SCOPE\_PROCESS Pick one or the other not both.)
  - guard size
  - stack address (See unistd.h and bits/posix\_opt.h \_POSIX\_THREAD\_ATTR\_STACKADDR)
  - stack size (default minimum PTHREAD\_STACK\_SIZE set in pthread.h),
- void  $^{\star}$  (\*start\_routine) pointer to the function to be threaded. Function has a single argument: pointer to void.
- \*arg pointer to argument of function. To pass multiple arguments, send a pointer to a structure.







### synchronization

- Pthread provides three synchronization mechanisms:
- mutexe Mutual exclusion lock is a blocking access to prevent racing condition. It enforces exclusive access by a thread to a variable or set of variables.
- 2. join Make a thread wait till others are complete (terminated).
- 3. condition variables data type pthread\_cond\_t







#### **MUTEX**

- A race condition often occurs when two or more threads competing on the same memory area
- Results of computations depends on the order in which the operations are executed.
- We can use mutex to access a critical section







## Example w/o mutex

```
int counter=0;
/* Function C */
void functionC() {
  counter++
}
```







## What could possibly be the problem?

• The statement "count++" for thread1 in machine language as:

```
register1 = counter
register1 = register1 + 1
counter = register1
```

The statement "count++" for thread2 implemented as:

```
register2 = counter
register2 = register2 + 1
counter = register2
```

- w/o mutex answers are unpredictable
- With mutex- we can protect the critical section and allow only one thread in CS at a time.
- Show pt2.c







## **Example with mutex**

```
pthread_mutex_t mutex1 =
   PTHREAD_MUTEX_INITIALIZER;
int counter=0;
/* Function C */
void functionC() {
   pthread_mutex_lock( &mutex1 );
   counter++;
   pthread_mutex_unlock( &mutex1 );
}
```



#### Join

- A function to wait for the completion of the threads with a join.
- A thread calling routine may launch multiple threads then wait for them to finish to get the results.
- Show pt3.c





# pthread\_cond\_t

- A condition variable is used with the appropriate functions for waiting and later, process continuation.
- allows threads to suspend execution and give up the processor until a given condition is true.
- must always be associated with a mutex to avoid a race condition.



## pthread\_cond\_t (continued)

- Functions used in conjunction with the condition variable:
- Creating/Destroying:
  - pthread\_cond\_init
  - pthread\_cond\_t cond = PTHREAD\_COND\_INITIALIZER;
  - pthread\_cond\_destroy
- Waiting on condition:
  - pthread cond wait
  - <u>pthread\_cond\_timedwait</u> place limit on how long it will block.
- · Waking thread based on condition:
  - pthread\_cond\_signal
  - <u>pthread\_cond\_broadcast</u> wake up all threads blocked by the specified condition variable.







# Example for pthread\_cond\_t

See pt4.c







### Thread Scheduling

- When the option is enabled, each thread may have its own scheduling properties.
   Scheduling attributes may be specified:
  - during thread creation
  - by dynamically by changing the attributes of a thread already created
  - by defining the effect of a mutex on the thread's scheduling when creating a mutex
  - by dynamically changing the scheduling of a thread during synchronization operations.
  - The threads library provides default values that are sufficient for most cases.







#### Thread safe

- is threaded routine code that must call functions which are "thread safe".
- no static or global variables which other threads may cause a racing condition
- If static or global variables are used then mutexes must be applied or the functions must be re-written to avoid the use of these variables.





